# libbfd

The Binary File Descriptor Library

First Edition—BFD version < 3.0 % Since no product is stable before version 3.0:-) Original Document Created: April 1991

Steve Chamberlain Cygnus Support

Free Software Foundation  $sac@www.gnu.org \\BFD, 1.5$  TeXinfo 2009-03-28.05

Copyright © 1991-2014 Free Software Foundation, Inc.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.3 or any later version published by the Free Software Foundation; with no Invariant Sections, with no Front-Cover Texts, and with no Back-Cover Texts. A copy of the license is included in the section entitled "GNU Free Documentation License".

# Table of Contents

1	Introduc	tion	1
	1.1 History		. 1
		Use BFD	
	1.3 What BF	D Version 2 Can Do	. 2
		rmation Loss	
		BFD canonical object-file format	
		v	
2	BFD Fro	ont End	5
		bfd	
		orting	
		e bfd_error_type	
	2.2.1.1	bfd_get_error	
	2.2.1.2	bfd_set_error	
	2.2.1.3	bfd_errmsg	
	2.2.1.4	bfd_perror	
	2.2.2 BFD	error handler	
	2.2.2.1	bfd_set_error_handler	
	2.2.2.2	bfd_set_error_program_name	
	2.2.2.3	bfd_get_error_handler	
	2.2.3 BFD	assert handler	
	2.2.3.1	bfd_set_assert_handler	13
	2.2.3.2	bfd_get_assert_handler	
	2.3 Miscellan	eous	14
	2.3.1 Misc	cellaneous functions	14
	2.3.1.1	bfd_get_reloc_upper_bound	14
	2.3.1.2	bfd_canonicalize_reloc	14
	2.3.1.3	bfd_set_reloc	14
	2.3.1.4	bfd_set_file_flags	14
	2.3.1.5	bfd_get_arch_size	15
	2.3.1.6	bfd_get_sign_extend_vma	15
	2.3.1.7	bfd_set_start_address	15
	2.3.1.8	bfd_get_gp_size	15
	2.3.1.9	bfd_set_gp_size	15
	2.3.1.10	bfd_scan_vma	16
	2.3.1.11	bfd_copy_private_header_data	16
	2.3.1.12	bfd_copy_private_bfd_data	16
	2.3.1.13	bfd_merge_private_bfd_data	16
	2.3.1.14	bfd_set_private_flags	17
	2.3.1.15	Other functions	17
	2.3.1.16	bfd_alt_mach_code	19
	2.3.1.17	bfd_emul_get_maxpagesize	19
	2.3.1.18	bfd_emul_set_maxpagesize	19
	2.3.1.19	bfd_emul_get_commonpagesize	19

2.3.1.20	bfd_emul_set_commonpagesize	20
2.3.1.21	bfd_demangle	. 20
2.3.1.22	struct bfd_iovec	20
2.3.1.23	bfd_get_mtime	21
2.3.1.24	_	
2.3.1.25	_	
2.4 Memory	Usage	
	$\operatorname{tion}^{\circ}$	
2.5.1 Initi	alization functions	22
2.5.1.1	bfd_init	22
2.6 Sections.		. 22
2.6.1 Sect	ion input	. 22
2.6.2 Sect	ion output	23
2.6.3 Link	g orders	23
2.6.4 type	edef asection	23
2.6.5 Sect	ion prototypes	36
	bfd_section_list_clear	
2.6.5.2		
2.6.5.3	bfd_get_next_section_by_name	
2.6.5.4	bfd_get_linker_section	
2.6.5.5	bfd_get_section_by_name_if	
2.6.5.6	bfd_get_unique_section_name	
2.6.5.7	bfd_make_section_old_way	
2.6.5.8	bfd_make_section_anyway_with_flags	
2.6.5.9	bfd_make_section_anyway	
2.6.5.10		
2.6.5.11		
2.6.5.12	bfd_set_section_flags	38
2.6.5.13		
2.6.5.14		
2.6.5.15	bfd_sections_find_if	
2.6.5.16	bfd_set_section_size	
2.6.5.17		
2.6.5.18		
2.6.5.19	bfd_malloc_and_get_section	
2.6.5.20	bfd_copy_private_section_data	
2.6.5.21	bfd_generic_is_group_section	
2.6.5.22	bfd_generic_discard_group	
·	ding symbols	
	ting symbols	
	i Symbols	
	edef asymbol	
	abol handling functions	
2.7.5.1	bfd_get_symtab_upper_bound	
2.7.5.2	bfd_is_local_label	
2.7.5.3	bfd_is_local_label_name	
	hfd is target special symbol	

2.7.5.5	bfd_canonicalize_symtab	47
2.7.5.6	bfd_set_symtab	. 47
2.7.5.7	bfd_print_symbol_vandf	. 47
2.7.5.8	bfd_make_empty_symbol	. 48
2.7.5.9	_bfd_generic_make_empty_symbol	
2.7.5.10		
2.7.5.11		
2.7.5.12	· · · · · · · · · · · · · · · · · · ·	
2.7.5.13	· · · · · · · · · · · · · · · · · · ·	
2.7.5.14		
2.8 Archives		
2.8.1 Arcl	nive functions	50
2.8.1.1	bfd_get_next_mapent	. 50
2.8.1.2	bfd_set_archive_head	. 50
2.8.1.3	bfd_openr_next_archived_file	. 50
2.9 File form	ats	
2.9.1 File	format functions	. 51
2.9.1.1	bfd_check_format	. 51
2.9.1.2		
2.9.1.3		
2.9.1.4	bfd_format_string	
2.10 Relocati	ions	
2.10.1 typ	pedef arelent	. 52
2.10.1.1	enum complain_overflow	. 55
2.10.1.2		
2.10.1.3		
2.10.1.4	bfd_get_reloc_size	. 58
2.10.1.5	<del>-</del>	
2.10.1.6	bfd_check_overflow	
2.10.1.7	bfd_perform_relocation	
2.10.1.8	bfd_install_relocation	
2.10.2 Th	e howto manager	
2.10.2.1	9	
2.10.2.2	· -	
2.10.2.3		
2.10.2.4	bfd_get_reloc_code_name	
2.10.2.5	bfd_generic_relax_section	
2.10.2.6	bfd_generic_gc_sections	
2.10.2.7		
2.10.2.8	bfd_generic_merge_sections	
2.10.2.9	bfd_generic_get_relocated_section_contents	
		129
2.11 Core file	es	
	re file functions	
2.11.1.1	bfd_core_file_failing_command	
2.11.1.2	bfd_core_file_failing_signal	
2.11.1.3		
	core file matches executable n	

2.11.1.5	5 generic_core_file_matches_executable_p	130
2.12 Targets	S	130
2.12.1 bf	fd_target	131
2.12.1.1	1 bfd_set_default_target	141
2.12.1.5	<u> </u>	
2.12.1.3	<u> </u>	
2.12.1.4		
2.12.1.5	-	
2.13 Archite	ectures	
	fd_architecture	
	fd_arch_info	
2.13.2.1		
2.13.2.5	-	
2.13.2.3		
2.13.2.4		
2.13.2.	0 _1	
2.13.2.0		
2.13.2.		
2.13.2.8		
2.13.2.9	_	
2.13.2.		
2.13.2.1		
2.13.2.		
2.13.2.1	<del>-</del>	
2.13.2.1		
2.13.2.1	_5	
2.13.2.1	_ 1 -	
2.13.2.1	-1	
2.13.2.1	1	
2.13.2.1		
	ng and closing BFDs	
	unctions for opening and closing	
2.14.1.1	. 0	
2.14.1.5	-	
	3 bfd_fdopenr	
2.14.1.4	-	
2.14.1.		
2.14.1.0	-	
2.14.1.	<b>– 1</b>	
2.14.1.8	<del>-</del>	
2.14.1.9		
2.14.1.	_	
2.14.1.1		
2.14.1.		
2.14.1.	· · · · · = · · · · · · · · · · · · · ·	
2.14.1.	_	
2.14.1.		
2.14.1.	_	
4.14.1.	IO DIU_CAIC_KIIU_UCDUKIIIK_CLCOZ	

2.14.1.17	bfd_get_debug_link_info	160
2.14.1.18	bfd_get_alt_debug_link_info	160
2.14.1.19	separate_debug_file_exists	161
2.14.1.20	separate_alt_debug_file_exists	161
2.14.1.21	find_separate_debug_file	161
2.14.1.22	bfd_follow_gnu_debuglink	
2.14.1.23	bfd_follow_gnu_debugaltlink	
2.14.1.24	bfd_create_gnu_debuglink_section	
2.14.1.25	bfd_fill_in_gnu_debuglink_section	
2.15 Implemen	ntation details	
2.15.1 Inter	rnal functions	162
2.15.1.1	bfd_write_bigendian_4byte_int	162
2.15.1.2	bfd_put_size	
2.15.1.3	bfd_get_size	
2.15.1.4	bfd_h_put_size	
2.15.1.5	bfd_log2	
2.16 File cachi	ing	
	hing functions	
2.16.1.1	bfd_cache_init	
2.16.1.2	bfd_cache_close	
2.16.1.3	bfd_cache_close_all	
2.16.1.4	bfd_open_file	
2.17 Linker Fu	ınctions	
2.17.1 Crea	ating a linker hash table	167
2.17.2 Add	ing symbols to the hash table	167
2.17.2.1	Differing file formats	
2.17.2.2	Adding symbols from an object file	168
2.17.2.3	Adding symbols from an archive	169
2.17.3 Perf	forming the final link	169
2.17.3.1	Information provided by the linker	170
2.17.3.2	Relocating the section contents	170
2.17.3.3	Writing the symbol table	170
2.17.3.4	bfd_link_split_section	
2.17.3.5	bfd_section_already_linked	
2.17.3.6	bfd_generic_define_common_symbol	171
2.17.3.7	bfd_find_version_for_sym	
2.17.3.8	bfd_hide_sym_by_version	172
2.18 Hash Tab	bles	
2.18.1 Crea	ating and freeing a hash table	172
	king up or entering a string	
	versing a hash table	
2.18.4 Deri	ving a new hash table type	173
2.18.4.1	Define the derived structures	
2.18.4.2	Write the derived creation routine	174
	Write other derived routines	

3	$\mathbf{BFD}$	) ba	ck ends	<b>176</b>
	3.1 Wh	at to	Put Where	176
	3.2 a.o.	ut bac	kends	176
	3.2.1	Relo	ocations	. 177
	3.2.2	Inte	rnal entry points	. 177
	3.2	2.2.1	aout_size_swap_exec_header_in	. 177
	3.5	2.2.2	aout_size_swap_exec_header_out	. 177
	3.5	2.2.3	aout_size_some_aout_object_p	. 178
	3.2	2.2.4	aout_size_mkobject	. 178
	3.2	2.2.5	aout_size_machine_type	178
	3.2	2.2.6	aout_size_set_arch_mach	. 178
	3.2	2.2.7	aout_size_new_section_hook	179
	3.3  coff	f back	ends	. 179
	3.3.1		ting to a new version of coff	
	3.3.2		the coff backend works	. 179
		3.2.1	File layout	
		3.2.2	Coff long section names	
		3.2.3	Bit twiddling	
		3.2.4	Symbol reading	
		3.2.5	Symbol writing	
		3.2.6	coff_symbol_type	
	3.3	3.2.7	bfd_coff_backend_data	
		3.2.8	Writing relocations	
		3.2.9	Reading linenumbers	
		3.2.10	Reading relocations	
			kends	
			kend	
	3.5.1		layout	
	3.5.2		abol table format	
	3.5.3	mme	o section mapping	195
R	FD In	dev		205

# 1 Introduction

BFD is a package which allows applications to use the same routines to operate on object files whatever the object file format. A new object file format can be supported simply by creating a new BFD back end and adding it to the library.

BFD is split into two parts: the front end, and the back ends (one for each object file format).

- The front end of BFD provides the interface to the user. It manages memory and various canonical data structures. The front end also decides which back end to use and when to call back end routines.
- The back ends provide BFD its view of the real world. Each back end provides a set of calls which the BFD front end can use to maintain its canonical form. The back ends also may keep around information for their own use, for greater efficiency.

# 1.1 History

One spur behind BFD was the desire, on the part of the GNU 960 team at Intel Oregon, for interoperability of applications on their COFF and bout file formats. Cygnus was providing GNU support for the team, and was contracted to provide the required functionality.

The name came from a conversation David Wallace was having with Richard Stallman about the library: RMS said that it would be quite hard—David said "BFD". Stallman was right, but the name stuck.

At the same time, Ready Systems wanted much the same thing, but for different object file formats: IEEE-695, Oasys, Srecords, a.out and 68k coff.

BFD was first implemented by members of Cygnus Support; Steve Chamberlain (sac@cygnus.com), John Gilmore (gnu@cygnus.com), K. Richard Pixley (rich@cygnus.com) and David Henkel-Wallace (gumby@cygnus.com).

# 1.2 How To Use BFD

To use the library, include 'bfd.h' and link with 'libbfd.a'.

BFD provides a common interface to the parts of an object file for a calling application.

When an application successfully opens a target file (object, archive, or whatever), a pointer to an internal structure is returned. This pointer points to a structure called bfd, described in 'bfd.h'. Our convention is to call this pointer a BFD, and instances of it within code abfd. All operations on the target object file are applied as methods to the BFD. The mapping is defined within bfd.h in a set of macros, all beginning with 'bfd\_' to reduce namespace pollution.

For example, this sequence does what you would probably expect: return the number of sections in an object file attached to a BFD abfd.

```
#include "bfd.h"
unsigned int number_of_sections (abfd)
bfd *abfd;
{
```

```
return bfd_count_sections (abfd);
}
```

The abstraction used within BFD is that an object file has:

- a header,
- a number of sections containing raw data (see Section 2.6 [Sections], page 22),
- a set of relocations (see Section 2.10 [Relocations], page 52), and
- some symbol information (see Section 2.7 [Symbols], page 41).

Also, BFDs opened for archives have the additional attribute of an index and contain subordinate BFDs. This approach is fine for a out and coff, but loses efficiency when applied to formats such as S-records and IEEE-695.

# 1.3 What BFD Version 2 Can Do

When an object file is opened, BFD subroutines automatically determine the format of the input object file. They then build a descriptor in memory with pointers to routines that will be used to access elements of the object file's data structures.

As different information from the object files is required, BFD reads from different sections of the file and processes them. For example, a very common operation for the linker is processing symbol tables. Each BFD back end provides a routine for converting between the object file's representation of symbols and an internal canonical format. When the linker asks for the symbol table of an object file, it calls through a memory pointer to the routine from the relevant BFD back end which reads and converts the table into a canonical form. The linker then operates upon the canonical form. When the link is finished and the linker writes the output file's symbol table, another BFD back end routine is called to take the newly created symbol table and convert it into the chosen output format.

### 1.3.1 Information Loss

Information can be lost during output. The output formats supported by BFD do not provide identical facilities, and information which can be described in one form has nowhere to go in another format. One example of this is alignment information in b.out. There is nowhere in an a.out format file to store alignment information on the contained data, so when a file is linked from b.out and an a.out image is produced, alignment information will not propagate to the output file. (The linker will still use the alignment information internally, so the link is performed correctly).

Another example is COFF section names. COFF files may contain an unlimited number of sections, each one with a textual section name. If the target of the link is a format which does not have many sections (e.g., a.out) or has sections without names (e.g., the Oasys format), the link cannot be done simply. You can circumvent this problem by describing the desired input-to-output section mapping with the linker command language.

Information can be lost during canonicalization. The BFD internal canonical form of the external formats is not exhaustive; there are structures in input formats for which there is no direct representation internally. This means that the BFD back ends cannot maintain all possible data richness through the transformation between external to internal and back to external formats.

This limitation is only a problem when an application reads one format and writes another. Each BFD back end is responsible for maintaining as much data as possible, and the internal BFD canonical form has structures which are opaque to the BFD core, and exported only to the back ends. When a file is read in one format, the canonical form is generated for BFD and the application. At the same time, the back end saves away any information which may otherwise be lost. If the data is then written back in the same format, the back end routine will be able to use the canonical form provided by the BFD core as well as the information it prepared earlier. Since there is a great deal of commonality between back ends, there is no information lost when linking or copying big endian COFF to little endian COFF, or a.out to b.out. When a mixture of formats is linked, the information is only lost from the files whose format differs from the destination.

# 1.3.2 The BFD canonical object-file format

The greatest potential for loss of information occurs when there is the least overlap between the information provided by the source format, that stored by the canonical format, and that needed by the destination format. A brief description of the canonical form may help you understand which kinds of data you can count on preserving across conversions.

files

Information stored on a per-file basis includes target machine architecture, particular implementation format type, a demand pageable bit, and a write protected bit. Information like Unix magic numbers is not stored here—only the magic numbers' meaning, so a ZMAGIC file would have both the demand pageable bit and the write protected text bit set. The byte order of the target is stored on a per-file basis, so that big- and little-endian object files may be used with one another.

sections

Each section in the input file contains the name of the section, the section's original address in the object file, size and alignment information, various flags, and pointers into other BFD data structures.

symbols

Each symbol contains a pointer to the information for the object file which originally defined it, its name, its value, and various flag bits. When a BFD back end reads in a symbol table, it relocates all symbols to make them relative to the base of the section where they were defined. Doing this ensures that each symbol points to its containing section. Each symbol also has a varying amount of hidden private data for the BFD back end. Since the symbol points to the original file, the private data format for that symbol is accessible. 1d can operate on a collection of symbols of wildly different formats without problems.

Normal global and simple local symbols are maintained on output, so an output file (no matter its format) will retain symbols pointing to functions and to global, static, and common variables. Some symbol information is not worth retaining; in a.out, type information is stored in the symbol table as long symbol names. This information would be useless to most COFF debuggers; the linker has command line switches to allow users to throw it away.

There is one word of type information within the symbol, so if the format supports symbol type information within symbols (for example, COFF, IEEE, Oasys) and the type is simple enough to fit within one word (nearly everything but aggregates), the information will be preserved.

#### relocation level

Each canonical BFD relocation record contains a pointer to the symbol to relocate to, the offset of the data to relocate, the section the data is in, and a pointer to a relocation type descriptor. Relocation is performed by passing messages through the relocation type descriptor and the symbol pointer. Therefore, relocations can be performed on output data using a relocation method that is only available in one of the input formats. For instance, Oasys provides a byte relocation format. A relocation record requesting this relocation type would point indirectly to a routine to perform this, so the relocation may be performed on a byte being written to a 68k COFF file, even though 68k COFF has no such relocation type.

#### line numbers

Object formats can contain, for debugging purposes, some form of mapping between symbols, source line numbers, and addresses in the output file. These addresses have to be relocated along with the symbol information. Each symbol with an associated list of line number records points to the first record of the list. The head of a line number list consists of a pointer to the symbol, which allows finding out the address of the function whose line number is being described. The rest of the list is made up of pairs: offsets into the section and line numbers. Any format which can simply derive this information can pass it successfully between formats (COFF, IEEE and Oasys).

# 2 BFD Front End

# 2.1 typedef bfd

A BFD has type bfd; objects of this type are the cornerstone of any application using BFD. Using BFD consists of making references though the BFD and to data in the BFD.

Here is the structure that defines the type bfd. It contains the major data about the file and pointers to the rest of the data.

```
enum bfd_direction
   no_direction = 0,
   read_direction = 1,
   write_direction = 2,
   both\_direction = 3
 };
struct bfd
 /* The filename the application opened the BFD with. */
 const char *filename;
 /* A pointer to the target jump table. */
 const struct bfd_target *xvec;
 /* The IOSTREAM, and corresponding IO vector that provide access
     to the file backing the BFD. */
 void *iostream;
 const struct bfd_iovec *iovec;
 /* The caching routines use these to maintain a
     least-recently-used list of BFDs. */
 struct bfd *lru_prev, *lru_next;
  /* When a file is closed by the caching routines, BFD retains
     state information on the file here... */
 ufile_ptr where;
 /* File modified time, if mtime_set is TRUE. */
 long mtime;
 /* A unique identifier of the BFD */
 unsigned int id;
 /* The format which belongs to the BFD. (object, core, etc.) */
 ENUM_BITFIELD (bfd_format) format : 3;
```

```
/* The direction with which the BFD was opened. */
 ENUM_BITFIELD (bfd_direction) direction : 2;
 /* Format_specific flags. */
 flagword flags : 17;
 /* Values that may appear in the flags field of a BFD. These also
     appear in the object_flags field of the bfd_target structure, where
     they indicate the set of flags used by that backend (not all flags
     are meaningful for all object file formats) (FIXME: at the moment,
     the object_flags values have mostly just been copied from backend
     to another, and are not necessarily correct). */
#define BFD_NO_FLAGS
                      0x00
 /* BFD contains relocation entries. */
#define HAS_RELOC
                      0x01
  /* BFD is directly executable. */
#define EXEC_P
                      0x02
 /* BFD has line number information (basically used for F_LNNO in a
     COFF header). */
#define HAS_LINENO
                      0x04
 /* BFD has debugging information. */
#define HAS_DEBUG
                      80x0
 /* BFD has symbols.
                       */
#define HAS_SYMS
                      0x10
 /* BFD has local symbols (basically used for F_LSYMS in a COFF
    header). */
#define HAS LOCALS
                      0x20
 /* BFD is a dynamic object. */
#define DYNAMIC
                      0x40
 /* Text section is write protected (if D_PAGED is not set, this is
     like an a.out NMAGIC file) (the linker sets this by default, but
     clears it for -r or -N). */
#define WP_TEXT
                      0x80
 /* BFD is dynamically paged (this is like an a.out ZMAGIC file) (the
     linker sets this by default, but clears it for -r or -n or -N). */
#define D_PAGED
                      0x100
```

```
/* BFD is relaxable (this means that bfd_relax_section may be able to
    do something) (sometimes bfd_relax_section can do something even if
    this is not set). */
#define BFD_IS_RELAXABLE 0x200
 /* This may be set before writing out a BFD to request using a
    traditional format. For example, this is used to request that when
    writing out an a.out object the symbols not be hashed to eliminate
    duplicates. */
#define BFD_TRADITIONAL_FORMAT 0x400
 /* This flag indicates that the BFD contents are actually cached
    in memory. If this is set, iostream points to a bfd_in_memory
    struct. */
#define BFD_IN_MEMORY 0x800
 /* This BFD has been created by the linker and doesn't correspond
    to any input file. */
#define BFD_LINKER_CREATED 0x1000
 /* This may be set before writing out a BFD to request that it
    be written using values for UIDs, GIDs, timestamps, etc. that
    will be consistent from run to run. */
#define BFD_DETERMINISTIC_OUTPUT 0x2000
 /* Compress sections in this BFD. */
#define BFD_COMPRESS 0x4000
 /* Decompress sections in this BFD. */
#define BFD_DECOMPRESS 0x8000
 /* BFD is a dummy, for plugins. */
#define BFD_PLUGIN 0x10000
  /* Flags bits to be saved in bfd_preserve_save. */
#define BFD_FLAGS_SAVED \
  (BFD_IN_MEMORY | BFD_COMPRESS | BFD_DECOMPRESS | BFD_PLUGIN)
  /* Flags bits which are for BFD use only. */
#define BFD_FLAGS_FOR_BFD_USE_MASK \
  (BFD_IN_MEMORY | BFD_COMPRESS | BFD_DECOMPRESS | BFD_LINKER_CREATED \
   | BFD_PLUGIN | BFD_TRADITIONAL_FORMAT | BFD_DETERMINISTIC_OUTPUT)
 /* Is the file descriptor being cached? That is, can it be closed as
    needed, and re-opened when accessed later? */
 unsigned int cacheable: 1;
```

```
/* Marks whether there was a default target specified when the
   BFD was opened. This is used to select which matching algorithm
   to use to choose the back end. */
unsigned int target_defaulted : 1;
/* ... and here: (''once'' means at least once). */
unsigned int opened_once : 1;
/* Set if we have a locally maintained mtime value, rather than
   getting it from the file each time. */
unsigned int mtime_set : 1;
/* Flag set if symbols from this BFD should not be exported. */
unsigned int no_export : 1;
/* Remember when output has begun, to stop strange things
   from happening.
                   */
unsigned int output_has_begun : 1;
/* Have archive map. */
unsigned int has_armap : 1;
/* Set if this is a thin archive. */
unsigned int is_thin_archive : 1;
/* Set if only required symbols should be added in the link hash table for
   this object. Used by VMS linkers. */
unsigned int selective_search : 1;
/* Set if this is the linker output BFD. */
unsigned int is_linker_output : 1;
/* Currently my_archive is tested before adding origin to
   anything. I believe that this can become always an add of
   origin, with origin set to 0 for non archive files. */
ufile_ptr origin;
/* The origin in the archive of the proxy entry. This will
   normally be the same as origin, except for thin archives,
   when it will contain the current offset of the proxy in the
   thin archive rather than the offset of the bfd in its actual
   container. */
ufile_ptr proxy_origin;
/* A hash table for section names. */
struct bfd_hash_table section_htab;
```

```
/* Pointer to linked list of sections. */
struct bfd_section *sections;
/* The last section on the section list. */
struct bfd_section *section_last;
/* The number of sections. */
unsigned int section_count;
/* A field used by _bfd_generic_link_add_archive_symbols. This will
   be used only for archive elements. */
int archive_pass;
/* Stuff only useful for object files:
   The start address. */
bfd_vma start_address;
/* Symbol table for output BFD (with symcount entries).
   Also used by the linker to cache input BFD symbols. */
struct bfd_symbol **outsymbols;
/* Used for input and output. */
unsigned int symcount;
/* Used for slurped dynamic symbol tables. */
unsigned int dynsymcount;
/* Pointer to structure which contains architecture information. */
const struct bfd_arch_info *arch_info;
/* Stuff only useful for archives. */
void *arelt_data;
struct bfd *my_archive; /* The containing archive BFD. */
struct bfd *archive_next; /* The next BFD in the archive. */
struct bfd *archive_head;
                            /* The first BFD in the archive. */
struct bfd *nested_archives; /* List of nested archive in a flattened
                               thin archive. */
union {
  /* For input BFDs, a chain of BFDs involved in a link. */
 struct bfd *next;
  /* For output BFD, the linker hash table. */
  struct bfd_link_hash_table *hash;
} link;
/* Used by the back end to hold private data. */
```

union {

```
struct aout_data_struct *aout_data;
    struct artdata *aout_ar_data;
    struct _oasys_data *oasys_obj_data;
    struct _oasys_ar_data *oasys_ar_data;
    struct coff_tdata *coff_obj_data;
    struct pe_tdata *pe_obj_data;
    struct xcoff_tdata *xcoff_obj_data;
    struct ecoff_tdata *ecoff_obj_data;
    struct ieee_data_struct *ieee_data;
    struct ieee_ar_data_struct *ieee_ar_data;
    struct srec_data_struct *srec_data;
    struct verilog_data_struct *verilog_data;
    struct ihex_data_struct *ihex_data;
    struct tekhex_data_struct *tekhex_data;
    struct elf_obj_tdata *elf_obj_data;
    struct nlm_obj_tdata *nlm_obj_data;
    struct bout_data_struct *bout_data;
    struct mmo_data_struct *mmo_data;
    struct sun_core_struct *sun_core_data;
    struct sco5_core_struct *sco5_core_data;
    struct trad_core_struct *trad_core_data;
    struct som_data_struct *som_data;
    struct hpux_core_struct *hpux_core_data;
    struct hppabsd_core_struct *hppabsd_core_data;
    struct sgi_core_struct *sgi_core_data;
    struct lynx_core_struct *lynx_core_data;
    struct osf_core_struct *osf_core_data;
    struct cisco_core_struct *cisco_core_data;
    struct versados_data_struct *versados_data;
    struct netbsd_core_struct *netbsd_core_data;
    struct mach_o_data_struct *mach_o_data;
    struct mach_o_fat_data_struct *mach_o_fat_data;
    struct plugin_data_struct *plugin_data;
    struct bfd_pef_data_struct *pef_data;
    struct bfd_pef_xlib_data_struct *pef_xlib_data;
    struct bfd_sym_data_struct *sym_data;
    void *any;
  }
tdata;
/* Used by the application to hold private data. */
void *usrdata;
/* Where all the allocated stuff under this BFD goes. This is a
   struct objalloc *, but we use void * to avoid requiring the inclusion■
```

```
of objalloc.h. */
void *memory;
};

/* See note beside bfd_set_section_userdata. */
static inline bfd_boolean
bfd_set_cacheable (bfd * abfd, bfd_boolean val)
{
   abfd->cacheable = val;
   return TRUE;
}
```

# 2.2 Error reporting

Most BFD functions return nonzero on success (check their individual documentation for precise semantics). On an error, they call bfd\_set\_error to set an error condition that callers can check by calling bfd\_get\_error. If that returns bfd\_error\_system\_call, then check error.

The easiest way to report a BFD error to the user is to use bfd\_perror.

# 2.2.1 Type bfd\_error\_type

The values returned by bfd\_get\_error are defined by the enumerated type bfd\_error\_type.

```
typedef enum bfd_error
 bfd_error_no_error = 0,
 bfd_error_system_call,
 bfd_error_invalid_target,
 bfd_error_wrong_format,
  bfd_error_wrong_object_format,
  bfd_error_invalid_operation,
  bfd_error_no_memory,
  bfd_error_no_symbols,
  bfd_error_no_armap,
  bfd_error_no_more_archived_files,
  bfd_error_malformed_archive,
  bfd_error_missing_dso,
  bfd_error_file_not_recognized,
  bfd_error_file_ambiguously_recognized,
  bfd_error_no_contents,
  bfd_error_nonrepresentable_section,
  bfd_error_no_debug_section,
  bfd_error_bad_value,
  bfd_error_file_truncated,
  bfd_error_file_too_big,
```

```
bfd_error_on_input,
  bfd_error_invalid_error_code
}
bfd_error_type;
```

# 2.2.1.1 bfd\_get\_error

# **Synopsis**

```
bfd_error_type bfd_get_error (void);
```

# Description

Return the current BFD error condition.

# 2.2.1.2 bfd\_set\_error

# **Synopsis**

```
void bfd_set_error (bfd_error_type error_tag, ...);
```

# Description

Set the BFD error condition to be *error\_tag*. If *error\_tag* is bfd\_error\_on\_input, then this function takes two more parameters, the input bfd where the error occurred, and the bfd\_error\_type error.

# 2.2.1.3 bfd\_errmsg

#### Synopsis

```
const char *bfd_errmsg (bfd_error_type error_tag);
```

# Description

Return a string describing the error error\_tag, or the system error if error\_tag is bfd\_error\_system\_call.

# 2.2.1.4 bfd\_perror

#### Synopsis

```
void bfd_perror (const char *message);
```

#### Description

Print to the standard error stream a string describing the last BFD error that occurred, or the last system error if the last BFD error was a system call failure. If message is non-NULL and non-empty, the error string printed is preceded by message, a colon, and a space. It is followed by a newline.

#### 2.2.2 BFD error handler

Some BFD functions want to print messages describing the problem. They call a BFD error handler function. This function may be overridden by the program.

The BFD error handler acts like printf.

```
typedef void (*bfd_error_handler_type) (const char *, ...);
```

# 2.2.2.1 bfd\_set\_error\_handler

# **Synopsis**

```
bfd_error_handler_type bfd_set_error_handler (bfd_error_handler_type);
```

# Description

Set the BFD error handler function. Returns the previous function.

# 2.2.2 bfd\_set\_error\_program\_name

#### **Synopsis**

```
void bfd_set_error_program_name (const char *);
```

#### Description

Set the program name to use when printing a BFD error. This is printed before the error message followed by a colon and space. The string must not be changed after it is passed to this function.

# 2.2.2.3 bfd\_get\_error\_handler

# **Synopsis**

```
bfd_error_handler_type bfd_get_error_handler (void);
```

#### Description

Return the BFD error handler function.

# 2.2.3 BFD assert handler

If BFD finds an internal inconsistency, the bfd assert handler is called with information on the BFD version, BFD source file and line. If this happens, most programs linked against BFD are expected to want to exit with an error, or mark the current BFD operation as failed, so it is recommended to override the default handler, which just calls \_bfd\_error\_handler and continues.

# 2.2.3.1 bfd\_set\_assert\_handler

#### **Synopsis**

```
bfd_assert_handler_type bfd_set_assert_handler (bfd_assert_handler_type);
```

#### Description

Set the BFD assert handler function. Returns the previous function.

# 2.2.3.2 bfd\_get\_assert\_handler

#### **Synopsis**

```
bfd_assert_handler_type bfd_get_assert_handler (void);
```

#### Description

Return the BFD assert handler function.

# 2.3 Miscellaneous

# 2.3.1 Miscellaneous functions

# 2.3.1.1 bfd\_get\_reloc\_upper\_bound

# **Synopsis**

```
long bfd_get_reloc_upper_bound (bfd *abfd, asection *sect);
```

# Description

Return the number of bytes required to store the relocation information associated with section sect attached to bfd abfd. If an error occurs, return -1.

# 2.3.1.2 bfd\_canonicalize\_reloc

# **Synopsis**

```
long bfd_canonicalize_reloc
  (bfd *abfd, asection *sec, arelent **loc, asymbol **syms);
```

# Description

Call the back end associated with the open BFD abfd and translate the external form of the relocation information attached to sec into the internal canonical form. Place the table into memory at loc, which has been preallocated, usually by a call to bfd\_get\_reloc\_upper\_bound. Returns the number of relocs, or -1 on error.

The syms table is also needed for horrible internal magic reasons.

# 2.3.1.3 bfd\_set\_reloc

#### **Synopsis**

```
void bfd_set_reloc
  (bfd *abfd, asection *sec, arelent **rel, unsigned int count);
```

#### Description

Set the relocation pointer and count within section sec to the values rel and count. The argument abfd is ignored.

# 2.3.1.4 bfd\_set\_file\_flags

# Synopsis

```
bfd_boolean bfd_set_file_flags (bfd *abfd, flagword flags);
```

#### Description

Set the flag word in the BFD abfd to the value flags.

Possible errors are:

- bfd\_error\_wrong\_format The target bfd was not of object format.
- bfd\_error\_invalid\_operation The target bfd was open for reading.
- bfd\_error\_invalid\_operation The flag word contained a bit which was not applicable to the type of file. E.g., an attempt was made to set the D\_PAGED bit on a BFD format which does not support demand paging.

# 2.3.1.5 bfd\_get\_arch\_size

# **Synopsis**

```
int bfd_get_arch_size (bfd *abfd);
```

# Description

Returns the normalized architecture address size, in bits, as determined by the object file's format. By normalized, we mean either 32 or 64. For ELF, this information is included in the header. Use bfd\_arch\_bits\_per\_address for number of bits in the architecture address.

#### Returns

Returns the arch size in bits if known, -1 otherwise.

# 2.3.1.6 bfd\_get\_sign\_extend\_vma

# **Synopsis**

```
int bfd_get_sign_extend_vma (bfd *abfd);
```

### Description

Indicates if the target architecture "naturally" sign extends an address. Some architectures implicitly sign extend address values when they are converted to types larger than the size of an address. For instance, bfd\_get\_start\_address() will return an address sign extended to fill a bfd\_vma when this is the case.

#### Returns

Returns 1 if the target architecture is known to sign extend addresses, 0 if the target architecture is known to not sign extend addresses, and -1 otherwise.

# 2.3.1.7 bfd\_set\_start\_address

#### **Synopsis**

```
bfd_boolean bfd_set_start_address (bfd *abfd, bfd_vma vma);
```

# Description

Make vma the entry point of output BFD abfd.

#### Returns

Returns TRUE on success, FALSE otherwise.

# 2.3.1.8 bfd\_get\_gp\_size

# Synopsis

```
unsigned int bfd_get_gp_size (bfd *abfd);
```

#### Description

Return the maximum size of objects to be optimized using the GP register under MIPS ECOFF. This is typically set by the -G argument to the compiler, assembler or linker.

# 2.3.1.9 bfd\_set\_gp\_size

#### **Synopsis**

```
void bfd_set_gp_size (bfd *abfd, unsigned int i);
```

# Description

Set the maximum size of objects to be optimized using the GP register under ECOFF or MIPS ELF. This is typically set by the -G argument to the compiler, assembler or linker.

# 2.3.1.10 bfd\_scan\_vma

# **Synopsis**

```
bfd_vma bfd_scan_vma (const char *string, const char **end, int base);
Description
```

Convert, like strtoul, a numerical expression *string* into a bfd\_vma integer, and return that integer. (Though without as many bells and whistles as strtoul.) The expression is assumed to be unsigned (i.e., positive). If given a *base*, it is used as the base for conversion. A base of 0 causes the function to interpret the string in hex if a leading "0x" or "0X" is found, otherwise in octal if a leading zero is found, otherwise in decimal.

If the value would overflow, the maximum bfd\_vma value is returned.

# 2.3.1.11 bfd\_copy\_private\_header\_data

### **Synopsis**

```
bfd_boolean bfd_copy_private_header_data (bfd *ibfd, bfd *obfd);
```

# Description

Copy private BFD header information from the BFD *ibfd* to the the BFD *obfd*. This copies information that may require sections to exist, but does not require symbol tables. Return true on success, false on error. Possible error returns are:

• bfd\_error\_no\_memory - Not enough memory exists to create private data for obfd.

# 2.3.1.12 bfd\_copy\_private\_bfd\_data

# **Synopsis**

```
bfd_boolean bfd_copy_private_bfd_data (bfd *ibfd, bfd *obfd);
```

### Description

Copy private BFD information from the BFD *ibfd* to the BFD *obfd*. Return TRUE on success, FALSE on error. Possible error returns are:

• bfd\_error\_no\_memory - Not enough memory exists to create private data for obfd.

# 2.3.1.13 bfd\_merge\_private\_bfd\_data

# **Synopsis**

```
bfd_boolean bfd_merge_private_bfd_data (bfd *ibfd, bfd *obfd);
```

#### Description

Merge private BFD information from the BFD *ibfd* to the the output file BFD *obfd* when linking. Return TRUE on success, FALSE on error. Possible error returns are:

• bfd\_error\_no\_memory - Not enough memory exists to create private data for obfd.

# 2.3.1.14 bfd\_set\_private\_flags

Synopsis

```
bfd_boolean bfd_set_private_flags (bfd *abfd, flagword flags);
```

#### Description

Set private BFD flag information in the BFD *abfd*. Return TRUE on success, FALSE on error. Possible error returns are:

 $\bullet$  bfd\_error\_no\_memory - Not enough memory exists to create private data for obfd.

```
#define bfd_set_private_flags(abfd, flags) \
    BFD_SEND (abfd, _bfd_set_private_flags, (abfd, flags))
```

#### 2.3.1.15 Other functions

#### Description

The following functions exist but have not yet been documented.

```
#define bfd_sizeof_headers(abfd, info) \
       BFD_SEND (abfd, _bfd_sizeof_headers, (abfd, info))
#define bfd_find_nearest_line(abfd, sec, syms, off, file, func, line) \
       BFD_SEND (abfd, _bfd_find_nearest_line, \
                 (abfd, syms, sec, off, file, func, line, NULL))
#define bfd_find_nearest_line_discriminator(abfd, sec, syms, off, file, func, \
                                            line, disc) \
       BFD_SEND (abfd, _bfd_find_nearest_line, \
                 (abfd, syms, sec, off, file, func, line, disc))
#define bfd_find_line(abfd, syms, sym, file, line) \
       BFD_SEND (abfd, _bfd_find_line, \
                 (abfd, syms, sym, file, line))
#define bfd_find_inliner_info(abfd, file, func, line) \
       BFD_SEND (abfd, _bfd_find_inliner_info, \
                 (abfd, file, func, line))
#define bfd_debug_info_start(abfd) \
       BFD_SEND (abfd, _bfd_debug_info_start, (abfd))
#define bfd_debug_info_end(abfd) \
       BFD_SEND (abfd, _bfd_debug_info_end, (abfd))
#define bfd_debug_info_accumulate(abfd, section) \
       BFD_SEND (abfd, _bfd_debug_info_accumulate, (abfd, section))
#define bfd_stat_arch_elt(abfd, stat) \
       BFD_SEND (abfd, _bfd_stat_arch_elt,(abfd, stat))
```

```
#define bfd_update_armap_timestamp(abfd) \
       BFD_SEND (abfd, _bfd_update_armap_timestamp, (abfd))
#define bfd_set_arch_mach(abfd, arch, mach)\
       BFD_SEND (abfd, _bfd_set_arch_mach, (abfd, arch, mach))
#define bfd_relax_section(abfd, section, link_info, again) \
       BFD_SEND (abfd, _bfd_relax_section, (abfd, section, link_info, again))
#define bfd_gc_sections(abfd, link_info) \
       BFD_SEND (abfd, _bfd_gc_sections, (abfd, link_info))
#define bfd_lookup_section_flags(link_info, flag_info, section) \
      BFD_SEND (abfd, _bfd_lookup_section_flags, (link_info, flag_info, section))
#define bfd_merge_sections(abfd, link_info) \
       BFD_SEND (abfd, _bfd_merge_sections, (abfd, link_info))
#define bfd_is_group_section(abfd, sec) \
       BFD_SEND (abfd, _bfd_is_group_section, (abfd, sec))
#define bfd_discard_group(abfd, sec) \
       BFD_SEND (abfd, _bfd_discard_group, (abfd, sec))
#define bfd_link_hash_table_create(abfd) \
       BFD_SEND (abfd, _bfd_link_hash_table_create, (abfd))
#define bfd_link_add_symbols(abfd, info) \
       BFD_SEND (abfd, _bfd_link_add_symbols, (abfd, info))
#define bfd_link_just_syms(abfd, sec, info) \
       BFD_SEND (abfd, _bfd_link_just_syms, (sec, info))
#define bfd_final_link(abfd, info) \
       BFD_SEND (abfd, _bfd_final_link, (abfd, info))
#define bfd_free_cached_info(abfd) \
       BFD_SEND (abfd, _bfd_free_cached_info, (abfd))
#define bfd_get_dynamic_symtab_upper_bound(abfd) \
       BFD_SEND (abfd, _bfd_get_dynamic_symtab_upper_bound, (abfd))
#define bfd_print_private_bfd_data(abfd, file)\
       BFD_SEND (abfd, _bfd_print_private_bfd_data, (abfd, file))
#define bfd_canonicalize_dynamic_symtab(abfd, asymbols) \
       BFD_SEND (abfd, _bfd_canonicalize_dynamic_symtab, (abfd, asymbols))
```

# 2.3.1.16 bfd alt mach code

# **Synopsis**

```
bfd_boolean bfd_alt_mach_code (bfd *abfd, int alternative);
```

#### Description

When more than one machine code number is available for the same machine type, this function can be used to switch between the preferred one (alternative ==0) and any others. Currently, only ELF supports this feature, with up to two alternate machine codes.

# 2.3.1.17 bfd\_emul\_get\_maxpagesize

# Synopsis

```
bfd_vma bfd_emul_get_maxpagesize (const char *);
```

#### Description

Returns the maximum page size, in bytes, as determined by emulation.

#### Returns

Returns the maximum page size in bytes for ELF, 0 otherwise.

# 2.3.1.18 bfd\_emul\_set\_maxpagesize

# **Synopsis**

```
void bfd_emul_set_maxpagesize (const char *, bfd_vma);
```

#### Description

For ELF, set the maximum page size for the emulation. It is a no-op for other formats.

# 2.3.1.19 bfd\_emul\_get\_commonpagesize

#### **Synopsis**

```
bfd_vma bfd_emul_get_commonpagesize (const char *);
```

#### Description

Returns the common page size, in bytes, as determined by emulation.

#### Returns

Returns the common page size in bytes for ELF, 0 otherwise.

# 2.3.1.20 bfd\_emul\_set\_commonpagesize

# **Synopsis**

```
void bfd_emul_set_commonpagesize (const char *, bfd_vma);
```

#### Description

For ELF, set the common page size for the emulation. It is a no-op for other formats.

# 2.3.1.21 bfd\_demangle

#### Synopsis

```
char *bfd_demangle (bfd *, const char *, int);
```

#### Description

Wrapper around cplus\_demangle. Strips leading underscores and other such chars that would otherwise confuse the demangler. If passed a g++ v3 ABI mangled name, returns a buffer allocated with malloc holding the demangled name. Returns NULL otherwise and on memory alloc failure.

### 2.3.1.22 struct bfd iovec

#### Description

The struct bfd\_iovec contains the internal file I/O class. Each BFD has an instance of this class and all file I/O is routed through it (it is assumed that the instance implements all methods listed below).

```
struct bfd_iovec
 /* To avoid problems with macros, a "b" rather than "f"
     prefix is prepended to each method name. */
 /* Attempt to read/write NBYTES on ABFD's IOSTREAM storing/fetching
     bytes starting at PTR. Return the number of bytes actually
     transfered (a read past end-of-file returns less than NBYTES),
     or -1 (setting bfd_error) if an error occurs. */
 file_ptr (*bread) (struct bfd *abfd, void *ptr, file_ptr nbytes);
 file_ptr (*bwrite) (struct bfd *abfd, const void *ptr,
                     file_ptr nbytes);
  /* Return the current IOSTREAM file offset, or -1 (setting bfd_error
     if an error occurs. */
 file_ptr (*btell) (struct bfd *abfd);
  /* For the following, on successful completion a value of 0 is returned. ■
     Otherwise, a value of -1 is returned (and bfd_error is set). */
  int (*bseek) (struct bfd *abfd, file_ptr offset, int whence);
  int (*bclose) (struct bfd *abfd);
  int (*bflush) (struct bfd *abfd);
  int (*bstat) (struct bfd *abfd, struct stat *sb);
  /* Mmap a part of the files. ADDR, LEN, PROT, FLAGS and OFFSET are the usual ▮
     mmap parameter, except that LEN and OFFSET do not need to be page
     aligned. Returns (void *)-1 on failure, mmapped address on success. ■
     Also write in MAP_ADDR the address of the page aligned buffer and in
     MAP_LEN the size mapped (a page multiple). Use unmap with MAP_ADDR and
```

# 2.3.1.23 bfd\_get\_mtime

# **Synopsis**

```
long bfd_get_mtime (bfd *abfd);
```

#### Description

Return the file modification time (as read from the file system, or from the archive header for archive members).

# 2.3.1.24 bfd\_get\_size

# **Synopsis**

```
file_ptr bfd_get_size (bfd *abfd);
```

#### Description

Return the file size (as read from file system) for the file associated with BFD abfd.

The initial motivation for, and use of, this routine is not so we can get the exact size of the object the BFD applies to, since that might not be generally possible (archive members for example). It would be ideal if someone could eventually modify it so that such results were guaranteed.

Instead, we want to ask questions like "is this NNN byte sized object I'm about to try read from file offset YYY reasonable?" As as example of where we might do this, some object formats use string tables for which the first sizeof (long) bytes of the table contain the size of the table itself, including the size bytes. If an application tries to read what it thinks is one of these string tables, without some way to validate the size, and for some reason the size is wrong (byte swapping error, wrong location for the string table, etc.), the only clue is likely to be a read error when it tries to read the table, or a "virtual memory exhausted" error when it tries to allocate 15 bazillon bytes of space for the 15 bazillon byte table it is about to read. This function at least allows us to answer the question, "is the size reasonable?".

# 2.3.1.25 bfd\_mmap

# **Synopsis**

```
void *bfd_mmap (bfd *abfd, void *addr, bfd_size_type len,
   int prot, int flags, file_ptr offset,
   void **map_addr, bfd_size_type *map_len);
```

# Description

Return mmap()ed region of the file, if possible and implemented. LEN and OFFSET do not need to be page aligned. The page aligned address and length are written to MAP\_ADDR and MAP\_LEN.

# 2.4 Memory Usage

BFD keeps all of its internal structures in obstacks. There is one obstack per open BFD file, into which the current state is stored. When a BFD is closed, the obstack is deleted, and so everything which has been allocated by BFD for the closing file is thrown away.

BFD does not free anything created by an application, but pointers into bfd structures become invalid on a bfd\_close; for example, after a bfd\_close the vector passed to bfd\_canonicalize\_symtab is still around, since it has been allocated by the application, but the data that it pointed to are lost.

The general rule is to not close a BFD until all operations dependent upon data from the BFD have been completed, or all the data from within the file has been copied. To help with the management of memory, there is a function (bfd\_alloc\_size) which returns the number of bytes in obstacks associated with the supplied BFD. This could be used to select the greediest open BFD, close it to reclaim the memory, perform some operation and reopen the BFD again, to get a fresh copy of the data structures.

# 2.5 Initialization

# 2.5.1 Initialization functions

These are the functions that handle initializing a BFD.

# 2.5.1.1 bfd\_init

### **Synopsis**

void bfd\_init (void);

# Description

This routine must be called before any other BFD function to initialize magical internal data structures.

# 2.6 Sections

The raw data contained within a BFD is maintained through the section abstraction. A single BFD may have any number of sections. It keeps hold of them by pointing to the first; each one points to the next in the list.

Sections are supported in BFD in section.c.

# 2.6.1 Section input

When a BFD is opened for reading, the section structures are created and attached to the BFD.

Each section has a name which describes the section in the outside world—for example, a.out would contain at least three sections, called .text, .data and .bss.

Names need not be unique; for example a COFF file may have several sections named .data. Sometimes a BFD will contain more than the "natural" number of sections. A back end may attach other sections containing constructor data, or an application may add a section (using bfd\_make\_section) to the sections attached to an already open BFD. For example, the linker creates an extra section COMMON for each input file's BFD to hold information about common storage.

The raw data is not necessarily read in when the section descriptor is created. Some targets may leave the data in place until a bfd\_get\_section\_contents call is made. Other back ends may read in all the data at once. For example, an S-record file has to be read once to determine the size of the data. An IEEE-695 file doesn't contain raw data in sections, but data and relocation expressions intermixed, so the data area has to be parsed to get out the data and relocations.

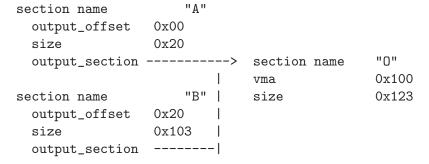
# 2.6.2 Section output

To write a new object style BFD, the various sections to be written have to be created. They are attached to the BFD in the same way as input sections; data is written to the sections using bfd\_set\_section\_contents.

Any program that creates or combines sections (e.g., the assembler and linker) must use the assection fields output\_section and output\_offset to indicate the file sections to which each section must be written. (If the section is being created from scratch, output\_section should probably point to the section itself and output\_offset should probably be zero.)

The data to be written comes from input sections attached (via output\_section pointers) to the output sections. The output section structure can be considered a filter for the input section: the output section determines the vma of the output data and the name, but the input section determines the offset into the output section of the data to be written.

E.g., to create a section "O", starting at 0x100, 0x123 long, containing two subsections, "A" at offset 0x0 (i.e., at vma 0x100) and "B" at offset 0x20 (i.e., at vma 0x120) the asection structures would look like:



# 2.6.3 Link orders

The data within a section is stored in a *link\_order*. These are much like the fixups in gas. The link\_order abstraction allows a section to grow and shrink within itself.

A link\_order knows how big it is, and which is the next link\_order and where the raw data for it is; it also points to a list of relocations which apply to it.

The link\_order is used by the linker to perform relaxing on final code. The compiler creates code which is as big as necessary to make it work without relaxing, and the user can select whether to relax. Sometimes relaxing takes a lot of time. The linker runs around the relocations to see if any are attached to data which can be shrunk, if so it does it on a link\_order by link\_order basis.

# 2.6.4 typedef asection

Here is the section structure:

```
typedef struct bfd_section
 /* The name of the section; the name isn't a copy, the pointer is
    the same as that passed to bfd_make_section. */
 const char *name;
 /* A unique sequence number. */
 int id;
 /* Which section in the bfd; 0..n-1 as sections are created in a bfd. */■
 int index;
 /* The next section in the list belonging to the BFD, or NULL. */
 struct bfd_section *next;
 /* The previous section in the list belonging to the BFD, or NULL. */
 struct bfd_section *prev;
 /* The field flags contains attributes of the section. Some
     flags are read in from the object file, and some are
     synthesized from other information. */
 flagword flags;
#define SEC_NO_FLAGS
                       0x000
 /* Tells the OS to allocate space for this section when loading.
     This is clear for a section containing debug information only. */
#define SEC ALLOC
                       0x001
 /* Tells the OS to load the section from the file when loading.
     This is clear for a .bss section. */
#define SEC_LOAD
                      0x002
 /* The section contains data still to be relocated, so there is
     some relocation information too. */
#define SEC_RELOC
                      0x004
 /* A signal to the OS that the section contains read only data. */
#define SEC_READONLY
                       0x008
  /* The section contains code only. */
#define SEC_CODE
                      0x010
 /* The section contains data only. */
#define SEC_DATA
                      0x020
```

/\* The section will reside in ROM. \*/ #define SEC\_ROM 0x040

/\* The section contains constructor information. This section type is used by the linker to create lists of constructors and destructors used by g++. When a back end sees a symbol which should be used in a constructor list, it creates a new section for the type of name (e.g., \_\_CTOR\_LIST\_\_), attaches the symbol to it, and builds a relocation. To build the lists of constructors, all the linker has to do is catenate all the sections called \_\_CTOR\_LIST\_\_ and relocate the data contained within - exactly the operations it would peform on standard data. \*/

#define SEC\_CONSTRUCTOR 0x080

- /\* The section has contents a data section could be SEC\_ALLOC | SEC\_HAS\_CONTENTS; a debug section could be SEC\_HAS\_CONTENTS \*/ #define SEC\_HAS\_CONTENTS 0x100
- /\* An instruction to the linker to not output the section
   even if it has information which would normally be written. \*/
  #define SEC\_NEVER\_LOAD 0x200
- /\* The section contains thread local data. \*/ #define SEC\_THREAD\_LOCAL 0x400
  - /\* The section has GOT references. This flag is only for the linker, and is currently only used by the elf32-hppa back end. It will be set if global offset table references were detected in this section, which indicate to the linker that the section contains PIC code, and must be handled specially when doing a static link. \*/

#define SEC\_HAS\_GOT\_REF 0x800

- /\* The section contains common symbols (symbols may be defined
   multiple times, the value of a symbol is the amount of
   space it requires, and the largest symbol value is the one
   used). Most targets have exactly one of these (which we
   translate to bfd\_com\_section\_ptr), but ECOFF has two. \*/
  #define SEC\_IS\_COMMON 0x1000
- /\* The section contains only debugging information. For example, this is set for ELF .debug and .stab sections. strip tests this flag to see if a section can be discarded. \*/ #define SEC\_DEBUGGING 0x2000

- /\* The contents of this section are held in memory pointed to
   by the contents field. This is checked by bfd\_get\_section\_contents,
   and the data is retrieved from memory if appropriate. \*/
  #define SEC\_IN\_MEMORY 0x4000
- /\* The contents of this section are to be excluded by the linker for executable and shared objects unless those objects are to be further relocated. \*/ #define SEC\_EXCLUDE 0x8000
- /\* The contents of this section are to be sorted based on the sum of
   the symbol and addend values specified by the associated relocation
   entries. Entries without associated relocation entries will be
   appended to the end of the section in an unspecified order. \*/
  #define SEC\_SORT\_ENTRIES 0x10000
- /\* When linking, duplicate sections of the same name should be
   discarded, rather than being combined into a single section as
   is usually done. This is similar to how common symbols are
   handled. See SEC\_LINK\_DUPLICATES below. \*/
  #define SEC\_LINK\_ONCE 0x20000
- /\* If SEC\_LINK\_ONCE is set, this bitfield describes how the linker
   should handle duplicate sections. \*/
  #define SEC\_LINK\_DUPLICATES 0xc0000
- /\* This value for SEC\_LINK\_DUPLICATES means that duplicate
   sections with the same name should simply be discarded. \*/
  #define SEC\_LINK\_DUPLICATES\_DISCARD 0x0
- /\* This value for SEC\_LINK\_DUPLICATES means that the linker
   should warn if there are any duplicate sections, although
   it should still only link one copy. \*/
  #define SEC\_LINK\_DUPLICATES\_ONE\_ONLY 0x40000
- /\* This value for SEC\_LINK\_DUPLICATES means that the linker
   should warn if any duplicate sections are a different size. \*/
  #define SEC\_LINK\_DUPLICATES\_SAME\_SIZE 0x80000
  - /\* This value for SEC\_LINK\_DUPLICATES means that the linker should warn if any duplicate sections contain different contents. \*/
- #define SEC\_LINK\_DUPLICATES\_SAME\_CONTENTS \
   (SEC\_LINK\_DUPLICATES\_ONE\_ONLY | SEC\_LINK\_DUPLICATES\_SAME\_SIZE)
  - /\* This section was created by the linker as part of dynamic

relocation or other arcane processing. It is skipped when going through the first-pass output, trusting that someone else up the line will take care of it later. \*/
#define SEC\_LINKER\_CREATED 0x100000

- /\* This section should not be subject to garbage collection.
   Also set to inform the linker that this section should not be
   listed in the link map as discarded. \*/
  #define SEC\_KEEP 0x200000
- /\* This section contains "short" data, and should be placed
   "near" the GP. \*/
  #define SEC\_SMALL\_DATA 0x400000
- /\* Attempt to merge identical entities in the section.
   Entity size is given in the entsize field. \*/
  #define SEC\_MERGE 0x800000
- /\* If given with SEC\_MERGE, entities to merge are zero terminated
   strings where entsize specifies character size instead of fixed
   size entries. \*/
  #define SEC\_STRINGS 0x1000000
- /\* This section contains data about section groups. \*/ #define SEC\_GROUP 0x2000000
  - /\* The section is a COFF shared library section. This flag is
    only for the linker. If this type of section appears in
    the input file, the linker must copy it to the output file
    without changing the vma or size. FIXME: Although this
    was originally intended to be general, it really is COFF
    specific (and the flag was renamed to indicate this). It
    might be cleaner to have some more general mechanism to
    allow the back end to control what the linker does with
    sections. \*/

#define SEC\_COFF\_SHARED\_LIBRARY 0x4000000

#define SEC\_COFF\_SHARED 0x8000000

- /\* This input section should be copied to output in reverse order
   as an array of pointers. This is for ELF linker internal use
   only. \*/
  #define SEC\_ELF\_REVERSE\_COPY 0x4000000
- /\* This section contains data which may be shared with other
  executables or shared objects. This is for COFF only. \*/
  - /\* When a section with this flag is being linked, then if the size of

```
the input section is less than a page, it should not cross a page
    boundary. If the size of the input section is one page or more,
    it should be aligned on a page boundary. This is for TI
    TMS320C54X only. */
#define SEC_TIC54X_BLOCK 0x10000000
 /* Conditionally link this section; do not link if there are no
    references found to any symbol in the section. This is for TI
    TMS320C54X only. */
#define SEC_TIC54X_CLINK 0x20000000
  /* Indicate that section has the no read flag set. This happens
    when memory read flag isn't set. */
#define SEC_COFF_NOREAD 0x40000000
 /* End of section flags. */
 /* Some internal packed boolean fields. */
 /* See the vma field. */
 unsigned int user_set_vma : 1;
 /* A mark flag used by some of the linker backends. */
 unsigned int linker_mark : 1;
 /* Another mark flag used by some of the linker backends. Set for
    output sections that have an input section. */
 unsigned int linker_has_input : 1;
  /* Mark flag used by some linker backends for garbage collection. */
 unsigned int gc_mark : 1;
 /* Section compression status. */
 unsigned int compress_status : 2;
#define COMPRESS_SECTION_NONE
#define COMPRESS_SECTION_DONE
#define DECOMPRESS_SECTION_SIZED 2
 /* The following flags are used by the ELF linker. */
 /* Mark sections which have been allocated to segments. */
 unsigned int segment_mark : 1;
 /* Type of sec_info information. */
 unsigned int sec_info_type:3;
#define SEC_INFO_TYPE_NONE
#define SEC_INFO_TYPE_STABS
```

```
#define SEC_INFO_TYPE_MERGE
#define SEC_INFO_TYPE_EH_FRAME 3
#define SEC_INFO_TYPE_JUST_SYMS 4
#define SEC_INFO_TYPE_TARGET
 /st Nonzero if this section uses RELA relocations, rather than REL. st/
 unsigned int use_rela_p:1;
 /* Bits used by various backends. The generic code doesn't touch
     these fields. */
 unsigned int sec_flg0:1;
 unsigned int sec_flg1:1;
 unsigned int sec_flg2:1;
 unsigned int sec_flg3:1;
 unsigned int sec_flg4:1;
 unsigned int sec_flg5:1;
 /* End of internal packed boolean fields. */
 /* The virtual memory address of the section - where it will be
     at run time. The symbols are relocated against this. The
     user_set_vma flag is maintained by bfd; if it's not set, the
     backend can assign addresses (for example, in a.out, where
      the default address for .data is dependent on the specific
      target and various flags). */
 bfd_vma vma;
  /* The load address of the section - where it would be in a
     rom image; really only used for writing section header
      information. */
 bfd_vma lma;
  /* The size of the section in octets, as it will be output.
     Contains a value even if the section has no contents (e.g., the
     size of .bss). */
 bfd_size_type size;
  /* For input sections, the original size on disk of the section, in
     octets. This field should be set for any section whose size is
     changed by linker relaxation. It is required for sections where
     the linker relaxation scheme doesn't cache altered section and
     reloc contents (stabs, eh_frame, SEC_MERGE, some coff relaxing
     targets), and thus the original size needs to be kept to read the
     section multiple times. For output sections, rawsize holds the
     section size calculated on a previous linker relaxation pass. */
 bfd_size_type rawsize;
```

```
/* The compressed size of the section in octets. */
bfd_size_type compressed_size;
/* Relaxation table. */
struct relax_table *relax;
/* Count of used relaxation table entries. */
int relax_count;
/* If this section is going to be output, then this value is the
   offset in *bytes* into the output section of the first byte in the
   input section (byte ==> smallest addressable unit on the
   target). In most cases, if this was going to start at the
   100th octet (8-bit quantity) in the output section, this value
   would be 100. However, if the target byte size is 16 bits
   (bfd_octets_per_byte is "2"), this value would be 50. */
bfd_vma output_offset;
/* The output section through which to map on output. */
struct bfd_section *output_section;
/* The alignment requirement of the section, as an exponent of 2 -
   e.g., 3 aligns to 2^3 (or 8). */
unsigned int alignment_power;
/* If an input section, a pointer to a vector of relocation
   records for the data in this section. */
struct reloc_cache_entry *relocation;
/* If an output section, a pointer to a vector of pointers to
   relocation records for the data in this section. */
struct reloc_cache_entry **orelocation;
/* The number of relocation records in one of the above. */
unsigned reloc_count;
/* Information below is back end specific - and not always used
   or updated. */
/* File position of section data. */
file_ptr filepos;
/* File position of relocation info. */
file_ptr rel_filepos;
```

```
/* File position of line data. */
file_ptr line_filepos;
/* Pointer to data for applications. */
void *userdata;
/* If the SEC_IN_MEMORY flag is set, this points to the actual
   contents. */
unsigned char *contents;
/* Attached line number information. */
alent *lineno;
/* Number of line number records. */
unsigned int lineno_count;
/* Entity size for merging purposes. */
unsigned int entsize;
/* Points to the kept section if this section is a link-once section,
   and is discarded. */
struct bfd_section *kept_section;
/* When a section is being output, this value changes as more
   linenumbers are written out. */
file_ptr moving_line_filepos;
/* What the section number is in the target world. */
int target_index;
void *used_by_bfd;
/* If this is a constructor section then here is a list of the
   relocations created to relocate items within it. */
struct relent_chain *constructor_chain;
/* The BFD which owns the section. */
bfd *owner;
/* A symbol which points at this section only. */
struct bfd_symbol *symbol;
struct bfd_symbol **symbol_ptr_ptr;
/* Early in the link process, map_head and map_tail are used to build
   a list of input sections attached to an output section. Later,
   output sections use these fields for a list of bfd_link_order
   structs. */
```

```
union {
   struct bfd_link_order *link_order;
   struct bfd_section *s;
 } map_head, map_tail;
} asection;
/* Relax table contains information about instructions which can
  be removed by relaxation -- replacing a long address with a
  short address. */
struct relax_table {
 /* Address where bytes may be deleted. */
 bfd_vma addr;
 /* Number of bytes to be deleted. */
 int size;
};
/* Note: the following are provided as inline functions rather than macros
  because not all callers use the return value. A macro implementation
  would use a comma expression, eg: "((ptr)->foo = val, TRUE)" and some
  compilers will complain about comma expressions that have no effect. */
static inline bfd_boolean
bfd_set_section_userdata (bfd * abfd ATTRIBUTE_UNUSED, asection * ptr, void * val)■
 ptr->userdata = val;
 return TRUE;
}
static inline bfd_boolean
bfd_set_section_vma (bfd * abfd ATTRIBUTE_UNUSED, asection * ptr, bfd_vma val)
 ptr->vma = ptr->lma = val;
 ptr->user_set_vma = TRUE;
 return TRUE;
}
static inline bfd_boolean
bfd_set_section_alignment (bfd * abfd ATTRIBUTE_UNUSED, asection * ptr, un-
signed int val)
 ptr->alignment_power = val;
 return TRUE;
}
/* These sections are global, and are managed by BFD. The application
   and target back end are not permitted to change the values in
  these sections. */
```

```
extern asection _bfd_std_section[4];
#define BFD_ABS_SECTION_NAME "*ABS*"
#define BFD_UND_SECTION_NAME "*UND*"
#define BFD_COM_SECTION_NAME "*COM*"
#define BFD_IND_SECTION_NAME "*IND*"
/* Pointer to the common section. */
#define bfd_com_section_ptr (&_bfd_std_section[0])
/* Pointer to the undefined section. */
#define bfd_und_section_ptr (&_bfd_std_section[1])
/* Pointer to the absolute section. */
#define bfd_abs_section_ptr (&_bfd_std_section[2])
/* Pointer to the indirect section. */
#define bfd_ind_section_ptr (&_bfd_std_section[3])
#define bfd_is_und_section(sec) ((sec) == bfd_und_section_ptr)
#define bfd_is_abs_section(sec) ((sec) == bfd_abs_section_ptr)
#define bfd_is_ind_section(sec) ((sec) == bfd_ind_section_ptr)
#define bfd_is_const_section(SEC)
 ( ((SEC) == bfd_abs_section_ptr)
 || ((SEC) == bfd_und_section_ptr)
                                               /
  || ((SEC) == bfd_com_section_ptr)
  || ((SEC) == bfd_ind_section_ptr))
/* Macros to handle insertion and deletion of a bfd's sections.
   only handle the list pointers, ie. do not adjust section_count,
  target_index etc. */
#define bfd_section_list_remove(ABFD, S) \
 do
   {
     asection *_s = S;
      asection *_next = _s->next;
      asection *_prev = _s->prev;
     if (_prev)
       _prev->next = _next;
     else
        (ABFD)->sections = _next;
      if (_next)
       _next->prev = _prev;
        (ABFD)->section_last = _prev;
   }
 while (0)
#define bfd_section_list_append(ABFD, S) \
 do
```

```
{
      asection *_s = S;
      bfd *_abfd = ABFD;
      _s->next = NULL;
      if (_abfd->section_last)
        {
          _s->prev = _abfd->section_last;
          _abfd->section_last->next = _s;
        }
      else
        {
          _s->prev = NULL;
          _abfd->sections = _s;
      _abfd->section_last = _s;
    }
 while (0)
#define bfd_section_list_prepend(ABFD, S) \
 do
   {
      asection *_s = S;
      bfd *_abfd = ABFD;
      _s->prev = NULL;
      if (_abfd->sections)
        {
          _s->next = _abfd->sections;
          _abfd->sections->prev = _s;
        }
      else
        {
          _s->next = NULL;
          _abfd->section_last = _s;
      _abfd->sections = _s;
   }
#define bfd_section_list_insert_after(ABFD, A, S) \
 do
    {
      asection *_a = A;
      asection *_s = S;
      asection *_next = _a->next;
      _s->next = _next;
      _s->prev = _a;
      a\rightarrow next = s;
      if (_next)
        _next->prev = _s;
```

```
else
       (ABFD)->section_last = _s;
   }
 while (0)
#define bfd_section_list_insert_before(ABFD, B, S) \
 do
   {
     asection *_b = B;
     asection *_s = S;
     asection *_prev = _b->prev;
     _s->prev = _prev;
     _s->next = _b;
     b\rightarrow prev = s;
     if (_prev)
       _prev->next = _s;
     else
       (ABFD)->sections = _s;
   }
 while (0)
#define bfd_section_removed_from_list(ABFD, S) \
 ((S)-\text{next} == \text{NULL ? (ABFD)}-\text{section_last }!=(S):(S)-\text{next}-\text{prev }!=(S))
#define BFD_FAKE_SECTION(SEC, FLAGS, SYM, NAME, IDX)
 /* name, id, index, next, prev, flags, user_set_vma,
 { NAME, IDX, O, NULL, NULL, FLAGS, O,
 /* linker_mark, linker_has_input, gc_mark, decompress_status,
                0,
                                1, 0,
 /* segment_mark, sec_info_type, use_rela_p,
                 0,
                               0,
 /* sec_flg0, sec_flg1, sec_flg2, sec_flg3, sec_flg4, sec_flg5,
         0, 0, 0, 0,
 /* vma, lma, size, rawsize, compressed_size, relax, relax_count, */
    0, 0, 0, 0, 0,
 /* output_offset, output_section, alignment_power,
                  &SEC,
    0,
 /* relocation, orelocation, reloc_count, filepos, rel_filepos,
    NULL,
               NULL,
                            0,
                                               Ο,
 /* line_filepos, userdata, contents, lineno, lineno_count,
    Ο,
                 NULL, NULL, O,
```

## 2.6.5 Section prototypes

These are the functions exported by the section handling part of BFD.

