

1.5

a. Instruction per second = instruction count / CPU time = instruction count / (instruction count * CPI / clock rate) = clock rate / CPI

Thus for P1, $3 \times 10^9 / 1.5 = 2 \times 10^9$ instructions/second

For P2, $2.5 \times 10^9 / 1.0 = 2.5 \times 10^9$ instructions/second

P3, $4 \times 10^9 / 2.2 = 1.82 \times 10^9$ instructions/second

b.

clock cycles = CPU time * clock rate

Thus for P1, $10 \times 3 \times 10^9 = 3 \times 10^{10}$ cycles

For P2, $10 \times 2.5 \times 10^9 = 2.5 \times 10^{10}$ cycles

For P3, $10 \times 4 \times 10^9 = 4 \times 10^{10}$ cycles

Instruction count = clock cycles / CPI

Thus for P1, $10 \times 3 \times 10^9 / 1.5 = 2 \times 10^{10}$ instructions

For P2, $10 \times 2.5 \times 10^9 / 1.0 = 2.5 \times 10^{10}$ instructions

For P3, $10 \times 4 \times 10^9 / 2.2 = 1.82 \times 10^{10}$ instructions

c.

CPU time = CPI * instruction count / clock rate

$0.7 \times \text{CPU time} = 1.2 \times \text{CPI} \times \text{instruction count} / \text{expected clock rate}$

Thus expected clock rate = $(1.2 / 0.7) \times \text{clock rate}$

Thus for P1, $(1.2 / 0.7) \times 3 = 5.14$ GHz

For P2, $(1.2 / 0.7) \times 2.5 = 4.29$ GHz

For P3, $(1.2 / 0.7) \times 4 = 6.86$ GHz

1.6

a.

given the program, for P1, global CPI = $1 \times 0.1 + 2 \times 0.2 + 3 \times 0.5 + 3 \times 0.2 = 2.6$ cycle/instruction

for P2, global CPI = $2 \times 0.1 + 2 \times 0.2 + 2 \times 0.5 + 2 \times 0.2 = 2.0$ cycle/instruction

b.

clock cycles = CPI * instruction count

For P1, $2.6 \times 10^6 = 2.6 \times 10^6$ clock cycles

For P2, $2.0 \times 10^6 = 2.0 \times 10^6$ clock cycles

To say which one is faster, we know that: CPU time = clock cycles / clock rate

Thus for P1, CPU time = $2.6 \times 10^6 / (2.5 \times 10^9) = 1.04 \times 10^{-3}$ seconds

for P2, CPU time = $2.0 \times 10^6 / (3 \times 10^9) = 6.67 \times 10^{-4}$ seconds

So in terms of CPU time, P2 is faster.

1.7

a.

average CPI = execution time / (clock cycle time * instruction count)

For A, average CPI= $1.1 / (10^{(-9)} * 10^9) = 1.1$ cycle/second

For B, average CPI= $1.5 / (10^{(-9)} * 1.2 * 10^9) = 1.25$ cycle/second

b.

Here we assume CPI is not affected by the new processor, otherwise we won't be able to solve this problem.

Since execution time= CPI * instruction count* clock cycle time

$1.1 * 10^9 * \text{clock time A} = 1.25 * 1.2 * 10^9 * \text{clock time B}$

So clock time A/ clock time B= $1.25 * 1.2 / 1.1 = 1.36$

Since clock rate= $1 / \text{clock time}$, clock rate A/ clock rate B = $1.1 / 1.5 = 0.73$

So the clock rate of compiler A is 0.73 times that of compiler B.

c.

Here we measure speed up with CPU time. CPU time= instruction count * CPI* clock cycle time.

And we assume that clock cycle time is the same for all compilers.

CPU performance new / CPU performance A = CPU time A/ CPU time new=
 $(1.1 * 10^9) / (1.1 * 6 * 10^8) = 1.67$

CPU performance new / CPU performance B = CPU time B/ CPU time new=
 $(1.25 * 1.2 * 10^9) / (1.1 * 6 * 10^8) = 2.27$

Thus, speed up of new compiler versus A is 1.67, speed up of new compiler versus B is 2.27.

1.13

1.13.1

$0.8 * 70 + 85 + 40 + 55 = 236$ seconds $1 - 236 / 250 = 5.6\%$

So total time is reduced by 5.6%.

1.13.2

Here we assume INT operations time is the unmentioned time, which is 55 seconds, otherwise we won't be able to solve this problem. And we assume all times except INT operation time remain the same.

$0.2 * 250 = 50$ seconds $50 / 55 = 90.9\%$

So INT operation time is reduced by 90.9% if total time is reduced by 20%.

1.13.3

$250 * 0.2 = 50$ seconds > 40 seconds

So the expected reduce time is greater than the total time of current branch instructions. Thus this can't be done.