



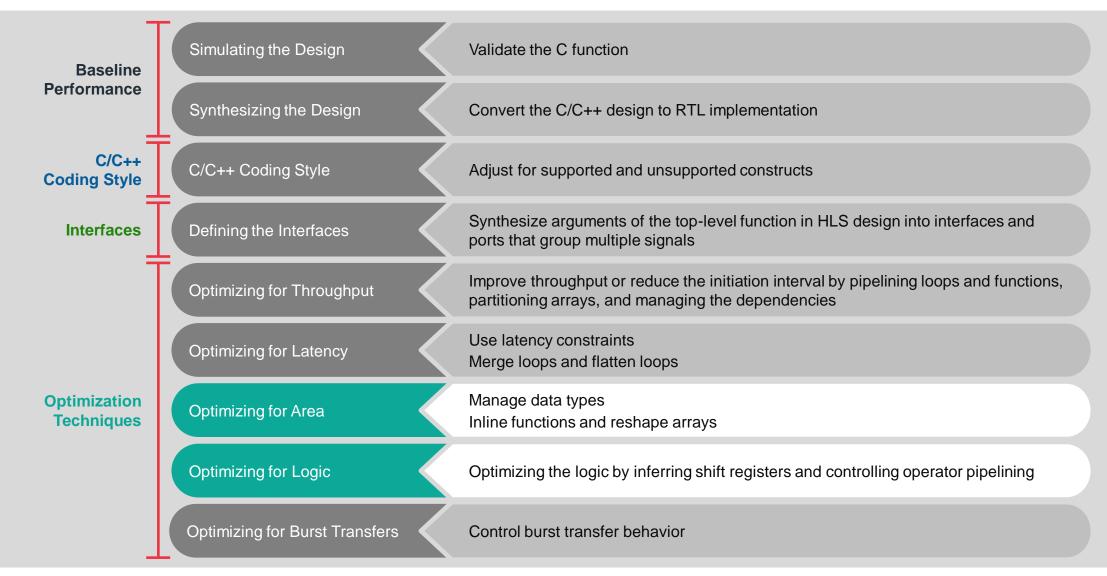
# Improving Area and Resources

### **Objectives**

- After completing this module, you will be able to:
  - Describe different methods for improving resource utilization
  - Control the structure of the design by using directives to improve the area
  - Explain how arbitrary precision types can result in optimal resource usage



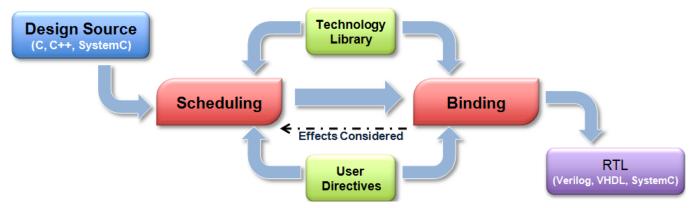
# Vitis HLS Design Methodology





### Review: Control Scheduling & Binding

- Scheduling & Binding
  - Scheduling and Binding are the processes at the heart of HLS



- BIND\_OP/ BIND\_STORAGE
  - Can be used to specify a device resource for implementation
  - Can be used to assign a specific memory type in the RTL
- ▶ The allocation directive
  - Can be used to limit the number of operation in scheduling & binding stages



# **Directives to improve Area**

<b>Directives and Configurations</b>	Description
ALLOCATION	Specify a limit for the number if operations, cores, or functions used. This can force the sharing of hardware resources and might increase latency
ARRAY_RESHAPE	Reshapes an array from one with many elements to one with greater word-width. Useful for improving block RAM accesses without using more block RAM
LOOP_MERGE	Merges consecutive loops to reduce overall latency, increase sharing, and improve logic optimization
OCCURRENCE	Used when pipelining functions or loops to specify that the code in a location is executed at a lesser rate than the code in the enclosing function or loop
STRAM	Specifies that a specific memory channel is to be implemented as a FIFO or RAM during dataflow optimization
Config Dataflow	This configuration specifies the default memory channel and FIFP depth in a dataflow region
BIND_STORAGE	Define a specific implementation for a storage element, or memory, in the RTL
BIND_OP	Define a specific implementation for an operation in the RTL