## 2.6.5.1 bfd\_section\_list\_clear

## **Synopsis**

```
void bfd_section_list_clear (bfd *);
```

## Description

Clears the section list, and also resets the section count and hash table entries.

# 2.6.5.2 bfd\_get\_section\_by\_name

## **Synopsis**

```
asection *bfd_get_section_by_name (bfd *abfd, const char *name);
```

### Description

Return the most recently created section attached to abfd named name. Return NULL if no such section exists.

## 2.6.5.3 bfd\_get\_next\_section\_by\_name

### **Synopsis**

```
asection *bfd_get_next_section_by_name (asection *sec);
```

### Description

Given sec is a section returned by bfd\_get\_section\_by\_name, return the next most recently created section attached to the same BFD with the same name. Return NULL if no such section exists.

## 2.6.5.4 bfd\_get\_linker\_section

#### Synopsis

```
asection *bfd_get_linker_section (bfd *abfd, const char *name);
```

### Description

Return the linker created section attached to abfd named name. Return NULL if no such section exists.

# 2.6.5.5 bfd\_get\_section\_by\_name\_if

## **Synopsis**

```
asection *bfd_get_section_by_name_if
  (bfd *abfd,
    const char *name,
    bfd_boolean (*func) (bfd *abfd, asection *sect, void *obj),
    void *obj);
```

## Description

Call the provided function func for each section attached to the BFD abfd whose name matches name, passing obj as an argument. The function will be called as if by

```
func (abfd, the_section, obj);
```

It returns the first section for which func returns true, otherwise NULL.

## 2.6.5.6 bfd\_get\_unique\_section\_name

### **Synopsis**

```
char *bfd_get_unique_section_name
  (bfd *abfd, const char *templat, int *count);
```

## Description

Invent a section name that is unique in *abfd* by tacking a dot and a digit suffix onto the original *templat*. If *count* is non-NULL, then it specifies the first number tried as a suffix to generate a unique name. The value pointed to by *count* will be incremented in this case.

## 2.6.5.7 bfd\_make\_section\_old\_way

## **Synopsis**

```
asection *bfd_make_section_old_way (bfd *abfd, const char *name);
```

#### Description

Create a new empty section called *name* and attach it to the end of the chain of sections for the BFD *abfd*. An attempt to create a section with a name which is already in use returns its pointer without changing the section chain.

It has the funny name since this is the way it used to be before it was rewritten....

Possible errors are:

- bfd\_error\_invalid\_operation If output has already started for this BFD.
- bfd\_error\_no\_memory If memory allocation fails.

## 2.6.5.8 bfd\_make\_section\_anyway\_with\_flags

### **Synopsis**

```
asection *bfd_make_section_anyway_with_flags
  (bfd *abfd, const char *name, flagword flags);
```

### Description

Create a new empty section called *name* and attach it to the end of the chain of sections for *abfd*. Create a new section even if there is already a section with that name. Also set the attributes of the new section to the value *flags*.

Return NULL and set bfd\_error on error; possible errors are:

- bfd\_error\_invalid\_operation If output has already started for abfd.
- bfd\_error\_no\_memory If memory allocation fails.

## 2.6.5.9 bfd\_make\_section\_anyway

**Synopsis** 

```
asection *bfd_make_section_anyway (bfd *abfd, const char *name);
```

## Description

Create a new empty section called *name* and attach it to the end of the chain of sections for *abfd*. Create a new section even if there is already a section with that name.

Return NULL and set bfd\_error on error; possible errors are:

- bfd\_error\_invalid\_operation If output has already started for abfd.
- bfd\_error\_no\_memory If memory allocation fails.

## 2.6.5.10 bfd\_make\_section\_with\_flags

**Synopsis** 

```
asection *bfd_make_section_with_flags
  (bfd *, const char *name, flagword flags);
```

### Description

Like bfd\_make\_section\_anyway, but return NULL (without calling bfd\_set\_error ()) without changing the section chain if there is already a section named name. Also set the attributes of the new section to the value flags. If there is an error, return NULL and set bfd\_error.

## 2.6.5.11 bfd\_make\_section

**Synopsis** 

```
asection *bfd_make_section (bfd *, const char *name);
```

### Description

Like bfd\_make\_section\_anyway, but return NULL (without calling bfd\_set\_error ()) without changing the section chain if there is already a section named name. If there is an error, return NULL and set bfd\_error.

## 2.6.5.12 bfd\_set\_section\_flags

Synopsis

```
bfd_boolean bfd_set_section_flags
   (bfd *abfd, asection *sec, flagword flags);
```

#### Description

Set the attributes of the section sec in the BFD abfd to the value flags. Return TRUE on success, FALSE on error. Possible error returns are:

• bfd\_error\_invalid\_operation - The section cannot have one or more of the attributes requested. For example, a .bss section in a.out may not have the SEC\_HAS\_CONTENTS field set.

## 2.6.5.13 bfd\_rename\_section

## **Synopsis**

```
void bfd_rename_section
  (bfd *abfd, asection *sec, const char *newname);
```

### Description

Rename section sec in abfd to newname.

## 2.6.5.14 bfd\_map\_over\_sections

## **Synopsis**

```
void bfd_map_over_sections
  (bfd *abfd,
    void (*func) (bfd *abfd, asection *sect, void *obj),
    void *obj);
```

### Description

Call the provided function func for each section attached to the BFD abfd, passing obj as an argument. The function will be called as if by

```
func (abfd, the_section, obj);
```

This is the preferred method for iterating over sections; an alternative would be to use a loop:

```
asection *p;
for (p = abfd->sections; p != NULL; p = p->next)
  func (abfd, p, ...)
```

## 2.6.5.15 bfd\_sections\_find\_if

## Synopsis

```
asection *bfd_sections_find_if
  (bfd *abfd,
   bfd_boolean (*operation) (bfd *abfd, asection *sect, void *obj),
   void *obj);
```

## Description

Call the provided function operation for each section attached to the BFD abfd, passing obj as an argument. The function will be called as if by

```
operation (abfd, the_section, obj);
```

It returns the first section for which operation returns true.

## 2.6.5.16 bfd\_set\_section\_size

### Synopsis

```
bfd_boolean bfd_set_section_size
   (bfd *abfd, asection *sec, bfd_size_type val);
```

## Description

Set sec to the size val. If the operation is ok, then TRUE is returned, else FALSE.

Possible error returns:

• bfd\_error\_invalid\_operation - Writing has started to the BFD, so setting the size is invalid.

## 2.6.5.17 bfd\_set\_section\_contents

## **Synopsis**

```
bfd_boolean bfd_set_section_contents
  (bfd *abfd, asection *section, const void *data,
    file_ptr offset, bfd_size_type count);
```

#### Description

Sets the contents of the section section in BFD abfd to the data starting in memory at data. The data is written to the output section starting at offset offset for count octets.

Normally TRUE is returned, else FALSE. Possible error returns are:

- bfd\_error\_no\_contents The output section does not have the SEC\_HAS\_CONTENTS attribute, so nothing can be written to it.
- and some more too

This routine is front end to the back end function \_bfd\_set\_section\_contents.

## 2.6.5.18 bfd\_get\_section\_contents

### Synopsis

```
bfd_boolean bfd_get_section_contents
  (bfd *abfd, asection *section, void *location, file_ptr offset,
    bfd_size_type count);
```

## Description

Read data from section in BFD abfd into memory starting at location. The data is read at an offset of offset from the start of the input section, and is read for count bytes.

If the contents of a constructor with the SEC\_CONSTRUCTOR flag set are requested or if the section does not have the SEC\_HAS\_CONTENTS flag set, then the *location* is filled with zeroes. If no errors occur, TRUE is returned, else FALSE.

## 2.6.5.19 bfd\_malloc\_and\_get\_section

#### **Synopsis**

```
bfd_boolean bfd_malloc_and_get_section
  (bfd *abfd, asection *section, bfd_byte **buf);
```

#### Description

Read all data from section in BFD abfd into a buffer, \*buf, malloc'd by this function.

## 2.6.5.20 bfd\_copy\_private\_section\_data

### **Synopsis**

```
bfd_boolean bfd_copy_private_section_data
  (bfd *ibfd, asection *isec, bfd *obfd, asection *osec);
```

### Description

Copy private section information from *isec* in the BFD *ibfd* to the section *osec* in the BFD *obfd*. Return TRUE on success, FALSE on error. Possible error returns are:

• bfd\_error\_no\_memory - Not enough memory exists to create private data for osec.

# 2.6.5.21 bfd\_generic\_is\_group\_section

**Synopsis** 

```
bfd_boolean bfd_generic_is_group_section (bfd *, const asection *sec);
Description
```

Returns TRUE if sec is a member of a group.

## 2.6.5.22 bfd\_generic\_discard\_group

**Synopsis** 

```
bfd_boolean bfd_generic_discard_group (bfd *abfd, asection *group);
```

## Description

Remove all members of group from the output.

# 2.7 Symbols

BFD tries to maintain as much symbol information as it can when it moves information from file to file. BFD passes information to applications though the asymbol structure. When the application requests the symbol table, BFD reads the table in the native form and translates parts of it into the internal format. To maintain more than the information passed to applications, some targets keep some information "behind the scenes" in a structure only the particular back end knows about. For example, the coff back end keeps the original symbol table structure as well as the canonical structure when a BFD is read in. On output, the coff back end can reconstruct the output symbol table so that no information is lost, even information unique to coff which BFD doesn't know or understand. If a coff symbol table were read, but were written through an a.out back end, all the coff specific information would be lost. The symbol table of a BFD is not necessarily read in until a canonicalize request is made. Then the BFD back end fills in a table provided by the application with pointers to the canonical information. To output symbols, the application provides BFD with a table of pointers to pointers to asymbols. This allows applications like the linker to output a symbol as it was read, since the "behind the scenes" information will be still available.

# 2.7.1 Reading symbols

There are two stages to reading a symbol table from a BFD: allocating storage, and the actual reading process. This is an excerpt from an application which reads the symbol table:

```
long storage_needed;
asymbol **symbol_table;
long number_of_symbols;
long i;
storage_needed = bfd_get_symtab_upper_bound (abfd);
if (storage_needed < 0)
   FAIL

if (storage_needed == 0)
   return;</pre>
```

```
symbol_table = xmalloc (storage_needed);
...
number_of_symbols =
   bfd_canonicalize_symtab (abfd, symbol_table);
if (number_of_symbols < 0)
   FAIL

for (i = 0; i < number_of_symbols; i++)
   process_symbol (symbol_table[i]);</pre>
```

All storage for the symbols themselves is in an objalloc connected to the BFD; it is freed when the BFD is closed.

# 2.7.2 Writing symbols

Writing of a symbol table is automatic when a BFD open for writing is closed. The application attaches a vector of pointers to pointers to symbols to the BFD being written, and fills in the symbol count. The close and cleanup code reads through the table provided and performs all the necessary operations. The BFD output code must always be provided with an "owned" symbol: one which has come from another BFD, or one which has been created using bfd\_make\_empty\_symbol. Here is an example showing the creation of a symbol table with only one element:

```
#include "sysdep.h"
#include "bfd.h"
int main (void)
{
  bfd *abfd;
  asymbol *ptrs[2];
  asymbol *new;
  abfd = bfd_openw ("foo", "a.out-sunos-big");
  bfd_set_format (abfd, bfd_object);
  new = bfd_make_empty_symbol (abfd);
  new->name = "dummy_symbol";
  new->section = bfd_make_section_old_way (abfd, ".text");
  new->flags = BSF_GLOBAL;
  new->value = 0x12345;
  ptrs[0] = new;
  ptrs[1] = 0;
  bfd_set_symtab (abfd, ptrs, 1);
  bfd_close (abfd);
  return 0;
}
```

```
./makesym
nm foo
00012345 A dummy_symbol
```

Many formats cannot represent arbitrary symbol information; for instance, the a.out object format does not allow an arbitrary number of sections. A symbol pointing to a section which is not one of .text, .data or .bss cannot be described.

# 2.7.3 Mini Symbols

Mini symbols provide read-only access to the symbol table. They use less memory space, but require more time to access. They can be useful for tools like nm or objdump, which may have to handle symbol tables of extremely large executables.

The bfd\_read\_minisymbols function will read the symbols into memory in an internal form. It will return a void \* pointer to a block of memory, a symbol count, and the size of each symbol. The pointer is allocated using malloc, and should be freed by the caller when it is no longer needed.

The function bfd\_minisymbol\_to\_symbol will take a pointer to a minisymbol, and a pointer to a structure returned by bfd\_make\_empty\_symbol, and return a asymbol structure. The return value may or may not be the same as the value from bfd\_make\_empty\_symbol which was passed in.

# 2.7.4 typedef asymbol

An asymbol has the form:

```
typedef struct bfd_symbol
 /* A pointer to the BFD which owns the symbol. This information
    is necessary so that a back end can work out what additional
    information (invisible to the application writer) is carried
    with the symbol.
    This field is *almost* redundant, since you can use section->owner
    instead, except that some symbols point to the global sections
    bfd_{abs,com,und}_section. This could be fixed by making
    these globals be per-bfd (or per-target-flavor). FIXME. */
  struct bfd *the_bfd; /* Use bfd_asymbol_bfd(sym) to access this field. */
  /* The text of the symbol. The name is left alone, and not copied; the
    application may not alter it. */
  const char *name;
  /* The value of the symbol. This really should be a union of a
    numeric value with a pointer, since some flags indicate that
    a pointer to another symbol is stored here. */
  symvalue value;
  /* Attributes of a symbol. */
```

```
#define BSF_NO_FLAGS
                              0x00
 /* The symbol has local scope; static in C. The value
    is the offset into the section of the data. */
#define BSF_LOCAL
                               (1 << 0)
 /* The symbol has global scope; initialized data in C. The
    value is the offset into the section of the data. */
#define BSF_GLOBAL
                               (1 << 1)
 /* The symbol has global scope and is exported. The value is
    the offset into the section of the data. */
#define BSF_EXPORT
                     BSF_GLOBAL /* No real difference. */
 /* A normal C symbol would be one of:
    BSF_LOCAL, BSF_COMMON, BSF_UNDEFINED or
    BSF_GLOBAL. */
 /* The symbol is a debugging record. The value has an arbitrary
    meaning, unless BSF_DEBUGGING_RELOC is also set. */
#define BSF_DEBUGGING
                              (1 << 2)
 /* The symbol denotes a function entry point. Used in ELF,
    perhaps others someday. */
#define BSF_FUNCTION
                              (1 << 3)
 /* Used by the linker. */
#define BSF_KEEP
                               (1 << 5)
                              (1 << 6)
#define BSF_KEEP_G
 /* A weak global symbol, overridable without warnings by
    a regular global symbol of the same name. */
#define BSF_WEAK
                               (1 << 7)
 /* This symbol was created to point to a section, e.g. ELF's
    STT_SECTION symbols. */
#define BSF_SECTION_SYM
                              (1 << 8)
 /* The symbol used to be a common symbol, but now it is
    allocated. */
#define BSF_OLD_COMMON
                              (1 << 9)
  /* In some files the type of a symbol sometimes alters its
    location in an output file - ie in coff a ISFCN symbol
    which is also C_EXT symbol appears where it was
    declared and not at the end of a section. This bit is set
    by the target BFD part to convey this information. */
```

```
#define BSF_NOT_AT_END
                              (1 << 10)
 /* Signal that the symbol is the label of constructor section. */
#define BSF_CONSTRUCTOR
                              (1 << 11)
 /* Signal that the symbol is a warning symbol. The name is a
    warning. The name of the next symbol is the one to warn about;
    if a reference is made to a symbol with the same name as the next
    symbol, a warning is issued by the linker. */
#define BSF_WARNING
                              (1 << 12)
  /* Signal that the symbol is indirect. This symbol is an indirect
    pointer to the symbol with the same name as the next symbol. */
#define BSF_INDIRECT
                              (1 << 13)
 /* BSF_FILE marks symbols that contain a file name. This is used
    for ELF STT_FILE symbols. */
#define BSF_FILE
                              (1 << 14)
  /* Symbol is from dynamic linking information. */
#define BSF_DYNAMIC
                              (1 << 15)
 /* The symbol denotes a data object. Used in ELF, and perhaps
    others someday. */
#define BSF_OBJECT
                              (1 << 16)
 /* This symbol is a debugging symbol. The value is the offset
    into the section of the data. BSF_DEBUGGING should be set
    as well. */
#define BSF_DEBUGGING_RELOC (1 << 17)</pre>
 /* This symbol is thread local. Used in ELF. */
#define BSF_THREAD_LOCAL
                          (1 << 18)
 /* This symbol represents a complex relocation expression,
    with the expression tree serialized in the symbol name. */
#define BSF_RELC
                              (1 << 19)
 /* This symbol represents a signed complex relocation expression,
    with the expression tree serialized in the symbol name. */
#define BSF_SRELC
                              (1 << 20)
 /* This symbol was created by bfd_get_synthetic_symtab. */
#define BSF_SYNTHETIC
                              (1 << 21)
 /* This symbol is an indirect code object. Unrelated to BSF_INDIRECT.
    The dynamic linker will compute the value of this symbol by
```

```
calling the function that it points to. BSF_FUNCTION must
     also be also set. */
#define BSF_GNU_INDIRECT_FUNCTION (1 << 22)</pre>
  /* This symbol is a globally unique data object. The dynamic linker
     will make sure that in the entire process there is just one symbol
     with this name and type in use. BSF_OBJECT must also be set. */
#define BSF_GNU_UNIQUE
                               (1 << 23)
 flagword flags;
 /* A pointer to the section to which this symbol is
     relative. This will always be non NULL, there are special
     sections for undefined and absolute symbols. */
 struct bfd_section *section;
 /* Back end special data. */
 union
   {
     void *p;
     bfd_vma i;
   }
 udata;
asymbol;
```

## 2.7.5 Symbol handling functions

## 2.7.5.1 bfd\_get\_symtab\_upper\_bound

## Description

Return the number of bytes required to store a vector of pointers to asymbols for all the symbols in the BFD *abfd*, including a terminal NULL pointer. If there are no symbols in the BFD, then return 0. If an error occurs, return -1.

```
#define bfd_get_symtab_upper_bound(abfd) \
    BFD_SEND (abfd, _bfd_get_symtab_upper_bound, (abfd))
```

## 2.7.5.2 bfd\_is\_local\_label

Synopsis

```
bfd_boolean bfd_is_local_label (bfd *abfd, asymbol *sym);
```

#### Description

Return TRUE if the given symbol sym in the BFD abfd is a compiler generated local label, else return FALSE.

## 2.7.5.3 bfd\_is\_local\_label\_name

Synopsis

bfd\_boolean bfd\_is\_local\_label\_name (bfd \*abfd, const char \*name);

### Description

Return TRUE if a symbol with the name name in the BFD abfd is a compiler generated local label, else return FALSE. This just checks whether the name has the form of a local label.

```
#define bfd_is_local_label_name(abfd, name) \
    BFD_SEND (abfd, _bfd_is_local_label_name, (abfd, name))
```

## 2.7.5.4 bfd\_is\_target\_special\_symbol

## **Synopsis**

```
bfd_boolean bfd_is_target_special_symbol (bfd *abfd, asymbol *sym);
```

## Description

Return TRUE iff a symbol sym in the BFD abfd is something special to the particular target represented by the BFD. Such symbols should normally not be mentioned to the user.

```
#define bfd_is_target_special_symbol(abfd, sym) \
   BFD_SEND (abfd, _bfd_is_target_special_symbol, (abfd, sym))
```

## 2.7.5.5 bfd\_canonicalize\_symtab

## Description

Read the symbols from the BFD *abfd*, and fills in the vector *location* with pointers to the symbols and a trailing NULL. Return the actual number of symbol pointers, not including the NULL.

```
#define bfd_canonicalize_symtab(abfd, location) \
   BFD_SEND (abfd, _bfd_canonicalize_symtab, (abfd, location))
```

## 2.7.5.6 bfd\_set\_symtab

### **Synopsis**

```
bfd_boolean bfd_set_symtab
   (bfd *abfd, asymbol **location, unsigned int count);
```

### Description

Arrange that when the output BFD abfd is closed, the table location of count pointers to symbols will be written.

## 2.7.5.7 bfd\_print\_symbol\_vandf

## **Synopsis**

```
void bfd_print_symbol_vandf (bfd *abfd, void *file, asymbol *symbol);
```

#### Description

Print the value and flags of the symbol supplied to the stream file.

## 2.7.5.8 bfd\_make\_empty\_symbol

## Description

Create a new asymbol structure for the BFD abfd and return a pointer to it.

This routine is necessary because each back end has private information surrounding the asymbol. Building your own asymbol and pointing to it will not create the private information, and will cause problems later on.

```
#define bfd_make_empty_symbol(abfd) \
BFD_SEND (abfd, _bfd_make_empty_symbol, (abfd))
```

## 2.7.5.9 \_bfd\_generic\_make\_empty\_symbol

### **Synopsis**

```
asymbol *_bfd_generic_make_empty_symbol (bfd *);
```

## Description

Create a new asymbol structure for the BFD abfd and return a pointer to it. Used by core file routines, binary back-end and anywhere else where no private info is needed.

## 2.7.5.10 bfd\_make\_debug\_symbol

### Description

Create a new asymbol structure for the BFD *abfd*, to be used as a debugging symbol. Further details of its use have yet to be worked out.

```
#define bfd_make_debug_symbol(abfd,ptr,size) \
BFD_SEND (abfd, _bfd_make_debug_symbol, (abfd, ptr, size))
```

## 2.7.5.11 bfd\_decode\_symclass

### Description

Return a character corresponding to the symbol class of symbol, or '?' for an unknown class.

## **Synopsis**

```
int bfd_decode_symclass (asymbol *symbol);
```

# 2.7.5.12 bfd\_is\_undefined\_symclass

### Description

Returns non-zero if the class symbol returned by bfd\_decode\_symclass represents an undefined symbol. Returns zero otherwise.

## **Synopsis**

```
bfd_boolean bfd_is_undefined_symclass (int symclass);
```

## 2.7.5.13 bfd\_symbol\_info

## Description

Fill in the basic info about symbol that nm needs. Additional info may be added by the back-ends after calling this function.

## **Synopsis**

```
void bfd_symbol_info (asymbol *symbol, symbol_info *ret);
2.7.5.14 bfd_copy_private_symbol_data
Synopsis
    bfd_boolean bfd_copy_private_symbol_data
        (bfd *ibfd, asymbol *isym, bfd *obfd, asymbol *osym);
```

### Description

Copy private symbol information from *isym* in the BFD *ibfd* to the symbol *osym* in the BFD *obfd*. Return TRUE on success, FALSE on error. Possible error returns are:

• bfd\_error\_no\_memory - Not enough memory exists to create private data for osec.

## 2.8 Archives

## Description

An archive (or library) is just another BFD. It has a symbol table, although there's not much a user program will do with it.

The big difference between an archive BFD and an ordinary BFD is that the archive doesn't have sections. Instead it has a chain of BFDs that are considered its contents. These BFDs can be manipulated like any other. The BFDs contained in an archive opened for reading will all be opened for reading. You may put either input or output BFDs into an archive opened for output; they will be handled correctly when the archive is closed.

Use bfd\_openr\_next\_archived\_file to step through the contents of an archive opened for input. You don't have to read the entire archive if you don't want to! Read it until you find what you want.

A BFD returned by bfd\_openr\_next\_archived\_file can be closed manually with bfd\_close. If you do not close it, then a second iteration through the members of an archive may return the same BFD. If you close the archive BFD, then all the member BFDs will automatically be closed as well.

Archive contents of output BFDs are chained through the archive\_next pointer in a BFD. The first one is findable through the archive\_head slot of the archive. Set it with bfd\_set\_archive\_head (q.v.). A given BFD may be in only one open output archive at a time.

As expected, the BFD archive code is more general than the archive code of any given environment. BFD archives may contain files of different formats (e.g., a.out and coff) and even different architectures. You may even place archives recursively into archives!

This can cause unexpected confusion, since some archive formats are more expressive than others. For instance, Intel COFF archives can preserve long filenames; SunOS a.out archives cannot. If you move a file from the first to the second format and back again, the filename may be truncated. Likewise, different a.out environments have different conventions as to how they truncate filenames, whether they preserve directory names in filenames, etc. When interoperating with native tools, be sure your files are homogeneous.

Beware: most of these formats do not react well to the presence of spaces in filenames. We do the best we can, but can't always handle this case due to restrictions in the format of archives. Many Unix utilities are braindead in regards to spaces and such in filenames anyway, so this shouldn't be much of a restriction.

Archives are supported in BFD in archive.c.

## 2.8.1 Archive functions

## 2.8.1.1 bfd\_get\_next\_mapent

## **Synopsis**

```
symindex bfd_get_next_mapent
   (bfd *abfd, symindex previous, carsym **sym);
```

### Description

Step through archive abfd's symbol table (if it has one). Successively update sym with the next symbol's information, returning that symbol's (internal) index into the symbol table.

Supply BFD\_NO\_MORE\_SYMBOLS as the *previous* entry to get the first one; returns BFD\_NO\_MORE\_SYMBOLS when you've already got the last one.

A carsym is a canonical archive symbol. The only user-visible element is its name, a null-terminated string.

## 2.8.1.2 bfd\_set\_archive\_head

#### **Synopsis**

```
bfd_boolean bfd_set_archive_head (bfd *output, bfd *new_head);
```

### Description

Set the head of the chain of BFDs contained in the archive output to new\_head.

## 2.8.1.3 bfd\_openr\_next\_archived\_file

## Synopsis

```
bfd *bfd_openr_next_archived_file (bfd *archive, bfd *previous);
```

#### Description

Provided a BFD, archive, containing an archive and NULL, open an input BFD on the first contained element and returns that. Subsequent calls should pass the archive and the previous return value to return a created BFD to the next contained element. NULL is returned when there are no more.

## 2.9 File formats

A format is a BFD concept of high level file contents type. The formats supported by BFD are:

• bfd\_object

The BFD may contain data, symbols, relocations and debug info.

• bfd\_archive

The BFD contains other BFDs and an optional index.

• bfd\_core

The BFD contains the result of an executable core dump.

## 2.9.1 File format functions

## 2.9.1.1 bfd\_check\_format

### **Synopsis**

```
bfd_boolean bfd_check_format (bfd *abfd, bfd_format format);
```

### Description

Verify if the file attached to the BFD *abfd* is compatible with the format *format* (i.e., one of bfd\_object, bfd\_archive or bfd\_core).

If the BFD has been set to a specific target before the call, only the named target and format combination is checked. If the target has not been set, or has been set to default, then all the known target backends is interrogated to determine a match. If the default target matches, it is used. If not, exactly one target must recognize the file, or an error results.

The function returns TRUE on success, otherwise FALSE with one of the following error codes:

- bfd\_error\_invalid\_operation if format is not one of bfd\_object, bfd\_archive or bfd\_core.
- bfd\_error\_system\_call if an error occurred during a read even some file mismatches can cause bfd\_error\_system\_calls.
- file\_not\_recognised none of the backends recognised the file format.
- bfd\_error\_file\_ambiguously\_recognized more than one backend recognised the file format.

## 2.9.1.2 bfd\_check\_format\_matches

### **Synopsis**

```
bfd_boolean bfd_check_format_matches
   (bfd *abfd, bfd_format format, char ***matching);
```

#### Description

Like bfd\_check\_format, except when it returns FALSE with bfd\_errno set to bfd\_error\_file\_ambiguously\_recognized. In that case, if *matching* is not NULL, it will be filled in with a NULL-terminated list of the names of the formats that matched, allocated with malloc. Then the user may choose a format and try again.

When done with the list that matching points to, the caller should free it.

## 2.9.1.3 bfd\_set\_format

## **Synopsis**

```
bfd_boolean bfd_set_format (bfd *abfd, bfd_format format);
```

### Description

This function sets the file format of the BFD abfd to the format format. If the target set in the BFD does not support the format requested, the format is invalid, or the BFD is not open for writing, then an error occurs.

## 2.9.1.4 bfd\_format\_string

## **Synopsis**

```
const char *bfd_format_string (bfd_format format);
```

## Description

Return a pointer to a const string invalid, object, archive, core, or unknown, depending upon the value of *format*.

## 2.10 Relocations

BFD maintains relocations in much the same way it maintains symbols: they are left alone until required, then read in en-masse and translated into an internal form. A common routine bfd\_perform\_relocation acts upon the canonical form to do the fixup.

Relocations are maintained on a per section basis, while symbols are maintained on a per BFD basis.

All that a back end has to do to fit the BFD interface is to create a struct reloc\_cache\_entry for each relocation in a particular section, and fill in the right bits of the structures.

# 2.10.1 typedef arelent

This is the structure of a relocation entry:

```
typedef enum bfd_reloc_status
 /* No errors detected. */
 bfd_reloc_ok,
 /* The relocation was performed, but there was an overflow. */
 bfd_reloc_overflow,
 /* The address to relocate was not within the section supplied. */
 bfd_reloc_outofrange,
 /* Used by special functions. */
 bfd_reloc_continue,
  /* Unsupported relocation size requested. */
 bfd_reloc_notsupported,
 /* Unused. */
 bfd_reloc_other,
  /* The symbol to relocate against was undefined. */
 bfd_reloc_undefined,
  /* The relocation was performed, but may not be ok - presently
     generated only when linking i960 coff files with i960 b.out
     symbols. If this type is returned, the error_message argument
```

```
to bfd_perform_relocation will be set. */
bfd_reloc_dangerous
}
bfd_reloc_status_type;

typedef struct reloc_cache_entry
{
   /* A pointer into the canonical table of pointers. */
   struct bfd_symbol **sym_ptr_ptr;

   /* offset in section. */
   bfd_size_type address;

   /* addend for relocation value. */
   bfd_vma addend;

   /* Pointer to how to perform the required relocation. */
   reloc_howto_type *howto;
}
arelent;
```

## Description

Here is a description of each of the fields within an arelent:

## • sym\_ptr\_ptr

The symbol table pointer points to a pointer to the symbol associated with the relocation request. It is the pointer into the table returned by the back end's canonicalize\_symtab action. See Section 2.7 [Symbols], page 41. The symbol is referenced through a pointer to a pointer so that tools like the linker can fix up all the symbols of the same name by modifying only one pointer. The relocation routine looks in the symbol and uses the base of the section the symbol is attached to and the value of the symbol as the initial relocation offset. If the symbol pointer is zero, then the section provided is looked up.

### • address

The address field gives the offset in bytes from the base of the section data which owns the relocation record to the first byte of relocatable information. The actual data relocated will be relative to this point; for example, a relocation type which modifies the bottom two bytes of a four byte word would not touch the first byte pointed to in a big endian world.

## • addend

The addend is a value provided by the back end to be added (!) to the relocation offset. Its interpretation is dependent upon the howto. For example, on the 68k the code:

```
char foo[];
main()
{
```

unlk fp rts

```
return foo[0x12345678];
}
Could be compiled into:
linkw fp,#-4
moveb @#12345678,d0
extbl d0
```

This could create a reloc pointing to foo, but leave the offset in the data, something like:

```
RELOCATION RECORDS FOR [.text]:
offset
                   value
         type
00000006 32
                    _foo
00000000 4e56 fffc
                             ; linkw fp,#-4
00000004 1039 1234 5678
                             ; moveb @#12345678,d0
0000000a 49c0
                             ; extbl d0
0000000c 4e5e
                             ; unlk fp
0000000e 4e75
                             ; rts
```

Using coff and an 88k, some instructions don't have enough space in them to represent the full address range, and pointers have to be loaded in two parts. So you'd get something like:

```
or.u r13,r0,hi16(_foo+0x12345678)
ld.b r2,r13,lo16(_foo+0x12345678)
jmp r1
```

This should create two relocs, both pointing to \_foo, and with 0x12340000 in their addend field. The data would consist of:

```
RELOCATION RECORDS FOR [.text]:

offset type value

000000002 HVRT16 _foo+0x12340000

00000006 LVRT16 _foo+0x12340000

00000000 5da05678 ; or.u r13,r0,0x5678

00000004 1c4d5678 ; ld.b r2,r13,0x5678

00000008 f400c001 ; jmp r1
```

The relocation routine digs out the value from the data, adds it to the addend to get the original offset, and then adds the value of \_foo. Note that all 32 bits have to be kept around somewhere, to cope with carry from bit 15 to bit 16.

One further example is the sparc and the a.out format. The sparc has a similar problem to the 88k, in that some instructions don't have room for an entire offset, but on the sparc the parts are created in odd sized lumps. The designers of the a.out format chose to not use the data within the section for storing part of the offset; all the offset is kept within the reloc. Anything in the data should be ignored.

```
save %sp,-112,%sp
sethi %hi(_foo+0x12345678),%g2
```

```
ldsb [%g2+%lo(_foo+0x12345678)],%i0
ret
restore
```

Both relocs contain a pointer to foo, and the offsets contain junk.

```
RELOCATION RECORDS FOR [.text]:
offset
         type
                   value
00000004 HI22
                   _foo+0x12345678
00000008 L010
                   _foo+0x12345678
                      ; save %sp,-112,%sp
00000000 9de3bf90
                      ; sethi %hi(_foo+0), %g2
00000004 05000000
                      ; ldsb [%g2+%lo(_foo+0)],%i0
00000008 f048a000
0000000c 81c7e008
                      ; ret
00000010 81e80000
                      ; restore
```

#### • howto

The howto field can be imagined as a relocation instruction. It is a pointer to a structure which contains information on what to do with all of the other information in the reloc record and data section. A back end would normally have a relocation instruction set and turn relocations into pointers to the correct structure on input - but it would be possible to create each howto field on demand.

## 2.10.1.1 enum complain\_overflow

Indicates what sort of overflow checking should be done when performing a relocation.

```
enum complain_overflow
{
    /* Do not complain on overflow. */
    complain_overflow_dont,

    /* Complain if the value overflows when considered as a signed
        number one bit larger than the field. ie. A bitfield of N bits
        is allowed to represent -2**n to 2**n-1. */
    complain_overflow_bitfield,

    /* Complain if the value overflows when considered as a signed
        number. */
    complain_overflow_signed,

    /* Complain if the value overflows when considered as an
        unsigned number. */
    complain_overflow_unsigned
};
```

## 2.10.1.2 reloc\_howto\_type

The reloc\_howto\_type is a structure which contains all the information that libbfd needs to know to tie up a back end's data.

```
struct bfd_symbol;
                              /* Forward declaration. */
struct reloc_howto_struct
 /* The type field has mainly a documentary use - the back end can
     do what it wants with it, though normally the back end's
     external idea of what a reloc number is stored
      in this field. For example, a PC relative word relocation
      in a coff environment has the type 023 - because that's
      what the outside world calls a R_PCRWORD reloc. */
 unsigned int type;
 /* The value the final relocation is shifted right by. This drops
     unwanted data from the relocation. */
 unsigned int rightshift;
 /* The size of the item to be relocated. This is *not* a
     power-of-two measure. To get the number of bytes operated
     on by a type of relocation, use bfd_get_reloc_size. */
 int size:
  /* The number of bits in the item to be relocated. This is used
      when doing overflow checking. */
 unsigned int bitsize;
  /* The relocation is relative to the field being relocated. */
 bfd_boolean pc_relative;
 /* The bit position of the reloc value in the destination.
      The relocated value is left shifted by this amount. */
 unsigned int bitpos;
 /* What type of overflow error should be checked for when
    relocating.
                 */
 enum complain_overflow complain_on_overflow;
 /* If this field is non null, then the supplied function is
    called rather than the normal function. This allows really
    strange relocation methods to be accommodated (e.g., i960 callj
    instructions). */
 bfd_reloc_status_type (*special_function)
    (bfd *, arelent *, struct bfd_symbol *, void *, asection *,
    bfd *, char **);
 /* The textual name of the relocation type. */
 char *name;
```

/\* Some formats record a relocation addend in the section contents rather than with the relocation. For ELF formats this is the distinction between USE\_REL and USE\_RELA (though the code checks for USE\_REL == 1/0). The value of this field is TRUE if the addend is recorded with the section contents; when performing a partial link (ld -r) the section contents (the data) will be modified. The value of this field is FALSE if addends are recorded with the relocation (in arelent.addend); when performing a partial link the relocation will be modified. All relocations for all ELF USE\_RELA targets should set this field to FALSE (values of TRUE should be looked on with suspicion). However, the converse is not true: not all relocations of all ELF USE\_REL targets set this field to TRUE. Why this is so is peculiar to each particular target. For relocs that aren't used in partial links (e.g. GOT stuff) it doesn't matter what this is set to. \*/ bfd\_boolean partial\_inplace;

/\* src\_mask selects the part of the instruction (or data) to be used
in the relocation sum. If the target relocations don't have an
addend in the reloc, eg. ELF USE\_REL, src\_mask will normally equal
dst\_mask to extract the addend from the section contents. If
relocations do have an addend in the reloc, eg. ELF USE\_RELA, this
field should be zero. Non-zero values for ELF USE\_RELA targets are
bogus as in those cases the value in the dst\_mask part of the
section contents should be treated as garbage. \*/
bfd\_vma src\_mask;

```
/* dst_mask selects which parts of the instruction (or data) are
  replaced with a relocated value. */
bfd_vma dst_mask;
```

/\* When some formats create PC relative instructions, they leave
 the value of the pc of the place being relocated in the offset
 slot of the instruction, so that a PC relative relocation can
 be made just by adding in an ordinary offset (e.g., sun3 a.out).
 Some formats leave the displacement part of an instruction
 empty (e.g., m88k bcs); this flag signals the fact. \*/
 bfd\_boolean pcrel\_offset;
};

#### 2.10.1.3 The HOWTO Macro

#### Description

The HOWTO define is horrible and will go away.

```
#define HOWTO(C, R, S, B, P, BI, O, SF, NAME, INPLACE, MASKSRC, MASKDST, PC) \■ { (unsigned) C, R, S, B, P, BI, O, SF, NAME, INPLACE, MASKSRC, MASKDST, PC }■
```

## Description

And will be replaced with the totally magic way. But for the moment, we are compatible, so do it this way.

### Description

This is used to fill in an empty how entry in an array.

```
#define EMPTY_HOWTO(C) \
HOWTO ((C), 0, 0, 0, FALSE, 0, complain_overflow_dont, NULL, \
NULL, FALSE, 0, 0, FALSE)
```

## Description

Helper routine to turn a symbol into a relocation value.

```
#define HOWTO_PREPARE(relocation, symbol)
{
   if (symbol != NULL)
   {
      if (bfd_is_com_section (symbol->section))
      {
        relocation = 0;
      }
      else
      {
        relocation = symbol->value;
      }
   }
}
```

## 2.10.1.4 bfd\_get\_reloc\_size

### **Synopsis**

```
unsigned int bfd_get_reloc_size (reloc_howto_type *);
```

#### Description

For a reloc\_howto\_type that operates on a fixed number of bytes, this returns the number of bytes operated on.

## 2.10.1.5 arelent\_chain

#### Description

How relocs are tied together in an asection:

```
typedef struct relent_chain
{
  arelent relent;
  struct relent_chain *next;
```

```
}
arelent_chain;
```

### 2.10.1.6 bfd check overflow

## **Synopsis**

```
bfd_reloc_status_type bfd_check_overflow
  (enum complain_overflow how,
    unsigned int bitsize,
    unsigned int rightshift,
    unsigned int addrsize,
    bfd_vma relocation);
```

### Description

Perform overflow checking on *relocation* which has *bitsize* significant bits and will be shifted right by *rightshift* bits, on a machine with addresses containing *addrsize* significant bits. The result is either of bfd\_reloc\_ok or bfd\_reloc\_overflow.

## 2.10.1.7 bfd\_perform\_relocation

## **Synopsis**

```
bfd_reloc_status_type bfd_perform_relocation
  (bfd *abfd,
    arelent *reloc_entry,
    void *data,
    asection *input_section,
    bfd *output_bfd,
    char **error_message);
```

### Description

If output\_bfd is supplied to this function, the generated image will be relocatable; the relocations are copied to the output file after they have been changed to reflect the new state of the world. There are two ways of reflecting the results of partial linkage in an output file: by modifying the output data in place, and by modifying the relocation record. Some native formats (e.g., basic a.out and basic coff) have no way of specifying an addend in the relocation type, so the addend has to go in the output data. This is no big deal since in these formats the output data slot will always be big enough for the addend. Complex reloc types with addends were invented to solve just this problem. The error\_message argument is set to an error message if this return bfd\_reloc\_dangerous.

### 2.10.1.8 bfd\_install\_relocation

## **Synopsis**

```
bfd_reloc_status_type bfd_install_relocation
  (bfd *abfd,
    arelent *reloc_entry,
    void *data, bfd_vma data_start,
    asection *input_section,
    char **error_message);
```

### Description

This looks remarkably like bfd\_perform\_relocation, except it does not expect that the section contents have been filled in. I.e., it's suitable for use when creating, rather than applying a relocation.

For now, this function should be considered reserved for the assembler.

## 2.10.2 The howto manager

When an application wants to create a relocation, but doesn't know what the target machine might call it, it can find out by using this bit of code.

## 2.10.2.1 bfd\_reloc\_code\_type

## Description

The insides of a reloc code. The idea is that, eventually, there will be one enumerator for every type of relocation we ever do. Pass one of these values to bfd\_reloc\_type\_lookup, and it'll return a howto pointer.

This does mean that the application must determine the correct enumerator value; you can't get a howto pointer from a random set of attributes.

Here are the possible values for enum bfd\_reloc\_code\_real:

```
BFD_RELOC_64
BFD_RELOC_32
BFD_RELOC_26
BFD_RELOC_24
BFD_RELOC_16
BFD_RELOC_14
BFD_RELOC_8
```

Basic absolute relocations of N bits.

```
BFD_RELOC_64_PCREL
BFD_RELOC_32_PCREL
BFD_RELOC_24_PCREL
BFD_RELOC_16_PCREL
BFD_RELOC_12_PCREL
BFD_RELOC_8_PCREL
```

PC-relative relocations. Sometimes these are relative to the address of the relocation itself; sometimes they are relative to the start of the section containing the relocation. It depends on the specific target.

The 24-bit relocation is used in some Intel 960 configurations.

#### BFD\_RELOC\_32\_SECREL

Section relative relocations. Some targets need this for DWARF2.

```
BFD_RELOC_32_GOT_PCREL
BFD_RELOC_16_GOT_PCREL
BFD_RELOC_8_GOT_PCREL
BFD_RELOC_32_GOTOFF
BFD_RELOC_16_GOTOFF
```

```
BFD_RELOC_LO16_GOTOFF
BFD_RELOC_HI16_GOTOFF
BFD_RELOC_HI16_S_GOTOFF
BFD_RELOC_8_GOTOFF
BFD_RELOC_64_PLT_PCREL
BFD_RELOC_32_PLT_PCREL
BFD_RELOC_24_PLT_PCREL
BFD_RELOC_16_PLT_PCREL
BFD_RELOC_8_PLT_PCREL
BFD_RELOC_64_PLTOFF
BFD_RELOC_32_PLTOFF
BFD_RELOC_16_PLTOFF
BFD_RELOC_LO16_PLTOFF
BFD_RELOC_HI16_PLT0FF
BFD_RELOC_HI16_S_PLTOFF
BFD_RELOC_8_PLTOFF
    For ELF.
```

BFD\_RELOC\_SIZE32 BFD\_RELOC\_SIZE64

Size relocations.

BFD\_RELOC\_68K\_GLOB\_DAT BFD\_RELOC\_68K\_JMP\_SLOT BFD\_RELOC\_68K\_RELATIVE BFD\_RELOC\_68K\_TLS\_GD32 BFD\_RELOC\_68K\_TLS\_GD16 BFD\_RELOC\_68K\_TLS\_GD8 BFD\_RELOC\_68K\_TLS\_LDM32 BFD\_RELOC\_68K\_TLS\_LDM16 BFD\_RELOC\_68K\_TLS\_LDM8 BFD\_RELOC\_68K\_TLS\_LD032 BFD\_RELOC\_68K\_TLS\_LD016 BFD\_RELOC\_68K\_TLS\_LD08 BFD\_RELOC\_68K\_TLS\_IE32 BFD\_RELOC\_68K\_TLS\_IE16 BFD\_RELOC\_68K\_TLS\_IE8 BFD\_RELOC\_68K\_TLS\_LE32 BFD\_RELOC\_68K\_TLS\_LE16 BFD\_RELOC\_68K\_TLS\_LE8 Relocations used by 68K ELF.

BFD\_RELOC\_32\_BASEREL BFD\_RELOC\_16\_BASEREL BFD\_RELOC\_LO16\_BASEREL BFD\_RELOC\_HI16\_BASEREL BFD\_RELOC\_HI16\_S\_BASEREL BFD\_RELOC\_8\_BASEREL

#### BFD\_RELOC\_RVA

Linkage-table relative.

#### BFD\_RELOC\_8\_FFnn

Absolute 8-bit relocation, but used to form an address like 0xFFnn.

BFD\_RELOC\_32\_PCREL\_S2

BFD\_RELOC\_16\_PCREL\_S2

BFD\_RELOC\_23\_PCREL\_S2

These PC-relative relocations are stored as word displacements – i.e., byte displacements shifted right two bits. The 30-bit word displacement (<<32\_PCREL\_S2>> – 32 bits, shifted 2) is used on the SPARC. (SPARC tools generally refer to this as <<WDISP30>>.) The signed 16-bit displacement is used on the MIPS, and the 23-bit displacement is used on the Alpha.

BFD\_RELOC\_HI22

BFD\_RELOC\_LO10

High 22 bits and low 10 bits of 32-bit value, placed into lower bits of the target word. These are used on the SPARC.

BFD\_RELOC\_GPREL16

BFD\_RELOC\_GPREL32

For systems that allocate a Global Pointer register, these are displacements off that register. These relocation types are handled specially, because the value the register will have is decided relatively late.

## BFD\_RELOC\_I960\_CALLJ

Reloc types used for i960/b.out.

BFD\_RELOC\_NONE

BFD\_RELOC\_SPARC\_WDISP22

BFD\_RELOC\_SPARC22

BFD\_RELOC\_SPARC13

BFD\_RELOC\_SPARC\_GOT10

BFD\_RELOC\_SPARC\_GOT13

BFD\_RELOC\_SPARC\_GOT22

BFD\_RELOC\_SPARC\_PC10

BFD\_RELOC\_SPARC\_PC22

BFD\_RELOC\_SPARC\_WPLT30

BFD\_RELOC\_SPARC\_COPY

BFD\_RELOC\_SPARC\_GLOB\_DAT

BFD\_RELOC\_SPARC\_JMP\_SLOT

BFD\_RELOC\_SPARC\_RELATIVE

BFD\_RELOC\_SPARC\_UA16

BFD\_RELOC\_SPARC\_UA32

BFD\_RELOC\_SPARC\_UA64

BFD\_RELOC\_SPARC\_GOTDATA\_HIX22

BFD\_RELOC\_SPARC\_GOTDATA\_LOX10

BFD\_RELOC\_SPARC\_GOTDATA\_OP\_HIX22

```
BFD_RELOC_SPARC_GOTDATA_OP_LOX10
BFD_RELOC_SPARC_GOTDATA_OP
BFD_RELOC_SPARC_JMP_IREL
BFD_RELOC_SPARC_IRELATIVE
     SPARC ELF relocations. There is probably some overlap with other relocation types
     already defined.
BFD_RELOC_SPARC_BASE13
BFD_RELOC_SPARC_BASE22
     I think these are specific to SPARC a.out (e.g., Sun 4).
BFD_RELOC_SPARC_64
BFD_RELOC_SPARC_10
BFD_RELOC_SPARC_11
BFD_RELOC_SPARC_OLO10
BFD_RELOC_SPARC_HH22
BFD_RELOC_SPARC_HM10
BFD_RELOC_SPARC_LM22
BFD_RELOC_SPARC_PC_HH22
BFD_RELOC_SPARC_PC_HM10
BFD_RELOC_SPARC_PC_LM22
BFD_RELOC_SPARC_WDISP16
BFD_RELOC_SPARC_WDISP19
BFD_RELOC_SPARC_7
BFD_RELOC_SPARC_6
BFD_RELOC_SPARC_5
BFD_RELOC_SPARC_DISP64
BFD_RELOC_SPARC_PLT32
BFD_RELOC_SPARC_PLT64
BFD_RELOC_SPARC_HIX22
BFD_RELOC_SPARC_LOX10
BFD_RELOC_SPARC_H44
BFD_RELOC_SPARC_M44
BFD_RELOC_SPARC_L44
BFD_RELOC_SPARC_REGISTER
BFD_RELOC_SPARC_H34
BFD_RELOC_SPARC_SIZE32
BFD_RELOC_SPARC_SIZE64
BFD_RELOC_SPARC_WDISP10
     SPARC64 relocations
BFD_RELOC_SPARC_REV32
     SPARC little endian relocation
BFD_RELOC_SPARC_TLS_GD_HI22
BFD_RELOC_SPARC_TLS_GD_LO10
BFD_RELOC_SPARC_TLS_GD_ADD
```

BFD\_RELOC\_SPARC\_TLS\_GD\_CALL

```
BFD_RELOC_SPARC_TLS_LDM_HI22
BFD_RELOC_SPARC_TLS_LDM_LO10
BFD_RELOC_SPARC_TLS_LDM_ADD
BFD_RELOC_SPARC_TLS_LDM_CALL
BFD_RELOC_SPARC_TLS_LDO_HIX22
BFD_RELOC_SPARC_TLS_LDO_LOX10
BFD_RELOC_SPARC_TLS_LDO_ADD
BFD_RELOC_SPARC_TLS_IE_HI22
BFD_RELOC_SPARC_TLS_IE_LO10
BFD_RELOC_SPARC_TLS_IE_LD
BFD_RELOC_SPARC_TLS_IE_LDX
BFD_RELOC_SPARC_TLS_IE_ADD
BFD_RELOC_SPARC_TLS_LE_HIX22
BFD_RELOC_SPARC_TLS_LE_LOX10
BFD_RELOC_SPARC_TLS_DTPMOD32
BFD_RELOC_SPARC_TLS_DTPMOD64
BFD_RELOC_SPARC_TLS_DTP0FF32
BFD_RELOC_SPARC_TLS_DTP0FF64
BFD_RELOC_SPARC_TLS_TPOFF32
BFD_RELOC_SPARC_TLS_TP0FF64
    SPARC TLS relocations
```

```
BFD_RELOC_SPU_IMM7
BFD_RELOC_SPU_IMM8
BFD_RELOC_SPU_IMM10
BFD_RELOC_SPU_IMM10W
BFD_RELOC_SPU_IMM16
BFD_RELOC_SPU_IMM16W
BFD_RELOC_SPU_IMM18
BFD_RELOC_SPU_PCREL9a
BFD_RELOC_SPU_PCREL9b
BFD_RELOC_SPU_PCREL16
BFD_RELOC_SPU_L016
BFD RELOC SPU HI16
BFD_RELOC_SPU_PPU32
BFD_RELOC_SPU_PPU64
BFD_RELOC_SPU_ADD_PIC
    SPU Relocations.
```

## BFD\_RELOC\_ALPHA\_GPDISP\_HI16

Alpha ECOFF and ELF relocations. Some of these treat the symbol or "addend" in some special way. For GPDISP\_HI16 ("gpdisp") relocations, the symbol is ignored when writing; when reading, it will be the absolute section symbol. The addend is the displacement in bytes of the "lda" instruction from the "ldah" instruction (which is at the address of this reloc).

## BFD\_RELOC\_ALPHA\_GPDISP\_L016

For GPDISP\_LO16 ("ignore") relocations, the symbol is handled as with GPDISP\_HI16 relocs. The addend is ignored when writing the relocations out, and is filled in with the file's GP value on reading, for convenience.

#### BFD RELOC ALPHA GPDISP

The ELF GPDISP relocation is exactly the same as the GPDISP\_HI16 relocation except that there is no accompanying GPDISP\_LO16 relocation.

#### BFD\_RELOC\_ALPHA\_LITERAL

BFD\_RELOC\_ALPHA\_ELF\_LITERAL

## BFD\_RELOC\_ALPHA\_LITUSE

The Alpha LITERAL/LITUSE relocs are produced by a symbol reference; the assembler turns it into a LDQ instruction to load the address of the symbol, and then fills in a register in the real instruction.

The LITERAL reloc, at the LDQ instruction, refers to the .lita section symbol. The addend is ignored when writing, but is filled in with the file's GP value on reading, for convenience, as with the GPDISP\_LO16 reloc.

The ELF\_LITERAL reloc is somewhere between 16\_GOTOFF and GPDISP\_LO16. It should refer to the symbol to be referenced, as with 16\_GOTOFF, but it generates output not based on the position within the .got section, but relative to the GP value chosen for the file during the final link stage.

The LITUSE reloc, on the instruction using the loaded address, gives information to the linker that it might be able to use to optimize away some literal section references. The symbol is ignored (read as the absolute section symbol), and the "addend" indicates the type of instruction using the register: 1 - "memory" fmt insn 2 - bytemanipulation (byte offset reg) 3 - jsr (target of branch)

## BFD\_RELOC\_ALPHA\_HINT

The HINT relocation indicates a value that should be filled into the "hint" field of a jmp/jsr/ret instruction, for possible branch- prediction logic which may be provided on some processors.

#### BFD\_RELOC\_ALPHA\_LINKAGE

The LINKAGE relocation outputs a linkage pair in the object file, which is filled by the linker.

## BFD\_RELOC\_ALPHA\_CODEADDR

The CODEADDR relocation outputs a STO\_CA in the object file, which is filled by the linker.

## BFD\_RELOC\_ALPHA\_GPREL\_HI16

#### BFD\_RELOC\_ALPHA\_GPREL\_L016

The GPREL\_HI/LO relocations together form a 32-bit offset from the GP register.

## BFD\_RELOC\_ALPHA\_BRSGP

Like BFD\_RELOC\_23\_PCREL\_S2, except that the source and target must share a common GP, and the target address is adjusted for STO\_ALPHA\_STD\_GPLOAD.

# BFD\_RELOC\_ALPHA\_NOP

The NOP relocation outputs a NOP if the longword displacement between two procedure entry points is  $< 2^21$ .

#### BFD\_RELOC\_ALPHA\_BSR

The BSR relocation outputs a BSR if the longword displacement between two procedure entry points is  $< 2^21$ .

## BFD\_RELOC\_ALPHA\_LDA

The LDA relocation outputs a LDA if the longword displacement between two procedure entry points is  $< 2^16$ .

# BFD\_RELOC\_ALPHA\_BOH

The BOH relocation outputs a BSR if the longword displacement between two procedure entry points is  $< 2^21$ , or else a hint.

BFD\_RELOC\_ALPHA\_TLSGD

BFD\_RELOC\_ALPHA\_TLSLDM

BFD\_RELOC\_ALPHA\_DTPMOD64

BFD\_RELOC\_ALPHA\_GOTDTPREL16

BFD\_RELOC\_ALPHA\_DTPREL64

BFD\_RELOC\_ALPHA\_DTPREL\_HI16

BFD\_RELOC\_ALPHA\_DTPREL\_L016

BFD\_RELOC\_ALPHA\_DTPREL16

BFD\_RELOC\_ALPHA\_GOTTPREL16

BFD\_RELOC\_ALPHA\_TPREL64

BFD\_RELOC\_ALPHA\_TPREL\_HI16

BFD\_RELOC\_ALPHA\_TPREL\_L016

BFD\_RELOC\_ALPHA\_TPREL16

Alpha thread-local storage relocations.

#### BFD\_RELOC\_MIPS\_JMP

#### BFD\_RELOC\_MICROMIPS\_JMP

The MIPS jump instruction.

#### BFD\_RELOC\_MIPS16\_JMP

The MIPS16 jump instruction.

#### BFD\_RELOC\_MIPS16\_GPREL

MIPS16 GP relative reloc.

#### BFD\_RELOC\_HI16

High 16 bits of 32-bit value; simple reloc.

#### BFD\_RELOC\_HI16\_S

High 16 bits of 32-bit value but the low 16 bits will be sign extended and added to form the final result. If the low 16 bits form a negative number, we need to add one to the high value to compensate for the borrow when the low bits are added.

#### BFD\_RELOC\_L016

Low 16 bits.

## BFD\_RELOC\_HI16\_PCREL

High 16 bits of 32-bit pc-relative value

### BFD\_RELOC\_HI16\_S\_PCREL

High 16 bits of 32-bit pc-relative value, adjusted

## BFD\_RELOC\_LO16\_PCREL

Low 16 bits of pc-relative value

#### BFD\_RELOC\_MIPS16\_GOT16

BFD\_RELOC\_MIPS16\_CALL16

Equivalent of BFD\_RELOC\_MIPS\_\*, but with the MIPS16 layout of 16-bit immediate fields

#### BFD\_RELOC\_MIPS16\_HI16

MIPS16 high 16 bits of 32-bit value.

## BFD\_RELOC\_MIPS16\_HI16\_S

MIPS16 high 16 bits of 32-bit value but the low 16 bits will be sign extended and added to form the final result. If the low 16 bits form a negative number, we need to add one to the high value to compensate for the borrow when the low bits are added.

## BFD\_RELOC\_MIPS16\_L016

MIPS16 low 16 bits.

BFD\_RELOC\_MIPS16\_TLS\_GD

BFD\_RELOC\_MIPS16\_TLS\_LDM

BFD\_RELOC\_MIPS16\_TLS\_DTPREL\_HI16

BFD\_RELOC\_MIPS16\_TLS\_DTPREL\_L016

BFD\_RELOC\_MIPS16\_TLS\_GOTTPREL

BFD\_RELOC\_MIPS16\_TLS\_TPREL\_HI16

BFD\_RELOC\_MIPS16\_TLS\_TPREL\_L016

MIPS16 TLS relocations

#### BFD\_RELOC\_MIPS\_LITERAL

BFD\_RELOC\_MICROMIPS\_LITERAL

Relocation against a MIPS literal section.

BFD\_RELOC\_MICROMIPS\_7\_PCREL\_S1

BFD\_RELOC\_MICROMIPS\_10\_PCREL\_S1

BFD\_RELOC\_MICROMIPS\_16\_PCREL\_S1

BFD\_RELOC\_MICROMIPS\_21\_PCREL\_S1

BFD\_RELOC\_MICROMIPS\_26\_PCREL\_S1

BFD\_RELOC\_MICROMIPS\_18\_PCREL\_S3

BFD\_RELOC\_MICROMIPS\_19\_PCREL\_S2 microMIPS PC-relative relocations.

BFD\_RELOC\_MIPS\_21\_PCREL\_S2

BFD\_RELOC\_MIPS\_26\_PCREL\_S2

BFD\_RELOC\_MIPS\_18\_PCREL\_S3

BFD\_RELOC\_MIPS\_19\_PCREL\_S2

MIPS PC-relative relocations.

BFD\_RELOC\_MIPS\_CALL\_HI16

BFD\_RELOC\_MIPS\_CALL\_L016

BFD\_RELOC\_MICROMIPS\_SUB BFD\_RELOC\_MIPS\_GOT\_PAGE

BFD\_RELOC\_MIPS\_GOT\_OFST

BFD\_RELOC\_MIPS\_GOT\_DISP

BFD\_RELOC\_MIPS\_SHIFT5 BFD\_RELOC\_MIPS\_SHIFT6 BFD\_RELOC\_MIPS\_INSERT\_A BFD\_RELOC\_MIPS\_INSERT\_B BFD\_RELOC\_MIPS\_DELETE BFD\_RELOC\_MIPS\_HIGHEST

BFD\_RELOC\_MIPS\_HIGHER

BFD\_RELOC\_MIPS\_REL16 BFD\_RELOC\_MIPS\_RELGOT BFD\_RELOC\_MIPS\_JALR

BFD\_RELOC\_MICROMIPS\_JALR BFD\_RELOC\_MIPS\_TLS\_DTPMOD32 BFD\_RELOC\_MIPS\_TLS\_DTPREL32 BFD\_RELOC\_MIPS\_TLS\_DTPMOD64 BFD\_RELOC\_MIPS\_TLS\_DTPREL64

BFD\_RELOC\_MIPS\_SUB

BFD\_RELOC\_MICROMIPS\_CALL\_HI16

BFD\_RELOC\_MICROMIPS\_CALL\_L016

BFD\_RELOC\_MICROMIPS\_GOT\_PAGE

BFD\_RELOC\_MICROMIPS\_GOT\_OFST

BFD\_RELOC\_MICROMIPS\_GOT\_DISP

BFD\_RELOC\_MICROMIPS\_HIGHEST

BFD\_RELOC\_MICROMIPS\_HIGHER BFD\_RELOC\_MIPS\_SCN\_DISP BFD\_RELOC\_MICROMIPS\_SCN\_DISP

```
BFD_RELOC_MICROMIPS_GPREL16
BFD_RELOC_MICROMIPS_HI16
BFD_RELOC_MICROMIPS_HI16_S
BFD_RELOC_MICROMIPS_LO16
BFD_RELOC_MICROMIPS_HI16_S_PCREL
BFD_RELOC_MICROMIPS_LO16_PCREL
microMIPS versions of generic BFD relocs.

BFD_RELOC_MIPS_GOT16
BFD_RELOC_MICROMIPS_GOT16
BFD_RELOC_MICROMIPS_CALL16
BFD_RELOC_MICROMIPS_CALL16
BFD_RELOC_MICROMIPS_CALL16
BFD_RELOC_MICROMIPS_GOT_HI16
BFD_RELOC_MICROMIPS_GOT_HI16
BFD_RELOC_MICROMIPS_GOT_LO16
BFD_RELOC_MICROMIPS_GOT_LO16
```

```
BFD_RELOC_MIPS_TLS_GD
BFD_RELOC_MICROMIPS_TLS_GD
BFD_RELOC_MIPS_TLS_LDM
BFD_RELOC_MICROMIPS_TLS_LDM
BFD_RELOC_MIPS_TLS_DTPREL_HI16
BFD_RELOC_MICROMIPS_TLS_DTPREL_HI16
BFD_RELOC_MIPS_TLS_DTPREL_L016
BFD_RELOC_MICROMIPS_TLS_DTPREL_L016
BFD_RELOC_MIPS_TLS_GOTTPREL
BFD_RELOC_MICROMIPS_TLS_GOTTPREL
BFD_RELOC_MIPS_TLS_TPREL32
BFD_RELOC_MIPS_TLS_TPREL64
BFD_RELOC_MIPS_TLS_TPREL_HI16
BFD_RELOC_MICROMIPS_TLS_TPREL_HI16
BFD_RELOC_MIPS_TLS_TPREL_L016
BFD_RELOC_MICROMIPS_TLS_TPREL_L016
BFD_RELOC_MIPS_EH
    MIPS ELF relocations.
BFD_RELOC_MIPS_COPY
BFD_RELOC_MIPS_JUMP_SLOT
    MIPS ELF relocations (VxWorks and PLT extensions).
BFD_RELOC_MOXIE_10_PCREL
    Moxie ELF relocations.
BFD_RELOC_FRV_LABEL16
BFD_RELOC_FRV_LABEL24
BFD_RELOC_FRV_L016
BFD_RELOC_FRV_HI16
BFD_RELOC_FRV_GPREL12
BFD_RELOC_FRV_GPRELU12
BFD_RELOC_FRV_GPREL32
BFD_RELOC_FRV_GPRELHI
BFD_RELOC_FRV_GPRELLO
BFD_RELOC_FRV_GOT12
BFD_RELOC_FRV_GOTHI
BFD_RELOC_FRV_GOTLO
BFD_RELOC_FRV_FUNCDESC
BFD_RELOC_FRV_FUNCDESC_GOT12
BFD_RELOC_FRV_FUNCDESC_GOTHI
BFD_RELOC_FRV_FUNCDESC_GOTLO
BFD_RELOC_FRV_FUNCDESC_VALUE
BFD_RELOC_FRV_FUNCDESC_GOTOFF12
BFD_RELOC_FRV_FUNCDESC_GOTOFFHI
BFD_RELOC_FRV_FUNCDESC_GOTOFFLO
BFD_RELOC_FRV_GOTOFF12
BFD_RELOC_FRV_GOTOFFHI
```

BFD\_RELOC\_FRV\_GOTOFFLO

BFD\_RELOC\_FRV\_GETTLSOFF

BFD\_RELOC\_FRV\_TLSDESC\_VALUE

BFD\_RELOC\_FRV\_GOTTLSDESC12

BFD\_RELOC\_FRV\_GOTTLSDESCHI

BFD\_RELOC\_FRV\_GOTTLSDESCLO

BFD\_RELOC\_FRV\_TLSMOFF12

BFD\_RELOC\_FRV\_TLSMOFFHI

BFD\_RELOC\_FRV\_TLSMOFFLO

BFD\_RELOC\_FRV\_GOTTLSOFF12

BFD\_RELOC\_FRV\_GOTTLSOFFHI

BFD\_RELOC\_FRV\_GOTTLSOFFLO

BFD\_RELOC\_FRV\_TLSOFF

BFD\_RELOC\_FRV\_TLSDESC\_RELAX

BFD\_RELOC\_FRV\_GETTLSOFF\_RELAX

BFD\_RELOC\_FRV\_TLSOFF\_RELAX

BFD\_RELOC\_FRV\_TLSMOFF

Fujitsu Frv Relocations.

