语义信息在领域编译中的角色

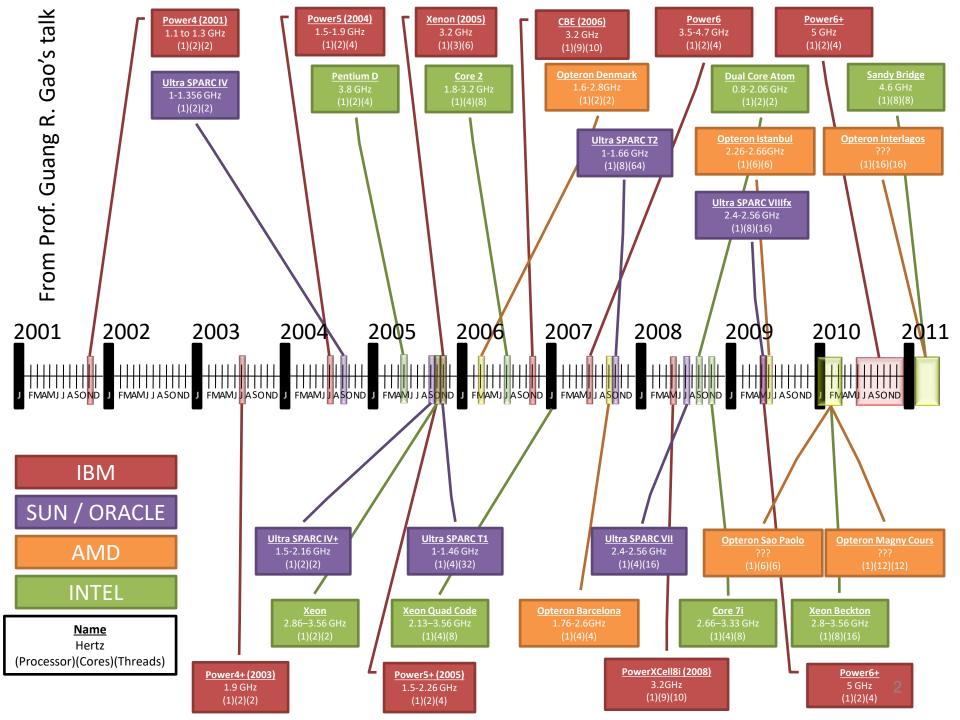
Using Semantic To Guide Compiler Optimizations

中科院计算所

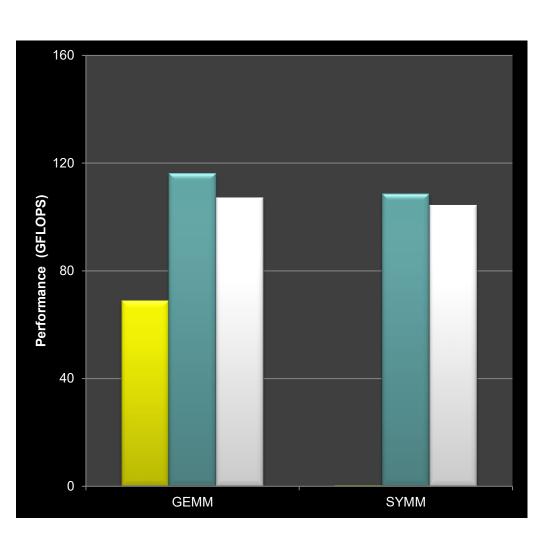
崔慧敏

2020年10月

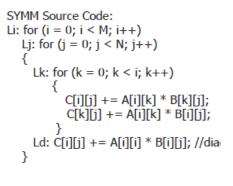




Semantic Case 1: SYMM



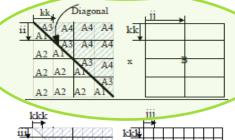
```
SYMM Source Code:
Li: for (i = 0; i < M; i++)
Lj: for (j = 0; j < N; j++)
{
    Lk: for (k = 0; k < i; k++)
    {
        C[i][j] += A[i][k] * B[k][j];
        C[k][j] += A[i][k] * B[i][j];
    }
Ld: C[i][j] += A[i][i] * B[i][j]; //diagonal
}
```

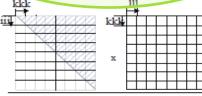


Triangular area

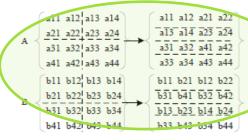
iii kk

Symmetry





Matrix-mul



第一步,线程划分:对循环迭代j进行strip-mining,并对循环次序进行调整,原始的循环顺序变为 j, i, jj, k。并将最外层循环j分配到不同的线程。

第二步,循环分块:对循环迭代i和k进行循环分块(tiling),改善L2缓存的局部性。这时新的循环顺序为**j, i, k, ii, jj, kk**.

