

## Let's get started

We want you to be successful.

We will work together to build an environment in CSE 105 that supports your learning in a way that respects your perspectives, experiences, and identities (including race, ethnicity, heritage, gender, sex, class, sexuality, religion, ability, age, educational background, etc.). Our goal is for you to engage with interesting and challenging concepts and feel comfortable exploring, asking questions, and thriving.

If you or someone you know is suffering from food and/or housing insecurities there are UCSD resources here to help:

Basic Needs Office: <https://basicneeds.ucsd.edu/>

Triton Food Pantry (in the old Student Center) is free and anonymous, and includes produce:

<https://www.facebook.com/tritonfoodpantry/>

Mutual Aid UCSD: <https://mutualaiducsd.wordpress.com/>

Financial aid resources, the possibility of emergency grant funding, and off-campus housing referral resources are available: see your College Dean of Student Affairs.

If you find yourself in an uncomfortable situation, ask for help. We are committed to upholding University policies regarding nondiscrimination, sexual violence and sexual harassment. Here are some campus contacts that could provide this help: Counseling and Psychological Services (CAPS) at 858 534-3755 or <http://caps.ucsd.edu> OPHD at 858 534-8298 or [ophd@ucsd.edu](mailto:ophd@ucsd.edu) , <http://ophd.ucsd.edu> CARE at Sexual Assault Resource Center at 858 534-5793 or [sarc@ucsd.edu](mailto:sarc@ucsd.edu) , <http://care.ucsd.edu>

Please reach out ([minnes@ucsd.edu](mailto:minnes@ucsd.edu)) if you need support with extenuating circumstances affecting CSE 105.

## Introductions

Class website on Canvas <https://canvas.ucsd.edu>

Instructor: Prof. Mia Minnes "Minnes" rhymes with Guinness, [minnes@ucsd.edu](mailto:minnes@ucsd.edu), <http://cseweb.ucsd.edu/~minnes>

Our team: One instructor + two TAs and eight tutors + all of you

Fill in contact info for students around you, if you'd like:

## CSE 105's Big Questions

- What problems are computers capable of solving?
- What resources are needed to solve a problem?
- Are some problems harder than others?

In this context, a **problem** is defined as: “Making a decision or computing a value based on some input”

Consider the following problems:

- Find a file on your computer
- Determine if your code will compile
- Find a run-time error in your code
- Certify that your system is un-hackable

file system organized,  
name.

Which of these is hardest?

Definition of problem

- what are we starting from?

Approach to solution

- e.g. finding bugs might require more  
than running test case.

\* Running solution

In Computer Science, we operationalize “hardest” as “requires most resources”, where resources might be memory, time, parallelism, randomness, power, etc.

To be able to compare “hardness” of problems, we use a consistent description of problems

**Input:** String

**Output:** Yes/ No, where Yes means that the input string matches the pattern or property described by the problem.

## Weeks 0 and 1 at a glance

### Textbook reading: Chapter 0, Sections 1.3, 1.1

Before Monday, review class syllabus on Canvas (<https://canvas.ucsd.edu/>) and read Example 1.51 and Definition 1.52 (definition of regular expressions) on page 64. *Notice:* we are jumping to Section 1.3 and then will come back to Section 1.1 on Wednesday.

Before Wednesday, read Figure 1.4 and Definition 1.5 (definition of finite automata) on pages 34-35.

Before Friday, read the definition of the union, concatenation, and star operations for languages is given as Definition 1.23 on page 44 and a useful example is Example 1.24.

For Week 2 Monday: Pages 41-43 (Figures 1.18, 1.19, 1.20) (examples of automata and languages).

*Textbook references:* Within a chapter, each item is numbered consecutively. Figure 1.22 is the twenty-second numbered item in chapter one; it comes right after Example 1.21 and right before Definition 1.23.