## BFD\_RELOC\_MN10300\_GOTOFF24

This is a 24bit GOT-relative reloc for the mn10300.

## BFD\_RELOC\_MN10300\_GOT32

This is a 32bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

## BFD\_RELOC\_MN10300\_GOT24

This is a 24bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

## BFD\_RELOC\_MN10300\_GOT16

This is a 16bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

## BFD\_RELOC\_MN10300\_COPY

Copy symbol at runtime.

## BFD\_RELOC\_MN10300\_GLOB\_DAT

Create GOT entry.

## BFD\_RELOC\_MN10300\_JMP\_SLOT

Create PLT entry.

## BFD\_RELOC\_MN10300\_RELATIVE

Adjust by program base.

## BFD\_RELOC\_MN10300\_SYM\_DIFF

Together with another reloc targeted at the same location, allows for a value that is the difference of two symbols in the same section.

# BFD\_RELOC\_MN10300\_ALIGN

The addend of this reloc is an alignment power that must be honoured at the offset's location, regardless of linker relaxation.

```
BFD_RELOC_MN10300_TLS_GD
BFD_RELOC_MN10300_TLS_LD
BFD_RELOC_MN10300_TLS_LD0
BFD_RELOC_MN10300_TLS_GOTIE
BFD_RELOC_MN10300_TLS_IE
BFD_RELOC_MN10300_TLS_LE
BFD_RELOC_MN10300_TLS_DTPMOD
BFD_RELOC_MN10300_TLS_DTPOFF
BFD_RELOC_MN10300_TLS_TPOFF
Various TLS-related relocations.
```

## BFD\_RELOC\_MN10300\_32\_PCREL

This is a 32bit perel reloc for the mn10300, offset by two bytes in the instruction.

#### BFD\_RELOC\_MN10300\_16\_PCREL

This is a 16bit perel reloc for the mn10300, offset by two bytes in the instruction.

```
BFD_RELOC_386_GOT32
BFD_RELOC_386_PLT32
BFD_RELOC_386_COPY
BFD_RELOC_386_GLOB_DAT
BFD_RELOC_386_JUMP_SLOT
BFD_RELOC_386_RELATIVE
BFD_RELOC_386_GOTOFF
BFD_RELOC_386_GOTPC
BFD_RELOC_386_TLS_TPOFF
BFD_RELOC_386_TLS_IE
BFD_RELOC_386_TLS_GOTIE
BFD_RELOC_386_TLS_LE
BFD_RELOC_386_TLS_GD
BFD_RELOC_386_TLS_LDM
BFD_RELOC_386_TLS_LDO_32
BFD_RELOC_386_TLS_IE_32
BFD_RELOC_386_TLS_LE_32
BFD_RELOC_386_TLS_DTPMOD32
BFD_RELOC_386_TLS_DTP0FF32
BFD_RELOC_386_TLS_TP0FF32
BFD_RELOC_386_TLS_GOTDESC
BFD_RELOC_386_TLS_DESC_CALL
BFD_RELOC_386_TLS_DESC
BFD_RELOC_386_IRELATIVE
    i386/elf relocations
BFD_RELOC_X86_64_GOT32
BFD_RELOC_X86_64_PLT32
```

```
BFD_RELOC_X86_64_COPY
BFD_RELOC_X86_64_GLOB_DAT
BFD_RELOC_X86_64_JUMP_SLOT
BFD_RELOC_X86_64_RELATIVE
BFD_RELOC_X86_64_GOTPCREL
BFD_RELOC_X86_64_32S
BFD_RELOC_X86_64_DTPMOD64
BFD_RELOC_X86_64_DTP0FF64
BFD_RELOC_X86_64_TP0FF64
BFD_RELOC_X86_64_TLSGD
BFD_RELOC_X86_64_TLSLD
BFD_RELOC_X86_64_DTP0FF32
BFD_RELOC_X86_64_GOTTPOFF
BFD_RELOC_X86_64_TP0FF32
BFD_RELOC_X86_64_GOTOFF64
BFD_RELOC_X86_64_GOTPC32
BFD_RELOC_X86_64_GOT64
BFD_RELOC_X86_64_GOTPCREL64
BFD_RELOC_X86_64_GOTPC64
BFD_RELOC_X86_64_GOTPLT64
BFD_RELOC_X86_64_PLT0FF64
BFD_RELOC_X86_64_GOTPC32_TLSDESC
BFD_RELOC_X86_64_TLSDESC_CALL
BFD_RELOC_X86_64_TLSDESC
BFD_RELOC_X86_64_IRELATIVE
BFD_RELOC_X86_64_PC32_BND
BFD_RELOC_X86_64_PLT32_BND
    x86-64/elf relocations
```

BFD\_RELOC\_NS32K\_IMM\_8
BFD\_RELOC\_NS32K\_IMM\_16
BFD\_RELOC\_NS32K\_IMM\_32
BFD\_RELOC\_NS32K\_IMM\_8\_PCREL
BFD\_RELOC\_NS32K\_IMM\_16\_PCREL
BFD\_RELOC\_NS32K\_IMM\_32\_PCREL
BFD\_RELOC\_NS32K\_DISP\_8
BFD\_RELOC\_NS32K\_DISP\_16
BFD\_RELOC\_NS32K\_DISP\_32
BFD\_RELOC\_NS32K\_DISP\_32
BFD\_RELOC\_NS32K\_DISP\_8\_PCREL
BFD\_RELOC\_NS32K\_DISP\_16\_PCREL
BFD\_RELOC\_NS32K\_DISP\_16\_PCREL
BFD\_RELOC\_NS32K\_DISP\_32\_PCREL
ns32k\_relocations

BFD\_RELOC\_PDP11\_DISP\_8\_PCREL BFD\_RELOC\_PDP11\_DISP\_6\_PCREL PDP11 relocations BFD\_RELOC\_PJ\_CODE\_HI16

```
BFD_RELOC_PJ_CODE_LO16
BFD_RELOC_PJ_CODE_DIR16
BFD_RELOC_PJ_CODE_DIR32
BFD_RELOC_PJ_CODE_REL16
BFD_RELOC_PJ_CODE_REL32
    Picojava relocs. Not all of these appear in object files.
BFD_RELOC_PPC_B26
BFD_RELOC_PPC_BA26
BFD_RELOC_PPC_TOC16
BFD_RELOC_PPC_B16
BFD_RELOC_PPC_B16_BRTAKEN
BFD_RELOC_PPC_B16_BRNTAKEN
BFD_RELOC_PPC_BA16
BFD_RELOC_PPC_BA16_BRTAKEN
BFD_RELOC_PPC_BA16_BRNTAKEN
BFD_RELOC_PPC_COPY
BFD_RELOC_PPC_GLOB_DAT
BFD_RELOC_PPC_JMP_SLOT
BFD_RELOC_PPC_RELATIVE
BFD_RELOC_PPC_LOCAL24PC
BFD_RELOC_PPC_EMB_NADDR32
BFD_RELOC_PPC_EMB_NADDR16
BFD_RELOC_PPC_EMB_NADDR16_LO
BFD_RELOC_PPC_EMB_NADDR16_HI
BFD_RELOC_PPC_EMB_NADDR16_HA
BFD_RELOC_PPC_EMB_SDAI16
BFD_RELOC_PPC_EMB_SDA2I16
BFD_RELOC_PPC_EMB_SDA2REL
BFD_RELOC_PPC_EMB_SDA21
BFD_RELOC_PPC_EMB_MRKREF
BFD_RELOC_PPC_EMB_RELSEC16
BFD_RELOC_PPC_EMB_RELST_LO
BFD_RELOC_PPC_EMB_RELST_HI
BFD_RELOC_PPC_EMB_RELST_HA
BFD_RELOC_PPC_EMB_BIT_FLD
BFD_RELOC_PPC_EMB_RELSDA
BFD_RELOC_PPC_VLE_REL8
BFD_RELOC_PPC_VLE_REL15
BFD_RELOC_PPC_VLE_REL24
BFD_RELOC_PPC_VLE_L016A
BFD_RELOC_PPC_VLE_L016D
BFD_RELOC_PPC_VLE_HI16A
BFD_RELOC_PPC_VLE_HI16D
```

BFD\_RELOC\_PPC\_VLE\_HA16A BFD\_RELOC\_PPC\_VLE\_HA16D

```
BFD_RELOC_PPC_VLE_SDA21
BFD_RELOC_PPC_VLE_SDA21_LO
BFD_RELOC_PPC_VLE_SDAREL_LO16A
BFD_RELOC_PPC_VLE_SDAREL_LO16D
BFD_RELOC_PPC_VLE_SDAREL_HI16A
BFD_RELOC_PPC_VLE_SDAREL_HI16D
BFD_RELOC_PPC_VLE_SDAREL_HA16A
BFD_RELOC_PPC_VLE_SDAREL_HA16D
BFD_RELOC_PPC64_HIGHER
BFD_RELOC_PPC64_HIGHER_S
BFD_RELOC_PPC64_HIGHEST
BFD_RELOC_PPC64_HIGHEST_S
BFD_RELOC_PPC64_TOC16_LO
BFD_RELOC_PPC64_TOC16_HI
BFD_RELOC_PPC64_TOC16_HA
BFD_RELOC_PPC64_TOC
BFD_RELOC_PPC64_PLTGOT16
BFD_RELOC_PPC64_PLTG0T16_L0
BFD_RELOC_PPC64_PLTGOT16_HI
BFD_RELOC_PPC64_PLTGOT16_HA
BFD_RELOC_PPC64_ADDR16_DS
BFD_RELOC_PPC64_ADDR16_LO_DS
BFD_RELOC_PPC64_GOT16_DS
BFD_RELOC_PPC64_GOT16_LO_DS
BFD_RELOC_PPC64_PLT16_LO_DS
BFD_RELOC_PPC64_SECTOFF_DS
BFD_RELOC_PPC64_SECTOFF_LO_DS
BFD_RELOC_PPC64_TOC16_DS
BFD_RELOC_PPC64_TOC16_LO_DS
BFD_RELOC_PPC64_PLTGOT16_DS
BFD_RELOC_PPC64_PLTGOT16_LO_DS
BFD_RELOC_PPC64_ADDR16_HIGH
BFD_RELOC_PPC64_ADDR16_HIGHA
BFD_RELOC_PPC64_ADDR64_LOCAL
    Power(rs6000) and PowerPC relocations.
```

BFD\_RELOC\_PPC\_TLS
BFD\_RELOC\_PPC\_TLSGD
BFD\_RELOC\_PPC\_TLSLD
BFD\_RELOC\_PPC\_DTPMOD
BFD\_RELOC\_PPC\_TPREL16
BFD\_RELOC\_PPC\_TPREL16\_LO
BFD\_RELOC\_PPC\_TPREL16\_HI
BFD\_RELOC\_PPC\_TPREL16\_HA
BFD\_RELOC\_PPC\_TPREL
BFD\_RELOC\_PPC\_DTPREL16
BFD\_RELOC\_PPC\_DTPREL16\_LO

```
BFD_RELOC_PPC_DTPREL16_HI
BFD_RELOC_PPC_DTPREL16_HA
BFD_RELOC_PPC_DTPREL
BFD_RELOC_PPC_GOT_TLSGD16
BFD_RELOC_PPC_GOT_TLSGD16_LO
BFD_RELOC_PPC_GOT_TLSGD16_HI
BFD_RELOC_PPC_GOT_TLSGD16_HA
BFD_RELOC_PPC_GOT_TLSLD16
BFD_RELOC_PPC_GOT_TLSLD16_LO
BFD_RELOC_PPC_GOT_TLSLD16_HI
BFD_RELOC_PPC_GOT_TLSLD16_HA
BFD_RELOC_PPC_GOT_TPREL16
BFD_RELOC_PPC_GOT_TPREL16_LO
BFD_RELOC_PPC_GOT_TPREL16_HI
BFD_RELOC_PPC_GOT_TPREL16_HA
BFD_RELOC_PPC_GOT_DTPREL16
BFD_RELOC_PPC_GOT_DTPREL16_LO
BFD_RELOC_PPC_GOT_DTPREL16_HI
BFD_RELOC_PPC_GOT_DTPREL16_HA
BFD_RELOC_PPC64_TPREL16_DS
BFD_RELOC_PPC64_TPREL16_LO_DS
BFD_RELOC_PPC64_TPREL16_HIGHER
BFD_RELOC_PPC64_TPREL16_HIGHERA
BFD_RELOC_PPC64_TPREL16_HIGHEST
BFD_RELOC_PPC64_TPREL16_HIGHESTA
BFD_RELOC_PPC64_DTPREL16_DS
BFD_RELOC_PPC64_DTPREL16_LO_DS
BFD_RELOC_PPC64_DTPREL16_HIGHER
BFD_RELOC_PPC64_DTPREL16_HIGHERA
BFD_RELOC_PPC64_DTPREL16_HIGHEST
BFD_RELOC_PPC64_DTPREL16_HIGHESTA
BFD_RELOC_PPC64_TPREL16_HIGH
BFD_RELOC_PPC64_TPREL16_HIGHA
BFD_RELOC_PPC64_DTPREL16_HIGH
BFD_RELOC_PPC64_DTPREL16_HIGHA
    PowerPC and PowerPC64 thread-local storage relocations.
```

#### BFD\_RELOC\_I370\_D12

IBM 370/390 relocations

## BFD\_RELOC\_CTOR

The type of reloc used to build a constructor table - at the moment probably a 32 bit wide absolute relocation, but the target can choose. It generally does map to one of the other relocation types.

# BFD\_RELOC\_ARM\_PCREL\_BRANCH

ARM 26 bit pc-relative branch. The lowest two bits must be zero and are not stored in the instruction.

## BFD\_RELOC\_ARM\_PCREL\_BLX

ARM 26 bit pc-relative branch. The lowest bit must be zero and is not stored in the instruction. The 2nd lowest bit comes from a 1 bit field in the instruction.

#### BFD\_RELOC\_THUMB\_PCREL\_BLX

Thumb 22 bit pc-relative branch. The lowest bit must be zero and is not stored in the instruction. The 2nd lowest bit comes from a 1 bit field in the instruction.

## BFD\_RELOC\_ARM\_PCREL\_CALL

ARM 26-bit pc-relative branch for an unconditional BL or BLX instruction.

#### BFD\_RELOC\_ARM\_PCREL\_JUMP

ARM 26-bit pc-relative branch for B or conditional BL instruction.

#### BFD\_RELOC\_THUMB\_PCREL\_BRANCH7

BFD\_RELOC\_THUMB\_PCREL\_BRANCH9

BFD\_RELOC\_THUMB\_PCREL\_BRANCH12

BFD\_RELOC\_THUMB\_PCREL\_BRANCH20

BFD\_RELOC\_THUMB\_PCREL\_BRANCH23

BFD\_RELOC\_THUMB\_PCREL\_BRANCH25

Thumb 7-, 9-, 12-, 20-, 23-, and 25-bit pc-relative branches. The lowest bit must be zero and is not stored in the instruction. Note that the corresponding ELF R\_ARM\_THM\_JUMPnn constant has an "nn" one smaller in all cases. Note further that BRANCH23 corresponds to R\_ARM\_THM\_CALL.

#### BFD\_RELOC\_ARM\_OFFSET\_IMM

12-bit immediate offset, used in ARM-format ldr and str instructions.

#### BFD\_RELOC\_ARM\_THUMB\_OFFSET

5-bit immediate offset, used in Thumb-format ldr and str instructions.

#### BFD\_RELOC\_ARM\_TARGET1

Pc-relative or absolute relocation depending on target. Used for entries in .init\_array sections.

## BFD\_RELOC\_ARM\_ROSEGREL32

Read-only segment base relative address.

#### BFD\_RELOC\_ARM\_SBREL32

Data segment base relative address.

#### BFD\_RELOC\_ARM\_TARGET2

This reloc is used for references to RTTI data from exception handling tables. The actual definition depends on the target. It may be a pc-relative or some form of GOT-indirect relocation.

#### BFD\_RELOC\_ARM\_PREL31

31-bit PC relative address.

BFD\_RELOC\_ARM\_LDRS\_PC\_G1 BFD\_RELOC\_ARM\_LDRS\_PC\_G2 BFD\_RELOC\_ARM\_LDC\_PC\_G0

```
BFD_RELOC_ARM_MOVW
BFD_RELOC_ARM_MOVT
BFD_RELOC_ARM_MOVW_PCREL
BFD_RELOC_ARM_MOVT_PCREL
BFD_RELOC_ARM_THUMB_MOVW
BFD_RELOC_ARM_THUMB_MOVT
BFD_RELOC_ARM_THUMB_MOVW_PCREL
BFD_RELOC_ARM_THUMB_MOVT_PCREL
    Low and High halfword relocations for MOVW and MOVT instructions.
BFD_RELOC_ARM_JUMP_SLOT
BFD_RELOC_ARM_GLOB_DAT
BFD_RELOC_ARM_GOT32
BFD_RELOC_ARM_PLT32
BFD_RELOC_ARM_RELATIVE
BFD_RELOC_ARM_GOTOFF
BFD_RELOC_ARM_GOTPC
BFD_RELOC_ARM_GOT_PREL
    Relocations for setting up GOTs and PLTs for shared libraries.
BFD_RELOC_ARM_TLS_GD32
BFD_RELOC_ARM_TLS_LD032
BFD_RELOC_ARM_TLS_LDM32
BFD_RELOC_ARM_TLS_DTPOFF32
BFD_RELOC_ARM_TLS_DTPMOD32
BFD_RELOC_ARM_TLS_TPOFF32
BFD_RELOC_ARM_TLS_IE32
BFD_RELOC_ARM_TLS_LE32
BFD_RELOC_ARM_TLS_GOTDESC
BFD_RELOC_ARM_TLS_CALL
BFD_RELOC_ARM_THM_TLS_CALL
BFD_RELOC_ARM_TLS_DESCSEQ
BFD_RELOC_ARM_THM_TLS_DESCSEQ
BFD_RELOC_ARM_TLS_DESC
    ARM thread-local storage relocations.
BFD_RELOC_ARM_ALU_PC_GO_NC
BFD_RELOC_ARM_ALU_PC_GO
BFD_RELOC_ARM_ALU_PC_G1_NC
BFD_RELOC_ARM_ALU_PC_G1
BFD_RELOC_ARM_ALU_PC_G2
BFD_RELOC_ARM_LDR_PC_GO
BFD_RELOC_ARM_LDR_PC_G1
BFD_RELOC_ARM_LDR_PC_G2
BFD_RELOC_ARM_LDRS_PC_GO
```

```
BFD_RELOC_ARM_LDC_PC_G1
BFD_RELOC_ARM_LDC_PC_G2
BFD_RELOC_ARM_ALU_SB_GO_NC
BFD_RELOC_ARM_ALU_SB_GO
BFD_RELOC_ARM_ALU_SB_G1_NC
BFD_RELOC_ARM_ALU_SB_G1
BFD_RELOC_ARM_ALU_SB_G2
BFD_RELOC_ARM_LDR_SB_GO
BFD_RELOC_ARM_LDR_SB_G1
BFD_RELOC_ARM_LDR_SB_G2
BFD_RELOC_ARM_LDRS_SB_GO
BFD_RELOC_ARM_LDRS_SB_G1
BFD_RELOC_ARM_LDRS_SB_G2
BFD_RELOC_ARM_LDC_SB_GO
BFD_RELOC_ARM_LDC_SB_G1
BFD_RELOC_ARM_LDC_SB_G2
    ARM group relocations.
```

## BFD\_RELOC\_ARM\_V4BX

Annotation of BX instructions.

# BFD\_RELOC\_ARM\_IRELATIVE ARM support for STT\_GNU\_IFUNC.

BFD\_RELOC\_ARM\_IMMEDIATE BFD\_RELOC\_ARM\_ADRL\_IMMEDIATE BFD\_RELOC\_ARM\_T32\_IMMEDIATE BFD\_RELOC\_ARM\_T32\_ADD\_IMM BFD\_RELOC\_ARM\_T32\_IMM12 BFD\_RELOC\_ARM\_T32\_ADD\_PC12 BFD\_RELOC\_ARM\_SHIFT\_IMM BFD\_RELOC\_ARM\_SMC BFD\_RELOC\_ARM\_HVC BFD\_RELOC\_ARM\_SWI BFD\_RELOC\_ARM\_MULTI BFD\_RELOC\_ARM\_CP\_OFF\_IMM BFD\_RELOC\_ARM\_CP\_OFF\_IMM\_S2 BFD\_RELOC\_ARM\_T32\_CP\_OFF\_IMM BFD\_RELOC\_ARM\_T32\_CP\_OFF\_IMM\_S2 BFD\_RELOC\_ARM\_ADR\_IMM BFD\_RELOC\_ARM\_LDR\_IMM BFD\_RELOC\_ARM\_LITERAL BFD\_RELOC\_ARM\_IN\_POOL BFD\_RELOC\_ARM\_OFFSET\_IMM8 BFD\_RELOC\_ARM\_T32\_OFFSET\_U8 BFD\_RELOC\_ARM\_T32\_OFFSET\_IMM BFD\_RELOC\_ARM\_HWLITERAL

BFD\_RELOC\_ARM\_THUMB\_ADD

```
BFD_RELOC_ARM_THUMB_IMM
BFD_RELOC_ARM_THUMB_SHIFT
These relocs are only used within the ARM assembler. They are not (at present) written to any object files.
```

```
BFD_RELOC_SH_PCDISP8BY2
BFD_RELOC_SH_PCDISP12BY2
BFD_RELOC_SH_IMM3
BFD_RELOC_SH_IMM3U
BFD_RELOC_SH_DISP12
BFD_RELOC_SH_DISP12BY2
BFD_RELOC_SH_DISP12BY4
BFD_RELOC_SH_DISP12BY8
BFD_RELOC_SH_DISP20
BFD_RELOC_SH_DISP20BY8
BFD_RELOC_SH_IMM4
BFD_RELOC_SH_IMM4BY2
BFD_RELOC_SH_IMM4BY4
BFD_RELOC_SH_IMM8
BFD_RELOC_SH_IMM8BY2
BFD_RELOC_SH_IMM8BY4
BFD_RELOC_SH_PCRELIMM8BY2
BFD_RELOC_SH_PCRELIMM8BY4
BFD_RELOC_SH_SWITCH16
BFD_RELOC_SH_SWITCH32
BFD_RELOC_SH_USES
BFD_RELOC_SH_COUNT
BFD_RELOC_SH_ALIGN
BFD_RELOC_SH_CODE
BFD_RELOC_SH_DATA
BFD_RELOC_SH_LABEL
BFD_RELOC_SH_LOOP_START
BFD_RELOC_SH_LOOP_END
BFD_RELOC_SH_COPY
BFD_RELOC_SH_GLOB_DAT
BFD_RELOC_SH_JMP_SLOT
BFD_RELOC_SH_RELATIVE
BFD_RELOC_SH_GOTPC
BFD_RELOC_SH_GOT_LOW16
BFD_RELOC_SH_GOT_MEDLOW16
BFD_RELOC_SH_GOT_MEDHI16
BFD_RELOC_SH_GOT_HI16
BFD_RELOC_SH_GOTPLT_LOW16
BFD_RELOC_SH_GOTPLT_MEDLOW16
BFD_RELOC_SH_GOTPLT_MEDHI16
BFD_RELOC_SH_GOTPLT_HI16
```

BFD\_RELOC\_SH\_PLT\_LOW16

- BFD\_RELOC\_SH\_PLT\_MEDLOW16
- BFD\_RELOC\_SH\_PLT\_MEDHI16
- BFD\_RELOC\_SH\_PLT\_HI16
- BFD\_RELOC\_SH\_GOTOFF\_LOW16
- BFD\_RELOC\_SH\_GOTOFF\_MEDLOW16
- BFD\_RELOC\_SH\_GOTOFF\_MEDHI16
- BFD\_RELOC\_SH\_GOTOFF\_HI16
- BFD\_RELOC\_SH\_GOTPC\_LOW16
- BFD\_RELOC\_SH\_GOTPC\_MEDLOW16
- BFD\_RELOC\_SH\_GOTPC\_MEDHI16
- BFD\_RELOC\_SH\_GOTPC\_HI16
- BFD\_RELOC\_SH\_COPY64
- BFD\_RELOC\_SH\_GLOB\_DAT64
- BFD\_RELOC\_SH\_JMP\_SLOT64
- BFD\_RELOC\_SH\_RELATIVE64
- BFD\_RELOC\_SH\_GOT10BY4
- BFD\_RELOC\_SH\_GOT10BY8
- BFD\_RELOC\_SH\_GOTPLT10BY4
- BFD\_RELOC\_SH\_GOTPLT10BY8
- BFD\_RELOC\_SH\_GOTPLT32
- BFD\_RELOC\_SH\_SHMEDIA\_CODE
- BFD\_RELOC\_SH\_IMMU5
- BFD\_RELOC\_SH\_IMMS6
- BFD\_RELOC\_SH\_IMMS6BY32
- BFD\_RELOC\_SH\_IMMU6
- BFD\_RELOC\_SH\_IMMS10
- BFD\_RELOC\_SH\_IMMS10BY2
- BFD\_RELOC\_SH\_IMMS10BY4
- BFD\_RELOC\_SH\_IMMS10BY8
- BFD\_RELOC\_SH\_IMMS16
- BFD\_RELOC\_SH\_IMMU16
- BFD\_RELOC\_SH\_IMM\_LOW16
- BFD\_RELOC\_SH\_IMM\_LOW16\_PCREL
- BFD\_RELOC\_SH\_IMM\_MEDLOW16
- BFD\_RELOC\_SH\_IMM\_MEDLOW16\_PCREL
- BFD\_RELOC\_SH\_IMM\_MEDHI16
- BFD\_RELOC\_SH\_IMM\_MEDHI16\_PCREL
- BFD\_RELOC\_SH\_IMM\_HI16
- BFD\_RELOC\_SH\_IMM\_HI16\_PCREL
- BFD\_RELOC\_SH\_PT\_16
- BFD\_RELOC\_SH\_TLS\_GD\_32
- BFD\_RELOC\_SH\_TLS\_LD\_32
- BFD\_RELOC\_SH\_TLS\_LDO\_32
- BFD\_RELOC\_SH\_TLS\_IE\_32
- BFD\_RELOC\_SH\_TLS\_LE\_32
- BFD\_RELOC\_SH\_TLS\_DTPMOD32
- BFD\_RELOC\_SH\_TLS\_DTPOFF32

BFD\_RELOC\_SH\_TLS\_TPOFF32

BFD\_RELOC\_SH\_GOT20

BFD\_RELOC\_SH\_GOTOFF20

BFD\_RELOC\_SH\_GOTFUNCDESC

BFD\_RELOC\_SH\_GOTFUNCDESC20

BFD\_RELOC\_SH\_GOTOFFFUNCDESC

BFD\_RELOC\_SH\_GOTOFFFUNCDESC20

BFD\_RELOC\_SH\_FUNCDESC

Renesas / SuperH SH relocs. Not all of these appear in object files.

## BFD\_RELOC\_ARC\_B22\_PCREL

ARC Cores relocs. ARC 22 bit pc-relative branch. The lowest two bits must be zero and are not stored in the instruction. The high 20 bits are installed in bits 26 through 7 of the instruction.

#### BFD\_RELOC\_ARC\_B26

ARC 26 bit absolute branch. The lowest two bits must be zero and are not stored in the instruction. The high 24 bits are installed in bits 23 through 0.

#### BFD\_RELOC\_BFIN\_16\_IMM

ADI Blackfin 16 bit immediate absolute reloc.

#### BFD\_RELOC\_BFIN\_16\_HIGH

ADI Blackfin 16 bit immediate absolute reloc higher 16 bits.

# BFD\_RELOC\_BFIN\_4\_PCREL

ADI Blackfin 'a' part of LSETUP.

## BFD\_RELOC\_BFIN\_5\_PCREL

ADI Blackfin.

## BFD\_RELOC\_BFIN\_16\_LOW

ADI Blackfin 16 bit immediate absolute reloc lower 16 bits.

#### BFD\_RELOC\_BFIN\_10\_PCREL

ADI Blackfin.

#### BFD\_RELOC\_BFIN\_11\_PCREL

ADI Blackfin 'b' part of LSETUP.

## BFD\_RELOC\_BFIN\_12\_PCREL\_JUMP

ADI Blackfin.

#### BFD\_RELOC\_BFIN\_12\_PCREL\_JUMP\_S

ADI Blackfin Short jump, pcrel.

## BFD\_RELOC\_BFIN\_24\_PCREL\_CALL\_X

ADI Blackfin Call.x not implemented.

#### BFD\_RELOC\_BFIN\_24\_PCREL\_JUMP\_L

ADI Blackfin Long Jump pcrel.

- BFD\_RELOC\_BFIN\_GOT17M4
- BFD\_RELOC\_BFIN\_GOTHI
- BFD\_RELOC\_BFIN\_GOTLO
- BFD\_RELOC\_BFIN\_FUNCDESC
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOT17M4
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOTHI
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOTLO
- BFD\_RELOC\_BFIN\_FUNCDESC\_VALUE
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOTOFF17M4
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOTOFFHI
- BFD\_RELOC\_BFIN\_FUNCDESC\_GOTOFFLO
- BFD\_RELOC\_BFIN\_GOTOFF17M4
- BFD\_RELOC\_BFIN\_GOTOFFHI
- BFD\_RELOC\_BFIN\_GOTOFFLO
  - ADI Blackfin FD-PIC relocations.
- BFD\_RELOC\_BFIN\_GOT
  - ADI Blackfin GOT relocation.
- BFD\_RELOC\_BFIN\_PLTPC
  - ADI Blackfin PLTPC relocation.
- BFD\_ARELOC\_BFIN\_PUSH
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_CONST
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_ADD
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_SUB
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_MULT
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_DIV
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_MOD
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_LSHIFT
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_RSHIFT
  - ADI Blackfin arithmetic relocation.
- BFD\_ARELOC\_BFIN\_AND
  - ADI Blackfin arithmetic relocation.

# BFD\_ARELOC\_BFIN\_OR

ADI Blackfin arithmetic relocation.

## BFD\_ARELOC\_BFIN\_XOR

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_LAND

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_LOR

ADI Blackfin arithmetic relocation.

## BFD\_ARELOC\_BFIN\_LEN

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_NEG

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_COMP

ADI Blackfin arithmetic relocation.

## BFD\_ARELOC\_BFIN\_PAGE

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_HWPAGE

ADI Blackfin arithmetic relocation.

#### BFD\_ARELOC\_BFIN\_ADDR

ADI Blackfin arithmetic relocation.

## BFD\_RELOC\_D10V\_10\_PCREL\_R

Mitsubishi D10V relocs. This is a 10-bit reloc with the right 2 bits assumed to be 0.

## BFD\_RELOC\_D10V\_10\_PCREL\_L

Mitsubishi D10V relocs. This is a 10-bit reloc with the right 2 bits assumed to be 0. This is the same as the previous reloc except it is in the left container, i.e., shifted left 15 bits.

# BFD\_RELOC\_D10V\_18

This is an 18-bit reloc with the right 2 bits assumed to be 0.

#### BFD\_RELOC\_D10V\_18\_PCREL

This is an 18-bit reloc with the right 2 bits assumed to be 0.

## BFD\_RELOC\_D30V\_6

Mitsubishi D30V relocs. This is a 6-bit absolute reloc.

# BFD\_RELOC\_D30V\_9\_PCREL

This is a 6-bit pc-relative reloc with the right 3 bits assumed to be 0.

#### BFD\_RELOC\_D30V\_9\_PCREL\_R

This is a 6-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.

## BFD\_RELOC\_D30V\_15

This is a 12-bit absolute reloc with the right 3 bits assumed to be 0.

#### BFD\_RELOC\_D30V\_15\_PCREL

This is a 12-bit pc-relative reloc with the right 3 bits assumed to be 0.

#### BFD\_RELOC\_D30V\_15\_PCREL\_R

This is a 12-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.

## BFD\_RELOC\_D30V\_21

This is an 18-bit absolute reloc with the right 3 bits assumed to be 0.

## BFD\_RELOC\_D30V\_21\_PCREL

This is an 18-bit pc-relative reloc with the right 3 bits assumed to be 0.

#### BFD\_RELOC\_D30V\_21\_PCREL\_R

This is an 18-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.

## BFD\_RELOC\_D30V\_32

This is a 32-bit absolute reloc.

## BFD\_RELOC\_D30V\_32\_PCREL

This is a 32-bit pc-relative reloc.

## BFD\_RELOC\_DLX\_HI16\_S

DLX relocs

#### BFD\_RELOC\_DLX\_L016

DLX relocs

## BFD\_RELOC\_DLX\_JMP26

DLX relocs

## BFD\_RELOC\_M32C\_H18

BFD\_RELOC\_M32C\_RL\_JUMP

BFD\_RELOC\_M32C\_RL\_1ADDR

BFD\_RELOC\_M32C\_RL\_2ADDR

Renesas M16C/M32C Relocations.

#### BFD\_RELOC\_M32R\_24

Renesas M32R (formerly Mitsubishi M32R) relocs. This is a 24 bit absolute address.

#### BFD\_RELOC\_M32R\_10\_PCREL

This is a 10-bit pc-relative reloc with the right 2 bits assumed to be 0.

## BFD\_RELOC\_M32R\_18\_PCREL

This is an 18-bit reloc with the right 2 bits assumed to be 0.

#### BFD\_RELOC\_M32R\_26\_PCREL

This is a 26-bit reloc with the right 2 bits assumed to be 0.

## BFD\_RELOC\_M32R\_HI16\_ULO

This is a 16-bit reloc containing the high 16 bits of an address used when the lower 16 bits are treated as unsigned.

#### BFD\_RELOC\_M32R\_HI16\_SLO

This is a 16-bit reloc containing the high 16 bits of an address used when the lower 16 bits are treated as signed.

#### BFD\_RELOC\_M32R\_L016

This is a 16-bit reloc containing the lower 16 bits of an address.

#### BFD\_RELOC\_M32R\_SDA16

This is a 16-bit reloc containing the small data area offset for use in add3, load, and store instructions.

# BFD\_RELOC\_M32R\_GOT24

BFD\_RELOC\_M32R\_26\_PLTREL

BFD\_RELOC\_M32R\_COPY

BFD\_RELOC\_M32R\_GLOB\_DAT

BFD\_RELOC\_M32R\_JMP\_SLOT

BFD\_RELOC\_M32R\_RELATIVE

BFD\_RELOC\_M32R\_GOTOFF

BFD\_RELOC\_M32R\_GOTOFF\_HI\_ULO

BFD\_RELOC\_M32R\_GOTOFF\_HI\_SLO

BFD\_RELOC\_M32R\_GOTOFF\_LO

BFD\_RELOC\_M32R\_GOTPC24

BFD\_RELOC\_M32R\_GOT16\_HI\_ULO

BFD\_RELOC\_M32R\_GOT16\_HI\_SLO

BFD\_RELOC\_M32R\_GOT16\_LO

BFD\_RELOC\_M32R\_GOTPC\_HI\_ULO

BFD\_RELOC\_M32R\_GOTPC\_HI\_SLO

BFD\_RELOC\_M32R\_GOTPC\_LO

For PIC.

## BFD\_RELOC\_NDS32\_20

NDS32 relocs. This is a 20 bit absolute address.

## BFD\_RELOC\_NDS32\_9\_PCREL

This is a 9-bit pc-relative reloc with the right 1 bit assumed to be 0.

# BFD\_RELOC\_NDS32\_WORD\_9\_PCREL

This is a 9-bit pc-relative reloc with the right 1 bit assumed to be 0.

## BFD\_RELOC\_NDS32\_15\_PCREL

This is an 15-bit reloc with the right 1 bit assumed to be 0.

#### BFD\_RELOC\_NDS32\_17\_PCREL

This is an 17-bit reloc with the right 1 bit assumed to be 0.

#### BFD\_RELOC\_NDS32\_25\_PCREL

This is a 25-bit reloc with the right 1 bit assumed to be 0.

## BFD\_RELOC\_NDS32\_HI20

This is a 20-bit reloc containing the high 20 bits of an address used with the lower 12 bits

#### BFD\_RELOC\_NDS32\_L012S3

This is a 12-bit reloc containing the lower 12 bits of an address then shift right by 3. This is used with ldi,sdi...

## BFD\_RELOC\_NDS32\_L012S2

This is a 12-bit reloc containing the lower 12 bits of an address then shift left by 2. This is used with lwi,swi...

#### BFD\_RELOC\_NDS32\_L012S1

This is a 12-bit reloc containing the lower 12 bits of an address then shift left by 1. This is used with lhi,shi...

## BFD\_RELOC\_NDS32\_L012S0

This is a 12-bit reloc containing the lower 12 bits of an address then shift left by 0. This is used with lbisbi...

#### BFD\_RELOC\_NDS32\_L012S0\_ORI

This is a 12-bit reloc containing the lower 12 bits of an address then shift left by 0. This is only used with branch relaxations

#### BFD\_RELOC\_NDS32\_SDA15S3

This is a 15-bit reloc containing the small data area 18-bit signed offset and shift left by 3 for use in ldi, sdi...

#### BFD\_RELOC\_NDS32\_SDA15S2

This is a 15-bit reloc containing the small data area 17-bit signed offset and shift left by 2 for use in lwi, swi...

## BFD\_RELOC\_NDS32\_SDA15S1

This is a 15-bit reloc containing the small data area 16-bit signed offset and shift left by 1 for use in lhi, shi...

#### BFD\_RELOC\_NDS32\_SDA15S0

This is a 15-bit reloc containing the small data area 15-bit signed offset and shift left by 0 for use in lbi, sbi...

#### BFD\_RELOC\_NDS32\_SDA16S3

This is a 16-bit reloc containing the small data area 16-bit signed offset and shift left by 3

#### BFD\_RELOC\_NDS32\_SDA17S2

This is a 17-bit reloc containing the small data area 17-bit signed offset and shift left by 2 for use in lwi.gp, swi.gp...

#### BFD\_RELOC\_NDS32\_SDA18S1

This is a 18-bit reloc containing the small data area 18-bit signed offset and shift left by 1 for use in lhi.gp, shi.gp...

# BFD\_RELOC\_NDS32\_SDA19S0

This is a 19-bit reloc containing the small data area 19-bit signed offset and shift left by 0 for use in lbi.gp, sbi.gp...

```
BFD_RELOC_NDS32_GOT20
BFD_RELOC_NDS32_9_PLTREL
BFD_RELOC_NDS32_25_PLTREL
BFD_RELOC_NDS32_COPY
BFD_RELOC_NDS32_GLOB_DAT
BFD_RELOC_NDS32_JMP_SLOT
BFD_RELOC_NDS32_RELATIVE
BFD_RELOC_NDS32_GOTOFF
BFD_RELOC_NDS32_GOTOFF_HI20
BFD_RELOC_NDS32_GOTOFF_L012
BFD_RELOC_NDS32_GOTPC20
BFD_RELOC_NDS32_GOT_HI20
BFD_RELOC_NDS32_GOT_L012
BFD_RELOC_NDS32_GOTPC_HI20
BFD_RELOC_NDS32_GOTPC_LO12
    for PIC
BFD_RELOC_NDS32_INSN16
BFD_RELOC_NDS32_LABEL
BFD_RELOC_NDS32_LONGCALL1
BFD_RELOC_NDS32_LONGCALL2
BFD_RELOC_NDS32_LONGCALL3
BFD_RELOC_NDS32_LONGJUMP1
BFD_RELOC_NDS32_LONGJUMP2
BFD_RELOC_NDS32_LONGJUMP3
BFD_RELOC_NDS32_LOADSTORE
BFD_RELOC_NDS32_9_FIXED
BFD_RELOC_NDS32_15_FIXED
BFD_RELOC_NDS32_17_FIXED
BFD_RELOC_NDS32_25_FIXED
BFD_RELOC_NDS32_LONGCALL4
BFD_RELOC_NDS32_LONGCALL5
BFD_RELOC_NDS32_LONGCALL6
BFD_RELOC_NDS32_LONGJUMP4
BFD_RELOC_NDS32_LONGJUMP5
BFD_RELOC_NDS32_LONGJUMP6
BFD_RELOC_NDS32_LONGJUMP7
    for relax
BFD_RELOC_NDS32_PLTREL_HI20
BFD_RELOC_NDS32_PLTREL_L012
BFD_RELOC_NDS32_PLT_GOTREL_HI20
BFD_RELOC_NDS32_PLT_GOTREL_L012
    for PIC
```

BFD\_RELOC\_NDS32\_DIFF16

```
BFD_RELOC_NDS32_SDA12S2_DP
BFD_RELOC_NDS32_SDA12S2_SP
BFD_RELOC_NDS32_L012S2_DP
BFD_RELOC_NDS32_L012S2_SP
     for floating point
BFD_RELOC_NDS32_DWARF2_OP1
BFD_RELOC_NDS32_DWARF2_OP2
BFD_RELOC_NDS32_DWARF2_LEB
     for dwarf2 debug_line.
BFD_RELOC_NDS32_UPDATE_TA
     for eliminate 16-bit instructions
BFD_RELOC_NDS32_PLT_GOTREL_LO20
BFD_RELOC_NDS32_PLT_GOTREL_L015
BFD_RELOC_NDS32_PLT_GOTREL_L019
BFD_RELOC_NDS32_GOT_L015
BFD_RELOC_NDS32_GOT_L019
BFD_RELOC_NDS32_GOTOFF_L015
BFD_RELOC_NDS32_GOTOFF_L019
BFD_RELOC_NDS32_GOT15S2
BFD_RELOC_NDS32_GOT17S2
     for PIC object relaxation
BFD RELOC NDS32 5
     NDS32 relocs. This is a 5 bit absolute address.
BFD_RELOC_NDS32_10_UPCREL
     This is a 10-bit unsigned pc-relative reloc with the right 1 bit assumed to be 0.
BFD_RELOC_NDS32_SDA_FP7U2_RELA
     If fp were omitted, fp can used as another gp.
BFD_RELOC_NDS32_RELAX_ENTRY
BFD_RELOC_NDS32_GOT_SUFF
BFD_RELOC_NDS32_GOTOFF_SUFF
BFD_RELOC_NDS32_PLT_GOT_SUFF
BFD_RELOC_NDS32_MULCALL_SUFF
BFD_RELOC_NDS32_PTR
BFD_RELOC_NDS32_PTR_COUNT
BFD_RELOC_NDS32_PTR_RESOLVED
BFD_RELOC_NDS32_PLTBLOCK
BFD_RELOC_NDS32_RELAX_REGION_BEGIN
BFD_RELOC_NDS32_RELAX_REGION_END
BFD_RELOC_NDS32_MINUEND
BFD_RELOC_NDS32_SUBTRAHEND
BFD_RELOC_NDS32_DIFF8
```

BFD\_RELOC\_NDS32\_DIFF32

BFD\_RELOC\_NDS32\_DIFF\_ULEB128

BFD\_RELOC\_NDS32\_EMPTY

relaxation relative relocation types

#### BFD\_RELOC\_NDS32\_25\_ABS

This is a 25 bit absolute address.

BFD\_RELOC\_NDS32\_DATA

BFD\_RELOC\_NDS32\_TRAN

BFD\_RELOC\_NDS32\_17IFC\_PCREL

BFD\_RELOC\_NDS32\_10IFCU\_PCREL

For ex9 and ifc using.

BFD\_RELOC\_NDS32\_TPOFF

BFD\_RELOC\_NDS32\_TLS\_LE\_HI20

BFD\_RELOC\_NDS32\_TLS\_LE\_LO12

BFD\_RELOC\_NDS32\_TLS\_LE\_ADD

BFD\_RELOC\_NDS32\_TLS\_LE\_LS

BFD\_RELOC\_NDS32\_GOTTPOFF

BFD\_RELOC\_NDS32\_TLS\_IE\_HI20

BFD\_RELOC\_NDS32\_TLS\_IE\_L012S2

BFD\_RELOC\_NDS32\_TLS\_TPOFF

BFD\_RELOC\_NDS32\_TLS\_LE\_20

BFD\_RELOC\_NDS32\_TLS\_LE\_15S0

BFD\_RELOC\_NDS32\_TLS\_LE\_15S1

BFD\_RELOC\_NDS32\_TLS\_LE\_15S2

For TLS.

## BFD\_RELOC\_V850\_9\_PCREL

This is a 9-bit reloc

# BFD\_RELOC\_V850\_22\_PCREL

This is a 22-bit reloc

#### BFD\_RELOC\_V850\_SDA\_16\_16\_OFFSET

This is a 16 bit offset from the short data area pointer.

#### BFD\_RELOC\_V850\_SDA\_15\_16\_OFFSET

This is a 16 bit offset (of which only 15 bits are used) from the short data area pointer.

## BFD\_RELOC\_V850\_ZDA\_16\_16\_0FFSET

This is a 16 bit offset from the zero data area pointer.

## BFD\_RELOC\_V850\_ZDA\_15\_16\_OFFSET

This is a 16 bit offset (of which only 15 bits are used) from the zero data area pointer.

#### BFD\_RELOC\_V850\_TDA\_6\_8\_OFFSET

This is an 8 bit offset (of which only 6 bits are used) from the tiny data area pointer.

# BFD\_RELOC\_V850\_TDA\_7\_8\_OFFSET

This is an 8bit offset (of which only 7 bits are used) from the tiny data area pointer.

#### BFD\_RELOC\_V850\_TDA\_7\_7\_OFFSET

This is a 7 bit offset from the tiny data area pointer.

# BFD\_RELOC\_V850\_TDA\_16\_16\_0FFSET

This is a 16 bit offset from the tiny data area pointer.

## BFD\_RELOC\_V850\_TDA\_4\_5\_OFFSET

This is a 5 bit offset (of which only 4 bits are used) from the tiny data area pointer.

# BFD\_RELOC\_V850\_TDA\_4\_4\_OFFSET

This is a 4 bit offset from the tiny data area pointer.

# BFD\_RELOC\_V850\_SDA\_16\_16\_SPLIT\_OFFSET

This is a 16 bit offset from the short data area pointer, with the bits placed non-contiguously in the instruction.

# BFD\_RELOC\_V850\_ZDA\_16\_16\_SPLIT\_OFFSET

This is a 16 bit offset from the zero data area pointer, with the bits placed non-contiguously in the instruction.

## BFD\_RELOC\_V850\_CALLT\_6\_7\_OFFSET

This is a 6 bit offset from the call table base pointer.

## BFD\_RELOC\_V850\_CALLT\_16\_16\_0FFSET

This is a 16 bit offset from the call table base pointer.

#### BFD\_RELOC\_V850\_LONGCALL

Used for relaxing indirect function calls.

## BFD\_RELOC\_V850\_LONGJUMP

Used for relaxing indirect jumps.

#### BFD\_RELOC\_V850\_ALIGN

Used to maintain alignment whilst relaxing.

## BFD\_RELOC\_V850\_L016\_SPLIT\_OFFSET

This is a variation of BFD\_RELOC\_LO16 that can be used in v850e ld.bu instructions.

# BFD\_RELOC\_V850\_16\_PCREL

This is a 16-bit reloc.

#### BFD\_RELOC\_V850\_17\_PCREL

This is a 17-bit reloc.

## BFD\_RELOC\_V850\_23

This is a 23-bit reloc.

#### BFD\_RELOC\_V850\_32\_PCREL

This is a 32-bit reloc.

# BFD\_RELOC\_V850\_32\_ABS

This is a 32-bit reloc.

## BFD\_RELOC\_V850\_16\_SPLIT\_OFFSET

This is a 16-bit reloc.

## BFD\_RELOC\_V850\_16\_S1

This is a 16-bit reloc.

# BFD\_RELOC\_V850\_L016\_S1

Low 16 bits. 16 bit shifted by 1.

## BFD\_RELOC\_V850\_CALLT\_15\_16\_OFFSET

This is a 16 bit offset from the call table base pointer.

# BFD\_RELOC\_V850\_32\_GOTPCREL

DSO relocations.

# BFD\_RELOC\_V850\_16\_GOT

DSO relocations.

## BFD\_RELOC\_V850\_32\_GOT

DSO relocations.

# BFD\_RELOC\_V850\_22\_PLT\_PCREL

DSO relocations.

# BFD\_RELOC\_V850\_32\_PLT\_PCREL

DSO relocations.

## BFD\_RELOC\_V850\_COPY

DSO relocations.

## BFD\_RELOC\_V850\_GLOB\_DAT

DSO relocations.

# BFD\_RELOC\_V850\_JMP\_SLOT

DSO relocations.

## BFD\_RELOC\_V850\_RELATIVE

DSO relocations.

# BFD\_RELOC\_V850\_16\_GOTOFF

DSO relocations.

# BFD\_RELOC\_V850\_32\_GOTOFF

DSO relocations.

# BFD\_RELOC\_V850\_CODE

start code.

## BFD\_RELOC\_V850\_DATA

start data in text.

# BFD\_RELOC\_TIC3O\_LDP

This is a 8bit DP reloc for the tms320c30, where the most significant 8 bits of a 24 bit word are placed into the least significant 8 bits of the opcode.

#### BFD\_RELOC\_TIC54X\_PARTLS7

This is a 7bit reloc for the tms320c54x, where the least significant 7 bits of a 16 bit word are placed into the least significant 7 bits of the opcode.

## BFD\_RELOC\_TIC54X\_PARTMS9

This is a 9bit DP reloc for the tms320c54x, where the most significant 9 bits of a 16 bit word are placed into the least significant 9 bits of the opcode.

# BFD\_RELOC\_TIC54X\_23

This is an extended address 23-bit reloc for the tms320c54x.

## BFD\_RELOC\_TIC54X\_16\_0F\_23

This is a 16-bit reloc for the tms320c54x, where the least significant 16 bits of a 23-bit extended address are placed into the opcode.

## BFD\_RELOC\_TIC54X\_MS7\_OF\_23

This is a reloc for the tms320c54x, where the most significant 7 bits of a 23-bit extended address are placed into the opcode.

```
BFD_RELOC_C6000_PCR_S21
BFD_RELOC_C6000_PCR_S12
BFD_RELOC_C6000_PCR_S10
BFD_RELOC_C6000_PCR_S7
BFD_RELOC_C6000_ABS_S16
BFD_RELOC_C6000_ABS_L16
BFD_RELOC_C6000_ABS_H16
BFD_RELOC_C6000_SBR_U15_B
BFD_RELOC_C6000_SBR_U15_H
BFD_RELOC_C6000_SBR_U15_W
BFD_RELOC_C6000_SBR_S16
BFD_RELOC_C6000_SBR_L16_B
BFD_RELOC_C6000_SBR_L16_H
BFD_RELOC_C6000_SBR_L16_W
BFD_RELOC_C6000_SBR_H16_B
BFD_RELOC_C6000_SBR_H16_H
BFD_RELOC_C6000_SBR_H16_W
BFD_RELOC_C6000_SBR_GOT_U15_W
BFD_RELOC_C6000_SBR_GOT_L16_W
BFD_RELOC_C6000_SBR_GOT_H16_W
BFD_RELOC_C6000_DSBT_INDEX
BFD_RELOC_C6000_PREL31
BFD_RELOC_C6000_COPY
BFD_RELOC_C6000_JUMP_SLOT
BFD_RELOC_C6000_EHTYPE
```

BFD\_RELOC\_C6000\_PCR\_H16

BFD\_RELOC\_C6000\_PCR\_L16

BFD\_RELOC\_C6000\_ALIGN

BFD\_RELOC\_C6000\_FPHEAD

BFD\_RELOC\_C6000\_NOCMP

TMS320C6000 relocations.

#### BFD\_RELOC\_FR30\_48

This is a 48 bit reloc for the FR30 that stores 32 bits.

## BFD\_RELOC\_FR30\_20

This is a 32 bit reloc for the FR30 that stores 20 bits split up into two sections.

#### BFD\_RELOC\_FR30\_6\_IN\_4

This is a 16 bit reloc for the FR30 that stores a 6 bit word offset in 4 bits.

#### BFD\_RELOC\_FR30\_8\_IN\_8

This is a 16 bit reloc for the FR30 that stores an 8 bit byte offset into 8 bits.

### BFD\_RELOC\_FR30\_9\_IN\_8

This is a 16 bit reloc for the FR30 that stores a 9 bit short offset into 8 bits.

## BFD\_RELOC\_FR30\_10\_IN\_8

This is a 16 bit reloc for the FR30 that stores a 10 bit word offset into 8 bits.

## BFD\_RELOC\_FR30\_9\_PCREL

This is a 16 bit reloc for the FR30 that stores a 9 bit pc relative short offset into 8 bits.

#### BFD\_RELOC\_FR30\_12\_PCREL

This is a 16 bit reloc for the FR30 that stores a 12 bit pc relative short offset into 11 bits.

BFD\_RELOC\_MCORE\_PCREL\_IMM8BY4

BFD\_RELOC\_MCORE\_PCREL\_IMM11BY2

BFD\_RELOC\_MCORE\_PCREL\_IMM4BY2

BFD\_RELOC\_MCORE\_PCREL\_32

BFD\_RELOC\_MCORE\_PCREL\_JSR\_IMM11BY2

BFD\_RELOC\_MCORE\_RVA

Motorola Mcore relocations.

BFD\_RELOC\_MEP\_8

BFD\_RELOC\_MEP\_16

BFD\_RELOC\_MEP\_32

BFD\_RELOC\_MEP\_PCREL8A2

BFD\_RELOC\_MEP\_PCREL12A2

BFD\_RELOC\_MEP\_PCREL17A2

BFD\_RELOC\_MEP\_PCREL24A2

BFD\_RELOC\_MEP\_PCABS24A2

BFD\_RELOC\_MEP\_LOW16

BFD\_RELOC\_MEP\_HI16U

```
BFD_RELOC_MEP_HI16S
BFD_RELOC_MEP_GPREL
BFD_RELOC_MEP_TPREL
BFD_RELOC_MEP_TPREL7
BFD_RELOC_MEP_TPREL7A2
BFD_RELOC_MEP_TPREL7A4
BFD_RELOC_MEP_UIMM24
BFD_RELOC_MEP_ADDR24A4
BFD_RELOC_MEP_GNU_VTINHERIT
BFD_RELOC_MEP_GNU_VTENTRY
    Toshiba Media Processor Relocations.
BFD_RELOC_METAG_HIADDR16
BFD_RELOC_METAG_LOADDR16
BFD_RELOC_METAG_RELBRANCH
BFD_RELOC_METAG_GETSETOFF
BFD_RELOC_METAG_HIOG
BFD_RELOC_METAG_LOOG
BFD_RELOC_METAG_REL8
BFD_RELOC_METAG_REL16
BFD_RELOC_METAG_HI16_GOTOFF
```

BFD\_RELOC\_METAG\_LO16\_GOTOFF BFD\_RELOC\_METAG\_GETSET\_GOTOFF BFD\_RELOC\_METAG\_HI16\_GOTPC BFD\_RELOC\_METAG\_LO16\_GOTPC BFD\_RELOC\_METAG\_HI16\_PLT BFD\_RELOC\_METAG\_LO16\_PLT

BFD\_RELOC\_METAG\_RELBRANCH\_PLT

BFD\_RELOC\_METAG\_TLS\_LDO\_HI16 BFD\_RELOC\_METAG\_TLS\_LDO\_L016

BFD\_RELOC\_METAG\_TLS\_IENONPIC

BFD\_RELOC\_METAG\_TLS\_TPOFF BFD\_RELOC\_METAG\_TLS\_DTPMOD BFD\_RELOC\_METAG\_TLS\_DTPOFF

BFD\_RELOC\_METAG\_TLS\_IENONPIC\_HI16 BFD\_RELOC\_METAG\_TLS\_IENONPIC\_L016

BFD\_RELOC\_METAG\_TLS\_LDO
BFD\_RELOC\_METAG\_TLS\_IE

BFD\_RELOC\_METAG\_GOTOFF
BFD\_RELOC\_METAG\_PLT
BFD\_RELOC\_METAG\_COPY
BFD\_RELOC\_METAG\_JMP\_SLOT
BFD\_RELOC\_METAG\_RELATIVE
BFD\_RELOC\_METAG\_GLOB\_DAT
BFD\_RELOC\_METAG\_TLS\_GD
BFD\_RELOC\_METAG\_TLS\_LDM

```
BFD_RELOC_METAG_TLS_LE
BFD_RELOC_METAG_TLS_LE_HI16
BFD_RELOC_METAG_TLS_LE_L016
     Imagination Technologies Meta relocations.
BFD_RELOC_MMIX_GETA
BFD_RELOC_MMIX_GETA_1
BFD_RELOC_MMIX_GETA_2
BFD_RELOC_MMIX_GETA_3
     These are relocations for the GETA instruction.
BFD_RELOC_MMIX_CBRANCH
BFD_RELOC_MMIX_CBRANCH_J
BFD_RELOC_MMIX_CBRANCH_1
BFD_RELOC_MMIX_CBRANCH_2
BFD_RELOC_MMIX_CBRANCH_3
     These are relocations for a conditional branch instruction.
BFD_RELOC_MMIX_PUSHJ
BFD_RELOC_MMIX_PUSHJ_1
BFD_RELOC_MMIX_PUSHJ_2
BFD_RELOC_MMIX_PUSHJ_3
BFD_RELOC_MMIX_PUSHJ_STUBBABLE
     These are relocations for the PUSHJ instruction.
BFD_RELOC_MMIX_JMP
BFD_RELOC_MMIX_JMP_1
```

These are relocations for the JMP instruction.

#### BFD\_RELOC\_MMIX\_ADDR19

BFD\_RELOC\_MMIX\_JMP\_2 BFD\_RELOC\_MMIX\_JMP\_3

This is a relocation for a relative address as in a GETA instruction or a branch.

#### BFD\_RELOC\_MMIX\_ADDR27

This is a relocation for a relative address as in a JMP instruction.

#### BFD\_RELOC\_MMIX\_REG\_OR\_BYTE

This is a relocation for an instruction field that may be a general register or a value 0..255.

#### BFD\_RELOC\_MMIX\_REG

This is a relocation for an instruction field that may be a general register.

#### BFD\_RELOC\_MMIX\_BASE\_PLUS\_OFFSET

This is a relocation for two instruction fields holding a register and an offset, the equivalent of the relocation.

## BFD\_RELOC\_MMIX\_LOCAL

This relocation is an assertion that the expression is not allocated as a global register. It does not modify contents.

#### BFD\_RELOC\_AVR\_7\_PCREL

This is a 16 bit reloc for the AVR that stores 8 bit pc relative short offset into 7 bits.

#### BFD\_RELOC\_AVR\_13\_PCREL

This is a 16 bit reloc for the AVR that stores 13 bit pc relative short offset into 12 bits.

# BFD\_RELOC\_AVR\_16\_PM

This is a 16 bit reloc for the AVR that stores 17 bit value (usually program memory address) into 16 bits.

## BFD\_RELOC\_AVR\_LO8\_LDI

This is a 16 bit reloc for the AVR that stores 8 bit value (usually data memory address) into 8 bit immediate value of LDI insn.

## BFD\_RELOC\_AVR\_HI8\_LDI

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of data memory address) into 8 bit immediate value of LDI insn.

#### BFD\_RELOC\_AVR\_HH8\_LDI

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of program memory address) into 8 bit immediate value of LDI insn.

## BFD\_RELOC\_AVR\_MS8\_LDI

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of 32 bit value) into 8 bit immediate value of LDI insn.

## BFD\_RELOC\_AVR\_LO8\_LDI\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (usually data memory address) into 8 bit immediate value of SUBI insn.

#### BFD\_RELOC\_AVR\_HI8\_LDI\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 8 bit of data memory address) into 8 bit immediate value of SUBI insn.

#### BFD\_RELOC\_AVR\_HH8\_LDI\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (most high 8 bit of program memory address) into 8 bit immediate value of LDI or SUBI insn.

#### BFD\_RELOC\_AVR\_MS8\_LDI\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (msb of 32 bit value) into 8 bit immediate value of LDI insn.

## BFD\_RELOC\_AVR\_LO8\_LDI\_PM

This is a 16 bit reloc for the AVR that stores 8 bit value (usually command address) into 8 bit immediate value of LDI insn.

# BFD\_RELOC\_AVR\_LO8\_LDI\_GS

This is a 16 bit reloc for the AVR that stores 8 bit value (command address) into 8 bit immediate value of LDI insn. If the address is beyond the 128k boundary, the linker inserts a jump stub for this reloc in the lower 128k.

## BFD\_RELOC\_AVR\_HI8\_LDI\_PM

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of command address) into 8 bit immediate value of LDI insn.

## BFD\_RELOC\_AVR\_HI8\_LDI\_GS

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of command address) into 8 bit immediate value of LDI insn. If the address is beyond the 128k boundary, the linker inserts a jump stub for this reloc below 128k.

#### BFD\_RELOC\_AVR\_HH8\_LDI\_PM

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of command address) into 8 bit immediate value of LDI insn.

#### BFD\_RELOC\_AVR\_LO8\_LDI\_PM\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (usually command address) into 8 bit immediate value of SUBI insn.

## BFD\_RELOC\_AVR\_HI8\_LDI\_PM\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 8 bit of 16 bit command address) into 8 bit immediate value of SUBI insn.

## BFD\_RELOC\_AVR\_HH8\_LDI\_PM\_NEG

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 6 bit of 22 bit command address) into 8 bit immediate value of SUBI insn.

## BFD\_RELOC\_AVR\_CALL

This is a 32 bit reloc for the AVR that stores 23 bit value into 22 bits.

#### BFD\_RELOC\_AVR\_LDI

This is a 16 bit reloc for the AVR that stores all needed bits for absolute addressing with ldi with overflow check to linktime

#### BFD\_RELOC\_AVR\_6

This is a 6 bit reloc for the AVR that stores offset for ldd/std instructions

#### BFD\_RELOC\_AVR\_6\_ADIW

This is a 6 bit reloc for the AVR that stores offset for adiw/sbiw instructions

#### BFD\_RELOC\_AVR\_8\_LO

This is a 8 bit reloc for the AVR that stores bits 0..7 of a symbol in .byte lo8(symbol)

#### BFD\_RELOC\_AVR\_8\_HI

This is a 8 bit reloc for the AVR that stores bits 8..15 of a symbol in .byte hi8(symbol)

## BFD\_RELOC\_AVR\_8\_HLO

This is a 8 bit reloc for the AVR that stores bits 16..23 of a symbol in .byte hlo8(symbol)

## BFD\_RELOC\_AVR\_DIFF8

BFD\_RELOC\_AVR\_DIFF16

# BFD\_RELOC\_AVR\_DIFF32

AVR relocations to mark the difference of two local symbols. These are only needed to support linker relaxation and can be ignored when not relaxing. The field is set to the value of the difference assuming no relaxation. The relocation encodes the position of the second symbol so the linker can determine whether to adjust the field value.

## BFD\_RELOC\_AVR\_LDS\_STS\_16

This is a 7 bit reloc for the AVR that stores SRAM address for 16bit lds and sts instructions supported only tiny core.

#### BFD\_RELOC\_AVR\_PORT6

This is a 6 bit reloc for the AVR that stores an I/O register number for the IN and OUT instructions

# BFD\_RELOC\_AVR\_PORT5

This is a 5 bit reloc for the AVR that stores an I/O register number for the SBIC, SBIS, SBI and CBI instructions

```
BFD_RELOC_RL78_NEG8
BFD_RELOC_RL78_NEG16
BFD_RELOC_RL78_NEG24
BFD_RELOC_RL78_NEG32
BFD_RELOC_RL78_16_OP
BFD_RELOC_RL78_24_OP
BFD_RELOC_RL78_32_OP
BFD_RELOC_RL78_8U
BFD_RELOC_RL78_16U
BFD_RELOC_RL78_24U
BFD_RELOC_RL78_DIR3U_PCREL
BFD_RELOC_RL78_DIFF
BFD_RELOC_RL78_GPRELB
BFD_RELOC_RL78_GPRELW
BFD_RELOC_RL78_GPRELL
BFD_RELOC_RL78_SYM
BFD_RELOC_RL78_OP_SUBTRACT
BFD_RELOC_RL78_OP_NEG
BFD_RELOC_RL78_OP_AND
BFD_RELOC_RL78_OP_SHRA
BFD_RELOC_RL78_ABS8
BFD_RELOC_RL78_ABS16
BFD_RELOC_RL78_ABS16_REV
BFD_RELOC_RL78_ABS32
BFD_RELOC_RL78_ABS32_REV
BFD_RELOC_RL78_ABS16U
BFD_RELOC_RL78_ABS16UW
```

BFD\_RELOC\_RL78\_ABS16UL BFD\_RELOC\_RL78\_RELAX BFD\_RELOC\_RL78\_HI16

BFD\_RELOC\_RL78\_HI8

BFD\_RELOC\_RL78\_L016

BFD\_RELOC\_RL78\_CODE

Renesas RL78 Relocations.

BFD\_RELOC\_RX\_NEG8

BFD\_RELOC\_RX\_NEG16

BFD\_RELOC\_RX\_NEG24

BFD\_RELOC\_RX\_NEG32

BFD\_RELOC\_RX\_16\_OP

BFD\_RELOC\_RX\_24\_OP

BFD\_RELOC\_RX\_32\_OP

BFD\_RELOC\_RX\_8U

BFD\_RELOC\_RX\_16U

BFD\_RELOC\_RX\_24U

BFD\_RELOC\_RX\_DIR3U\_PCREL

BFD\_RELOC\_RX\_DIFF

BFD\_RELOC\_RX\_GPRELB

BFD\_RELOC\_RX\_GPRELW

BFD\_RELOC\_RX\_GPRELL

BFD\_RELOC\_RX\_SYM

BFD\_RELOC\_RX\_OP\_SUBTRACT

BFD\_RELOC\_RX\_OP\_NEG

BFD\_RELOC\_RX\_ABS8

BFD\_RELOC\_RX\_ABS16

BFD\_RELOC\_RX\_ABS16\_REV

BFD\_RELOC\_RX\_ABS32

BFD\_RELOC\_RX\_ABS32\_REV

BFD\_RELOC\_RX\_ABS16U

BFD\_RELOC\_RX\_ABS16UW

BFD\_RELOC\_RX\_ABS16UL

BFD\_RELOC\_RX\_RELAX

Renesas RX Relocations.

BFD\_RELOC\_390\_12

Direct 12 bit.

BFD\_RELOC\_390\_GOT12

12 bit GOT offset.

