

Lecture 6

Binary Exponential Backoff Algorithm

How should nodes that collide schedule their reattempt?

- If they simply use CSMA, then they will likely collide again
- They could independently choose a random delay to wait before reattempting
- The more the collisions, the random delay should be chosen over a larger range to reduce the chance of collision

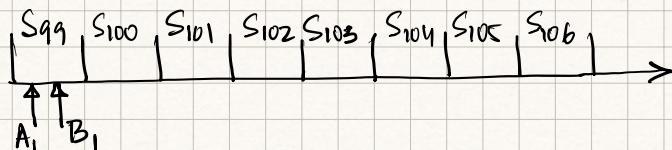
n : number of reattempt for the pkt being transmitted

$n = 0$ when a new pkt is transmitted (initial value)

We consider a slotted system for simplicity

Algorithm: On detecting collision, the number of slots to delay to reattempt, τ is chosen randomly in the range $(0, 2^k)$ where $k = \min(n, 16)$

Example: We consider 2 nodes & a slotted system $(A \& B)$



- A_1 and B_1 refers to pkt arrivals in nodes A & B respectively. $n=0$ for both these pkts
- Both nodes transmit in S_{100} and they collide
 Node A chooses randomly from $(0, 1)$
 say it choose 0 n is incremented by 1
- Node B chooses randomly from $(0, 1)$
 say it choose 0 n is incremented by 1
- Both nodes transmit in S_{101} and they collide
 Node A chooses randomly from $(0, 2)$
 say it chooses 1 n is now 2
- Node B chooses randomly from $(0, 2)$
 say it chooses 1 n is now 2
- Both nodes transmit in S_{103} and they collide
 Node A chooses randomly from $(0, 4)$
 say it chooses 0 n is now 3
- Node B chooses randomly from $(0, 4)$
 say it chooses 2 n is now 3
- Node A successfully transmits in S_{104}
- Node B successfully transmits in S_{106}

Notes:

- The range over which τ (the number slots to delay) is randomly picked expands exponentially with collisions

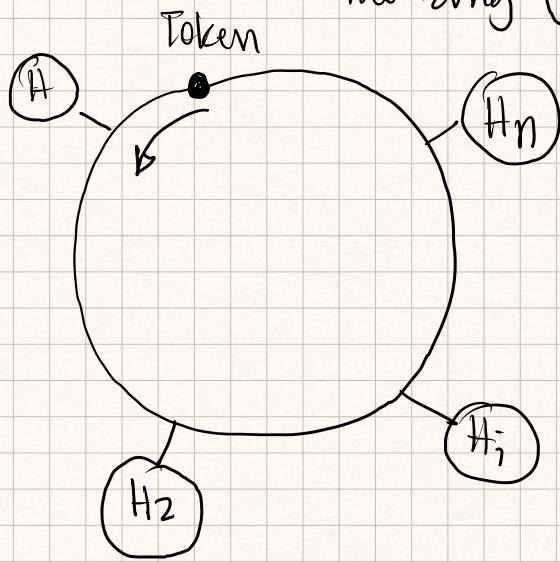
- n is the current number of retransmits and it is incremented after determining τ
- n is reset to 0 after successful transmission
- n is capped at 10 through the function
 $k = \min(n, 10)$ 2^{10} is 1024 which is large
- In the real algorithm, the time delay until retransmit is $\tau * (\text{transmission time of } 512 \frac{\text{bytes}}{\text{bit}})$
or $?$ bits

Project 1: We will see that this algorithm works scalably well.

Taking Turns Medium Access Protocol

- Nodes share a broadcast channel and they take turns in accessing the channels and so only 1 node is transmitting at a time
- Many ways to implement this. The most well known method is called the Token Ring Protocol
 - known as the FDDI (Fiber Distributed Data Interface)
 - Nodes are connected in a ring network

→ A token (a special bit pattern) goes around the ring (in one direction)



When a node has a pkt. to transmit

→ Waits for the token to arrive

→ Hold the token

→ Transmit pkt(s)

→ Releases token

Service Policies:

Limited : only 1 pkt is transmitted per token

Gated : All pkts in the queue when the token arrives are transmitted.
New pkts that arrive during service are buffered

Exhaustive : The token is released when the queue becomes empty

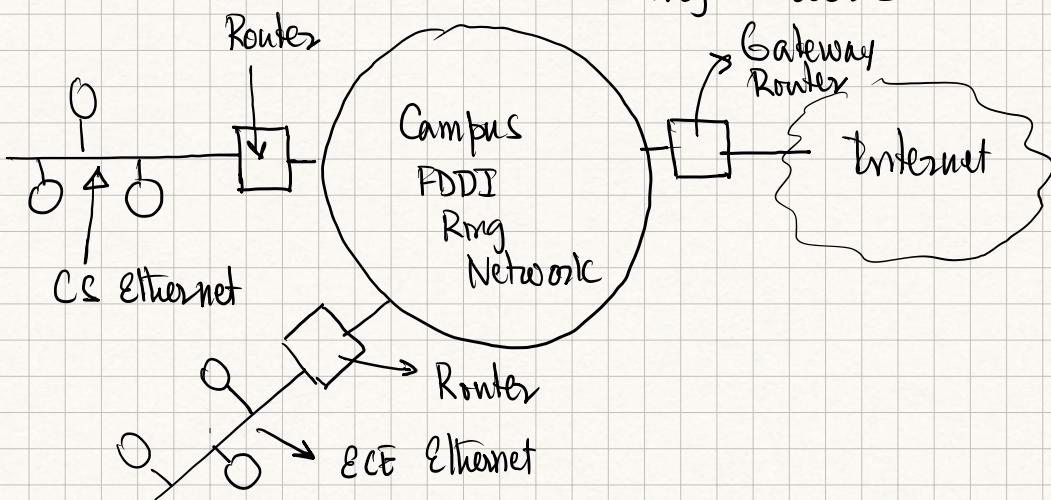
Token Rotation Time : Time between two consecutive arrivals of the token at a node

Minimum TRT : Propagation delay + bit delays at the repeaters

Maximum TRT : Depends on the service policy

For the limited service policy
Min TRT + a transmission time at each node

Example: UC Davis campus backbone was a FDDI token ring network



Slotted Ring Network

- Instead of a token, slots go around the ring
- Slots are marked Free or Busy
- A node wishing to transmit waits for a Free slot, marks the slot as Busy, and writes data into the data part of the slot
- The destination node reads the data and marks the slot as Free (Destination Removal)

