

Reachability Analysis and Revision of Dynamics of Biological Regulatory Networks

Xinwei Chai

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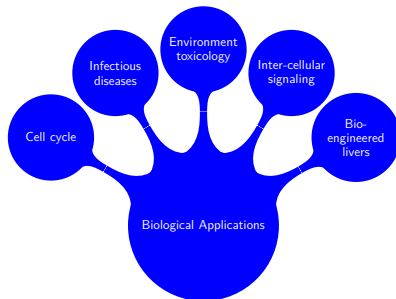
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Pascale LE GALL, Professeur des universités, Centrale Supélec

Examineurs : Béatrice DUVAL, Professeur des universités, Université d'Angers
Loïc PAULEVÉ, Chargé de recherche, LaBRI, UMR CNRS

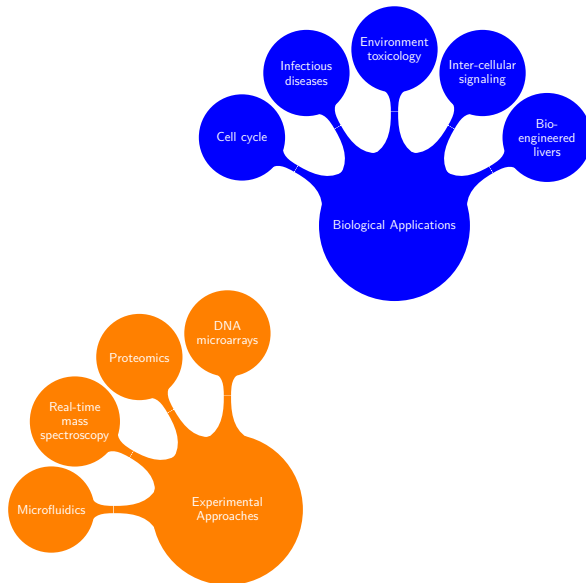
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Co-encadrant de thèse : Morgan MAGNIN, Professeur des universités, École Centrale de Nantes

May 24, 2019

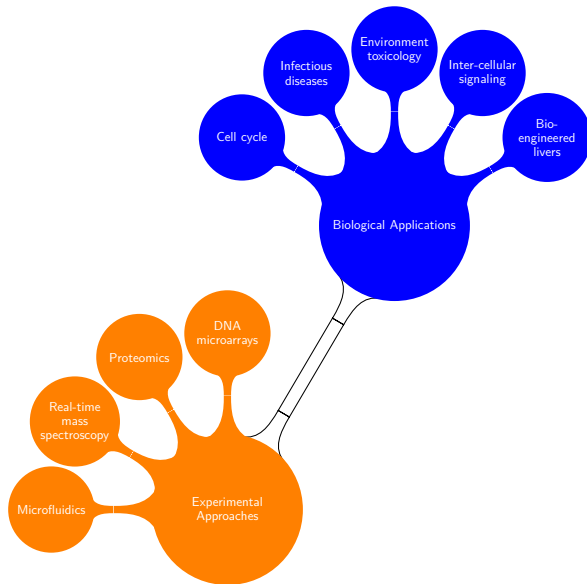
Positioning of Our Work



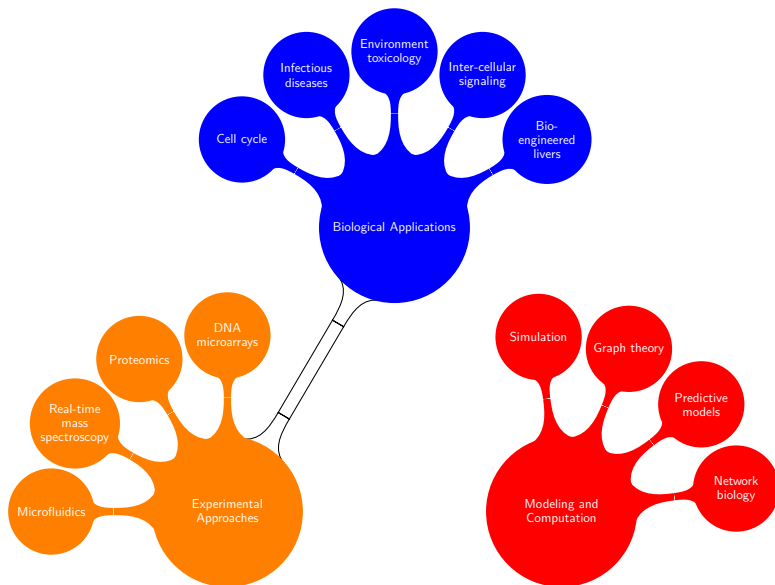
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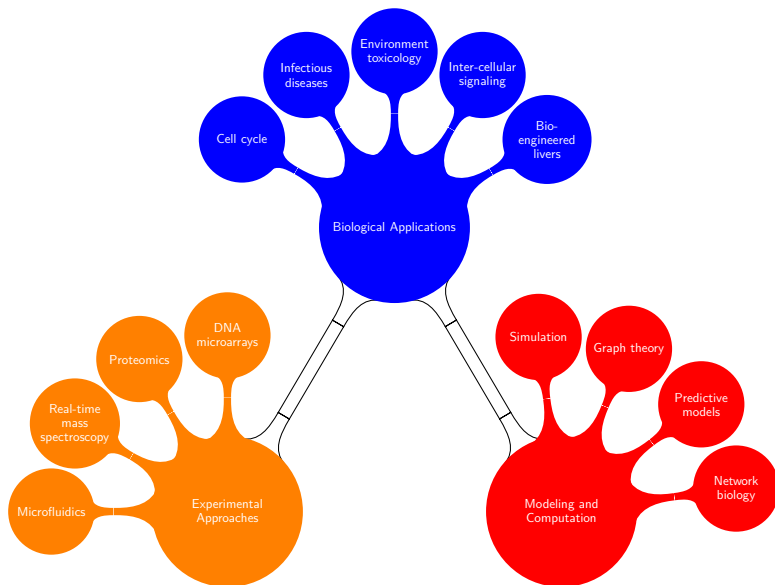
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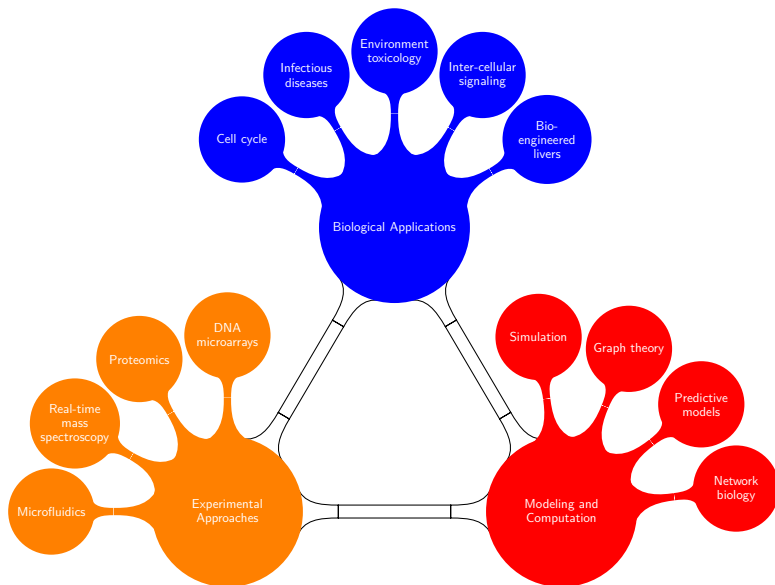
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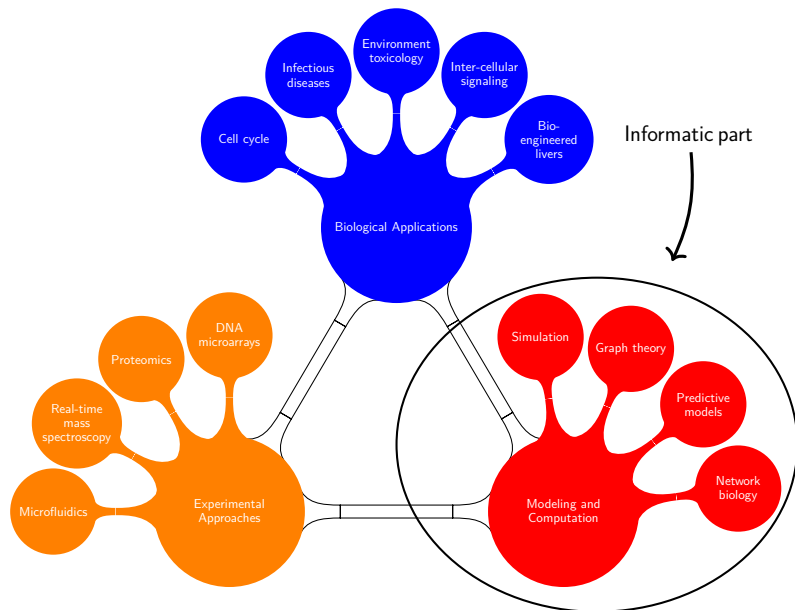
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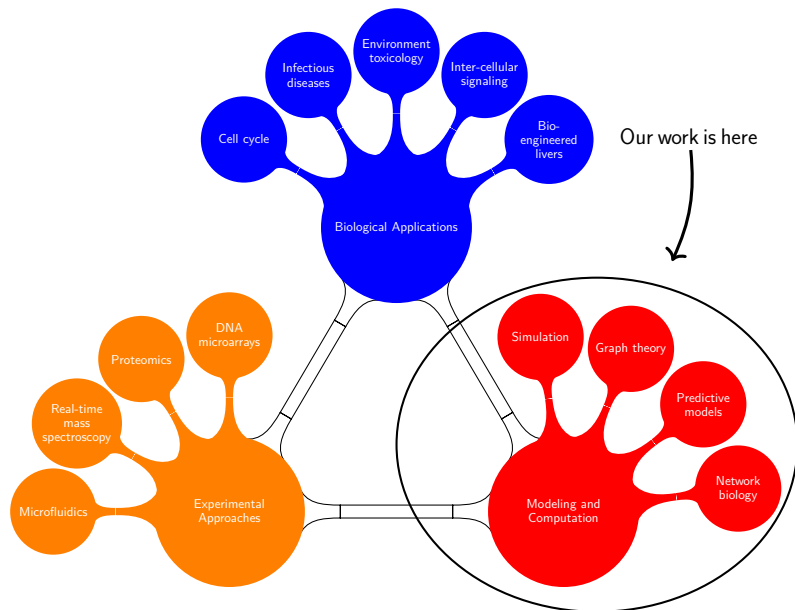
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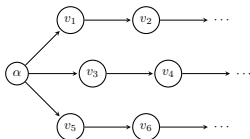


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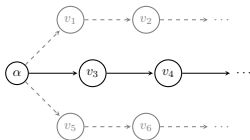
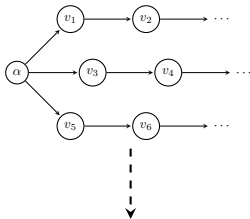
Problematics

Real system dynamics



Problematics

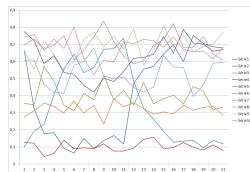
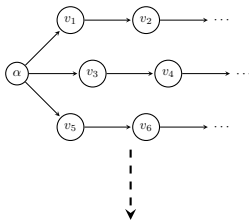
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Partial observation

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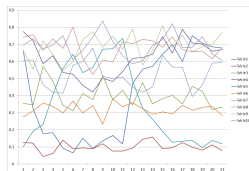
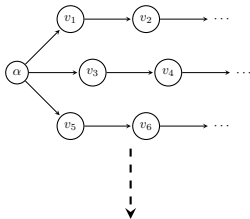
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Time series data

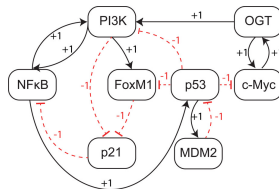
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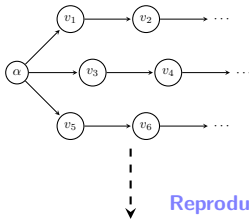
Learning



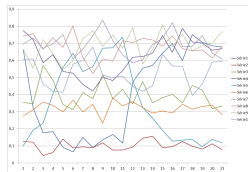
Biological network

Problematics

Real system dynamics

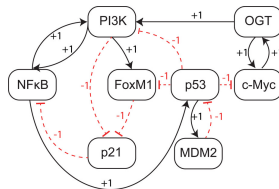


Reproducible?