### We will be learning and practicing to:

- Clearly and unambiguously communicate computational ideas using appropriate formalism. Translate across levels of abstraction.
  - Translate a decision problem to a set of strings coding the problem.
    - \* **Distinguish between alphabet, language, sets, and strings**
  - Use regular expressions and relate them to languages and automata.
    - \* **Write and debug regular expressions using correct syntax**
    - \* **Determine if a given string is in the language described by a regular expression**
  - Use precise notation to formally define the state diagram of finite automata.
  - Use clear English to describe computations of finite automata TM informally.
    - \* **State the formal definition of (deterministic) finite automata**
    - \* **Trace the computation of a finite automaton on a given string using its state diagram**
    - \* **Translate between a state diagram and a formal definition**
    - \* **Determine if a given string is in the language recognized by a finite automaton**

### TODO:

#FinAid Assignment on Canvas (complete as soon as possible) and read syllabus on Canvas

Schedule your Test 1 Attempt 1, Test 2 Attempt 1, Test 1 Attempt 2, and Test 2 Attempt 2 times at PrairieTest (<http://us.prairietest.com>)

Review Quiz on PrairieLearn (<http://us.prairielearn.com>), complete by Sunday 10/6/2024

Create a homework group, possibly by using the Piazza (<https://piazza.com/>) find-a-teammate tool

Homework 1 submitted via Gradescope (<https://www.gradescope.com/>), due Tuesday 10/8/2024

# Week 0 Friday: Terminology and Notation

The CSE 105 vocabulary and notation build on discrete math and introduction to proofs classes. Some of the conventions may be a bit different from what you saw before so we'll draw your attention to them.

For consistency, we will use the notation from this class' textbook<sup>1</sup>.

These definitions are on pages 3, 4, 6, 13, 14, 53.

Term	Typical symbol or Notation	Meaning
Alphabet	<i>Sigma</i> $\Sigma, \Gamma$	A non-empty finite set
Symbol over $\Sigma$	<i>Gamma</i> $\sigma, b, x$	An element of the alphabet $\Sigma$
String over $\Sigma$	$u, v, w$	A finite list of symbols from $\Sigma$
(The) empty string	$\varepsilon$	The (only) string of length 0
The set of all strings over $\Sigma$	$\Sigma^*$	The collection of all possible strings formed from symbols from $\Sigma$
(Some) language over $\Sigma$	$L$	(Some) set of strings over $\Sigma$
(The) empty language	$\emptyset$	The empty set, i.e. the set that has no strings (and no other elements either)
The power set of a set $X$	$\mathcal{P}(X)$	The set of all subsets of $X$
(The set of) natural numbers	$\mathcal{N}$	The set of positive integers
(Some) finite set		The empty set or a set whose distinct elements can be counted by a natural number
(Some) infinite set		A set that is not finite.
Reverse of a string $w$	$w^{\mathcal{R}}$	write $w$ in the opposite order, if $w = w_1 \cdots w_n$ then $w^{\mathcal{R}} = w_n \cdots w_1$ . Note: $\varepsilon^{\mathcal{R}} = \varepsilon$
Concatenating strings $x$ and $y$	$xy$	take $x = x_1 \cdots x_m$ , $y = y_1 \cdots y_n$ and form $xy = x_1 \cdots x_m y_1 \cdots y_n$
String $z$ is a substring of string $w$		there are strings $u, v$ such that $w = uzv$
String $x$ is a prefix of string $y$		there is a string $z$ such that $y = xz$
String $x$ is a proper prefix of string $y$		$x$ is a prefix of $y$ and $x \neq y$
Shortlex order, also known as string order over alphabet $\Sigma$		Order strings over $\Sigma$ first by length and then according to the dictionary order, assuming symbols in $\Sigma$ have an ordering

<sup>1</sup>Page references are to the 3rd edition of Sipser's Introduction to the Theory of Computation, available through various sources for approximately \$30. You may be able to opt in to purchase a digital copy through Canvas. Copies of the book are also available for those who can't access the book to borrow from the course instructor, while supplies last (minnes@ucsd.edu)

Write out in words the meaning of the symbols below:

$$\{a, b, c\}$$

The set whose three elements are a and b and c

$$|\{a, b, a\}| = 2$$

There are (exactly) 2 distinct elements in the set  $\{a, b, a\}$ .

$$|aba| = 3$$

The string aba has length 3.

Circle the correct choice:

A **string** over an alphabet  $\Sigma$  is an element of  $\Sigma^*$  ~~OR a subset of  $\Sigma^*$~~ .

A **language** over an alphabet  $\Sigma$  is ~~an element of  $\Sigma^*$~~  ~~OR~~ a subset of  $\Sigma^*$ .

