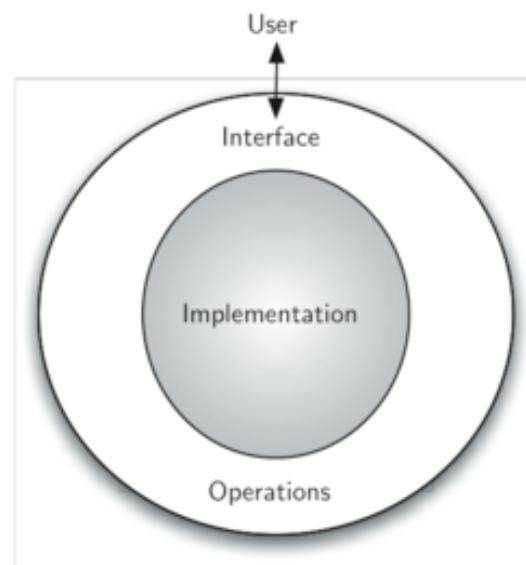


# Abstract Data Types (ADTs)

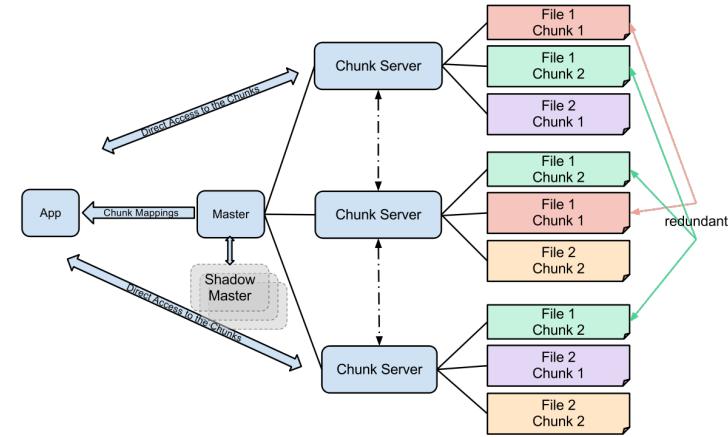


# Abstraction



- Abstraction: hiding unnecessary low-level details
  - An abstraction is a simplified view of an entity, which omits *unimportant* details
- An abstraction can go wrong if
  - It includes details that are not important
    - Increases the cognitive load of the programmer
  - It omits important details
    - Programmer doesn't have necessary information to use the module, class, function, etc.

# Abstraction



- A file system is an abstraction
  - Omits many details, such as the mechanism for choosing which blocks on a storage device to use for the data in a given file
  - Most users don't care about these details
  - Some of the details of a file system's implementation are important to some users
    - Most file systems cache data in main memory, and they may delay writing new data to the storage device in order to improve performance.
    - Some applications, such as databases, need to know exactly when data is written through to storage, so they can ensure that data will be preserved after system crashes.
    - Thus, the rules for flushing data to secondary storage must be visible in the file system's interface.
  - Abstraction should provide appropriate detail to appropriate users

# Control Abstraction

- Control abstraction (procedural abstraction)
  - A procedure (method) implements the details of an algorithm
  - One part of abstraction: **signature**, provides name, parameter types, return type.
    - E.g., **int binarySearch(int[] a, int key)**
    - Another part: **specification**, provides details about behavior and effects
  - Reasoning about code connects implementation to specification