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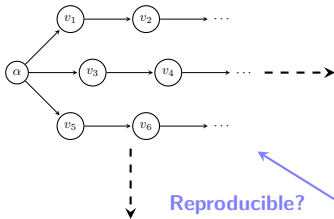


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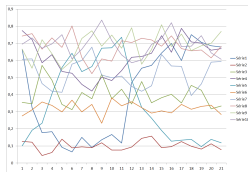
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Biological *a priori* knowledge

Real system dynamics

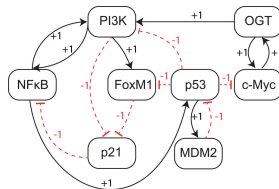


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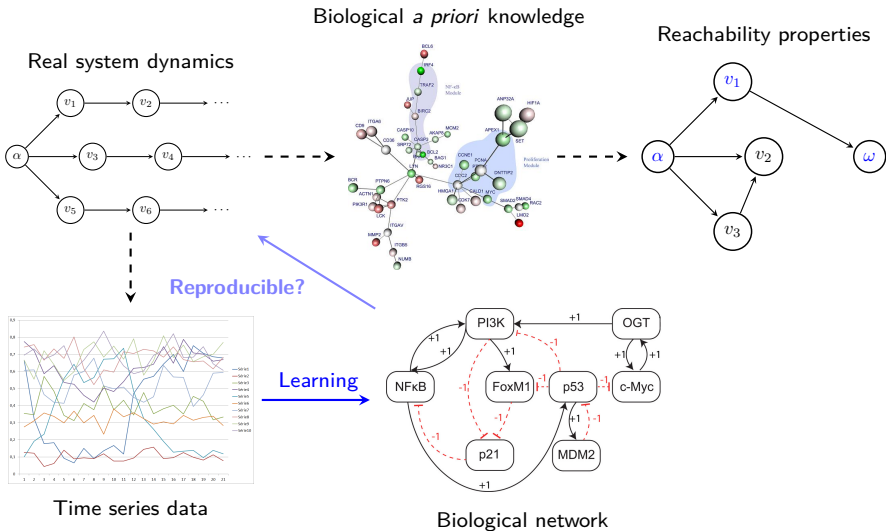
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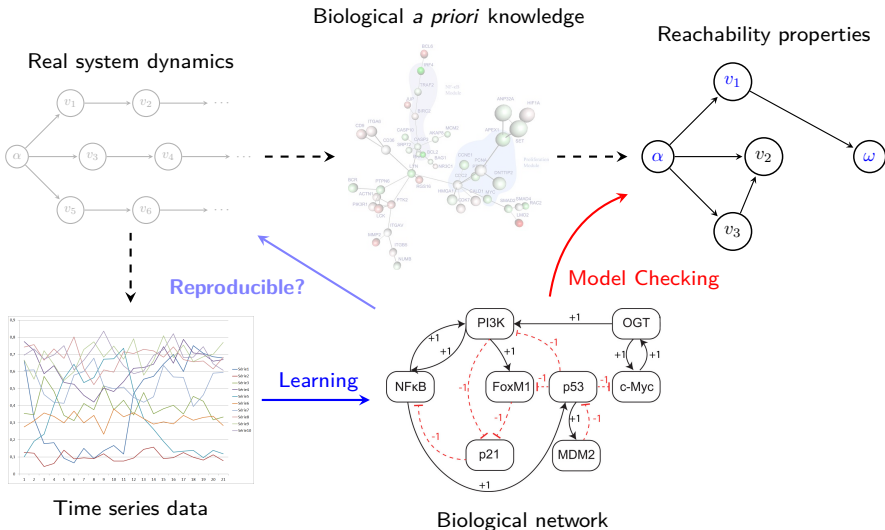


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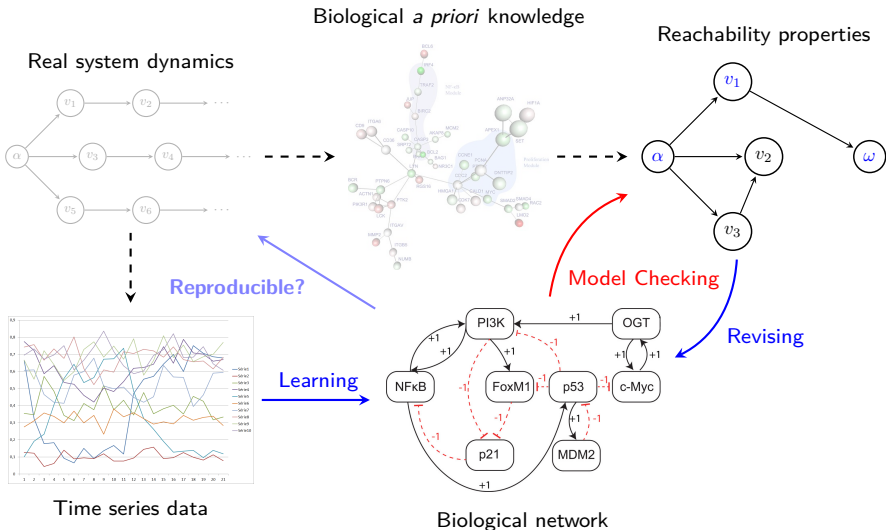
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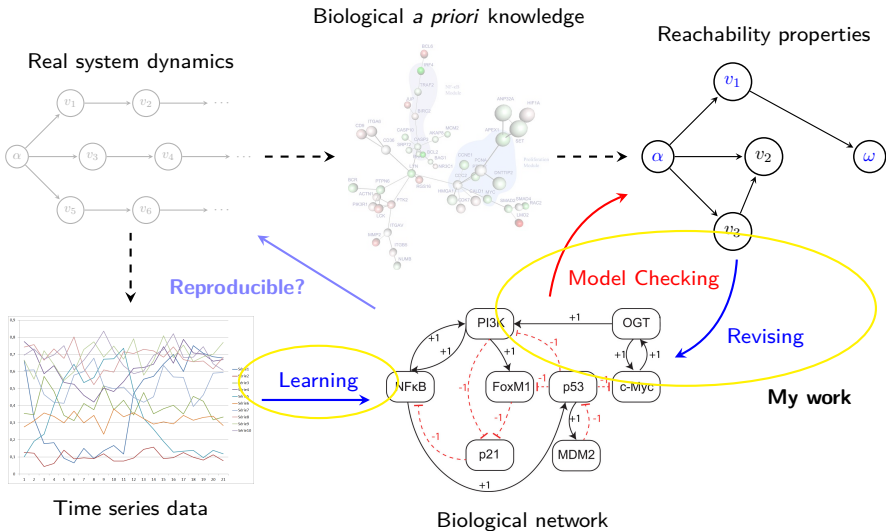
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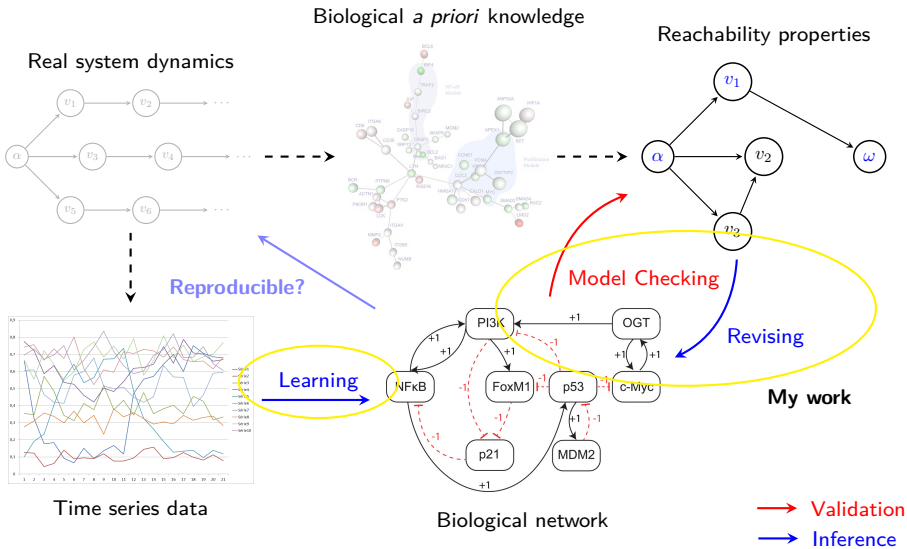
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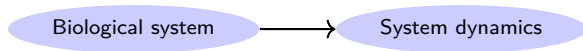
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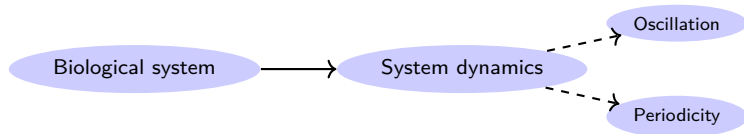
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Outline: Reachability problem (model checking) → model learning → model revising based on reachability properties