BFD\_RELOC\_390\_PLT32

32 bit PC relative PLT address.

BFD\_RELOC\_390\_COPY

Copy symbol at runtime.

BFD\_RELOC\_390\_GLOB\_DAT

Create GOT entry.

# BFD\_RELOC\_390\_JMP\_SLOT Create PLT entry.

# BFD\_RELOC\_390\_RELATIVE Adjust by program base.

BFD\_RELOC\_390\_GOTPC
32 bit PC relative offset to GOT.

BFD\_RELOC\_390\_GOT16 16 bit GOT offset.

BFD\_RELOC\_390\_PC12DBL PC relative 12 bit shifted by 1.

BFD\_RELOC\_390\_PLT12DBL 12 bit PC rel. PLT shifted by 1.

BFD\_RELOC\_390\_PC16DBL PC relative 16 bit shifted by 1.

BFD\_RELOC\_390\_PLT16DBL 16 bit PC rel. PLT shifted by 1.

BFD\_RELOC\_390\_PC24DBL PC relative 24 bit shifted by 1.

BFD\_RELOC\_390\_PLT24DBL 24 bit PC rel. PLT shifted by 1.

BFD\_RELOC\_390\_PC32DBL PC relative 32 bit shifted by 1.

BFD\_RELOC\_390\_GOTPCDBL 32 bit PC rel. GOT shifted by 1.

BFD\_RELOC\_390\_G0T64 64 bit GOT offset.

BFD\_RELOC\_390\_PLT64 64 bit PC relative PLT address.

BFD\_RELOC\_390\_GOTENT
32 bit rel. offset to GOT entry.

BFD\_RELOC\_390\_GOTOFF64 64 bit offset to GOT.

BFD\_RELOC\_390\_GOTPLT12 12-bit offset to symbol-entry within GOT, with PLT handling.

## BFD\_RELOC\_390\_GOTPLT16

16-bit offset to symbol-entry within GOT, with PLT handling.

## BFD\_RELOC\_390\_GOTPLT32

32-bit offset to symbol-entry within GOT, with PLT handling.

#### BFD\_RELOC\_390\_GOTPLT64

64-bit offset to symbol-entry within GOT, with PLT handling.

## BFD\_RELOC\_390\_GOTPLTENT

32-bit rel. offset to symbol-entry within GOT, with PLT handling.

## BFD\_RELOC\_390\_PLT0FF16

16-bit rel. offset from the GOT to a PLT entry.

#### BFD\_RELOC\_390\_PLT0FF32

32-bit rel. offset from the GOT to a PLT entry.

## BFD\_RELOC\_390\_PLT0FF64

64-bit rel. offset from the GOT to a PLT entry.

## BFD\_RELOC\_390\_TLS\_LOAD

BFD\_RELOC\_390\_TLS\_GDCALL

BFD\_RELOC\_390\_TLS\_LDCALL

BFD\_RELOC\_390\_TLS\_GD32

BFD\_RELOC\_390\_TLS\_GD64

BFD\_RELOC\_390\_TLS\_GOTIE12

BFD\_RELOC\_390\_TLS\_GOTIE32

BFD\_RELOC\_390\_TLS\_GOTIE64

BFD\_RELOC\_390\_TLS\_LDM32

BFD\_RELOC\_390\_TLS\_LDM64

BFD\_RELOC\_390\_TLS\_IE32

BFD\_RELOC\_390\_TLS\_IE64

BFD\_RELOC\_390\_TLS\_IEENT

BFD\_RELOC\_390\_TLS\_LE32

BFD\_RELOC\_390\_TLS\_LE64

BFD\_RELOC\_390\_TLS\_LD032

BFD\_RELOC\_390\_TLS\_LD064

BFD\_RELOC\_390\_TLS\_DTPMOD

BFD\_RELOC\_390\_TLS\_DTPOFF

BFD\_RELOC\_390\_TLS\_TPOFF

s390 tls relocations.

BFD\_RELOC\_390\_20

BFD\_RELOC\_390\_GOT20

BFD\_RELOC\_390\_GOTPLT20

BFD\_RELOC\_390\_TLS\_GOTIE20

Long displacement extension.

#### BFD\_RELOC\_390\_IRELATIVE

STT\_GNU\_IFUNC relocation.

## BFD\_RELOC\_SCORE\_GPREL15

Score relocations Low 16 bit for load/store

#### BFD\_RELOC\_SCORE\_DUMMY2

## BFD\_RELOC\_SCORE\_JMP

This is a 24-bit reloc with the right 1 bit assumed to be 0

#### BFD\_RELOC\_SCORE\_BRANCH

This is a 19-bit reloc with the right 1 bit assumed to be 0

#### BFD\_RELOC\_SCORE\_IMM30

This is a 32-bit reloc for 48-bit instructions.

## BFD\_RELOC\_SCORE\_IMM32

This is a 32-bit reloc for 48-bit instructions.

## BFD\_RELOC\_SCORE16\_JMP

This is a 11-bit reloc with the right 1 bit assumed to be 0

#### BFD\_RELOC\_SCORE16\_BRANCH

This is a 8-bit reloc with the right 1 bit assumed to be 0

## BFD\_RELOC\_SCORE\_BCMP

This is a 9-bit reloc with the right 1 bit assumed to be 0

## BFD\_RELOC\_SCORE\_GOT15

BFD\_RELOC\_SCORE\_GOT\_L016

BFD\_RELOC\_SCORE\_CALL15

#### BFD\_RELOC\_SCORE\_DUMMY\_HI16

Undocumented Score relocs

## BFD\_RELOC\_IP2K\_FR9

Scenix IP2K - 9-bit register number / data address

#### BFD\_RELOC\_IP2K\_BANK

Scenix IP2K - 4-bit register/data bank number

## BFD\_RELOC\_IP2K\_ADDR16CJP

Scenix IP2K - low 13 bits of instruction word address

#### BFD\_RELOC\_IP2K\_PAGE3

Scenix IP2K - high 3 bits of instruction word address

## BFD\_RELOC\_IP2K\_LO8DATA

BFD\_RELOC\_IP2K\_HI8DATA

#### BFD\_RELOC\_IP2K\_EX8DATA

Scenix IP2K - ext/low/high 8 bits of data address

## BFD\_RELOC\_IP2K\_LO8INSN

#### BFD\_RELOC\_IP2K\_HI8INSN

Scenix IP2K - low/high 8 bits of instruction word address

BFD\_RELOC\_IP2K\_PC\_SKIP

Scenix IP2K - even/odd PC modifier to modify snb pcl.0

BFD\_RELOC\_IP2K\_TEXT

Scenix IP2K - 16 bit word address in text section.

BFD\_RELOC\_IP2K\_FR\_OFFSET

Scenix IP2K - 7-bit sp or dp offset

BFD\_RELOC\_VPE4KMATH\_DATA

BFD\_RELOC\_VPE4KMATH\_INSN

Scenix VPE4K coprocessor - data/insn-space addressing

BFD\_RELOC\_VTABLE\_INHERIT

BFD\_RELOC\_VTABLE\_ENTRY

These two relocations are used by the linker to determine which of the entries in a C++ virtual function table are actually used. When the -gc-sections option is given, the linker will zero out the entries that are not used, so that the code for those functions need not be included in the output.

VTABLE\_INHERIT is a zero-space relocation used to describe to the linker the inheritance tree of a C++ virtual function table. The relocation's symbol should be the parent class' vtable, and the relocation should be located at the child vtable.

VTABLE\_ENTRY is a zero-space relocation that describes the use of a virtual function table entry. The reloc's symbol should refer to the table of the class mentioned in the code. Off of that base, an offset describes the entry that is being used. For Rela hosts, this offset is stored in the reloc's addend. For Rel hosts, we are forced to put this offset in the reloc's section offset.

BFD\_RELOC\_IA64\_IMM14

BFD\_RELOC\_IA64\_IMM22

BFD\_RELOC\_IA64\_IMM64

BFD\_RELOC\_IA64\_DIR32MSB

BFD\_RELOC\_IA64\_DIR32LSB

BFD\_RELOC\_IA64\_DIR64MSB

BFD\_RELOC\_IA64\_DIR64LSB

BFD\_RELOC\_IA64\_GPREL22

BFD\_RELOC\_IA64\_GPREL64I

BFD\_RELOC\_IA64\_GPREL32MSB

BFD\_RELOC\_IA64\_GPREL32LSB BFD\_RELOC\_IA64\_GPREL64MSB

BFD\_RELOC\_IA64\_GPREL64LSB

BFD\_RELOC\_IA64\_LT0FF22

BFD\_RELOC\_IA64\_LT0FF64I

BFD\_RELOC\_IA64\_PLT0FF22

BFD\_RELOC\_IA64\_PLT0FF64I

BFD\_RELOC\_IA64\_PLT0FF64MSB

BFD\_RELOC\_IA64\_PLT0FF64LSB

BFD\_RELOC\_IA64\_FPTR64I

```
BFD_RELOC_IA64_FPTR32MSB
BFD_RELOC_IA64_FPTR32LSB
BFD_RELOC_IA64_FPTR64MSB
BFD_RELOC_IA64_FPTR64LSB
BFD_RELOC_IA64_PCREL21B
BFD_RELOC_IA64_PCREL21BI
BFD_RELOC_IA64_PCREL21M
BFD_RELOC_IA64_PCREL21F
BFD_RELOC_IA64_PCREL22
BFD_RELOC_IA64_PCREL60B
BFD_RELOC_IA64_PCREL64I
BFD_RELOC_IA64_PCREL32MSB
BFD_RELOC_IA64_PCREL32LSB
BFD_RELOC_IA64_PCREL64MSB
BFD_RELOC_IA64_PCREL64LSB
BFD_RELOC_IA64_LT0FF_FPTR22
BFD_RELOC_IA64_LTOFF_FPTR64I
BFD_RELOC_IA64_LTOFF_FPTR32MSB
BFD_RELOC_IA64_LT0FF_FPTR32LSB
BFD_RELOC_IA64_LTOFF_FPTR64MSB
BFD_RELOC_IA64_LTOFF_FPTR64LSB
BFD_RELOC_IA64_SEGREL32MSB
BFD_RELOC_IA64_SEGREL32LSB
BFD_RELOC_IA64_SEGREL64MSB
BFD_RELOC_IA64_SEGREL64LSB
BFD_RELOC_IA64_SECREL32MSB
BFD_RELOC_IA64_SECREL32LSB
BFD_RELOC_IA64_SECREL64MSB
BFD_RELOC_IA64_SECREL64LSB
BFD_RELOC_IA64_REL32MSB
BFD_RELOC_IA64_REL32LSB
BFD_RELOC_IA64_REL64MSB
BFD_RELOC_IA64_REL64LSB
BFD_RELOC_IA64_LTV32MSB
BFD_RELOC_IA64_LTV32LSB
BFD_RELOC_IA64_LTV64MSB
BFD_RELOC_IA64_LTV64LSB
BFD_RELOC_IA64_IPLTMSB
BFD_RELOC_IA64_IPLTLSB
BFD_RELOC_IA64_COPY
BFD_RELOC_IA64_LT0FF22X
BFD_RELOC_IA64_LDXMOV
BFD_RELOC_IA64_TPREL14
BFD_RELOC_IA64_TPREL22
BFD_RELOC_IA64_TPREL64I
BFD_RELOC_IA64_TPREL64MSB
```

BFD\_RELOC\_IA64\_TPREL64LSB

```
BFD_RELOC_IA64_LT0FF_TPREL22
```

BFD\_RELOC\_IA64\_DTPMOD64MSB

BFD\_RELOC\_IA64\_DTPMOD64LSB

BFD\_RELOC\_IA64\_LTOFF\_DTPMOD22

BFD\_RELOC\_IA64\_DTPREL14

BFD\_RELOC\_IA64\_DTPREL22

BFD\_RELOC\_IA64\_DTPREL64I

BFD\_RELOC\_IA64\_DTPREL32MSB

BFD\_RELOC\_IA64\_DTPREL32LSB

BFD\_RELOC\_IA64\_DTPREL64MSB

BFD\_RELOC\_IA64\_DTPREL64LSB

BFD\_RELOC\_IA64\_LTOFF\_DTPREL22

Intel IA64 Relocations.

## BFD\_RELOC\_M68HC11\_HI8

Motorola 68HC11 reloc. This is the 8 bit high part of an absolute address.

#### BFD\_RELOC\_M68HC11\_L08

Motorola 68HC11 reloc. This is the 8 bit low part of an absolute address.

## BFD\_RELOC\_M68HC11\_3B

Motorola 68HC11 reloc. This is the 3 bit of a value.

#### BFD\_RELOC\_M68HC11\_RL\_JUMP

Motorola 68HC11 reloc. This reloc marks the beginning of a jump/call instruction. It is used for linker relaxation to correctly identify beginning of instruction and change some branches to use PC-relative addressing mode.

#### BFD\_RELOC\_M68HC11\_RL\_GROUP

Motorola 68HC11 reloc. This reloc marks a group of several instructions that gcc generates and for which the linker relaxation pass can modify and/or remove some of them.

#### BFD\_RELOC\_M68HC11\_L016

Motorola 68HC11 reloc. This is the 16-bit lower part of an address. It is used for 'call' instruction to specify the symbol address without any special transformation (due to memory bank window).

## BFD\_RELOC\_M68HC11\_PAGE

Motorola 68HC11 reloc. This is a 8-bit reloc that specifies the page number of an address. It is used by 'call' instruction to specify the page number of the symbol.

## BFD\_RELOC\_M68HC11\_24

Motorola 68HC11 reloc. This is a 24-bit reloc that represents the address with a 16-bit value and a 8-bit page number. The symbol address is transformed to follow the 16K memory bank of 68HC12 (seen as mapped in the window).

#### BFD\_RELOC\_M68HC12\_5B

Motorola 68HC12 reloc. This is the 5 bits of a value.

## BFD\_RELOC\_XGATE\_RL\_JUMP

Freescale XGATE reloc. This reloc marks the beginning of a bra/jal instruction.

#### BFD\_RELOC\_XGATE\_RL\_GROUP

Freescale XGATE reloc. This reloc marks a group of several instructions that gcc generates and for which the linker relaxation pass can modify and/or remove some of them.

## BFD\_RELOC\_XGATE\_LO16

Freescale XGATE reloc. This is the 16-bit lower part of an address. It is used for the '16-bit' instructions.

#### BFD\_RELOC\_XGATE\_GPAGE

Freescale XGATE reloc.

## BFD\_RELOC\_XGATE\_24

Freescale XGATE reloc.

#### BFD\_RELOC\_XGATE\_PCREL\_9

Freescale XGATE reloc. This is a 9-bit pc-relative reloc.

## BFD\_RELOC\_XGATE\_PCREL\_10

Freescale XGATE reloc. This is a 10-bit pc-relative reloc.

#### BFD\_RELOC\_XGATE\_IMM8\_LO

Freescale XGATE reloc. This is the 16-bit lower part of an address. It is used for the '16-bit' instructions.

#### BFD\_RELOC\_XGATE\_IMM8\_HI

Freescale XGATE reloc. This is the 16-bit higher part of an address. It is used for the '16-bit' instructions.

#### BFD\_RELOC\_XGATE\_IMM3

Freescale XGATE reloc. This is a 3-bit pc-relative reloc.

#### BFD\_RELOC\_XGATE\_IMM4

Freescale XGATE reloc. This is a 4-bit pc-relative reloc.

## BFD\_RELOC\_XGATE\_IMM5

Freescale XGATE reloc. This is a 5-bit pc-relative reloc.

#### BFD\_RELOC\_M68HC12\_9B

Motorola 68HC12 reloc. This is the 9 bits of a value.

#### BFD\_RELOC\_M68HC12\_16B

Motorola 68HC12 reloc. This is the 16 bits of a value.

## BFD\_RELOC\_M68HC12\_9\_PCREL

Motorola 68HC12/XGATE reloc. This is a PCREL9 branch.

#### BFD\_RELOC\_M68HC12\_10\_PCREL

Motorola 68HC12/XGATE reloc. This is a PCREL10 branch.

## BFD\_RELOC\_M68HC12\_L08XG

Motorola 68HC12/XGATE reloc. This is the 8 bit low part of an absolute address and immediately precedes a matching HI8XG part.

## BFD\_RELOC\_M68HC12\_HI8XG

Motorola 68HC12/XGATE reloc. This is the 8 bit high part of an absolute address and immediately follows a matching LO8XG part.

```
BFD_RELOC_16C_NUMO8
BFD_RELOC_16C_NUMO8_C
BFD_RELOC_16C_NUM16
BFD_RELOC_16C_NUM16_C
BFD_RELOC_16C_NUM32
BFD_RELOC_16C_NUM32_C
BFD_RELOC_16C_DISPO4
BFD_RELOC_16C_DISPO4_C
BFD_RELOC_16C_DISP08
BFD_RELOC_16C_DISP08_C
BFD_RELOC_16C_DISP16
BFD_RELOC_16C_DISP16_C
BFD_RELOC_16C_DISP24
BFD_RELOC_16C_DISP24_C
BFD_RELOC_16C_DISP24a
BFD_RELOC_16C_DISP24a_C
BFD_RELOC_16C_REG04
BFD_RELOC_16C_REGO4_C
BFD_RELOC_16C_REG04a
BFD_RELOC_16C_REG04a_C
BFD_RELOC_16C_REG14
BFD_RELOC_16C_REG14_C
BFD_RELOC_16C_REG16
BFD_RELOC_16C_REG16_C
BFD_RELOC_16C_REG20
BFD_RELOC_16C_REG2O_C
BFD_RELOC_16C_ABS20
BFD_RELOC_16C_ABS20_C
BFD_RELOC_16C_ABS24
BFD_RELOC_16C_ABS24_C
BFD_RELOC_16C_IMMO4
BFD_RELOC_16C_IMMO4_C
BFD_RELOC_16C_IMM16
BFD_RELOC_16C_IMM16_C
BFD_RELOC_16C_IMM20
BFD_RELOC_16C_IMM2O_C
BFD_RELOC_16C_IMM24
BFD_RELOC_16C_IMM24_C
```

BFD\_RELOC\_16C\_IMM32

BFD\_RELOC\_16C\_IMM32\_C NS CR16C Relocations.

BFD\_RELOC\_CR16\_NUM8

BFD\_RELOC\_CR16\_NUM16

BFD\_RELOC\_CR16\_NUM32

BFD\_RELOC\_CR16\_NUM32a

BFD\_RELOC\_CR16\_REGRELO

BFD\_RELOC\_CR16\_REGREL4

BFD\_RELOC\_CR16\_REGREL4a

BFD\_RELOC\_CR16\_REGREL14

BFD\_RELOC\_CR16\_REGREL14a

BFD\_RELOC\_CR16\_REGREL16

BFD\_RELOC\_CR16\_REGREL20

BFD\_RELOC\_CR16\_REGREL20a

BFD\_RELOC\_CR16\_ABS20

BFD\_RELOC\_CR16\_ABS24

BFD\_RELOC\_CR16\_IMM4

BFD\_RELOC\_CR16\_IMM8

BFD\_RELOC\_CR16\_IMM16

BFD\_RELOC\_CR16\_IMM20

BFD\_RELOC\_CR16\_IMM24

BFD\_RELOC\_CR16\_IMM32

BFD\_RELOC\_CR16\_IMM32a

BFD\_RELOC\_CR16\_DISP4

BFD\_RELOC\_CR16\_DISP8

BFD\_RELOC\_CR16\_DISP16

BFD\_RELOC\_CR16\_DISP20

BFD\_RELOC\_CR16\_DISP24

BFD\_RELOC\_CR16\_DISP24a

BFD\_RELOC\_CR16\_SWITCH8

BFD\_RELOC\_CR16\_SWITCH16

BFD\_RELOC\_CR16\_SWITCH32

BFD\_RELOC\_CR16\_GOT\_REGREL20

BFD\_RELOC\_CR16\_GOTC\_REGREL20

BFD\_RELOC\_CR16\_GLOB\_DAT

NS CR16 Relocations.

BFD\_RELOC\_CRX\_REL4

BFD\_RELOC\_CRX\_REL8

BFD\_RELOC\_CRX\_REL8\_CMP

BFD\_RELOC\_CRX\_REL16

BFD\_RELOC\_CRX\_REL24

BFD\_RELOC\_CRX\_REL32

BFD\_RELOC\_CRX\_REGREL12

BFD\_RELOC\_CRX\_REGREL22

BFD\_RELOC\_CRX\_REGREL28

BFD\_RELOC\_CRX\_REGREL32

```
BFD_RELOC_CRX_ABS16
BFD_RELOC_CRX_ABS32
BFD_RELOC_CRX_NUM8
BFD_RELOC_CRX_NUM16
BFD_RELOC_CRX_NUM32
BFD_RELOC_CRX_IMM16
BFD_RELOC_CRX_IMM32
BFD_RELOC_CRX_SWITCH8
BFD_RELOC_CRX_SWITCH16
BFD_RELOC_CRX_SWITCH32
     NS CRX Relocations.
BFD_RELOC_CRIS_BDISP8
BFD_RELOC_CRIS_UNSIGNED_5
BFD_RELOC_CRIS_SIGNED_6
BFD_RELOC_CRIS_UNSIGNED_6
BFD_RELOC_CRIS_SIGNED_8
BFD_RELOC_CRIS_UNSIGNED_8
BFD_RELOC_CRIS_SIGNED_16
BFD_RELOC_CRIS_UNSIGNED_16
BFD_RELOC_CRIS_LAPCQ_OFFSET
BFD_RELOC_CRIS_UNSIGNED_4
     These relocs are only used within the CRIS assembler. They are not (at present)
     written to any object files.
BFD_RELOC_CRIS_COPY
BFD_RELOC_CRIS_GLOB_DAT
BFD_RELOC_CRIS_JUMP_SLOT
BFD_RELOC_CRIS_RELATIVE
     Relocs used in ELF shared libraries for CRIS.
BFD_RELOC_CRIS_32_GOT
     32-bit offset to symbol-entry within GOT.
BFD_RELOC_CRIS_16_GOT
     16-bit offset to symbol-entry within GOT.
BFD_RELOC_CRIS_32_GOTPLT
     32-bit offset to symbol-entry within GOT, with PLT handling.
BFD_RELOC_CRIS_16_GOTPLT
     16-bit offset to symbol-entry within GOT, with PLT handling.
BFD_RELOC_CRIS_32_GOTREL
     32-bit offset to symbol, relative to GOT.
BFD_RELOC_CRIS_32_PLT_GOTREL
     32-bit offset to symbol with PLT entry, relative to GOT.
BFD_RELOC_CRIS_32_PLT_PCREL
```

32-bit offset to symbol with PLT entry, relative to this relocation.

```
BFD_RELOC_CRIS_32_GOT_GD
BFD_RELOC_CRIS_16_GOT_GD
BFD_RELOC_CRIS_32_GD
BFD_RELOC_CRIS_DTP
BFD_RELOC_CRIS_32_DTPREL
BFD_RELOC_CRIS_16_DTPREL
BFD_RELOC_CRIS_32_GOT_TPREL
BFD_RELOC_CRIS_16_GOT_TPREL
BFD_RELOC_CRIS_32_TPREL
BFD_RELOC_CRIS_16_TPREL
BFD_RELOC_CRIS_DTPMOD
BFD_RELOC_CRIS_32_IE
    Relocs used in TLS code for CRIS.
BFD_RELOC_860_COPY
BFD_RELOC_860_GLOB_DAT
BFD_RELOC_860_JUMP_SLOT
BFD_RELOC_860_RELATIVE
BFD_RELOC_860_PC26
BFD_RELOC_860_PLT26
BFD_RELOC_860_PC16
BFD_RELOC_860_LOWO
BFD_RELOC_860_SPLITO
BFD_RELOC_860_LOW1
BFD_RELOC_860_SPLIT1
BFD_RELOC_860_LOW2
BFD_RELOC_860_SPLIT2
BFD_RELOC_860_LOW3
BFD_RELOC_860_LOGOTO
BFD_RELOC_860_SPGOTO
BFD_RELOC_860_LOGOT1
BFD_RELOC_860_SPGOT1
BFD_RELOC_860_LOGOTOFFO
BFD_RELOC_860_SPGOTOFFO
BFD_RELOC_860_LOGOTOFF1
BFD_RELOC_860_SPGOTOFF1
BFD_RELOC_860_LOGOTOFF2
BFD_RELOC_860_LOGOTOFF3
```

BFD\_RELOC\_860\_LOPC
BFD\_RELOC\_860\_HIGHADJ
BFD\_RELOC\_860\_HAGOT
BFD\_RELOC\_860\_HAGOTOFF
BFD\_RELOC\_860\_HAPC
BFD\_RELOC\_860\_HIGH
BFD\_RELOC\_860\_HIGOT
BFD\_RELOC\_860\_HIGOTOFF
Intel i860 Relocations.

```
BFD_RELOC_OR1K_REL_26
BFD_RELOC_OR1K_GOTPC_HI16
BFD_RELOC_OR1K_GOTPC_LO16
BFD_RELOC_OR1K_GOT16
BFD_RELOC_OR1K_PLT26
BFD_RELOC_OR1K_GOTOFF_HI16
BFD_RELOC_OR1K_GOTOFF_LO16
BFD_RELOC_OR1K_COPY
BFD_RELOC_OR1K_GLOB_DAT
BFD_RELOC_OR1K_JMP_SLOT
BFD_RELOC_OR1K_RELATIVE
BFD_RELOC_OR1K_TLS_GD_HI16
BFD_RELOC_OR1K_TLS_GD_LO16
BFD_RELOC_OR1K_TLS_LDM_HI16
BFD_RELOC_OR1K_TLS_LDM_LO16
BFD_RELOC_OR1K_TLS_LDO_HI16
BFD_RELOC_OR1K_TLS_LDO_LO16
BFD_RELOC_OR1K_TLS_IE_HI16
BFD_RELOC_OR1K_TLS_IE_LO16
BFD_RELOC_OR1K_TLS_LE_HI16
BFD_RELOC_OR1K_TLS_LE_LO16
BFD_RELOC_OR1K_TLS_TPOFF
BFD_RELOC_OR1K_TLS_DTPOFF
BFD_RELOC_OR1K_TLS_DTPMOD
    OpenRISC 1000 Relocations.
```

BFD\_RELOC\_H8\_DIR16A8
BFD\_RELOC\_H8\_DIR16R8
BFD\_RELOC\_H8\_DIR24A8
BFD\_RELOC\_H8\_DIR24R8
BFD\_RELOC\_H8\_DIR32A16
BFD\_RELOC\_H8\_DISP32A16
H8 elf Relocations.

BFD\_RELOC\_XSTORMY16\_REL\_12

BFD\_RELOC\_XSTORMY16\_12 BFD\_RELOC\_XSTORMY16\_24

BFD\_RELOC\_XSTORMY16\_FPTR16

Sony Xstormy16 Relocations.

BFD\_RELOC\_RELC

Self-describing complex relocations.

BFD\_RELOC\_XC16X\_PAG

BFD\_RELOC\_XC16X\_POF

BFD\_RELOC\_XC16X\_SEG

BFD\_RELOC\_XC16X\_SOF

Infineon Relocations.

BFD\_RELOC\_VAX\_GLOB\_DAT

BFD\_RELOC\_VAX\_JMP\_SLOT

BFD\_RELOC\_VAX\_RELATIVE

Relocations used by VAX ELF.

## BFD\_RELOC\_MT\_PC16

Morpho MT - 16 bit immediate relocation.

## BFD\_RELOC\_MT\_HI16

Morpho MT - Hi 16 bits of an address.

## BFD\_RELOC\_MT\_L016

Morpho MT - Low 16 bits of an address.

## BFD\_RELOC\_MT\_GNU\_VTINHERIT

Morpho MT - Used to tell the linker which vtable entries are used.

## BFD\_RELOC\_MT\_GNU\_VTENTRY

Morpho MT - Used to tell the linker which vtable entries are used.

#### BFD\_RELOC\_MT\_PCINSN8

Morpho MT - 8 bit immediate relocation.

BFD\_RELOC\_MSP430\_10\_PCREL

BFD\_RELOC\_MSP430\_16\_PCREL

BFD\_RELOC\_MSP430\_16

BFD\_RELOC\_MSP430\_16\_PCREL\_BYTE

BFD\_RELOC\_MSP430\_16\_BYTE

BFD\_RELOC\_MSP430\_2X\_PCREL

BFD\_RELOC\_MSP430\_RL\_PCREL

BFD\_RELOC\_MSP430\_ABS8

BFD\_RELOC\_MSP430X\_PCR20\_EXT\_SRC

BFD\_RELOC\_MSP430X\_PCR20\_EXT\_DST

BFD\_RELOC\_MSP430X\_PCR20\_EXT\_ODST

BFD\_RELOC\_MSP430X\_ABS20\_EXT\_SRC

BFD\_RELOC\_MSP430X\_ABS20\_EXT\_DST

BFD\_RELOC\_MSP430X\_ABS20\_EXT\_ODST

BFD\_RELOC\_MSP430X\_ABS20\_ADR\_SRC

BFD\_RELOC\_MSP430X\_ABS20\_ADR\_DST

BFD\_RELOC\_MSP430X\_PCR16

BFD\_RELOC\_MSP430X\_PCR20\_CALL

BFD\_RELOC\_MSP430X\_ABS16

BFD\_RELOC\_MSP430\_ABS\_HI16

BFD\_RELOC\_MSP430\_PREL31

BFD\_RELOC\_MSP430\_SYM\_DIFF

msp430 specific relocation codes

BFD\_RELOC\_NIOS2\_S16

BFD\_RELOC\_NIOS2\_U16

BFD\_RELOC\_NIOS2\_CALL26

```
BFD_RELOC_NIOS2_IMM5
BFD_RELOC_NIOS2_CACHE_OPX
BFD_RELOC_NIOS2_IMM6
BFD_RELOC_NIOS2_IMM8
BFD_RELOC_NIOS2_HI16
BFD_RELOC_NIOS2_LO16
BFD_RELOC_NIOS2_HIADJ16
BFD_RELOC_NIOS2_GPREL
BFD_RELOC_NIOS2_UJMP
BFD_RELOC_NIOS2_CJMP
BFD_RELOC_NIOS2_CALLR
BFD_RELOC_NIOS2_ALIGN
BFD_RELOC_NIOS2_GOT16
BFD_RELOC_NIOS2_CALL16
BFD_RELOC_NIOS2_GOTOFF_LO
BFD_RELOC_NIOS2_GOTOFF_HA
BFD_RELOC_NIOS2_PCREL_LO
BFD_RELOC_NIOS2_PCREL_HA
BFD_RELOC_NIOS2_TLS_GD16
BFD_RELOC_NIOS2_TLS_LDM16
BFD_RELOC_NIOS2_TLS_LD016
BFD_RELOC_NIOS2_TLS_IE16
BFD_RELOC_NIOS2_TLS_LE16
BFD_RELOC_NIOS2_TLS_DTPMOD
BFD_RELOC_NIOS2_TLS_DTPREL
BFD_RELOC_NIOS2_TLS_TPREL
BFD_RELOC_NIOS2_COPY
BFD_RELOC_NIOS2_GLOB_DAT
BFD_RELOC_NIOS2_JUMP_SLOT
BFD_RELOC_NIOS2_RELATIVE
BFD_RELOC_NIOS2_GOTOFF
BFD_RELOC_NIOS2_CALL26_NOAT
BFD_RELOC_NIOS2_GOT_LO
BFD_RELOC_NIOS2_GOT_HA
BFD_RELOC_NIOS2_CALL_LO
BFD_RELOC_NIOS2_CALL_HA
    Relocations used by the Altera Nios II core.
BFD_RELOC_IQ2000_OFFSET_16
BFD_RELOC_IQ2000_OFFSET_21
BFD_RELOC_IQ2000_UHI16
```

## BFD\_RELOC\_XTENSA\_RTLD

IQ2000 Relocations.

Special Xtensa relocation used only by PLT entries in ELF shared objects to indicate that the runtime linker should set the value to one of its own internal functions or data structures.

```
BFD_RELOC_XTENSA_GLOB_DAT
BFD_RELOC_XTENSA_JMP_SLOT
BFD_RELOC_XTENSA_RELATIVE
```

Xtensa relocations for ELF shared objects.

#### BFD\_RELOC\_XTENSA\_PLT

Xtensa relocation used in ELF object files for symbols that may require PLT entries. Otherwise, this is just a generic 32-bit relocation.

```
BFD_RELOC_XTENSA_DIFF8
BFD_RELOC_XTENSA_DIFF16
BFD_RELOC_XTENSA_DIFF32
```

Xtensa relocations to mark the difference of two local symbols. These are only needed to support linker relaxation and can be ignored when not relaxing. The field is set to the value of the difference assuming no relaxation. The relocation encodes the position of the first symbol so the linker can determine whether to adjust the field value.

```
BFD_RELOC_XTENSA_SLOTO_OP
BFD_RELOC_XTENSA_SLOT1_OP
BFD_RELOC_XTENSA_SLOT2_OP
BFD_RELOC_XTENSA_SLOT3_OP
BFD_RELOC_XTENSA_SLOT4_OP
BFD_RELOC_XTENSA_SLOT5_OP
BFD_RELOC_XTENSA_SLOT6_OP
BFD_RELOC_XTENSA_SLOT7_OP
BFD_RELOC_XTENSA_SLOT7_OP
BFD_RELOC_XTENSA_SLOT9_OP
BFD_RELOC_XTENSA_SLOT10_OP
BFD_RELOC_XTENSA_SLOT11_OP
BFD_RELOC_XTENSA_SLOT11_OP
BFD_RELOC_XTENSA_SLOT112_OP
BFD_RELOC_XTENSA_SLOT13_OP
BFD_RELOC_XTENSA_SLOT13_OP
BFD_RELOC_XTENSA_SLOT14_OP
```

Generic Xtensa relocations for instruction operands. Only the slot number is encoded in the relocation. The relocation applies to the last PC-relative immediate operand, or if there are no PC-relative immediates, to the last immediate operand.

```
BFD_RELOC_XTENSA_SLOTO_ALT
BFD_RELOC_XTENSA_SLOT1_ALT
BFD_RELOC_XTENSA_SLOT2_ALT
BFD_RELOC_XTENSA_SLOT3_ALT
BFD_RELOC_XTENSA_SLOT4_ALT
BFD_RELOC_XTENSA_SLOT5_ALT
BFD_RELOC_XTENSA_SLOT6_ALT
BFD_RELOC_XTENSA_SLOT6_ALT
BFD_RELOC_XTENSA_SLOT7_ALT
BFD_RELOC_XTENSA_SLOT8_ALT
BFD_RELOC_XTENSA_SLOT9_ALT
BFD_RELOC_XTENSA_SLOT9_ALT
BFD_RELOC_XTENSA_SLOT10_ALT
```

BFD\_RELOC\_XTENSA\_SLOT11\_ALT

BFD\_RELOC\_XTENSA\_SLOT12\_ALT

BFD\_RELOC\_XTENSA\_SLOT13\_ALT

BFD\_RELOC\_XTENSA\_SLOT14\_ALT

Alternate Xtensa relocations. Only the slot is encoded in the relocation. The meaning of these relocations is opcode-specific.

BFD\_RELOC\_XTENSA\_OPO

BFD\_RELOC\_XTENSA\_OP1

BFD\_RELOC\_XTENSA\_OP2

Xtensa relocations for backward compatibility. These have all been replaced by BFD\_RELOC\_XTENSA\_SLOT0\_OP.

#### BFD\_RELOC\_XTENSA\_ASM\_EXPAND

Xtensa relocation to mark that the assembler expanded the instructions from an original target. The expansion size is encoded in the reloc size.

## BFD\_RELOC\_XTENSA\_ASM\_SIMPLIFY

Xtensa relocation to mark that the linker should simplify assembler-expanded instructions. This is commonly used internally by the linker after analysis of a BFD\_RELOC\_XTENSA\_ASM\_EXPAND.

BFD\_RELOC\_XTENSA\_TLSDESC\_FN

BFD\_RELOC\_XTENSA\_TLSDESC\_ARG

BFD\_RELOC\_XTENSA\_TLS\_DTPOFF

BFD\_RELOC\_XTENSA\_TLS\_TPOFF

BFD\_RELOC\_XTENSA\_TLS\_FUNC

BFD\_RELOC\_XTENSA\_TLS\_ARG

BFD\_RELOC\_XTENSA\_TLS\_CALL

Xtensa TLS relocations.

## BFD\_RELOC\_Z80\_DISP8

8 bit signed offset in (ix+d) or (iy+d).

BFD\_RELOC\_Z8K\_DISP7

DJNZ offset.

BFD\_RELOC\_Z8K\_CALLR

CALR offset.

BFD\_RELOC\_Z8K\_IMM4L

4 bit value.

BFD\_RELOC\_LM32\_CALL

BFD\_RELOC\_LM32\_BRANCH

BFD\_RELOC\_LM32\_16\_GOT

BFD\_RELOC\_LM32\_GOTOFF\_HI16

BFD\_RELOC\_LM32\_GOTOFF\_L016

BFD\_RELOC\_LM32\_COPY

BFD\_RELOC\_LM32\_GLOB\_DAT

## BFD\_RELOC\_LM32\_JMP\_SLOT

## BFD\_RELOC\_LM32\_RELATIVE

Lattice Mico32 relocations.

## BFD\_RELOC\_MACH\_O\_SECTDIFF

Difference between two section addreses. Must be followed by a BFD\_RELOC\_MACH\_O\_PAIR.

## BFD\_RELOC\_MACH\_O\_LOCAL\_SECTDIFF

Like BFD\_RELOC\_MACH\_O\_SECTDIFF but with a local symbol.

#### BFD\_RELOC\_MACH\_O\_PAIR

Pair of relocation. Contains the first symbol.

#### BFD\_RELOC\_MACH\_O\_X86\_64\_BRANCH32

## BFD\_RELOC\_MACH\_O\_X86\_64\_BRANCH8

PCREL relocations. They are marked as branch to create PLT entry if required.

#### BFD\_RELOC\_MACH\_O\_X86\_64\_GOT

Used when referencing a GOT entry.

## BFD\_RELOC\_MACH\_O\_X86\_64\_GOT\_LOAD

Used when loading a GOT entry with movq. It is specially marked so that the linker could optimize the movq to a leaq if possible.

## BFD\_RELOC\_MACH\_O\_X86\_64\_SUBTRACTOR32

Symbol will be substracted. Must be followed by a BFD\_RELOC\_64.

#### BFD\_RELOC\_MACH\_O\_X86\_64\_SUBTRACTOR64

Symbol will be substracted. Must be followed by a BFD\_RELOC\_64.

#### BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_1

Same as BFD\_RELOC\_32\_PCREL but with an implicit -1 addend.

## BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_2

Same as BFD\_RELOC\_32\_PCREL but with an implicit -2 addend.

## BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_4

Same as BFD\_RELOC\_32\_PCREL but with an implicit -4 addend.

#### BFD\_RELOC\_MICROBLAZE\_32\_LO

This is a 32 bit reloc for the microblaze that stores the low 16 bits of a value

## BFD\_RELOC\_MICROBLAZE\_32\_LO\_PCREL

This is a 32 bit pc-relative reloc for the microblaze that stores the low 16 bits of a value

## BFD\_RELOC\_MICROBLAZE\_32\_ROSDA

This is a 32 bit reloc for the microblaze that stores a value relative to the read-only small data area anchor

#### BFD\_RELOC\_MICROBLAZE\_32\_RWSDA

This is a 32 bit reloc for the microblaze that stores a value relative to the read-write small data area anchor

## BFD\_RELOC\_MICROBLAZE\_32\_SYM\_OP\_SYM

This is a 32 bit reloc for the microblaze to handle expressions of the form "Symbol Op Symbol"

#### BFD\_RELOC\_MICROBLAZE\_64\_NONE

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). No relocation is done here - only used for relaxing

### BFD\_RELOC\_MICROBLAZE\_64\_GOTPC

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is PC-relative GOT offset

## BFD\_RELOC\_MICROBLAZE\_64\_GOT

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is GOT offset

### BFD\_RELOC\_MICROBLAZE\_64\_PLT

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is PC-relative offset into PLT

#### BFD\_RELOC\_MICROBLAZE\_64\_GOTOFF

This is a 64 bit reloc that stores the 32 bit GOT relative value in two words (with an imm instruction). The relocation is relative offset from \_GLOBAL\_OFFSET\_TABLE\_

## BFD\_RELOC\_MICROBLAZE\_32\_GOTOFF

This is a 32 bit reloc that stores the 32 bit GOT relative value in a word. The relocation is relative offset from

## BFD\_RELOC\_MICROBLAZE\_COPY

This is used to tell the dynamic linker to copy the value out of the dynamic object into the runtime process image.

## BFD\_RELOC\_MICROBLAZE\_64\_TLS

Unused Reloc

#### BFD\_RELOC\_MICROBLAZE\_64\_TLSGD

This is a 64 bit reloc that stores the 32 bit GOT relative value of the GOT TLS GD info entry in two words (with an imm instruction). The relocation is GOT offset.

## BFD\_RELOC\_MICROBLAZE\_64\_TLSLD

This is a 64 bit reloc that stores the 32 bit GOT relative value of the GOT TLS LD info entry in two words (with an imm instruction). The relocation is GOT offset.

#### BFD\_RELOC\_MICROBLAZE\_32\_TLSDTPMOD

This is a 32 bit reloc that stores the Module ID to GOT(n).

## BFD\_RELOC\_MICROBLAZE\_32\_TLSDTPREL

This is a 32 bit reloc that stores TLS offset to GOT(n+1).

#### BFD\_RELOC\_MICROBLAZE\_64\_TLSDTPREL

This is a 32 bit reloc for storing TLS offset to two words (uses imm instruction)

## BFD\_RELOC\_MICROBLAZE\_64\_TLSGOTTPREL

This is a 64 bit reloc that stores 32-bit thread pointer relative offset to two words (uses imm instruction).

## BFD\_RELOC\_MICROBLAZE\_64\_TLSTPREL

This is a 64 bit reloc that stores 32-bit thread pointer relative offset to two words (uses imm instruction).

## BFD\_RELOC\_AARCH64\_RELOC\_START

AArch64 pseudo relocation code to mark the start of the AArch64 relocation enumerators. N.B. the order of the enumerators is important as several tables in the AArch64 bfd backend are indexed by these enumerators; make sure they are all synced.

#### BFD\_RELOC\_AARCH64\_NONE

AArch64 null relocation code.

## BFD\_RELOC\_AARCH64\_64

BFD\_RELOC\_AARCH64\_32

## BFD\_RELOC\_AARCH64\_16

Basic absolute relocations of N bits. These are equivalent to BFD\_RELOC\_N and they were added to assist the indexing of the howto table.

## BFD\_RELOC\_AARCH64\_64\_PCREL

BFD\_RELOC\_AARCH64\_32\_PCREL

#### BFD\_RELOC\_AARCH64\_16\_PCREL

PC-relative relocations. These are equivalent to BFD\_RELOC\_N\_PCREL and they were added to assist the indexing of the howto table.

#### BFD\_RELOC\_AARCH64\_MOVW\_GO

AArch64 MOV[NZK] instruction with most significant bits 0 to 15 of an unsigned address/value.

## BFD\_RELOC\_AARCH64\_MOVW\_GO\_NC

AArch64 MOV[NZK] instruction with less significant bits 0 to 15 of an address/value. No overflow checking.

## BFD\_RELOC\_AARCH64\_MOVW\_G1

AArch64 MOV[NZK] instruction with most significant bits 16 to 31 of an unsigned address/value.

## BFD\_RELOC\_AARCH64\_MOVW\_G1\_NC

AArch64 MOV[NZK] instruction with less significant bits 16 to 31 of an address/value. No overflow checking.

## BFD\_RELOC\_AARCH64\_MOVW\_G2

AArch64 MOV[NZK] instruction with most significant bits 32 to 47 of an unsigned address/value.

#### BFD\_RELOC\_AARCH64\_MOVW\_G2\_NC

AArch64 MOV[NZK] instruction with less significant bits 32 to 47 of an address/value. No overflow checking.

### BFD\_RELOC\_AARCH64\_MOVW\_G3

AArch64 MOV[NZK] instruction with most signficant bits 48 to 64 of a signed or unsigned address/value.

## BFD\_RELOC\_AARCH64\_MOVW\_GO\_S

AArch64 MOV[NZ] instruction with most significant bits 0 to 15 of a signed value. Changes instruction to MOVZ or MOVN depending on the value's sign.

## BFD\_RELOC\_AARCH64\_MOVW\_G1\_S

AArch64 MOV[NZ] instruction with most significant bits 16 to 31 of a signed value. Changes instruction to MOVZ or MOVN depending on the value's sign.

#### BFD\_RELOC\_AARCH64\_MOVW\_G2\_S

AArch64 MOV[NZ] instruction with most significant bits 32 to 47 of a signed value. Changes instruction to MOVZ or MOVN depending on the value's sign.

#### BFD\_RELOC\_AARCH64\_LD\_L019\_PCREL

AArch64 Load Literal instruction, holding a 19 bit pc-relative word offset. The lowest two bits must be zero and are not stored in the instruction, giving a 21 bit signed byte offset.

## BFD\_RELOC\_AARCH64\_ADR\_LO21\_PCREL

AArch64 ADR instruction, holding a simple 21 bit pc-relative byte offset.

## BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL

AArch64 ADRP instruction, with bits 12 to 32 of a pc-relative page offset, giving a 4KB aligned page base address.

## BFD\_RELOC\_AARCH64\_ADR\_HI21\_NC\_PCREL

AArch64 ADRP instruction, with bits 12 to 32 of a pc-relative page offset, giving a 4KB aligned page base address, but with no overflow checking.

## BFD\_RELOC\_AARCH64\_ADD\_L012

AArch64 ADD immediate instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

## BFD\_RELOC\_AARCH64\_LDST8\_L012

AArch64 8-bit load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

## BFD\_RELOC\_AARCH64\_TSTBR14

AArch64 14 bit pc-relative test bit and branch. The lowest two bits must be zero and are not stored in the instruction, giving a 16 bit signed byte offset.

## BFD\_RELOC\_AARCH64\_BRANCH19

AArch64 19 bit pc-relative conditional branch and compare & branch. The lowest two bits must be zero and are not stored in the instruction, giving a 21 bit signed byte offset.

#### BFD\_RELOC\_AARCH64\_JUMP26

AArch64 26 bit pc-relative unconditional branch. The lowest two bits must be zero and are not stored in the instruction, giving a 28 bit signed byte offset.

#### BFD\_RELOC\_AARCH64\_CALL26

AArch64 26 bit pc-relative unconditional branch and link. The lowest two bits must be zero and are not stored in the instruction, giving a 28 bit signed byte offset.

## BFD\_RELOC\_AARCH64\_LDST16\_L012

AArch64 16-bit load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

#### BFD\_RELOC\_AARCH64\_LDST32\_L012

AArch64 32-bit load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

### BFD\_RELOC\_AARCH64\_LDST64\_L012

AArch64 64-bit load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

#### BFD\_RELOC\_AARCH64\_LDST128\_L012

AArch64 128-bit load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

#### BFD\_RELOC\_AARCH64\_GOT\_LD\_PREL19

AArch64 Load Literal instruction, holding a 19 bit PC relative word offset of the global offset table entry for a symbol. The lowest two bits must be zero and are not stored in the instruction, giving a 21 bit signed byte offset. This relocation type requires signed overflow checking.

## BFD\_RELOC\_AARCH64\_ADR\_GOT\_PAGE

Get to the page base of the global offset table entry for a symbol as part of an ADRP instruction using a 21 bit PC relative value. Used in conjunction with BFD\_RELOC\_AARCH64\_LD64\_GOT\_LO12\_NC.

#### BFD\_RELOC\_AARCH64\_LD64\_GOT\_LO12\_NC

Unsigned 12 bit byte offset for 64 bit load/store from the page of the GOT entry for this symbol. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_GOTPAGE. Valid in LP64 ABI only.

#### BFD\_RELOC\_AARCH64\_LD32\_GOT\_LO12\_NC

Unsigned 12 bit byte offset for 32 bit load/store from the page of the GOT entry for this symbol. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_GOTPAGE. Valid in ILP32 ABI only.

- BFD\_RELOC\_AARCH64\_TLSGD\_ADR\_PAGE21
  - Get to the page base of the global offset table entry for a symbols tls\_index structure as part of an adrp instruction using a 21 bit PC relative value. Used in conjunction with BFD\_RELOC\_AARCH64\_TLSGD\_ADD\_LO12\_NC.
- BFD\_RELOC\_AARCH64\_TLSGD\_ADD\_LO12\_NC

  Unsigned 12 bit byte offset to global offset table entry for a symbols tls\_index structure. Used in conjunction with BFD\_RELOC\_AARCH64\_TLSGD\_ADR\_PAGE21.
- BFD\_RELOC\_AARCH64\_TLSIE\_MOVW\_GOTTPREL\_G1
   AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSIE\_MOVW\_GOTTPREL\_GO\_NC AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSIE\_ADR\_GOTTPREL\_PAGE21 AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSIE\_LD64\_GOTTPREL\_LO12\_NC AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSIE\_LD32\_GOTTPREL\_LO12\_NC AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSIE\_LD\_GOTTPREL\_PREL19
  AArch64 TLS INITIAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_MOVW\_TPREL\_G2 AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_MOVW\_TPREL\_G1 AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_MOVW\_TPREL\_G1\_NC AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_MOVW\_TPREL\_GO AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_MOVW\_TPREL\_GO\_NC AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_ADD\_TPREL\_HI12 AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_ADD\_TPREL\_L012 AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSLE\_ADD\_TPREL\_L012\_NC AArch64 TLS LOCAL EXEC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_LD\_PREL19 AArch64 TLS DESC relocation.

- BFD\_RELOC\_AARCH64\_TLSDESC\_ADR\_PREL21 AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_ADR\_PAGE21 AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_LD64\_LO12\_NC AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_LD32\_L012\_NC AArch64 TLS DESC relocation.
- ${\tt BFD\_RELOC\_AARCH64\_TLSDESC\_ADD\_L012\_NC} \\ {\tt AArch64\ TLS\ DESC\ relocation}.$
- BFD\_RELOC\_AARCH64\_TLSDESC\_OFF\_G1 AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_OFF\_GO\_NC AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_LDR AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_ADD AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC\_CALL AArch64 TLS DESC relocation.
- BFD\_RELOC\_AARCH64\_COPY
   AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_GLOB\_DAT AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_JUMP\_SLOT AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_RELATIVE AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_TLS\_DTPMOD AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_TLS\_DTPREL AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_TLS\_TPREL AArch64 TLS relocation.
- BFD\_RELOC\_AARCH64\_TLSDESC AArch64 TLS relocation.

## BFD\_RELOC\_AARCH64\_IRELATIVE

AArch64 support for STT\_GNU\_IFUNC.

#### BFD\_RELOC\_AARCH64\_RELOC\_END

AArch64 pseudo relocation code to mark the end of the AArch64 relocation enumerators that have direct mapping to ELF reloc codes. There are a few more enumerators after this one; those are mainly used by the AArch64 assembler for the internal fixup or to select one of the above enumerators.

#### BFD\_RELOC\_AARCH64\_GAS\_INTERNAL\_FIXUP

AArch64 pseudo relocation code to be used internally by the AArch64 assembler and not (currently) written to any object files.

## BFD\_RELOC\_AARCH64\_LDST\_L012

AArch64 unspecified load/store instruction, holding bits 0 to 11 of the address. Used in conjunction with BFD\_RELOC\_AARCH64\_ADR\_HI21\_PCREL.

#### BFD\_RELOC\_AARCH64\_LD\_GOT\_LO12\_NC

AArch64 pseudo relocation code to be used internally by the AArch64 assembler and not (currently) written to any object files.

## BFD\_RELOC\_AARCH64\_TLSIE\_LD\_GOTTPREL\_LO12\_NC

AArch64 pseudo relocation code to be used internally by the AArch64 assembler and not (currently) written to any object files.

## BFD\_RELOC\_AARCH64\_TLSDESC\_LD\_L012\_NC

AArch64 pseudo relocation code to be used internally by the AArch64 assembler and not (currently) written to any object files.

```
BFD_RELOC_TILEPRO_COPY
BFD_RELOC_TILEPRO_GLOB_DAT
BFD_RELOC_TILEPRO_JMP_SLOT
BFD_RELOC_TILEPRO_RELATIVE
BFD_RELOC_TILEPRO_BROFF_X1
BFD_RELOC_TILEPRO_JOFFLONG_X1
BFD_RELOC_TILEPRO_JOFFLONG_X1_PLT
BFD_RELOC_TILEPRO_IMM8_XO
BFD_RELOC_TILEPRO_IMM8_YO
BFD_RELOC_TILEPRO_IMM8_X1
BFD_RELOC_TILEPRO_IMM8_Y1
BFD_RELOC_TILEPRO_DEST_IMM8_X1
BFD_RELOC_TILEPRO_MT_IMM15_X1
BFD_RELOC_TILEPRO_MF_IMM15_X1
BFD_RELOC_TILEPRO_IMM16_XO
BFD_RELOC_TILEPRO_IMM16_X1
BFD_RELOC_TILEPRO_IMM16_XO_LO
BFD_RELOC_TILEPRO_IMM16_X1_LO
BFD_RELOC_TILEPRO_IMM16_XO_HI
BFD_RELOC_TILEPRO_IMM16_X1_HI
```

```
BFD_RELOC_TILEPRO_IMM16_XO_HA
BFD_RELOC_TILEPRO_IMM16_X1_HA
BFD_RELOC_TILEPRO_IMM16_XO_PCREL
BFD_RELOC_TILEPRO_IMM16_X1_PCREL
BFD_RELOC_TILEPRO_IMM16_XO_LO_PCREL
BFD_RELOC_TILEPRO_IMM16_X1_LO_PCREL
BFD_RELOC_TILEPRO_IMM16_XO_HI_PCREL
BFD_RELOC_TILEPRO_IMM16_X1_HI_PCREL
BFD_RELOC_TILEPRO_IMM16_XO_HA_PCREL
BFD_RELOC_TILEPRO_IMM16_X1_HA_PCREL
BFD_RELOC_TILEPRO_IMM16_XO_GOT
BFD_RELOC_TILEPRO_IMM16_X1_GOT
BFD_RELOC_TILEPRO_IMM16_XO_GOT_LO
BFD_RELOC_TILEPRO_IMM16_X1_GOT_LO
BFD_RELOC_TILEPRO_IMM16_XO_GOT_HI
BFD_RELOC_TILEPRO_IMM16_X1_GOT_HI
BFD_RELOC_TILEPRO_IMM16_XO_GOT_HA
BFD_RELOC_TILEPRO_IMM16_X1_GOT_HA
BFD_RELOC_TILEPRO_MMSTART_XO
BFD_RELOC_TILEPRO_MMEND_XO
BFD_RELOC_TILEPRO_MMSTART_X1
BFD_RELOC_TILEPRO_MMEND_X1
BFD_RELOC_TILEPRO_SHAMT_XO
BFD_RELOC_TILEPRO_SHAMT_X1
BFD_RELOC_TILEPRO_SHAMT_YO
BFD_RELOC_TILEPRO_SHAMT_Y1
BFD_RELOC_TILEPRO_TLS_GD_CALL
BFD_RELOC_TILEPRO_IMM8_XO_TLS_GD_ADD
BFD_RELOC_TILEPRO_IMM8_X1_TLS_GD_ADD
BFD_RELOC_TILEPRO_IMM8_YO_TLS_GD_ADD
BFD_RELOC_TILEPRO_IMM8_Y1_TLS_GD_ADD
BFD_RELOC_TILEPRO_TLS_IE_LOAD
BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD
BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD
BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_LO
BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_LO
BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_HI
BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_HI
BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_HA
BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_HA
BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE
BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE
BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_LO
BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_LO
BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_HI
BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_HI
BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_HA
```

```
BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_HA
BFD_RELOC_TILEPRO_TLS_DTPMOD32
BFD_RELOC_TILEPRO_TLS_DTPOFF32
BFD_RELOC_TILEPRO_TLS_TPOFF32
BFD_RELOC_TILEPRO_IMM16_X0_TLS_LE
BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE
BFD_RELOC_TILEPRO_IMM16_X0_TLS_LE_LO
BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_LO
BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_LO
BFD_RELOC_TILEPRO_IMM16_X0_TLS_LE_HI
BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_HI
BFD_RELOC_TILEPRO_IMM16_X0_TLS_LE_HA
BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_HA
Tilera_TILEPRO_IMM16_X1_TLS_LE_HA
```

```
BFD_RELOC_TILEGX_HWO
BFD_RELOC_TILEGX_HW1
BFD_RELOC_TILEGX_HW2
BFD_RELOC_TILEGX_HW3
BFD_RELOC_TILEGX_HWO_LAST
BFD_RELOC_TILEGX_HW1_LAST
BFD_RELOC_TILEGX_HW2_LAST
BFD_RELOC_TILEGX_COPY
BFD_RELOC_TILEGX_GLOB_DAT
BFD_RELOC_TILEGX_JMP_SLOT
BFD_RELOC_TILEGX_RELATIVE
BFD_RELOC_TILEGX_BROFF_X1
BFD_RELOC_TILEGX_JUMPOFF_X1
BFD_RELOC_TILEGX_JUMPOFF_X1_PLT
BFD_RELOC_TILEGX_IMM8_XO
BFD_RELOC_TILEGX_IMM8_YO
BFD_RELOC_TILEGX_IMM8_X1
BFD_RELOC_TILEGX_IMM8_Y1
BFD_RELOC_TILEGX_DEST_IMM8_X1
BFD_RELOC_TILEGX_MT_IMM14_X1
BFD_RELOC_TILEGX_MF_IMM14_X1
BFD_RELOC_TILEGX_MMSTART_XO
BFD_RELOC_TILEGX_MMEND_XO
BFD_RELOC_TILEGX_SHAMT_XO
BFD_RELOC_TILEGX_SHAMT_X1
BFD_RELOC_TILEGX_SHAMT_YO
BFD_RELOC_TILEGX_SHAMT_Y1
BFD_RELOC_TILEGX_IMM16_XO_HWO
BFD_RELOC_TILEGX_IMM16_X1_HWO
BFD_RELOC_TILEGX_IMM16_XO_HW1
BFD_RELOC_TILEGX_IMM16_X1_HW1
BFD_RELOC_TILEGX_IMM16_XO_HW2
BFD_RELOC_TILEGX_IMM16_X1_HW2
```

```
BFD_RELOC_TILEGX_IMM16_XO_HW3
BFD_RELOC_TILEGX_IMM16_X1_HW3
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST
BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST
BFD_RELOC_TILEGX_IMM16_XO_HWO_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HWO_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW1_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW1_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW2_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW2_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW3_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW3_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_GOT
BFD_RELOC_TILEGX_IMM16_X1_HWO_GOT
BFD_RELOC_TILEGX_IMM16_XO_HWO_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HWO_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW1_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW1_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW2_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW2_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_GOT
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_GOT
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_GOT
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_GOT
BFD_RELOC_TILEGX_IMM16_XO_HW3_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW3_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_TLS_GD
BFD_RELOC_TILEGX_IMM16_X1_HWO_TLS_GD
BFD_RELOC_TILEGX_IMM16_XO_HWO_TLS_LE
BFD_RELOC_TILEGX_IMM16_X1_HWO_TLS_LE
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_LE
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_LE
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_LE
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_LE
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_GD
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_GD
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_GD
```

```
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_GD
BFD_RELOC_TILEGX_IMM16_XO_HWO_TLS_IE
BFD_RELOC_TILEGX_IMM16_X1_HWO_TLS_IE
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_IE
BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_IE
BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_IE
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_IE
BFD_RELOC_TILEGX_TLS_DTPMOD64
BFD_RELOC_TILEGX_TLS_DTP0FF64
BFD_RELOC_TILEGX_TLS_TP0FF64
BFD_RELOC_TILEGX_TLS_DTPMOD32
BFD_RELOC_TILEGX_TLS_DTPOFF32
BFD_RELOC_TILEGX_TLS_TP0FF32
BFD_RELOC_TILEGX_TLS_GD_CALL
BFD_RELOC_TILEGX_IMM8_XO_TLS_GD_ADD
BFD_RELOC_TILEGX_IMM8_X1_TLS_GD_ADD
BFD_RELOC_TILEGX_IMM8_YO_TLS_GD_ADD
BFD_RELOC_TILEGX_IMM8_Y1_TLS_GD_ADD
BFD_RELOC_TILEGX_TLS_IE_LOAD
BFD_RELOC_TILEGX_IMM8_XO_TLS_ADD
BFD_RELOC_TILEGX_IMM8_X1_TLS_ADD
BFD_RELOC_TILEGX_IMM8_YO_TLS_ADD
BFD_RELOC_TILEGX_IMM8_Y1_TLS_ADD
    Tilera TILE-Gx Relocations.
BFD_RELOC_EPIPHANY_SIMM8
    Adapteva EPIPHANY - 8 bit signed pc-relative displacement
BFD_RELOC_EPIPHANY_SIMM24
    Adapteva EPIPHANY - 24 bit signed pc-relative displacement
BFD RELOC EPIPHANY HIGH
    Adapteva EPIPHANY - 16 most-significant bits of absolute address
BFD_RELOC_EPIPHANY_LOW
    Adapteva EPIPHANY - 16 least-significant bits of absolute address
BFD_RELOC_EPIPHANY_SIMM11
    Adapteva EPIPHANY - 11 bit signed number - add/sub immediate
BFD_RELOC_EPIPHANY_IMM11
    Adapteva EPIPHANY - 11 bit sign-magnitude number (ld/st displacement)
```

## BFD\_RELOC\_EPIPHANY\_IMM8

Adapteva EPIPHANY - 8 bit immediate for 16 bit mov instruction.

```
typedef enum bfd_reloc_code_real bfd_reloc_code_real_type;
```

## 2.10.2.2 bfd\_reloc\_type\_lookup

## **Synopsis**

```
reloc_howto_type *bfd_reloc_type_lookup
   (bfd *abfd, bfd_reloc_code_real_type code);
reloc_howto_type *bfd_reloc_name_lookup
   (bfd *abfd, const char *reloc_name);
```

#### Description

Return a pointer to a howto structure which, when invoked, will perform the relocation code on data from the architecture noted.

## 2.10.2.3 bfd\_default\_reloc\_type\_lookup

#### **Synopsis**

```
reloc_howto_type *bfd_default_reloc_type_lookup
  (bfd *abfd, bfd_reloc_code_real_type code);
```

## Description

Provides a default relocation lookup routine for any architecture.

## 2.10.2.4 bfd\_get\_reloc\_code\_name

#### **Synopsis**

```
const char *bfd_get_reloc_code_name (bfd_reloc_code_real_type code);
```

#### Description

Provides a printable name for the supplied relocation code. Useful mainly for printing error messages.

## 2.10.2.5 bfd\_generic\_relax\_section

#### **Synopsis**

```
bfd_boolean bfd_generic_relax_section
  (bfd *abfd,
    asection *section,
    struct bfd_link_info *,
    bfd_boolean *);
```

### Description

Provides default handling for relaxing for back ends which don't do relaxing.

## 2.10.2.6 bfd\_generic\_gc\_sections

## **Synopsis**

```
bfd_boolean bfd_generic_gc_sections
   (bfd *, struct bfd_link_info *);
```

### Description

Provides default handling for relaxing for back ends which don't do section gc - i.e., does nothing.

## 2.10.2.7 bfd\_generic\_lookup\_section\_flags

### **Synopsis**

```
bfd_boolean bfd_generic_lookup_section_flags
  (struct bfd_link_info *, struct flag_info *, asection *);
```

#### Description

Provides default handling for section flags lookup – i.e., does nothing. Returns FALSE if the section should be omitted, otherwise TRUE.

## 2.10.2.8 bfd\_generic\_merge\_sections

## **Synopsis**

```
bfd_boolean bfd_generic_merge_sections
  (bfd *, struct bfd_link_info *);
```

#### Description

Provides default handling for SEC\_MERGE section merging for back ends which don't have SEC\_MERGE support – i.e., does nothing.

## 2.10.2.9 bfd\_generic\_get\_relocated\_section\_contents

#### **Synopsis**

```
bfd_byte *bfd_generic_get_relocated_section_contents
  (bfd *abfd,
    struct bfd_link_info *link_info,
    struct bfd_link_order *link_order,
    bfd_byte *data,
    bfd_boolean relocatable,
    asymbol **symbols);
```

#### Description

Provides default handling of relocation effort for back ends which can't be bothered to do it efficiently.

## 2.11 Core files

# 2.11.1 Core file functions

## Description

These are functions pertaining to core files.

## 2.11.1.1 bfd\_core\_file\_failing\_command

#### Synopsis

```
const char *bfd_core_file_failing_command (bfd *abfd);
```

## Description

Return a read-only string explaining which program was running when it failed and produced the core file *abfd*.

## 2.11.1.2 bfd\_core\_file\_failing\_signal

## **Synopsis**

```
int bfd_core_file_failing_signal (bfd *abfd);
```

#### Description

Returns the signal number which caused the core dump which generated the file the BFD abfd is attached to.

## 2.11.1.3 bfd\_core\_file\_pid

## **Synopsis**

```
int bfd_core_file_pid (bfd *abfd);
```

#### Description

Returns the PID of the process the core dump the BFD *abfd* is attached to was generated from.

## 2.11.1.4 core\_file\_matches\_executable\_p

## **Synopsis**

```
bfd_boolean core_file_matches_executable_p
   (bfd *core_bfd, bfd *exec_bfd);
```

#### Description

Return TRUE if the core file attached to  $core_{-}bfd$  was generated by a run of the executable file attached to  $exec_{-}bfd$ , FALSE otherwise.

## 2.11.1.5 generic\_core\_file\_matches\_executable\_p

#### **Synopsis**

```
bfd_boolean generic_core_file_matches_executable_p
  (bfd *core_bfd, bfd *exec_bfd);
```

## Description

Return TRUE if the core file attached to *core\_bfd* was generated by a run of the executable file attached to *exec\_bfd*. The match is based on executable basenames only.

Note: When not able to determine the core file failing command or the executable name, we still return TRUE even though we're not sure that core file and executable match. This is to avoid generating a false warning in situations where we really don't know whether they match or not.

# 2.12 Targets

#### Description

Each port of BFD to a different machine requires the creation of a target back end. All the back end provides to the root part of BFD is a structure containing pointers to functions which perform certain low level operations on files. BFD translates the applications's requests through a pointer into calls to the back end routines.

When a file is opened with bfd\_openr, its format and target are unknown. BFD uses various mechanisms to determine how to interpret the file. The operations performed are:

- Create a BFD by calling the internal routine \_bfd\_new\_bfd, then call bfd\_find\_target with the target string supplied to bfd\_openr and the new BFD pointer.
- If a null target string was provided to bfd\_find\_target, look up the environment variable GNUTARGET and use that as the target string.
- If the target string is still NULL, or the target string is default, then use the first item in the target vector as the target type, and set target\_defaulted in the BFD to cause bfd\_check\_format to loop through all the targets. See Section 2.12.1 [bfd\_target], page 131. See Section 2.9 [Formats], page 50.
- Otherwise, inspect the elements in the target vector one by one, until a match on target name is found. When found, use it.
- Otherwise return the error bfd\_error\_invalid\_target to bfd\_openr.
- bfd\_openr attempts to open the file using bfd\_open\_file, and returns the BFD.

