

# Quicksort

$$\text{QUICKSORT}(A, p, r)$$

```

1   if  $p < r$ 
2       then  $q \leftarrow \text{PARTITION}(A, p, r)$ 
3        $\text{QUICKSORT}(A, p, q - 1)$ 
4        $\text{QUICKSORT}(A, q + 1, r)$ 

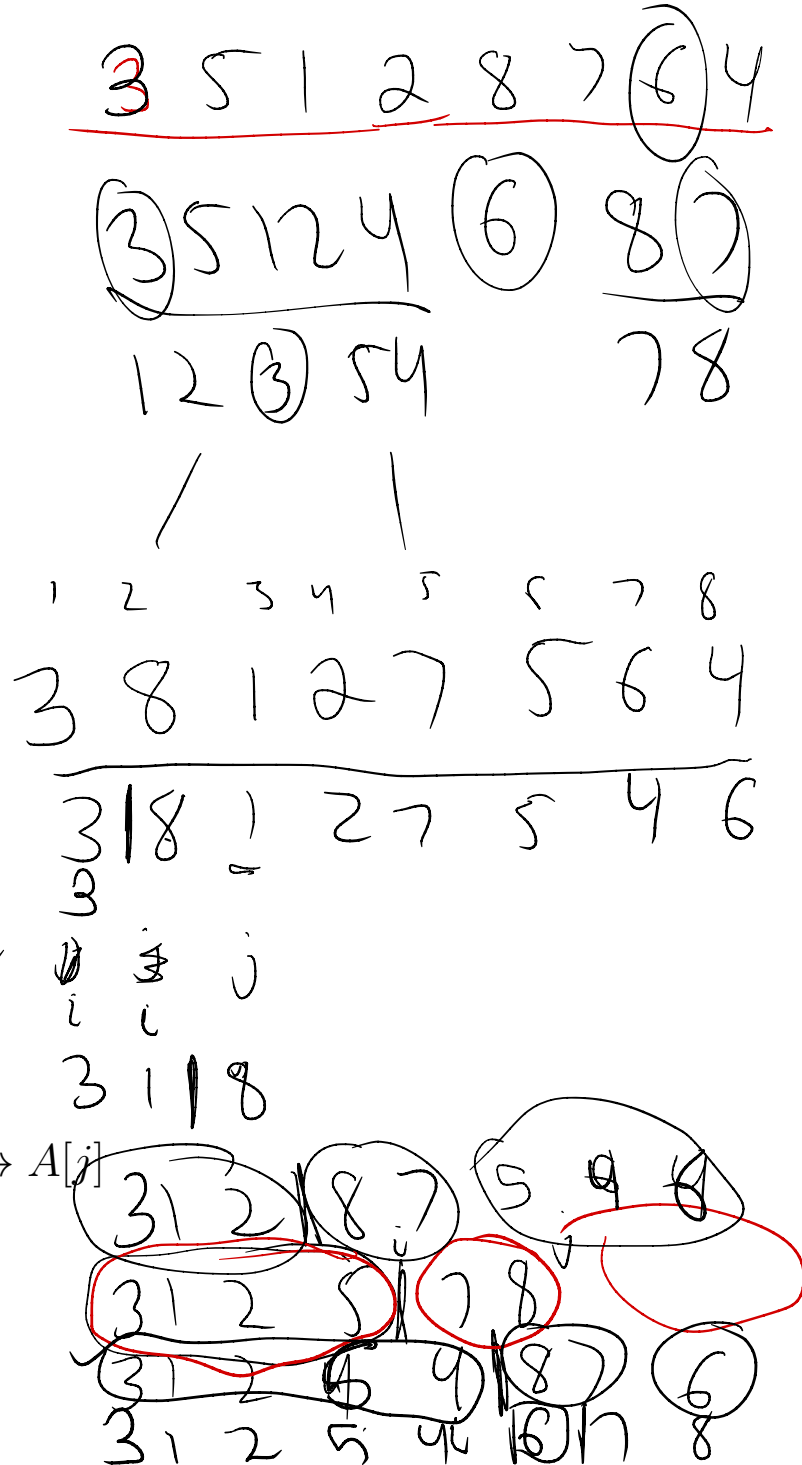
```

$$\text{PARTITION}(A, p, r)$$

```

1   $y \leftarrow \text{RANDOM}(p, r)$ 
2  Exchange  $A[y]$  and  $A[r]$ 
3   $x \leftarrow A[r]$ 
4   $i \leftarrow p - 1$ 
5  for  $j \leftarrow p$  to  $r - 1$ 
6      do if  $A[j] \leq x$ 
7          then  $i \leftarrow i + 1$ 
8              exchange  $A[i] \leftrightarrow A[j]$ 
9  exchange  $A[i + 1] \leftrightarrow A[r]$ 
10 return  $i + 1$ 

