#### 1. Problem

K-th Element of Two Sorted Arrays

**Input**: Two sorted arrays of size m and n respectively

**Output**: The element that would be at the k'th position of the final sorted array.

**Example**: A1 = [2, 3, 6, 7] A2 = [1, 4, 5] k = 5

Final sorted array: [1, 2, 3, 4, 5, 6, 7],  $5^{th}$  element = 5

	Time complexity	Space complexity
Merge 2 sorted arrays:	O(k)	O(1)
Decrease and conquer:	$O(\min(m, n, k))$	O(1)
Decrease and conquer 2:	$O(\log(\min(m, n, k)))$	O(1)

m and n are the size of A1 and A2, respectively k is the position of the required element

# 1) Merge 2 sorted arrays



**2 3** ... A2

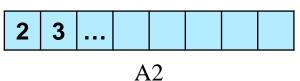
k = 3

Compare: 1st element of A1

4



1st element of A2



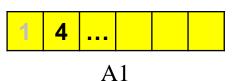
1st of final array: 1

Remaining k = 2

Compare: 2<sup>nd</sup> element of A1



1st element of A2



**A**1

2 3 ....

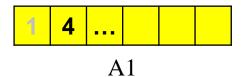
2<sup>nd</sup> of final array: 2

Remaining k = 1

Compare: 2<sup>nd</sup> element of A1



2<sup>nd</sup> element of A2



**2 3 ...** A2

A2

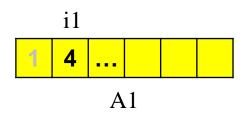
Remaining k = 0

3<sup>rd</sup> of final array: 3

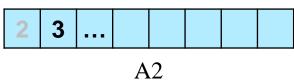
### 3. Algorithms

# K-th Element of Two Sorted Arrays

### 1) Merge 2 sorted arrays



**i**2



Time complexity: O(k)

Space complexity: O(1)

#### Pseudocode:

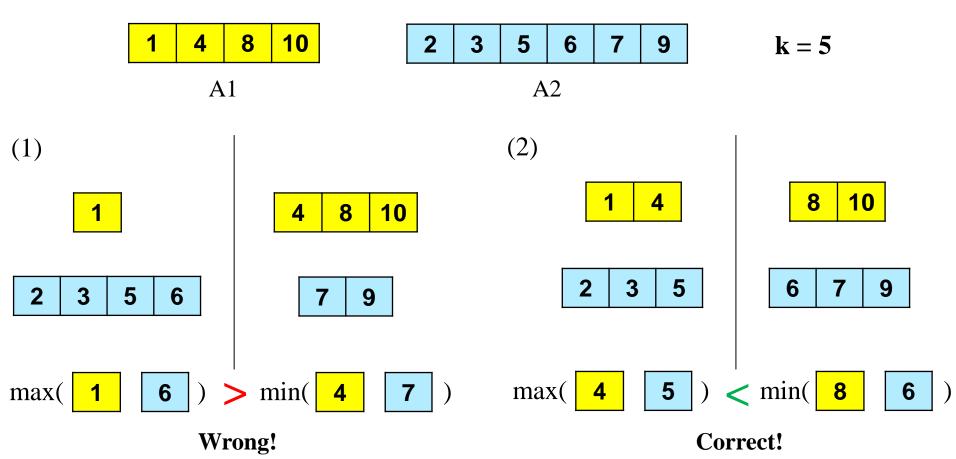
#### kthElement (A1[0..m-1], A2[0..n-1], k)

**Input:** : Two sorted arrays of size m and n respectively

**Output:** : The element that would be at the k'th position of the final sorted array.

- 1.  $i1 \leftarrow 0$ ;  $i2 \leftarrow 0$ ; res = null
- 2. **while** i1 < m and i2 < n **do**
- 3. **if** k = 0 **then return** res
- 4. **if** A1[i1] < A2[i2] **then**
- 5.  $res \leftarrow A1[i1]; i1 \leftarrow i1 + 1; k = k 1$
- 6. **else**
- 7.  $res \leftarrow A2[i2]; i2 \leftarrow i2 + 1; k = k 1$
- 8. **while** i1 < m do
- 9. **if** k = 0 **then return** res
- 10. res  $\leftarrow$  A1[i1]; i1  $\leftarrow$  i1 + 1; k = k 1
- 11. **while** i2 < n **do**
- 12. **if** k = 0 **then return** res
- 13. res  $\leftarrow$  A2[i2]; i2  $\leftarrow$  i2 + 1; k = k 1