# Data Abstraction

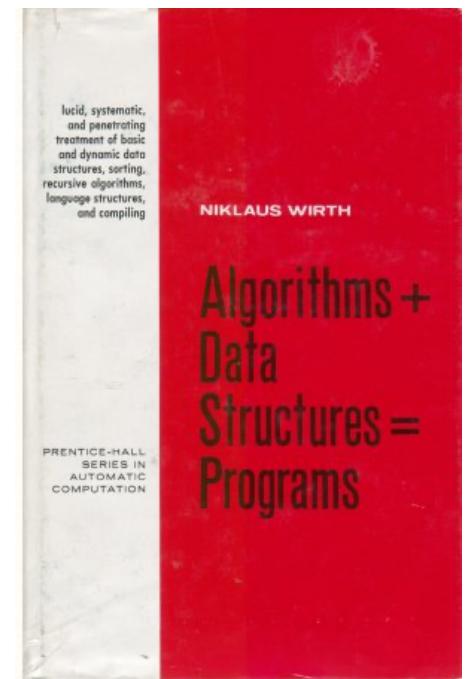
- Data abstraction
  - Types: abstract away from the details of data representation
  - E.g., type **String** is an abstraction
  - E.g., C type **struct Person** is an abstraction
- Abstract Data Type (ADT): high-level data abstraction
  - The ADT is operations + state
  - A specification mechanism
  - A way of thinking about programs and design

# Abstract Data Types are Important

- ADTs are about organizing and manipulating data
- Organizing and manipulating data is pervasive.
  - Inventing and describing algorithms is not
- Start your design by **thinking about the data model.**
  - What is the data?
  - Consider what operations are needed.
  - Implementation
    - Choose data structures carefully.
      - Will they allow necessary operations to be efficient?
    - Write code to access and manipulate data.

*Data dominates. If you've chosen the right data structures and organized things well, the algorithms will almost always be self-evident. Data structures, not algorithms, are central to programming.*

~ Rob Pike



# ADT is a way of thinking about programs and design

- From **domain concept**
  - E.g., the math concept of a polynomial, a set of integers, the concept of a library item, etc.
    - abstract
- through **ADT**
  - Describes domain concept in terms of **specification fields** and **abstract operations**
    - Abstract
    - Specification fields – what are the important pieces of information?
    - Operations – what manipulations will be needed
- to **implementation**
  - Implements ADT with representation fields and concrete operations
    - Concrete
  - Abstraction is how programmers deal with complexity

# Example: Polynomial with Integer Coefficients - ADT

ADT:

Overview description:

A Poly is an **immutable polynomial with integer coefficients**. A Poly is:

$$c_0 + c_1x + c_2x^2 + \dots$$

- Specification fields  $c_0, c_1$  etc.

Set of abstract operations:

**add, mul, eval, etc.** with PSoft style specs referencing **abstract** specification fields

# Example: Polynomial with Integer Coefficients

## - Implementation

```
class Poly {  
    // rep. invariant: d = coeffs.length-1  
    private int d; // degree of the polynomial  
    private int[] coeffs; // coefficients  
    // operations add, sub, mul,  
    // eval, in terms of rep. fields coeffs, d.  
}
```

```
class Poly {  
    // rep. invariant: ...  
    private List<Term> terms; // terms of poly  
    // operations add, sub, mul, eval, etc. in terms of terms  
    // term - object specifying coefficient and exponent  
}
```

# Another Example: A Meeting: Domain Concept & ADT

ADT:

Overview description:

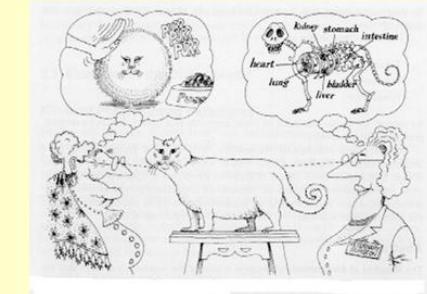
**An appointment for a meeting.**

**With date, room, attendees**

Specification fields.