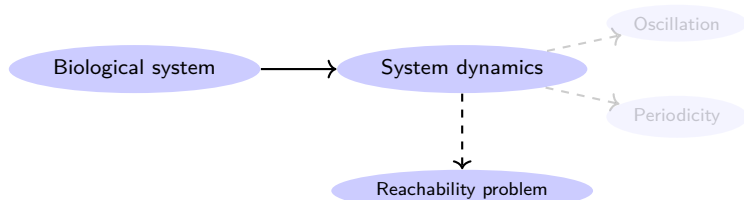
Problematic of Reachability Problem



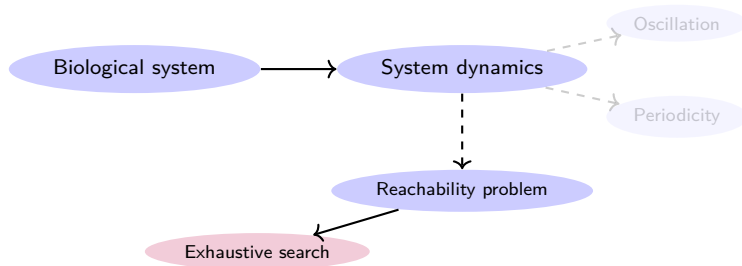
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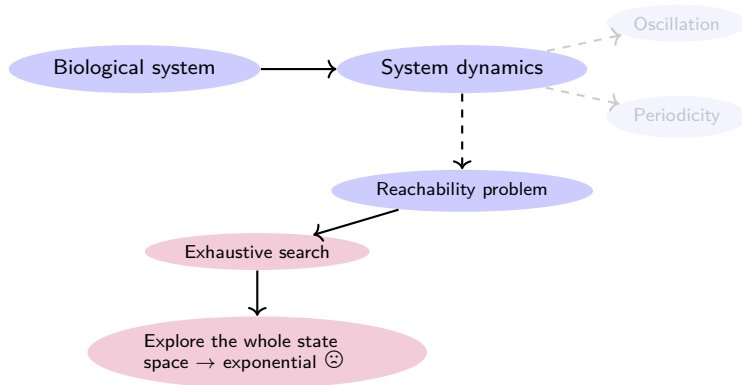
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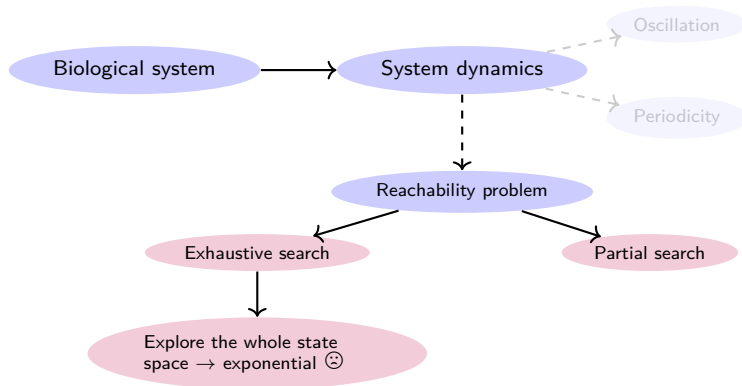
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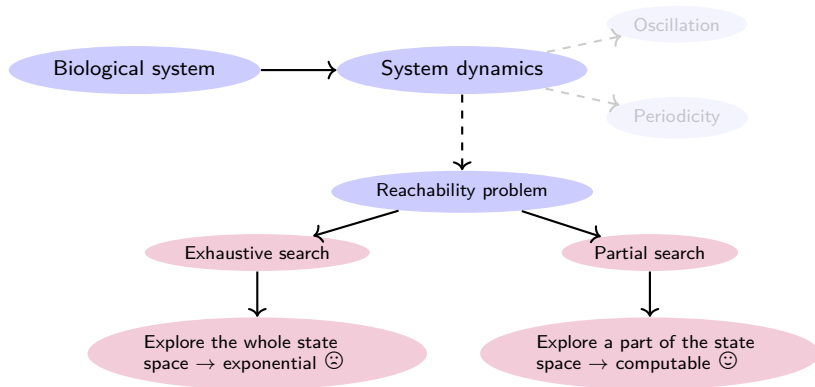
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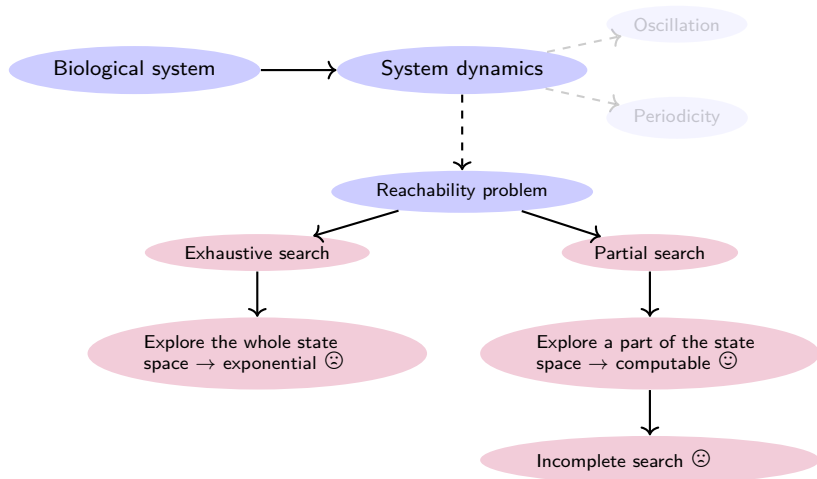
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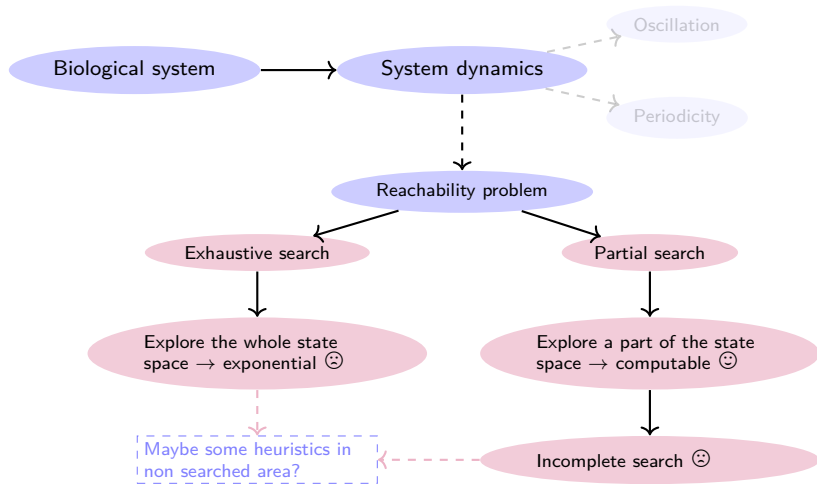
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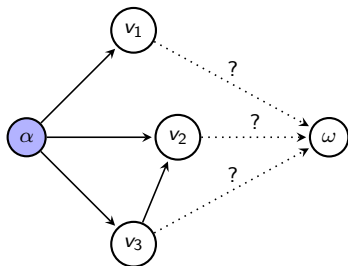
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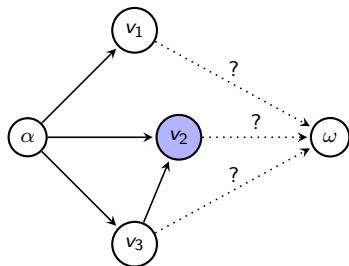


Reachability Problem Illustrated by Transition Graph (Exhaustive)



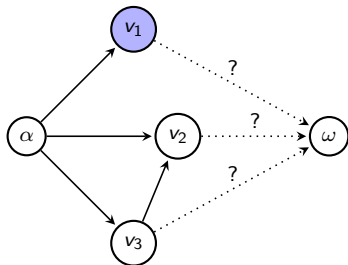
- A digraph representing state space
 - Nodes = system states
 - Edges = state transitions
 - α = initial state
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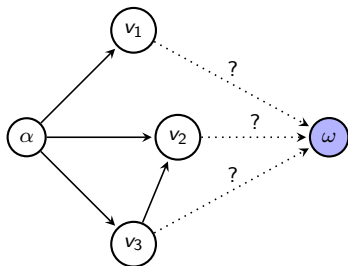
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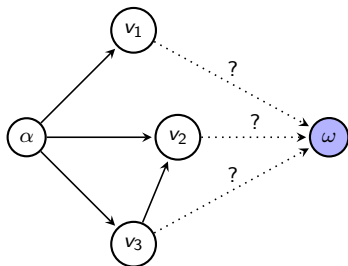
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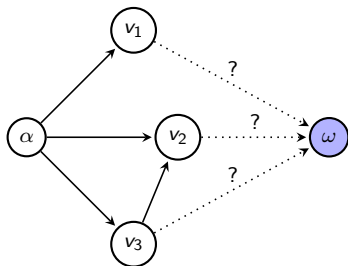
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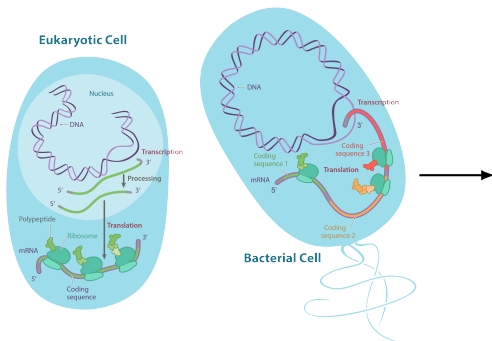
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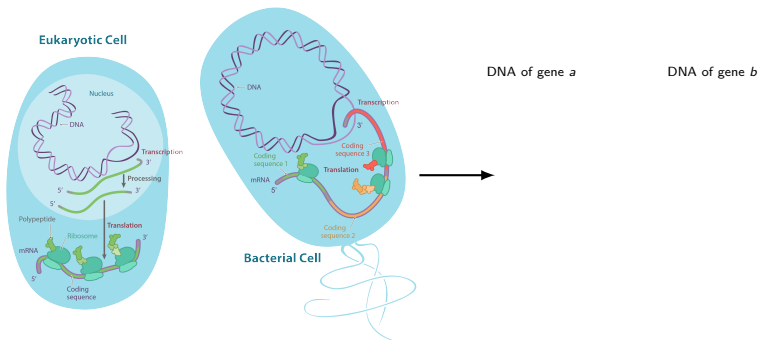


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- A pertinent modeling framework is necessary to describe system dynamics

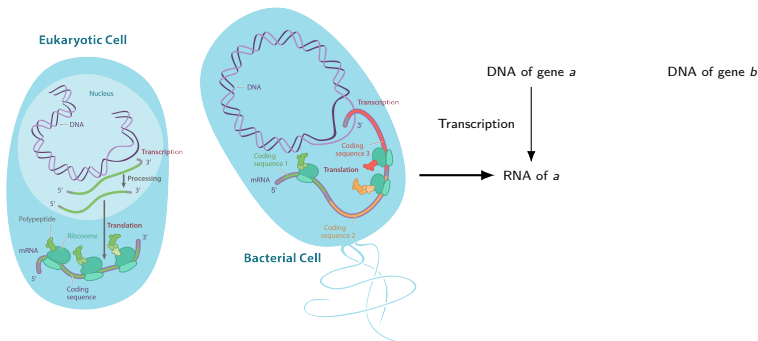
Biological Regulatory Networks (BRN)