# Reducing Area Usage



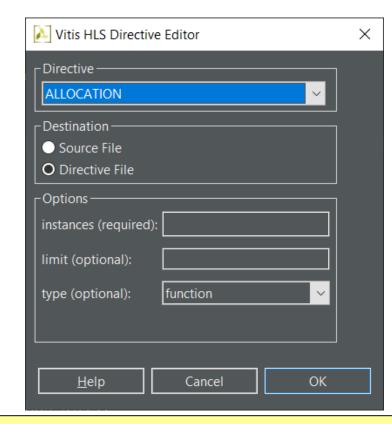
### Improving Area/Resource Utilization

- Control the number of elements
  - Directives can be used to control scheduling and binding phases, which in turn control the number of elements used
- Control the design hierarchy
  - Like RTL synthesis, removing the hierarchy can help optimize across function and loop boundaries
    - Functions can be inlined
    - Loops can be unrolled
- Array implementation
  - VITIS HLS provides directives for combining memories
    - Allowing a single large memory to be used instead of multiple smaller memories
- Bit-width optimization
  - Bit width impact the size of the storage elements and operators
  - Arbitrary precision types ensure correct operator sizing



#### **Allocation: Limit the Numbers**

- Allocation directive limits different types
  - Type: Operation
    - The instances are the operators
    - Add, mul, urem, etc.
  - Type: Core
    - The instances are the cores
    - Adder, Addsub, PipeMult2s, etc
  - Type: Functions
    - The functions in the code
    - Discussed in more detail later
- Allocations are defined for a scope
  - Like all directives, allocations are set for the scope they are applied in
  - If the directive is applied to a function, loop or region, it does not include objects outside that scope

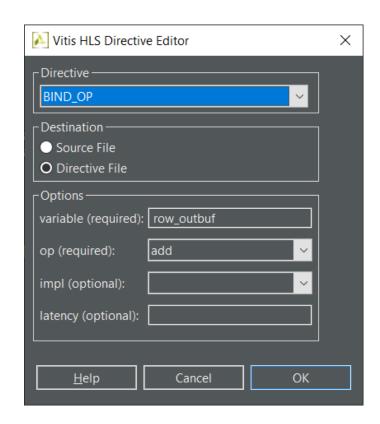


Operators and Cores are listed in the VITIS HLS Library Guide



#### BIND\_OP: Specific a device resource

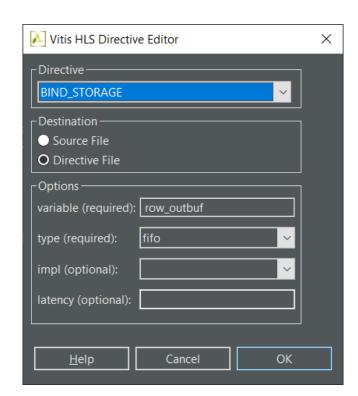
- Specifies that for a given variable, an operation(mul, add,sub) should be mapped to a device resource for implementation
  - If not, Vitis HLS tool automatically determines the resource to use
- You can also specify:
  - op: Operation to bind to a specific implementation resource
  - -impl: Implementation to use for the specified operation
  - -latency: Latency for the operation





## BIND\_STORAGE: Assign a specific memory

- Assigns a variable(array or function argument) to a specific memory type in the RTL
  - If not, Vitis HLS tool determines the memory type to assign
  - HLS implements the memory using specified implementations(impl) in the hardware
  - Controls the array implementation as a single/dual-port RAM
- You can also specify:
  - -type: Type of memory to bind to the specified variable
    - Support types: fifo, ram\_1p, ram\_1wnr, ram\_2p, ram\_s2p, ram\_t2p, rom\_1p, rom\_2p, and rom\_np
  - -impl: Implementation for the specified memory type
    - Support impementations: bram, bram\_ecc, lutram, uram, uram\_ecc and srl
  - -latency: Default latency for the binding of the storage type to the implementation

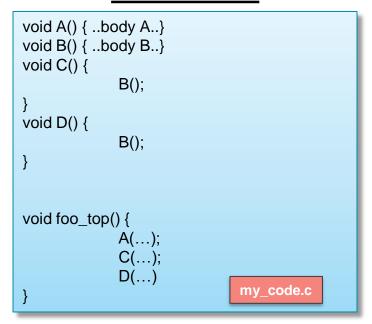




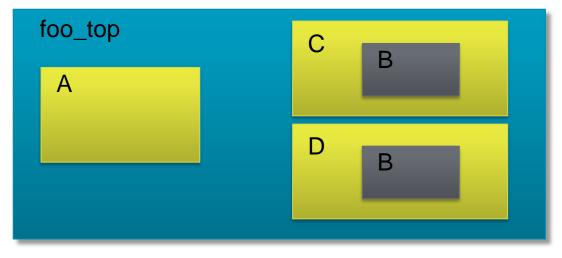
## **Review: Functions & RTL Hierarchy**