Once the BFD has been opened and the target selected, the file format may be determined. This is done by calling bfd\_check\_format on the BFD with a suggested format. If target\_defaulted has been set, each possible target type is tried to see if it recognizes the specified format. bfd\_check\_format returns TRUE when the caller guesses right.

## 2.12.1 bfd\_target

## Description

This structure contains everything that BFD knows about a target. It includes things like its byte order, name, and which routines to call to do various operations.

Every BFD points to a target structure with its xvec member.

The macros below are used to dispatch to functions through the bfd\_target vector. They are used in a number of macros further down in 'bfd.h', and are also used when calling various routines by hand inside the BFD implementation. The arglist argument must be parenthesized; it contains all the arguments to the called function.

They make the documentation (more) unpleasant to read, so if someone wants to fix this and not break the above, please do.

```
#define BFD_SEND(bfd, message, arglist) \
    ((*((bfd)->xvec->message)) arglist)

#ifdef DEBUG_BFD_SEND
#undef BFD_SEND
#define BFD_SEND(bfd, message, arglist) \
    (((bfd) && (bfd)->xvec && (bfd)->xvec->message) ? \
        ((*((bfd)->xvec->message)) arglist) : \
        (bfd_assert (__FILE__,__LINE__), NULL))
#endif

For operations which index on the BFD format:
    #define BFD_SEND_FMT(bfd, message, arglist) \
        (((bfd)->xvec->message[(int) ((bfd)->format)]) arglist)

#ifdef DEBUG_BFD_SEND
#undef BFD_SEND_FMT
```

```
#define BFD_SEND_FMT(bfd, message, arglist) \
  (((bfd) && (bfd)->xvec && (bfd)->xvec->message) ? \
  (((bfd)->xvec->message[(int) ((bfd)->format)]) arglist) : \
  (bfd_assert (__FILE__,__LINE__), NULL))
#endif
```

This is the structure which defines the type of BFD this is. The xvec member of the struct bfd itself points here. Each module that implements access to a different target under BFD, defines one of these.

FIXME, these names should be rationalised with the names of the entry points which call them. Too bad we can't have one macro to define them both!

```
enum bfd_flavour
{
  bfd_target_unknown_flavour,
 bfd_target_aout_flavour,
 bfd_target_coff_flavour,
 bfd_target_ecoff_flavour,
  bfd_target_xcoff_flavour,
 bfd_target_elf_flavour,
  bfd_target_ieee_flavour,
 bfd_target_nlm_flavour,
  bfd_target_oasys_flavour,
 bfd_target_tekhex_flavour,
  bfd_target_srec_flavour,
 bfd_target_verilog_flavour,
  bfd_target_ihex_flavour,
 bfd_target_som_flavour,
 bfd_target_os9k_flavour,
  bfd_target_versados_flavour,
  bfd_target_msdos_flavour,
  bfd_target_ovax_flavour,
  bfd_target_evax_flavour,
  bfd_target_mmo_flavour,
 bfd_target_mach_o_flavour,
  bfd_target_pef_flavour,
 bfd_target_pef_xlib_flavour,
 bfd_target_sym_flavour
};
enum bfd_endian { BFD_ENDIAN_BIG, BFD_ENDIAN_LITTLE, BFD_ENDIAN_UNKNOWN };
/* Forward declaration. */
typedef struct bfd_link_info _bfd_link_info;
/* Forward declaration. */
typedef struct flag_info flag_info;
```

```
typedef struct bfd_target
 /* Identifies the kind of target, e.g., SunOS4, Ultrix, etc. */
 char *name:
 /* The "flavour" of a back end is a general indication about
   the contents of a file. */
 enum bfd_flavour flavour;
 /* The order of bytes within the data area of a file. */
 enum bfd_endian byteorder;
 /* The order of bytes within the header parts of a file. */
 enum bfd_endian header_byteorder;
 /* A mask of all the flags which an executable may have set -
    from the set BFD_NO_FLAGS, HAS_RELOC, ...D_PAGED. */
 flagword object_flags;
 /* A mask of all the flags which a section may have set - from
   the set SEC_NO_FLAGS, SEC_ALLOC, ...SET_NEVER_LOAD. */
 flagword section_flags;
 /* The character normally found at the front of a symbol.
    (if any), perhaps '_{-}'. */
 char symbol_leading_char;
 /* The pad character for file names within an archive header. */
 char ar_pad_char;
 /* The maximum number of characters in an archive header. */
 unsigned char ar_max_namelen;
 /* How well this target matches, used to select between various
    possible targets when more than one target matches. */
 unsigned char match_priority;
  /* Entries for byte swapping for data. These are different from the
    other entry points, since they don't take a BFD as the first argument.
    Certain other handlers could do the same. */
 bfd_uint64_t (*bfd_getx64) (const void *);
 bfd_int64_t (*bfd_getx_signed_64) (const void *);
                (*bfd_putx64) (bfd_uint64_t, void *);
 void
              (*bfd_getx32) (const void *);
 bfd_vma
 bfd_signed_vma (*bfd_getx_signed_32) (const void *);
                (*bfd_putx32) (bfd_vma, void *);
 void
```

```
bfd_vma
              (*bfd_getx16) (const void *);
bfd_signed_vma (*bfd_getx_signed_16) (const void *);
void
               (*bfd_putx16) (bfd_vma, void *);
/* Byte swapping for the headers. */
bfd_uint64_t (*bfd_h_getx64) (const void *);
bfd_int64_t
               (*bfd_h_getx_signed_64) (const void *);
void
              (*bfd_h_putx64) (bfd_uint64_t, void *);
bfd_vma
             (*bfd_h_getx32) (const void *);
bfd_signed_vma (*bfd_h_getx_signed_32) (const void *);
void
              (*bfd_h_putx32) (bfd_vma, void *);
bfd_vma
              (*bfd_h_getx16) (const void *);
bfd_signed_vma (*bfd_h_getx_signed_16) (const void *);
void
              (*bfd_h_putx16) (bfd_vma, void *);
/* Format dependent routines: these are vectors of entry points
   within the target vector structure, one for each format to check. */
/* Check the format of a file being read. Return a bfd_target * or zero. */
const struct bfd_target *(*_bfd_check_format[bfd_type_end]) (bfd *);
/* Set the format of a file being written. */
bfd_boolean (*_bfd_set_format[bfd_type_end]) (bfd *);
/* Write cached information into a file being written, at bfd_close. */■
bfd_boolean (*_bfd_write_contents[bfd_type_end]) (bfd *);
```

The general target vector. These vectors are initialized using the BFD\_JUMP\_TABLE macros.

```
/* Generic entry points. */
#define BFD_JUMP_TABLE_GENERIC(NAME) \
    NAME##_close_and_cleanup, \
    NAME##_bfd_free_cached_info, \
    NAME##_new_section_hook, \
    NAME##_get_section_contents, \
    NAME##_get_section_contents_in_window

/* Called when the BFD is being closed to do any necessary cleanup. */
bfd_boolean (*_close_and_cleanup) (bfd *);
/* Ask the BFD to free all cached information. */
bfd_boolean (*_bfd_free_cached_info) (bfd *);
/* Called when a new section is created. */
bfd_boolean (*_new_section_hook) (bfd *, sec_ptr);
/* Read the contents of a section. */
bfd_boolean (*_bfd_get_section_contents)
```

```
(bfd *, sec_ptr, void *, file_ptr, bfd_size_type);
 bfd_boolean (*_bfd_get_section_contents_in_window)
    (bfd *, sec_ptr, bfd_window *, file_ptr, bfd_size_type);
 /* Entry points to copy private data. */
#define BFD_JUMP_TABLE_COPY(NAME) \
 NAME##_bfd_copy_private_bfd_data, \
 NAME##_bfd_merge_private_bfd_data, \
  _bfd_generic_init_private_section_data, \
 NAME##_bfd_copy_private_section_data, \
 NAME##_bfd_copy_private_symbol_data, \
 NAME##_bfd_copy_private_header_data, \
 NAME##_bfd_set_private_flags, \
 NAME##_bfd_print_private_bfd_data
 /* Called to copy BFD general private data from one object file
    to another. */
 bfd_boolean (*_bfd_copy_private_bfd_data) (bfd *, bfd *);
 /* Called to merge BFD general private data from one object file
     to a common output file when linking. */
 bfd_boolean (*_bfd_merge_private_bfd_data) (bfd *, bfd *);
 /* Called to initialize BFD private section data from one object file
     to another. */
#define bfd_init_private_section_data(ibfd, isec, obfd, osec, link_info) \
 BFD_SEND (obfd, _bfd_init_private_section_data, (ibfd, isec, obfd, osec, link_info))
 bfd_boolean (*_bfd_init_private_section_data)
    (bfd *, sec_ptr, bfd *, sec_ptr, struct bfd_link_info *);
 /* Called to copy BFD private section data from one object file
    to another. */
 bfd_boolean (*_bfd_copy_private_section_data)
    (bfd *, sec_ptr, bfd *, sec_ptr);
  /* Called to copy BFD private symbol data from one symbol
     to another. */
 bfd_boolean (*_bfd_copy_private_symbol_data)
    (bfd *, asymbol *, bfd *, asymbol *);
  /* Called to copy BFD private header data from one object file
     to another. */
 bfd_boolean (*_bfd_copy_private_header_data)
    (bfd *, bfd *);
  /* Called to set private backend flags. */
 bfd_boolean (*_bfd_set_private_flags) (bfd *, flagword);
 /* Called to print private BFD data. */
 bfd_boolean (*_bfd_print_private_bfd_data) (bfd *, void *);
 /* Core file entry points. */
#define BFD_JUMP_TABLE_CORE(NAME) \
```

```
NAME##_core_file_failing_command, \
 NAME##_core_file_failing_signal, \
 NAME##_core_file_matches_executable_p, \
 NAME##_core_file_pid
 char *
              (*_core_file_failing_command) (bfd *);
 int
              (*_core_file_failing_signal) (bfd *);
 bfd_boolean (*_core_file_matches_executable_p) (bfd *, bfd *);
 int
             (*_core_file_pid) (bfd *);
 /* Archive entry points. */
#define BFD_JUMP_TABLE_ARCHIVE(NAME) \
 NAME##_slurp_armap, \
 NAME##_slurp_extended_name_table, \
 NAME##_construct_extended_name_table, \
 NAME##_truncate_arname, \
 NAME##_write_armap, \
 NAME##_read_ar_hdr, \
 NAME##_write_ar_hdr, \
 NAME##_openr_next_archived_file, \
 NAME##_get_elt_at_index, \
 NAME##_generic_stat_arch_elt, \
 NAME##_update_armap_timestamp
 bfd_boolean (*_bfd_slurp_armap) (bfd *);
 bfd_boolean (*_bfd_slurp_extended_name_table) (bfd *);
 bfd_boolean (*_bfd_construct_extended_name_table)
    (bfd *, char **, bfd_size_type *, const char **);
             (*_bfd_truncate_arname) (bfd *, const char *, char *);
 void
 bfd_boolean (*write_armap)
    (bfd *, unsigned int, struct orl *, unsigned int, int);
              (*_bfd_read_ar_hdr_fn) (bfd *);
 void *
 bfd_boolean (*_bfd_write_ar_hdr_fn) (bfd *, bfd *);
             (*openr_next_archived_file) (bfd *, bfd *);
#define bfd_get_elt_at_index(b,i) BFD_SEND (b, _bfd_get_elt_at_index, (b,i))
             (*_bfd_get_elt_at_index) (bfd *, symindex);
 bfd *
 int
              (*_bfd_stat_arch_elt) (bfd *, struct stat *);
 bfd_boolean (*_bfd_update_armap_timestamp) (bfd *);
  /* Entry points used for symbols. */
#define BFD_JUMP_TABLE_SYMBOLS(NAME) \
 NAME##_get_symtab_upper_bound, \
 NAME##_canonicalize_symtab, \
 NAME##_make_empty_symbol, \
 NAME##_print_symbol, \
 NAME##_get_symbol_info, \
 NAME##_bfd_is_local_label_name, \
```

```
NAME##_bfd_is_target_special_symbol, \
 NAME##_get_lineno, \
 NAME##_find_nearest_line, \
 NAME##_find_line, \
 NAME##_find_inliner_info, \
 NAME##_bfd_make_debug_symbol, \
 NAME##_read_minisymbols, \
 NAME##_minisymbol_to_symbol
             (*_bfd_get_symtab_upper_bound) (bfd *);
 long
 long
              (*_bfd_canonicalize_symtab)
    (bfd *, struct bfd_symbol **);
 struct bfd_symbol *
              (*_bfd_make_empty_symbol) (bfd *);
              (*_bfd_print_symbol)
 void
    (bfd *, void *, struct bfd_symbol *, bfd_print_symbol_type);
#define bfd_print_symbol(b,p,s,e) BFD_SEND (b, _bfd_print_symbol, (b,p,s,e))
             (*_bfd_get_symbol_info)
 void
    (bfd *, struct bfd_symbol *, symbol_info *);
#define bfd_get_symbol_info(b,p,e) BFD_SEND (b, _bfd_get_symbol_info, (b,p,e))
 bfd_boolean (*_bfd_is_local_label_name) (bfd *, const char *);
 bfd_boolean (*_bfd_is_target_special_symbol) (bfd *, asymbol *);
            (*_get_lineno) (bfd *, struct bfd_symbol *);
 bfd_boolean (*_bfd_find_nearest_line)
    (bfd *, struct bfd_symbol **, struct bfd_section *, bfd_vma,
     const char **, const char **, unsigned int *, unsigned int *);
 bfd_boolean (*_bfd_find_line)
    (bfd *, struct bfd_symbol **, struct bfd_symbol *,
     const char **, unsigned int *);
 bfd_boolean (*_bfd_find_inliner_info)
    (bfd *, const char **, const char **, unsigned int *);
 /* Back-door to allow format-aware applications to create debug symbols
   while using BFD for everything else. Currently used by the assembler
   when creating COFF files. */
            (*_bfd_make_debug_symbol)
 asymbol *
    (bfd *, void *, unsigned long size);
#define bfd_read_minisymbols(b, d, m, s) \
 BFD_SEND (b, _read_minisymbols, (b, d, m, s))
             (*_read_minisymbols)
    (bfd *, bfd_boolean, void **, unsigned int *);
#define bfd_minisymbol_to_symbol(b, d, m, f) \
 BFD_SEND (b, _minisymbol_to_symbol, (b, d, m, f))
 asymbol * (*_minisymbol_to_symbol)
    (bfd *, bfd_boolean, const void *, asymbol *);
 /* Routines for relocs. */
#define BFD_JUMP_TABLE_RELOCS(NAME) \
```

```
NAME##_get_reloc_upper_bound, \
 NAME##_canonicalize_reloc, \
 NAME##_bfd_reloc_type_lookup, \
 NAME##_bfd_reloc_name_lookup
              (*_get_reloc_upper_bound) (bfd *, sec_ptr);
 long
 long
              (*_bfd_canonicalize_reloc)
    (bfd *, sec_ptr, arelent **, struct bfd_symbol **);
 /* See documentation on reloc types. */
 reloc_howto_type *
              (*reloc_type_lookup) (bfd *, bfd_reloc_code_real_type);
 reloc_howto_type *
              (*reloc_name_lookup) (bfd *, const char *);
  /* Routines used when writing an object file. */
#define BFD_JUMP_TABLE_WRITE(NAME) \
 NAME##_set_arch_mach, \
 NAME##_set_section_contents
 bfd_boolean (*_bfd_set_arch_mach)
    (bfd *, enum bfd_architecture, unsigned long);
 bfd_boolean (*_bfd_set_section_contents)
    (bfd *, sec_ptr, const void *, file_ptr, bfd_size_type);
  /* Routines used by the linker. */
#define BFD_JUMP_TABLE_LINK(NAME) \
 NAME##_sizeof_headers, \
 NAME##_bfd_get_relocated_section_contents, \
 NAME##_bfd_relax_section, \
 NAME##_bfd_link_hash_table_create, \
 NAME##_bfd_link_add_symbols, \
 NAME##_bfd_link_just_syms, \
 NAME##_bfd_copy_link_hash_symbol_type, \
 NAME##_bfd_final_link, \
 NAME##_bfd_link_split_section, \
 NAME##_bfd_gc_sections, \
 NAME##_bfd_lookup_section_flags, \
 NAME##_bfd_merge_sections, \
 NAME##_bfd_is_group_section, \
 NAME##_bfd_discard_group, \
 NAME##_section_already_linked, \
 NAME##_bfd_define_common_symbol
              (*_bfd_sizeof_headers) (bfd *, struct bfd_link_info *);
 int
 bfd_byte * (*_bfd_get_relocated_section_contents)
    (bfd *, struct bfd_link_info *, struct bfd_link_order *,
```

```
bfd_byte *, bfd_boolean, struct bfd_symbol **);
 bfd_boolean (*_bfd_relax_section)
    (bfd *, struct bfd_section *, struct bfd_link_info *, bfd_boolean *);
 /* Create a hash table for the linker. Different backends store
    different information in this table. */
 struct bfd_link_hash_table *
              (*_bfd_link_hash_table_create) (bfd *);
 /* Add symbols from this object file into the hash table. */
 bfd_boolean (*_bfd_link_add_symbols) (bfd *, struct bfd_link_info *);
 /* Indicate that we are only retrieving symbol values from this section. */■
              (*_bfd_link_just_syms) (asection *, struct bfd_link_info *);
 void
 /* Copy the symbol type and other attributes for a linker script
    assignment of one symbol to another. */
#define bfd_copy_link_hash_symbol_type(b, t, f) \
 BFD_SEND (b, _bfd_copy_link_hash_symbol_type, (b, t, f))
 void (*_bfd_copy_link_hash_symbol_type)
    (bfd *, struct bfd_link_hash_entry *, struct bfd_link_hash_entry *);
 /* Do a link based on the link_order structures attached to each
    section of the BFD. */
 bfd_boolean (*_bfd_final_link) (bfd *, struct bfd_link_info *);
 /* Should this section be split up into smaller pieces during linking. */
 bfd_boolean (*_bfd_link_split_section) (bfd *, struct bfd_section *);
 /* Remove sections that are not referenced from the output. */
 bfd_boolean (*_bfd_gc_sections) (bfd *, struct bfd_link_info *);
  /* Sets the bitmask of allowed and disallowed section flags. */
 bfd_boolean (*_bfd_lookup_section_flags) (struct bfd_link_info *,
                                           struct flag_info *,
                                           asection *);
 /* Attempt to merge SEC_MERGE sections. */
 bfd_boolean (*_bfd_merge_sections) (bfd *, struct bfd_link_info *);
 /* Is this section a member of a group? */
 bfd_boolean (*_bfd_is_group_section) (bfd *, const struct bfd_section *);
 /* Discard members of a group. */
 bfd_boolean (*_bfd_discard_group) (bfd *, struct bfd_section *);
```

```
/* Check if SEC has been already linked during a reloceatable or
     final link. */
 bfd_boolean (*_section_already_linked) (bfd *, asection *,
                                          struct bfd_link_info *);
 /* Define a common symbol. */
 bfd_boolean (*_bfd_define_common_symbol) (bfd *, struct bfd_link_info *,
                                            struct bfd_link_hash_entry *);
  /* Routines to handle dynamic symbols and relocs. */
#define BFD_JUMP_TABLE_DYNAMIC(NAME) \
 NAME##_get_dynamic_symtab_upper_bound, \
 NAME##_canonicalize_dynamic_symtab, \
 NAME##_get_synthetic_symtab, \
 NAME##_get_dynamic_reloc_upper_bound, \
 NAME##_canonicalize_dynamic_reloc
 /* Get the amount of memory required to hold the dynamic symbols. */
             (*_bfd_get_dynamic_symtab_upper_bound) (bfd *);
  /* Read in the dynamic symbols. */
             (*_bfd_canonicalize_dynamic_symtab)
    (bfd *, struct bfd_symbol **);
 /* Create synthetized symbols. */
              (*_bfd_get_synthetic_symtab)
    (bfd *, long, struct bfd_symbol **, long, struct bfd_symbol **,
    struct bfd_symbol **);
 /* Get the amount of memory required to hold the dynamic relocs. */
              (*_bfd_get_dynamic_reloc_upper_bound) (bfd *);
 /* Read in the dynamic relocs. */
              (*_bfd_canonicalize_dynamic_reloc)
 long
    (bfd *, arelent **, struct bfd_symbol **);
```

A pointer to an alternative bfd\_target in case the current one is not satisfactory. This can happen when the target cpu supports both big and little endian code, and target chosen by the linker has the wrong endianness. The function open\_output() in ld/ldlang.c uses this field to find an alternative output format that is suitable.

```
/* Opposite endian version of this target. */
const struct bfd_target * alternative_target;

/* Data for use by back-end routines, which isn't
    generic enough to belong in this structure. */
const void *backend_data;
} bfd_target;
```

# 2.12.1.1 bfd\_set\_default\_target

# **Synopsis**

```
bfd_boolean bfd_set_default_target (const char *name);
```

### Description

Set the default target vector to use when recognizing a BFD. This takes the name of the target, which may be a BFD target name or a configuration triplet.

# 2.12.1.2 bfd\_find\_target

# Synopsis

```
const bfd_target *bfd_find_target (const char *target_name, bfd *abfd);
```

# Description

Return a pointer to the transfer vector for the object target named <code>target\_name</code>. If <code>target\_name</code> is <code>NULL</code>, choose the one in the environment variable <code>GNUTARGET</code>; if that is null or not defined, then choose the first entry in the target list. Passing in the string "default" or setting the environment variable to "default" will cause the first entry in the target list to be returned, and "target\_defaulted" will be set in the BFD if <code>abfd</code> isn't <code>NULL</code>. This causes <code>bfd\_check\_format</code> to loop over all the targets to find the one that matches the file being read.

# 2.12.1.3 bfd\_get\_target\_info

# **Synopsis**

```
const bfd_target *bfd_get_target_info (const char *target_name,
    bfd *abfd,
    bfd_boolean *is_bigendian,
    int *underscoring,
    const char **def_target_arch);
```

### Description

Return a pointer to the transfer vector for the object target named <code>target\_name</code>. If <code>target\_name</code> is <code>NULL</code>, choose the one in the environment variable <code>GNUTARGET</code>; if that is null or not defined, then choose the first entry in the target list. Passing in the string "default" or setting the environment variable to "default" will cause the first entry in the target list to be returned, and "target\_defaulted" will be set in the BFD if <code>abfd</code> isn't <code>NULL</code>. This causes <code>bfd\_check\_format</code> to loop over all the targets to find the one that matches the file being read. If <code>is\_bigendian</code> is not <code>NULL</code>, then set this value to target's endian mode. True for big-endian, <code>FALSE</code> for little-endian or for invalid target. If <code>underscoring</code> is not <code>NULL</code>, then set this value to target's underscoring mode. Zero for none-underscoring, -1 for invalid target, else the value of target vector's symbol underscoring. If <code>def\_target\_arch</code> is not <code>NULL</code>, then set it to the architecture string specified by the target\_name.

# 2.12.1.4 bfd\_target\_list

#### **Synopsis**

```
const char ** bfd_target_list (void);
```

# Description

Return a freshly malloced NULL-terminated vector of the names of all the valid BFD targets. Do not modify the names.

# 2.12.1.5 bfd\_seach\_for\_target

# **Synopsis**

```
const bfd_target *bfd_search_for_target
  (int (*search_func) (const bfd_target *, void *),
    void *);
```

### Description

Return a pointer to the first transfer vector in the list of transfer vectors maintained by BFD that produces a non-zero result when passed to the function search\_func. The parameter data is passed, unexamined, to the search function.

# 2.13 Architectures

BFD keeps one atom in a BFD describing the architecture of the data attached to the BFD: a pointer to a bfd\_arch\_info\_type.

Pointers to structures can be requested independently of a BFD so that an architecture's information can be interrogated without access to an open BFD.

The architecture information is provided by each architecture package. The set of default architectures is selected by the macro SELECT\_ARCHITECTURES. This is normally set up in the 'config/target.mt' file of your choice. If the name is not defined, then all the architectures supported are included.

When BFD starts up, all the architectures are called with an initialize method. It is up to the architecture back end to insert as many items into the list of architectures as it wants to; generally this would be one for each machine and one for the default case (an item with a machine field of 0).

BFD's idea of an architecture is implemented in 'archures.c'.

### 2.13.1 bfd architecture

#### Description

This enum gives the object file's CPU architecture, in a global sense—i.e., what processor family does it belong to? Another field indicates which processor within the family is in use. The machine gives a number which distinguishes different versions of the architecture, containing, for example, 2 and 3 for Intel i960 KA and i960 KB, and 68020 and 68030 for Motorola 68020 and 68030.

```
enum bfd_architecture
{
   bfd_arch_unknown,    /* File arch not known. */
   bfd_arch_obscure,    /* Arch known, not one of these. */
   bfd_arch_m68k,    /* Motorola 68xxx */
#define bfd_mach_m68000 1
#define bfd_mach_m68010 3
#define bfd_mach_m68010 3
#define bfd_mach_m68020 4
#define bfd_mach_m68030 5
#define bfd_mach_m68040 6
#define bfd_mach_m68060 7
```

```
#define bfd_mach_cpu32 8
#define bfd_mach_fido
#define bfd_mach_mcf_isa_a_nodiv 10
#define bfd_mach_mcf_isa_a 11
#define bfd_mach_mcf_isa_a_mac 12
#define bfd_mach_mcf_isa_a_emac 13
#define bfd_mach_mcf_isa_aplus 14
#define bfd_mach_mcf_isa_aplus_mac 15
#define bfd_mach_mcf_isa_aplus_emac 16
#define bfd_mach_mcf_isa_b_nousp 17
#define bfd_mach_mcf_isa_b_nousp_mac 18
#define bfd_mach_mcf_isa_b_nousp_emac 19
#define bfd_mach_mcf_isa_b 20
#define bfd_mach_mcf_isa_b_mac 21
#define bfd_mach_mcf_isa_b_emac 22
#define bfd_mach_mcf_isa_b_float 23
#define bfd_mach_mcf_isa_b_float_mac 24
#define bfd_mach_mcf_isa_b_float_emac 25
#define bfd_mach_mcf_isa_c 26
#define bfd_mach_mcf_isa_c_mac 27
#define bfd_mach_mcf_isa_c_emac 28
#define bfd_mach_mcf_isa_c_nodiv 29
#define bfd_mach_mcf_isa_c_nodiv_mac 30
#define bfd_mach_mcf_isa_c_nodiv_emac 31
                     /* DEC Vax */
  bfd_arch_vax,
                      /* Intel 960 */
  bfd_arch_i960,
    /* The order of the following is important.
       lower number indicates a machine type that
       only accepts a subset of the instructions
       available to machines with higher numbers.
       The exception is the "ca", which is
       incompatible with all other machines except
       "core". */
#define bfd_mach_i960_core
#define bfd_mach_i960_ka_sa
#define bfd_mach_i960_kb_sb
                                3
                                4
#define bfd_mach_i960_mc
                                5
#define bfd_mach_i960_xa
                                6
#define bfd_mach_i960_ca
                                7
#define bfd_mach_i960_jx
#define bfd_mach_i960_hx
                                8
                      /* OpenRISC 1000 */
  bfd_arch_or1k,
#define bfd_mach_or1k
                                1
#define bfd_mach_or1knd
                                2
```

```
bfd_arch_sparc,
                      /* SPARC */
#define bfd_mach_sparc
                                        1
/* The difference between v8plus and v9 is that v9 is a true 64 bit env.
#define bfd_mach_sparc_sparclet
                                        2
#define bfd_mach_sparc_sparclite
                                        3
#define bfd_mach_sparc_v8plus
                                        4
#define bfd_mach_sparc_v8plusa
                                        5 /* with ultrasparc add'ns. */
#define bfd_mach_sparc_sparclite_le
                                        6
#define bfd_mach_sparc_v9
                                       7
#define bfd_mach_sparc_v9a
                                        8 /* with ultrasparc add'ns.
#define bfd_mach_sparc_v8plusb
                                       9 /* with cheetah add'ns.
#define bfd_mach_sparc_v9b
                                        10 /* with cheetah add'ns. */
/* Nonzero if MACH has the v9 instruction set. */
#define bfd_mach_sparc_v9_p(mach) \
  ((mach) >= bfd_mach_sparc_v8plus && (mach) <= bfd_mach_sparc_v9b \</pre>
  && (mach) != bfd_mach_sparc_sparclite_le)
/* Nonzero if MACH is a 64 bit sparc architecture. */
#define bfd_mach_sparc_64bit_p(mach) \
  ((mach) >= bfd_mach_sparc_v9 && (mach) != bfd_mach_sparc_v8plusb)
  bfd_arch_spu,
                      /* PowerPC SPU */
#define bfd_mach_spu
                               256
  bfd_arch_mips,
                      /* MIPS Rxxxx */
#define bfd_mach_mips3000
                                        3000
#define bfd_mach_mips3900
                                        3900
#define bfd_mach_mips4000
                                        4000
#define bfd_mach_mips4010
                                        4010
#define bfd_mach_mips4100
                                        4100
#define bfd_mach_mips4111
                                        4111
#define bfd_mach_mips4120
                                        4120
#define bfd_mach_mips4300
                                        4300
#define bfd_mach_mips4400
                                        4400
#define bfd_mach_mips4600
                                        4600
#define bfd_mach_mips4650
                                        4650
#define bfd_mach_mips5000
                                        5000
#define bfd_mach_mips5400
                                        5400
#define bfd_mach_mips5500
                                        5500
#define bfd_mach_mips5900
                                        5900
#define bfd_mach_mips6000
                                        6000
#define bfd_mach_mips7000
                                        7000
#define bfd_mach_mips8000
                                        8000
#define bfd_mach_mips9000
                                        9000
#define bfd_mach_mips10000
                                        10000
#define bfd_mach_mips12000
                                        12000
#define bfd_mach_mips14000
                                        14000
#define bfd_mach_mips16000
                                        16000
#define bfd_mach_mips16
                                        16
#define bfd_mach_mips5
                                        5
```

```
#define bfd_mach_mips_loongson_2e
                                       3001
#define bfd_mach_mips_loongson_2f
                                       3002
#define bfd_mach_mips_loongson_3a
                                       3003
#define bfd_mach_mips_sb1
                                       12310201 /* octal 'SB', 01 */
#define bfd_mach_mips_octeon
                                       6501
#define bfd_mach_mips_octeonp
                                       6601
#define bfd_mach_mips_octeon2
                                       6502
#define bfd_mach_mips_xlr
                                       887682
                                                 /* decimal 'XLR' */
#define bfd_mach_mipsisa32
                                       32
#define bfd_mach_mipsisa32r2
                                       33
#define bfd_mach_mipsisa32r3
                                       34
#define bfd_mach_mipsisa32r5
                                       36
#define bfd_mach_mipsisa32r6
                                       37
#define bfd_mach_mipsisa64
                                       64
#define bfd_mach_mipsisa64r2
                                       65
#define bfd_mach_mipsisa64r3
                                       66
                                       68
#define bfd_mach_mipsisa64r5
#define bfd_mach_mipsisa64r6
                                       69
#define bfd_mach_mips_micromips
                                       96
  bfd_arch_i386,
                      /* Intel 386 */
#define bfd_mach_i386_intel_syntax
                                        (1 << 0)
#define bfd_mach_i386_i8086
                                        (1 << 1)
#define bfd_mach_i386_i386
                                        (1 << 2)
#define bfd_mach_x86_64
                                        (1 << 3)
#define bfd_mach_x64_32
                                        (1 << 4)
#define bfd_mach_i386_i386_intel_syntax (bfd_mach_i386_i386 | bfd_mach_i386_intel_synt
                                        (bfd_mach_x86_64 | bfd_mach_i386_intel_syntax)
#define bfd_mach_x86_64_intel_syntax
#define bfd_mach_x64_32_intel_syntax
                                        (bfd_mach_x64_32 | bfd_mach_i386_intel_syntax)
  bfd_arch_l1om,
                 /* Intel L10M */
                                        (1 << 5)
#define bfd_mach_l1om
#define bfd_mach_l1om_intel_syntax
                                        (bfd_mach_l1om | bfd_mach_i386_intel_syntax)
                   /* Intel K10M */
  bfd_arch_k1om,
#define bfd_mach_k1om
#define bfd_mach_k1om_intel_syntax
                                        (bfd_mach_k1om | bfd_mach_i386_intel_syntax)
#define bfd_mach_i386_nacl
                                        (1 << 7)
#define bfd_mach_i386_i386_nacl
                                        (bfd_mach_i386_i386 | bfd_mach_i386_nacl)
#define bfd_mach_x86_64_nacl
                                        (bfd_mach_x86_64 | bfd_mach_i386_nacl)
#define bfd_mach_x64_32_nacl
                                        (bfd_mach_x64_32 | bfd_mach_i386_nacl)
  bfd_arch_we32k,
                     /* AT&T WE32xxx */
                      /* CCI/Harris Tahoe */
  bfd_arch_tahoe,
                     /* Intel 860 */
  bfd_arch_i860,
                      /* IBM 360/370 Mainframes */
  bfd_arch_i370,
                      /* IBM ROMP PC/RT */
  bfd_arch_romp,
                      /* Convex */
  bfd_arch_convex,
                      /* Motorola 88xxx */
  bfd_arch_m88k,
  bfd_arch_m98k,
                      /* Motorola 98xxx */
                      /* Pyramid Technology */
  bfd_arch_pyramid,
```

```
bfd_arch_h8300,
                      /* Renesas H8/300 (formerly Hitachi H8/300) */
#define bfd_mach_h8300
#define bfd_mach_h8300h
                          2
#define bfd_mach_h8300s
                          3
#define bfd_mach_h8300hn
#define bfd_mach_h8300sn
#define bfd_mach_h8300sx 6
#define bfd_mach_h8300sxn 7
  bfd_arch_pdp11,
                      /* DEC PDP-11 */
  bfd_arch_plugin,
  bfd_arch_powerpc,
                      /* PowerPC */
#define bfd_mach_ppc
                                32
#define bfd_mach_ppc64
                                64
#define bfd_mach_ppc_403
                                403
#define bfd_mach_ppc_403gc
                                4030
#define bfd_mach_ppc_405
                                405
#define bfd_mach_ppc_505
                                505
#define bfd_mach_ppc_601
                                601
#define bfd_mach_ppc_602
                                602
#define bfd_mach_ppc_603
                                603
#define bfd_mach_ppc_ec603e
                                6031
#define bfd_mach_ppc_604
                                604
#define bfd_mach_ppc_620
                                620
#define bfd_mach_ppc_630
                                630
#define bfd_mach_ppc_750
                                750
#define bfd_mach_ppc_860
                                860
#define bfd_mach_ppc_a35
                                35
#define bfd_mach_ppc_rs64ii
                                642
#define bfd_mach_ppc_rs64iii
                                643
#define bfd_mach_ppc_7400
                                7400
#define bfd_mach_ppc_e500
                                500
#define bfd_mach_ppc_e500mc
                                5001
#define bfd_mach_ppc_e500mc64
                                5005
#define bfd_mach_ppc_e5500
                                5006
#define bfd_mach_ppc_e6500
                                5007
#define bfd_mach_ppc_titan
                                83
#define bfd_mach_ppc_vle
                                84
  bfd_arch_rs6000,
                      /* IBM RS/6000 */
#define bfd_mach_rs6k
                                6000
#define bfd_mach_rs6k_rs1
                                6001
#define bfd_mach_rs6k_rsc
                                6003
#define bfd_mach_rs6k_rs2
                                6002
  bfd_arch_hppa,
                      /* HP PA RISC */
#define bfd_mach_hppa10
                                10
#define bfd_mach_hppa11
                                11
#define bfd_mach_hppa20
                                20
#define bfd_mach_hppa20w
                                25
```

```
bfd_arch_d10v,
                      /* Mitsubishi D10V */
#define bfd_mach_d10v
                               1
#define bfd_mach_d10v_ts2
                               2
#define bfd_mach_d10v_ts3
                               3
  bfd_arch_d30v,
                    /* Mitsubishi D30V */
                      /* DLX */
  bfd_arch_dlx,
  bfd_arch_m68hc11,
                      /* Motorola 68HC11 */
  bfd_arch_m68hc12,
                      /* Motorola 68HC12 */
#define bfd_mach_m6812_default 0
#define bfd_mach_m6812
#define bfd_mach_m6812s
                               2
  bfd_arch_m9s12x,
                     /* Freescale S12X */
  bfd_arch_m9s12xg, /* Freescale XGATE */
                      /* Zilog Z8000 */
  bfd_arch_z8k,
#define bfd_mach_z8001
                               1
#define bfd_mach_z8002
                               2
  bfd_arch_h8500,
                      /* Renesas H8/500 (formerly Hitachi H8/500) */
                      /* Renesas / SuperH SH (formerly Hitachi SH) */
  bfd_arch_sh,
#define bfd_mach_sh
                               1
#define bfd_mach_sh2
                            0x20
#define bfd_mach_sh_dsp
                            0x2d
#define bfd_mach_sh2a
                            0x2a
#define bfd_mach_sh2a_nofpu 0x2b
#define bfd_mach_sh2a_nofpu_or_sh4_nommu_nofpu 0x2a1
#define bfd_mach_sh2a_nofpu_or_sh3_nommu 0x2a2
#define bfd_mach_sh2a_or_sh4 0x2a3
#define bfd_mach_sh2a_or_sh3e 0x2a4
#define bfd_mach_sh2e
                            0x2e
#define bfd_mach_sh3
                            0x30
#define bfd_mach_sh3_nommu
                            0x31
#define bfd_mach_sh3_dsp
                            0x3d
#define bfd_mach_sh3e
                            0x3e
#define bfd_mach_sh4
                            0x40
#define bfd_mach_sh4_nofpu 0x41
#define bfd_mach_sh4_nommu_nofpu
#define bfd_mach_sh4a
#define bfd_mach_sh4a_nofpu 0x4b
#define bfd_mach_sh4al_dsp
                            0x4d
#define bfd_mach_sh5
                            0x50
  bfd_arch_alpha,
                      /* Dec Alpha */
#define bfd_mach_alpha_ev4 0x10
#define bfd_mach_alpha_ev5 0x20
#define bfd_mach_alpha_ev6 0x30
                      /* Advanced Risc Machines ARM.
  bfd_arch_arm,
#define bfd_mach_arm_unknown
                               0
#define bfd_mach_arm_2
                               1
#define bfd_mach_arm_2a
                               2
```

```
#define bfd_mach_arm_3
                               3
                               4
#define bfd_mach_arm_3M
#define bfd_mach_arm_4
                               5
                               6
#define bfd_mach_arm_4T
#define bfd_mach_arm_5
                               7
#define bfd_mach_arm_5T
                               8
#define bfd_mach_arm_5TE
                               9
#define bfd_mach_arm_XScale
                               10
#define bfd_mach_arm_ep9312
                               11
#define bfd_mach_arm_iWMMXt
                               12
#define bfd_mach_arm_iWMMXt2
                               13
                      /* Andes NDS32 */
 bfd_arch_nds32,
#define bfd_mach_n1
                               1
                               2
#define bfd_mach_n1h
                               3
#define bfd_mach_n1h_v2
                               4
#define bfd_mach_n1h_v3
                               5
#define bfd_mach_n1h_v3m
 bfd_arch_ns32k,
                    /* National Semiconductors ns32000 */
                     /* WDC 65816 */
 bfd_arch_w65,
                    /* Texas Instruments TMS320C30 */
 bfd_arch_tic30,
 bfd_arch_tic4x,
                    /* Texas Instruments TMS320C3X/4X */
#define bfd_mach_tic3x
                               30
                               40
#define bfd_mach_tic4x
                      /* Texas Instruments TMS320C54X */
 bfd_arch_tic54x,
                     /* Texas Instruments TMS320C6X */
 bfd_arch_tic6x,
 bfd_arch_tic80,
                     /* TI TMS320c80 (MVP) */
                   /* NEC V850 */
 bfd_arch_v850,
 bfd_arch_v850_rh850,/* NEC V850 (using RH850 ABI) */
#define bfd_mach_v850
                               1
                               'E'
#define bfd_mach_v850e
#define bfd_mach_v850e1
                               11
#define bfd_mach_v850e2
                               0x4532
#define bfd_mach_v850e2v3
                               0x45325633
#define bfd_mach_v850e3v5
                               0x45335635 /* ('E'|'3'|'V'|'5') */
 bfd_arch_arc,
                      /* ARC Cores */
#define bfd_mach_arc_5
                               5
                               6
#define bfd_mach_arc_6
                               7
#define bfd_mach_arc_7
#define bfd_mach_arc_8
                               8
                    /* Renesas M16C/M32C. */
bfd_arch_m32c,
#define bfd_mach_m16c
                             0x75
#define bfd_mach_m32c
                             0x78
 bfd_arch_m32r,
                     /* Renesas M32R (formerly Mitsubishi M32R/D) */
#define bfd_mach_m32r
                               1 /* For backwards compatibility. */
                               , x ,
#define bfd_mach_m32rx
#define bfd_mach_m32r2
                               22
 bfd_arch_mn10200, /* Matsushita MN10200 */
```

```
bfd_arch_mn10300,
                      /* Matsushita MN10300 */
#define bfd_mach_mn10300
                                        300
#define bfd_mach_am33
                               330
#define bfd_mach_am33_2
                               332
  bfd_arch_fr30,
#define bfd_mach_fr30
                               0x46523330
  bfd_arch_frv,
#define bfd_mach_frv
                               1
#define bfd_mach_frvsimple
                               2
#define bfd_mach_fr300
                               300
#define bfd_mach_fr400
                               400
                               450
#define bfd_mach_fr450
#define bfd_mach_frvtomcat
                               499
                                        /* fr500 prototype */
                               500
#define bfd_mach_fr500
                               550
#define bfd_mach_fr550
                        /* The moxie processor */
  bfd_arch_moxie,
#define bfd_mach_moxie
                               1
  bfd_arch_mcore,
  bfd_arch_mep,
#define bfd_mach_mep
#define bfd_mach_mep_h1
                               0x6831
#define bfd_mach_mep_c5
                               0x6335
  bfd_arch_metag,
#define bfd_mach_metag
                               1
                      /* HP/Intel ia64 */
  bfd_arch_ia64,
#define bfd_mach_ia64_elf64
                               64
#define bfd_mach_ia64_elf32
                               32
  bfd_arch_ip2k,
                      /* Ubicom IP2K microcontrollers. */
#define bfd_mach_ip2022
                               1
#define bfd_mach_ip2022ext
bfd_arch_iq2000,
                      /* Vitesse IQ2000.
#define bfd_mach_iq2000
                               1
#define bfd_mach_iq10
                       /* Adapteva EPIPHANY */
  bfd_arch_epiphany,
#define bfd_mach_epiphany16
#define bfd_mach_epiphany32
  bfd_arch_mt,
#define bfd_mach_ms1
                               1
#define bfd_mach_mrisc2
                               2
                               3
#define bfd_mach_ms2
  bfd_arch_pj,
                      /* Atmel AVR microcontrollers.
  bfd_arch_avr,
#define bfd_mach_avr1
                               1
#define bfd_mach_avr2
                               25
#define bfd_mach_avr25
#define bfd_mach_avr3
                               3
                               31
#define bfd_mach_avr31
```

```
#define bfd_mach_avr35
                               35
#define bfd_mach_avr4
                               4
#define bfd_mach_avr5
                               5
#define bfd_mach_avr51
                               51
#define bfd_mach_avr6
                               6
#define bfd_mach_avrtiny
#define bfd_mach_avrxmega1 101
#define bfd_mach_avrxmega2 102
#define bfd_mach_avrxmega3 103
#define bfd_mach_avrxmega4 104
#define bfd_mach_avrxmega5 105
#define bfd_mach_avrxmega6 106
#define bfd_mach_avrxmega7 107
                       /* ADI Blackfin */
 bfd_arch_bfin,
#define bfd_mach_bfin
                               1
 bfd_arch_cr16,
                       /* National Semiconductor CompactRISC (ie CR16). */
#define bfd_mach_cr16
 bfd_arch_cr16c,
                        /* National Semiconductor CompactRISC. */
#define bfd_mach_cr16c
 bfd_arch_crx,
                      /* National Semiconductor CRX. */
#define bfd_mach_crx
                      /* Axis CRIS */
 bfd_arch_cris,
#define bfd_mach_cris_v0_v10
                               255
#define bfd_mach_cris_v32
                               32
#define bfd_mach_cris_v10_v32 1032
 bfd_arch_r178,
#define bfd_mach_rl78 0x75
                      /* Renesas RX.
 bfd_arch_rx,
#define bfd_mach_rx
                               0x75
                      /* IBM s390 */
 bfd_arch_s390,
#define bfd_mach_s390_31
                               31
#define bfd_mach_s390_64
 bfd_arch_score,
                      /* Sunplus score */
#define bfd_mach_score3
                                3
#define bfd_mach_score7
                                7
 bfd_arch_mmix,
                      /* Donald Knuth's educational processor. */
 bfd_arch_xstormy16,
#define bfd_mach_xstormy16
                               1
 bfd_arch_msp430,
                      /* Texas Instruments MSP430 architecture. */
#define bfd_mach_msp11
#define bfd_mach_msp110
                                110
#define bfd_mach_msp12
                                12
#define bfd_mach_msp13
                                13
#define bfd_mach_msp14
                                14
#define bfd_mach_msp15
                                15
#define bfd_mach_msp16
                                16
#define bfd_mach_msp20
                                20
```

```
#define bfd_mach_msp21
                                21
#define bfd_mach_msp22
                                22
#define bfd_mach_msp23
                                23
#define bfd_mach_msp24
                                24
#define bfd_mach_msp26
                                26
#define bfd_mach_msp31
                                31
#define bfd_mach_msp32
                                32
#define bfd_mach_msp33
                                33
#define bfd_mach_msp41
                                41
#define bfd_mach_msp42
                                42
#define bfd_mach_msp43
                                43
#define bfd_mach_msp44
                                44
#define bfd_mach_msp430x
                                45
#define bfd_mach_msp46
                                46
#define bfd_mach_msp47
                                47
#define bfd_mach_msp54
                                54
  bfd_arch_xc16x,
                      /* Infineon's XC16X Series.
                                                                 */
#define bfd_mach_xc16x
                               1
#define bfd_mach_xc16xl
                               2
#define bfd_mach_xc16xs
                               3
  bfd_arch_xgate, /* Freescale XGATE */
#define bfd_mach_xgate
  bfd_arch_xtensa,
                      /* Tensilica's Xtensa cores. */
#define bfd_mach_xtensa
  bfd_arch_z80,
#define bfd_mach_z80strict
                                1 /* No undocumented opcodes. */
#define bfd_mach_z80
                                3 /* With ixl, ixh, iyl, and iyh.
#define bfd_mach_z80full
                               7 /* All undocumented instructions. */
                                11 /* R800: successor with multiplication. */
#define bfd_mach_r800
  bfd_arch_lm32,
                      /* Lattice Mico32 */
#define bfd_mach_lm32
  bfd_arch_microblaze,/* Xilinx MicroBlaze. */
  bfd_arch_tilepro, /* Tilera TILEPro */
  bfd_arch_tilegx, /* Tilera TILE-Gx */
#define bfd_mach_tilepro
#define bfd_mach_tilegx
#define bfd_mach_tilegx32 2
                      /* AArch64 */
  bfd_arch_aarch64,
#define bfd_mach_aarch64 0
#define bfd_mach_aarch64_ilp32 32
  bfd_arch_nios2,
#define bfd_mach_nios2 0
  bfd_arch_last
  };
```

#### 2.13.2 bfd\_arch\_info

### Description

This structure contains information on architectures for use within BFD.

```
typedef struct bfd_arch_info
       int bits_per_word;
       int bits_per_address;
       int bits_per_byte;
       enum bfd_architecture arch;
       unsigned long mach;
       const char *arch_name;
       const char *printable_name;
       unsigned int section_align_power;
       /* TRUE if this is the default machine for the architecture.
          The default arch should be the first entry for an arch so that
          all the entries for that arch can be accessed via next. */
      bfd_boolean the_default;
       const struct bfd_arch_info * (*compatible)
         (const struct bfd_arch_info *a, const struct bfd_arch_info *b);
      bfd_boolean (*scan) (const struct bfd_arch_info *, const char *);
       /* Allocate via bfd_malloc and return a fill buffer of size COUNT. If
          IS_BIGENDIAN is TRUE, the order of bytes is big endian. If CODE is
          TRUE, the buffer contains code. */
       void *(*fill) (bfd_size_type count, bfd_boolean is_bigendian,
                      bfd_boolean code);
       const struct bfd_arch_info *next;
     }
     bfd_arch_info_type;
2.13.2.1 bfd_printable_name
```

Synopsis

```
const char *bfd_printable_name (bfd *abfd);
```

#### Description

Return a printable string representing the architecture and machine from the pointer to the architecture info structure.

### 2.13.2.2 bfd\_scan\_arch

#### **Synopsis**

```
const bfd_arch_info_type *bfd_scan_arch (const char *string);
```

### Description

Figure out if BFD supports any cpu which could be described with the name *string*. Return a pointer to an arch\_info structure if a machine is found, otherwise NULL.

### 2.13.2.3 bfd\_arch\_list

# **Synopsis**

```
const char **bfd_arch_list (void);
```

# Description

Return a freshly malloced NULL-terminated vector of the names of all the valid BFD architectures. Do not modify the names.

# 2.13.2.4 bfd\_arch\_get\_compatible

### **Synopsis**

```
const bfd_arch_info_type *bfd_arch_get_compatible
  (const bfd *abfd, const bfd *bbfd, bfd_boolean accept_unknowns);
```

### Description

Determine whether two BFDs' architectures and machine types are compatible. Calculates the lowest common denominator between the two architectures and machine types implied by the BFDs and returns a pointer to an arch\_info structure describing the compatible machine.

# 2.13.2.5 bfd\_default\_arch\_struct

#### Description

The bfd\_default\_arch\_struct is an item of bfd\_arch\_info\_type which has been initialized to a fairly generic state. A BFD starts life by pointing to this structure, until the correct back end has determined the real architecture of the file.

```
extern const bfd_arch_info_type bfd_default_arch_struct;
```

#### 2.13.2.6 bfd set arch info

#### **Synopsis**

```
void bfd_set_arch_info (bfd *abfd, const bfd_arch_info_type *arg);
```

# Description

Set the architecture info of abfd to arg.

# 2.13.2.7 bfd\_default\_set\_arch\_mach

# Synopsis

```
bfd_boolean bfd_default_set_arch_mach
   (bfd *abfd, enum bfd_architecture arch, unsigned long mach);
```

### Description

Set the architecture and machine type in BFD abfd to arch and mach. Find the correct pointer to a structure and insert it into the arch\_info pointer.

# 2.13.2.8 bfd\_get\_arch

# **Synopsis**

```
enum bfd_architecture bfd_get_arch (bfd *abfd);
```

# Description

Return the enumerated type which describes the BFD abfd's architecture.

### 2.13.2.9 bfd\_get\_mach

# **Synopsis**

```
unsigned long bfd_get_mach (bfd *abfd);
```

# Description

Return the long type which describes the BFD abfd's machine.

# 2.13.2.10 bfd\_arch\_bits\_per\_byte

### **Synopsis**

```
unsigned int bfd_arch_bits_per_byte (bfd *abfd);
```

### Description

Return the number of bits in one of the BFD abfd's architecture's bytes.

# 2.13.2.11 bfd\_arch\_bits\_per\_address

### **Synopsis**

```
unsigned int bfd_arch_bits_per_address (bfd *abfd);
```

#### Description

Return the number of bits in one of the BFD abfd's architecture's addresses.

### 2.13.2.12 bfd\_default\_compatible

### **Synopsis**

```
const bfd_arch_info_type *bfd_default_compatible
  (const bfd_arch_info_type *a, const bfd_arch_info_type *b);
```

#### Description

The default function for testing for compatibility.

### 2.13.2.13 bfd\_default\_scan

### **Synopsis**

```
bfd_boolean bfd_default_scan
  (const struct bfd_arch_info *info, const char *string);
```

#### Description

The default function for working out whether this is an architecture hit and a machine hit.

# 2.13.2.14 bfd\_get\_arch\_info

### **Synopsis**

```
const bfd_arch_info_type *bfd_get_arch_info (bfd *abfd);
```

#### Description

Return the architecture info struct in abfd.

# 2.13.2.15 bfd\_lookup\_arch

# **Synopsis**

```
const bfd_arch_info_type *bfd_lookup_arch
  (enum bfd_architecture arch, unsigned long machine);
```

### Description

Look for the architecture info structure which matches the arguments arch and machine. A machine of 0 matches the machine/architecture structure which marks itself as the default.

# 2.13.2.16 bfd\_printable\_arch\_mach

#### **Synopsis**

```
const char *bfd_printable_arch_mach
  (enum bfd_architecture arch, unsigned long machine);
```

#### Description

Return a printable string representing the architecture and machine type.

This routine is depreciated.

# 2.13.2.17 bfd\_octets\_per\_byte

# **Synopsis**

```
unsigned int bfd_octets_per_byte (bfd *abfd);
```

### Description

Return the number of octets (8-bit quantities) per target byte (minimum addressable unit). In most cases, this will be one, but some DSP targets have 16, 32, or even 48 bits per byte.

# 2.13.2.18 bfd\_arch\_mach\_octets\_per\_byte

### **Synopsis**

```
unsigned int bfd_arch_mach_octets_per_byte
  (enum bfd_architecture arch, unsigned long machine);
```

### Description

See bfd\_octets\_per\_byte.

This routine is provided for those cases where a bfd \* is not available

### 2.13.2.19 bfd\_arch\_default\_fill

#### Synopsis

```
void *bfd_arch_default_fill (bfd_size_type count,
    bfd_boolean is_bigendian,
    bfd_boolean code);
```

# Description

Allocate via bfd\_malloc and return a fill buffer of size COUNT. If IS\_BIGENDIAN is TRUE, the order of bytes is big endian. If CODE is TRUE, the buffer contains code.

```
/* Set to N to open the next N BFDs using an alternate id space. */
extern unsigned int bfd_use_reserved_id;
```

# 2.14 Opening and closing BFDs

# 2.14.1 Functions for opening and closing

# 2.14.1.1 bfd\_fopen

# **Synopsis**

```
bfd *bfd_fopen (const char *filename, const char *target,
      const char *mode, int fd);
```

# Description

Open the file *filename* with the target *target*. Return a pointer to the created BFD. If *fd* is not -1, then **fdopen** is used to open the file; otherwise, **fopen** is used. *mode* is passed directly to **fopen** or **fdopen**.

Calls bfd\_find\_target, so target is interpreted as by that function.

The new BFD is marked as cacheable iff fd is -1.

If NULL is returned then an error has occured. Possible errors are bfd\_error\_no\_memory, bfd\_error\_invalid\_target or system\_call error.

On error, fd is always closed.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.2 bfd\_openr

# **Synopsis**

```
bfd *bfd_openr (const char *filename, const char *target);
```

# Description

Open the file *filename* (using fopen) with the target target. Return a pointer to the created BFD.

Calls bfd\_find\_target, so target is interpreted as by that function.

If NULL is returned then an error has occured. Possible errors are bfd\_error\_no\_memory, bfd\_error\_invalid\_target or system\_call error.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

### 2.14.1.3 bfd\_fdopenr

# **Synopsis**

```
bfd *bfd_fdopenr (const char *filename, const char *target, int fd);
```

#### Description

bfd\_fdopenr is to bfd\_fopenr much like fdopen is to fopen. It opens a BFD on a file already described by the fd supplied.

When the file is later bfd\_closed, the file descriptor will be closed. If the caller desires that this file descriptor be cached by BFD (opened as needed, closed as needed to free descriptors for other opens), with the supplied fd used as an initial file descriptor (but subject to closure at any time), call bfd\_set\_cacheable(bfd, 1) on the returned BFD. The default is to assume

no caching; the file descriptor will remain open until bfd\_close, and will not be affected by BFD operations on other files.

Possible errors are bfd\_error\_no\_memory, bfd\_error\_invalid\_target and bfd\_error\_system\_call.

On error, fd is closed.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.4 bfd\_openstreamr

# **Synopsis**

```
bfd *bfd_openstreamr (const char * filename, const char * target, void * stream);
```

# Description

Open a BFD for read access on an existing stdio stream. When the BFD is passed to bfd\_close, the stream will be closed.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.5 bfd\_openr\_iovec

# **Synopsis**

```
bfd *bfd_openr_iovec (const char *filename, const char *target,
    void *(*open_func) (struct bfd *nbfd,
    void *open_closure),
    void *open_closure,
    file_ptr (*pread_func) (struct bfd *nbfd,
    void *stream,
    void *buf,
    file_ptr nbytes,
    file_ptr offset),
    int (*close_func) (struct bfd *nbfd,
    void *stream),
    int (*stat_func) (struct bfd *abfd,
    void *stream,
    struct stat *sb));
```

#### Description

Create and return a BFD backed by a read-only stream. The stream is created using open\_func, accessed using pread\_func and destroyed using close\_func.

Calls bfd\_find\_target, so target is interpreted as by that function.

Calls open\_func (which can call bfd\_zalloc and bfd\_get\_filename) to obtain the readonly stream backing the BFD. open\_func either succeeds returning the non-NULL stream, or fails returning NULL (setting bfd\_error).

Calls pread\_func to request nbytes of data from stream starting at offset (e.g., via a call to bfd\_read). pread\_func either succeeds returning the number of bytes read (which can be less than nbytes when end-of-file), or fails returning -1 (setting bfd\_error).

Calls *close\_func* when the BFD is later closed using bfd\_close. *close\_func* either succeeds returning 0, or fails returning -1 (setting bfd\_error).

Calls  $stat\_func$  to fill in a stat structure for bfd\_stat, bfd\_get\_size, and bfd\_get\_mtime calls.  $stat\_func$  returns 0 on success, or returns -1 on failure (setting bfd\_error).

If bfd\_openr\_iovec returns NULL then an error has occurred. Possible errors are bfd\_error\_no\_memory, bfd\_error\_invalid\_target and bfd\_error\_system\_call.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.6 bfd\_openw

### **Synopsis**

```
bfd *bfd_openw (const char *filename, const char *target);
```

### Description

Create a BFD, associated with file *filename*, using the file format *target*, and return a pointer to it.

Possible errors are bfd\_error\_system\_call, bfd\_error\_no\_memory, bfd\_error\_invalid\_target.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.7 bfd\_close

# **Synopsis**

```
bfd_boolean bfd_close (bfd *abfd);
```

#### Description

Close a BFD. If the BFD was open for writing, then pending operations are completed and the file written out and closed. If the created file is executable, then chmod is called to mark it as such.

All memory attached to the BFD is released.

The file descriptor associated with the BFD is closed (even if it was passed in to BFD by bfd\_fdopenr).

#### Returns

TRUE is returned if all is ok, otherwise FALSE.

# 2.14.1.8 bfd\_close\_all\_done

### **Synopsis**

```
bfd_boolean bfd_close_all_done (bfd *);
```

# Description

Close a BFD. Differs from bfd\_close since it does not complete any pending operations. This routine would be used if the application had just used BFD for swapping and didn't want to use any of the writing code.

If the created file is executable, then chmod is called to mark it as such.

All memory attached to the BFD is released.

#### Returns

TRUE is returned if all is ok, otherwise FALSE.

# 2.14.1.9 bfd\_create

# **Synopsis**

```
bfd *bfd_create (const char *filename, bfd *templ);
```

#### Description

Create a new BFD in the manner of bfd\_openw, but without opening a file. The new BFD takes the target from the target used by templ. The format is always set to bfd\_object.

A copy of the *filename* argument is stored in the newly created BFD. It can be accessed via the bfd\_get\_filename() macro.

# 2.14.1.10 bfd\_make\_writable

# **Synopsis**

```
bfd_boolean bfd_make_writable (bfd *abfd);
```

# Description

Takes a BFD as created by bfd\_create and converts it into one like as returned by bfd\_openw. It does this by converting the BFD to BFD\_IN\_MEMORY. It's assumed that you will call bfd\_make\_readable on this bfd later.

#### Returns

TRUE is returned if all is ok, otherwise FALSE.

# 2.14.1.11 bfd\_make\_readable

# **Synopsis**

```
bfd_boolean bfd_make_readable (bfd *abfd);
```

#### Description

Takes a BFD as created by bfd\_create and bfd\_make\_writable and converts it into one like as returned by bfd\_openr. It does this by writing the contents out to the memory buffer, then reversing the direction.

#### Returns

TRUE is returned if all is ok, otherwise FALSE.

# 2.14.1.12 bfd\_alloc

# Synopsis

```
void *bfd_alloc (bfd *abfd, bfd_size_type wanted);
```

#### Description

Allocate a block of wanted bytes of memory attached to abfd and return a pointer to it.

# 2.14.1.13 bfd\_alloc2

### Synopsis

```
void *bfd_alloc2 (bfd *abfd, bfd_size_type nmemb, bfd_size_type size);
```

#### Description

Allocate a block of *nmemb* elements of size bytes each of memory attached to abfd and return a pointer to it.

# 2.14.1.14 bfd\_zalloc

# **Synopsis**

```
void *bfd_zalloc (bfd *abfd, bfd_size_type wanted);
```

#### Description

Allocate a block of wanted bytes of zeroed memory attached to abfd and return a pointer to it.

# 2.14.1.15 bfd\_zalloc2

# **Synopsis**

```
void *bfd_zalloc2 (bfd *abfd, bfd_size_type nmemb, bfd_size_type size);
```

### Description

Allocate a block of *nmemb* elements of size bytes each of zeroed memory attached to abfd and return a pointer to it.

# 2.14.1.16 bfd\_calc\_gnu\_debuglink\_crc32

# **Synopsis**

```
unsigned long bfd_calc_gnu_debuglink_crc32
  (unsigned long crc, const unsigned char *buf, bfd_size_type len);
```

#### Description

Computes a CRC value as used in the .gnu\_debuglink section. Advances the previously computed *crc* value by computing and adding in the crc32 for *len* bytes of *buf* .

#### Returns

Return the updated CRC32 value.

# 2.14.1.17 bfd\_get\_debug\_link\_info

# Synopsis

```
char *bfd_get_debug_link_info (bfd *abfd, unsigned long *crc32_out);
```

# Description

Fetch the filename and CRC32 value for any separate debuginfo associated with abfd. Return NULL if no such info found, otherwise return filename and update crc32\_out. The returned filename is allocated with malloc; freeing it is the responsibility of the caller.

# 2.14.1.18 bfd\_get\_alt\_debug\_link\_info

# **Synopsis**

```
char *bfd_get_alt_debug_link_info (bfd * abfd,
    bfd_size_type *buildid_len,
    bfd_byte **buildid_out);
```

#### Description

Fetch the filename and BuildID value for any alternate debuginfo associated with abfd. Return NULL if no such info found, otherwise return filename and update buildid\_len and buildid\_out. The returned filename and build\_id are allocated with malloc; freeing them is the responsibility of the caller.

# 2.14.1.19 separate\_debug\_file\_exists

# **Synopsis**

```
bfd_boolean separate_debug_file_exists
  (char *name, unsigned long crc32);
```

# Description

Checks to see if name is a file and if its contents match crc32.

# 2.14.1.20 separate\_alt\_debug\_file\_exists

### **Synopsis**

```
bfd_boolean separate_alt_debug_file_exists
  (char *name, unsigned long crc32);
```

### Description

Checks to see if name is a file and if its BuildID matches buildid.

# 2.14.1.21 find\_separate\_debug\_file

### **Synopsis**

```
char *find_separate_debug_file (bfd *abfd);
```

### Description

Searches abfd for a section called section\_name which is expected to contain a reference to a file containing separate debugging information. The function scans various locations in the filesystem, including the file tree rooted at debug\_file\_directory, and returns the first matching filename that it finds. If check\_crc is TRUE then the contents of the file must also match the CRC value contained in section\_name. Returns NULL if no valid file could be found.

# 2.14.1.22 bfd\_follow\_gnu\_debuglink

# **Synopsis**

```
char *bfd_follow_gnu_debuglink (bfd *abfd, const char *dir);
```

### Description

Takes a BFD and searches it for a .gnu\_debuglink section. If this section is found, it examines the section for the name and checksum of a '.debug' file containing auxiliary debugging information. It then searches the filesystem for this .debug file in some standard locations, including the directory tree rooted at dir, and if found returns the full filename. If dir is NULL, it will search a default path configured into libbfd at build time. [XXX this

feature is not currently implemented].

#### Returns

NULL on any errors or failure to locate the .debug file, otherwise a pointer to a heap-allocated string containing the filename. The caller is responsible for freeing this string.

# 2.14.1.23 bfd\_follow\_gnu\_debugaltlink

### **Synopsis**

```
char *bfd_follow_gnu_debugaltlink (bfd *abfd, const char *dir);
```

#### Description

Takes a BFD and searches it for a .gnu\_debugaltlink section. If this section is found, it

examines the section for the name of a file containing auxiliary debugging information. It then searches the filesystem for this file in a set of standard locations, including the directory tree rooted at *dir*, and if found returns the full filename.

If dir is NULL, it will search a default path configured into libbfd at build time. [FIXME: This feature is not currently implemented].

#### Returns

NULL on any errors or failure to locate the debug file, otherwise a pointer to a heap-allocated string containing the filename. The caller is responsible for freeing this string.

# 2.14.1.24 bfd\_create\_gnu\_debuglink\_section

### **Synopsis**

```
struct bfd_section *bfd_create_gnu_debuglink_section
  (bfd *abfd, const char *filename);
```

#### Description

Takes a BFD and adds a .gnu\_debuglink section to it. The section is sized to be big enough to contain a link to the specified filename.

#### Returns

A pointer to the new section is returned if all is ok. Otherwise NULL is returned and bfd\_error is set.

# 2.14.1.25 bfd\_fill\_in\_gnu\_debuglink\_section

### **Synopsis**

```
bfd_boolean bfd_fill_in_gnu_debuglink_section
  (bfd *abfd, struct bfd_section *sect, const char *filename);
```

#### Description

Takes a BFD and containing a .gnu\_debuglink section SECT and fills in the contents of the section to contain a link to the specified filename. The filename should be relative to the current directory.

#### Returns

TRUE is returned if all is ok. Otherwise FALSE is returned and bfd\_error is set.

# 2.15 Implementation details

# 2.15.1 Internal functions

### Description

These routines are used within BFD. They are not intended for export, but are documented here for completeness.

# 2.15.1.1 bfd\_write\_bigendian\_4byte\_int

#### **Synopsis**

```
bfd_boolean bfd_write_bigendian_4byte_int (bfd *, unsigned int);
```

#### Description

Write a 4 byte integer i to the output BFD abfd, in big endian order regardless of what else is going on. This is useful in archives.

