

mimalloc研究报告



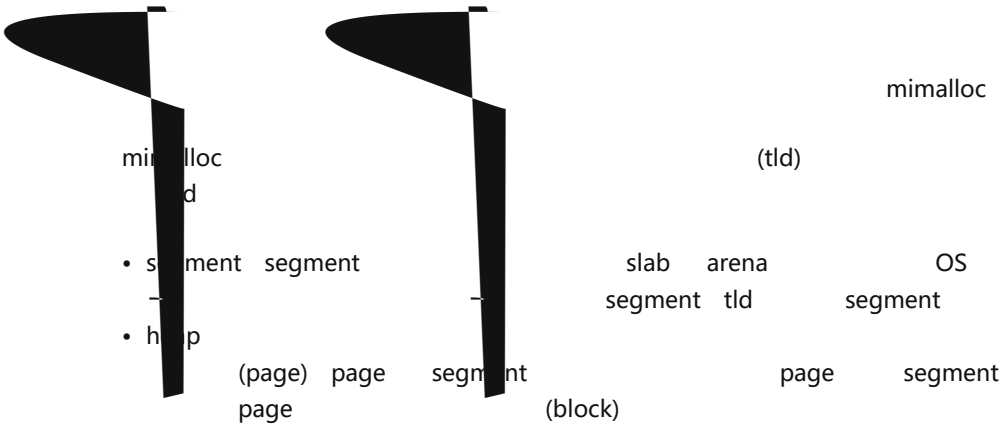
co lin

242

前言

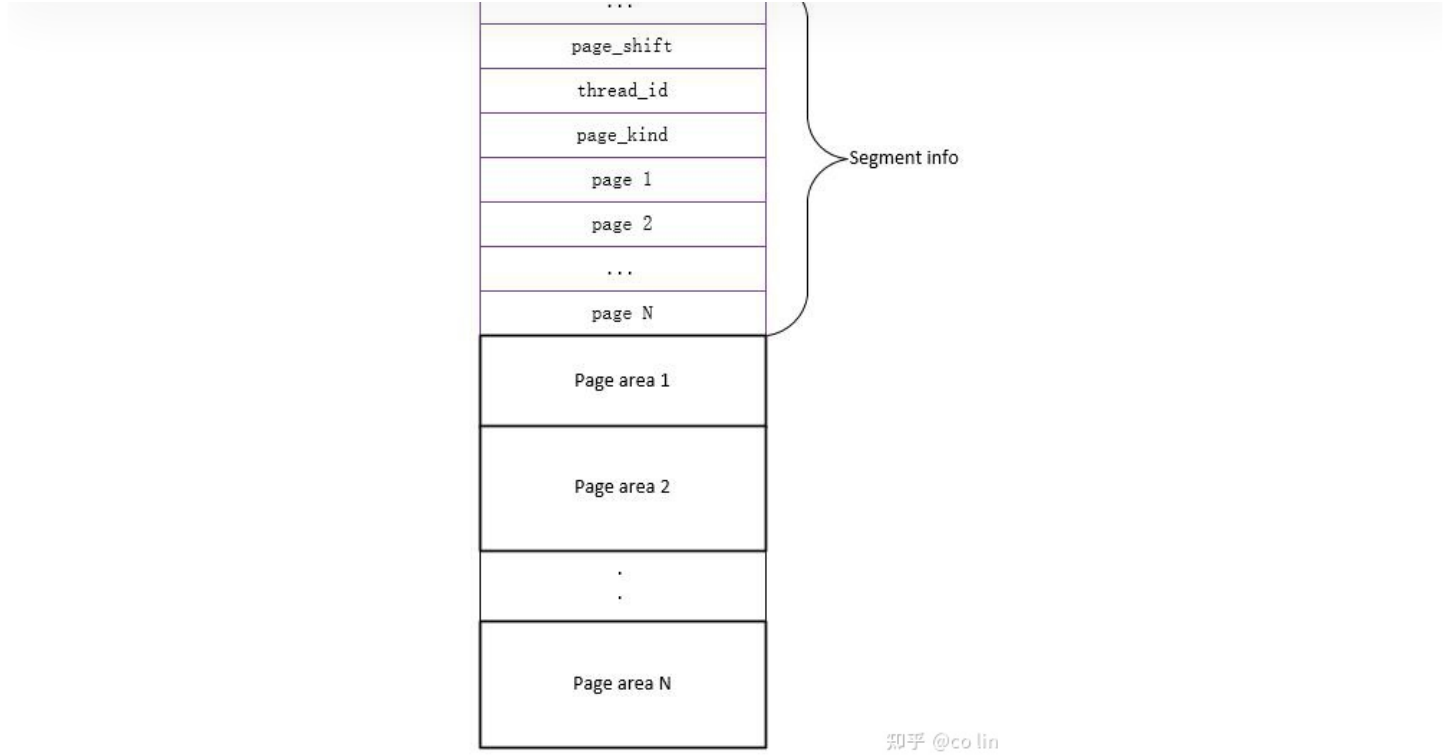
mimalloc

线程本地数据、



内存段(segment)

segment



segment

(segment info)

N

page

mi_segment_t

• thread_id

• page_kind

• page_shift

• page1~pageN

1

~

