

1. 解 对于每一组 (k_i, c_i) , 都可计算得

$$C^*(k_i, c_i) = \begin{cases} \frac{k_i}{c_i} + \frac{h\alpha_i}{2} + c_i \lambda & Q_{i+1} > \sqrt{\frac{2k_i \lambda}{h}} \\ \sqrt{2\lambda h k_i} + c_i \lambda & Q_{i+1} \leq \sqrt{\frac{2k_i \lambda}{h}} < Q_{i+1} \\ \frac{k_i \lambda}{c_{i+1}} + \frac{h\alpha_{i+1}}{2} + c_i \lambda & \sqrt{\frac{2k_i \lambda}{h}} \geq Q_{i+1} \end{cases}$$

取 $\arg\max_i C^*(k_i, c_i)$, 然后再找出最优决策 Q_i 即可. \square

2. Suppose, in the newsvendor model (basic model), that demand D is discrete in the sense that there exists a sequence of real numbers, $\{x_1, x_2, \dots\}$, representing all potential demand quantities, and $0 < x_1 < x_2 < \dots$. Assume that $\mathbb{P}(D = x_i) = q_i$, $i = 1, 2, \dots$. Furthermore, suppose all these quantities are admissible stock levels. Prove that the smallest x_i such that

$$\Phi(x_i) \geq \frac{p - c}{p}$$

is an optimal stock level.

Hint: In two steps:

- (a) First show that for any $y \in (x_i, x_{i+1})$, $V(y) \leq \max\{V(x_i), V(x_{i+1})\}$, $i = 0, 1, 2, \dots$, where $x_0 = 0$;
- (b) Second show that the smallest x_i such that $V(x_{i+1}) - V(x_i) \leq 0$ is an optimal stock level.

(a) 解 当 $y \in (x_i, x_{i+1})$ 时 $V(y) = (p - c)E(D|y) + pE(D|y)^-$

$$\begin{aligned} &= \sum_{j=1}^{\infty} (p - c)q_j(x_j - y) + \sum_{j=1}^i p q_j(y - x_j) \\ &= \left(\sum_{j=1}^i (p - c)q_j - \sum_{j=1}^i p q_j \right) x_j + \left(\sum_{j=1}^i p q_j - \sum_{j=1}^{\infty} (p - c)q_j \right) y \\ &= (p - c - \sum_{j=1}^i p q_j) x_j + (\sum_{j=1}^i p q_j - (p - c)) y \end{aligned}$$

若 $\sum_{j=1}^i p q_j - (p - c) > 0$ 则 $V(y) \leq V(x_{i+1})$ 否则 $V(y) \leq V(x_i)$

(b) 解 由于 $\sum_{j=1}^i p q_j - (p - c) \downarrow$ 故 \downarrow 得证.

若 $\sum_{j=1}^i p q_j - (p - c) \text{ 前项} \geq 0$ 则 $\sum_{j=1}^i p q_j - (p - c) = p(x_0) - (p - c) \geq 0$
 $\Rightarrow x_0 \geq \frac{p - c}{p}$ \square

3. Suppose that g is a convex function defined on \mathbb{R} and the real valued function G is defined on \mathbb{R} by $G(y) := \mathbb{E}[g(y - D)]$, where D is a random variable with a density ϕ . Prove, without assuming that g is differentiable everywhere, that G is convex on \mathbb{R} .

$$\begin{aligned} G(\lambda_1 y + \lambda_2 y) &= E[g(\lambda_1 y + \lambda_2 y - D)] = E[g(\lambda_1 y - \lambda_1 D + \lambda_2 y - \lambda_2 D)] \\ &= E[\lambda_1 g(y - D) + \lambda_2 g(y - D)] \\ &= \lambda_1 E[g(y - D)] + \lambda_2 E[g(y - D)]. \quad \square \end{aligned}$$

4. The dynamic inventory problem of Section 4.1 assume that all shortages were backlogged. Assume now that the unsatisfied demand are lost (i.e., lost-sales model). A unit penalty cost of p is charged for each unit that is lost in each period. Note that the salvage value function $v_T(x) = -cx$ is now defined only for $x \geq 0$. All other aspects of the problem remain unchanged.

(a) Write down the optimality equations for this problem.

(b) Show that when $p + h > \alpha c$, a base stock policy with base stock level S is optimal in each period, where S is the following fractile

$$\Phi(S) = \frac{p - c}{p + h - \alpha c}. \quad (4.8)$$

(We assume all functions that will be used are differentiable)

(a) 解一个周期内的费用为

$$L(y) = h \int_0^y (y - \xi) \phi(\xi) d\xi + p \int_y^\infty (\xi - y) \phi(\xi) d\xi$$

求最优方程为

$$f_t(x) = \min_{y \geq x} \left\{ c(y - x) + h \int_0^y (y - \xi) \phi(\xi) d\xi + p \int_y^\infty (\xi - y) \phi(\xi) d\xi + \alpha \int_0^y f_{t+1}(\xi) \phi(\xi) d\xi \right\}$$

$$t = 1, 2, \dots, N$$

$$f_{N+1}(x) = \begin{cases} -cx & x \geq 0 \\ 0 & x < 0 \end{cases}$$

(b) 考虑第 N 个状态有

$$\begin{aligned} f_t(x_N) &= \min_{y \geq x_N} \left\{ c(y_N - x_N) + h \int_0^{y_N} (y_N - \xi) \phi(\xi) d\xi + p \int_{y_N}^\infty (\xi - y_N) \phi(\xi) d\xi \right. \\ &\quad \left. - \alpha \int_0^{y_N} f_{t+1}(\xi) \phi(\xi) d\xi \right\} \end{aligned}$$

$$\text{设 } g(y) = c(y-x) + \int_0^y (y-\xi) \phi(\xi) d\xi + p \int_y^\infty (\xi-y) \phi(\xi) d\xi - \alpha \int_0^y c(y-\xi) \phi(\xi) d\xi$$

$$\text{于是 } g'(y) = c + h\bar{\pi}(y) - p(1-\bar{\pi}(y)) - \alpha c(\bar{\pi}(y))$$

$$g'(y) = (h+p-\alpha c)\bar{\pi}(y) \quad \text{由于 } h+p-\alpha c > 0 \Rightarrow g'(y) \geq 0$$