**Set of abstract operations:**

e.g., **addAttendee**, **scheduleRoom**, etc. with PSoft  
style specs referencing **abstract** specification fields



# Why ADTs?

- Bridges gap between **domain concept** and **implementation**
- Formalizes domain concept, provides basis for reasoning about correctness of the implementation
- **Shields client from implementation. Implementation can vary without affecting client!**
  - Information hiding
  - Client code interacts with object through a set of abstract operations

# An ADT has a Set of Operations

- Operations act on data
- ADT is about the **meaning** of data
- ADT is about the **use** of data
- Data representation (implementation) may change without causing problems for the client.

```
class RightTriangle {  
    float base, altitude;  
}
```

```
class RightTriangle {  
    float base, angle;  
}
```

- Instead, think of a type as a **set of operations**: `create`, `getBase`, `getAltitude`, `getBottomAngle`, etc.
- Force clients to call operations to access data
  - Client doesn't need to know the details of representation

# Are These Types Same or Different?

```
class Point {  
    float x;  
    float y;  
}
```

```
class Point {  
    float r;  
    float theta;  
}
```

- They are **different**! Different implementations
- They are the **same**! Both implement the concept of a 2D point.
- Goal of ADT methodology is to **express sameness**
  - Clients depend only on the set of operations: x(), y(), r(), theta(), etc.
  - Data representation can be changed: to change algorithms, to fix bugs, etc.

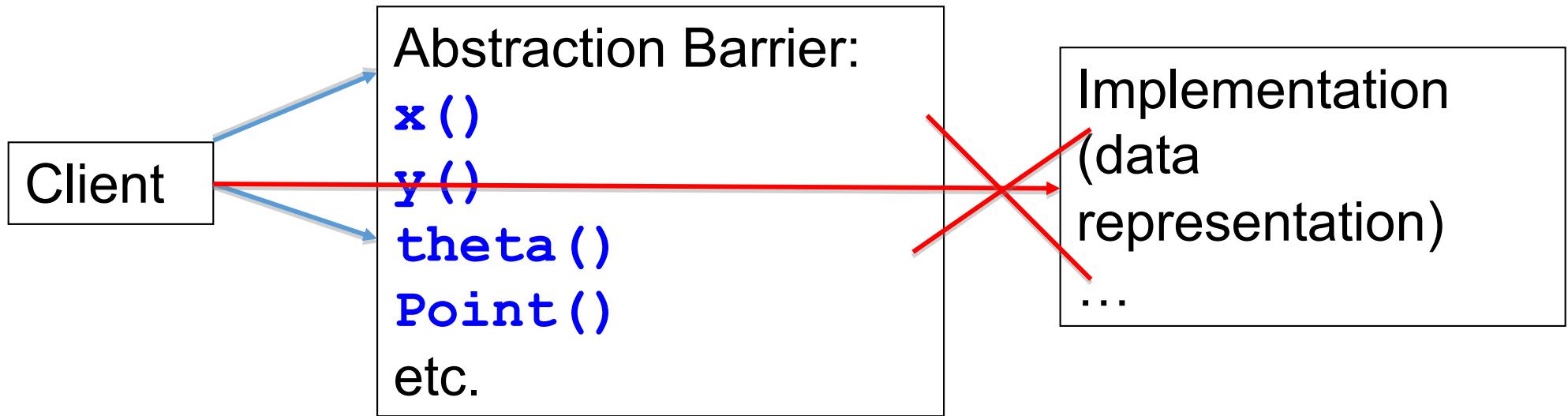
# Are These Types Same or Different?

```
class Poly {  
    private int d;  
    private int[] coeffs;  
}
```

```
class Poly {  
    private List<Term> terms;  
}
```

- Same ADT
- Clients depend only on the set of operations: `add(Poly)`, `mul(Poly)`, etc.

# Abstraction Barrier



ADT provides an **abstraction barrier**  
Clients access the ADT through its operations.  
They never access the data representation.