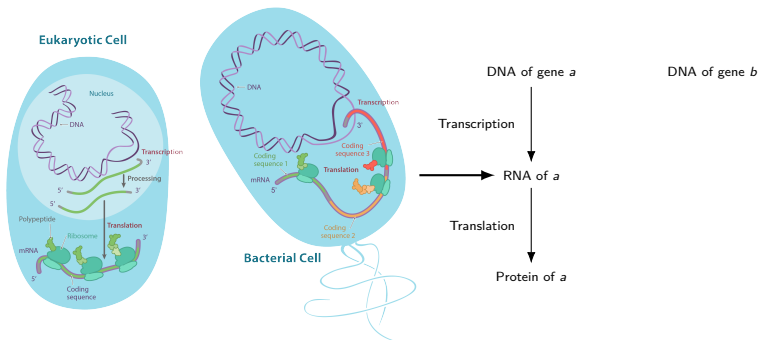
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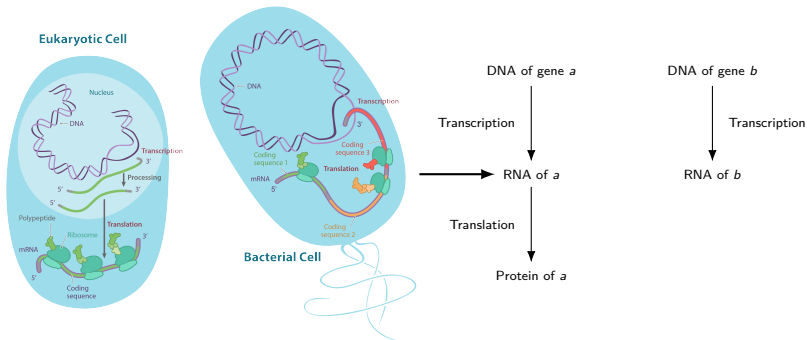
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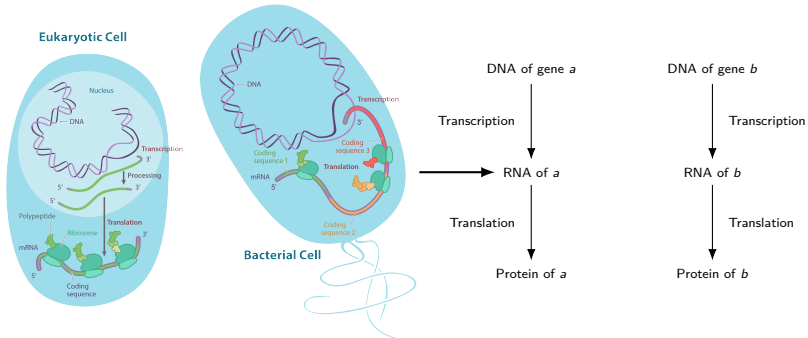
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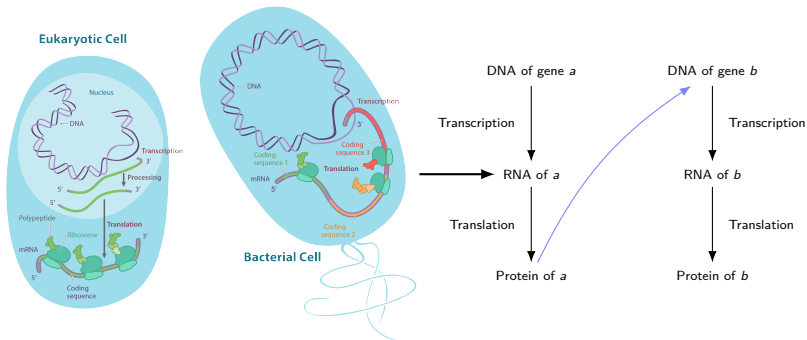
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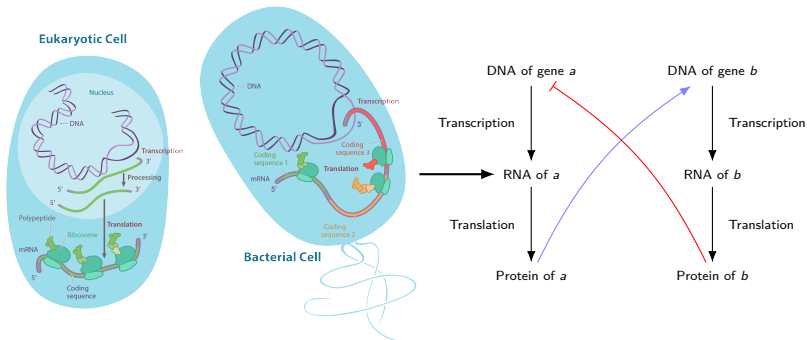
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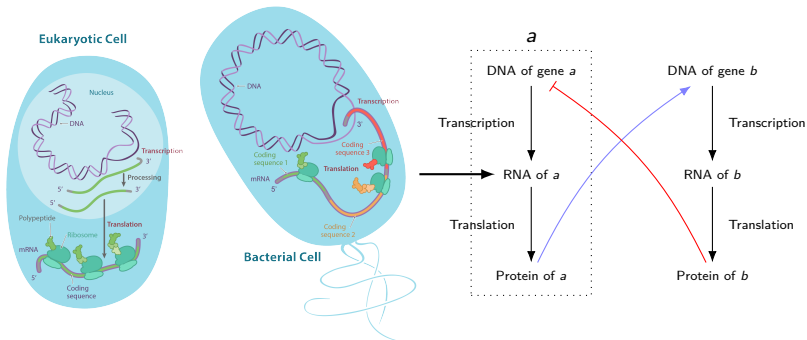
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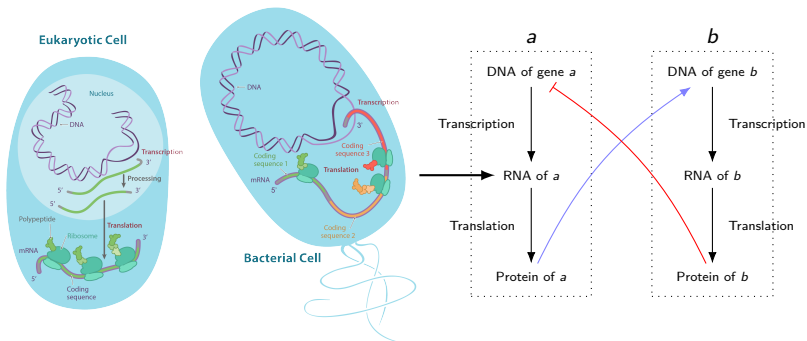
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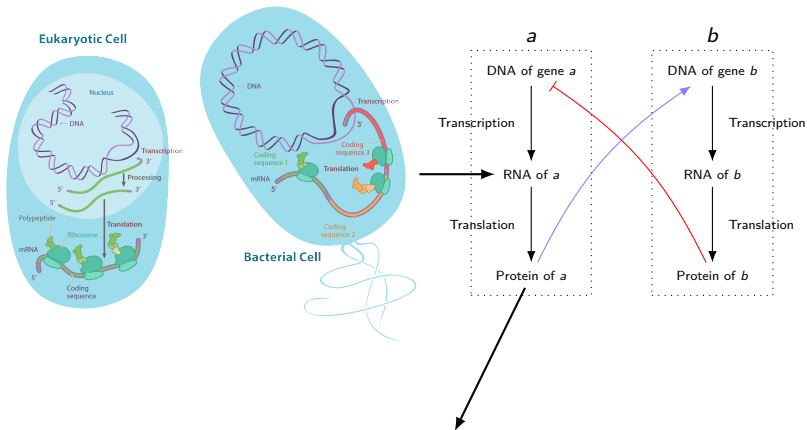
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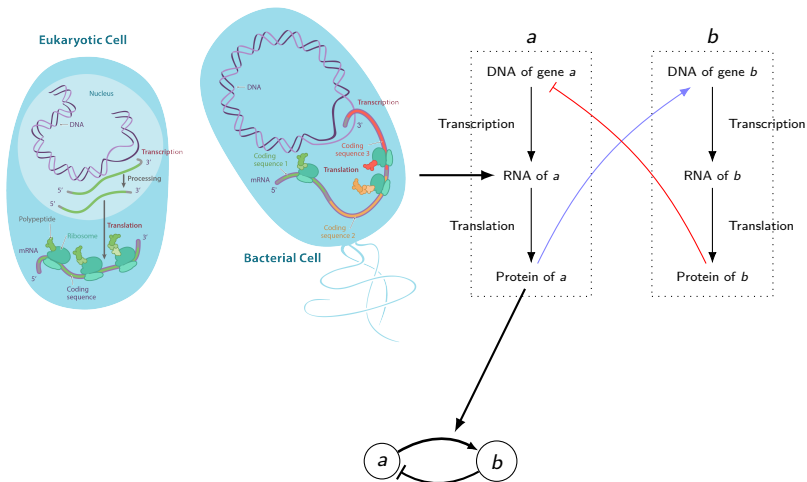
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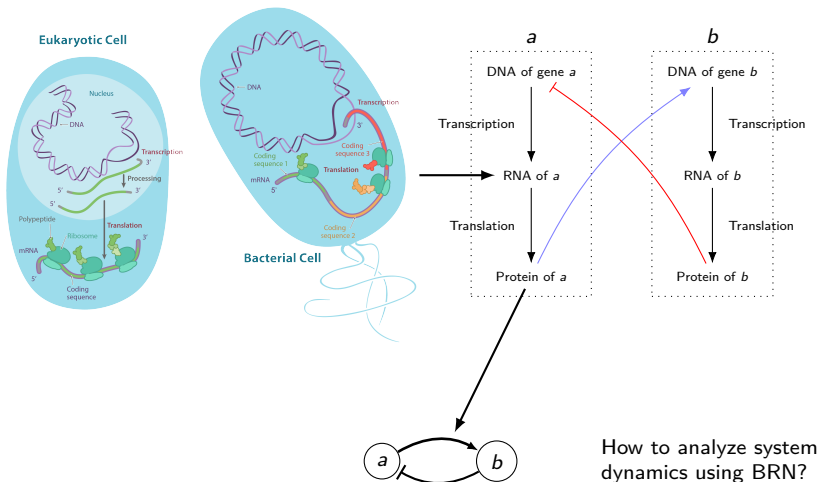
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Discrete Modeling [BT09]



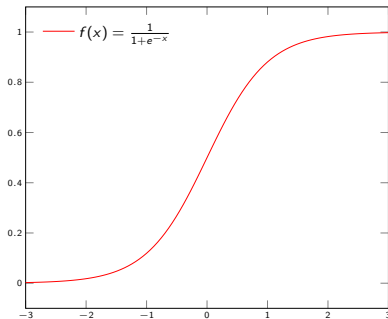
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Discrete Modeling [BT09]



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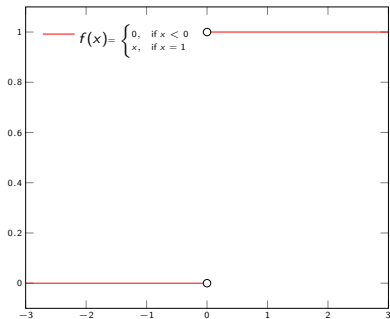


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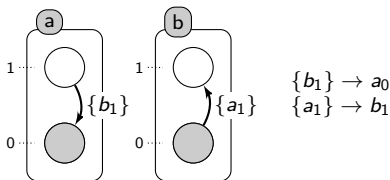


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Automata Network (AN)

A modeling framework representing state transitions and using $O(\#nodes)$ memory



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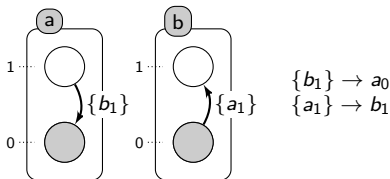


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Automata Network (AN)

A modeling framework representing state transitions and using $O(\#nodes)$ memory



The global reachability of a system state $(a, b) = (1, 1)$ is hard to compute

Discrete Modeling [BT09]

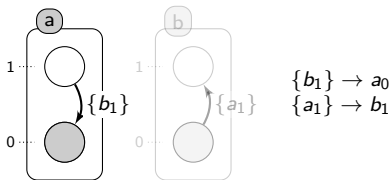


What are the possible values for a and b ?