# 2.15.1.2 bfd\_put\_size

# 2.15.1.3 bfd\_get\_size

# Description

These macros as used for reading and writing raw data in sections; each access (except for bytes) is vectored through the target format of the BFD and mangled accordingly. The mangling performs any necessary endian translations and removes alignment restrictions. Note that types accepted and returned by these macros are identical so they can be swapped around in macros—for example, 'libaout.h' defines GET\_WORD to either bfd\_get\_32 or bfd\_get\_64.

In the put routines, val must be a bfd\_vma. If we are on a system without prototypes, the caller is responsible for making sure that is true, with a cast if necessary. We don't cast them in the macro definitions because that would prevent lint or gcc -Wall from detecting sins such as passing a pointer. To detect calling these with less than a bfd\_vma, use gcc -Wconversion on a host with 64 bit bfd\_vma's.

```
/* Byte swapping macros for user section data. */
#define bfd_put_8(abfd, val, ptr) \
  ((void) (*((unsigned char *) (ptr)) = (val) & 0xff))
#define bfd_put_signed_8 \
  bfd_put_8
#define bfd_get_8(abfd, ptr) \
  (*(const unsigned char *) (ptr) & 0xff)
#define bfd_get_signed_8(abfd, ptr) \
  (((*(const unsigned char *) (ptr) & 0xff) ^ 0x80) - 0x80)
#define bfd_put_16(abfd, val, ptr) \
  BFD_SEND (abfd, bfd_putx16, ((val),(ptr)))
#define bfd_put_signed_16 \
  bfd_put_16
#define bfd_get_16(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx16, (ptr))
#define bfd_get_signed_16(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx_signed_16, (ptr))
#define bfd_put_32(abfd, val, ptr) \
  BFD_SEND (abfd, bfd_putx32, ((val),(ptr)))
#define bfd_put_signed_32 \
  bfd_put_32
#define bfd_get_32(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx32, (ptr))
#define bfd_get_signed_32(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx_signed_32, (ptr))
#define bfd_put_64(abfd, val, ptr) \
```

```
BFD_SEND (abfd, bfd_putx64, ((val), (ptr)))
#define bfd_put_signed_64 \
  bfd_put_64
#define bfd_get_64(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx64, (ptr))
#define bfd_get_signed_64(abfd, ptr) \
  BFD_SEND (abfd, bfd_getx_signed_64, (ptr))
#define bfd_get(bits, abfd, ptr)
  ((bits) == 8 ? (bfd_vma) bfd_get_8 (abfd, ptr)
   : (bits) == 16 ? bfd_get_16 (abfd, ptr)
   : (bits) == 32 ? bfd_get_32 (abfd, ptr)
   : (bits) == 64 ? bfd_get_64 (abfd, ptr)
   : (abort (), (bfd_vma) - 1))
#define bfd_put(bits, abfd, val, ptr)
  ((bits) == 8 ? bfd_put_8 (abfd, val, ptr)
   : (bits) == 16 ? bfd_put_16 (abfd, val, ptr)
   : (bits) == 32 ? bfd_put_32 (abfd, val, ptr)
   : (bits) == 64 ? bfd_put_64 (abfd, val, ptr)
   : (abort (), (void) 0))
```

# 2.15.1.4 bfd\_h\_put\_size

### Description

These macros have the same function as their bfd\_get\_x brethren, except that they are used for removing information for the header records of object files. Believe it or not, some object files keep their header records in big endian order and their data in little endian order.

```
/* Byte swapping macros for file header data. */
#define bfd_h_put_8(abfd, val, ptr) \
   bfd_put_8 (abfd, val, ptr)
#define bfd_h_put_signed_8(abfd, val, ptr) \
   bfd_put_8 (abfd, val, ptr)
#define bfd_h_get_8(abfd, ptr) \
   bfd_get_8 (abfd, ptr)
#define bfd_h_get_signed_8(abfd, ptr) \
   bfd_get_signed_8 (abfd, ptr)

#define bfd_h_put_16(abfd, val, ptr) \
   BFD_SEND (abfd, bfd_h_putx16, (val, ptr))
#define bfd_h_put_signed_16 \
   bfd_h_put_16
#define bfd_h_get_16(abfd, ptr) \
```

```
BFD_SEND (abfd, bfd_h_getx16, (ptr))
#define bfd_h_get_signed_16(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_16, (ptr))
#define bfd_h_put_32(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_h_putx32, (val, ptr))
#define bfd_h_put_signed_32 \
 bfd_h_put_32
#define bfd_h_get_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx32, (ptr))
#define bfd_h_get_signed_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_32, (ptr))
#define bfd_h_put_64(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_h_putx64, (val, ptr))
#define bfd_h_put_signed_64 \
 bfd_h_put_64
#define bfd_h_get_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx64, (ptr))
#define bfd_h_get_signed_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_64, (ptr))
/* Aliases for the above, which should eventually go away. */
#define H_PUT_64 bfd_h_put_64
#define H_PUT_32 bfd_h_put_32
#define H_PUT_16 bfd_h_put_16
#define H_PUT_8
                 bfd_h_put_8
#define H_PUT_S64 bfd_h_put_signed_64
#define H_PUT_S32 bfd_h_put_signed_32
#define H_PUT_S16 bfd_h_put_signed_16
#define H_PUT_S8 bfd_h_put_signed_8
#define H_GET_64 bfd_h_get_64
#define H_GET_32 bfd_h_get_32
#define H_GET_16 bfd_h_get_16
#define H_GET_8
                bfd_h_get_8
#define H_GET_S64 bfd_h_get_signed_64
#define H_GET_S32 bfd_h_get_signed_32
#define H_GET_S16 bfd_h_get_signed_16
#define H_GET_S8 bfd_h_get_signed_8
```

# 2.15.1.5 bfd\_log2

### Synopsis

```
unsigned int bfd_log2 (bfd_vma x);
```

#### Description

Return the log base 2 of the value supplied, rounded up. E.g., an x of 1025 returns 11. A x of 0 returns 0.

# 2.16 File caching

The file caching mechanism is embedded within BFD and allows the application to open as many BFDs as it wants without regard to the underlying operating system's file descriptor limit (often as low as 20 open files). The module in cache.c maintains a least recently used list of bfd\_cache\_max\_open files, and exports the name bfd\_cache\_lookup, which runs around and makes sure that the required BFD is open. If not, then it chooses a file to close, closes it and opens the one wanted, returning its file handle.

# 2.16.1 Caching functions

# 2.16.1.1 bfd\_cache\_init

### Synopsis

```
bfd_boolean bfd_cache_init (bfd *abfd);
```

#### Description

Add a newly opened BFD to the cache.

### 2.16.1.2 bfd\_cache\_close

# **Synopsis**

```
bfd_boolean bfd_cache_close (bfd *abfd);
```

#### Description

Remove the BFD abfd from the cache. If the attached file is open, then close it too.

#### Returns

FALSE is returned if closing the file fails, TRUE is returned if all is well.

### 2.16.1.3 bfd\_cache\_close\_all

#### Synopsis

```
bfd_boolean bfd_cache_close_all (void);
```

#### Description

Remove all BFDs from the cache. If the attached file is open, then close it too.

#### Returns

FALSE is returned if closing one of the file fails, TRUE is returned if all is well.

### 2.16.1.4 bfd\_open\_file

#### **Synopsis**

```
FILE* bfd_open_file (bfd *abfd);
```

#### Description

Call the OS to open a file for *abfd*. Return the FILE \* (possibly NULL) that results from this operation. Set up the BFD so that future accesses know the file is open. If the FILE \* returned is NULL, then it won't have been put in the cache, so it won't have to be removed from it.

# 2.17 Linker Functions

The linker uses three special entry points in the BFD target vector. It is not necessary to write special routines for these entry points when creating a new BFD back end, since generic versions are provided. However, writing them can speed up linking and make it use significantly less runtime memory.

The first routine creates a hash table used by the other routines. The second routine adds the symbols from an object file to the hash table. The third routine takes all the object files and links them together to create the output file. These routines are designed so that the linker proper does not need to know anything about the symbols in the object files that it is linking. The linker merely arranges the sections as directed by the linker script and lets BFD handle the details of symbols and relocs.

The second routine and third routines are passed a pointer to a struct bfd\_link\_info structure (defined in bfdlink.h) which holds information relevant to the link, including the linker hash table (which was created by the first routine) and a set of callback functions to the linker proper.

The generic linker routines are in linker.c, and use the header file genlink.h. As of this writing, the only back ends which have implemented versions of these routines are a.out (in aoutx.h) and ECOFF (in ecoff.c). The a.out routines are used as examples throughout this section.

# 2.17.1 Creating a linker hash table

The linker routines must create a hash table, which must be derived from struct bfd\_link\_hash\_table described in bfdlink.c. See Section 2.18 [Hash Tables], page 172, for information on how to create a derived hash table. This entry point is called using the target vector of the linker output file.

The \_bfd\_link\_hash\_table\_create entry point must allocate and initialize an instance of the desired hash table. If the back end does not require any additional information to be stored with the entries in the hash table, the entry point may simply create a struct bfd\_link\_hash\_table. Most likely, however, some additional information will be needed.

For example, with each entry in the hash table the a.out linker keeps the index the symbol has in the final output file (this index number is used so that when doing a relocatable link the symbol index used in the output file can be quickly filled in when copying over a reloc). The a.out linker code defines the required structures and functions for a hash table derived from struct bfd\_link\_hash\_table. The a.out linker hash table is created by the function NAME(aout,link\_hash\_table\_create); it simply allocates space for the hash table, initializes it, and returns a pointer to it.

When writing the linker routines for a new back end, you will generally not know exactly which fields will be required until you have finished. You should simply create a new hash table which defines no additional fields, and then simply add fields as they become necessary.

# 2.17.2 Adding symbols to the hash table

The linker proper will call the \_bfd\_link\_add\_symbols entry point for each object file or archive which is to be linked (typically these are the files named on the command line, but some may also come from the linker script). The entry point is responsible for examining the file. For an object file, BFD must add any relevant symbol information to the hash

table. For an archive, BFD must determine which elements of the archive should be used and adding them to the link.

The alout version of this entry point is NAME(aout,link\_add\_symbols).

# 2.17.2.1 Differing file formats

Normally all the files involved in a link will be of the same format, but it is also possible to link together different format object files, and the back end must support that. The <code>\_bfd\_link\_add\_symbols</code> entry point is called via the target vector of the file to be added. This has an important consequence: the function may not assume that the hash table is the type created by the corresponding <code>\_bfd\_link\_hash\_table\_create</code> vector. All the <code>\_bfd\_link\_add\_symbols</code> function can assume about the hash table is that it is derived from <code>struct bfd\_link\_hash\_table</code>.

Sometimes the \_bfd\_link\_add\_symbols function must store some information in the hash table entry to be used by the \_bfd\_final\_link function. In such a case the output bfd xvec must be checked to make sure that the hash table was created by an object file of the same format.

The \_bfd\_final\_link routine must be prepared to handle a hash entry without any extra information added by the \_bfd\_link\_add\_symbols function. A hash entry without extra information will also occur when the linker script directs the linker to create a symbol. Note that, regardless of how a hash table entry is added, all the fields will be initialized to some sort of null value by the hash table entry initialization function.

See ecoff\_link\_add\_externals for an example of how to check the output bfd before saving information (in this case, the ECOFF external symbol debugging information) in a hash table entry.

# 2.17.2.2 Adding symbols from an object file

When the \_bfd\_link\_add\_symbols routine is passed an object file, it must add all externally visible symbols in that object file to the hash table. The actual work of adding the symbol to the hash table is normally handled by the function \_bfd\_generic\_link\_add\_one\_symbol. The \_bfd\_link\_add\_symbols routine is responsible for reading all the symbols from the object file and passing the correct information to \_bfd\_generic\_link\_add\_one\_symbol.

The \_bfd\_link\_add\_symbols routine should not use bfd\_canonicalize\_symtab to read the symbols. The point of providing this routine is to avoid the overhead of converting the symbols into generic asymbol structures.

\_bfd\_generic\_link\_add\_one\_symbol handles the details of combining common symbols, warning about multiple definitions, and so forth. It takes arguments which describe the symbol to add, notably symbol flags, a section, and an offset. The symbol flags include such things as BSF\_WEAK or BSF\_INDIRECT. The section is a section in the object file, or something like bfd\_und\_section\_ptr for an undefined symbol or bfd\_com\_section\_ptr for a common symbol.

If the \_bfd\_final\_link routine is also going to need to read the symbol information, the \_bfd\_link\_add\_symbols routine should save it somewhere attached to the object file BFD. However, the information should only be saved if the keep\_memory field of the info argument is TRUE, so that the -no-keep-memory linker switch is effective.

The alout function which adds symbols from an object file is aout\_link\_add\_object\_symbols, and most of the interesting work is in aout\_link\_add\_symbols. The latter saves pointers to the hash tables entries created by \_bfd\_generic\_link\_add\_one\_symbol indexed by symbol number, so that the \_bfd\_final\_link routine does not have to call the hash table lookup routine to locate the entry.

# 2.17.2.3 Adding symbols from an archive

When the \_bfd\_link\_add\_symbols routine is passed an archive, it must look through the symbols defined by the archive and decide which elements of the archive should be included in the link. For each such element it must call the add\_archive\_element linker callback, and it must add the symbols from the object file to the linker hash table. (The callback may in fact indicate that a replacement BFD should be used, in which case the symbols from that BFD should be added to the linker hash table instead.)

In most cases the work of looking through the symbols in the archive should be done by the \_bfd\_generic\_link\_add\_archive\_symbols function. \_bfd\_generic\_link\_add\_archive\_symbols is passed a function to call to make the final decision about adding an archive element to the link and to do the actual work of adding the symbols to the linker hash table. If the element is to be included, the add\_archive\_element linker callback routine must be called with the element as an argument, and the element's symbols must be added to the linker hash table just as though the element had itself been passed to the \_bfd\_link\_add\_symbols function.

When the a.out \_bfd\_link\_add\_symbols function receives an archive, it calls \_bfd\_generic\_link\_add\_archive\_symbols passing aout\_link\_check\_archive\_element as the function argument. aout\_link\_check\_archive\_element calls aout\_link\_check\_ar\_symbols. If the latter decides to add the element (an element is only added if it provides a real, non-common, definition for a previously undefined or common symbol) it calls the add\_archive\_element callback and then aout\_link\_check\_archive\_element calls aout\_link\_add\_symbols to actually add the symbols to the linker hash table possibly those of a substitute BFD, if the add\_archive\_element callback avails itself of that option.

The ECOFF back end is unusual in that it does not normally call \_bfd\_generic\_link\_add\_archive\_symbols, because ECOFF archives already contain a hash table of symbols. The ECOFF back end searches the archive itself to avoid the overhead of creating a new hash table.

# 2.17.3 Performing the final link

When all the input files have been processed, the linker calls the <code>\_bfd\_final\_link</code> entry point of the output BFD. This routine is responsible for producing the final output file, which has several aspects. It must relocate the contents of the input sections and copy the data into the output sections. It must build an output symbol table including any local symbols from the input files and the global symbols from the hash table. When producing relocatable output, it must modify the input relocs and write them into the output file. There may also be object format dependent work to be done.

The linker will also call the write\_object\_contents entry point when the BFD is closed. The two entry points must work together in order to produce the correct output file.

The details of how this works are inevitably dependent upon the specific object file format. The a.out \_bfd\_final\_link routine is NAME(aout,final\_link).

# 2.17.3.1 Information provided by the linker

Before the linker calls the \_bfd\_final\_link entry point, it sets up some data structures for the function to use.

The input\_bfds field of the bfd\_link\_info structure will point to a list of all the input files included in the link. These files are linked through the link.next field of the bfd structure.

Each section in the output file will have a list of link\_order structures attached to the map\_head.link\_order field (the link\_order structure is defined in bfdlink.h). These structures describe how to create the contents of the output section in terms of the contents of various input sections, fill constants, and, eventually, other types of information. They also describe relocs that must be created by the BFD backend, but do not correspond to any input file; this is used to support -Ur, which builds constructors while generating a relocatable object file.

# 2.17.3.2 Relocating the section contents

The \_bfd\_final\_link function should look through the link\_order structures attached to each section of the output file. Each link\_order structure should either be handled specially, or it should be passed to the function \_bfd\_default\_link\_order which will do the right thing (\_bfd\_default\_link\_order is defined in linker.c).

For efficiency, a link\_order of type bfd\_indirect\_link\_order whose associated section belongs to a BFD of the same format as the output BFD must be handled specially. This type of link\_order describes part of an output section in terms of a section belonging to one of the input files. The \_bfd\_final\_link function should read the contents of the section and any associated relocs, apply the relocs to the section contents, and write out the modified section contents. If performing a relocatable link, the relocs themselves must also be modified and written out.

The functions \_bfd\_relocate\_contents and \_bfd\_final\_link\_relocate provide some general support for performing the actual relocations, notably overflow checking. Their arguments include information about the symbol the relocation is against and a reloc\_howto\_type argument which describes the relocation to perform. These functions are defined in reloc.c.

The alout function which handles reading, relocating, and writing section contents is aout\_link\_input\_section. The actual relocation is done in aout\_link\_input\_section\_std and aout\_link\_input\_section\_ext.

# 2.17.3.3 Writing the symbol table

The \_bfd\_final\_link function must gather all the symbols in the input files and write them out. It must also write out all the symbols in the global hash table. This must be controlled by the strip and discard fields of the bfd\_link\_info structure.

The local symbols of the input files will not have been entered into the linker hash table. The \_bfd\_final\_link routine must consider each input file and include the symbols in

the output file. It may be convenient to do this when looking through the link\_order structures, or it may be done by stepping through the input\_bfds list.

The \_bfd\_final\_link routine must also traverse the global hash table to gather all the externally visible symbols. It is possible that most of the externally visible symbols may be written out when considering the symbols of each input file, but it is still necessary to traverse the hash table since the linker script may have defined some symbols that are not in any of the input files.

The strip field of the bfd\_link\_info structure controls which symbols are written out. The possible values are listed in bfdlink.h. If the value is strip\_some, then the keep\_hash field of the bfd\_link\_info structure is a hash table of symbols to keep; each symbol should be looked up in this hash table, and only symbols which are present should be included in the output file.

If the strip field of the bfd\_link\_info structure permits local symbols to be written out, the discard field is used to further controls which local symbols are included in the output file. If the value is discard\_1, then all local symbols which begin with a certain prefix are discarded; this is controlled by the bfd\_is\_local\_label\_name entry point.

The alout backend handles symbols by calling aout\_link\_write\_symbols on each input BFD and then traversing the global hash table with the function aout\_link\_write\_other\_symbol. It builds a string table while writing out the symbols, which is written to the output file at the end of NAME(aout,final\_link).

# 2.17.3.4 bfd\_link\_split\_section

#### **Synopsis**

```
bfd_boolean bfd_link_split_section (bfd *abfd, asection *sec);
```

### Description

Return nonzero if sec should be split during a reloceatable or final link.

```
#define bfd_link_split_section(abfd, sec) \
    BFD_SEND (abfd, _bfd_link_split_section, (abfd, sec))
```

# 2.17.3.5 bfd\_section\_already\_linked

### **Synopsis**

#### Description

Check if data has been already linked during a reloceatable or final link. Return TRUE if it has.

# 2.17.3.6 bfd\_generic\_define\_common\_symbol

#### **Synopsis**

```
bfd_boolean bfd_generic_define_common_symbol
   (bfd *output_bfd, struct bfd_link_info *info,
        struct bfd_link_hash_entry *h);
```

### Description

Convert common symbol h into a defined symbol. Return TRUE on success and FALSE on failure.

```
#define bfd_define_common_symbol(output_bfd, info, h) \
BFD_SEND (output_bfd, _bfd_define_common_symbol, (output_bfd, info, h))
```

# 2.17.3.7 bfd\_find\_version\_for\_sym

#### **Synopsis**

```
struct bfd_elf_version_tree * bfd_find_version_for_sym
  (struct bfd_elf_version_tree *verdefs,
    const char *sym_name, bfd_boolean *hide);
```

#### Description

Search an elf version script tree for symbol versioning info and export / don't-export status for a given symbol. Return non-NULL on success and NULL on failure; also sets the output 'hide' boolean parameter.

# 2.17.3.8 bfd\_hide\_sym\_by\_version

# **Synopsis**

```
bfd_boolean bfd_hide_sym_by_version
  (struct bfd_elf_version_tree *verdefs, const char *sym_name);
```

#### Description

Search an elf version script tree for symbol versioning info for a given symbol. Return TRUE if the symbol is hidden.

#### 2.18 Hash Tables

BFD provides a simple set of hash table functions. Routines are provided to initialize a hash table, to free a hash table, to look up a string in a hash table and optionally create an entry for it, and to traverse a hash table. There is currently no routine to delete an string from a hash table.

The basic hash table does not permit any data to be stored with a string. However, a hash table is designed to present a base class from which other types of hash tables may be derived. These derived types may store additional information with the string. Hash tables were implemented in this way, rather than simply providing a data pointer in a hash table entry, because they were designed for use by the linker back ends. The linker may create thousands of hash table entries, and the overhead of allocating private data and storing and following pointers becomes noticeable.

The basic hash table code is in hash.c.

# 2.18.1 Creating and freeing a hash table

To create a hash table, create an instance of a struct bfd\_hash\_table (defined in bfd.h) and call bfd\_hash\_table\_init (if you know approximately how many entries you will

need, the function bfd\_hash\_table\_init\_n, which takes a *size* argument, may be used). bfd\_hash\_table\_init returns FALSE if some sort of error occurs.

The function bfd\_hash\_table\_init take as an argument a function to use to create new entries. For a basic hash table, use the function bfd\_hash\_newfunc. See Section 2.18.4 [Deriving a New Hash Table Type], page 173, for why you would want to use a different value for this argument.

bfd\_hash\_table\_init will create an objalloc which will be used to allocate new entries. You may allocate memory on this objalloc using bfd\_hash\_allocate.

Use bfd\_hash\_table\_free to free up all the memory that has been allocated for a hash table. This will not free up the struct bfd\_hash\_table itself, which you must provide.

Use bfd\_hash\_set\_default\_size to set the default size of hash table to use.

## 2.18.2 Looking up or entering a string

The function bfd\_hash\_lookup is used both to look up a string in the hash table and to create a new entry.

If the *create* argument is FALSE, bfd\_hash\_lookup will look up a string. If the string is found, it will returns a pointer to a struct bfd\_hash\_entry. If the string is not found in the table bfd\_hash\_lookup will return NULL. You should not modify any of the fields in the returns struct bfd\_hash\_entry.

If the *create* argument is TRUE, the string will be entered into the hash table if it is not already there. Either way a pointer to a struct bfd\_hash\_entry will be returned, either to the existing structure or to a newly created one. In this case, a NULL return means that an error occurred.

If the *create* argument is TRUE, and a new entry is created, the *copy* argument is used to decide whether to copy the string onto the hash table objalloc or not. If *copy* is passed as FALSE, you must be careful not to deallocate or modify the string as long as the hash table exists.

## 2.18.3 Traversing a hash table

The function bfd\_hash\_traverse may be used to traverse a hash table, calling a function on each element. The traversal is done in a random order.

bfd\_hash\_traverse takes as arguments a function and a generic void \* pointer. The function is called with a hash table entry (a struct bfd\_hash\_entry \*) and the generic pointer passed to bfd\_hash\_traverse. The function must return a boolean value, which indicates whether to continue traversing the hash table. If the function returns FALSE, bfd\_hash\_traverse will stop the traversal and return immediately.

## 2.18.4 Deriving a new hash table type

Many uses of hash tables want to store additional information which each entry in the hash table. Some also find it convenient to store additional information with the hash table itself. This may be done using a derived hash table.

Since C is not an object oriented language, creating a derived hash table requires sticking together some boilerplate routines with a few differences specific to the type of hash table you want to create.

An example of a derived hash table is the linker hash table. The structures for this are defined in bfdlink.h. The functions are in linker.c.

You may also derive a hash table from an already derived hash table. For example, the a.out linker backend code uses a hash table derived from the linker hash table.

## 2.18.4.1 Define the derived structures

You must define a structure for an entry in the hash table, and a structure for the hash table itself.

The first field in the structure for an entry in the hash table must be of the type used for an entry in the hash table you are deriving from. If you are deriving from a basic hash table this is struct bfd\_hash\_entry, which is defined in bfd.h. The first field in the structure for the hash table itself must be of the type of the hash table you are deriving from itself. If you are deriving from a basic hash table, this is struct bfd\_hash\_table.

For example, the linker hash table defines struct bfd\_link\_hash\_entry (in bfdlink.h). The first field, root, is of type struct bfd\_hash\_entry. Similarly, the first field in struct bfd\_link\_hash\_table, table, is of type struct bfd\_hash\_table.

#### 2.18.4.2 Write the derived creation routine

You must write a routine which will create and initialize an entry in the hash table. This routine is passed as the function argument to bfd\_hash\_table\_init.

In order to permit other hash tables to be derived from the hash table you are creating, this routine must be written in a standard way.

The first argument to the creation routine is a pointer to a hash table entry. This may be NULL, in which case the routine should allocate the right amount of space. Otherwise the space has already been allocated by a hash table type derived from this one.

After allocating space, the creation routine must call the creation routine of the hash table type it is derived from, passing in a pointer to the space it just allocated. This will initialize any fields used by the base hash table.

Finally the creation routine must initialize any local fields for the new hash table type.

Here is a boilerplate example of a creation routine. *function\_name* is the name of the routine. *entry\_type* is the type of an entry in the hash table you are creating. *base\_newfunc* is the name of the creation routine of the hash table type your hash table is derived from.

## Description

The creation routine for the linker hash table, which is in linker.c, looks just like this example. function\_name is \_bfd\_link\_hash\_newfunc. entry\_type is struct bfd\_link\_hash\_entry. base\_newfunc is bfd\_hash\_newfunc, the creation routine for a basic hash table.

\_bfd\_link\_hash\_newfunc also initializes the local fields in a linker hash table entry: type, written and next.

## 2.18.4.3 Write other derived routines

You will want to write other routines for your new hash table, as well.

You will want an initialization routine which calls the initialization routine of the hash table you are deriving from and initializes any other local fields. For the linker hash table, this is <code>\_bfd\_link\_hash\_table\_init</code> in <code>linker.c</code>.

You will want a lookup routine which calls the lookup routine of the hash table you are deriving from and casts the result. The linker hash table uses bfd\_link\_hash\_lookup in linker.c (this actually takes an additional argument which it uses to decide how to return the looked up value).

You may want a traversal routine. This should just call the traversal routine of the hash table you are deriving from with appropriate casts. The linker hash table uses bfd\_link\_hash\_traverse in linker.c.

These routines may simply be defined as macros. For example, the a.out backend linker hash table, which is derived from the linker hash table, uses macros for the lookup and traversal routines. These are aout\_link\_hash\_lookup and aout\_link\_hash\_traverse in aoutx.h.

# 3 BFD back ends

## 3.1 What to Put Where

All of BFD lives in one directory.

## 3.2 a.out backends

## Description

BFD supports a number of different flavours of a out format, though the major differences are only the sizes of the structures on disk, and the shape of the relocation information.

The support is split into a basic support file 'aoutx.h' and other files which derive functions from the base. One derivation file is 'aoutf1.h' (for a.out flavour 1), and adds to the basic a.out functions support for sun3, sun4, 386 and 29k a.out files, to create a target jump vector for a specific target.

This information is further split out into more specific files for each machine, including 'sunos.c' for sun3 and sun4, 'newsos3.c' for the Sony NEWS, and 'demo64.c' for a demonstration of a 64 bit a.out format.

The base file 'aoutx.h' defines general mechanisms for reading and writing records to and from disk and various other methods which BFD requires. It is included by 'aout32.c' and 'aout64.c' to form the names aout\_32\_swap\_exec\_header\_in, aout\_64\_swap\_exec\_header\_in, etc.

As an example, this is what goes on to make the back end for a sun4, from 'aout32.c':

```
#define ARCH_SIZE 32
#include "aoutx.h"
```

Which exports names:

```
aout_32_canonicalize_reloc
aout_32_find_nearest_line
aout_32_get_lineno
aout_32_get_reloc_upper_bound
...
```

from 'sunos.c':

```
#define TARGET_NAME "a.out-sunos-big"
#define VECNAME sparc_aout_sunos_be_vec
#include "aoutf1.h"
```

requires all the names from 'aout32.c', and produces the jump vector

```
sparc_aout_sunos_be_vec
```

The file 'host-aout.c' is a special case. It is for a large set of hosts that use "more or less standard" a.out files, and for which cross-debugging is not interesting. It uses the standard 32-bit a.out support routines, but determines the file offsets and addresses of the text, data, and BSS sections, the machine architecture and machine type, and the entry point address, in a host-dependent manner. Once these values have been determined, generic code is used to handle the object file.

When porting it to run on a new system, you must supply:

```
HOST_PAGE_SIZE
HOST_SEGMENT_SIZE
HOST_MACHINE_ARCH (optional)
HOST_MACHINE_MACHINE (optional)
HOST_TEXT_START_ADDR
HOST_STACK_END_ADDR
```

in the file '../include/sys/h-XXX.h' (for your host). These values, plus the structures and macros defined in 'a.out.h' on your host system, will produce a BFD target that will access ordinary a.out files on your host. To configure a new machine to use 'host-aout.c', specify:

```
TDEFAULTS = -DDEFAULT_VECTOR=host_aout_big_vec
TDEPFILES= host-aout.o trad-core.o
```

in the 'config/XXX.mt' file, and modify 'configure.ac' to use the 'XXX.mt' file (by setting "bfd\_target=XXX") when your configuration is selected.

## 3.2.1 Relocations

### Description

The file 'aoutx.h' provides for both the *standard* and *extended* forms of a.out relocation records.

The standard records contain only an address, a symbol index, and a type field. The extended records (used on 29ks and sparcs) also have a full integer for an addend.

## 3.2.2 Internal entry points

#### Description

'aoutx.h' exports several routines for accessing the contents of an a.out file, which are gathered and exported in turn by various format specific files (eg sunos.c).

## 3.2.2.1 aout\_size\_swap\_exec\_header\_in

## **Synopsis**

```
void aout_size_swap_exec_header_in,
  (bfd *abfd,
    struct external_exec *bytes,
    struct internal_exec *execp);
```

#### Description

Swap the information in an executable header  $raw_bytes$  taken from a raw byte stream memory image into the internal exec header structure execp.

## 3.2.2.2 aout\_size\_swap\_exec\_header\_out

## **Synopsis**

```
void aout_size_swap_exec_header_out
  (bfd *abfd,
    struct internal_exec *execp,
    struct external_exec *raw_bytes);
```

## Description

Swap the information in an internal exec header structure execp into the buffer raw\_bytes ready for writing to disk.

## 3.2.2.3 aout\_size\_some\_aout\_object\_p

## **Synopsis**

```
const bfd_target *aout_size_some_aout_object_p
  (bfd *abfd,
    struct internal_exec *execp,
    const bfd_target *(*callback_to_real_object_p) (bfd *));
```

### Description

Some alout variant thinks that the file open in *abfd* checking is an alout file. Do some more checking, and set up for access if it really is. Call back to the calling environment's "finish up" function just before returning, to handle any last-minute setup.

## 3.2.2.4 aout\_size\_mkobject

## **Synopsis**

```
bfd_boolean aout_size_mkobject, (bfd *abfd);
```

## Description

Initialize BFD abfd for use with a out files.

## 3.2.2.5 aout\_size\_machine\_type

## Synopsis

```
enum machine_type aout_size_machine_type
  (enum bfd_architecture arch,
   unsigned long machine,
   bfd_boolean *unknown);
```

#### Description

Keep track of machine architecture and machine type for a.out's. Return the machine\_type for a particular architecture and machine, or M\_UNKNOWN if that exact architecture and machine can't be represented in a.out format.

If the architecture is understood, machine type 0 (default) is always understood.

## 3.2.2.6 aout\_size\_set\_arch\_mach

#### **Synopsis**

```
bfd_boolean aout_size_set_arch_mach,
   (bfd *,
      enum bfd_architecture arch,
      unsigned long machine);
```

#### Description

Set the architecture and the machine of the BFD *abfd* to the values *arch* and *machine*. Verify that *abfd*'s format can support the architecture required.

## 3.2.2.7 aout\_size\_new\_section\_hook

## **Synopsis**

```
bfd_boolean aout_size_new_section_hook,
    (bfd *abfd,
     asection *newsect);
```

### Description

Called by the BFD in response to a bfd\_make\_section request.

## 3.3 coff backends

BFD supports a number of different flavours of coff format. The major differences between formats are the sizes and alignments of fields in structures on disk, and the occasional extra field.

Coff in all its varieties is implemented with a few common files and a number of implementation specific files. For example, The 88k bcs coff format is implemented in the file 'coff-m88k.c'. This file #includes 'coff/m88k.h' which defines the external structure of the coff format for the 88k, and 'coff/internal.h' which defines the internal structure. 'coff-m88k.c' also defines the relocations used by the 88k format See Section 2.10 [Relocations], page 52.

The Intel i960 processor version of coff is implemented in 'coff-i960.c'. This file has the same structure as 'coff-m88k.c', except that it includes 'coff/i960.h' rather than 'coff-m88k.h'.

## 3.3.1 Porting to a new version of coff

The recommended method is to select from the existing implementations the version of coff which is most like the one you want to use. For example, we'll say that i386 coff is the one you select, and that your coff flavour is called foo. Copy 'i386coff.c' to 'foocoff.c', copy '../include/coff/i386.h' to '../include/coff/foo.h', and add the lines to 'targets.c' and 'Makefile.in' so that your new back end is used. Alter the shapes of the structures in '../include/coff/foo.h' so that they match what you need. You will probably also have to add #ifdefs to the code in 'coff/internal.h' and 'coffcode.h' if your version of coff is too wild.

You can verify that your new BFD backend works quite simply by building 'objdump' from the 'binutils' directory, and making sure that its version of what's going on and your host system's idea (assuming it has the pretty standard coff dump utility, usually called att-dump or just dump) are the same. Then clean up your code, and send what you've done to Cygnus. Then your stuff will be in the next release, and you won't have to keep integrating it.

## 3.3.2 How the coff backend works

## 3.3.2.1 File layout

The Coff backend is split into generic routines that are applicable to any Coff target and routines that are specific to a particular target. The target-specific routines are further split into ones which are basically the same for all Coff targets except that they use the external symbol format or use different values for certain constants.

The generic routines are in 'coffgen.c'. These routines work for any Coff target. They use some hooks into the target specific code; the hooks are in a bfd\_coff\_backend\_data structure, one of which exists for each target.

The essentially similar target-specific routines are in 'coffcode.h'. This header file includes executable C code. The various Coff targets first include the appropriate Coff header file, make any special defines that are needed, and then include 'coffcode.h'.

Some of the Coff targets then also have additional routines in the target source file itself.

For example, 'coff-i960.c' includes 'coff/internal.h' and 'coff/i960.h'. It then defines a few constants, such as I960, and includes 'coffcode.h'. Since the i960 has complex relocation types, 'coff-i960.c' also includes some code to manipulate the i960 relocs. This code is not in 'coffcode.h' because it would not be used by any other target.

## 3.3.2.2 Coff long section names

In the standard Coff object format, section names are limited to the eight bytes available in the s\_name field of the SCNHDR section header structure. The format requires the field to be NUL-padded, but not necessarily NUL-terminated, so the longest section names permitted are a full eight characters.

The Microsoft PE variants of the Coff object file format add an extension to support the use of long section names. This extension is defined in section 4 of the Microsoft PE/COFF specification (rev 8.1). If a section name is too long to fit into the section header's s\_name field, it is instead placed into the string table, and the s\_name field is filled with a slash ("/") followed by the ASCII decimal representation of the offset of the full name relative to the string table base.

Note that this implies that the extension can only be used in object files, as executables do not contain a string table. The standard specifies that long section names from objects emitted into executable images are to be truncated.

However, as a GNU extension, BFD can generate executable images that contain a string table and long section names. This would appear to be technically valid, as the standard only says that Coff debugging information is deprecated, not forbidden, and in practice it works, although some tools that parse PE files expecting the MS standard format may become confused; 'PEview' is one known example.

The functionality is supported in BFD by code implemented under the control of the macro COFF\_LONG\_SECTION\_NAMES. If not defined, the format does not support long section names in any way. If defined, it is used to initialise a flag, \_bfd\_coff\_long\_section\_names, and a hook function pointer, \_bfd\_coff\_set\_long\_section\_names, in the Coff backend data structure. The flag controls the generation of long section names in output BFDs at runtime; if it is false, as it will be by default when generating an executable image, long section names are truncated; if true, the long section names extension is employed. The hook points to a function that allows the value of the flag to be altered at runtime, on formats that support long section names at all; on other formats it points to a stub that returns an error indication.

With input BFDs, the flag is set according to whether any long section names are detected while reading the section headers. For a completely new BFD, the flag is set to the default for the target format. This information can be used by a client of the BFD library when deciding what output format to generate, and means that a BFD that is opened for read

and subsequently converted to a writeable BFD and modified in-place will retain whatever format it had on input.

If COFF\_LONG\_SECTION\_NAMES is simply defined (blank), or is defined to the value "1", then long section names are enabled by default; if it is defined to the value zero, they are disabled by default (but still accepted in input BFDs). The header 'coffcode.h' defines a macro, COFF\_DEFAULT\_LONG\_SECTION\_NAMES, which is used in the backends to initialise the backend data structure fields appropriately; see the comments for further detail.

## 3.3.2.3 Bit twiddling

Each flavour of coff supported in BFD has its own header file describing the external layout of the structures. There is also an internal description of the coff layout, in 'coff/internal.h'. A major function of the coff backend is swapping the bytes and twiddling the bits to translate the external form of the structures into the normal internal form. This is all performed in the bfd\_swap\_thing\_direction routines. Some elements are different sizes between different versions of coff; it is the duty of the coff version specific include file to override the definitions of various packing routines in 'coffcode.h'. E.g., the size of line number entry in coff is sometimes 16 bits, and sometimes 32 bits. #defineing PUT\_LNSZ\_LNNO and GET\_LNSZ\_ LNNO will select the correct one. No doubt, some day someone will find a version of coff which has a varying field size not catered to at the moment. To port BFD, that person will have to add more #defines. Three of the bit twiddling routines are exported to gdb; coff\_swap\_aux\_in, coff\_swap\_sym\_in and coff\_swap\_lineno\_in. GDB reads the symbol table on its own, but uses BFD to fix things up. More of the bit twiddlers are exported for gas; coff\_swap\_aux\_out, coff\_swap\_sym\_out, coff\_swap\_lineno\_out, coff\_swap\_ reloc\_out, coff\_swap\_filehdr\_out, coff\_swap\_aouthdr\_out, coff\_swap\_scnhdr\_out. Gas currently keeps track of all the symbol table and reloc drudgery itself, thereby saving the internal BFD overhead, but uses BFD to swap things on the way out, making cross ports much safer. Doing so also allows BFD (and thus the linker) to use the same header files as gas, which makes one avenue to disaster disappear.

## 3.3.2.4 Symbol reading

The simple canonical form for symbols used by BFD is not rich enough to keep all the information available in a coff symbol table. The back end gets around this problem by keeping the original symbol table around, "behind the scenes".

When a symbol table is requested (through a call to bfd\_canonicalize\_symtab), a request gets through to coff\_get\_normalized\_symtab. This reads the symbol table from the coff file and swaps all the structures inside into the internal form. It also fixes up all the pointers in the table (represented in the file by offsets from the first symbol in the table) into physical pointers to elements in the new internal table. This involves some work since the meanings of fields change depending upon context: a field that is a pointer to another structure in the symbol table at one moment may be the size in bytes of a structure at the next. Another pass is made over the table. All symbols which mark file names (C\_FILE symbols) are modified so that the internal string points to the value in the auxent (the real filename) rather than the normal text associated with the symbol (".file").

At this time the symbol names are moved around. Coff stores all symbols less than nine characters long physically within the symbol table; longer strings are kept at the end of the

file in the string table. This pass moves all strings into memory and replaces them with pointers to the strings.

The symbol table is massaged once again, this time to create the canonical table used by the BFD application. Each symbol is inspected in turn, and a decision made (using the sclass field) about the various flags to set in the asymbol. See Section 2.7 [Symbols], page 41. The generated canonical table shares strings with the hidden internal symbol table.

Any linenumbers are read from the coff file too, and attached to the symbols which own the functions the linenumbers belong to.

## 3.3.2.5 Symbol writing

Writing a symbol to a coff file which didn't come from a coff file will lose any debugging information. The asymbol structure remembers the BFD from which the symbol was taken, and on output the back end makes sure that the same destination target as source target is present.

When the symbols have come from a coff file then all the debugging information is preserved. Symbol tables are provided for writing to the back end in a vector of pointers to pointers. This allows applications like the linker to accumulate and output large symbol tables without having to do too much byte copying.

This function runs through the provided symbol table and patches each symbol marked as a file place holder (C\_FILE) to point to the next file place holder in the list. It also marks each offset field in the list with the offset from the first symbol of the current symbol.

Another function of this procedure is to turn the canonical value form of BFD into the form used by coff. Internally, BFD expects symbol values to be offsets from a section base; so a symbol physically at 0x120, but in a section starting at 0x100, would have the value 0x20. Coff expects symbols to contain their final value, so symbols have their values changed at this point to reflect their sum with their owning section. This transformation uses the output\_section field of the asymbol's asection See Section 2.6 [Sections], page 22.

#### • coff\_mangle\_symbols

This routine runs though the provided symbol table and uses the offsets generated by the previous pass and the pointers generated when the symbol table was read in to create the structured hierarchy required by coff. It changes each pointer to a symbol into the index into the symbol table of the asymbol.

## • coff\_write\_symbols

This routine runs through the symbol table and patches up the symbols from their internal form into the coff way, calls the bit twiddlers, and writes out the table to the file.

## 3.3.2.6 coff\_symbol\_type

#### Description

The hidden information for an asymbol is described in a combined\_entry\_type:

```
typedef struct coff_ptr_struct
{
   /* Remembers the offset from the first symbol in the file for
     this symbol. Generated by coff_renumber_symbols. */
```

```
unsigned int offset;
 /* Should the value of this symbol be renumbered. Used for
     XCOFF C_BSTAT symbols. Set by coff_slurp_symbol_table. */
 unsigned int fix_value : 1;
 /* Should the tag field of this symbol be renumbered.
     Created by coff_pointerize_aux. */
 unsigned int fix_tag : 1;
 /* Should the endidx field of this symbol be renumbered.
     Created by coff_pointerize_aux. */
 unsigned int fix_end : 1;
 /* Should the x_csect.x_scnlen field be renumbered.
     Created by coff_pointerize_aux. */
 unsigned int fix_scnlen : 1;
 /* Fix up an XCOFF C_BINCL/C_EINCL symbol. The value is the
     index into the line number entries. Set by coff_slurp_symbol_table. */
 unsigned int fix_line : 1;
 /* The container for the symbol structure as read and translated
     from the file. */
 union
   union internal_auxent auxent;
   struct internal_syment syment;
 } u:
/* Selector for the union above. */
bfd_boolean is_sym;
} combined_entry_type;
/* Each canonical asymbol really looks like this: */
typedef struct coff_symbol_struct
 /* The actual symbol which the rest of BFD works with */
 asymbol symbol;
 /* A pointer to the hidden information for this symbol */
 combined_entry_type *native;
 /* A pointer to the linenumber information for this symbol */
 struct lineno_cache_entry *lineno;
```

```
/* Have the line numbers been relocated yet ? */
      bfd_boolean done_lineno;
     } coff_symbol_type;
3.3.2.7 bfd_coff_backend_data
     /* COFF symbol classifications. */
     enum coff_symbol_classification
       /* Global symbol. */
      COFF_SYMBOL_GLOBAL,
       /* Common symbol. */
      COFF_SYMBOL_COMMON,
       /* Undefined symbol. */
      COFF_SYMBOL_UNDEFINED,
       /* Local symbol. */
      COFF_SYMBOL_LOCAL,
       /* PE section symbol. */
      COFF_SYMBOL_PE_SECTION
    };
Special entry points for gdb to swap in coff symbol table parts:
     typedef struct
      void (*_bfd_coff_swap_aux_in)
         (bfd *, void *, int, int, int, void *);
      void (*_bfd_coff_swap_sym_in)
         (bfd *, void *, void *);
      void (*_bfd_coff_swap_lineno_in)
         (bfd *, void *, void *);
       unsigned int (*_bfd_coff_swap_aux_out)
         (bfd *, void *, int, int, int, void *);
       unsigned int (*_bfd_coff_swap_sym_out)
         (bfd *, void *, void *);
       unsigned int (*_bfd_coff_swap_lineno_out)
         (bfd *, void *, void *);
       unsigned int (*_bfd_coff_swap_reloc_out)
         (bfd *, void *, void *);
```

```
unsigned int (*_bfd_coff_swap_filehdr_out)
  (bfd *, void *, void *);
unsigned int (*_bfd_coff_swap_aouthdr_out)
  (bfd *, void *, void *);
unsigned int (*_bfd_coff_swap_scnhdr_out)
  (bfd *, void *, void *);
unsigned int _bfd_filhsz;
unsigned int _bfd_aoutsz;
unsigned int _bfd_scnhsz;
unsigned int _bfd_symesz;
unsigned int _bfd_auxesz;
unsigned int _bfd_relsz;
unsigned int _bfd_linesz;
unsigned int _bfd_filnmlen;
bfd_boolean _bfd_coff_long_filenames;
bfd_boolean _bfd_coff_long_section_names;
bfd_boolean (*_bfd_coff_set_long_section_names)
  (bfd *, int);
unsigned int _bfd_coff_default_section_alignment_power;
bfd_boolean _bfd_coff_force_symnames_in_strings;
unsigned int _bfd_coff_debug_string_prefix_length;
unsigned int _bfd_coff_max_nscns;
void (*_bfd_coff_swap_filehdr_in)
  (bfd *, void *, void *);
void (*_bfd_coff_swap_aouthdr_in)
  (bfd *, void *, void *);
void (*_bfd_coff_swap_scnhdr_in)
  (bfd *, void *, void *);
void (*_bfd_coff_swap_reloc_in)
  (bfd *abfd, void *, void *);
bfd_boolean (*_bfd_coff_bad_format_hook)
  (bfd *, void *);
bfd_boolean (*_bfd_coff_set_arch_mach_hook)
  (bfd *, void *);
void * (*_bfd_coff_mkobject_hook)
```

```
(bfd *, void *, void *);
bfd_boolean (*_bfd_styp_to_sec_flags_hook)
  (bfd *, void *, const char *, asection *, flagword *);
void (*_bfd_set_alignment_hook)
  (bfd *, asection *, void *);
bfd_boolean (*_bfd_coff_slurp_symbol_table)
  (bfd *);
bfd_boolean (*_bfd_coff_symname_in_debug)
  (bfd *, struct internal_syment *);
bfd_boolean (*_bfd_coff_pointerize_aux_hook)
  (bfd *, combined_entry_type *, combined_entry_type *,
          unsigned int, combined_entry_type *);
bfd_boolean (*_bfd_coff_print_aux)
  (bfd *, FILE *, combined_entry_type *, combined_entry_type *,
          combined_entry_type *, unsigned int);
void (*_bfd_coff_reloc16_extra_cases)
  (bfd *, struct bfd_link_info *, struct bfd_link_order *, arelent *,
         bfd_byte *, unsigned int *, unsigned int *);
int (*_bfd_coff_reloc16_estimate)
  (bfd *, asection *, arelent *, unsigned int,
          struct bfd_link_info *);
enum coff_symbol_classification (*_bfd_coff_classify_symbol)
  (bfd *, struct internal_syment *);
bfd_boolean (*_bfd_coff_compute_section_file_positions)
  (bfd *);
bfd_boolean (*_bfd_coff_start_final_link)
  (bfd *, struct bfd_link_info *);
bfd_boolean (*_bfd_coff_relocate_section)
  (bfd *, struct bfd_link_info *, bfd *, asection *, bfd_byte *,
          struct internal_reloc *, struct internal_syment *, asection **);
reloc_howto_type *(*_bfd_coff_rtype_to_howto)
  (bfd *, asection *, struct internal_reloc *,
          struct coff_link_hash_entry *, struct internal_syment *,
          bfd_vma *);
```

```
bfd_boolean (*_bfd_coff_adjust_symndx)
    (bfd *, struct bfd_link_info *, bfd *, asection *,
            struct internal_reloc *, bfd_boolean *);
 bfd_boolean (*_bfd_coff_link_add_one_symbol)
    (struct bfd_link_info *, bfd *, const char *, flagword,
            asection *, bfd_vma, const char *, bfd_boolean, bfd_boolean,
            struct bfd_link_hash_entry **);
 bfd_boolean (*_bfd_coff_link_output_has_begun)
    (bfd *, struct coff_final_link_info *);
 bfd_boolean (*_bfd_coff_final_link_postscript)
    (bfd *, struct coff_final_link_info *);
 bfd_boolean (*_bfd_coff_print_pdata)
    (bfd *, void *);
} bfd_coff_backend_data;
#define coff_backend_info(abfd) \
  ((bfd_coff_backend_data *) (abfd)->xvec->backend_data)
#define bfd_coff_swap_aux_in(a,e,t,c,ind,num,i) \
  ((coff_backend_info (a)->_bfd_coff_swap_aux_in) (a,e,t,c,ind,num,i))
#define bfd_coff_swap_sym_in(a,e,i) \
  ((coff_backend_info (a)->_bfd_coff_swap_sym_in) (a,e,i))
#define bfd_coff_swap_lineno_in(a,e,i) \
  ((coff_backend_info ( a)->_bfd_coff_swap_lineno_in) (a,e,i))
#define bfd_coff_swap_reloc_out(abfd, i, o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_reloc_out) (abfd, i, o))
#define bfd_coff_swap_lineno_out(abfd, i, o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_lineno_out) (abfd, i, o))
#define bfd_coff_swap_aux_out(a,i,t,c,ind,num,o) \
  ((coff_backend_info (a)->_bfd_coff_swap_aux_out) (a,i,t,c,ind,num,o))
#define bfd_coff_swap_sym_out(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_sym_out) (abfd, i, o))
#define bfd_coff_swap_scnhdr_out(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_scnhdr_out) (abfd, i, o))
```

```
#define bfd_coff_swap_filehdr_out(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_filehdr_out) (abfd, i, o))
#define bfd_coff_swap_aouthdr_out(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_aouthdr_out) (abfd, i, o))
#define bfd_coff_filhsz(abfd) (coff_backend_info (abfd)->_bfd_filhsz)
#define bfd_coff_aoutsz(abfd) (coff_backend_info (abfd)->_bfd_aoutsz)
#define bfd_coff_scnhsz(abfd) (coff_backend_info (abfd)->_bfd_scnhsz)
#define bfd_coff_symesz(abfd) (coff_backend_info (abfd)->_bfd_symesz)
#define bfd_coff_auxesz(abfd) (coff_backend_info (abfd)->_bfd_auxesz)
#define bfd_coff_relsz(abfd) (coff_backend_info (abfd)->_bfd_relsz)
#define bfd_coff_linesz(abfd) (coff_backend_info (abfd)->_bfd_linesz)
#define bfd_coff_filnmlen(abfd) (coff_backend_info (abfd)->_bfd_filnmlen)
#define bfd_coff_long_filenames(abfd) \
  (coff_backend_info (abfd)->_bfd_coff_long_filenames)
#define bfd_coff_long_section_names(abfd) \
  (coff_backend_info (abfd)->_bfd_coff_long_section_names)
#define bfd_coff_set_long_section_names(abfd, enable) \
  ((coff_backend_info (abfd)->_bfd_coff_set_long_section_names) (abfd, enable))
#define bfd_coff_default_section_alignment_power(abfd) \
  (coff_backend_info (abfd)->_bfd_coff_default_section_alignment_power)
#define bfd_coff_max_nscns(abfd) \
  (coff_backend_info (abfd)->_bfd_coff_max_nscns)
#define bfd_coff_swap_filehdr_in(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_filehdr_in) (abfd, i, o))
#define bfd_coff_swap_aouthdr_in(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_aouthdr_in) (abfd, i, o))
#define bfd_coff_swap_scnhdr_in(abfd, i,o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_scnhdr_in) (abfd, i, o))
#define bfd_coff_swap_reloc_in(abfd, i, o) \
  ((coff_backend_info (abfd)->_bfd_coff_swap_reloc_in) (abfd, i, o))
#define bfd_coff_bad_format_hook(abfd, filehdr) \
  ((coff_backend_info (abfd)->_bfd_coff_bad_format_hook) (abfd, filehdr))
#define bfd_coff_set_arch_mach_hook(abfd, filehdr)\
  ((coff_backend_info (abfd)->_bfd_coff_set_arch_mach_hook) (abfd, filehdr))
#define bfd_coff_mkobject_hook(abfd, filehdr, aouthdr)\
  ((coff_backend_info (abfd)->_bfd_coff_mkobject_hook)\
   (abfd, filehdr, aouthdr))
```

```
#define bfd_coff_styp_to_sec_flags_hook(abfd, scnhdr, name, section, flags_ptr)\
■
  ((coff_backend_info (abfd)->_bfd_styp_to_sec_flags_hook)\
   (abfd, scnhdr, name, section, flags_ptr))
#define bfd_coff_set_alignment_hook(abfd, sec, scnhdr)\
  ((coff_backend_info (abfd)->_bfd_set_alignment_hook) (abfd, sec, scnhdr))
#define bfd_coff_slurp_symbol_table(abfd)\
  ((coff_backend_info (abfd)->_bfd_coff_slurp_symbol_table) (abfd))
#define bfd_coff_symname_in_debug(abfd, sym)\
  ((coff_backend_info (abfd)->_bfd_coff_symname_in_debug) (abfd, sym))
#define bfd_coff_force_symnames_in_strings(abfd)\
  (coff_backend_info (abfd)->_bfd_coff_force_symnames_in_strings)
#define bfd_coff_debug_string_prefix_length(abfd)\
  (coff_backend_info (abfd)->_bfd_coff_debug_string_prefix_length)
#define bfd_coff_print_aux(abfd, file, base, symbol, aux, indaux)\
  ((coff_backend_info (abfd)->_bfd_coff_print_aux)\
   (abfd, file, base, symbol, aux, indaux))
#define bfd_coff_reloc16_extra_cases(abfd, link_info, link_order,\
                                     reloc, data, src_ptr, dst_ptr)\
  ((coff_backend_info (abfd)->_bfd_coff_reloc16_extra_cases)\
   (abfd, link_info, link_order, reloc, data, src_ptr, dst_ptr))
#define bfd_coff_reloc16_estimate(abfd, section, reloc, shrink, link_info)\■
  ((coff_backend_info (abfd)->_bfd_coff_reloc16_estimate)\
   (abfd, section, reloc, shrink, link_info))
#define bfd_coff_classify_symbol(abfd, sym)\
  ((coff_backend_info (abfd)->_bfd_coff_classify_symbol)\
   (abfd, sym))
#define bfd_coff_compute_section_file_positions(abfd)\
  ((coff\_backend\_info\ (abfd) ->\_bfd\_coff\_compute\_section\_file\_positions) \setminus \\
   (abfd))
#define bfd_coff_start_final_link(obfd, info)
  ((coff_backend_info (obfd)->_bfd_coff_start_final_link)\
   (obfd, info))
#define bfd_coff_relocate_section(obfd,info,ibfd,o,con,rel,isyms,secs)\
  ((coff_backend_info (ibfd)->_bfd_coff_relocate_section)\
   (obfd, info, ibfd, o, con, rel, isyms, secs))
#define bfd_coff_rtype_to_howto(abfd, sec, rel, h, sym, addendp)\
```

```
((coff_backend_info (abfd)->_bfd_coff_rtype_to_howto)\
   (abfd, sec, rel, h, sym, addendp))
#define bfd_coff_adjust_symndx(obfd, info, ibfd, sec, rel, adjustedp)
  ((coff_backend_info (abfd)->_bfd_coff_adjust_symndx)\
   (obfd, info, ibfd, sec, rel, adjustedp))
#define bfd_coff_link_add_one_symbol(info, abfd, name, flags, section,\
                                     value, string, cp, coll, hashp)\
  ((coff_backend_info (abfd)->_bfd_coff_link_add_one_symbol)\
   (info, abfd, name, flags, section, value, string, cp, coll, hashp))
#define bfd_coff_link_output_has_begun(a,p) \
  ((coff_backend_info (a)->_bfd_coff_link_output_has_begun) (a, p))
#define bfd_coff_final_link_postscript(a,p) \
  ((coff_backend_info (a)->_bfd_coff_final_link_postscript) (a, p))
#define bfd_coff_have_print_pdata(a) \
  (coff_backend_info (a)->_bfd_coff_print_pdata)
#define bfd_coff_print_pdata(a,p) \
  ((coff_backend_info (a)->_bfd_coff_print_pdata) (a, p))
/* Macro: Returns true if the bfd is a PE executable as opposed to a
  PE object file. */
#define bfd_pei_p(abfd) \
  (CONST_STRNEQ ((abfd)->xvec->name, "pei-"))
```

## 3.3.2.8 Writing relocations

To write relocations, the back end steps though the canonical relocation table and create an internal\_reloc. The symbol index to use is removed from the offset field in the symbol table supplied. The address comes directly from the sum of the section base address and the relocation offset; the type is dug directly from the howto field. Then the internal\_reloc is swapped into the shape of an external\_reloc and written out to disk.

## 3.3.2.9 Reading linenumbers

Creating the linenumber table is done by reading in the entire coff linenumber table, and creating another table for internal use.

A coff linenumber table is structured so that each function is marked as having a line number of 0. Each line within the function is an offset from the first line in the function. The base of the line number information for the table is stored in the symbol associated with the function.

Note: The PE format uses line number 0 for a flag indicating a new source file.

The information is copied from the external to the internal table, and each symbol which marks a function is marked by pointing its...

How does this work?

## 3.3.2.10 Reading relocations

Coff relocations are easily transformed into the internal BFD form (arelent).

Reading a coff relocation table is done in the following stages:

- Read the entire coff relocation table into memory.
- Process each relocation in turn; first swap it from the external to the internal form.
- Turn the symbol referenced in the relocation's symbol index into a pointer into the canonical symbol table. This table is the same as the one returned by a call to bfd\_canonicalize\_symtab. The back end will call that routine and save the result if a canonicalization hasn't been done.
- The reloc index is turned into a pointer to a howto structure, in a back end specific way. For instance, the 386 and 960 use the r\_type to directly produce an index into a howto table vector; the 88k subtracts a number from the r\_type field and creates an addend field.

## 3.4 ELF backends

BFD support for ELF formats is being worked on. Currently, the best supported back ends are for sparc and i386 (running svr4 or Solaris 2).

Documentation of the internals of the support code still needs to be written. The code is changing quickly enough that we haven't bothered yet.

## 3.5 mmo backend

The mmo object format is used exclusively together with Professor Donald E. Knuth's educational 64-bit processor MMIX. The simulator mmix which is available at http://mmix.cs.hm.edu/src/index.html understands this format. That package also includes a combined assembler and linker called mmixal. The mmo format has no advantages feature-wise compared to e.g. ELF. It is a simple non-relocatable object format with no support for archives or debugging information, except for symbol value information and line numbers (which is not yet implemented in BFD). See http://mmix.cs.hm.edu/for more information about MMIX. The ELF format is used for intermediate object files in the BFD implementation.

# 3.5.1 File layout

The mmo file contents is not partitioned into named sections as with e.g. ELF. Memory areas is formed by specifying the location of the data that follows. Only the memory area '0x0000...00' to '0x01ff...ff' is executable, so it is used for code (and constants) and the area '0x2000...00' to '0x20ff...ff' is used for writable data. See Section 3.5.3 [mmo section mapping], page 195.

There is provision for specifying "special data" of 65536 different types. We use type 80 (decimal), arbitrarily chosen the same as the ELF e\_machine number for MMIX, filling it with section information normally found in ELF objects. See Section 3.5.3 [mmo section mapping], page 195.

Contents is entered as 32-bit words, xor:ed over previous contents, always zero-initialized. A word that starts with the byte '0x98' forms a command called a 'lopcode', where the next byte distinguished between the thirteen lopcodes. The two remaining bytes, called the 'Y' and 'Z' fields, or the 'YZ' field (a 16-bit big-endian number), are used for various purposes different for each lopcode. As documented in http://mmix.cs.hm.edu/doc/mmixal.pdf, the lopcodes are:

lop\_quote

0x98000001. The next word is contents, regardless of whether it starts with 0x98 or not.

lop\_loc 0x9801YYZZ, where 'Z' is 1 or 2. This is a location directive, setting the location for the next data to the next 32-bit word (for Z=1) or 64-bit word (for Z=2), plus  $Y*2^56$ . Normally 'Y' is 0 for the text segment and 2 for the data segment. Beware that the low bits of non- tetrabyte-aligned values are silently discarded when being automatically incremented and when storing contents (in contrast to e.g. its use as current location when followed by lop\_fixo et al before the next possibly-quoted tetrabyte contents).

lop\_skip 0x9802YYZZ. Increase the current location by 'YZ' bytes.

lop\_fixo 0x9803YYZZ, where 'Z' is 1 or 2. Store the current location as 64 bits into the location pointed to by the next 32-bit (Z = 1) or 64-bit (Z = 2) word, plus  $Y * 2^56$ .

lop\_fixr 0x9804YYZZ. 'YZ' is stored into the current location plus 2-4\*YZ.

lop\_fixrx

0x980500ZZ. 'Z' is 16 or 24. A value 'L' derived from the following 32-bit word are used in a manner similar to 'YZ' in lop\_fixr: it is xor:ed into the current location minus 4\*L. The first byte of the word is 0 or 1. If it is 1, then  $L = (lowest24bitsofword) - 2^{Z}$ , if 0, then L = (lowest24bitsofword).

lop\_file 0x9806YYZZ. 'Y' is the file number, 'Z' is count of 32-bit words. Set the file number to 'Y' and the line counter to 0. The next Z\*4 bytes contain the file name, padded with zeros if the count is not a multiple of four. The same 'Y' may occur multiple times, but 'Z' must be 0 for all but the first occurrence.

lop\_line 0x9807YYZZ. 'YZ' is the line number. Together with lop\_file, it forms the source location for the next 32-bit word. Note that for each non-lopcode 32-bit word, line numbers are assumed incremented by one.

lop\_spec 0x9808YYZZ. 'YZ' is the type number. Data until the next lopcode other than lop\_quote forms special data of type 'YZ'. See Section 3.5.3 [mmo section mapping], page 195.

Other types than 80, (or type 80 with a content that does not parse) is stored in sections named .MMIX.spec\_data.n where n is the 'YZ'-type. The flags for such a sections say not to allocate or load the data. The vma is 0. Contents of multiple occurrences of special data n is concatenated to the data of the previous lop\_spec ns. The location in data or code at which the lop\_spec occurred is lost.

lop\_pre 0x980901ZZ. The first lopcode in a file. The 'Z' field forms the length of header information in 32-bit words, where the first word tells the time in seconds since '00:00:00 GMT Jan 1 1970'.

lop\_post 0x980a00ZZ. Z > 32. This lopcode follows after all content-generating lopcodes in a program. The 'Z' field denotes the value of 'rG' at the beginning of the program. The following 256 - Z big-endian 64-bit words are loaded into global registers '\$G' . . . '\$255'.

lop\_stab 0x980b0000. The next-to-last lopcode in a program. Must follow immediately after the lop\_post lopcode and its data. After this lopcode follows all symbols in a compressed format (see Section 3.5.2 [Symbol-table], page 193).

lop\_end 0x980cYYZZ. The last lopcode in a program. It must follow the lop\_stab lopcode and its data. The 'YZ' field contains the number of 32-bit words of symbol table information after the preceding lop\_stab lopcode.

Note that the lopcode "fixups"; lop\_fixr, lop\_fixrx and lop\_fixo are not generated by BFD, but are handled. They are generated by mmixal.

This trivial one-label, one-instruction file:

```
:Main TRAP 1,2,3
```

can be represented this way in mmo:

```
0x98090101 - lop_pre, one 32-bit word with timestamp.
<timestamp>
0x98010002 - lop_loc, text segment, using a 64-bit address.
             Note that mmixal does not emit this for the file above.
0x00000000 - Address, high 32 bits.
0x00000000 - Address, low 32 bits.
0x98060002 - lop_file, 2 32-bit words for file-name.
0x74657374 - "test"
0x2e730000 - ".s\0\0"
0x98070001 - lop_line, line 1.
0x00010203 - TRAP 1,2,3
0x980a00ff - lop_post, setting $255 to 0.
0x0000000
0x0000000
0x980b0000 - lop_stab for ":Main" = 0, serial 1.
0x203a4040
             See Section 3.5.2 [Symbol-table], page 193.
0x10404020
0x4d206120
0x69016e00
0x81000000
0x980c0005 - lop_end; symbol table contained five 32-bit words.
```

## 3.5.2 Symbol table format

From mmixal.w (or really, the generated mmixal.tex) in the MMIXware package which also contains the mmix simulator: "Symbols are stored and retrieved by means of a 'ternary search trie', following ideas of Bentley and Sedgewick. (See ACM-SIAM Symp. on Discrete Algorithms '8' (1997), 360–369; R.Sedgewick, 'Algorithms in C' (Reading, Mass. Addison-Wesley, 1998), '15.4'.) Each trie node stores a character, and there are branches to subtries for the cases where a given character is less than, equal to, or greater than the character in the trie. There also is a pointer to a symbol table entry if a symbol ends at the current node."

So it's a tree encoded as a stream of bytes. The stream of bytes acts on a single virtual global symbol, adding and removing characters and signalling complete symbol points. Here, we read the stream and create symbols at the completion points.

First, there's a control byte m. If any of the listed bits in m is nonzero, we execute what stands at the right, in the listed order:

## (MMO3\_LEFT)

#### (MMO3\_SYMBITS)

0x2f - Read the next byte as a character and store it in the current character position; increment character position. Test the bits of m:

## (MMO3\_WCHAR)

### (MMO3\_TYPEBITS)

Oxf - We have a complete symbol; parse the type, value and serial number and do what should be done with a symbol. The type and length information is in j = (m & Oxf).

#### (MMO3\_REGQUAL\_BITS)

- j <= 8: An absolute symbol. Read j bytes as the big-endian number the symbol equals. A j = 2 with two zero bytes denotes an unknown symbol.
- j > 8: As with j <= 8, but add (0x20 << 56) to the value in the following j 8 bytes.

Then comes the serial number, as a variant of uleb128, but better named ubeb128: Read bytes and shift the previous value left 7 (multiply by 128). Add in the new byte, repeat until a byte has bit 7 set. The serial number is the computed value minus 128.

#### (MMO3\_MIDDLE)

0x20 - Traverse middle trie. (Read a new command byte and recurse.) Decrement character position.