# ADT Methods

- We group the access methods of an ADT into
  - Creators
    - Create a brand new object
    - Constructors are creators
  - Observers
    - Return information about *this* object
  - Producers
    - Return a new object of *this* type by performing operations on the current object
    - Poly add(Poly p) – return a new Poly, this + p
  - Mutators
    - Change the current object

# 2D Point abstract methods

```
class Point {  
    // A 2D point in the plane  
    public float x();  
    public float y();  
    public float r();  
    public float theta();
```

Observers

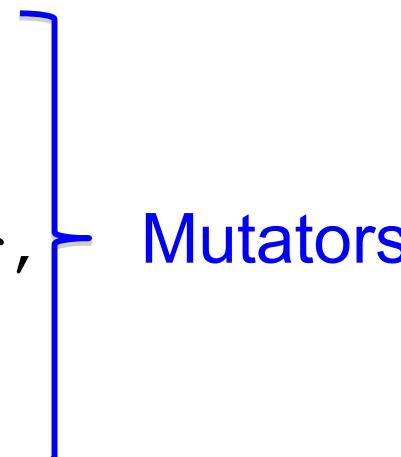
```
// ... can be created  
public Point(); // (0,0)  
public Point(float x, float y);  
  
public Point swapXY();
```

Creators

Producer

# 2D Point as an ADT

```
// class Point continued  
...  
  
// ... can be moved  
public void translate(float delta_x,  
                      float delta_y);  
public void scaleAndRotate(float delta_r,  
                           float delta_theta);  
}  
}
```



Mutators

# Specifying an ADT

**immutable**

**class TypeName**

1. overview
2. specification fields
3. creators
4. observers
5. producers
6. no ~~mutators~~

**mutable**

**class TypeName**

1. overview
2. specification fields
3. creators
4. observers
5. producers (**mutators are more common**)
6. mutators

# Poly, an immutable datatype: overview

```
/**  
 * A Poly is an immutable polynomial with  
 * integer coefficients. A Poly is:  
 *       $c_0 + c_1x + c_2x^2 + \dots$   
 */  
class Poly {
```

Abstract state (specification fields).  
More on this later.

**Overview:** Always state whether **mutable** or **immutable**  
Define abstract model for use in specification of  
operations. In ADTs, state is **abstract**, not concrete

I.e., these are NOT actual, implementation fields of Poly,  
just what we call specification fields.

Poly, an immutable datatype: **creators**

```
// modifies: none
// effects: makes a new Poly = 0
public Poly()

// modifies: none
// effect: makes a new Poly = cxn
// throws: NegExponentException if n < 0
public Poly(int c, int n)
```

**Creators:** This is an example of *overloading*, two **Poly** constructors with different signatures.

New object is part of effects, not preexisting state. Hence, modifies is none.

# Poly, an immutable datatype: **observers**

```
// returns: degree of this polynomial
public int degree()

// returns: the coefficient of the term of
// this polynomial, whose power is d
public int coeff(int d)
```

**Observers:** Used to obtain information about *this* polynomial.

Return values of **other types**.

Observers should **not** modify the abstract state!

**this**: the current Poly object. Also called the **receiver**

```
Poly x = new Poly(...) // creator
int c = x.coeff(3); // observer
```

# Poly, an immutable datatype: **producers**

```
// modifies: none
// returns: a new Poly with value this + q
public Poly add(Poly q)

// modifies: none
// returns: a new Poly with value this * q
public Poly mul(Poly q)
```

**Producers:** Operations on a type that create other objects of the same type. Common in immutable types.

**Should have no side effects**

i.e., should not change abstract values of any existing object

# Poly, an immutable datatype: **Mutators**

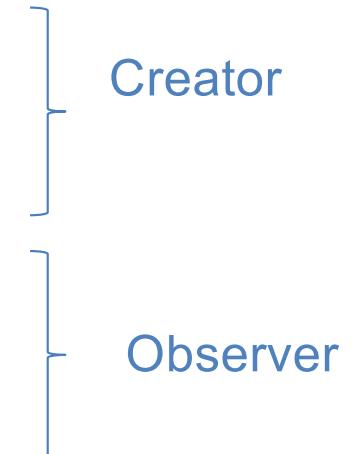
- Poly is an immutable ADT
- It has no mutators