1

segment

N

page

1

4Mb

512kb

4Mb

4Mb

64kb

segment

4Mb

segment

segment

static inline

return

*

=

*

& ~ *

* _mi_ptr_segment

const void*

SMALL, MEDIUM, LARGE

HUGE

segment

segment

4Mb

HUGE

segment

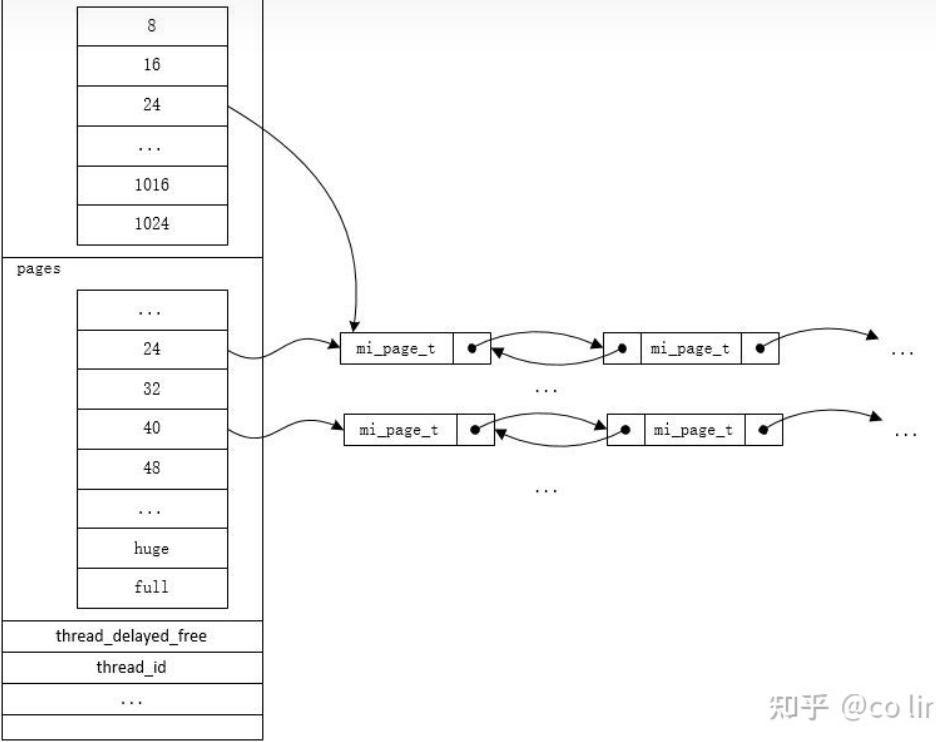
线程堆(heap)