- Behaviors of elements in BRN can be approximated to sigmoid functions [BT09]
- Approximated to piecewise functions. When the concentration of an entity is below the regulation threshold, the regulation is activated, otherwise inactivated. The value of a and b is in $\{0, 1, 2, \dots\}$
- A model containing dynamic information

Automata Network (AN)

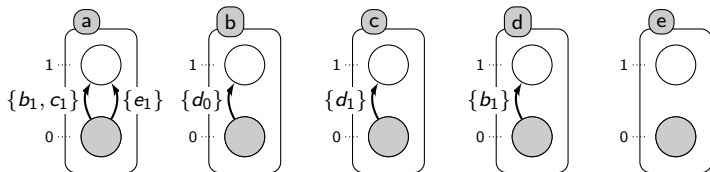
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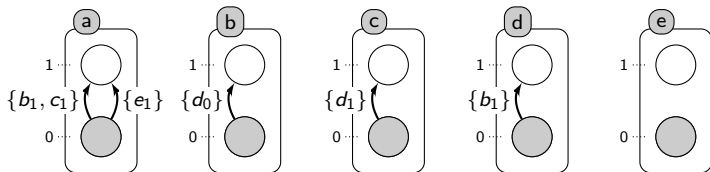
\Rightarrow The complexity of the reachability of the value of one variable $a = 1$ can be smaller?

Automata Network (AN)



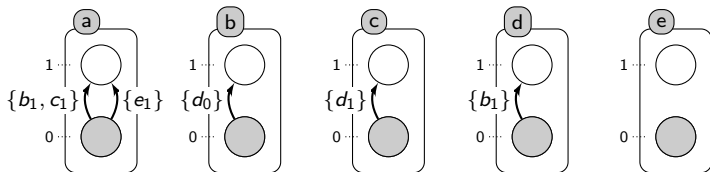
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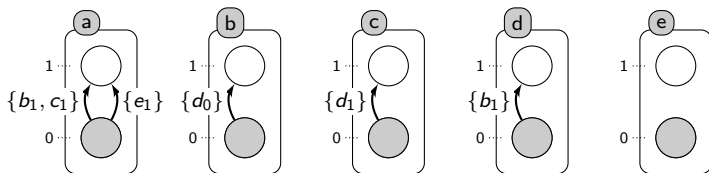
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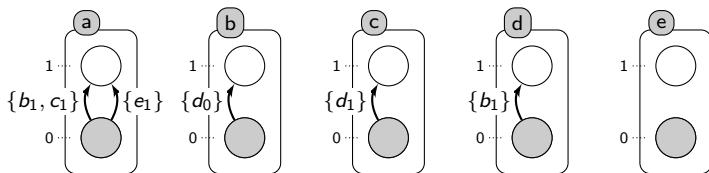
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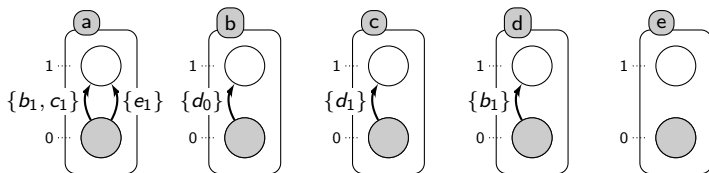
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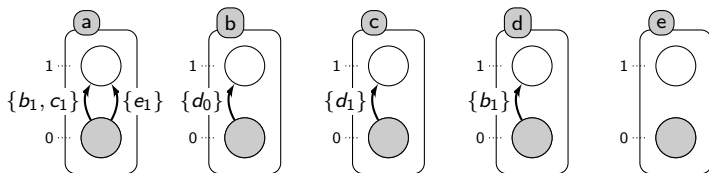
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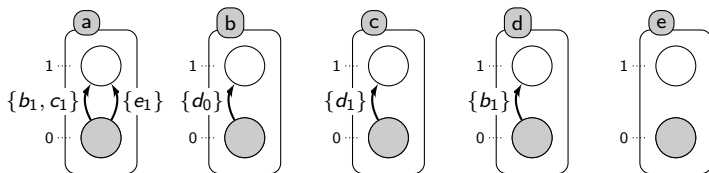
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Automata Network (AN)



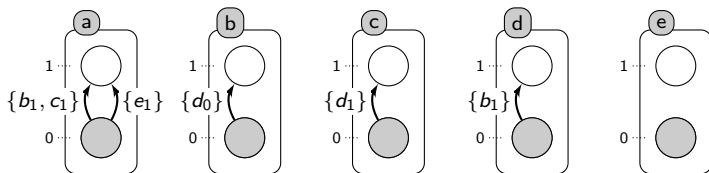
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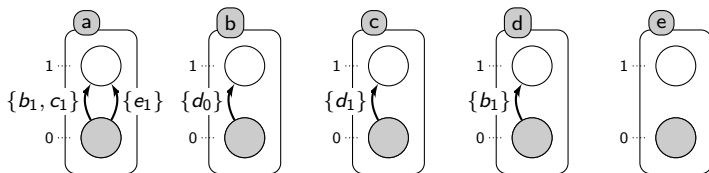
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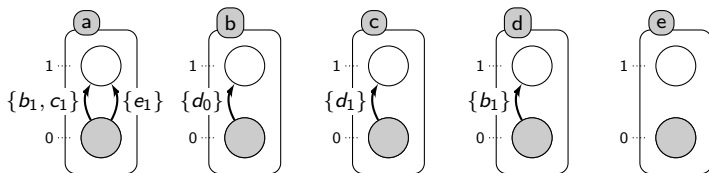
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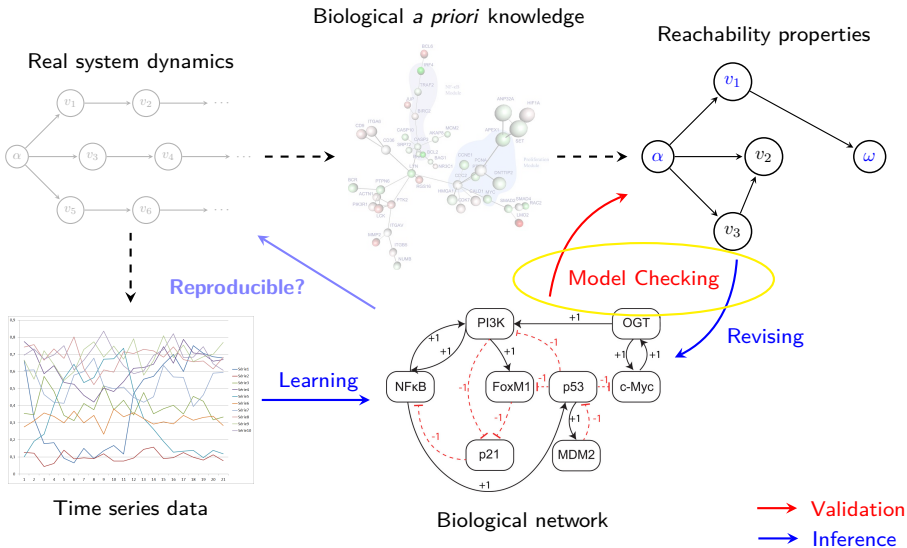
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Automata Network (AN)

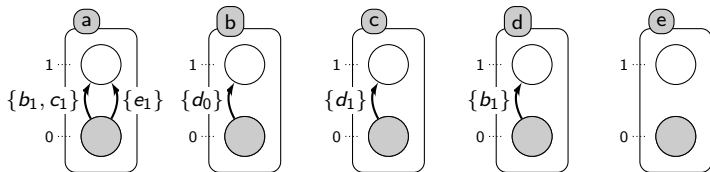


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- To study reachability problem → Simplified Local Causality Graph (SLCG) based on Local Causality Graph (LCG) [PMR12]

Problematics



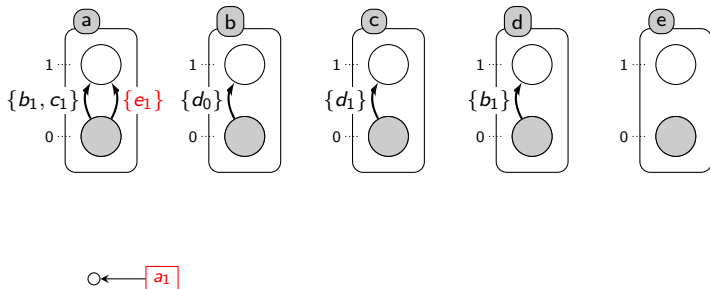
Simplified Local Causality Graph (SLCG)



a_1

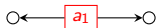
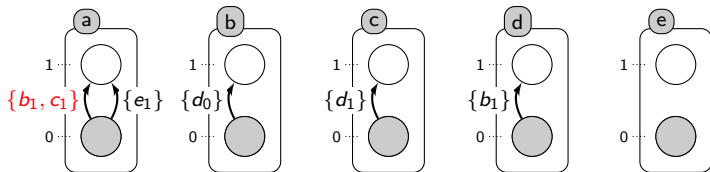
Small circles stand for transition nodes, squares for state nodes

Simplified Local Causality Graph (SLCG)



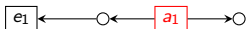
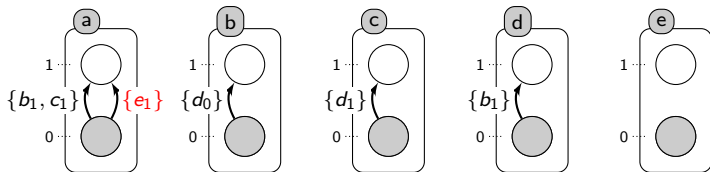
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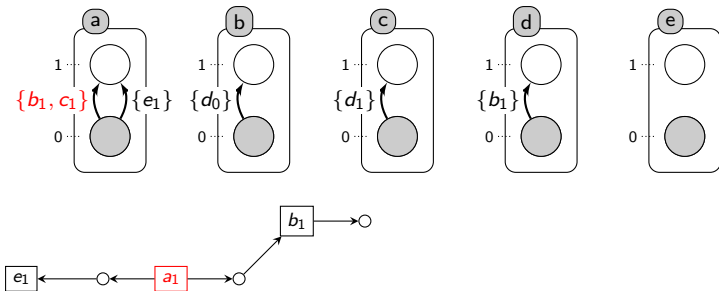
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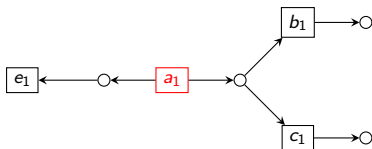
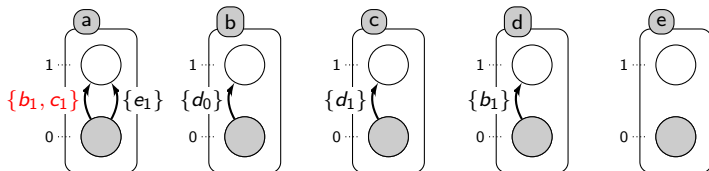
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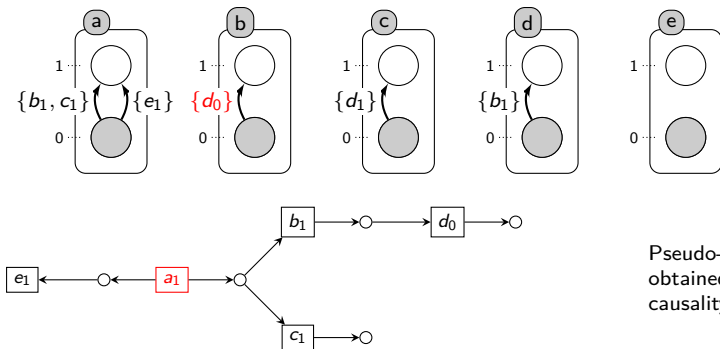
Simplified Local Causality Graph (SLCG)