```



# Partition Loop Invariant

PARTITION( $A, p, r$ )

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9  exchange  $A[i + 1] \leftrightarrow A[r]$ 
10 return  $i + 1$ 
```

**Loop Invariant** At the beginning of each iteration of the for loop in Partition

1.  $A[p \dots i] \leq x$
2.  $A[i + 1 \dots j - 1] > x$
3.  $A[j \dots r - 1]$  is unexamined
4.  $A[r] = x$

# Quicksort Analysis

- $T(n)$  is the expected running time of quicksort
- Partition takes  $O(n)$  time.
- If partition is  $x$  th smallest, then

$$T(n) = T(x) + T(n - x) + O(n)$$

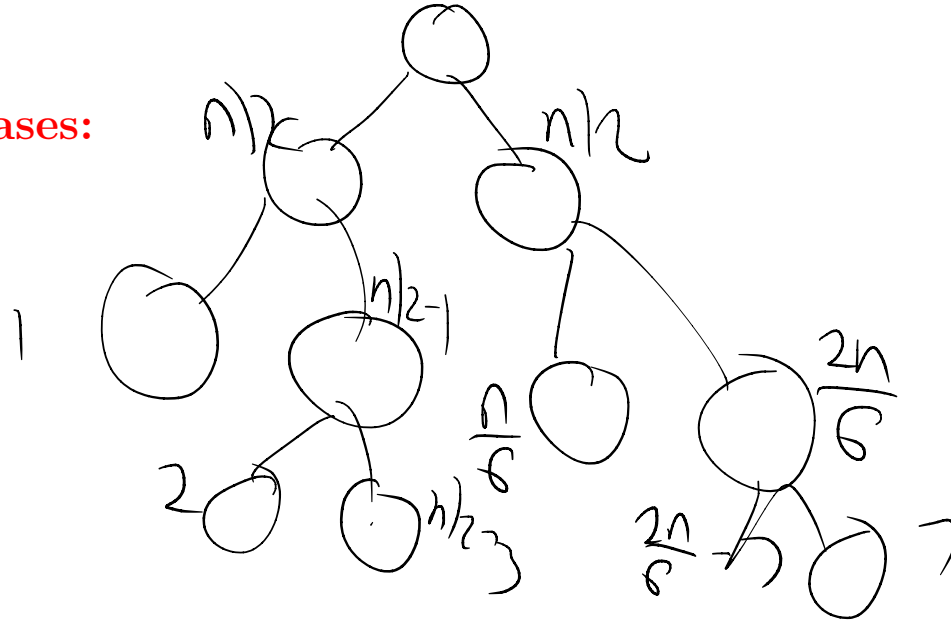
## Quicksort Analysis

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For intuition consider cases:

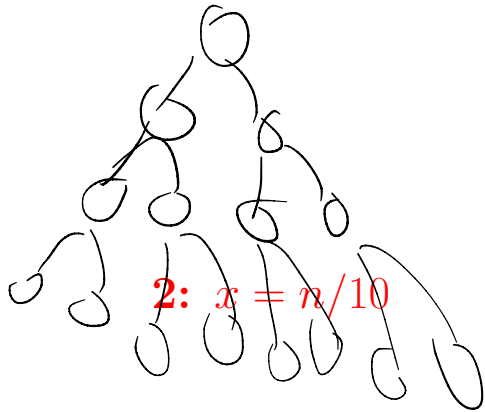
- $x = n/2$
- $x = n/10$
- $x = 1$



# Cases

$$T(n) = T(x) + T(n - x) + O(n)$$

**1:**  $x = n/2$



$$\begin{aligned} T(n) &= T(n/2) + T(n/2) + O(n) \\ &= 2T(n/2) + O(n) \\ &= O(n \log n) \end{aligned}$$

$$\begin{aligned} T(n) &= T(n/10) + T(9n/10) + O(n) \\ &= O(n \log n) \end{aligned}$$

**3:**  $x = 1$

$$\begin{aligned} T(n) &= T(1) + T(n - 1) + O(n) \\ &= T(n - 1) + O(n) \\ &= O(n^2) \end{aligned}$$

**What might this make us guess the answer is?**

## Following the Selection Analysis

$$T(n) = \sum_{i=1}^n \frac{1}{n} (T(i) + T(n-i) + O(n))$$

could continue as in Selection.

# Alternative Analysis

- We will count comparisons of data elements.
- Claim 1: The running time is dominated by comparison of data items.
- Claim 2: All comparisons are in line 6 of PARTITION, and compare some item  $A[j]$  to the pivot element.
- Claim 3: Once an element is chosen as a pivot, it is never compared to any other element again.
- Claim 4: Each pair of elements is compared to each other at most once.

## Analysis:

- Let the data be renamed  $Z_1, \dots, Z_n$  in sorted order.
- Use  $Z_{ij}$  to denote  $Z_i, Z_{i+1}, \dots, Z_j$
- Let  $X_{ij}$  be the indicator random variable for the comparison of  $Z_i$  to  $Z_j$ .
- Let  $X$  be the random variable counting the number of comparisons.