故 g 是一个凸函数. 当 $g'(y)=0$ 即 $\bar{\pi}(y)=\frac{p-c}{p+h-\alpha c}$ 时有最小值.

即第 N 个周期的订货量 $\bar{\pi}(S_N) = \frac{p-c}{p+h-\alpha c}$. 下考虑前 $N-1$ 个周期

$$\text{注意到 } \left(h \int_0^y (y-\xi) \phi(\xi) d\xi + p \int_y^\infty (\xi-y) \phi(\xi) d\xi \right)' = (h+p)\bar{\pi}(y)$$

$$h+p-\alpha c \Rightarrow h+p > 0 \Rightarrow h \int_0^y (y-\xi) \phi(\xi) d\xi + p \int_y^\infty (\xi-y) \phi(\xi) d\xi \in \mathbb{R}^+$$

于是若 f_{t+1} 凸, $g_t(y) := c(y-x) + h \int_0^y (y-\xi) \phi(\xi) d\xi + p \int_y^\infty (\xi-y) \phi(\xi) d\xi + \alpha \int_0^y f_{t+1}(t+\xi) \phi(\xi) d\xi$ 凸. 又 $f_{t+1} = \min_{y \in \mathbb{R}} \{g_t(y)\}$ 凸, 由 Lemma 4.1, g_t 的 $\min, \max (= s)$ 是 base stock policy. 由 Theorem 4.1 知 A base stock policy 是最优的.

$$\text{若 } g_t'(s) = c - p + h\bar{\pi}(s) + p\bar{\pi}(s) + \alpha \int_0^s f_{t+1}'(s-\xi) \phi(\xi) d\xi$$

$$\text{又 } f_{t+1}'(s-\xi) = (g_t(s) - c(s-\xi))' = -c \Rightarrow$$

$$g_t'(s) = c - p + h\bar{\pi}(s) + p\bar{\pi}(s) - \alpha c\bar{\pi}(s) = 0 \Rightarrow \text{满足 } \bar{\pi}(s) = \frac{p-c}{h+p-\alpha c} \text{ 的 } s.$$

是 +1 周期的最低策略



5. Proof Lemma 4.3

$$(1) f \text{ is } \mathbb{B} \Rightarrow f((1-\lambda)x_1 + \lambda x_2) \leq (1-\lambda)f(x_1) + \lambda f(x_2) + \lambda k$$

$$\begin{aligned} \Rightarrow \alpha f((1-\lambda)x_1 + \lambda x_2) &\leq \alpha(1-\lambda)f(x_1) + \alpha\lambda f(x_2) + \alpha\lambda k \\ &= (1-\lambda)\alpha f(x_1) + \lambda\alpha f(x_2) + \alpha\lambda k. \end{aligned}$$

$\Rightarrow f$ is k -凸的 $\forall k \geq \alpha k$

$$(2) f \text{ is } \mathbb{B}, g \text{ is } \mathbb{B} \Rightarrow f((1-\lambda)x_1 + \lambda x_2) \leq (1-\lambda)f(x_1) + \lambda f(x_2) + \lambda k \quad (1)$$

$$g((1-\lambda)x_1 + \lambda x_2) \leq (1-\lambda)g(x_1) + \lambda g(x_2) + \lambda k \quad (2)$$

(1) + (2) 之得

$$(3) G((1-\lambda)y_1 + \lambda y_2) = E[V((1-\lambda)y_1 + \lambda y_2 - D)] \leq E[(1-\lambda)V(y_1 - D) + \lambda V(y_2 - D) + \lambda k] \\ = (1-\lambda)G(y_1) + \lambda G(y_2) + \lambda k$$

$$(4) \text{ if } z = (1-\lambda)x + \lambda y$$

$$\Rightarrow f(z) \leq (1-\lambda)f(x) + \lambda f(y) + k = (1-\lambda)[k + f(y)] + \lambda f(y) + k \leq f(y) + k$$

$$(5) (a) \text{ 由定义 } \alpha f(x \vee x') + \alpha f(x \wedge x') = \alpha(f(x \vee x') + f(x \wedge x')) \\ = \alpha(f(x) + f(x')) = \alpha f(x) + \alpha f(x') \quad \forall \alpha \geq 0.$$

$$(b) \text{ 由定义 } f(x \vee x') + f(x \wedge x') \leq f(x) + f(x')$$

$$g(x \vee x') + g(x \wedge x') \leq g(x) + g(x')$$

$$(f+g)(x \vee x') + (f+g)(x \wedge x') \leq (f+g)(x) + (f+g)(x')$$

$$(6) (a) \text{ 由定义 } h(x \vee x') + h(x \wedge x') = E[f(x \vee x' - D) + f(x \wedge x' - D)] \\ = E[f(x - D) \vee f(x' - D)] + E[f(x - D) \wedge f(x' - D)] \\ \leq E[f(x - D) + f(x' - D)] = E[f(x - D)] + E[f(x' - D)] \\ = h(x) + h(x')$$