# IntSet, a mutable datatype: overview, creators and observers

```
/*
 * Overview: An IntSet is a mutable,
 * unbounded set of integers. E.g.,
 * { x1, x2 , ... xn } with no duplicates
 */
class IntSet {

    // effects: makes a new empty IntSet
    public IntSet()

    // returns: true if x in this IntSet,
    //           else false
    public boolean contains(int x)
```



# IntSet, a mutable datatype: **mutators**

```
// modifies: this
// effects: thispost = thispre U { x }
public void add(int x)

// modifies: this
// effects: thispost = thispre - { x }
public void remove(int x)
```



Mutators

**Mutators:** operations that modify receiver, ***this***. Rarely modify anything other than ***this***.

**effects:** describe how ***this*** changes.

Often, mutators have no return value.  
Although they may return true or false to indicate success

# Exercise: String, an immutable datatype

- Overview?
  - an immutable sequence of characters, example “abc”
  - don’t say array or list of characters
- Creators?
  - **String(), String(char[] value), String(String original), ...**
- Observers?
  - **charAt, compareTo, contains, endsWith, ...**
- Producers?
  - **concat, format, substring, ...**
- Mutators?
  - None!

# Exercise: The Stack datatype

**public Stack()**

**public boolean empty()**

**public E peek()**

**public void push(E item)**

**public E pop()**

# ADTs and Java Language Features

- Java classes
  - ADT operations are **public**
  - Other operations are private
  - Clients can only access ADT operations
- Java interfaces
  - Clients only see the ADT operations
  - Cannot see fields
  - Implementing classes must provide all operations

# Reasoning About ADTs

- ADT is a specification, a set of operations
  - E.g., contains(), add(), etc., (the IntSet ADT)
  - add(Poly q), mul(Poly q), etc., (the Poly ADT)
- When **specifying** ADTs, there is no mention of data representation!
  - Basic information hiding principle
- When **implementing** the ADT, we must select a specific data representation
- Reasoning connects implementation to specification:
  - is our implementation correct?

# Implementation of an ADT is Provided by a Class

- To implement the ADT
  - We must select the **representation**
  - Implement concrete operations in terms of that rep
  - E.g., the **rep** of our Poly can be
    - a) `int[] coeffs` or
    - b) `List<Term> terms`
- Choose representation such that
  - It is possible to implement all operations
  - Most frequently used operations are efficient

# Connecting Implementation to Specification

- **Representation invariant:** Object → Boolean
  - Maps the object to a true or false value
  - Indicates whether data representation is **well-formed**.
    - If rep is not-well formed, operations may produce incorrect results
  - Only well-formed representations are meaningful
  - Defines the set of **valid** values
- **Abstraction function:** Object → abstract value
  - What the data structure really **means**
    - E.g., array [2, 3, -1] represents  $-x^2 + 3x + 2$
  - How the data structure is to be interpreted

# The Representation Invariant

- States data structure **well-formedness**
  - E.g., IntSet objects, whose data array contains duplicates are not valid values
- Must hold before and after every method!
  - Not necessarily in the middle of a method.
  - Correctness of operation implementation (methods in the class) depends on it!

# IntSet ADT

```
/** Abstraction: An IntSet is a mutable set
 * of integers with no duplicates.
 * E.g., { x1, x2 , ... xn }, {}.
 */
// rep invariant: data has no duplicates
// effects: makes a new empty IntSet
public IntSet();

// modifies: this
// effects: thispost = thispre U { x }
public void add(Integer x);

// modifies: this
// effects: thispost = thispre - { x }
public void remove(Integer x);

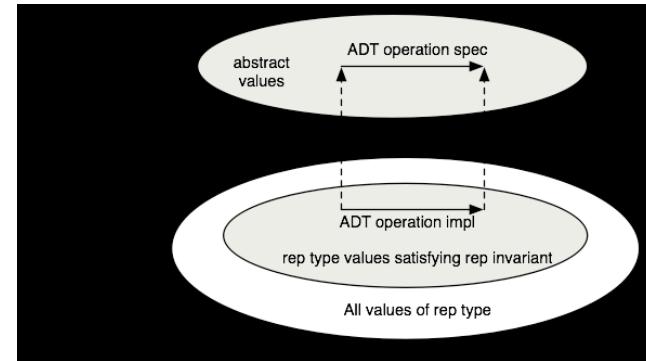
// returns: (x in this)
public boolean contains(Integer x)

// returns: cardinality of this
public int size()
```