Each function in the design gets synthesized into an RTL block, maintaining the design hierarchy

#### **Source Code**



#### **RTL** hierarchy

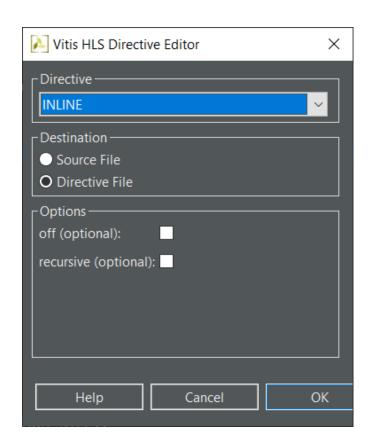


Functions can be inlined – the hierarchy removed & the function dissolved into the surrounding function



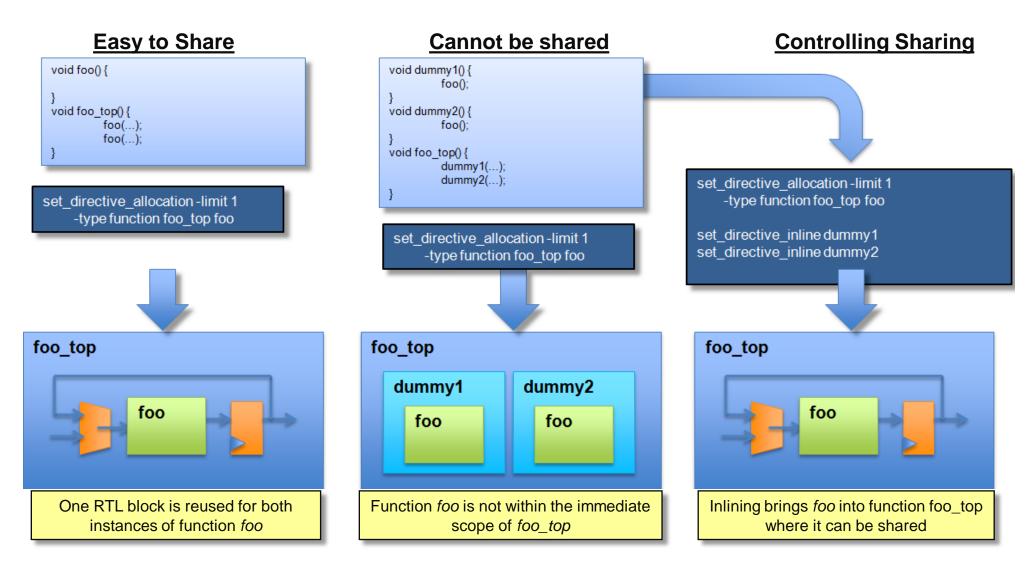
### **Function Inlining**

- Inlining a function may improve area by allowing the components within the function to be better shared or optimized with the logic in the calling function.
- Vitis HLS performs some inlining automatically
  - This is performed on small logic functions if Vitis HLS determines area or performance will benefit
- User Control
  - Optionally recursively down the hierarchy
  - Optionally inlining can be explicitly prevented
    - Turn inlining off
- Inlining functions allows for greater optimization
  - Like ungrouping RTL hierarchies: optimization across boundaries
  - Like ungrouping RTL hierarchies it can result in lots of operations & impact run time





### Inline and Allocation: Shape the Hierarchy





### Loops

- ▶ By default, loops are rolled and pipelined to improve performance
  - Each C loop iteration → Implemented in the same state with same resources

```
void foo_top (...) {
    ...
Add: for (i=3;i>=0;i--) {
        b = a[i] + b;
    ...
}
Synthesis
```

For Area optimization

Keeping loops rolled maximizes sharing across loop iterations: each <u>iteration</u> of the loop uses the same hardware resources



## **Loop Merging**

- Loop merging can remove the redundant computation among multiple (related) loops
  - Improving area and performance

```
for (i = 0; i < N; ++i) {
    A[i] = B[i] + 1;
    C[i] = A[i] / 2;
}
```