With  $\Sigma_1 = \{0, 1\}$  and  $\Sigma_2 = \{a, b, c, d, e, f, g, h, i, j, k, l, m, n, o, p, q, r, s, t, u, v, w, x, y, z\}$  and  $\Gamma = \{0, 1, x, y, z\}$

~~True~~ or False:  $\varepsilon \in \Sigma_1$

True or ~~False~~:  $\varepsilon$  is a string over  $\Sigma_1$

~~True~~ or False:  $\varepsilon$  is a language over  $\Sigma_1$

True or ~~False~~:  $\varepsilon$  is a prefix of some string over  $\Sigma_1$

~~True~~ or False: There is a string over  $\Sigma_1$  that is a proper prefix of  $\varepsilon$

The first five strings over  $\Sigma_1$  in string order, using the ordering  $0 < 1$ :

$\varepsilon, 0, 1, 00, 01$

The first five strings over  $\Sigma_2$  in string order, using the usual alphabetical ordering for single letters:

$\varepsilon, a, b, c, d$

# Week 1 Monday: Regular expressions

Our motivation in studying sets of strings is that they can be used to encode problems. To calibrate how difficult a problem is to solve, we describe how complicated the set of strings that encodes it is. How do we define sets of strings?

Set notation - roster  $\{-, -, -, -\}$  finite sets only

Set builder notation  $\{x \mid P(x)\}$  or  $\{f(x) \mid \dots\}$

Recursive definition  
- Basis step  
- Recursive step  
The collection of all values for which the property P(.) holds

Named sets Greek alphabet,  $\mathbb{N}$ ,  $\emptyset$ ,  $\mathbb{R}$

Subsets of these might be defined with set operations

How would you describe the language that has no elements at all?

↳ set of strings over a given alphabet

$\emptyset$

$\{\}$

$\{x \in \Sigma^* \mid x \neq x\}$

set of  
all strings  
over alphabet  $\Sigma$

How would you describe the language that has all strings over  $\{0,1\}$  as its elements?

Not  
roster  
definitions  
because infinite

Set operation: Kleene star

for any set  $A$   $A^* = \{w_1 \dots w_k \mid k \geq 0, w_i \in A\}$

Set of strings over  $\{0,1\}$   
is  $\{0,1\}^*$

$\{x \in \Sigma^* \mid x = x\}$

Recursively: (Basis step)  $\epsilon \in \{0,1\}^*$   
(Recursive step) For  $w \in \{0,1\}^*$   $w0 \in \{0,1\}^*$   
and  $w1 \in \{0,1\}^*$

**Definition 1.52:** A **regular expression** over alphabet  $\Sigma$  is a syntactic expression that can describe a language over  $\Sigma$ . The collection of all regular expressions over  $\Sigma$  is defined recursively:

*Basis steps of recursive definition*

- $a$  is a regular expression, for  $a \in \Sigma$
- $\varepsilon$  is a regular expression
- $\emptyset$  is a regular expression

*Recursive steps of recursive definition*

- $(R_1 \cup R_2)$  is a regular expression when  $R_1, R_2$  are regular expressions
- $(R_1 \circ R_2)$  is a regular expression when  $R_1, R_2$  are regular expressions
- $(R_1^*)$  is a regular expression when  $R_1$  is a regular expression

The *semantics* (or meaning) of the syntactic regular expression is the **language described by the regular expression**. The function that assigns a language to a regular expression over  $\Sigma$  is defined recursively, using familiar set operations:

*Basis steps of recursive definition*

- The language described by  $a$ , for  $a \in \Sigma$ , is  $\{a\}$  and we write  $L(a) = \{a\}$
- The language described by  $\varepsilon$  is  $\{\varepsilon\}$  and we write  $L(\varepsilon) = \{\varepsilon\}$
- The language described by  $\emptyset$  is  $\{\}$  and we write  $L(\emptyset) = \emptyset$ .

Handwritten notes illustrating the semantics of the basis steps:

- For  $\{ \emptyset \}$ : "Not a language", "the set that contains the empty set".
- For  $\{\varepsilon\}$ : "Language", "The set that contains the empty string".
- For  $\emptyset$ : "Language", "The set that has no elements".