Pseudo-reachability r' is obtained *via* pure recursive causality reasoning

$$r'(a_1) = r'(e_1) \vee (r'(b_1) \wedge r'(c_1))$$

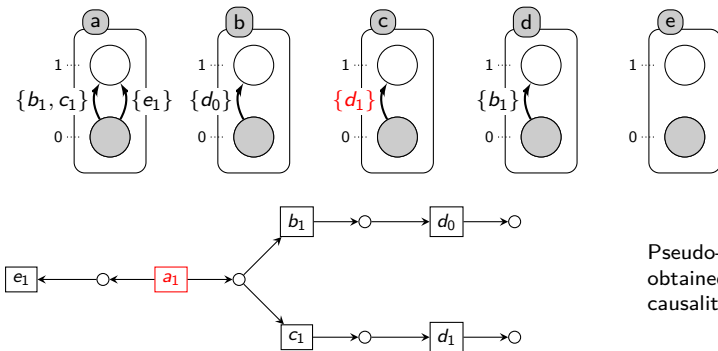
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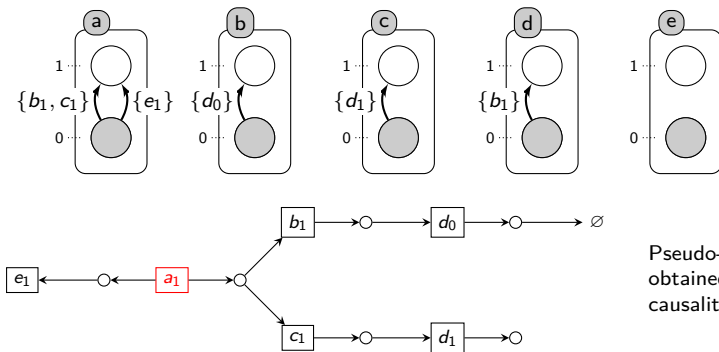
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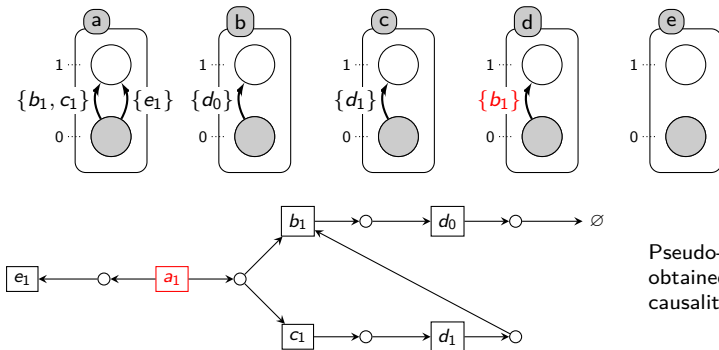
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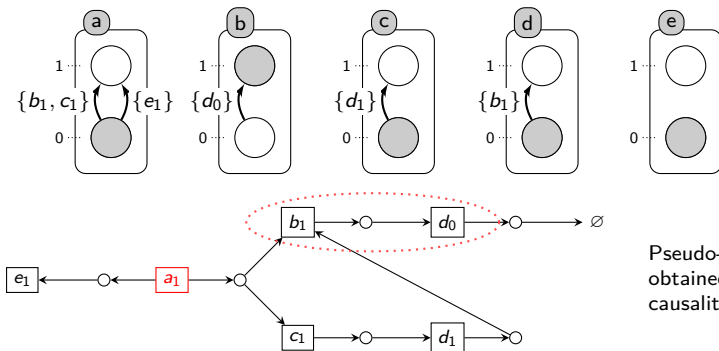


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$r'(d_0) = \text{True}$ because d_0 is at initial state

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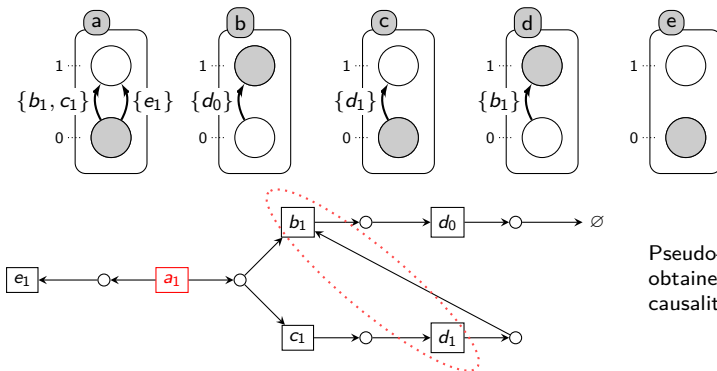


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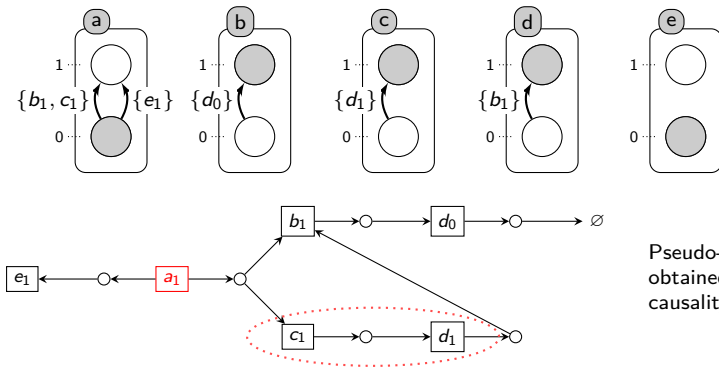
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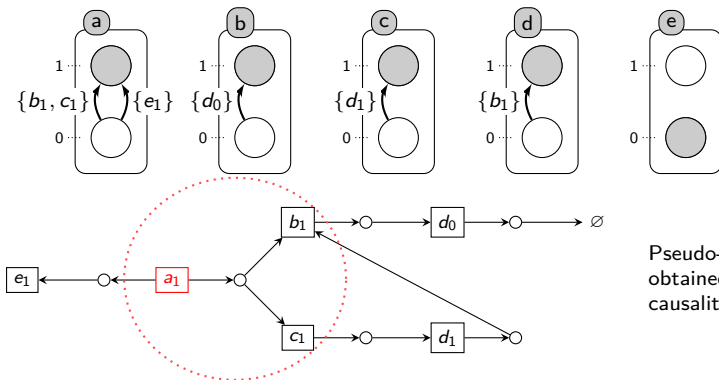


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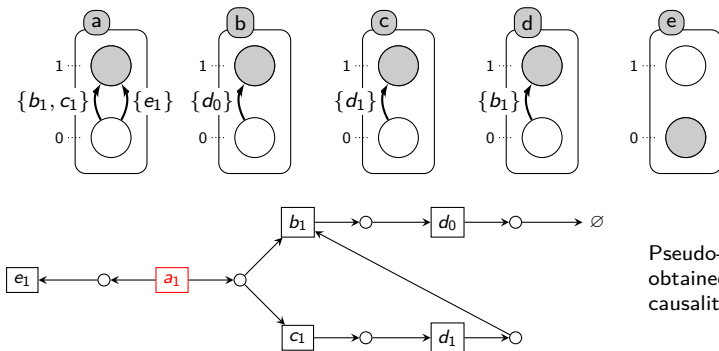
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$r'(a_i)$ is **not equivalent** to the reachability of a_1 because the state space is not fully explored

Counterexample of SLCG

- SLCG is **exact** for unreachability

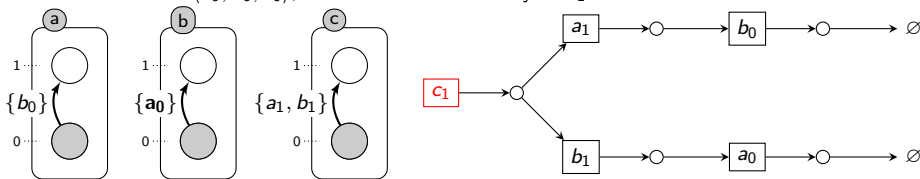
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Given initial state: $\langle a_0, b_0, c_0 \rangle$, consider the reachability of c_1 :

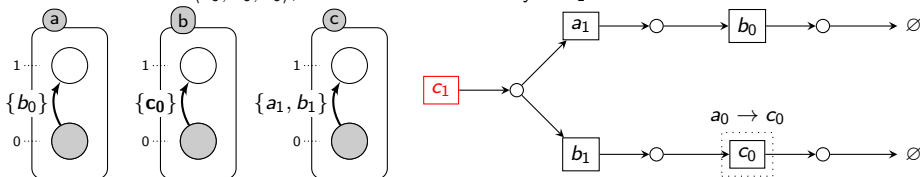


When a reaches level 1, b_1 is no longer reachable as the condition of $\{a_0\} \rightarrow b_1$ is not satisfied.
 \rightarrow Reaching a_1 disables the reachability of b_1 , *vice versa*, $\{a_1, b_1\}$ is not reachable, *i.e.* c_1 is unreachable

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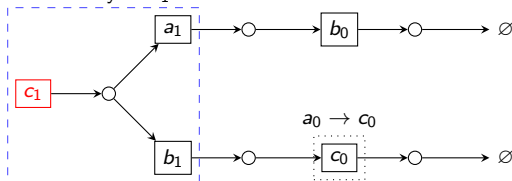
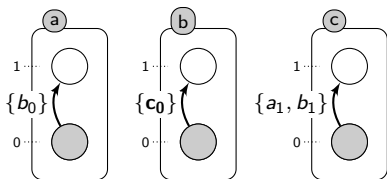


b_1 is not blocked by a , c_1 is reachable via $a_1 :: b_1 :: c_1$

Counterexample of SLCG

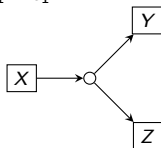
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SLCG does not show **orders** of local states
 \Rightarrow How to deal with this structure?