segment

heap

size

class



```
pages
size page
mi_page_queue_t
pages
mi_page_queue_t
segment
page
page

pages_free_direct
(64
pages
mi_page_t*
8)
pages
1024
mi_page_t
page
pages_free_direct
pages_free_direct

pages_free_direct
heap
size page

* = -> + >> //

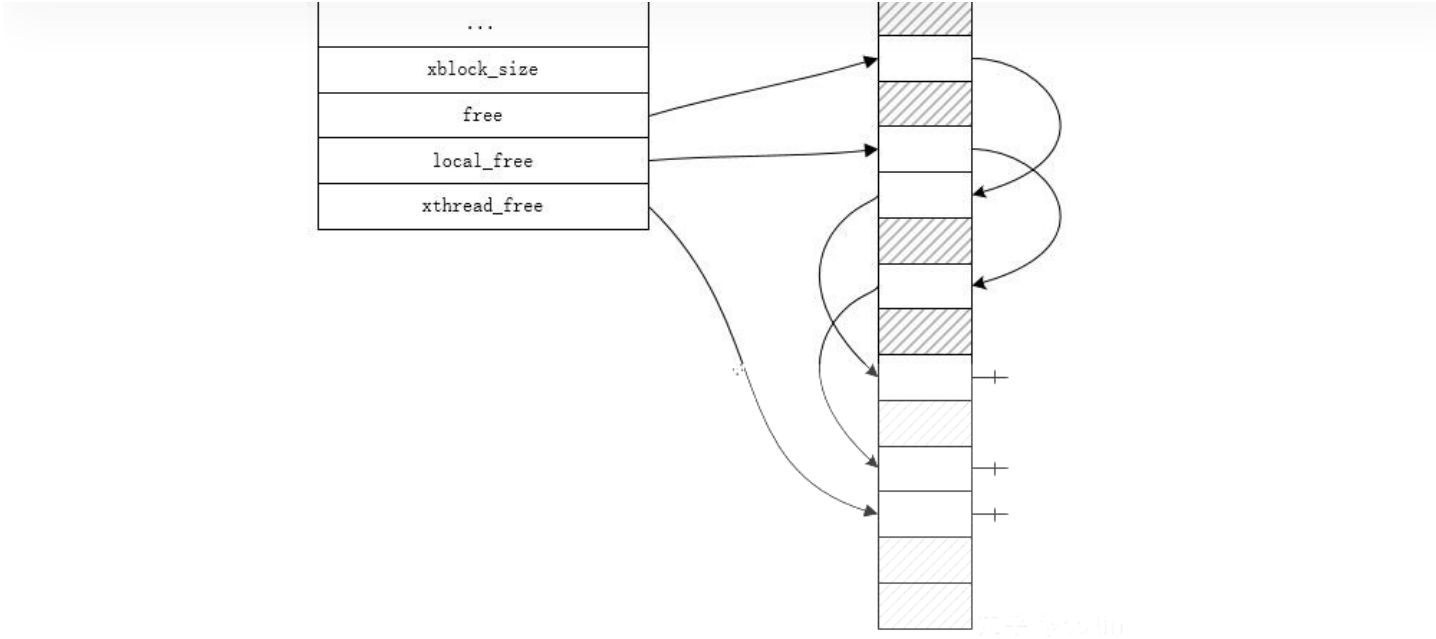
1024
mi_page_queue_t
size class
pages
mi_page_t
size
pages_free_direct

pages
)
full page(
full

mimalloc
thread_delayed_free
```

内存页(page)

(mi_page_t)



mi_page_t	(xblock_size)	freelist	mimalloc	(block)
• free				
• local_free			free	local_free
• local_free	free	free		
• xthread_free			free	local_free
	xthread_free			mimalloc

```
void* _mi_page_malloc
{
    * const = ->
    if ==
        return // slow path

    // pop from the free list
    -> ++
    -> = ->
    return
}
```

free

(_mi_malloc_generic)

free

local_free

" "

local_free thread_free

```
void _mi_page_free_collect *
```

```
// and the local free list
```

```

if      ->      !=
if      ->      ==
    // usual case
        ->      =      ->
        ->      =

```



```

thread_free      free      local_free
local_free      local_free      free
                                thread_free

void atomic_push      **      *
do      ->      = *
while !      ->

&      ->

lock free      _mi_page_thread_free_collect      thread_free

=      &      ->
      ->

thread_free      free      thread_free
      thread_free      page

mimalloc      freelist

```

分配和释放

mimalloc

```
void* mi_malloc
    * = // 取线程相关的堆
return
```

```
void* mi_heap_malloc      *
if      <=                // 如果<=1024, 进入小对象分配
    return
else    // 否则进行通用分配
    return
```

```

        *      =      ->
if      !=      // fast path
        ->      =      ->
        ->      ++
        return
else
        return      // slow path

void* mi_malloc_generic      *
        =      // 计算得到 size class
        *      =      ->      // 取相应的page队列

        // 收集页里的可用内存
if      ->      ==      // 整个页都空闲，回收掉

        else if      ->      !=      // 收集完如果有可用内存，则重分配入口
        return

// 到这儿表明找不到可用的page，从segment分配一个新鲜的page