# One Possible Implementation

```
class IntSet {  
    private List<Integer> data  
        = new ArrayList<Integer>();  
    public void add(Integer x) { data.add(x); }  
    public void remove(Integer x) {  
        data.remove(x);  
    }  
    public boolean contains(Integer x) {  
        return data.contains(x)  
    }  
    public int size() { return data.size(); }  
}
```

# The Representation Invariant



- ```
s = new IntSet();
s.add(1); s.add(1); s.remove(1);
System.out.println(s.contains(1));
```

- What is wrong with this code?
- Representation invariant tells us

```
class IntSet {
    // Rep invariant: data has no duplicates
    private List<Integer> data; ...
```

# The Representation Invariant

```
class IntSet {  
    // Rep invariant:  
    // data no duplicates  
    private List<Integer> data  
        = new ArrayList<Integer>();  
  
    public void add(Integer x) {  
        data.add(x);  
    }  
    // Rep invariant does not hold after second  
    // add!  
}
```

# The Representation Invariant

```
class IntSet {  
    // Rep invariant:  
    // data has no duplicates  
    private List<Integer> data  
        = new ArrayList<Integer>();  
    public void add(Integer x) {  
        if (!data.contains(x))  
            data.add(x);  
    }  
    // Rep invariant now holds after add  
}
```

# The Representation Invariant

- Rep invariant excludes values that do not correspond to abstract values

```
class LineSegment {  
    //Rep invariant: !(x1 == x2 && y1 == y2)  
    private float x1, y1; // start point  
    private float x2, y2; // end point
```

- Conceptually, a line segment is defined by two distinct points.
- Thus, values with the same start and end point (e.g.,  $x1=x2=1$ ,  $y1=y2=2$ ), are meaningless.
- Rep invariant excludes them

# The Representation Invariant

- Rep invariant excludes data representation values that do not correspond to **abstract values**

```
class RightTriangle {  
    // Rep invariant: 0 < angle < 90 &  
    // base > 0  
    float base, angle;  
    // Objects that don't meet the above constraints are  
    // not right triangles
```

## Additionally...

- Rep invariant states constraints imposed by specific data structures and algorithms
  - E.g., Tree has no cycles, array must be sorted, at most one term per exponent in Poly (if we choose the term representation)
- Rep invariant states **constraints** between fields that are synchronized with each other
  - E.g., **degree** and **coefficients** fields in Poly (if we choose the array data representation)
- In general, **rep invariant states correctness constraints**
  - if not met, implementation may behave incorrectly!

# Rep Invariant Example

```
class Account {  
    // Rep invariant:  
    // balance >= 0  
    // balance = Σitransactions.get(i).amount  
    // transactions != null  
    // no nulls in transactions  
  
    private int balance;  
    // history of transactions  
    private List<Transaction> transactions;  
  
    ...  
}
```

# More Rep Invariant Examples

```
class Polynomial {  
    //Rep. invariant: degree = coeffs.length-1  
    // Poly is the sum of coeffs[i]*x^i for i ranging from 0 to  
    // degree  
    private int degree;  
    private int[] coeffs;  
    // operations add, sub, mul, eval, etc.  
}
```

```
class Polynomial {  
    //Rep. invariant: for any 0 <= i,j < terms.size(), i != j =>  
    // termsi.getExp() != termsj.getExp()  
    // Poly is sum of terms...  
    private List<Term> terms;  
    // operations add, sub, mul, eval, etc.  
}
```