第三步,循环分裂: 将循环体分裂为三个,分别对应 real area (A1,…, A2,…), shadow area (A3,…, A4,…), 和对角线元素.

第四步,循环剥离(peeling):将三角形区域从矩形区域中剥离出来,即将左图中的A1和A3分别从A2和A4中剥离。

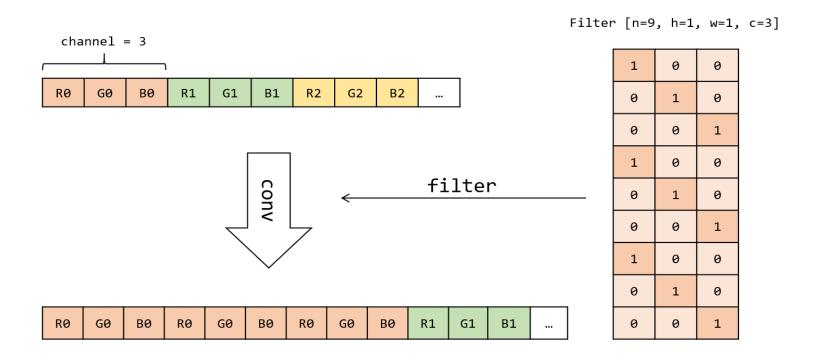
*第五步:循环分块:*对剥离出的三角形区域A1和A3分别实施循环分块。对每个矩形区域(其循环顺序为ii, jj, kk)实施循环分块,这时新的循环顺序为**j**, i, k, ii, jj, kk, iii, jjj, kkk

第六步:数据布局重组:将源矩阵A和B进行数据布局重组,从行优先或者列优先转换为块优先,块内和块间的顺序则遵从原有的行/列优先顺序。如左图所示,这一变换是为了提高L1数据缓存的局部性。



Semantic Case2: resize

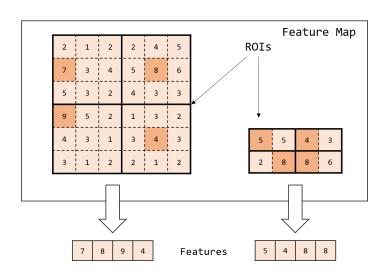
■We can use conv. to compute resize

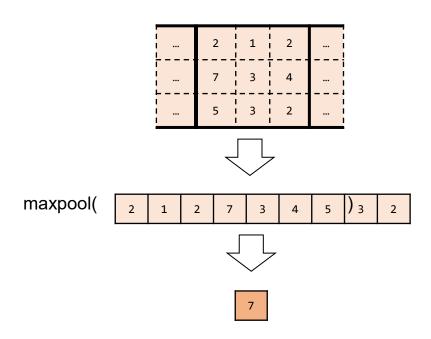




Semantic Case3: ROIPooling

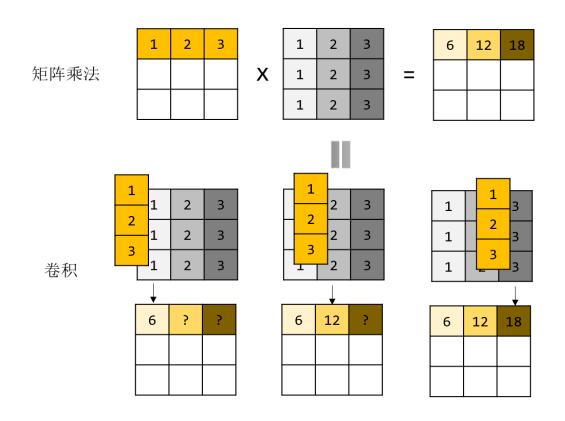
We can use Maxpooling to compute ROI pooling.