```

```

void page_collect
// 先收集thread free list
if      !=

// 然后才是local free list
if      ->      !=
if      ->      ==
// usual case
        ->      =      ->
        ->      =

```

```

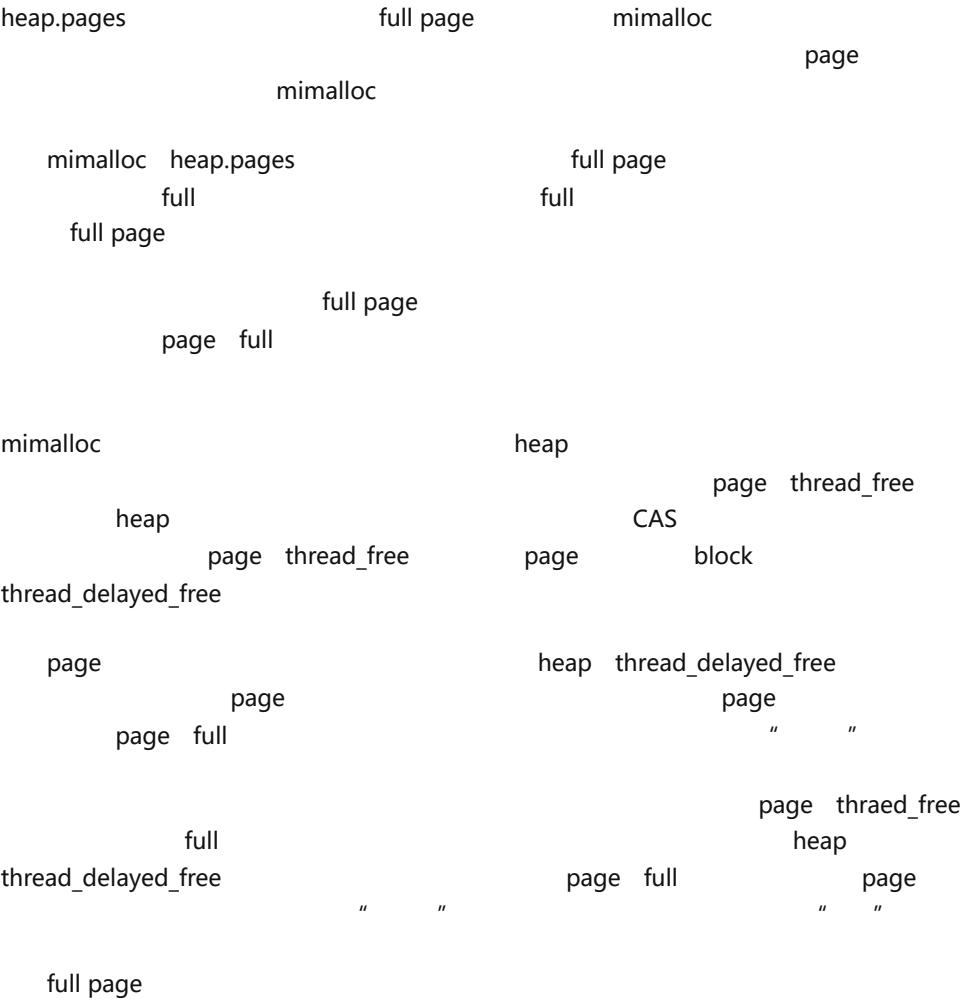
void mi_free void*
        *      =      *      & ~ *      // 找到对应的segment
if      ==      return
// 找到对应的page，这是简化过的，第1个page要特殊处理。
// 因为segment等分成N个page，这里只需要取相对地址，然后除去page的大小，即得到page的索引。
        *      = &      ->      -      >>      ->
        *      =      *

if      ==      ->      // 相同线程，释放到local_free
        ->      =      ->
        ->      =
        ->      --
if      ->      ==

else // 不同线程，释放到 thread_free

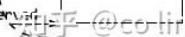
```

满页的处理(full page)



总结

- mimalloc
1. freelist
freelist
CPU
size class
 2. free
mimalloc lock-
 3. paper Heap Layout



- 2021-05-06 23:36

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Tools boy

2021-05-31

mimalloc v2.0.1 slice page



co lin () Tools boy

2021-05-31



2021-05-11