*Recursive steps of recursive definition*

When  $R_1, R_2$  are regular expressions, the language described by the regular expression  $(R_1 \cup R_2)$  is the union of the languages described by  $R_1$  and  $R_2$ , and we write

$$L(R_1 \cup R_2) = L(R_1) \cup L(R_2) = \{w \mid w \in L(R_1) \vee w \in L(R_2)\}$$

When  $R_1, R_2$  are regular expressions, the language described by the regular expression  $(R_1 \circ R_2)$  is the concatenation of the languages described by  $R_1$  and  $R_2$ , and we write

$$L(R_1 \circ R_2) = L(R_1) \circ L(R_2) = \{uv \mid u \in L(R_1) \wedge v \in L(R_2)\}$$

When  $R_1$  is a regular expression, the language described by the regular expression  $(R_1^*)$  is the **Kleene star** of the language described by  $R_1$  and we write

$$L(R_1^*) = (L(R_1))^* = \{w_1 \cdots w_k \mid k \geq 0 \text{ and each } w_i \in L(R_1)\}$$

For the following examples assume the alphabet is  $\Sigma_1 = \{0, 1\}$ :

The language described by the regular expression 0 is  $L(0) = \{0\}$

The language described by the regular expression 1 is  $L(1) = \{1\}$

The language described by the regular expression  $\varepsilon$  is  $L(\varepsilon) = \{\varepsilon\}$

The language described by the regular expression  $\emptyset$  is  $L(\emptyset) = \emptyset$

The language described by the regular expression  $(\Sigma_1 \Sigma_1 \Sigma_1)^*$  is  $L((\Sigma_1 \Sigma_1 \Sigma_1)^*) =$

$$\Sigma_1 \text{ is the regular expression } (0 \cup 1)$$

symbol for concatenation

$$L((\underline{(0 \cup 1)} \circ \underline{(0 \cup 1)} \circ \underline{(0 \cup 1)})^*)$$

$$= \left( L((0 \cup 1) \circ (0 \cup 1) \circ (0 \cup 1)) \right)^*$$

$$= \{ \text{length 3 strings} \}^*$$

$$= \{ \underbrace{w_1 \dots w_k}_{k \text{ many slots}} \mid w_i \text{ is a length 3 string} \}$$

The language described by the regular expression  $1^* \circ 1$  is  $L(1^* \circ 1) =$

(Added after class)

$$L(\underline{1^*} \circ \underline{1}) = L(\underline{1^*}) \circ L(\underline{1})$$

symbol      symbol      set operation

$$= L(\underline{1})^* \circ L(\underline{1})$$

set operation

$$= \underline{\{1\}^*} \circ \underline{\{1\}}$$

$$= \underline{\{1^i \mid i \geq 0\}} \circ \underline{\{1\}}$$

$$= \underline{\{1^{i+1} \mid i \geq 0\}}$$

= The set of nonempty strings of all 1s.

$$L((0 \cup 1)) =$$

$$L(0) \cup L(1) =$$

$$\{0\} \cup \{1\} = \{0, 1\}$$

$$\{0, 1\} \cdot \{0, 1\} \cdot \{0, 1\}$$

$$= \{ \underline{u} \underline{v} \underline{w} \mid u \in \{0, 1\} \text{ and } v \in \{0, 1\} \text{ and } w \in \{0, 1\} \}$$

/// : regular expressions

/// sets

by basis step of semantics of regular expressions



## Week 1 Wednesday: Regular expressions conventions

**Review:** Determine whether each statement below about regular expressions over the alphabet  $\{a, b, c\}$  is true or false:

True or False:  $ab \in L((a \cup b)^*)$

True or False:  $ba \in L(a^*b^*)$

True or False:  $\varepsilon \in L(a \cup b \cup c)$

True or False:  $\varepsilon \in L((a \cup b)^*)$

True or False:  $\varepsilon \in L(aa^* \cup bb^*)$

*Shorthand and conventions* (Sipser pages 63-65)

Assuming  $\Sigma$  is the alphabet, we use the following conventions

$\Sigma$	regular expression describing language consisting of all strings of length 1 over $\Sigma$
$*$ then $\circ$ then $\cup$	precedence order, unless parentheses are used to change it
$R_1R_2$	shorthand for $R_1 \circ R_2$ (concatenation symbol is implicit)
$R^+$	shorthand for $R^* \circ R$
$R^k$	shorthand for $R$ concatenated with itself $k$ times, where $k$ is a (specific) natural number

**Caution:** many programming languages that support regular expressions build in functionality that is more powerful than the “pure” definition of regular expressions given here.