Solution of Inconclusive Structure

Complete search on global states

Solution of Inconclusive Structure

Complete search on global states \rightarrow State space explosion problem 😞

Solution of Inconclusive Structure

Complete search on global states \rightarrow State space explosion problem ☹

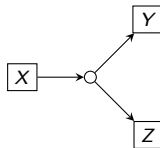
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Solution of Inconclusive Structure

Complete search on global states \rightarrow State space explosion problem ☹

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\Rightarrow Apply heuristic search on branches to find a **sequence** of local states in the form $Z :: Y :: X$ in order to avoid the inconclusiveness of the reachability of X

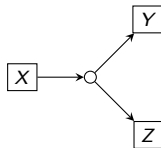


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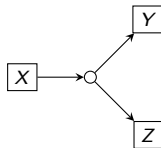
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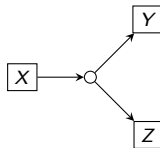
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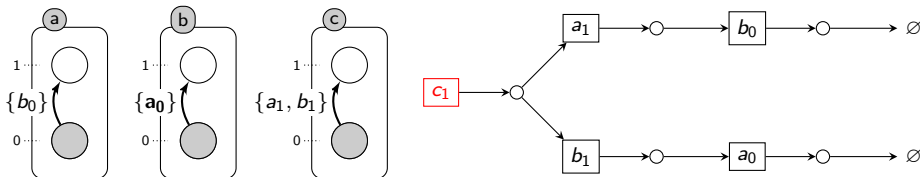
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- PermReach: searching all permutations in branches
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	PermReach	ASPReach [CRM ⁺ 18]
Method	Search all the permutations of branches	Search all the possible order of branches
Runtime	+	—
Conclusiveness	—	+

Solve the Counterexample Using ASPReach



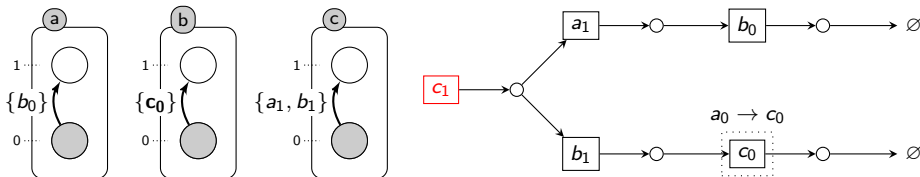
Notation: $a \triangleright b$ means a appears in the sequence before b

Order constraints in SLCG $\Rightarrow b_0 \triangleright a_1 \triangleright c_1$ and $a_0 \triangleright b_1 \triangleright c_1$

Additional constraint $\Rightarrow a_1 \triangleright b_1$ and $b_1 \triangleright a_1$

Contradiction in order, c_1 unreachable

Solve the Counterexample Using ASPReach



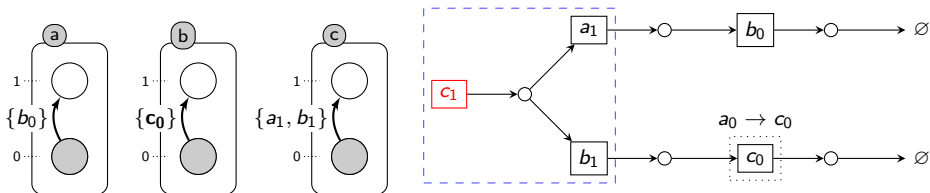
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Order constraints in SLCG $\Rightarrow b_0 \triangleright a_1 \triangleright c_1$ and $c_0 \triangleright b_1 \triangleright c_1$

Additional constraint $\Rightarrow a_1 \triangleright b_1$

The only admissible order is $a_1 :: b_1 :: c_1$

Solve the Counterexample Using ASPReach



Notation: $a \triangleright b$ means a appears in the sequence before b

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Additional constraint $\Rightarrow a_1 \triangleright b_1$

The only admissible order is $a_1 :: b_1 :: c_1$

\Rightarrow Problematic structure solved, reachable if sequences found, unreachable if not found

Benchmarks: on Biological Examples

- Traditional model checkers: Mole, NuSMV → **memory-out**, not listed in the benchmarks

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- Small example: λ -phage, 4 components

Model	λ -phage		
Inputs	4	Outputs	4
Total tests	$2^4 \times 4 = 64$		
Analyzer	Pint	PermReach	ASPREach
Reachable	36(56%)	38(59%)	
Unreachable	26(41%)		
Inconclusive	2(3%)	0(0%)	
Total time	< 1s		

Benchmarks: on Biological Examples

- Traditional model checkers: Mole, NuSMV → **memory-out**, not listed in the benchmarks
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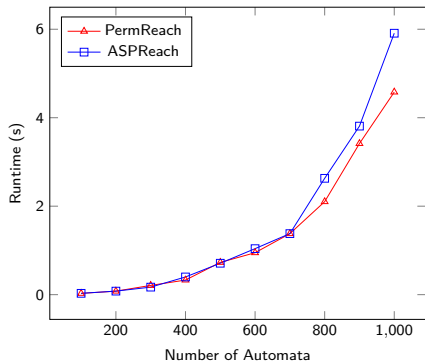
Model	λ-phage		
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Total tests	$2^4 \times 4 = 64$		
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Inconclusive	2(3%)	0(0%)	
Total time	< 1s		
Model	TCR		
Inputs	3	Outputs	5
Total tests	$2^3 \times 5 = 40$		
Analyzer	Pint	PermReach	ASPREach
Reachable	16(40%)		
Unreachable	24(60%)		
Inconclusive	0(0%)		
Total time	7s	0.85s	40s

Benchmarks: on Biological Examples

- Traditional model checkers: Mole, NuSMV → **memory-out**, not listed in the benchmarks
- Pure static analyzer: Pint [FPMR15]
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- Big examples: TCR (T-Cell Receptor, 95 components)
- EGFR (Epidermal Growth Factor Receptor, 104 components)

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Total tests	$2^4 \times 4 = 64$		
Analyzer	Pint	PermReach	ASPReach
Reachable	36(56%)	38(59%)	
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Model	TCR		
Inputs	3	Outputs	5
Total tests	$2^3 \times 5 = 40$		
Analyzer	Pint	PermReach	ASPReach
Reachable	16(40%)		
Unreachable	24(60%)		
Inconclusive	0(0%)		
Total time	7s	0.85s	40s
Model	EGFR		
Inputs	13	Outputs	12
Total tests	$2^{13} \times 12 = 98,304$		
Analyzer	Pint	PermReach	ASPReach
Reachable	64,282(65.4%)	74,268(75.5%)	
Unreachable	24,036(24.5%)		
Inconclusive	9,986(10.1%)	0(0%)	
Total time	9h50min	15min31s	3h46min

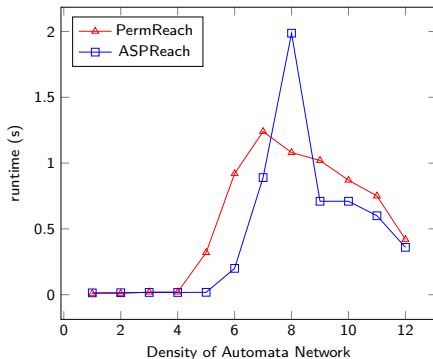
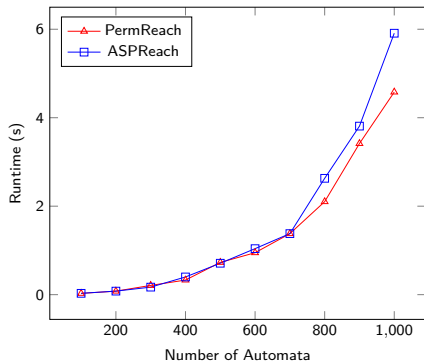
Benchmarks: on Random Examples



$$\text{density} = \frac{\# \text{transitions}}{\# \text{automata}}$$

Fixing density = 3

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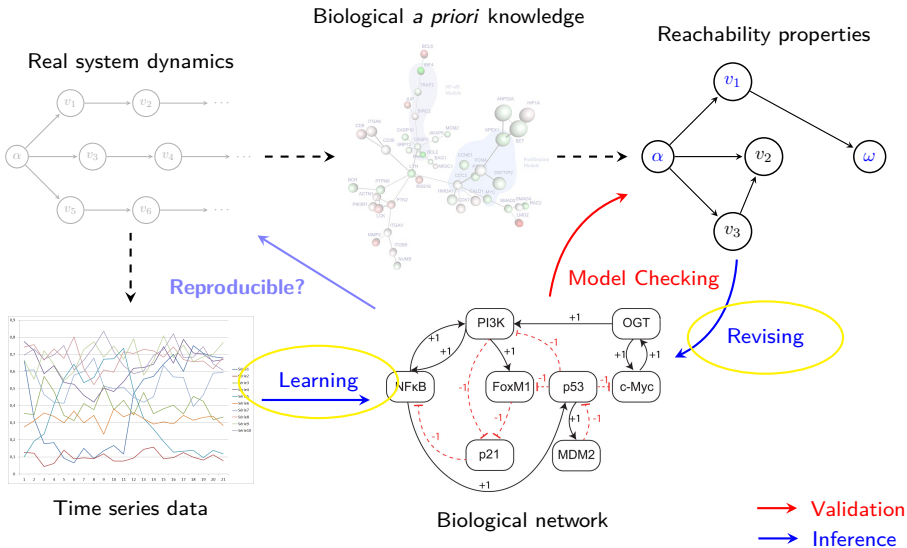


$$\text{density} = \frac{\# \text{transitions}}{\# \text{automata}}$$

Fixing density = 3

Fixing #automata = 20

Problematics



Learning and Revising Models

Time-series data

→ Validation
→ Reproduction

Learning and Revising Models



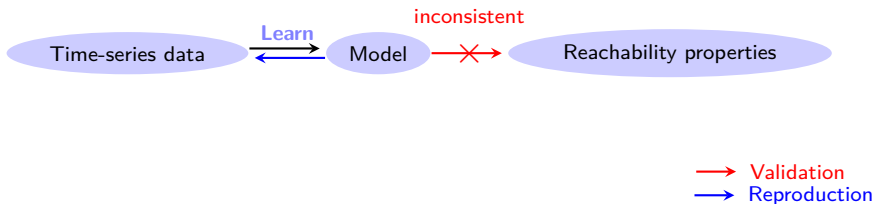
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Learning and Revising Models