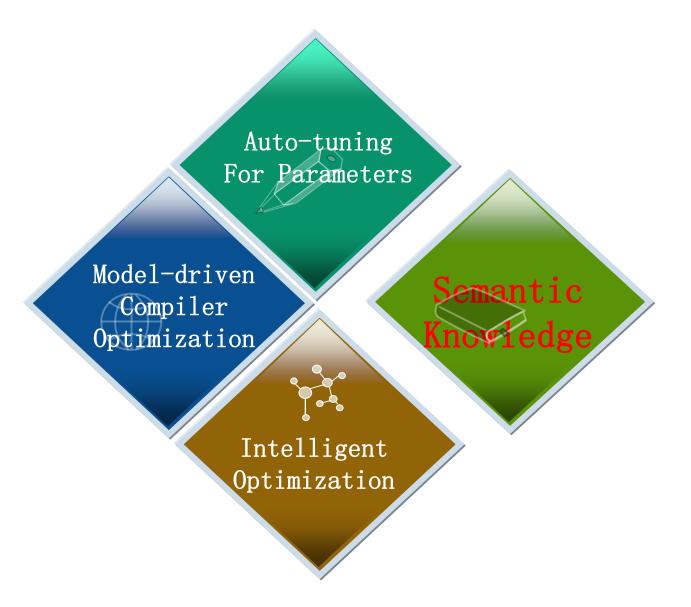


Semantic Case4: GEMM

We can use Conv to compute GEMM.



How to get Peak Performance?





Key Challenges

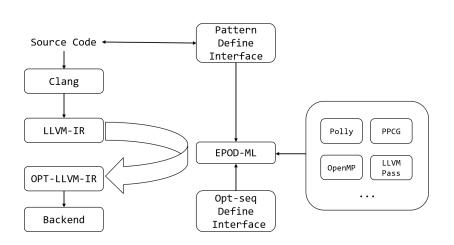
• How to connect into the workflow of compilers?

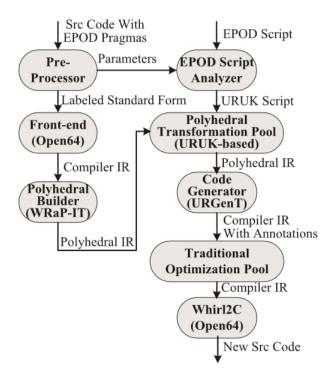
• How to specify an optimization sequence?

• How to apply specific opt to specific code region?



Extendable Pattern for Compilers







```
stmts {
stmt 1 {
   domain { i: (0, 63), j:(0, 63)},
  position {i, j, 0, 0},
  mem acc {
     store 0 { mem {C}, domain {i:(0, 63), j:(0, 63)}}
  },
  dep { waw { stmt 1, stmt 2 } },
  expr { st 0(const(0)) }
stmt 2 {
   domain { i: (0, 63), j:(0, 63), k(0, 63)},
  positon \{i, j, 1, k\},
  mem acc {
     load 0 { mem {C}, domain {i:(0, 63), j:(0, 63)}},
     load 1 { mem {A}, domain {i:(0, 63), k:(0, 63)}},
     load 2 { mem {B}, domain \{k:(0, 63), j:(0, 63)\}\},
     store_0 { mem {C}, domain {i:(0, 63), j:(0, 63)}}
  },
  dep { waw { stmt 1, stmt 2 } },
   expr { st 0(add(1d(0), mul(1d(1), 1d(2))))}
```

1. Define a semantic pattern

- Similar with Polyhedral SCoP
- Each statement includes:
 - Loop domain
 - Position
 - Memory Reference
 - Dependency
 - Expression



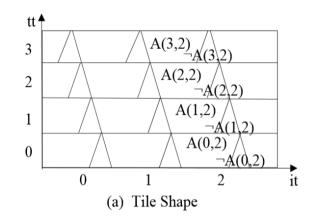
```
def rep conv(loop) {
  s = loop.get stmt();
  d = s.get domain(0);
  col = d.range();
  row = d.next().range();
 dst = s.get store();
  src = s.get_load(d);
 filter = s.get load();
  s.replace_with("||vm_conv(dst, src, filter,
                col, 1, row, 1, 1, 1, 1, 1);");
def matmul {
 I = 16, J = 16
  11 = polly.loop_tiling($$, [0, 1], [I, J])
  1 0, l_1 = polly.loop_fission(ll, 3);
  Ilvm.auto vec(1 0);
  [lym.rep\_conv(l_1);
```

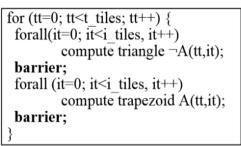
2. Define an optimization seq.