Effective code after compiler transformation

- Allows the logic within the loops to be optimized together
- Optimization cannot occur across loop boundaries
- Allows for more efficient architecture explorations

```
for (i = 0; i < N; ++i)

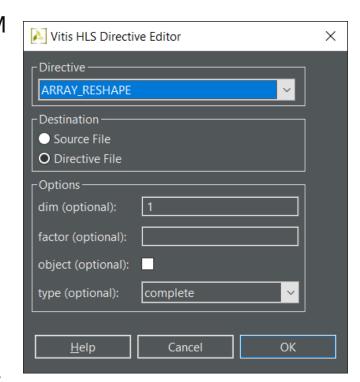
C[i] = (B[i] + 1) / 2;
```

Removes A[i], any address logic and any potential memory accesses



### **Mapping Arrays**

- ▶ The arrays in the C model may not be ideal for the available RAMs
  - The code may have many small arrays, each array is mapped into a block RAM or UltraRAM
  - Basic block RAM unit is 18K, the array may not utilize the RAMs very well
- Array Mapping
  - Mapping combines smaller arrays into larger arrays
    - Reduce the number or block RAMs required
  - Specify the array variable to be mapped
  - Give all arrays to be combined the same instance name
  - If an array is larger than 18K, they are automatically mapped into multiple 18K units.
- ARRAY\_RESHAPE directive supports mapping small arrays into a larger one





### **Arbitrary Precision Integers**

- ▶ C and C++ have standard types created on the 8-bit boundary
  - char (8-bit), short (16-bit), int (32-bit), long long (64-bit)
    - Also provides stdint.h (for C), and stdint.h and cstdint (for C++)
    - Types: int8\_t, uint16\_t, uint32\_t, int\_64\_t etc.
  - Implemented hardware sometimes need different sizes of data types
  - They result in hardware which is not bit-accurate and can give sub-standard QoR
  - Sigals with arbitrary widths are also converted into fixed sized widths when native C data types are used
- Vitis HLS provides both integer and fixed-point arbitrary precision data types for C++
  - Allow any arbitrary bit-width to be specified
  - Will simulate with bit-accuracy

```
#include ap cint.h
                                        my_code.c
void foo_top (...) {
 int1
                 var1:
                               // 1-bit
  uint1
                 var1u;
                               // 1-bit unsigned
 int2
                 var2;
                               // 2-bit
 int1024
                 var1024;
                             // 1024-bit
 uint1024
                             // 1024-bit unsigned
                 var1024;
```

```
#include ap_int.h
                                     my_code.cpp
void foo_top (...) {
  ap int<1>
                         var1:
                                          // 1-bit
  ap_uint<1>
                         var1u;
                                          // 1-bit unsigned
  ap_int<2>
                         var2;
                                          // 2-bit
                                          // 1024-bit
  ap_int<1024>
                         var1024:
                                          // 1024-bit unsigned
  ap_int<1024>
                         var1024u;
```



## Why are Arbitrary Precision types Needed?

Code using native C int type

```
int foo_top(int a, int b, int c)
{
  int sum, mult;
  sum=a+b;
  mult=sum*c;
  return mult;
}

Synthesis

a
b
c

32-bit Add & Mult

return

return
```

- ▶ However, if the inputs will only have a max range of 8-bit
  - Arbitrary precision data-types should be used

```
int17 foo_top(int8 a, int8 b, int8 c)
{
  int9 sum;
  int17 mult
  sum=a+b;
  mult=sum*c;
  return mult;
}

Synthesis

Synthesis

foo_top

9-bit Add, 17-bit Mult

return

return
```

It will result in smaller & faster hardware with full precision, full precision can be simulated/validated with C simulation and hardware will behave the same



#### **DEPEDENCE** Directive

#### **Loop-Independent Dependence**

Same element is accessed in a single loop iteration

```
for (i=0; i<N; i++) {
    A[i]=x;
    y=A[i];
}</pre>
```

#### **Loop-Carried Dependence**

Same element is accessed from a different loop iteration

```
for (i=1; i<N; i++) {
    A[0]=0;
    A[i]=A[i-1]*2;
}</pre>
```

DEPENDENCE directive allows you to explicitly define the dependencies and eliminate a false dependence