Regular expressions are everywhere (once you start looking for them).

Software tools and languages often have built-in support for regular expressions to describe **patterns** that we want to match (e.g. Excel/ Sheets, grep, Perl, python, Java, Ruby).

Under the hood, the first phase of **compilers** is to transform the strings we write in code to tokens (keywords, operators, identifiers, literals). Compilers use regular expressions to describe the sets of strings that can be used for each token type.

Next time: we’ll start to see how to build machines that decide whether strings match the pattern described by a regular expression.

Practice with the regular expressions over  $\{a, b\}$  below.

For example: Which regular expression(s) below describe a language that includes the string  $a$  as an element?

$$a^*b^*$$

$$a(ba)^*b$$

$$a^* \cup b^*$$

$$(aaa)^*$$

$$(\varepsilon \cup a)b$$

## Week 1 Friday: Finite automata

**\*\*This definition was in the pre-class reading\*\*** A finite automaton (FA) is specified by  $M = (Q, \Sigma, \delta, q_0, F)$ . This 5-tuple is called the **formal definition** of the FA. The FA can also be represented by its state diagram: with nodes for the state, labelled edges specifying the transition function, and decorations on nodes denoting the start and accept states.

Finite set of states  $Q$  can be labelled by any collection of distinct names. Often we use default state labels  $q_0, q_1, \dots$

The alphabet  $\Sigma$  determines the possible inputs to the automaton. Each input to the automaton is a string over  $\Sigma$ , and the automaton “processes” the input one symbol (or character) at a time.

The transition function  $\delta$  gives the next state of the automaton based on the current state of the machine and on the next input symbol.

The start state  $q_0$  is an element of  $Q$ . Each computation of the machine starts at the start state.

The accept (final) states  $F$  form a subset of the states of the automaton,  $F \subseteq Q$ . These states are used to flag if the machine accepts or rejects an input string.

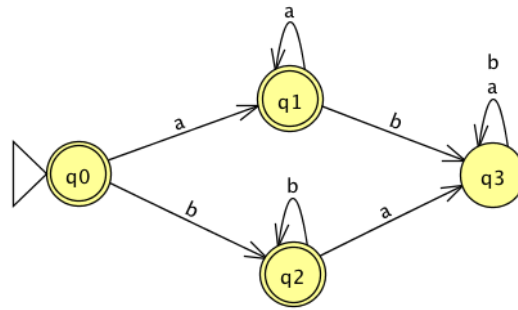
The computation of a machine on an input string is a sequence of states in the machine, starting with the start state, determined by transitions of the machine as it reads successive input symbols.

The finite automaton  $M$  accepts the given input string exactly when the computation of  $M$  on the input string ends in an accept state.  $M$  rejects the given input string exactly when the computation of  $M$  on the input string ends in a nonaccept state, that is, a state that is not in  $F$ .

The language of  $M$ ,  $L(M)$ , is defined as the set of all strings that are each accepted by the machine  $M$ . Each string that is rejected by  $M$  is not in  $L(M)$ . The language of  $M$  is also called the language recognized by  $M$ .

What is **finite** about all finite automata? (Select all that apply)

- ☐ The size of the machine (number of states, number of arrows)
- ☐ The length of each computation of the machine
- ☐ The number of strings that are accepted by the machine

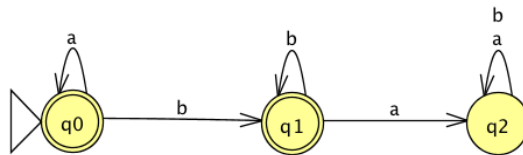


The formal definition of this FA is

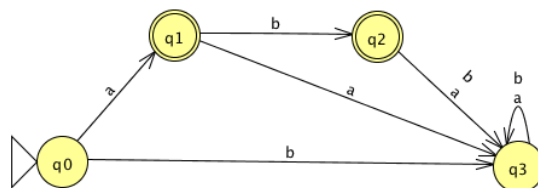
Classify each string  $a$ ,  $aa$ ,  $ab$ ,  $ba$ ,  $bb$ ,  $\varepsilon$  as accepted by the FA or rejected by the FA.

*Why are these the only two options?*

The language recognized by this automaton is



The language recognized by this automaton is



The language recognized by this automaton is