2) Decrease and conquer



**Condition:** max\_left <= min\_right

### 2) Decrease and conquer

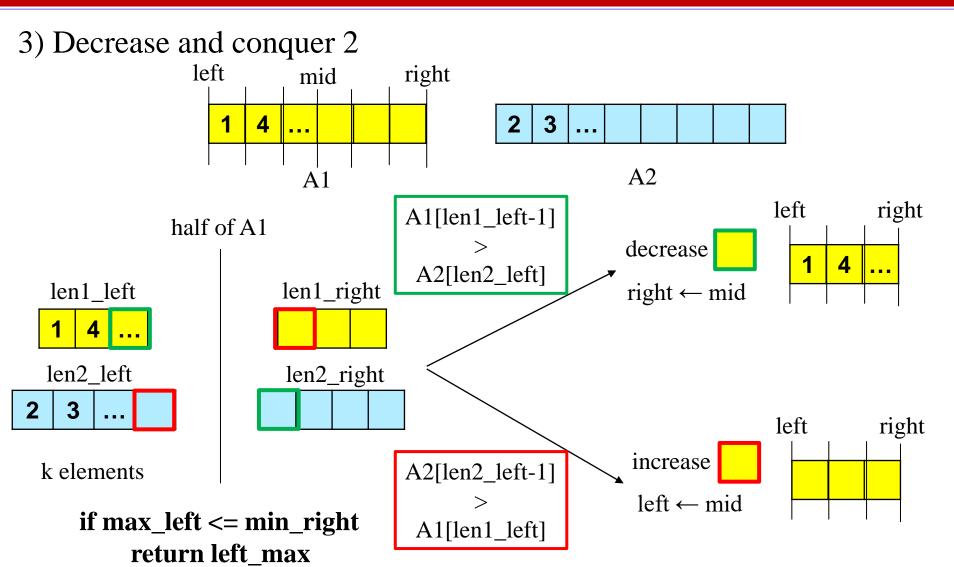
#### Pseudocode:

```
kthElement (A1[0..m-1], A2[0..n-1], k)
Input: : Two sorted arrays of size m and n. m \le n.
Output: : The element that would be at the k'th position of
the final sorted array.
   for len1 left \leftarrow 0 to min(m, k) do
      len1 right \leftarrow m – len1_left;
3.
      len2_left = k - len1_left; len2_right = n - len2_left
      if len1_left = 0 then left_max = A2[len2_left - 1]
5.
      else if len2_left = 0 then left_max = A1[len1_left - 1]
      else left_max = max(A1[len1_left-1],
6.
7.
                             A2[len2 left - 1])
8.
      if len1_right = 0 then right_min = A2[len2_left]
9.
      else if len2_right = 0 then right_min = A1[len1_left]
10.
      else right_min = min(A1[len1_left],
11.
                                 A2[len2_left])
12.
      if left_max <= right_min then</pre>
13.
         return left_max
```

```
Why m \le n:
Ex. m = 10, n = 1, k = 8
len1_left = 5 \Longrightarrow len2_left = 3
                      len1_right
  len1_left
                          8
                               10
 = len1 left
                   = m - len1 left
                      len2_right
  len2_left
2
         5
     3
              6
                            9
                   = n - len2_left
= k - len1 left
```

Time complexity:  $O(\min(m, n, k))$ 

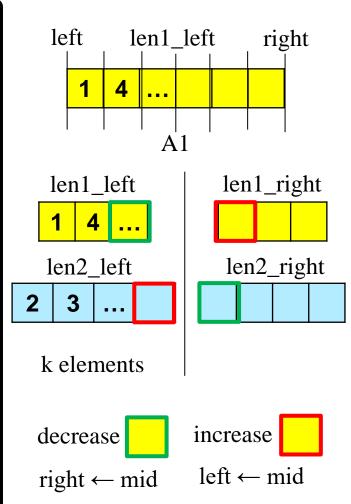
Space complexity: O(1)



# 3. Algorithms

## 3) Decrease and conquer 2

```
kthElement (A1[0..m-1], A2[0..n-1], k)
Input: : Two sorted arrays of size m and n. m \le n.
Output: : The k-th element of the final sorted array.
   left \leftarrow 0; right \leftarrow min(k, m)
   while left + 1 < end do
3.
      len1 left \leftarrow (left + right) / 2; len1 left = m - i
4.
      len2_left = k - len1_left; len2_right = n - len2_left
5.
      if len1_left = 0 then left_max = A2[len2_left - 1]
      else if len2_left = 0 then left_max = A1[len1_left-1]
6.
7.
      else left_max = max(A1[len1_left-1],
8.
                                 A2[len2_left - 1])
      if len1_right = 0 then right_min = A2[len2_left]
9.
10.
      else if len2_right = 0 then right_min = A1[len1_left]
11.
      else right_min = min(A1[len1_left],
12.
                                 A2[len2_left])
13.
      if left_max <= right_min then</pre>
14.
         return left_max
```



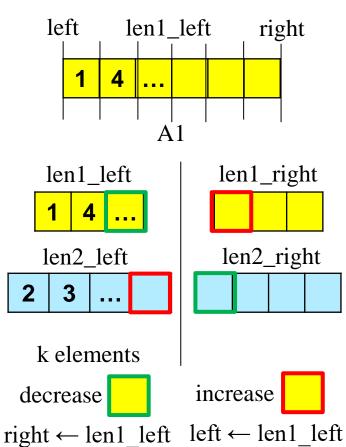
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# 3. Algorithms

### 3) Decrease and conquer 2

```
kthElement (A1[0..m-1], A2[0..n-1], k)
Input: : Two sorted arrays of size m and n. m \le n.
Output: : The k-th element of the final sorted array.
15.
      else
16.
         if len1_left = 0 then left \leftarrow len1_left
         else if len1_right = 0 then right \leftarrow len1_left
17.
18.
         else if A1[len1\_right - 1] > A2[len2\_left] then
            right \leftarrow len1 left
19.
         else if A2[len2_left - 1] > A1[len1_left] then
20.
21.
            left ← len1 left
22. if leftMax <= rightMin then
      return leftMax
23.
24. else
25.
      return rightMin
```

Time complexity:  $O(\log(\min(m, n, k)))$  Space complexity: O(1)



	Time complexity	Space complexity
Merge 2 sorted arrays:	O(k)	O(1)
Decrease and conquer:	$O(\min(m, n, k))$	O(1)
Decrease and conquer 2:	$O(\log(\min(m, n, k)))$	O(1)

m and n are the size of A1 and A2, respectively k is the position of the required element