- Polly for loop transformations.
 - Easy for extension
 - e.g. split tiling/ overlap tiling/ diamond tiling
- Backend opt. on LLVM/Open64.
 - Vectorization
 - On-chip memory usage

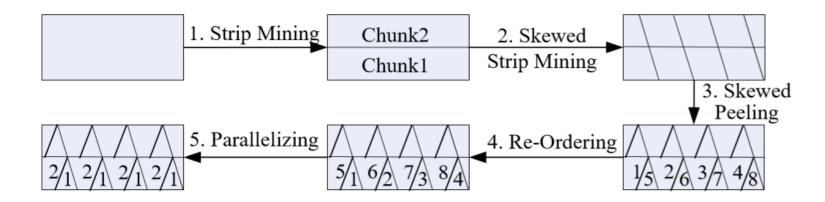


```
 \begin{aligned} & \text{Original Loop:} \\ & \text{for } (t=1;\,t < T;\,t++) \; \{ \\ & \text{for } (i=1;\,i < N;\,i++) \\ & B[i] = (A[i-1] + A[i] + A[i+1])/3; \\ & \text{for } (i=1;\,i < N;\,i++) \\ & A[i] = B[i]; \\ \} \\ & \text{Single Statement Form:} \\ & \text{for } (t=1;\,t < T;\,t++) \{ \\ & \text{for } (i=1;\,i < N;\,i++) \\ & A[t,\,i] = (A[t-1,i-1] + A[t-1,i] + A[t-1,i+1])/3; \\ \} \end{aligned}
```





(b) Execution Order



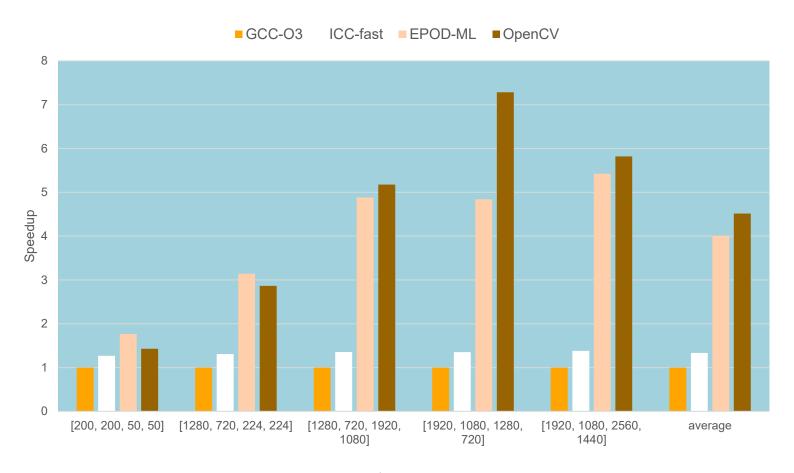


3. Specify a semantic region

```
#pragma EPOD-ML jacobi
for (t=0; t<T; t++) {
    for (i=0; i<N; i++) {
        for (j=0;j<N;j++) {
            if (i!=j) {
                sum += A[i][j]*x1[i];
            }
        }
        x[i] = (b[i]-sum)/A[i][i];
    }
    for (int i=0; i<N; i++) {
        x1[i] = x[i];
    }
}
#pragma EPOD-ML end</pre>
```



Evaluation

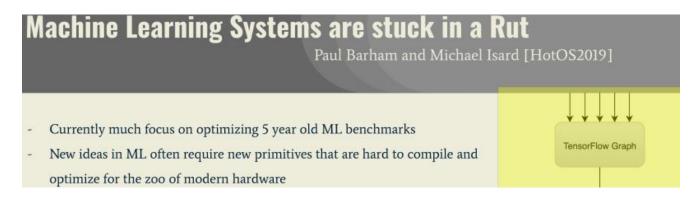


Resize加速比-CPU



人工智能编程语言与编译器

■问题1:人工智能需要通用的编程语言吗?



■问题2: 芯片敏捷化 vs. 编译器敏捷化

■问题3:编译器敏捷化 vs. 峰值性能