#### (MMO3\_RIGHT)

0x10 - Traverse right trie. (Read a new command byte and recurse.)

```
Let's look again at the lop_stab for the trivial file (see Section 3.5.1 [File layout], page 191).
```

```
0x980b0000 - lop_stab for ":Main" = 0, serial 1.
0x203a4040
0x10404020
0x4d206120
0x69016e00
0x81000000
```

This forms the trivial trie (note that the path between ":" and "M" is redundant):

```
203a
40
40
10
40
40
      "M"
204d
      "a"
2061
2069
      "i"
016e
      "n" is the last character in a full symbol, and
      with a value represented in one byte.
      The value is 0.
00
81
      The serial number is 1.
```

# 3.5.3 mmo section mapping

The implementation in BFD uses special data type 80 (decimal) to encapsulate and describe named sections, containing e.g. debug information. If needed, any datum in the encapsulation will be quoted using lop\_quote. First comes a 32-bit word holding the number of 32-bit words containing the zero-terminated zero-padded segment name. After the name there's a 32-bit word holding flags describing the section type. Then comes a 64-bit big-endian word with the section length (in bytes), then another with the section start address. Depending on the type of section, the contents might follow, zero-padded to 32-bit boundary. For a loadable section (such as data or code), the contents might follow at some later point, not necessarily immediately, as a lop\_loc with the same start address as in the section description, followed by the contents. This in effect forms a descriptor that must be emitted before the actual contents. Sections described this way must not overlap.

For areas that don't have such descriptors, synthetic sections are formed by BFD. Consecutive contents in the two memory areas '0x0000...00' to '0x01ff...ff' and '0x2000...00' to '0x20ff...ff' are entered in sections named .text and .data respectively. If an area is not otherwise described, but would together with a neighboring lower area be less than '0x4000000' bytes long, it is joined with the lower area and the gap is zero-filled. For other cases, a new section is formed, named .MMIX.sec.n. Here, n is a number, a running count through the mmo file, starting at 0.

A loadable section specified as:

```
.section secname,"ax"
TETRA 1,2,3,4,-1,-2009
BYTE 80
```

and linked to address '0x4', is represented by the sequence:

```
0x98080050 - lop_spec 80
0x00000002 - two 32-bit words for the section name
0x7365636e - "secn"
0x616d6500 - "ame \0"
0x00000033 - flags CODE, READONLY, LOAD, ALLOC
0x00000000 - high 32 bits of section length
0x0000001c - section length is 28 bytes; 6 * 4 + 1 + alignment to 32 bits
0x00000000 - high 32 bits of section address
0x00000004 - section address is 4
0x98010002 - 64 bits with address of following data
0x00000000 - high 32 bits of address
0x00000004 - low 32 bits: data starts at address 4
0x00000001 - 1
0x00000002 - 2
0x00000003 - 3
0x00000004 - 4
Oxfffffff - -1
0xfffff827 - -2009
0x50000000 - 80 as a byte, padded with zeros.
```

Note that the lop\_spec wrapping does not include the section contents. Compare this to a non-loaded section specified as:

```
.section thirdsec
TETRA 200001,100002
BYTE 38,40
```

This, when linked to address '0x2000000000001c', is represented by:

```
0x98080050 - lop_spec 80
0x00000002 - two 32-bit words for the section name
0x7365636e - "thir"
0x616d6500 - "dsec"
0x00000010 - flag READONLY
0x00000000 - high 32 bits of section length
0x0000000c - section length is 12 bytes; 2 * 4 + 2 + alignment to 32 bits
0x20000000 - high 32 bits of address
0x0000001c - low 32 bits of address
0x0000001c - low 32 bits of address 0x2000000000001c
0x00030d41 - 200001
0x000186a2 - 100002
0x26280000 - 38, 40 as bytes, padded with zeros
```

For the latter example, the section contents must not be loaded in memory, and is therefore specified as part of the special data. The address is usually unimportant but might provide information for e.g. the DWARF 2 debugging format.

```
Version 1.3, 3 November 2008
```

```
Copyright © 2000, 2001, 2002, 2007, 2008 Free Software Foundation, Inc. http://fsf.org/
```

Everyone is permitted to copy and distribute verbatim copies of this license document, but changing it is not allowed.

#### 0. PREAMBLE

The purpose of this License is to make a manual, textbook, or other functional and useful document free in the sense of freedom: to assure everyone the effective freedom to copy and redistribute it, with or without modifying it, either commercially or non-commercially. Secondarily, this License preserves for the author and publisher a way to get credit for their work, while not being considered responsible for modifications made by others.

This License is a kind of "copyleft", which means that derivative works of the document must themselves be free in the same sense. It complements the GNU General Public License, which is a copyleft license designed for free software.

We have designed this License in order to use it for manuals for free software, because free software needs free documentation: a free program should come with manuals providing the same freedoms that the software does. But this License is not limited to software manuals; it can be used for any textual work, regardless of subject matter or whether it is published as a printed book. We recommend this License principally for works whose purpose is instruction or reference.

#### 1. APPLICABILITY AND DEFINITIONS

This License applies to any manual or other work, in any medium, that contains a notice placed by the copyright holder saying it can be distributed under the terms of this License. Such a notice grants a world-wide, royalty-free license, unlimited in duration, to use that work under the conditions stated herein. The "Document", below, refers to any such manual or work. Any member of the public is a licensee, and is addressed as "you". You accept the license if you copy, modify or distribute the work in a way requiring permission under copyright law.

A "Modified Version" of the Document means any work containing the Document or a portion of it, either copied verbatim, or with modifications and/or translated into another language.

A "Secondary Section" is a named appendix or a front-matter section of the Document that deals exclusively with the relationship of the publishers or authors of the Document to the Document's overall subject (or to related matters) and contains nothing that could fall directly within that overall subject. (Thus, if the Document is in part a textbook of mathematics, a Secondary Section may not explain any mathematics.) The relationship could be a matter of historical connection with the subject or with related matters, or of legal, commercial, philosophical, ethical or political position regarding them.

The "Invariant Sections" are certain Secondary Sections whose titles are designated, as being those of Invariant Sections, in the notice that says that the Document is released under this License. If a section does not fit the above definition of Secondary then it is not allowed to be designated as Invariant. The Document may contain zero Invariant Sections. If the Document does not identify any Invariant Sections then there are none.

The "Cover Texts" are certain short passages of text that are listed, as Front-Cover Texts or Back-Cover Texts, in the notice that says that the Document is released under

this License. A Front-Cover Text may be at most 5 words, and a Back-Cover Text may be at most 25 words.

A "Transparent" copy of the Document means a machine-readable copy, represented in a format whose specification is available to the general public, that is suitable for revising the document straightforwardly with generic text editors or (for images composed of pixels) generic paint programs or (for drawings) some widely available drawing editor, and that is suitable for input to text formatters or for automatic translation to a variety of formats suitable for input to text formatters. A copy made in an otherwise Transparent file format whose markup, or absence of markup, has been arranged to thwart or discourage subsequent modification by readers is not Transparent. An image format is not Transparent if used for any substantial amount of text. A copy that is not "Transparent" is called "Opaque".

Examples of suitable formats for Transparent copies include plain ASCII without markup, Texinfo input format, LaT<sub>E</sub>X input format, SGML or XML using a publicly available DTD, and standard-conforming simple HTML, PostScript or PDF designed for human modification. Examples of transparent image formats include PNG, XCF and JPG. Opaque formats include proprietary formats that can be read and edited only by proprietary word processors, SGML or XML for which the DTD and/or processing tools are not generally available, and the machine-generated HTML, PostScript or PDF produced by some word processors for output purposes only.

The "Title Page" means, for a printed book, the title page itself, plus such following pages as are needed to hold, legibly, the material this License requires to appear in the title page. For works in formats which do not have any title page as such, "Title Page" means the text near the most prominent appearance of the work's title, preceding the beginning of the body of the text.

The "publisher" means any person or entity that distributes copies of the Document to the public.

A section "Entitled XYZ" means a named subunit of the Document whose title either is precisely XYZ or contains XYZ in parentheses following text that translates XYZ in another language. (Here XYZ stands for a specific section name mentioned below, such as "Acknowledgements", "Dedications", "Endorsements", or "History".) To "Preserve the Title" of such a section when you modify the Document means that it remains a section "Entitled XYZ" according to this definition.

The Document may include Warranty Disclaimers next to the notice which states that this License applies to the Document. These Warranty Disclaimers are considered to be included by reference in this License, but only as regards disclaiming warranties: any other implication that these Warranty Disclaimers may have is void and has no effect on the meaning of this License.

## 2. VERBATIM COPYING

You may copy and distribute the Document in any medium, either commercially or noncommercially, provided that this License, the copyright notices, and the license notice saying this License applies to the Document are reproduced in all copies, and that you add no other conditions whatsoever to those of this License. You may not use technical measures to obstruct or control the reading or further copying of the copies you make or distribute. However, you may accept compensation in exchange for copies.

If you distribute a large enough number of copies you must also follow the conditions in section 3.

You may also lend copies, under the same conditions stated above, and you may publicly display copies.

### 3. COPYING IN QUANTITY

If you publish printed copies (or copies in media that commonly have printed covers) of the Document, numbering more than 100, and the Document's license notice requires Cover Texts, you must enclose the copies in covers that carry, clearly and legibly, all these Cover Texts: Front-Cover Texts on the front cover, and Back-Cover Texts on the back cover. Both covers must also clearly and legibly identify you as the publisher of these copies. The front cover must present the full title with all words of the title equally prominent and visible. You may add other material on the covers in addition. Copying with changes limited to the covers, as long as they preserve the title of the Document and satisfy these conditions, can be treated as verbatim copying in other respects.

If the required texts for either cover are too voluminous to fit legibly, you should put the first ones listed (as many as fit reasonably) on the actual cover, and continue the rest onto adjacent pages.

If you publish or distribute Opaque copies of the Document numbering more than 100, you must either include a machine-readable Transparent copy along with each Opaque copy, or state in or with each Opaque copy a computer-network location from which the general network-using public has access to download using public-standard network protocols a complete Transparent copy of the Document, free of added material. If you use the latter option, you must take reasonably prudent steps, when you begin distribution of Opaque copies in quantity, to ensure that this Transparent copy will remain thus accessible at the stated location until at least one year after the last time you distribute an Opaque copy (directly or through your agents or retailers) of that edition to the public.

It is requested, but not required, that you contact the authors of the Document well before redistributing any large number of copies, to give them a chance to provide you with an updated version of the Document.

#### 4. MODIFICATIONS

You may copy and distribute a Modified Version of the Document under the conditions of sections 2 and 3 above, provided that you release the Modified Version under precisely this License, with the Modified Version filling the role of the Document, thus licensing distribution and modification of the Modified Version to whoever possesses a copy of it. In addition, you must do these things in the Modified Version:

- A. Use in the Title Page (and on the covers, if any) a title distinct from that of the Document, and from those of previous versions (which should, if there were any, be listed in the History section of the Document). You may use the same title as a previous version if the original publisher of that version gives permission.
- B. List on the Title Page, as authors, one or more persons or entities responsible for authorship of the modifications in the Modified Version, together with at least five of the principal authors of the Document (all of its principal authors, if it has fewer than five), unless they release you from this requirement.

- C. State on the Title page the name of the publisher of the Modified Version, as the publisher.
- D. Preserve all the copyright notices of the Document.
- E. Add an appropriate copyright notice for your modifications adjacent to the other copyright notices.
- F. Include, immediately after the copyright notices, a license notice giving the public permission to use the Modified Version under the terms of this License, in the form shown in the Addendum below.
- G. Preserve in that license notice the full lists of Invariant Sections and required Cover Texts given in the Document's license notice.
- H. Include an unaltered copy of this License.
- I. Preserve the section Entitled "History", Preserve its Title, and add to it an item stating at least the title, year, new authors, and publisher of the Modified Version as given on the Title Page. If there is no section Entitled "History" in the Document, create one stating the title, year, authors, and publisher of the Document as given on its Title Page, then add an item describing the Modified Version as stated in the previous sentence.
- J. Preserve the network location, if any, given in the Document for public access to a Transparent copy of the Document, and likewise the network locations given in the Document for previous versions it was based on. These may be placed in the "History" section. You may omit a network location for a work that was published at least four years before the Document itself, or if the original publisher of the version it refers to gives permission.
- K. For any section Entitled "Acknowledgements" or "Dedications", Preserve the Title of the section, and preserve in the section all the substance and tone of each of the contributor acknowledgements and/or dedications given therein.
- L. Preserve all the Invariant Sections of the Document, unaltered in their text and in their titles. Section numbers or the equivalent are not considered part of the section titles.
- M. Delete any section Entitled "Endorsements". Such a section may not be included in the Modified Version.
- N. Do not retitle any existing section to be Entitled "Endorsements" or to conflict in title with any Invariant Section.
- O. Preserve any Warranty Disclaimers.

If the Modified Version includes new front-matter sections or appendices that qualify as Secondary Sections and contain no material copied from the Document, you may at your option designate some or all of these sections as invariant. To do this, add their titles to the list of Invariant Sections in the Modified Version's license notice. These titles must be distinct from any other section titles.

You may add a section Entitled "Endorsements", provided it contains nothing but endorsements of your Modified Version by various parties—for example, statements of peer review or that the text has been approved by an organization as the authoritative definition of a standard.

You may add a passage of up to five words as a Front-Cover Text, and a passage of up to 25 words as a Back-Cover Text, to the end of the list of Cover Texts in the Modified Version. Only one passage of Front-Cover Text and one of Back-Cover Text may be added by (or through arrangements made by) any one entity. If the Document already includes a cover text for the same cover, previously added by you or by arrangement made by the same entity you are acting on behalf of, you may not add another; but you may replace the old one, on explicit permission from the previous publisher that added the old one.

The author(s) and publisher(s) of the Document do not by this License give permission to use their names for publicity for or to assert or imply endorsement of any Modified Version.

### 5. COMBINING DOCUMENTS

You may combine the Document with other documents released under this License, under the terms defined in section 4 above for modified versions, provided that you include in the combination all of the Invariant Sections of all of the original documents, unmodified, and list them all as Invariant Sections of your combined work in its license notice, and that you preserve all their Warranty Disclaimers.

The combined work need only contain one copy of this License, and multiple identical Invariant Sections may be replaced with a single copy. If there are multiple Invariant Sections with the same name but different contents, make the title of each such section unique by adding at the end of it, in parentheses, the name of the original author or publisher of that section if known, or else a unique number. Make the same adjustment to the section titles in the list of Invariant Sections in the license notice of the combined work.

In the combination, you must combine any sections Entitled "History" in the various original documents, forming one section Entitled "History"; likewise combine any sections Entitled "Acknowledgements", and any sections Entitled "Dedications". You must delete all sections Entitled "Endorsements."

## 6. COLLECTIONS OF DOCUMENTS

You may make a collection consisting of the Document and other documents released under this License, and replace the individual copies of this License in the various documents with a single copy that is included in the collection, provided that you follow the rules of this License for verbatim copying of each of the documents in all other respects.

You may extract a single document from such a collection, and distribute it individually under this License, provided you insert a copy of this License into the extracted document, and follow this License in all other respects regarding verbatim copying of that document.

## 7. AGGREGATION WITH INDEPENDENT WORKS

A compilation of the Document or its derivatives with other separate and independent documents or works, in or on a volume of a storage or distribution medium, is called an "aggregate" if the copyright resulting from the compilation is not used to limit the legal rights of the compilation's users beyond what the individual works permit. When the Document is included in an aggregate, this License does not apply to the other works in the aggregate which are not themselves derivative works of the Document.

If the Cover Text requirement of section 3 is applicable to these copies of the Document, then if the Document is less than one half of the entire aggregate, the Document's Cover Texts may be placed on covers that bracket the Document within the aggregate, or the electronic equivalent of covers if the Document is in electronic form. Otherwise they must appear on printed covers that bracket the whole aggregate.

#### 8. TRANSLATION

Translation is considered a kind of modification, so you may distribute translations of the Document under the terms of section 4. Replacing Invariant Sections with translations requires special permission from their copyright holders, but you may include translations of some or all Invariant Sections in addition to the original versions of these Invariant Sections. You may include a translation of this License, and all the license notices in the Document, and any Warranty Disclaimers, provided that you also include the original English version of this License and the original versions of those notices and disclaimers. In case of a disagreement between the translation and the original version of this License or a notice or disclaimer, the original version will prevail.

If a section in the Document is Entitled "Acknowledgements", "Dedications", or "History", the requirement (section 4) to Preserve its Title (section 1) will typically require changing the actual title.

## 9. TERMINATION

You may not copy, modify, sublicense, or distribute the Document except as expressly provided under this License. Any attempt otherwise to copy, modify, sublicense, or distribute it is void, and will automatically terminate your rights under this License.

However, if you cease all violation of this License, then your license from a particular copyright holder is reinstated (a) provisionally, unless and until the copyright holder explicitly and finally terminates your license, and (b) permanently, if the copyright holder fails to notify you of the violation by some reasonable means prior to 60 days after the cessation.

Moreover, your license from a particular copyright holder is reinstated permanently if the copyright holder notifies you of the violation by some reasonable means, this is the first time you have received notice of violation of this License (for any work) from that copyright holder, and you cure the violation prior to 30 days after your receipt of the notice.

Termination of your rights under this section does not terminate the licenses of parties who have received copies or rights from you under this License. If your rights have been terminated and not permanently reinstated, receipt of a copy of some or all of the same material does not give you any rights to use it.

### 10. FUTURE REVISIONS OF THIS LICENSE

The Free Software Foundation may publish new, revised versions of the GNU Free Documentation License from time to time. Such new versions will be similar in spirit to the present version, but may differ in detail to address new problems or concerns. See http://www.gnu.org/copyleft/.

Each version of the License is given a distinguishing version number. If the Document specifies that a particular numbered version of this License "or any later version" applies to it, you have the option of following the terms and conditions either of that

specified version or of any later version that has been published (not as a draft) by the Free Software Foundation. If the Document does not specify a version number of this License, you may choose any version ever published (not as a draft) by the Free Software Foundation. If the Document specifies that a proxy can decide which future versions of this License can be used, that proxy's public statement of acceptance of a version permanently authorizes you to choose that version for the Document.

#### 11. RELICENSING

"Massive Multiauthor Collaboration Site" (or "MMC Site") means any World Wide Web server that publishes copyrightable works and also provides prominent facilities for anybody to edit those works. A public wiki that anybody can edit is an example of such a server. A "Massive Multiauthor Collaboration" (or "MMC") contained in the site means any set of copyrightable works thus published on the MMC site.

"CC-BY-SA" means the Creative Commons Attribution-Share Alike 3.0 license published by Creative Commons Corporation, a not-for-profit corporation with a principal place of business in San Francisco, California, as well as future copyleft versions of that license published by that same organization.

"Incorporate" means to publish or republish a Document, in whole or in part, as part of another Document.

An MMC is "eligible for relicensing" if it is licensed under this License, and if all works that were first published under this License somewhere other than this MMC, and subsequently incorporated in whole or in part into the MMC, (1) had no cover texts or invariant sections, and (2) were thus incorporated prior to November 1, 2008.

The operator of an MMC Site may republish an MMC contained in the site under CC-BY-SA on the same site at any time before August 1, 2009, provided the MMC is eligible for relicensing.

# ADDENDUM: How to use this License for your documents

To use this License in a document you have written, include a copy of the License in the document and put the following copyright and license notices just after the title page:

Copyright (C) year your name. Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.3 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled ''GNU Free Documentation License''.

If you have Invariant Sections, Front-Cover Texts and Back-Cover Texts, replace the "with...Texts." line with this:

with the Invariant Sections being  $list\ their\ titles$ , with the Front-Cover Texts being list, and with the Back-Cover Texts being list.

If you have Invariant Sections without Cover Texts, or some other combination of the three, merge those two alternatives to suit the situation.

If your document contains nontrivial examples of program code, we recommend releasing these examples in parallel under your choice of free software license, such as the GNU General Public License, to permit their use in free software.

_		bfd_cache_close	166
_bfd_final_link_relocate	70	bfd_cache_close_all	166
_bfd_generic_link_add_archive_symbols 16	9	bfd_cache_init	
_bfd_generic_link_add_one_symbol 16		bfd_calc_gnu_debuglink_crc32	160
_bfd_generic_make_empty_symbol		bfd_canonicalize_reloc	14
_bfd_link_add_symbols in target vector 16	7	${\it bfd\_canonicalize\_symtab}$	
_bfd_link_final_link in target vector 16		bfd_check_format	51
_bfd_link_hash_table_create in target vector 16		bfd_check_format_matches	51
_bfd_relocate_contents		bfd_check_overflow	. 59
_bid_relocate_contents	0	bfd_close	158
		bfd_close_all_done	158
$\mathbf{A}$		bfd_coff_backend_data	
		bfd_copy_private_bfd_data	
aout_size_machine_type	8	bfd_copy_private_header_data	
aout_size_mkobject	8	bfd_copy_private_section_data	
aout_size_new_section_hook	8	bfd_copy_private_symbol_data	
aout_size_set_arch_mach	8	bfd_core_file_failing_command	
aout_size_some_aout_object_p	8	bfd_core_file_failing_signal	
aout_size_swap_exec_header_in	(	bfd_core_file_pid	
aout_size_swap_exec_header_out	1	bfd_create	
arelent_chain		bfd_create_gnu_debuglink_section	
		bfd_decode_symclass	
В		bfd_default_arch_struct	
Б		bfd_default_compatible	
BFD	1	bfd_default_reloc_type_lookup	
BFD canonical format		bfd_default_scan	
bfd_alloc		bfd_default_set_arch_mach	
bfd_alloc2		bfd_demangle	
$bfd\_alt\_mach\_code$		bfd_emul_get_commonpagesize	
bfd_arch_bits_per_address		bfd_emul_get_maxpagesize	
bfd_arch_bits_per_byte		bfd_emul_set_commonpagesizebfd_emul_set_commonpagesize	
bfd_arch_default_fill		bfd_emul_set_maxpagesize	
bfd_arch_get_compatible		bfd_errmsg	
bfd_arch_list		bfd_fdopenr	
bfd_arch_mach_octets_per_byte 15		bfd_fill_in_gnu_debuglink_section	
BFD_ARELOC_BFIN_ADD 8	_	bfd_find_target	
BFD_ARELOC_BFIN_ADDR8	_	bfd_find_version_for_sym	
BFD_ARELOC_BFIN_AND 8	_	bfd_follow_gnu_debugaltlink	
BFD_ARELOC_BFIN_COMP8		bfd_follow_gnu_debuglink	
BFD_ARELOC_BFIN_CONST		bfd_fopen	
BFD_ARELOC_BFIN_DIV		bfd_format_string	
BFD_ARELOC_BFIN_HWPAGE		bfd_generic_define_common_symbol	
BFD_ARELOC_BFIN_LAND		bfd_generic_discard_group	
BFD_ARELOC_BFIN_LEN		bfd_generic_gc_sections	
BFD_ARELOC_BFIN_LOR		bfd_generic_get_relocated_section_contents	
BFD_ARELOC_BFIN_LSHIFT 8		bid_generic_is_group_sectionbid_generic_is_group_section	
BFD_ARELOC_BFIN_MOD		bid_generic_lookup_section_flagsbid_generic_lookup_section_flags	
BFD_ARELOC_BFIN_MULT		bid_generic_nerge_sectionsbid_generic_merge_sections	
BFD_ARELOC_BFIN_NEG		bid_generic_merge_sectionsbid_generic_relax_section	
BFD_ARELOC_BFIN_OR		bid_get_alt_debug_link_info	
BFD_ARELOC_BFIN_PAGE			
BFD_ARELOC_BFIN_PUSH		bfd_get_archbfd_get_arch_info	
BFD_ARELOC_BFIN_RSHIFT8		bid_get_arch_size	
BFD_ARELOC_BFIN_SUB		bfd_get_arcn_size bfd_get_assert_handler	
BFD_ARELOC_BFIN_XOR 8	<b>.</b>	bfd_get_debug_link_info	100

bfd_get_error	bfd_perform_relocation	59
bfd_get_error_handler	bfd_perror	
bfd_get_gp_size	bfd_print_symbol_vandf	
bfd_get_linker_section	bfd_printable_arch_mach 1	
bfd_get_mach	bfd_printable_name 1	
bfd_get_mtime	bfd_put_size	
bfd_get_next_mapent	BFD_RELOC_12_PCREL	
bfd_get_next_section_by_name	BFD_RELOC_14	
bfd_get_reloc_code_name	BFD_RELOC_16	
bfd_get_reloc_size	BFD_RELOC_16_BASEREL	
bfd_get_reloc_upper_bound	BFD_RELOC_16_GOT_PCREL	
bfd_get_section_by_name	BFD_RELOC_16_GOTOFF	
bfd_get_section_by_name_if	BFD_RELOC_16_PCREL	
bfd_get_section_contents	BFD_RELOC_16_PCREL_S2	
bfd_get_sign_extend_vma	BFD_RELOC_16_PLT_PCREL	
bfd_get_size	BFD_RELOC_16_PLTOFF	
bfd_get_symtab_upper_bound	BFD_RELOC_16C_ABS20	
bfd_get_target_info	BFD_RELOC_16C_ABS2O_C	
bfd_get_unique_section_name	BFD_RELOC_16C_ABS24	
bfd_h_put_size	BFD_RELOC_16C_ABS24_C	
bfd_hash_allocate	BFD_RELOC_16C_DISP04	
bfd_hash_lookup	BFD_RELOC_16C_DISPO4_C	
bfd_hash_newfunc	BFD_RELOC_16C_DISP08	
bfd_hash_set_default_size	BFD_RELOC_16C_DISPO8_C	
bfd_hash_table_free	BFD_RELOC_16C_DISP16	
bfd_hash_table_init	BFD_RELOC_16C_DISP16_C	
bfd_hash_table_init_n	BFD_RELOC_16C_DISP24	
bfd_hash_traverse	BFD_RELOC_16C_DISP24_C	
bfd_hide_sym_by_version	BFD_RELOC_16C_DISP24a	
bfd_init	BFD_RELOC_16C_DISP24a_C	
bfd_install_relocation	BFD_RELOC_16C_IMMO4	
bfd_is_local_label	BFD_RELOC_16C_IMMO4_C	
bfd_is_local_label_name	BFD_RELOC_16C_IMM16	
bfd_is_target_special_symbol	BFD_RELOC_16C_IMM16_C	
bfd_is_undefined_symclass	BFD_RELOC_16C_IMM20	
bfd_link_split_section	BFD_RELOC_16C_IMM20_C	
bfd_log2	BFD_RELOC_16C_IMM24	
bfd_lookup_arch	BFD_RELOC_16C_IMM24_C	
bfd_make_debug_symbol	BFD_RELOC_16C_IMM32	
bfd_make_empty_symbol	BFD_RELOC_16C_IMM32_C	
bfd_make_readable	BFD_RELOC_16C_NUMO8	
bfd_make_section	BFD_RELOC_16C_NUMO8_C	
bfd_make_section_anyway	BFD_RELOC_16C_NUM16	
bfd_make_section_anyway_with_flags	BFD_RELOC_16C_NUM16_C	
bfd_make_section_old_way	BFD_RELOC_16C_NUM32	
bfd_make_section_with_flags	BFD_RELOC_16C_NUM32_C	
bfd_make_writable	BFD_RELOC_16C_REG04	
bfd_malloc_and_get_section	BFD_RELOC_16C_REG04_C	
bfd_map_over_sections	BFD_RELOC_16C_REG04a	
bfd_merge_private_bfd_data	BFD_RELOC_16C_REG04a_C	
bfd_mmap	BFD_RELOC_16C_REG14	
bfd_octets_per_byte	BFD_RELOC_16C_REG14_C	
bfd_open_file	BFD_RELOC_16C_REG14_C	
bfd_openr	BFD_RELOC_16C_REG16_C	
bid_openr_iovec	BFD_RELOC_16C_REG10_C	
bfd_openr_next_archived_file	BFD_RELOC_16C_REG20_C	
bfd_openstreamr	BFD_RELOC_23_PCREL_S2	
bfd_openw	BFD RELOC 24	
DIG-05011W 100	∪∪_∟∪∪∪∪∪∪∪∪∪∟	UU

BFD_RELOC_24_PCREL	BFD_RELOC_390_PC24DBL	
BFD_RELOC_24_PLT_PCREL	BFD_RELOC_390_PC32DBL	
BFD_RELOC_26	BFD_RELOC_390_PLT12DBL	
BFD_RELOC_32	BFD_RELOC_390_PLT16DBL	
BFD_RELOC_32_BASEREL	BFD_RELOC_390_PLT24DBL	
BFD_RELOC_32_GOT_PCREL	BFD_RELOC_390_PLT32	
BFD_RELOC_32_GOTOFF	BFD_RELOC_390_PLT32DBL	
BFD_RELOC_32_PCREL	BFD_RELOC_390_PLT64	
BFD_RELOC_32_PCREL_S2	BFD_RELOC_390_PLTOFF16BFD_RELOC_390_PLTOFF32	
BFD_RELOC_32_PLT_PCREL	BFD_RELOC_390_PLT0FF64	
BFD_RELOC_32_PLT0FF	BFD_RELOC_390_RELATIVE	
BFD_RELOC_386_COPY	BFD_RELOC_390_TLS_DTPMOD	
BFD_RELOC_386_GLOB_DAT	BFD_RELOC_390_TLS_DTPNOD	
BFD_RELOC_386_GOT32	BFD_RELOC_390_TLS_DIPOFF	
BFD_RELOC_386_GOTOFF	BFD_RELOC_390_TLS_GD64	
BFD_RELOC_386_GOTPC	BFD_RELOC_390_TLS_GDCALL	
BFD_RELOC_386_IRELATIVE	BFD_RELOC_390_TLS_GOTIE12	
BFD_RELOC_386_JUMP_SLOT	BFD_RELOC_390_TLS_GOTIE20	
BFD_RELOC_386_PLT32	BFD_RELOC_390_TLS_GOTIE32	
BFD_RELOC_386_RELATIVE	BFD_RELOC_390_TLS_GOTIE64	
BFD_RELOC_386_TLS_DESC	BFD_RELOC_390_TLS_IE32	
BFD_RELOC_386_TLS_DESC_CALL	BFD_RELOC_390_TLS_IE64	
BFD_RELOC_386_TLS_DTPMOD32	BFD_RELOC_390_TLS_IEENT	
BFD_RELOC_386_TLS_DTP0FF32	BFD_RELOC_390_TLS_LDCALL	
BFD_RELOC_386_TLS_GD	BFD_RELOC_390_TLS_LDM32	
BFD_RELOC_386_TLS_GOTDESC	BFD_RELOC_390_TLS_LDM64	
BFD_RELOC_386_TLS_GOTIE	BFD_RELOC_390_TLS_LD032	
BFD_RELOC_386_TLS_IE	BFD_RELOC_390_TLS_LD064	
BFD_RELOC_386_TLS_IE_32	BFD_RELOC_390_TLS_LE32	
BFD_RELOC_386_TLS_LDM	BFD_RELOC_390_TLS_LE64	
BFD_RELOC_386_TLS_LD0_32	BFD_RELOC_390_TLS_LOAD	
BFD_RELOC_386_TLS_LE	BFD_RELOC_390_TLS_TPOFF	
BFD_RELOC_386_TLS_LE_32	BFD_RELOC_64	
BFD_RELOC_386_TLS_TPOFF.	BFD_RELOC_64_PCREL	
BFD_RELOC_386_TLS_TP0FF32	BFD_RELOC_64_PLT_PCREL	
BFD_RELOC_390_12	BFD_RELOC_64_PLTOFF	
BFD_RELOC_390_20	BFD_RELOC_68K_GLOB_DAT	
BFD_RELOC_390_COPY	BFD_RELOC_68K_JMP_SLOT	
BFD_RELOC_390_GLOB_DAT	BFD_RELOC_68K_RELATIVE	
BFD_RELOC_390_GOT12	BFD_RELOC_68K_TLS_GD16	
BFD_RELOC_390_GOT16	BFD_RELOC_68K_TLS_GD32	
BFD_RELOC_390_GOT20	BFD_RELOC_68K_TLS_GD8	
BFD_RELOC_390_GOT64	BFD_RELOC_68K_TLS_IE16	
BFD_RELOC_390_GOTENT	BFD_RELOC_68K_TLS_IE32	
BFD_RELOC_390_GOTOFF64	BFD_RELOC_68K_TLS_IE8	
BFD_RELOC_390_GOTPC	BFD_RELOC_68K_TLS_LDM16	
BFD_RELOC_390_GOTPCDBL	BFD_RELOC_68K_TLS_LDM32	
BFD_RELOC_390_GOTPLT12.	BFD_RELOC_68K_TLS_LDM8	
BFD_RELOC_390_GOTPLT16.	BFD_RELOC_68K_TLS_LD016	
BFD_RELOC_390_GOTPLT20	BFD_RELOC_68K_TLS_LD032	
BFD_RELOC_390_GOTPLT32	BFD_RELOC_68K_TLS_LD08	
BFD_RELOC_390_GOTPLT64.	BFD_RELOC_68K_TLS_LE16	
BFD_RELOC_390_GOTPLTENT	BFD_RELOC_68K_TLS_LE32	
BFD_RELOC_390_IRELATIVE	BFD_RELOC_68K_TLS_LE8	
BFD_RELOC_390_JMP_SLOT.	BFD_RELOC_8	
BFD_RELOC_390_PC12DBL	BFD_RELOC_8_BASEREL	
BFD RELOC 390 PC16DBL	BFD_RELOC_8_FFnn	

BFD_RELOC_8_GOT_PCREL 60	BFD_RELOC_AARCH64_LD_L019_PCREL
BFD_RELOC_8_GOTOFF	BFD_RELOC_AARCH64_LD32_GOT_LO12_NC 120
BFD_RELOC_8_PCREL60	BFD_RELOC_AARCH64_LD64_GOT_LO12_NC 120
BFD_RELOC_8_PLT_PCREL61	BFD_RELOC_AARCH64_LDST_L012
BFD_RELOC_8_PLTOFF	BFD_RELOC_AARCH64_LDST128_L012 120
BFD_RELOC_860_COPY 110	BFD_RELOC_AARCH64_LDST16_L012120
BFD_RELOC_860_GLOB_DAT	BFD_RELOC_AARCH64_LDST32_L012
BFD_RELOC_860_HAGOT	BFD_RELOC_AARCH64_LDST64_L012
BFD_RELOC_860_HAGOTOFF	BFD_RELOC_AARCH64_LDST8_LO12
BFD_RELOC_860_HAPC 110	BFD_RELOC_AARCH64_MOVW_GO
BFD_RELOC_860_HIGH 110	BFD_RELOC_AARCH64_MOVW_GO_NC
BFD_RELOC_860_HIGHADJ110	BFD_RELOC_AARCH64_MOVW_GO_S
BFD_RELOC_860_HIGOT	BFD_RELOC_AARCH64_MOVW_G1
BFD_RELOC_860_HIGOTOFF110	BFD_RELOC_AARCH64_MOVW_G1_NC
BFD_RELOC_860_JUMP_SLOT 110	BFD_RELOC_AARCH64_MOVW_G1_S
BFD_RELOC_860_LOGOTO110	BFD_RELOC_AARCH64_MOVW_G2
BFD_RELOC_860_LOGOT1110	
BFD_RELOC_860_LOGOTOFF0	BFD_RELOC_AARCH64_MOVW_G2_NC
BFD_RELOC_860_LOGOTOFF1 110	BFD_RELOC_AARCH64_MOVW_G2_S
BFD_RELOC_860_LOGOTOFF2 110	BFD_RELOC_AARCH64_MOVW_G3
BFD_RELOC_860_LOGOTOFF3 110	BFD_RELOC_AARCH64_NONE
BFD_RELOC_860_LOPC	BFD_RELOC_AARCH64_RELATIVE
BFD_RELOC_860_LOW0	BFD_RELOC_AARCH64_RELOC_END
BFD_RELOC_860_LOW1	BFD_RELOC_AARCH64_RELOC_START
BFD_RELOC_860_LOW2	BFD_RELOC_AARCH64_TLS_DTPMOD
BFD_RELOC_860_LOW3	BFD_RELOC_AARCH64_TLS_DTPREL
BFD_RELOC_860_PC16	BFD_RELOC_AARCH64_TLS_TPREL 122
BFD_RELOC_860_PC26	BFD_RELOC_AARCH64_TLSDESC 122
BFD_RELOC_860_PLT26	BFD_RELOC_AARCH64_TLSDESC_ADD
BFD_RELOC_860_RELATIVE	BFD_RELOC_AARCH64_TLSDESC_ADD_L012_NC 122
BFD_RELOC_860_SPGOT0	BFD_RELOC_AARCH64_TLSDESC_ADR_PAGE21 122
BFD_RELOC_860_SPGOT1110	BFD_RELOC_AARCH64_TLSDESC_ADR_PREL21 122
BFD_RELOC_860_SPGOTOFF0	BFD_RELOC_AARCH64_TLSDESC_CALL 122
BFD_RELOC_860_SPGOTOFF1 110	BFD_RELOC_AARCH64_TLSDESC_LD_L012_NC 123
BFD_RELOC_860_SPLIT0110	BFD_RELOC_AARCH64_TLSDESC_LD_PREL19 121
BFD_RELOC_860_SPLIT1110	BFD_RELOC_AARCH64_TLSDESC_LD32_LO12_NC 122
BFD_RELOC_860_SPLIT2110	BFD_RELOC_AARCH64_TLSDESC_LD64_LO12_NC 122
BFD_RELOC_AARCH64_16	BFD_RELOC_AARCH64_TLSDESC_LDR 122
BFD_RELOC_AARCH64_16_PCREL	BFD_RELOC_AARCH64_TLSDESC_OFF_GO_NC 122
BFD_RELOC_AARCH64_32	BFD_RELOC_AARCH64_TLSDESC_OFF_G1 122
BFD_RELOC_AARCH64_32_PCREL 118	BFD_RELOC_AARCH64_TLSGD_ADD_L012_NC 121
BFD_RELOC_AARCH64_64	BFD_RELOC_AARCH64_TLSGD_ADR_PAGE21 121
BFD_RELOC_AARCH64_64_PCREL	BFD_RELOC_AARCH64_TLSIE_ADR_GOTTPREL_PAGE21
BFD_RELOC_AARCH64_ADD_LO12	
BFD_RELOC_AARCH64_ADR_GOT_PAGE 120	BFD_RELOC_AARCH64_TLSIE_LD_GOTTPREL_L012_NC
BFD_RELOC_AARCH64_ADR_HI21_NC_PCREL 119	
BFD_RELOC_AARCH64_ADR_HI21_PCREL	BFD_RELOC_AARCH64_TLSIE_LD_GOTTPREL_PREL19
BFD_RELOC_AARCH64_ADR_LO21_PCREL	
BFD_RELOC_AARCH64_BRANCH19	BFD_RELOC_AARCH64_TLSIE_LD32_GOTTPREL_L012_
BFD_RELOC_AARCH64_CALL26	NC
BFD_RELOC_AARCH64_COPY	BFD_RELOC_AARCH64_TLSIE_LD64_GOTTPREL_L012_
BFD_RELOC_AARCH64_GAS_INTERNAL_FIXUP 123	NC
BFD_RELOC_AARCH64_GLOB_DAT	BFD_RELOC_AARCH64_TLSIE_MOVW_GOTTPREL_GO_NC
BFD_RELOC_AARCH64_GOT_LD_PREL19	
BFD_RELOC_AARCH64_IRELATIVE	BFD_RELOC_AARCH64_TLSIE_MOVW_GOTTPREL_G1
BFD_RELOC_AARCH64_JUMP_SLOT	
BFD_RELOC_AARCH64_JUMP26	BFD_RELOC_AARCH64_TLSLE_ADD_TPREL_HI12 121
BFD RELOC AARCH64 LD GOT LO12 NC 123	BFD RELOC AARCH64 TLSLE ADD TPREL L012 121

	BFD_RELOC_ARM_GOTOFF	77
21	BFD_RELOC_ARM_GOTPC	77
21		
21		
	BFD_RELOC_ARM_LDR_SB_G2	78
	BFD_RELOC_ARM_LDRS_PC_GO	77
65	BFD_RELOC_ARM_LDRS_PC_G1	77
64	BFD_RELOC_ARM_LDRS_PC_G2	77
65	BFD_RELOC_ARM_LDRS_SB_GO	78
65	BFD_RELOC_ARM_LDRS_SB_G1	78
	RED REIOC ARM MILITI	78
78		
78		
	BFD_RELOC_ARM_SMC	78
	BFD_RELOC_ARM_SWI	
	BFD_RELOC_ARM_T32_ADD_IMM	78
78		
78	BFD_RELOC_ARM_T32_CP_OFF_IMM	78
78	BFD_RELOC_ARM_T32_CP_OFF_IMM_S2	
78		
77		77
	64 65 65 65 66 66 66 66 66 66 66 66 66 66	BFD_RELOC_ARM_HVC.

BFD_RELOC_ARM_THM_TLS_DESCSEQ	77	BFD_RELOC_BFIN_16_HIGH	81
BFD_RELOC_ARM_THUMB_ADD	78	BFD_RELOC_BFIN_16_IMM	81
BFD_RELOC_ARM_THUMB_IMM	78	BFD_RELOC_BFIN_16_LOW	81
BFD_RELOC_ARM_THUMB_MOVT	77	BFD_RELOC_BFIN_24_PCREL_CALL_X	81
BFD_RELOC_ARM_THUMB_MOVT_PCREL	77	BFD_RELOC_BFIN_24_PCREL_JUMP_L	81
BFD_RELOC_ARM_THUMB_MOVW		BFD_RELOC_BFIN_4_PCREL	
BFD_RELOC_ARM_THUMB_MOVW_PCREL	77	BFD_RELOC_BFIN_5_PCREL	
BFD_RELOC_ARM_THUMB_OFFSET		BFD_RELOC_BFIN_FUNCDESC	
BFD_RELOC_ARM_THUMB_SHIFT		BFD_RELOC_BFIN_FUNCDESC_GOT17M4	
BFD_RELOC_ARM_TLS_CALL		BFD_RELOC_BFIN_FUNCDESC_GOTHI	
BFD_RELOC_ARM_TLS_DESC		BFD_RELOC_BFIN_FUNCDESC_GOTLO	
BFD_RELOC_ARM_TLS_DESCSEQ		BFD_RELOC_BFIN_FUNCDESC_GOTOFF17M4	
BFD_RELOC_ARM_TLS_DTPMOD32		BFD_RELOC_BFIN_FUNCDESC_GOTOFFHI	
BFD_RELOC_ARM_TLS_DTPOFF32		BFD_RELOC_BFIN_FUNCDESC_GOTOFFLO	
BFD_RELOC_ARM_TLS_GD32		BFD_RELOC_BFIN_FUNCDESC_VALUE	
BFD_RELOC_ARM_TLS_GOTDESC		BFD_RELOC_BFIN_GOT	
BFD_RELOC_ARM_TLS_IE32		BFD_RELOC_BFIN_GOT17M4	
BFD_RELOC_ARM_TLS_LDM32		BFD_RELOC_BFIN_GOTHI	
BFD_RELOC_ARM_TLS_LD032		BFD_RELOC_BFIN_GOTLO	
BFD_RELOC_ARM_TLS_LE32		BFD_RELOC_BFIN_GOTOFF17M4	
BFD_RELOC_ARM_TLS_TP0FF32		BFD_RELOC_BFIN_GOTOFFHI	
BFD_RELOC_ARM_V4BX		BFD_RELOC_BFIN_GOTOFFLO	82
BFD_RELOC_AVR_13_PCREL		BFD_RELOC_BFIN_PLTPC	
BFD_RELOC_AVR_16_PM		BFD_RELOC_C6000_ABS_H16	
BFD_RELOC_AVR_6		BFD_RELOC_C6000_ABS_L16	
BFD_RELOC_AVR_6_ADIW		BFD_RELOC_C6000_ABS_S16	
BFD_RELOC_AVR_7_PCREL		BFD_RELOC_C6000_ALIGN	
BFD_RELOC_AVR_8_HI		BFD_RELOC_C6000_COPY	
BFD_RELOC_AVR_8_HLO		BFD_RELOC_C6000_DSBT_INDEX	
BFD_RELOC_AVR_8_LO		BFD_RELOC_C6000_EHTYPE	
BFD_RELOC_AVR_CALL		BFD_RELOC_C6000_FPHEAD	
BFD_RELOC_AVR_DIFF16		BFD_RELOC_C6000_JUMP_SLOT	
BFD_RELOC_AVR_DIFF32		BFD_RELOC_C6000_NOCMP	
BFD_RELOC_AVR_DIFF8		BFD_RELOC_C6000_PCR_H16	
BFD_RELOC_AVR_HH8_LDI		BFD_RELOC_C6000_PCR_L16	
BFD_RELOC_AVR_HH8_LDI_NEG		BFD_RELOC_C6000_PCR_S10	
BFD_RELOC_AVR_HH8_LDI_PM		BFD_RELOC_C6000_PCR_S12	
BFD_RELOC_AVR_HH8_LDI_PM_NEG		BFD_RELOC_C6000_PCR_S21	
BFD_RELOC_AVR_HI8_LDI		BFD_RELOC_C6000_PCR_S7	
BFD_RELOC_AVR_HI8_LDI_GS		BFD_RELOC_C6000_PREL31	
BFD_RELOC_AVR_HI8_LDI_NEG		BFD_RELOC_C6000_SBR_GOT_H16_W	
BFD_RELOC_AVR_HI8_LDI_PM	97	BFD_RELOC_C6000_SBR_GOT_L16_W	
BFD_RELOC_AVR_HI8_LDI_PM_NEG		BFD_RELOC_C6000_SBR_GOT_U15_W	
BFD_RELOC_AVR_LDI		BFD_RELOC_C6000_SBR_H16_B	
BFD_RELOC_AVR_LDS_STS_16	98	BFD_RELOC_C6000_SBR_H16_H	92
BFD_RELOC_AVR_LO8_LDI		BFD_RELOC_C6000_SBR_H16_W	92
BFD_RELOC_AVR_LO8_LDI_GS	96	BFD_RELOC_C6000_SBR_L16_B	
BFD_RELOC_AVR_LO8_LDI_NEG	96	BFD_RELOC_C6000_SBR_L16_H	
BFD_RELOC_AVR_LO8_LDI_PM	96	BFD_RELOC_C6000_SBR_L16_W	92
BFD_RELOC_AVR_LO8_LDI_PM_NEG	97	BFD_RELOC_C6000_SBR_S16	92
BFD_RELOC_AVR_MS8_LDI	96	BFD_RELOC_C6000_SBR_U15_B	
BFD_RELOC_AVR_MS8_LDI_NEG	96	BFD_RELOC_C6000_SBR_U15_H	92
BFD_RELOC_AVR_PORT5		BFD_RELOC_C6000_SBR_U15_W	
BFD_RELOC_AVR_PORT6	98	bfd_reloc_code_type	
BFD_RELOC_BFIN_10_PCREL		BFD_RELOC_CR16_ABS201	
BFD_RELOC_BFIN_11_PCREL		BFD_RELOC_CR16_ABS24	
BFD_RELOC_BFIN_12_PCREL_JUMP	81	BFD_RELOC_CR16_DISP161	
RED RELOC BEIN 12 POREL HIMP S			08

BFD_RELOC_CR16_DISP24		BFD_RELOC_CRIS_UNSIGNED_4	
BFD_RELOC_CR16_DISP24a		BFD_RELOC_CRIS_UNSIGNED_5	
BFD_RELOC_CR16_DISP4		BFD_RELOC_CRIS_UNSIGNED_6	
BFD_RELOC_CR16_DISP8		BFD_RELOC_CRIS_UNSIGNED_8	
BFD_RELOC_CR16_GLOB_DAT		BFD_RELOC_CRX_ABS16	
BFD_RELOC_CR16_GOT_REGREL20	108	BFD_RELOC_CRX_ABS32	
BFD_RELOC_CR16_GOTC_REGREL20	108	BFD_RELOC_CRX_IMM16	109
BFD_RELOC_CR16_IMM16	108	BFD_RELOC_CRX_IMM32	109
BFD_RELOC_CR16_IMM20	108	BFD_RELOC_CRX_NUM16	109
BFD_RELOC_CR16_IMM24	108	BFD_RELOC_CRX_NUM32	109
BFD_RELOC_CR16_IMM32	108	BFD_RELOC_CRX_NUM8	109
BFD_RELOC_CR16_IMM32a	108	BFD_RELOC_CRX_REGREL12	108
BFD_RELOC_CR16_IMM4	108	BFD_RELOC_CRX_REGREL22	108
BFD_RELOC_CR16_IMM8	108	BFD_RELOC_CRX_REGREL28	108
BFD_RELOC_CR16_NUM16	108	BFD_RELOC_CRX_REGREL32	108
BFD_RELOC_CR16_NUM32		BFD_RELOC_CRX_REL16	
BFD_RELOC_CR16_NUM32a		BFD_RELOC_CRX_REL24	
BFD_RELOC_CR16_NUM8		BFD_RELOC_CRX_REL32	
BFD_RELOC_CR16_REGRELO		BFD_RELOC_CRX_REL4	
BFD_RELOC_CR16_REGREL14		BFD_RELOC_CRX_REL8	
BFD_RELOC_CR16_REGREL14a		BFD_RELOC_CRX_REL8_CMP	
BFD_RELOC_CR16_REGREL16		BFD_RELOC_CRX_SWITCH16	
BFD_RELOC_CR16_REGREL20		BFD_RELOC_CRX_SWITCH32	
BFD_RELOC_CR16_REGREL20a		BFD_RELOC_CRX_SWITCH8	
BFD_RELOC_CR16_REGREL4		BFD_RELOC_CTOR	
BFD_RELOC_CR16_REGREL4a		BFD_RELOC_D10V_10_PCREL_L	
		BFD_RELOC_D10V_10_PCREL_R	
BFD_RELOC_CR16_SWITCH32		BFD_RELOC_D10V_18	
BFD_RELOC_CR16_SWITCH8		BFD_RELOC_D10V_18_PCREL	
BFD_RELOC_CRIS_16_DTPREL		BFD_RELOC_D30V_15	
BFD_RELOC_CRIS_16_GOT		BFD_RELOC_D30V_15_PCREL	
BFD_RELOC_CRIS_16_GOT_GD		BFD_RELOC_D30V_15_PCREL_R	
BFD_RELOC_CRIS_16_GOT_TPREL		BFD_RELOC_D30V_21	
BFD_RELOC_CRIS_16_GOTPLT		BFD_RELOC_D30V_21_PCREL	
BFD_RELOC_CRIS_16_TPREL		BFD_RELOC_D30V_21_PCREL_R	
BFD_RELOC_CRIS_32_DTPREL		BFD_RELOC_D30V_32	
BFD_RELOC_CRIS_32_GD		BFD_RELOC_D30V_32_PCREL	
BFD_RELOC_CRIS_32_GOT	109	BFD_RELOC_D30V_6	
BFD_RELOC_CRIS_32_GOT_GD		BFD_RELOC_D30V_9_PCREL	
BFD_RELOC_CRIS_32_GOT_TPREL	110	BFD_RELOC_D3OV_9_PCREL_R	83
BFD_RELOC_CRIS_32_GOTPLT	109	BFD_RELOC_DLX_HI16_S	84
BFD_RELOC_CRIS_32_GOTREL	109	BFD_RELOC_DLX_JMP26	84
BFD_RELOC_CRIS_32_IE	110	BFD_RELOC_DLX_L016	84
BFD_RELOC_CRIS_32_PLT_GOTREL	109	BFD_RELOC_EPIPHANY_HIGH	127
BFD_RELOC_CRIS_32_PLT_PCREL	109	BFD_RELOC_EPIPHANY_IMM11	
BFD_RELOC_CRIS_32_TPREL	110	BFD_RELOC_EPIPHANY_IMM8	128
BFD_RELOC_CRIS_BDISP8	109	BFD_RELOC_EPIPHANY_LOW	127
BFD_RELOC_CRIS_COPY		BFD_RELOC_EPIPHANY_SIMM11	127
BFD_RELOC_CRIS_DTP		BFD_RELOC_EPIPHANY_SIMM24	
BFD_RELOC_CRIS_DTPMOD		BFD_RELOC_EPIPHANY_SIMM8	
BFD_RELOC_CRIS_GLOB_DAT		BFD_RELOC_FR30_10_IN_8	
BFD_RELOC_CRIS_JUMP_SLOT		BFD_RELOC_FR30_12_PCREL	
BFD_RELOC_CRIS_LAPCQ_OFFSET		BFD_RELOC_FR30_20	
BFD_RELOC_CRIS_RELATIVE		BFD_RELOC_FR30_48.	
BFD_RELOC_CRIS_SIGNED_16		BFD_RELOC_FR30_46	
BFD_RELOC_CRIS_SIGNED_16		BFD_RELOC_FR30_8_IN_8	
BFD_RELOC_CRIS_SIGNED_8		BFD_RELOC_FR30_9_IN_8	
BFD RELOC CRIS UNSIGNED 16		BFD RELOC_FR30_9_IN_8BFD RELOC_FR30_9 PCREL	
DLD VEPOC CUTS ONSTRUCH 10	TOA	DFU_RELUU_FROU_9_FUREL	. ჟა

BFD_RELOC_FRV_FUNCDESC6	39	BFD_RELOC_I370_D12	75
BFD_RELOC_FRV_FUNCDESC_GOT126		BFD_RELOC_I960_CALLJ	62
BFD_RELOC_FRV_FUNCDESC_GOTHI6	39	BFD_RELOC_IA64_COPY	104
BFD_RELOC_FRV_FUNCDESC_GOTLO6	39	BFD_RELOC_IA64_DIR32LSB	103
BFD_RELOC_FRV_FUNCDESC_GOTOFF12 6	39	BFD_RELOC_IA64_DIR32MSB	103
BFD_RELOC_FRV_FUNCDESC_GOTOFFHI 6		BFD_RELOC_IA64_DIR64LSB	103
BFD_RELOC_FRV_FUNCDESC_GOTOFFLO 6		BFD_RELOC_IA64_DIR64MSB	
BFD_RELOC_FRV_FUNCDESC_VALUE6	<del>3</del> 9	BFD_RELOC_IA64_DTPMOD64LSB	105
BFD_RELOC_FRV_GETTLSOFF		BFD_RELOC_IA64_DTPMOD64MSB	
BFD_RELOC_FRV_GETTLSOFF_RELAX		BFD_RELOC_IA64_DTPREL14	
BFD_RELOC_FRV_GOT12		BFD_RELOC_IA64_DTPREL22	
BFD_RELOC_FRV_GOTHI		BFD_RELOC_IA64_DTPREL32LSB	
BFD_RELOC_FRV_GOTLO		BFD_RELOC_IA64_DTPREL32MSB	
BFD_RELOC_FRV_GOTOFF126		BFD_RELOC_IA64_DTPREL64I	
BFD_RELOC_FRV_GOTOFFHI		BFD_RELOC_IA64_DTPREL64LSB	
BFD_RELOC_FRV_GOTOFFLO6		BFD_RELOC_IA64_DTPREL64MSB	
BFD_RELOC_FRV_GOTTLSDESC12		BFD_RELOC_IA64_FPTR32LSB	
BFD_RELOC_FRV_GOTTLSDESCHI		BFD_RELOC_IA64_FPTR32MSB	
BFD_RELOC_FRV_GOTTLSDESCLO		BFD_RELOC_IA64_FPTR64I	
BFD_RELOC_FRV_GOTTLSOFF12		BFD_RELOC_IA64_FPTR64LSB	
BFD_RELOC_FRV_GOTTLSOFFHI		BFD_RELOC_IA64_FPTR64MSB	
BFD_RELOC_FRV_GOTTLSOFFLO		BFD_RELOC_IA64_GPREL22	
BFD_RELOC_FRV_GPREL126		BFD_RELOC_IA64_GPREL32LSB	
BFD_RELOC_FRV_GPREL32		BFD_RELOC_IA64_GPREL32MSB	
BFD_RELOC_FRV_GPRELHI		BFD_RELOC_IA64_GPREL64I	
BFD_RELOC_FRV_GPRELLO		BFD_RELOC_IA64_GPREL64LSB	
BFD_RELOC_FRV_GPRELU126		BFD_RELOC_IA64_GPREL64MSB	
BFD_RELOC_FRV_HI16		BFD_RELOC_IA64_IMM14	
BFD_RELOC_FRV_LABEL16		BFD_RELOC_IA64_IMM22	
BFD_RELOC_FRV_LABEL246		BFD_RELOC_IA64_IMM64	
BFD_RELOC_FRV_L016		BFD_RELOC_IA64_IPLTLSB	
BFD_RELOC_FRV_TLSDESC_RELAX		BFD_RELOC_IA64_IPLTMSB	
BFD_RELOC_FRV_TLSDESC_VALUE		BFD_RELOC_IA64_LDXMOV	
BFD_RELOC_FRV_TLSMOFF		BFD_RELOC_IA64_LTOFF_DTPMOD22	
BFD_RELOC_FRV_TLSMOFF127		BFD_RELOC_IA64_LTOFF_DTPREL22	
BFD_RELOC_FRV_TLSMOFFHI		BFD_RELOC_IA64_LT0FF_FPTR22	
BFD_RELOC_FRV_TLSMOFFLO		BFD_RELOC_IA64_LTOFF_FPTR32LSB	
BFD_RELOC_FRV_TLSOFF 7		BFD_RELOC_IA64_LTOFF_FPTR32MSB	
BFD_RELOC_FRV_TLSOFF_RELAX 7		BFD_RELOC_IA64_LTOFF_FPTR64I	
BFD_RELOC_GPREL166		BFD_RELOC_IA64_LTOFF_FPTR64LSB	
BFD_RELOC_GPREL326		BFD_RELOC_IA64_LTOFF_FPTR64MSB	
BFD_RELOC_H8_DIR16A811		BFD_RELOC_IA64_LTOFF_TPREL22	
BFD_RELOC_H8_DIR16R811	l1	BFD_RELOC_IA64_LT0FF22	103
BFD_RELOC_H8_DIR24A811	l1	BFD_RELOC_IA64_LTOFF22X	104
BFD_RELOC_H8_DIR24R811	l1	BFD_RELOC_IA64_LTOFF64I	103
BFD_RELOC_H8_DIR32A1611	l1	BFD_RELOC_IA64_LTV32LSB	104
BFD_RELOC_H8_DISP32A1611	l1	BFD_RELOC_IA64_LTV32MSB	104
BFD_RELOC_HI166	66	BFD_RELOC_IA64_LTV64LSB	104
BFD_RELOC_HI16_BASEREL6	31	BFD_RELOC_IA64_LTV64MSB	104
BFD_RELOC_HI16_GOTOFF6	31	BFD_RELOC_IA64_PCREL21B	104
BFD_RELOC_HI16_PCREL 6		BFD_RELOC_IA64_PCREL21BI	
BFD_RELOC_HI16_PLTOFF6	51	BFD_RELOC_IA64_PCREL21F	104
BFD_RELOC_HI16_S6		BFD_RELOC_IA64_PCREL21M	
BFD_RELOC_HI16_S_BASEREL 6	31	BFD_RELOC_IA64_PCREL22	104
BFD_RELOC_HI16_S_GOTOFF6		BFD_RELOC_IA64_PCREL32LSB	
BFD_RELOC_HI16_S_PCREL6		BFD_RELOC_IA64_PCREL32MSB	
BFD_RELOC_HI16_S_PLT0FF6		BFD_RELOC_IA64_PCREL60B	
			104

BFD_RELOC_IA64_PCREL64LSB	104	BFD_RELOC_M32R_18_PCREL84
BFD_RELOC_IA64_PCREL64MSB	104	BFD_RELOC_M32R_2484
BFD_RELOC_IA64_PLT0FF22	103	BFD_RELOC_M32R_26_PCREL84
BFD_RELOC_IA64_PLT0FF64I		BFD_RELOC_M32R_26_PLTREL
BFD_RELOC_IA64_PLT0FF64LSB		BFD_RELOC_M32R_COPY85
BFD_RELOC_IA64_PLTOFF64MSB	103	BFD_RELOC_M32R_GLOB_DAT85
BFD_RELOC_IA64_REL32LSB		BFD_RELOC_M32R_GOT16_HI_SLO85
BFD_RELOC_IA64_REL32MSB		BFD_RELOC_M32R_GOT16_HI_UL085
BFD_RELOC_IA64_REL64LSB		BFD_RELOC_M32R_GOT16_L085
BFD_RELOC_IA64_REL64MSB		BFD_RELOC_M32R_GOT24
BFD_RELOC_IA64_SECREL32LSB		BFD_RELOC_M32R_GOTOFF85
BFD_RELOC_IA64_SECREL32MSB		BFD_RELOC_M32R_GOTOFF_HI_SLO
BFD_RELOC_IA64_SECREL64LSB		BFD_RELOC_M32R_GOTOFF_HI_ULO
BFD_RELOC_IA64_SECREL64MSB		BFD_RELOC_M32R_GOTOFF_LO
BFD_RELOC_IA64_SEGREL32LSB		BFD_RELOC_M32R_GOTPC_HI_SLO85
BFD_RELOC_IA64_SEGREL32MSB		BFD_RELOC_M32R_GOTPC_HI_ULO
BFD_RELOC_IA64_SEGREL64LSB		BFD_RELOC_M32R_GOTPC_LO85
BFD_RELOC_IA64_SEGREL64MSB		BFD_RELOC_M32R_GOTPC2485
BFD_RELOC_IA64_TPREL14		BFD_RELOC_M32R_HI16_SL0
BFD_RELOC_IA64_TPREL14		BFD_RELOC_M32R_HI16_UL0
		BFD_RELOC_M32R_JMP_SLOT
BFD_RELOC_IA64_TPREL64I		
BFD_RELOC_IA64_TPREL64LSB		BFD_RELOC_M32R_L016
BFD_RELOC_IA64_TPREL64MSB		BFD_RELOC_M32R_RELATIVE
BFD_RELOC_IP2K_ADDR16CJP		BFD_RELOC_M32R_SDA16
BFD_RELOC_IP2K_BANK		BFD_RELOC_M68HC11_24
BFD_RELOC_IP2K_EX8DATA		BFD_RELOC_M68HC11_3B
BFD_RELOC_IP2K_FR_OFFSET		BFD_RELOC_M68HC11_HI8
BFD_RELOC_IP2K_FR9		BFD_RELOC_M68HC11_L016
BFD_RELOC_IP2K_HI8DATA		BFD_RELOC_M68HC11_L08
BFD_RELOC_IP2K_HI8INSN		${\tt BFD\_RELOC\_M68HC11\_PAGE$
BFD_RELOC_IP2K_LO8DATA		${\tt BFD\_RELOC\_M68HC11\_RL\_GROUP} \dots 105$
BFD_RELOC_IP2K_LO8INSN		${\tt BFD\_RELOC\_M68HC11\_RL\_JUMP} \dots 105$
BFD_RELOC_IP2K_PAGE3		${\tt BFD\_RELOC\_M68HC12\_10\_PCREL} \dots 106$
BFD_RELOC_IP2K_PC_SKIP		BFD_RELOC_M68HC12_16B
BFD_RELOC_IP2K_TEXT		${\tt BFD\_RELOC\_M68HC12\_5B} \dots \dots 105$
BFD_RELOC_IQ2000_OFFSET_16	113	${\tt BFD\_RELOC\_M68HC12\_9\_PCREL} \dots 106$
BFD_RELOC_IQ2000_OFFSET_21	113	BFD_RELOC_M68HC12_9B
BFD_RELOC_IQ2000_UHI16	113	BFD_RELOC_M68HC12_HI8XG 107
BFD_RELOC_LM32_16_GOT	115	BFD_RELOC_M68HC12_L08XG 107
BFD_RELOC_LM32_BRANCH	115	${\tt BFD\_RELOC\_MACH\_O\_LOCAL\_SECTDIFF} \dots 116$
BFD_RELOC_LM32_CALL	115	BFD_RELOC_MACH_O_PAIR116
BFD_RELOC_LM32_COPY		BFD_RELOC_MACH_O_SECTDIFF 116
BFD_RELOC_LM32_GLOB_DAT	115	BFD_RELOC_MACH_0_X86_64_BRANCH32 116
BFD_RELOC_LM32_GOTOFF_HI16	115	BFD_RELOC_MACH_O_X86_64_BRANCH8 116
BFD_RELOC_LM32_GOTOFF_L016	115	BFD_RELOC_MACH_O_X86_64_GOT116
BFD_RELOC_LM32_JMP_SLOT	115	BFD_RELOC_MACH_O_X86_64_GOT_LOAD 116
BFD_RELOC_LM32_RELATIVE	116	BFD_RELOC_MACH_O_X86_64_PCREL32_1 116
BFD_RELOC_LO10	62	BFD_RELOC_MACH_O_X86_64_PCREL32_2 116
BFD_RELOC_L016		BFD_RELOC_MACH_0_X86_64_PCREL32_4 116
BFD_RELOC_LO16_BASEREL		BFD_RELOC_MACH_O_X86_64_SUBTRACTOR32 116
BFD_RELOC_L016_GOTOFF		BFD_RELOC_MACH_O_X86_64_SUBTRACTOR64 116
BFD_RELOC_L016_PCREL		BFD_RELOC_MCORE_PCREL_32
BFD_RELOC_L016_PLT0FF		BFD_RELOC_MCORE_PCREL_IMM11BY293
BFD_RELOC_M32C_HI8		BFD_RELOC_MCORE_PCREL_IMM4BY293
BFD_RELOC_M32C_RL_1ADDR		BFD_RELOC_MCORE_PCREL_IMM8BY493
BFD_RELOC_M32C_RL_2ADDR		BFD_RELOC_MCORE_PCREL_JSR_IMM11BY293
BFD_RELOC_M32C_RL_JUMP		BFD_RELOC_MCORE_RVA
RED RELOC M32R 10 PCREI		BED RELOC MEP 16