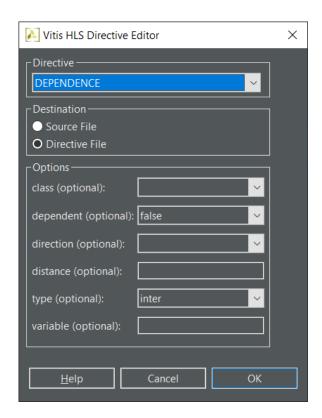


#### **DEPEDENCE** Directive

- DEPENDENCE directive allows to explicitly define the dependencies and eliminate a false dependence
  - Provides the tool with additional information about the dependencies
  - Allows designers to inform tool of "external" conditions
  - Makes sure there are no dependencies between loop iterations

```
for (row = 0; row < rows + 1; row++) {
    L2: for (col = 0; col < cols + 1; col++) {
    #pragma HLS PIPELINE II=1

    #pragma HLS dependence variable=buff_A type=inter dependent=false
    #pragma HLS dependence variable=buff_B type=inter dependent=false
    if (col < cols) {
        buff_A[2][col] = buff_A[1][col]; // read from buff_A[1][col]
        buff_A[1][col] = buff_A[0][col]; // write to buff_A[1][col]
        buff_B[1][col] = buff_B[0][col];
        temp = buff_A[0][col];} }</pre>
```



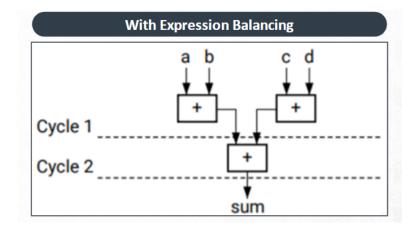


### **Optimizing Logic Expressions**

- Expression balancing rearranges operators to construct a balanced tree and reduce latency
  - For integer operations expression balancing is on by default but may be disabled using the EXPRESSION\_BALANCE pragma or directive.
  - For floating-point operations, expression balancing is off by default but may be enabled using using the config\_compile unsafe\_math\_optimizations command.

```
data_t foo_top (data_t a, data_t b, data_t c, data_t d) {
   data_t sum;
   sum = 0;
   sum += a;
   sum += b;
   sum += c;
   sum += d;
   return sum;
}
```

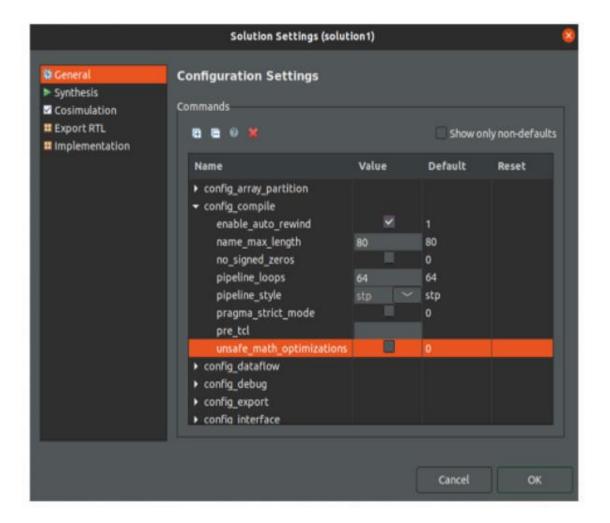
# Cycle 1 Cycle 2 Cycle 3 Cycle 4 Sum





### **Optimizing Logic Expressions**

- To configure the tool to enable expression balancing with float and double types
- Place the pragma in the C source within the boundaries of the required location.
  - #pragma HLS expression\_balance off
  - Turns off expression balancing at this location.
  - Specifying #pragma HLS expression\_balance enables expression balancing in the specified scope. Adding off disables it.
- Solution settings > General > Config\_compile>unsafe\_math\_optimizations





# Summary



## **Summary**

- Resource utilization can be reduced using allocation and binding controls
  - Allocation can define how many resource are used
- ▶ The design structure can be controlled by
  - Inlining functions: direct impact on RTL hierarchy & optimization possibilities
  - Loops: direct impact on reuse of resources
  - Arrays: direct impact on the RAM
- Major area optimization techniques
  - Minimize bit widths
  - Map smaller arrays into larger arrays
    - Make better use of existing RAMs
  - Control loop hierarchy
  - Control function call hierarchy
  - Control the number of operators and cores
- Arbitrary precision data types help controlling both the area and resource utilization



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