- By claim 4, we have

$$X = \sum_{i=1}^n \sum_{j=i+1}^n X_{ij}$$

# Analysis

$$X = \sum_{i=1}^n \sum_{j=i+1}^n X_{ij}$$

Taking expectations

$$\begin{aligned} E[X] &= E\left[\sum_{i=1}^n \sum_{j=i+1}^n X_{ij}\right] \\ &= \sum_{i=1}^n \sum_{j=i+1}^n E[X_{ij}] \quad \text{linearity of expectation} \\ &= \sum_{i=1}^n \sum_{j=i+1}^n \Pr(Z_i \text{ is compared to } Z_j) \end{aligned}$$



## What is the probability that $Z_i$ is compared to $Z_j$

When is  $Z_i$  compared to  $Z_j$  ?

- When either  $Z_i$  or  $Z_j$  is chosen as a pivot, and the other one is still in the same recursive problem.

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When is  $Z_i$  compared to  $Z_j$  ?

- When either  $Z_i$  or  $Z_j$  is chosen as a pivot, and the other one is still in the same recursive problem.
- Equivalently, when either  $Z_i$  or  $Z_j$  is the first element from  $Z_{ij}$  to be chosen as a pivot.
- What is the probability that  $Z_i$  or  $Z_j$  is the first element from  $Z_{ij}$  to be chosen as a pivot?

12 11 3 4 5 6 11 7 8 9

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- What is the probability that  $Z_i$  or  $Z_j$  is the first element from  $Z_{ij}$  to be chosen as a pivot?
  - $Z_{ij}$  has  $j - i + 1$  elements.
  - pivots are always chosen uniformly at random
  - $Pr(Z_i \text{ is compared to } Z_j) = 2/(j - i + 1)$

3 5 1 2 9 8 7 6 4

2 8-2+1

$\frac{2}{5}$

## Finishing analysis

$$\begin{aligned} E[X] &= E\left[\sum_{i=1}^n \sum_{j=i+1}^n X_{ij}\right] \\ &= \sum_{i=1}^n \sum_{j=i+1}^n E[X_{ij}] \quad \text{linearity of expectation} \\ &= \sum_{i=1}^n \sum_{j=i+1}^n \Pr(Z_i \text{ is compared to } Z_j) \\ &= \sum_{i=1}^n \sum_{j=i+1}^n \frac{2}{j-i+1} \end{aligned}$$

make the transformation of variables  $k = j - i + 1$

$$\begin{aligned} &= \sum_{i=1}^n \sum_{j=i+1}^n \frac{2}{j-i+1} \\ &= \sum_{i=1}^n \sum_{k=2}^{n-i+1} \frac{2}{k} \\ &\leq 2 \sum_{i=1}^n \ln(n-i+1) \\ &\leq 2 \sum_{i=1}^n \ln(n) \\ &= O(n \log n) \end{aligned}$$