BFD_RELOC_MEP_32	93	BFD_RELOC_MICROBLAZE_32_LO1	16
BFD_RELOC_MEP_8	93	BFD_RELOC_MICROBLAZE_32_LO_PCREL 1	
BFD_RELOC_MEP_ADDR24A4	94	BFD_RELOC_MICROBLAZE_32_ROSDA1	16
BFD_RELOC_MEP_GNU_VTENTRY	94	BFD_RELOC_MICROBLAZE_32_RWSDA 1	17
BFD_RELOC_MEP_GNU_VTINHERIT		BFD_RELOC_MICROBLAZE_32_SYM_OP_SYM 1	
BFD_RELOC_MEP_GPREL	94	BFD_RELOC_MICROBLAZE_32_TLSDTPMOD 1	17
BFD_RELOC_MEP_HI16S	93	BFD_RELOC_MICROBLAZE_32_TLSDTPREL 1	18
BFD_RELOC_MEP_HI16U	93	BFD_RELOC_MICROBLAZE_64_GOT1	17
BFD_RELOC_MEP_LOW16	93	BFD_RELOC_MICROBLAZE_64_GOTOFF 1	17
BFD_RELOC_MEP_PCABS24A2	93	BFD_RELOC_MICROBLAZE_64_GOTPC 1	17
BFD_RELOC_MEP_PCREL12A2	93	BFD_RELOC_MICROBLAZE_64_NONE 1	17
BFD_RELOC_MEP_PCREL17A2		BFD_RELOC_MICROBLAZE_64_PLT1	17
BFD_RELOC_MEP_PCREL24A2	93	BFD_RELOC_MICROBLAZE_64_TLS1	17
BFD_RELOC_MEP_PCREL8A2	93	BFD_RELOC_MICROBLAZE_64_TLSDTPREL 1	
BFD_RELOC_MEP_TPREL		BFD_RELOC_MICROBLAZE_64_TLSGD1	
BFD_RELOC_MEP_TPREL7	94	BFD_RELOC_MICROBLAZE_64_TLSGOTTPREL 1	18
BFD_RELOC_MEP_TPREL7A2	94	BFD_RELOC_MICROBLAZE_64_TLSLD1	17
BFD_RELOC_MEP_TPREL7A4		BFD_RELOC_MICROBLAZE_64_TLSTPREL 1	
BFD_RELOC_MEP_UIMM24		BFD_RELOC_MICROBLAZE_COPY 1	
BFD_RELOC_METAG_COPY		BFD_RELOC_MICROMIPS_10_PCREL_S1	
BFD_RELOC_METAG_GETSET_GOT		BFD_RELOC_MICROMIPS_16_PCREL_S1	
BFD_RELOC_METAG_GETSET_GOTOFF		BFD_RELOC_MICROMIPS_18_PCREL_S3	
BFD_RELOC_METAG_GETSETOFF		BFD_RELOC_MICROMIPS_19_PCREL_S2	
BFD_RELOC_METAG_GLOB_DAT		BFD_RELOC_MICROMIPS_21_PCREL_S1	
BFD_RELOC_METAG_GOTOFF		BFD_RELOC_MICROMIPS_26_PCREL_S1	
BFD_RELOC_METAG_HI16_GOTOFF		BFD_RELOC_MICROMIPS_7_PCREL_S1	
BFD_RELOC_METAG_HI16_GOTPC		BFD_RELOC_MICROMIPS_CALL_HI16	
BFD_RELOC_METAG_HI16_PLT		BFD_RELOC_MICROMIPS_CALL_LO16	
BFD_RELOC_METAG_HIADDR16		BFD_RELOC_MICROMIPS_CALL16	
BFD_RELOC_METAG_HIOG		BFD_RELOC_MICROMIPS_GOT_DISP	
BFD_RELOC_METAG_JMP_SLOT		BFD_RELOC_MICROMIPS_GOT_HI16	
BFD_RELOC_METAG_LO16_GOTOFF		BFD_RELOC_MICROMIPS_GOT_L016	
BFD_RELOC_METAG_LO16_GOTPC		BFD_RELOC_MICROMIPS_GOT_OFST	
BFD_RELOC_METAG_LO16_PLT		BFD_RELOC_MICROMIPS_GOT_PAGE	
BFD_RELOC_METAG_LOADDR16		BFD_RELOC_MICROMIPS_GOT16	
BFD_RELOC_METAG_LOOG		BFD_RELOC_MICROMIPS_GPREL16	
BFD_RELOC_METAG_PLT		BFD_RELOC_MICROMIPS_HI16	
BFD_RELOC_METAG_REL16		BFD_RELOC_MICROMIPS_HI16_S	
BFD_RELOC_METAG_REL8		BFD_RELOC_MICROMIPS_HI16_S_PCREL	
BFD_RELOC_METAG_RELATIVE		BFD_RELOC_MICROMIPS_HIGHER	
BFD_RELOC_METAG_RELBRANCH		BFD_RELOC_MICROMIPS_HIGHEST	
BFD_RELOC_METAG_RELBRANCH_PLT		BFD_RELOC_MICROMIPS_JALR	
BFD_RELOC_METAG_TLS_DTPMOD		BFD_RELOC_MICROMIPS_JMP	
BFD_RELOC_METAG_TLS_DTPOFF		BFD_RELOC_MICROMIPS_LITERAL	
BFD_RELOC_METAG_TLS_GD		BFD_RELOC_MICROMIPS_L016	
BFD_RELOC_METAG_TLS_IE		BFD_RELOC_MICROMIPS_L016_PCREL	
BFD_RELOC_METAG_TLS_IENONPIC		BFD_RELOC_MICROMIPS_SCN_DISP	
BFD_RELOC_METAG_TLS_IENONPIC_HI16		BFD_RELOC_MICROMIPS_SUB	
BFD_RELOC_METAG_TLS_IENONPIC_L016		BFD_RELOC_MICROMIPS_TLS_DTPREL_HI16	
BFD_RELOC_METAG_TLS_IENONFIC_E010		BFD_RELOC_MICROMIPS_TLS_DTPREL_L016	
BFD_RELOC_METAG_TLS_LDM		BFD_RELOC_MICROMIPS_TLS_DIFFREL_LOTO  BFD_RELOC_MICROMIPS_TLS_GD	
BFD_RELOC_METAG_TLS_LDO.HI16		BFD_RELOC_MICROMIPS_TLS_GD BFD_RELOC_MICROMIPS_TLS_GOTTPREL	
BFD_RELOC_METAG_TLS_LDO_HIT6BFD_RELOC_METAG_TLS_LDO_L016		BFD_RELOC_MICROMIPS_TLS_GOTTPREEL BFD_RELOC_MICROMIPS_TLS_LDM	
BFD_RELOC_METAG_TLS_LE		BFD_RELOC_MICROMIPS_TLS_EDM	
BFD_RELOC_METAG_TLS_LE_HI16		BFD_RELOC_MICROMIPS_TLS_TPREL_L016	
BFD_RELOC_METAG_TLS_LE_L016		BFD_RELOC_MIPS_18_PCREL_S3	
BFD_RELOC_METAG_TLS_LE_LUTOBFD_RELOC_METAG_TLS_TPOFF		BFD_RELOC_MIPS_19_PCREL_S2	
RED RELOC MICRORIATE 32 COTOFF		RED RELOC MIPS 21 PCREL S2	

BFD_RELOC_MIPS_26_PCREL_S2		BFD_RELOC_MMIX_CBRANCH_1	
BFD_RELOC_MIPS_CALL_HI16		BFD_RELOC_MMIX_CBRANCH_2	
BFD_RELOC_MIPS_CALL_L016		BFD_RELOC_MMIX_CBRANCH_3	
BFD_RELOC_MIPS_CALL16		BFD_RELOC_MMIX_CBRANCH_J	
BFD_RELOC_MIPS_COPY	69	BFD_RELOC_MMIX_GETA	
BFD_RELOC_MIPS_DELETE		BFD_RELOC_MMIX_GETA_1	
BFD_RELOC_MIPS_EH		BFD_RELOC_MMIX_GETA_2	
BFD_RELOC_MIPS_GOT_DISP		BFD_RELOC_MMIX_GETA_3	
BFD_RELOC_MIPS_GOT_HI16		BFD_RELOC_MMIX_JMP	
BFD_RELOC_MIPS_GOT_L016		BFD_RELOC_MMIX_JMP_1	
BFD_RELOC_MIPS_GOT_OFST		BFD_RELOC_MMIX_JMP_2	
BFD_RELOC_MIPS_GOT_PAGE		BFD_RELOC_MMIX_JMP_3	
BFD_RELOC_MIPS_GOT16		BFD_RELOC_MMIX_LOCAL	
BFD_RELOC_MIPS_HIGHER		BFD_RELOC_MMIX_PUSHJ	
BFD_RELOC_MIPS_HIGHEST		BFD_RELOC_MMIX_PUSHJ_1	
BFD_RELOC_MIPS_INSERT_A		BFD_RELOC_MMIX_PUSHJ_2	
BFD_RELOC_MIPS_INSERT_B		BFD_RELOC_MMIX_PUSHJ_3	
BFD_RELOC_MIPS_JALR		BFD_RELOC_MMIX_PUSHJ_STUBBABLE	
BFD_RELOC_MIPS_JMP		BFD_RELOC_MMIX_REG	
BFD_RELOC_MIPS_JUMP_SLOT	69	BFD_RELOC_MMIX_REG_OR_BYTE	
BFD_RELOC_MIPS_LITERAL	67	BFD_RELOC_MN10300_16_PCREL	71
BFD_RELOC_MIPS_REL16	68	BFD_RELOC_MN10300_32_PCREL	71
BFD_RELOC_MIPS_RELGOT	68	BFD_RELOC_MN10300_ALIGN	71
BFD_RELOC_MIPS_SCN_DISP	68	BFD_RELOC_MN10300_COPY	70
BFD_RELOC_MIPS_SHIFT5	68	BFD_RELOC_MN10300_GLOB_DAT	70
BFD_RELOC_MIPS_SHIFT6	68	BFD_RELOC_MN10300_GOT16	70
BFD_RELOC_MIPS_SUB	68	BFD_RELOC_MN10300_GOT24	70
BFD_RELOC_MIPS_TLS_DTPMOD32		BFD_RELOC_MN10300_GOT32	
BFD_RELOC_MIPS_TLS_DTPMOD64		BFD_RELOC_MN10300_GOTOFF24	
BFD_RELOC_MIPS_TLS_DTPREL_HI16		BFD_RELOC_MN10300_JMP_SLOT	
BFD_RELOC_MIPS_TLS_DTPREL_L016		BFD_RELOC_MN10300_RELATIVE	
BFD_RELOC_MIPS_TLS_DTPREL32		BFD_RELOC_MN10300_SYM_DIFF	
BFD_RELOC_MIPS_TLS_DTPREL64		BFD_RELOC_MN10300_TLS_DTPMOD	
BFD_RELOC_MIPS_TLS_GD		BFD_RELOC_MN10300_TLS_DTP0FF	
BFD_RELOC_MIPS_TLS_GOTTPREL		BFD_RELOC_MN10300_TLS_GD	
BFD_RELOC_MIPS_TLS_LDM		BFD_RELOC_MN10300_TLS_GOTIE	
BFD_RELOC_MIPS_TLS_TPREL_HI16		BFD_RELOC_MN10300_TLS_IE	71
BFD_RELOC_MIPS_TLS_TPREL_L016		BFD_RELOC_MN10300_TLS_LD	71
BFD_RELOC_MIPS_TLS_TPREL32		BFD_RELOC_MN10300_TLS_LD0	
BFD_RELOC_MIPS_TLS_TPREL64		BFD_RELOC_MN10300_TLS_LE	
BFD_RELOC_MIPS16_CALL16		BFD_RELOC_MN10300_TLS_TP0FF	
BFD_RELOC_MIPS16_GOT16		BFD_RELOC_MOXIE_10_PCREL	
BFD_RELOC_MIPS16_GPREL		BFD_RELOC_MSP430_10_PCREL	
BFD_RELOC_MIPS16_HI16		BFD_RELOC_MSP430_16	
BFD_RELOC_MIPS16_HI16_S		BFD_RELOC_MSP430_16_BYTE	
BFD_RELOC_MIPS16_JMP		BFD_RELOC_MSP430_16_PCREL	
BFD_RELOC_MIPS16_L016		BFD_RELOC_MSP430_16_PCREL_BYTE	
BFD_RELOC_MIPS16_TLS_DTPREL_HI16		BFD_RELOC_MSP430_2X_PCREL	
BFD_RELOC_MIPS16_TLS_DTPREL_LO16 BFD_RELOC_MIPS16_TLS_GD		BFD_RELOC_MSP430_ABS_HI16	
BFD_RELOC_MIPS16_TLS_GOTTPREL		BFD_RELOC_MSP430_PREL31	
BFD_RELOC_MIPS16_TLS_LDM			
BFD_RELOC_MIPS16_TLS_TPREL_HI16		BFD_RELOC_MSP430_SYM_DIFF	
BFD_RELOC_MIPS16_TLS_TPREL_L016		BFD_RELOC_MSP430X_ABS16	
BFD_RELOC_MMIX_ADDR19		BFD_RELOC_MSP430X_ABS20_ADR_DST	
BFD_RELOC_MMIX_ADDR27		BFD_RELOC_MSP430X_ABS20_ADR_SRC	
BFD_RELOC_MMIX_BASE_PLUS_OFFSET		BFD_RELOC_MSP430X_ABS20_EXT_DST	
BFD RELOC MMIX CBRANCH	95	BFD RELOC MSP430X ABS20 EXT ODST	112

BFD_RELOC_MSP430X_ABS20_EXT_SRC 11		BFD_RELOC_NDS32_INSN16	
BFD_RELOC_MSP430X_PCR16 11		BFD_RELOC_NDS32_JMP_SLOT	
BFD_RELOC_MSP430X_PCR20_CALL11		BFD_RELOC_NDS32_LABEL	
BFD_RELOC_MSP430X_PCR20_EXT_DST 11		BFD_RELOC_NDS32_L012S0	
BFD_RELOC_MSP430X_PCR20_EXT_ODST 11		BFD_RELOC_NDS32_L012S0_ORI	
BFD_RELOC_MSP430X_PCR20_EXT_SRC 11		BFD_RELOC_NDS32_L012S1	
BFD_RELOC_MT_GNU_VTENTRY 11	12	BFD_RELOC_NDS32_L012S2	86
BFD_RELOC_MT_GNU_VTINHERIT11	12	BFD_RELOC_NDS32_L012S2_DP	88
BFD_RELOC_MT_HI16 11	12	BFD_RELOC_NDS32_L012S2_SP	88
BFD_RELOC_MT_L016 11	12	BFD_RELOC_NDS32_L012S3	86
BFD_RELOC_MT_PC16 11	12	BFD_RELOC_NDS32_LOADSTORE	87
BFD_RELOC_MT_PCINSN811		BFD_RELOC_NDS32_LONGCALL1	
BFD_RELOC_NDS32_10_UPCREL 8	88	BFD_RELOC_NDS32_LONGCALL2	87
BFD_RELOC_NDS32_10IFCU_PCREL 8		BFD_RELOC_NDS32_LONGCALL3	
BFD_RELOC_NDS32_15_FIXED 8		BFD_RELOC_NDS32_LONGCALL4	
BFD_RELOC_NDS32_15_PCREL 8		BFD_RELOC_NDS32_LONGCALL5	
BFD_RELOC_NDS32_17_FIXED		BFD_RELOC_NDS32_LONGCALL6	
BFD_RELOC_NDS32_17_PCREL 8		BFD_RELOC_NDS32_LONGJUMP1	
BFD_RELOC_NDS32_17IFC_PCREL		BFD_RELOC_NDS32_LONGJUMP2	
BFD_RELOC_NDS32_20		BFD_RELOC_NDS32_LONGJUMP3	
BFD_RELOC_NDS32_25_ABS		BFD_RELOC_NDS32_LONGJUMP4	
BFD_RELOC_NDS32_25_FIXED		BFD_RELOC_NDS32_LONGJUMP5	
BFD_RELOC_NDS32_25_PCREL		BFD_RELOC_NDS32_LONGJUMP6	
BFD_RELOC_NDS32_25_PLTREL		BFD_RELOC_NDS32_LONGJUMP7	
BFD_RELOC_NDS32_5		BFD_RELOC_NDS32_MINUEND	
BFD_RELOC_NDS32_9_FIXED		BFD_RELOC_NDS32_MULCALL_SUFF	
BFD_RELOC_NDS32_9_PCREL		BFD_RELOC_NDS32_PLT_GOT_SUFF	
BFD_RELOC_NDS32_9_PLTREL		BFD_RELOC_NDS32_PLT_GOTREL_HI20	
BFD_RELOC_NDS32_COPY			
BFD_RELOC_NDS32_DATA		BFD_RELOC_NDS32_PLT_GOTREL_L012BFD_RELOC_NDS32_PLT_GOTREL_L015	
BFD_RELOC_NDS32_DIFF_ULEB128		BFD_RELOC_NDS32_PLT_GOTREL_L019	
BFD_RELOC_NDS32_DIFF16		BFD_RELOC_NDS32_PLT_GOTREL_LO20	
BFD_RELOC_NDS32_DIFF32		BFD_RELOC_NDS32_PLTBLOCK	
BFD_RELOC_NDS32_DIFF8		BFD_RELOC_NDS32_PLTREL_HI20	
BFD_RELOC_NDS32_DWARF2_LEB		BFD_RELOC_NDS32_PLTREL_L012	
BFD_RELOC_NDS32_DWARF2_OP1		BFD_RELOC_NDS32_PTR	
BFD_RELOC_NDS32_DWARF2_OP2		BFD_RELOC_NDS32_PTR_COUNT	
BFD_RELOC_NDS32_EMPTY		BFD_RELOC_NDS32_PTR_RESOLVED	
BFD_RELOC_NDS32_GLOB_DAT		BFD_RELOC_NDS32_RELATIVE	
BFD_RELOC_NDS32_GOT_HI20 8		BFD_RELOC_NDS32_RELAX_ENTRY	
BFD_RELOC_NDS32_GOT_LO12 8		BFD_RELOC_NDS32_RELAX_REGION_BEGIN	
BFD_RELOC_NDS32_GOT_LO15 8		BFD_RELOC_NDS32_RELAX_REGION_END	
BFD_RELOC_NDS32_GOT_L019 8		BFD_RELOC_NDS32_SDA_FP7U2_RELA	
BFD_RELOC_NDS32_GOT_SUFF 8		BFD_RELOC_NDS32_SDA12S2_DP	
BFD_RELOC_NDS32_GOT15S2 8		BFD_RELOC_NDS32_SDA12S2_SP	
BFD_RELOC_NDS32_GOT17S2		BFD_RELOC_NDS32_SDA15S0	
BFD_RELOC_NDS32_GOT20 8	87	BFD_RELOC_NDS32_SDA15S1	86
BFD_RELOC_NDS32_GOTOFF		BFD_RELOC_NDS32_SDA15S2	86
BFD_RELOC_NDS32_GOTOFF_HI20	87	BFD_RELOC_NDS32_SDA15S3	86
BFD_RELOC_NDS32_GOTOFF_LO12		BFD_RELOC_NDS32_SDA16S3	86
BFD_RELOC_NDS32_GOTOFF_LO15	88	BFD_RELOC_NDS32_SDA17S2	86
BFD_RELOC_NDS32_GOTOFF_LO19		BFD_RELOC_NDS32_SDA18S1	86
BFD_RELOC_NDS32_GOTOFF_SUFF		BFD_RELOC_NDS32_SDA19S0	
BFD_RELOC_NDS32_GOTPC_HI20 8		BFD_RELOC_NDS32_SUBTRAHEND	
BFD_RELOC_NDS32_GOTPC_LO12 8		BFD_RELOC_NDS32_TLS_IE_HI20	
BFD_RELOC_NDS32_GOTPC20		BFD_RELOC_NDS32_TLS_IE_L012S2	
BFD_RELOC_NDS32_GOTTPOFF		BFD_RELOC_NDS32_TLS_LE_15S0	
BFD_RELOC_NDS32_HI20 8		BFD_RELOC_NDS32_TLS_LE_15S1	

BFD_RELOC_NDS32_TLS_LE_15S289	BFD_RELOC_NS32K_IMM_16_PCREL72
BFD_RELOC_NDS32_TLS_LE_20 89	BFD_RELOC_NS32K_IMM_3272
BFD_RELOC_NDS32_TLS_LE_ADD 89	BFD_RELOC_NS32K_IMM_32_PCREL
BFD_RELOC_NDS32_TLS_LE_HI20 89	BFD_RELOC_NS32K_IMM_872
BFD_RELOC_NDS32_TLS_LE_LO1289	BFD_RELOC_NS32K_IMM_8_PCREL
BFD_RELOC_NDS32_TLS_LE_LS 89	BFD_RELOC_OR1K_COPY 111
BFD_RELOC_NDS32_TLS_TPOFF 89	BFD_RELOC_OR1K_GLOB_DAT 111
BFD_RELOC_NDS32_TPOFF89	BFD_RELOC_OR1K_GOT16111
BFD_RELOC_NDS32_TRAN	BFD_RELOC_OR1K_GOTOFF_HI16111
BFD_RELOC_NDS32_UPDATE_TA 88	BFD_RELOC_OR1K_GOTOFF_LO16111
BFD_RELOC_NDS32_WORD_9_PCREL85	BFD_RELOC_OR1K_GOTPC_HI16 111
BFD_RELOC_NIOS2_ALIGN	BFD_RELOC_OR1K_GOTPC_LO16 111
BFD_RELOC_NIOS2_CACHE_OPX 113	BFD_RELOC_OR1K_JMP_SLOT 111
BFD_RELOC_NIOS2_CALL_HA 113	BFD_RELOC_OR1K_PLT26111
BFD_RELOC_NIOS2_CALL_LO 113	BFD_RELOC_OR1K_REL_26111
BFD_RELOC_NIOS2_CALL16	BFD_RELOC_OR1K_RELATIVE 111
BFD_RELOC_NIOS2_CALL26	BFD_RELOC_OR1K_TLS_DTPMOD 111
BFD_RELOC_NIOS2_CALL26_NOAT	BFD_RELOC_OR1K_TLS_DTPOFF 111
BFD_RELOC_NIOS2_CALLR	BFD_RELOC_OR1K_TLS_GD_HI16111
BFD_RELOC_NIOS2_CJMP	BFD_RELOC_OR1K_TLS_GD_LO16
BFD_RELOC_NIOS2_COPY	BFD_RELOC_OR1K_TLS_IE_HI16111
BFD_RELOC_NIOS2_GLOB_DAT	BFD_RELOC_OR1K_TLS_IE_LO16111
BFD_RELOC_NIOS2_GOT_HA	BFD_RELOC_OR1K_TLS_LDM_HI16111
BFD_RELOC_NIOS2_GOT_LO	BFD_RELOC_OR1K_TLS_LDM_L016
BFD_RELOC_NIOS2_GOT16	BFD_RELOC_OR1K_TLS_LDO_HI16
BFD_RELOC_NIOS2_GOTOFF	BFD_RELOC_OR1K_TLS_LDO_L016
BFD_RELOC_NIOS2_GOTOFF_HA 113	BFD_RELOC_OR1K_TLS_LE_HI16111
BFD_RELOC_NIOS2_GOTOFF_LO	BFD_RELOC_OR1K_TLS_LE_L016
BFD_RELOC_NIOS2_GPREL	BFD_RELOC_OR1K_TLS_TPOFF
BFD_RELOC_NIOS2_HI16	BFD_RELOC_PDP11_DISP_6_PCREL
BFD_RELOC_NIOS2_HIADJ16	BFD_RELOC_PDP11_DISP_8_PCREL
BFD_RELOC_NIOS2_IMM5	BFD_RELOC_PJ_CODE_DIR16
BFD_RELOC_NIOS2_IMM6	BFD_RELOC_PJ_CODE_DIR32
BFD_RELOC_NIOS2_IMM8	BFD_RELOC_PJ_CODE_HI16
BFD_RELOC_NIOS2_JUMP_SLOT	BFD_RELOC_PJ_CODE_LO16
BFD_RELOC_NIOS2_L016	BFD_RELOC_PJ_CODE_REL16
BFD_RELOC_NIOS2_PCREL_HA	BFD_RELOC_PJ_CODE_REL32
BFD_RELOC_NIOS2_PCREL_LO	BFD_RELOC_PPC_B16
BFD_RELOC_NIOS2_RELATIVE	BFD_RELOC_PPC_B16_BRNTAKEN73
BFD_RELOC_NIOS2_S16	BFD_RELOC_PPC_B16_BRTAKEN
BFD_RELOC_NIOS2_TLS_DTPMOD	BFD_RELOC_PPC_B26
BFD_RELOC_NIOS2_TLS_DTPREL	BFD_RELOC_PPC_BA16
BFD_RELOC_NIOS2_TLS_GD16	BFD_RELOC_PPC_BA16_BRNTAKEN73
BFD_RELOC_NIOS2_TLS_IE16	BFD_RELOC_PPC_BA16_BRTAKEN73
BFD_RELOC_NIOS2_TLS_LDM16	BFD_RELOC_PPC_BA26
BFD_RELOC_NIOS2_TLS_LD016	BFD_RELOC_PPC_COPY
BFD_RELOC_NIOS2_TLS_LE16	BFD_RELOC_PPC_DTPMOD
BFD_RELOC_NIOS2_TLS_TPREL	BFD_RELOC_PPC_DTPREL
BFD_RELOC_NIOS2_U16	BFD_RELOC_PPC_DTPREL1674
BFD_RELOC_NIOS2_UJMP	
	BFD_RELOC_PPC_DTPREL16_HA
BFD_RELOC_NONE	BFD_RELOC_PPC_DTPREL16_HI
BFD_RELOC_NS32K_DISP_16	BFD_RELOC_PPC_DTPREL16_L0
BFD_RELOC_NS32K_DISP_16_PCREL	BFD_RELOC_PPC_EMB_BIT_FLD
BFD_RELOC_NS32K_DISP_32	BFD_RELOC_PPC_EMB_MRKREF
BFD_RELOC_NS32K_DISP_32_PCREL	BFD_RELOC_PPC_EMB_NADDR16
BFD_RELOC_NS32K_DISP_8	BFD_RELOC_PPC_EMB_NADDR16_HA
BFD_RELOC_NS32K_DISP_8_PCREL         72           BFD_RELOC_NS32K_IMM_16         72	BFD_RELOC_PPC_EMB_NADDR16_HI         73           BFD_RELOC_PPC_EMB_NADDR16_LO         73
DFD RELUC NO32N ITM 10	DFU_RELUC_PPC_EMB_NAUUK16_LU

BFD_RELOC_PPC_EMB_NADDR32	73	BFD_RELOC_PPC64_ADDR16_HIGHA	74
BFD_RELOC_PPC_EMB_RELSDA	73	BFD_RELOC_PPC64_ADDR16_LO_DS	74
BFD_RELOC_PPC_EMB_RELSEC16	73	BFD_RELOC_PPC64_ADDR64_LOCAL	74
BFD_RELOC_PPC_EMB_RELST_HA	73	BFD_RELOC_PPC64_DTPREL16_DS	75
BFD_RELOC_PPC_EMB_RELST_HI	73	BFD_RELOC_PPC64_DTPREL16_HIGH	75
BFD_RELOC_PPC_EMB_RELST_LO	73	BFD_RELOC_PPC64_DTPREL16_HIGHA	75
BFD_RELOC_PPC_EMB_SDA21	73	BFD_RELOC_PPC64_DTPREL16_HIGHER	75
BFD_RELOC_PPC_EMB_SDA2I16		BFD_RELOC_PPC64_DTPREL16_HIGHERA	75
BFD_RELOC_PPC_EMB_SDA2REL		BFD_RELOC_PPC64_DTPREL16_HIGHEST	75
BFD_RELOC_PPC_EMB_SDAI16		BFD_RELOC_PPC64_DTPREL16_HIGHESTA	
BFD_RELOC_PPC_GLOB_DAT		BFD_RELOC_PPC64_DTPREL16_LO_DS	
BFD_RELOC_PPC_GOT_DTPREL16		BFD_RELOC_PPC64_GOT16_DS	
BFD_RELOC_PPC_GOT_DTPREL16_HA		BFD_RELOC_PPC64_GOT16_LO_DS	
BFD_RELOC_PPC_GOT_DTPREL16_HI		BFD_RELOC_PPC64_HIGHER	
BFD_RELOC_PPC_GOT_DTPREL16_L0		BFD_RELOC_PPC64_HIGHER_S	
BFD_RELOC_PPC_GOT_TLSGD16		BFD_RELOC_PPC64_HIGHEST	
BFD_RELOC_PPC_GOT_TLSGD16_HA	75	BFD_RELOC_PPC64_HIGHEST_S	
BFD_RELOC_PPC_GOT_TLSGD16_HI		BFD_RELOC_PPC64_PLT16_LO_DS	
BFD_RELOC_PPC_GOT_TLSGD16_L0		BFD_RELOC_PPC64_PLTGOT16	
BFD_RELOC_PPC_GOT_TLSLD16		BFD_RELOC_PPC64_PLTGOT16_DS	
BFD_RELOC_PPC_GOT_TLSLD16_HA		BFD_RELOC_PPC64_PLTGOT16_HA	
BFD_RELOC_PPC_GOT_TLSLD16_HI		BFD_RELOC_PPC64_PLTGOT16_HI	
BFD_RELOC_PPC_GOT_TLSLD16_L0		BFD_RELOC_PPC64_PLTG0T16_L0	
BFD_RELOC_PPC_GOT_TPREL16		BFD_RELOC_PPC64_PLTGOT16_LO_DS	
BFD_RELOC_PPC_GOT_TPREL16_HA		BFD_RELOC_PPC64_SECTOFF_DS	
BFD_RELOC_PPC_GOT_TPREL16_HI		BFD_RELOC_PPC64_SECTOFF_LO_DS	
BFD_RELOC_PPC_GOT_TPREL16_LO		BFD_RELOC_PPC64_TOC	
BFD_RELOC_PPC_JMP_SLOT		BFD_RELOC_PPC64_TOC16_DS	
BFD_RELOC_PPC_LOCAL24PC		BFD_RELOC_PPC64_TOC16_HA	
BFD_RELOC_PPC_RELATIVE		BFD_RELOC_PPC64_TOC16_HI	
BFD_RELOC_PPC_TLS		BFD_RELOC_PPC64_TOC16_LO	
BFD_RELOC_PPC_TLSGD		BFD_RELOC_PPC64_TOC16_LO_DS	
BFD_RELOC_PPC_TLSLD		BFD_RELOC_PPC64_TPREL16_DS	
BFD_RELOC_PPC_TOC16		BFD_RELOC_PPC64_TPREL16_HIGH	
BFD_RELOC_PPC_TPREL		BFD_RELOC_PPC64_TPREL16_HIGHA	
BFD_RELOC_PPC_TPREL16		BFD_RELOC_PPC64_TPREL16_HIGHER	
BFD_RELOC_PPC_TPREL16_HA		BFD_RELOC_PPC64_TPREL16_HIGHERA	
BFD_RELOC_PPC_TPREL16_HI		BFD_RELOC_PPC64_TPREL16_HIGHEST	
BFD_RELOC_PPC_TPREL16_L0		BFD_RELOC_PPC64_TPREL16_HIGHESTA	
BFD_RELOC_PPC_VLE_HA16A		BFD_RELOC_PPC64_TPREL16_LO_DS	
BFD_RELOC_PPC_VLE_HA16D		BFD_RELOC_RELC	
BFD_RELOC_PPC_VLE_HI16A		BFD_RELOC_RL78_16_0P	
BFD_RELOC_PPC_VLE_HI16D		BFD_RELOC_RL78_16U	
BFD_RELOC_PPC_VLE_L016A		BFD_RELOC_RL78_24_OP	
BFD_RELOC_PPC_VLE_L016D		BFD_RELOC_RL78_24U	
BFD_RELOC_PPC_VLE_REL15		BFD_RELOC_RL78_32_OP	
BFD_RELOC_PPC_VLE_REL24		BFD_RELOC_RL78_8U	
BFD_RELOC_PPC_VLE_REL8		BFD_RELOC_RL78_ABS16	
BFD_RELOC_PPC_VLE_RELO		BFD_RELOC_RL78_ABS16_REV	
BFD_RELOC_PPC_VLE_SDA21_LO		BFD_RELOC_RL78_ABS16U	
BFD_RELOC_PPC_VLE_SDAZI_LU		BFD_RELOC_RL78_ABS16UL	
BFD_RELOC_PPC_VLE_SDAREL_HA16A BFD_RELOC_PPC_VLE_SDAREL_HA16D		BFD_RELOC_RL78_ABS16UW	
		BFD_RELOC_RL78_ABS160W	
BFD_RELOC_PPC_VLE_SDAREL_HI16A			
BFD_RELOC_PPC_VLE_SDAREL_HI16DBFD_RELOC_PPC_VLE_SDAREL_L016A		BFD_RELOC_RL78_ABS32_REV	
BFD_RELOC_PPC_VLE_SDAREL_L016A BFD_RELOC_PPC_VLE_SDAREL_L016D		BFD_RELOC_RL78_ABS8	
BFD_RELOC_PPC64_ADDR16_DS		BFD_RELOC_RL78_CODE	99
BFD_RELUC_PPC64_ADDR16_DS	74		98
	14	DELL DELLIG BLACK DARAU PUBEL.	20

BFD_RELOC_RL78_GPRELB	98	BFD_RELOC_SH_CODE	79
BFD_RELOC_RL78_GPRELL	98	BFD_RELOC_SH_COPY	79
BFD_RELOC_RL78_GPRELW	98	BFD_RELOC_SH_COPY64	80
BFD_RELOC_RL78_HI16	98	BFD_RELOC_SH_COUNT	79
BFD_RELOC_RL78_HI8		BFD_RELOC_SH_DATA	79
BFD_RELOC_RL78_L016		BFD_RELOC_SH_DISP12	
BFD_RELOC_RL78_NEG16		BFD_RELOC_SH_DISP12BY2	
BFD_RELOC_RL78_NEG24		BFD_RELOC_SH_DISP12BY4	
BFD_RELOC_RL78_NEG32		BFD_RELOC_SH_DISP12BY8	
BFD_RELOC_RL78_NEG8		BFD_RELOC_SH_DISP20	
BFD_RELOC_RL78_OP_AND		BFD_RELOC_SH_DISP20BY8	
BFD_RELOC_RL78_OP_NEG		BFD_RELOC_SH_FUNCDESC	
BFD_RELOC_RL78_OP_SHRA		BFD_RELOC_SH_GLOB_DAT	
BFD_RELOC_RL78_OP_SUBTRACT		BFD_RELOC_SH_GLOB_DAT64	
BFD_RELOC_RL78_RELAX		BFD_RELOC_SH_GOT_HI16	
BFD_RELOC_RL78_SYM		BFD_RELOC_SH_GOT_LOW16	
BFD_RELOC_RVA		BFD_RELOC_SH_GOT_MEDHI16	
BFD_RELOC_RX_16_OP		BFD_RELOC_SH_GOT_MEDLOW16	
BFD_RELOC_RX_16U		BFD_RELOC_SH_GOT10BY4	
BFD_RELOC_RX_160	99		
		BFD_RELOC_SH_GOT10BY8	
BFD_RELOC_RX_24U		BFD_RELOC_SH_GOT20	
BFD_RELOC_RX_32_OP		BFD_RELOC_SH_GOTFUNCDESC	
BFD_RELOC_RX_8U		BFD_RELOC_SH_GOTFUNCDESC20	
BFD_RELOC_RX_ABS16		BFD_RELOC_SH_GOTOFF_HI16	
BFD_RELOC_RX_ABS16_REV		BFD_RELOC_SH_GOTOFF_LOW16	
BFD_RELOC_RX_ABS16U		BFD_RELOC_SH_GOTOFF_MEDHI16	
BFD_RELOC_RX_ABS16UL		BFD_RELOC_SH_GOTOFF_MEDLOW16	
BFD_RELOC_RX_ABS16UW		BFD_RELOC_SH_GOTOFF20	
BFD_RELOC_RX_ABS32		BFD_RELOC_SH_GOTOFFFUNCDESC	
BFD_RELOC_RX_ABS32_REV		BFD_RELOC_SH_GOTOFFFUNCDESC20	
BFD_RELOC_RX_ABS8		BFD_RELOC_SH_GOTPC	
BFD_RELOC_RX_DIFF		BFD_RELOC_SH_GOTPC_HI16	
BFD_RELOC_RX_DIR3U_PCREL	99	BFD_RELOC_SH_GOTPC_LOW16	80
BFD_RELOC_RX_GPRELB		BFD_RELOC_SH_GOTPC_MEDHI16	
BFD_RELOC_RX_GPRELL		BFD_RELOC_SH_GOTPC_MEDLOW16	80
BFD_RELOC_RX_GPRELW		BFD_RELOC_SH_GOTPLT_HI16	
BFD_RELOC_RX_NEG16	99	BFD_RELOC_SH_GOTPLT_LOW16	79
BFD_RELOC_RX_NEG24	99	BFD_RELOC_SH_GOTPLT_MEDHI16	
BFD_RELOC_RX_NEG32		BFD_RELOC_SH_GOTPLT_MEDLOW16	79
BFD_RELOC_RX_NEG8	99	BFD_RELOC_SH_GOTPLT10BY4	80
BFD_RELOC_RX_OP_NEG	99	BFD_RELOC_SH_GOTPLT10BY8	80
BFD_RELOC_RX_OP_SUBTRACT	99	BFD_RELOC_SH_GOTPLT32	80
BFD_RELOC_RX_RELAX	99	BFD_RELOC_SH_IMM_HI16	80
BFD_RELOC_RX_SYM		BFD_RELOC_SH_IMM_HI16_PCREL	
BFD_RELOC_SCORE_BCMP	102	BFD_RELOC_SH_IMM_LOW16	
BFD_RELOC_SCORE_BRANCH		BFD_RELOC_SH_IMM_LOW16_PCREL	
BFD_RELOC_SCORE_CALL15		BFD_RELOC_SH_IMM_MEDHI16	
BFD_RELOC_SCORE_DUMMY_HI16		BFD_RELOC_SH_IMM_MEDHI16_PCREL	
BFD_RELOC_SCORE_DUMMY2		BFD_RELOC_SH_IMM_MEDLOW16	
BFD_RELOC_SCORE_GOT_LO16		BFD_RELOC_SH_IMM_MEDLOW16_PCREL	
BFD_RELOC_SCORE_GOT15		BFD_RELOC_SH_IMM3	
BFD_RELOC_SCORE_GPREL15		BFD_RELOC_SH_IMM3U	
BFD_RELOC_SCORE_IMM30		BFD_RELOC_SH_IMM4	
BFD_RELOC_SCORE_IMM32		BFD_RELOC_SH_IMM4BY2	
BFD_RELOC_SCORE_JMP		BFD_RELOC_SH_IMM4BY4	
BFD_RELOC_SCORE16_BRANCH		BFD_RELOC_SH_IMM8	
BFD_RELOC_SCORE16_JMP		BFD_RELOC_SH_IMM8BY2	
	79		79
	1.21	DEL DELLE OF THEODIA	1.21

BFD_RELOC_SH_IMMS10		BFD_RELOC_SPARC_GOTDATA_OP_LOX10	
BFD_RELOC_SH_IMMS10BY2		BFD_RELOC_SPARC_H34	
BFD_RELOC_SH_IMMS10BY4		BFD_RELOC_SPARC_H44	
BFD_RELOC_SH_IMMS10BY8		BFD_RELOC_SPARC_HH22	
BFD_RELOC_SH_IMMS16		BFD_RELOC_SPARC_HIX22	
BFD_RELOC_SH_IMMS6		BFD_RELOC_SPARC_HM10	
BFD_RELOC_SH_IMMS6BY32		BFD_RELOC_SPARC_IRELATIVE	
BFD_RELOC_SH_IMMU16	80	BFD_RELOC_SPARC_JMP_IREL	63
BFD_RELOC_SH_IMMU5		BFD_RELOC_SPARC_JMP_SLOT	
BFD_RELOC_SH_IMMU6		BFD_RELOC_SPARC_L44	63
BFD_RELOC_SH_JMP_SLOT	79	BFD_RELOC_SPARC_LM22	63
BFD_RELOC_SH_JMP_SLOT64	80	BFD_RELOC_SPARC_LOX10	
BFD_RELOC_SH_LABEL		BFD_RELOC_SPARC_M44	
BFD_RELOC_SH_LOOP_END	79	BFD_RELOC_SPARC_OLO10	63
BFD_RELOC_SH_LOOP_START		BFD_RELOC_SPARC_PC_HH22	63
BFD_RELOC_SH_PCDISP12BY2	79	BFD_RELOC_SPARC_PC_HM10	63
BFD_RELOC_SH_PCDISP8BY2	79	BFD_RELOC_SPARC_PC_LM22	63
BFD_RELOC_SH_PCRELIMM8BY2	79	BFD_RELOC_SPARC_PC10	62
BFD_RELOC_SH_PCRELIMM8BY4		BFD_RELOC_SPARC_PC22	
BFD_RELOC_SH_PLT_HI16		BFD_RELOC_SPARC_PLT32	
BFD_RELOC_SH_PLT_LOW16		BFD_RELOC_SPARC_PLT64	
BFD_RELOC_SH_PLT_MEDHI16		BFD_RELOC_SPARC_REGISTER	
BFD_RELOC_SH_PLT_MEDLOW16		BFD_RELOC_SPARC_RELATIVE	
BFD_RELOC_SH_PT_16		BFD_RELOC_SPARC_REV32	
BFD_RELOC_SH_RELATIVE		BFD_RELOC_SPARC_SIZE32	
BFD_RELOC_SH_RELATIVE64		BFD_RELOC_SPARC_SIZE64	
BFD_RELOC_SH_SHMEDIA_CODE		BFD_RELOC_SPARC_TLS_DTPMOD32	
BFD_RELOC_SH_SWITCH16		BFD_RELOC_SPARC_TLS_DTPMOD64	
BFD_RELOC_SH_SWITCH32		BFD_RELOC_SPARC_TLS_DTPOFF32	
BFD_RELOC_SH_TLS_DTPMOD32		BFD_RELOC_SPARC_TLS_DTP0FF64	
BFD_RELOC_SH_TLS_DTP0FF32		BFD_RELOC_SPARC_TLS_GD_ADD	
BFD_RELOC_SH_TLS_GD_32		BFD_RELOC_SPARC_TLS_GD_CALL	
BFD_RELOC_SH_TLS_IE_32		BFD_RELOC_SPARC_TLS_GD_HI22	63
BFD_RELOC_SH_TLS_LD_32		BFD_RELOC_SPARC_TLS_GD_L010	
BFD_RELOC_SH_TLS_LDO_32		BFD_RELOC_SPARC_TLS_IE_ADD	
BFD_RELOC_SH_TLS_LE_32		BFD_RELOC_SPARC_TLS_IE_HI22	
BFD_RELOC_SH_TLS_TP0FF32		BFD_RELOC_SPARC_TLS_IE_LD	64
BFD_RELOC_SH_USES		BFD_RELOC_SPARC_TLS_IE_LDX	
BFD_RELOC_SIZE32		BFD_RELOC_SPARC_TLS_IE_LO10	
		BFD_RELOC_SPARC_TLS_LDM_ADD	
BFD_RELOC_SPARC_10		BFD_RELOC_SPARC_TLS_LDM_CALL	
BFD_RELOC_SPARC_11		BFD_RELOC_SPARC_TLS_LDM_HI22	
BFD_RELOC_SPARC_5		BFD_RELOC_SPARC_TLS_LDM_LO10	
BFD_RELOC_SPARC_6		BFD_RELOC_SPARC_TLS_LDO_ADD	
BFD_RELOC_SPARC_64		BFD_RELOC_SPARC_TLS_LDO_HIX22	
BFD_RELOC_SPARC_7		BFD_RELOC_SPARC_TLS_LDO_LOX10	
BFD_RELOC_SPARC_BASE13		BFD_RELOC_SPARC_TLS_LE_HIX22	
BFD_RELOC_SPARC_BASE22		BFD_RELOC_SPARC_TLS_LE_LOX10	
BFD_RELOC_SPARC_COPY		BFD_RELOC_SPARC_TLS_TPOFF32	
BFD_RELOC_SPARC_DISP64		BFD_RELOC_SPARC_TLS_TP0FF64	
BFD_RELOC_SPARC_GLOB_DAT		BFD_RELOC_SPARC_UA16	
BFD_RELOC_SPARC_GOT10		BFD_RELOC_SPARC_UA32	
BFD_RELOC_SPARC_GOT13		BFD_RELOC_SPARC_UA64	
BFD_RELOC_SPARC_GOT22		BFD_RELOC_SPARC_WDISP10	
BFD_RELOC_SPARC_GOTDATA_HIX22		BFD_RELOC_SPARC_WDISP16	
BFD_RELOC_SPARC_GOTDATA_LOX10		BFD_RELOC_SPARC_WDISP19	
BFD_RELOC_SPARC_GOTDATA_OP		BFD_RELOC_SPARC_WDISP22	
BFD RELOC SPARC GOTDATA OP HIX22	62	BFD RELOC SPARC WPLT30	62

BFD_RELOC_SPARC13	62	${\tt BFD\_RELOC\_TILEGX\_IMM16\_XO\_HWO\_TLS\_GD} \ 126$
BFD_RELOC_SPARC22	62	BFD_RELOC_TILEGX_IMM16_XO_HWO_TLS_IE 127
BFD_RELOC_SPU_ADD_PIC	64	BFD_RELOC_TILEGX_IMM16_X0_HW0_TLS_LE 126
BFD_RELOC_SPU_HI16	64	BFD_RELOC_TILEGX_IMM16_XO_HW1
BFD_RELOC_SPU_IMM10	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST 126
BFD_RELOC_SPU_IMM10W	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_GOT 126
BFD_RELOC_SPU_IMM16	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_PCREL
BFD_RELOC_SPU_IMM16W	64	
BFD_RELOC_SPU_IMM18	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_PLT_
BFD_RELOC_SPU_IMM7	64	PCREL
BFD_RELOC_SPU_IMM8	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_GD
BFD_RELOC_SPU_L016	64	
BFD_RELOC_SPU_PCREL16	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_IE
BFD_RELOC_SPU_PCREL9a	64	
BFD_RELOC_SPU_PCREL9b		BFD_RELOC_TILEGX_IMM16_XO_HW1_LAST_TLS_LE
BFD_RELOC_SPU_PPU32		
BFD_RELOC_SPU_PPU64	64	BFD_RELOC_TILEGX_IMM16_XO_HW1_PCREL 126
BFD_RELOC_THUMB_PCREL_BLX		BFD_RELOC_TILEGX_IMM16_XO_HW1_PLT_PCREL
BFD_RELOC_THUMB_PCREL_BRANCH12		
BFD_RELOC_THUMB_PCREL_BRANCH20		BFD_RELOC_TILEGX_IMM16_XO_HW2
BFD_RELOC_THUMB_PCREL_BRANCH23	76	
BFD_RELOC_THUMB_PCREL_BRANCH25	76	BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST 126
BFD_RELOC_THUMB_PCREL_BRANCH7		BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST_PCREL
BFD_RELOC_THUMB_PCREL_BRANCH9	76	
BFD_RELOC_TIC30_LDP		BFD_RELOC_TILEGX_IMM16_XO_HW2_LAST_PLT_
BFD_RELOC_TIC54X_16_0F_23	09	PCREL
BFD_RELOC_TIC54X_23	00	BFD_RELOC_TILEGX_IMM16_XO_HW2_PCREL 126
BFD_RELOC_TIC54X_MS7_OF_23		BFD_RELOC_TILEGX_IMM16_XO_HW2_PLT_PCREL
BFD_RELOC_TIC54X_PARTLS7		
BFD_RELOC_TIC54X_PARTMS9		BFD_RELOC_TILEGX_IMM16_XO_HW3
BFD_RELOC_TILEGX_BROFF_X1 1		BFD_RELOC_TILEGX_IMM16_XO_HW3_PCREL 126
BFD_RELOC_TILEGX_COPY		BFD_RELOC_TILEGX_IMM16_XO_HW3_PLT_PCREL
BFD_RELOC_TILEGX_DEST_IMM8_X1	25	
BFD_RELOC_TILEGX_GLOB_DAT 1	25	BFD_RELOC_TILEGX_IMM16_X1_HW0
BFD_RELOC_TILEGX_HWO		BFD_RELOC_TILEGX_IMM16_X1_HWO_GOT 126
BFD_RELOC_TILEGX_HWO_LAST 1		BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST 126
BFD_RELOC_TILEGX_HW1		${\tt BFD\_RELOC\_TILEGX\_IMM16\_X1\_HW0\_LAST\_GOT~126}$
BFD_RELOC_TILEGX_HW1_LAST		BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_PCREL
BFD_RELOC_TILEGX_HW2	25	
BFD_RELOC_TILEGX_HW2_LAST 1		BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_PLT_
BFD_RELOC_TILEGX_HW3		PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO		BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_GD
BFD_RELOC_TILEGX_IMM16_XO_HWO_GOT1		
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST 1		BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_IE
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_GOT 1		
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_PCREL	20	BFD_RELOC_TILEGX_IMM16_X1_HWO_LAST_TLS_LE
1	26	
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_PLT_	20	BFD_RELOC_TILEGX_IMM16_X1_HWO_PCREL 126
PCREL1	27	BFD_RELOC_TILEGX_IMM16_X1_HWO_PLT_PCREL
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_GD	21	
	26	BFD_RELOC_TILEGX_IMM16_X1_HWO_TLS_GD 126
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_IE	20	BFD_RELOC_TILEGX_IMM16_X1_HWO_TLS_IE 127
bru_reloc_filegx_fmffo_xo_nwo_lasf_fls_fe	27	BFD_RELOC_TILEGX_IMM16_X1_HW0_TLS_LE 126
BFD_RELOC_TILEGX_IMM16_XO_HWO_LAST_TLS_LE		BFD_RELOC_TILEGX_IMM16_X1_HW1
bru_reloc_filegx_friffo_xo_nwo_lasf_fls_le		BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST 126
BFD_RELOC_TILEGX_IMM16_XO_HWO_PCREL1		BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_GOT 126
BFD_RELOC_TILEGX_IMM16_XO_HWO_PLT_PCREL		BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_PCREL
brb_reloc_filegx_fmffo_xo_nwo_fil_rcref		126
	40	120

BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_PLT_	BFD_RELOC_TILEPRO_DEST_IMM8_X1	
PCREL	BFD_RELOC_TILEPRO_GLOB_DAT	
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_GD	BFD_RELOC_TILEPRO_IMM16_X0	
126	BFD_RELOC_TILEPRO_IMM16_XO_GOT	
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_IE	BFD_RELOC_TILEPRO_IMM16_XO_GOT_HA	
	BFD_RELOC_TILEPRO_IMM16_XO_GOT_HI	124
BFD_RELOC_TILEGX_IMM16_X1_HW1_LAST_TLS_LE	BFD_RELOC_TILEPRO_IMM16_XO_GOT_LO	124
	BFD_RELOC_TILEPRO_IMM16_XO_HA	123
BFD_RELOC_TILEGX_IMM16_X1_HW1_PCREL 126	BFD_RELOC_TILEPRO_IMM16_XO_HA_PCREL	124
BFD_RELOC_TILEGX_IMM16_X1_HW1_PLT_PCREL	BFD_RELOC_TILEPRO_IMM16_XO_HI	
	BFD_RELOC_TILEPRO_IMM16_XO_HI_PCREL	
BFD_RELOC_TILEGX_IMM16_X1_HW2	BFD_RELOC_TILEPRO_IMM16_XO_LO	
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST 126	BFD_RELOC_TILEPRO_IMM16_XO_LO_PCREL	
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST_PCREL	BFD_RELOC_TILEPRO_IMM16_XO_PCREL	
	BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD	
BFD_RELOC_TILEGX_IMM16_X1_HW2_LAST_PLT_	BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_HA	
PCREL	BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_HI	
BFD_RELOC_TILEGX_IMM16_X1_HW2_PCREL 126	BFD_RELOC_TILEPRO_IMM16_XO_TLS_GD_LO	
BFD_RELOC_TILEGX_IMM16_X1_HW2_PLT_PCREL	BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE	
	BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_HA	
BFD_RELOC_TILEGX_IMM16_X1_HW3	BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_HI	
BFD_RELOC_TILEGX_IMM16_X1_HW3_PCREL 126	BFD_RELOC_TILEPRO_IMM16_XO_TLS_IE_LO	
BFD_RELOC_TILEGX_IMM16_X1_HW3_PLT_PCREL	BFD_RELOC_TILEPRO_IMM16_XO_TLS_LE	
	BFD_RELOC_TILEPRO_IMM16_XO_TLS_LE_HA	
BFD_RELOC_TILEGX_IMM8_X0	BFD_RELOC_TILEPRO_IMM16_XO_TLS_LE_HI	
BFD_RELOC_TILEGX_IMM8_XO_TLS_ADD	BFD_RELOC_TILEPRO_IMM16_XO_TLS_LE_LO	
BFD_RELOC_TILEGX_IMM8_XO_TLS_GD_ADD 127	BFD_RELOC_TILEPRO_IMM16_X1	
BFD_RELOC_TILEGX_IMM8_X1	BFD_RELOC_TILEPRO_IMM16_X1_GOT	
BFD_RELOC_TILEGX_IMM8_X1_TLS_ADD	BFD_RELOC_TILEPRO_IMM16_X1_GOT_HA	
BFD_RELOC_TILEGX_IMM8_X1_TLS_GD_ADD		
	BFD_RELOC_TILEPRO_IMM16_X1_GOT_HI	
BFD_RELOC_TILEGX_IMM8_YO	BFD_RELOC_TILEPRO_IMM16_X1_GOT_LO	
BFD_RELOC_TILEGX_IMM8_YO_TLS_ADD 127	BFD_RELOC_TILEPRO_IMM16_X1_HA	
BFD_RELOC_TILEGX_IMM8_YO_TLS_GD_ADD 127	BFD_RELOC_TILEPRO_IMM16_X1_HA_PCREL	
BFD_RELOC_TILEGX_IMM8_Y1	BFD_RELOC_TILEPRO_IMM16_X1_HI	
BFD_RELOC_TILEGX_IMM8_Y1_TLS_ADD 127	BFD_RELOC_TILEPRO_IMM16_X1_HI_PCREL	
BFD_RELOC_TILEGX_IMM8_Y1_TLS_GD_ADD 127	BFD_RELOC_TILEPRO_IMM16_X1_LO	
BFD_RELOC_TILEGX_JMP_SLOT	BFD_RELOC_TILEPRO_IMM16_X1_LO_PCREL	
BFD_RELOC_TILEGX_JUMPOFF_X1	BFD_RELOC_TILEPRO_IMM16_X1_PCREL	
BFD_RELOC_TILEGX_JUMPOFF_X1_PLT 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD	
BFD_RELOC_TILEGX_MF_IMM14_X1	BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_HA	
BFD_RELOC_TILEGX_MMEND_XO	BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_HI	
BFD_RELOC_TILEGX_MMSTART_XO	BFD_RELOC_TILEPRO_IMM16_X1_TLS_GD_LO	
BFD_RELOC_TILEGX_MT_IMM14_X1	BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE	124
BFD_RELOC_TILEGX_RELATIVE 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_HA	124
BFD_RELOC_TILEGX_SHAMT_XO 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_HI	124
BFD_RELOC_TILEGX_SHAMT_X1 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_IE_LO	124
BFD_RELOC_TILEGX_SHAMT_YO 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE	125
BFD_RELOC_TILEGX_SHAMT_Y1 125	BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_HA	125
BFD_RELOC_TILEGX_TLS_DTPMOD32127	BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_HI	125
BFD_RELOC_TILEGX_TLS_DTPMOD64127	BFD_RELOC_TILEPRO_IMM16_X1_TLS_LE_LO	125
BFD_RELOC_TILEGX_TLS_DTPOFF32127	BFD_RELOC_TILEPRO_IMM8_XO	
BFD_RELOC_TILEGX_TLS_DTP0FF64	BFD_RELOC_TILEPRO_IMM8_XO_TLS_GD_ADD	
BFD_RELOC_TILEGX_TLS_GD_CALL	BFD_RELOC_TILEPRO_IMM8_X1	
BFD_RELOC_TILEGX_TLS_IE_LOAD	BFD_RELOC_TILEPRO_IMM8_X1_TLS_GD_ADD	
BFD_RELOC_TILEGX_TLS_TPOFF32	BFD_RELOC_TILEPRO_IMM8_YO	
BFD_RELOC_TILEGX_TLS_TPOFF64	BFD_RELOC_TILEPRO_IMM8_YO_TLS_GD_ADD	
BFD_RELOC_TILEPRO_BROFF_X1	BFD_RELOC_TILEPRO_IMM8_Y1	
BED RELOC THEFPRO COPY 123	RED RELOC TILEPRO IMMS V1 TIS CD ADD	

DED DELOG MELEDDO TAD GLOW	DED DELOG HOLD ED LE CONTROLE
BFD_RELOC_TILEPRO_JMP_SLOT 123	BFD_RELOC_V850_TDA_7_8_OFFSET90
BFD_RELOC_TILEPRO_JOFFLONG_X1	BFD_RELOC_V850_ZDA_15_16_OFFSET 89
BFD_RELOC_TILEPRO_JOFFLONG_X1_PLT 123	BFD_RELOC_V850_ZDA_16_16_0FFSET
BFD_RELOC_TILEPRO_MF_IMM15_X1	BFD_RELOC_V850_ZDA_16_16_SPLIT_OFFSET 90
BFD_RELOC_TILEPRO_MMEND_XO 124	BFD_RELOC_VAX_GLOB_DAT112
BFD_RELOC_TILEPRO_MMEND_X1 124	BFD_RELOC_VAX_JMP_SLOT112
BFD_RELOC_TILEPRO_MMSTART_XO	BFD_RELOC_VAX_RELATIVE112
BFD_RELOC_TILEPRO_MMSTART_X1	BFD_RELOC_VPE4KMATH_DATA 103
BFD_RELOC_TILEPRO_MT_IMM15_X1	BFD_RELOC_VPE4KMATH_INSN 103
BFD_RELOC_TILEPRO_RELATIVE	BFD_RELOC_VTABLE_ENTRY
BFD_RELOC_TILEPRO_SHAMT_XO	BFD_RELOC_VTABLE_INHERIT
BFD_RELOC_TILEPRO_SHAMT_X1	BFD_RELOC_X86_64_32S
BFD_RELOC_TILEPRO_SHAMT_YO	BFD_RELOC_X86_64_COPY
BFD_RELOC_TILEPRO_SHAMT_Y1	BFD_RELOC_X86_64_DTPMOD64
BFD_RELOC_TILEPRO_TLS_DTPMOD32 125	BFD_RELOC_X86_64_DTPOFF32
BFD_RELOC_TILEPRO_TLS_DTPOFF32	BFD_RELOC_X86_64_DTP0FF64
BFD_RELOC_TILEPRO_TLS_GD_CALL	BFD_RELOC_X86_64_GLOB_DAT
BFD_RELOC_TILEPRO_TLS_IE_LOAD	BFD_RELOC_X86_64_GOT3271
BFD_RELOC_TILEPRO_TLS_TPOFF32	BFD_RELOC_X86_64_GOT64
bfd_reloc_type_lookup	BFD_RELOC_X86_64_GOTOFF64
BFD_RELOC_V850_16_GOT91	BFD_RELOC_X86_64_GOTPC32 72
BFD_RELOC_V850_16_GOTOFF 91	BFD_RELOC_X86_64_GOTPC32_TLSDESC 72
BFD_RELOC_V850_16_PCREL90	BFD_RELOC_X86_64_GOTPC64 72
BFD_RELOC_V850_16_S191	BFD_RELOC_X86_64_GOTPCREL 72
BFD_RELOC_V850_16_SPLIT_OFFSET91	BFD_RELOC_X86_64_GOTPCREL6472
BFD_RELOC_V850_17_PCREL90	BFD_RELOC_X86_64_GOTPLT64
BFD_RELOC_V850_22_PCREL89	BFD_RELOC_X86_64_GOTTPOFF
BFD_RELOC_V850_22_PLT_PCREL 91	BFD_RELOC_X86_64_IRELATIVE
BFD_RELOC_V850_23	BFD_RELOC_X86_64_JUMP_SLOT
BFD_RELOC_V850_32_ABS	BFD_RELOC_X86_64_PC32_BND
BFD_RELOC_V850_32_GOT	BFD_RELOC_X86_64_PLT3271
BFD_RELOC_V850_32_G0T0FF	BFD_RELOC_X86_64_PLT32_BND
BFD_RELOC_V850_32_GOTPCREL 91	BFD_RELOC_X86_64_PLT0FF64
	BFD_RELOC_X86_64_RELATIVE
BFD_RELOC_V850_32_PCREL	
BFD_RELOC_V850_32_PLT_PCREL	BFD_RELOC_X86_64_TLSDESC
BFD_RELOC_V850_9_PCREL	BFD_RELOC_X86_64_TLSDESC_CALL
BFD_RELOC_V850_ALIGN	BFD_RELOC_X86_64_TLSGD
BFD_RELOC_V850_CALLT_15_16_OFFSET 91	BFD_RELOC_X86_64_TLSLD
BFD_RELOC_V850_CALLT_16_16_OFFSET 90	BFD_RELOC_X86_64_TPOFF32 72
BFD_RELOC_V850_CALLT_6_7_OFFSET 90	BFD_RELOC_X86_64_TP0FF64 72
BFD_RELOC_V850_CODE	BFD_RELOC_XC16X_PAG 111
BFD_RELOC_V850_COPY 91	BFD_RELOC_XC16X_POF
BFD_RELOC_V850_DATA 91	BFD_RELOC_XC16X_SEG 111
BFD_RELOC_V850_GLOB_DAT91	BFD_RELOC_XC16X_SOF
BFD_RELOC_V850_JMP_SLOT91	BFD_RELOC_XGATE_24
BFD_RELOC_V850_L016_S1	BFD_RELOC_XGATE_GPAGE
BFD_RELOC_V850_L016_SPLIT_OFFSET 90	BFD_RELOC_XGATE_IMM3106
BFD_RELOC_V850_LONGCALL90	BFD_RELOC_XGATE_IMM4
BFD_RELOC_V850_LONGJUMP90	BFD_RELOC_XGATE_IMM5106
BFD_RELOC_V850_RELATIVE91	BFD_RELOC_XGATE_IMM8_HI 106
BFD_RELOC_V850_SDA_15_16_0FFSET	BFD_RELOC_XGATE_IMM8_LO
BFD_RELOC_V850_SDA_16_16_0FFSET	BFD_RELOC_XGATE_L016
BFD_RELOC_V850_SDA_16_16_SPLIT_OFFSET 90	BFD_RELOC_XGATE_PCREL_10
BFD_RELOC_V850_TDA_16_16_0FFSET90	BFD_RELOC_XGATE_PCREL_9
BFD_RELOC_V850_TDA_4_4_0FFSET	BFD_RELOC_XGATE_RL_GROUP
BFD_RELOC_V850_TDA_4_5_0FFSET	BFD_RELOC_XGATE_RL_JUMP
BFD_RELOC_V850_TDA_6_8_0FFSET	BFD_RELOC_XSTORMY16_12
BFD RELOC V850 TDA 7 7 OFFSET	BFD_RELOC_ASTORMY16_12
DED RELUC VOOU IDA / / UPPOEI	DED RELUC ADJURTITO 24

	1.01
BFD_RELOC_XSTORMY16_FPTR16 111	bfd_scan_vma
BFD_RELOC_XSTORMY16_REL_12111	bfd_seach_for_target
BFD_RELOC_XTENSA_ASM_EXPAND	bfd_section_already_linked
BFD_RELOC_XTENSA_ASM_SIMPLIFY 115	bfd_section_list_clear
BFD_RELOC_XTENSA_DIFF16 114	bfd_sections_find_if
BFD_RELOC_XTENSA_DIFF32 114	bfd_set_arch_info
BFD_RELOC_XTENSA_DIFF8	bfd_set_archive_head
BFD_RELOC_XTENSA_GLOB_DAT 114	bfd_set_assert_handler
BFD_RELOC_XTENSA_JMP_SLOT	bfd_set_default_target
BFD_RELOC_XTENSA_OPO	bfd_set_error
BFD_RELOC_XTENSA_OP1	bfd_set_error_handler
BFD_RELOC_XTENSA_OP2	bfd_set_error_program_name
BFD_RELOC_XTENSA_PLT	bfd_set_file_flags
BFD_RELOC_XTENSA_RELATIVE	bfd_set_format
BFD_RELOC_XTENSA_RTLD	bfd_set_gp_size
BFD_RELOC_XTENSA_SLOTO_ALT	bfd_set_private_flags
BFD_RELOC_XTENSA_SLOTO_OP	bfd_set_reloc
BFD_RELOC_XTENSA_SLOT1_ALT	bfd_set_section_contents
BFD_RELOC_XTENSA_SLOT1_OP	bfd_set_section_flags
BFD_RELOC_XTENSA_SLOT10_ALT	bfd_set_section_size
BFD_RELOC_XTENSA_SLOT10_OP	bfd_set_start_address
BFD_RELOC_XTENSA_SLOT11_ALT	bfd_set_symtab
BFD_RELOC_XTENSA_SLOT11_OP 114	bfd_symbol_info
BFD_RELOC_XTENSA_SLOT12_ALT	bfd_target_list
BFD_RELOC_XTENSA_SLOT12_OP 114	bfd_write_bigendian_4byte_int 162
BFD_RELOC_XTENSA_SLOT13_ALT	bfd_zalloc
BFD_RELOC_XTENSA_SLOT13_OP 114	bfd_zalloc2
BFD_RELOC_XTENSA_SLOT14_ALT	
BFD_RELOC_XTENSA_SLOT14_OP 114	
BFD_RELOC_XTENSA_SLOT2_ALT 114	$\mathbf{C}$
BFD_RELOC_XTENSA_SLOT2_OP 114	coff_symbol_type
BFD_RELOC_XTENSA_SLOT3_ALT	core_file_matches_executable_p
BFD_RELOC_XTENSA_SLOT3_OP 114	
BFD_RELOC_XTENSA_SLOT4_ALT 114	
BFD_RELOC_XTENSA_SLOT4_OP 114	F
BFD_RELOC_XTENSA_SLOT5_ALT 114	
BFD_RELOC_XTENSA_SLOT5_OP	find_separate_debug_file
BFD_RELOC_XTENSA_SLOT6_ALT 114	
BFD_RELOC_XTENSA_SLOT6_OP	G
BFD_RELOC_XTENSA_SLOT7_ALT 114	G
BFD_RELOC_XTENSA_SLOT7_OP 114	generic_core_file_matches_executable_p 130
BFD_RELOC_XTENSA_SLOT8_ALT	
BFD_RELOC_XTENSA_SLOT8_OP	TT
BFD_RELOC_XTENSA_SLOT9_ALT	H
BFD_RELOC_XTENSA_SLOT9_OP	Hash tables
BFD_RELOC_XTENSA_TLS_ARG	114311 (40)165
BFD_RELOC_XTENSA_TLS_CALL	
BFD_RELOC_XTENSA_TLS_DTPOFF	I
BFD_RELOC_XTENSA_TLS_DIFFOFF	internal object-file format
	internal object-file format
BFD_RELOC_XTENSA_TLS_TPOFF	
BFD_RELOC_XTENSA_TLSDESC_ARG	$\mathbf{L}$
BFD_RELOC_XTENSA_TLSDESC_FN	
BFD_RELOC_Z80_DISP8	Linker
BFD_RELOC_Z8K_CALLR	
BFD_RELOC_Z8K_DISP7	0
BFD_RELOC_Z8K_IMM4L	U
bfd_rename_section	Other functions
bfd_scan_arch	

S separate_alt_debug_file_exists	target vector (_bfd_link_add_symbols) 167 target vector (_bfd_link_hash_table_create) 167 The HOWTO Macro
${f T}$	$\mathbf{W}$
target vector (_bfd_final_link) 169	what is it? 1

The body of this manual is set in cmr10 at 10.95pt, with headings in cmb10 at 10.95pt and examples in cmtt10 at 10.95pt. cmti10 at 10.95pt and cmsl10 at 10.95pt are used for emphasis.