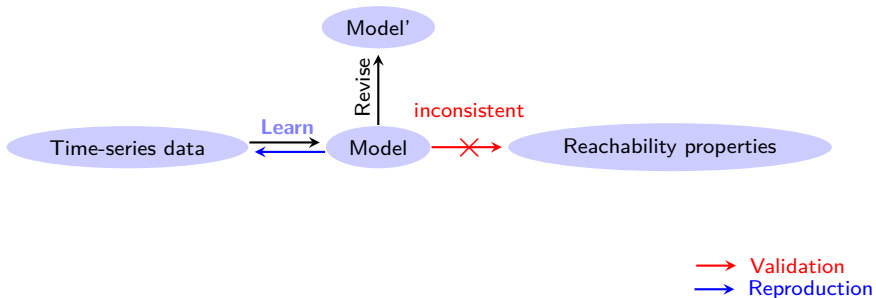


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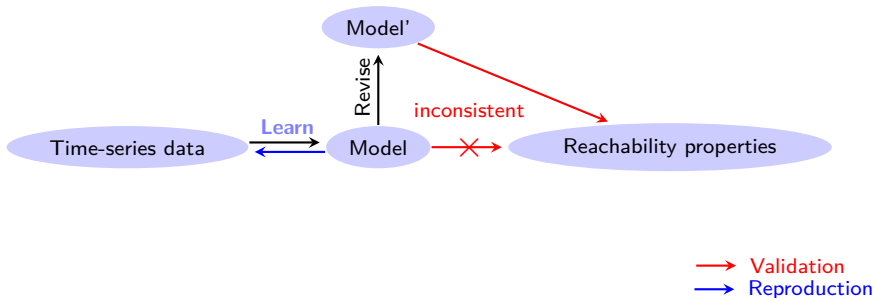
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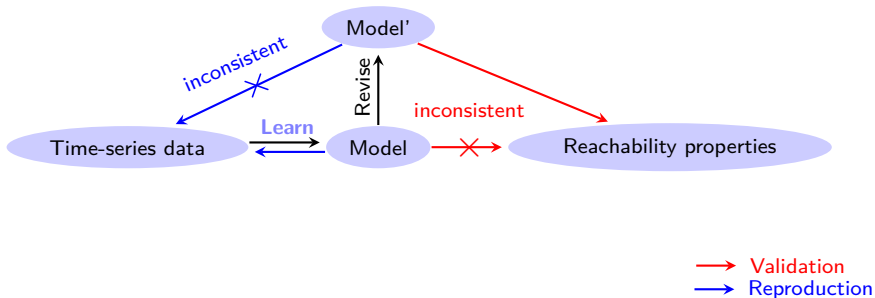
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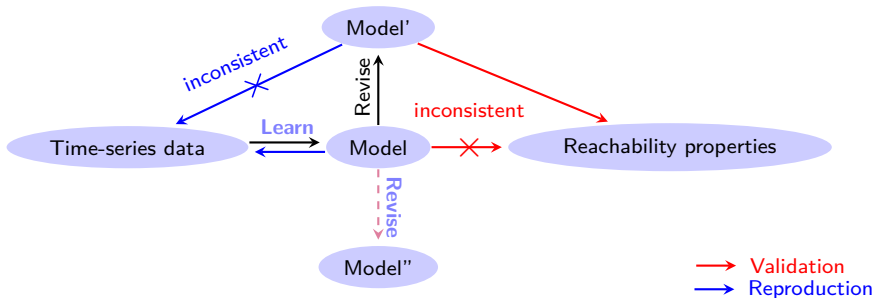
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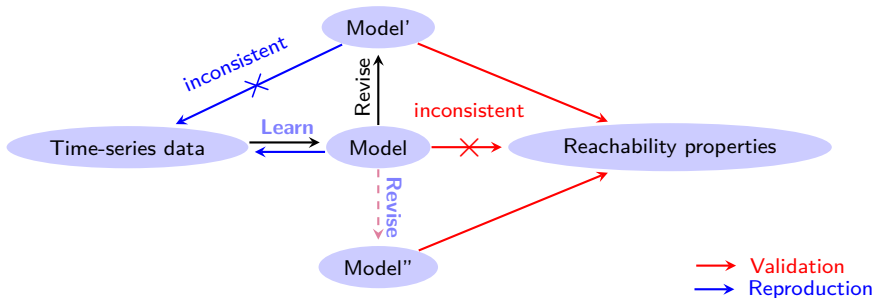
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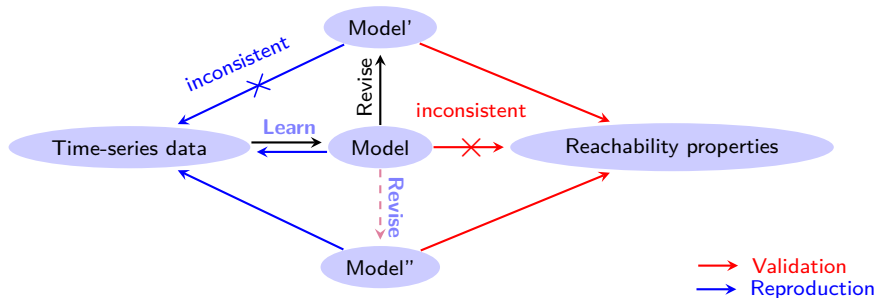
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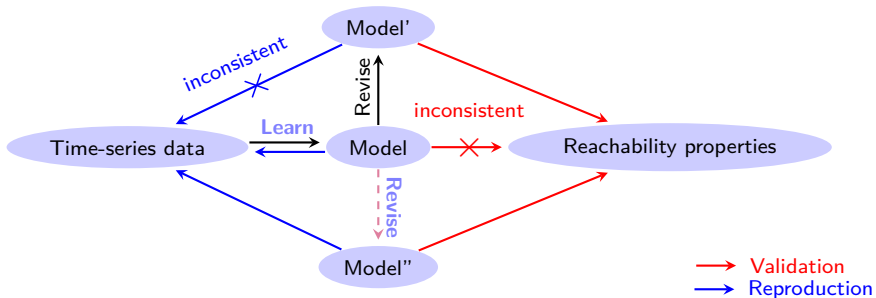
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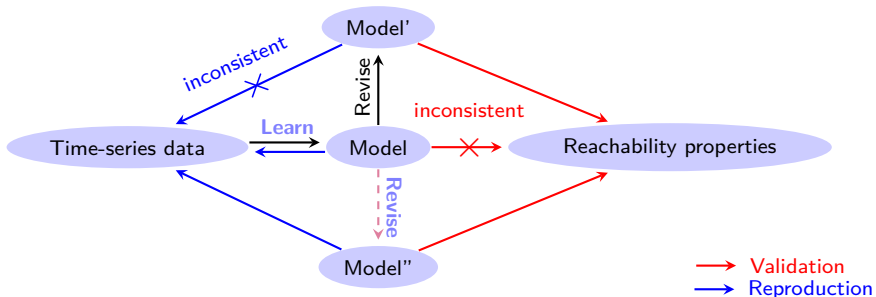


Learning and Revising Models



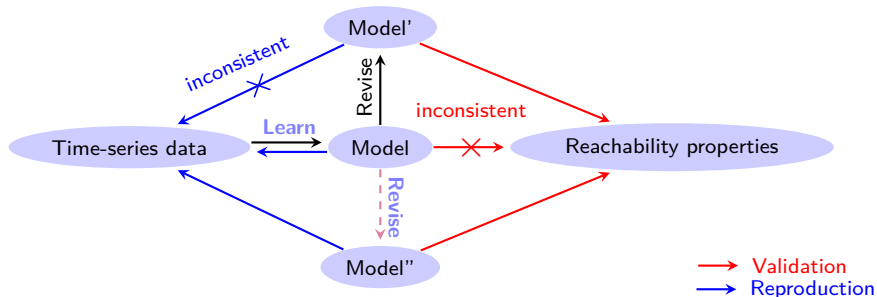
- CRAC: Completion *via* Reachability And Correlations

Learning and Revising Models



- CRAC: Completion *via* Reachability And Correlations
- M2RIT: Model Revision *via* Reachability and Interpretation Transitions

Learning and Revising Models



- CRAC: Completion *via* Reachability And Correlations
- M2RIT: Model Revision *via* Reachability and Interpretation Transitions

	CRAC	M2RIT
Learning phase	Correlation Coefficients	Learning from Interpretation Transitions (LFIT) [RFM ⁺ 18]
Revising phase	Reachability+ candidate transitions	Reachability+ time-series data

Learning Phase

CRAC: using correlation coefficient $r_{x,y} = \frac{\text{cov}(x,y)}{\sigma_x \sigma_y}$

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1.

t	0	1	2	3
a	2.01	2.51	1.97	1.17
b	0.74	0.87	0.78	0.33

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Δa	0.5	-0.54	-0.8
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M2RIT: using LFIT algorithm on discretized time-series data

Global transitions obtained from time-series data

$$\langle a_1, b_0, c_1 \rangle(t = T_1) \rightarrow \langle a_1, b_1, c_1 \rangle(t = T_1 + 1)$$

$$\langle a_1, b_0, c_0 \rangle(t = T_2) \rightarrow \langle a_1, b_1, c_0 \rangle(t = T_2 + 1)$$

$$\langle a_0, b_0, c_0 \rangle(t = T_3) \rightarrow \langle a_0, b_0, c_1 \rangle(t = T_3 + 1)$$

classified into

partial transitions

\Rightarrow

$$\{a_1\} \rightarrow b_1$$

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classified into

partial transitions

\Rightarrow

$$\{a_1\} \rightarrow b_1$$

The classification is sensitive to input \rightarrow revise the result

Revising Phase

Methodology:

	Reachable	Unreachable
Knowledge	R_K	U_K

Revising Phase

Methodology:

	Reachable	Unreachable
Knowledge	R_K	U_K
Inferred model	R_I	U_I

Revising Phase

Methodology:

	Reachable	Unreachable
Knowledge	R_K	U_K
Inferred model	R_I	U_I
Inconsistency	$R'_K = R_K \cap U_I$	$U'_K = R_I \cap U_K$

Revising Phase

Methodology:

	Reachable	Unreachable
Knowledge	R_K	U_K
Inferred model	R_I	U_I
Inconsistency	$R'_K = R_K \cap U_I$	$U'_K = R_I \cap U_K$
Keep consistent with	U_K	R_K

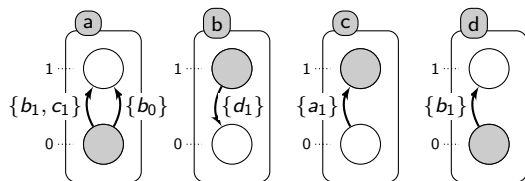
Revising Phase

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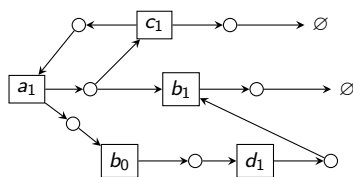
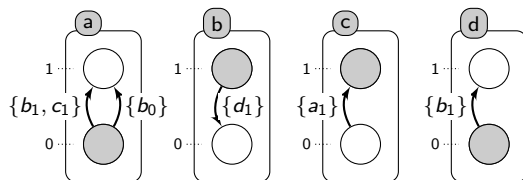
	Reachable	Unreachable
Knowledge	R_K	U_K
Inferred model	R_I	U_I
Inconsistency	$R'_K = R_K \cap U_I$	$U'_K = R_I \cap U_K$
Keep consistent with	U_K	R_K

- Generic solution: computes all the revisions \rightarrow complete but costly
 - Ensure unreachability by inhibiting elements in corresponding **cut set**
 - Ensure reachability by guaranteeing elements in corresponding **completion set**
- Heuristics: explores the saturated SLCG, aiming at revising transitions one by one \rightarrow memory-saving, fast but not complete

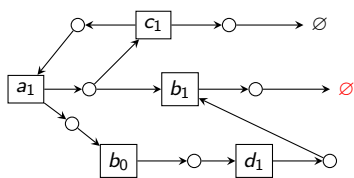
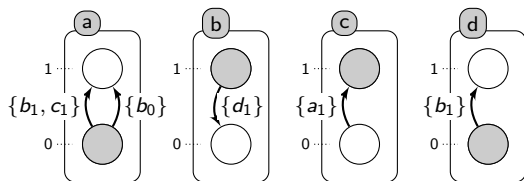
Cut Set



Cut Set



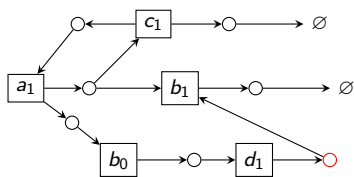
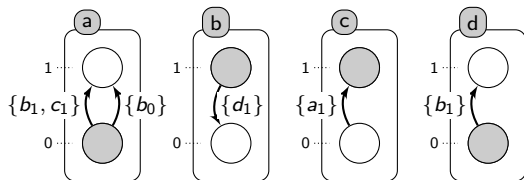
Cut Set



Node	Rank	\mathbb{V}
\emptyset (of b_1)	1	\emptyset
b_1	2	$\{\{b_1\}\}$
$\{b_1\} \rightarrow d_1$	3	$\{\{b_1\}\}$
d_1	4	$\{\{b_1\}, \{d_1\}\}$
$\{d_1\} \rightarrow b_0$	5	$\{\{b_1\}, \{d_1\}\}$
b_0	6	$\{\{b_0\}, \{b_1\}, \{d_1\}\}$
$\{b_0\} \rightarrow a_1$	7	$\{\{b_0\}, \{b_1\}, \{d_1\}\}$
\emptyset (of c_1)	8	\emptyset
c_1	9	$\{\{c_1\}\}$
$\{b_1, c_1\} \rightarrow a_1$	9	$\{\{b_1\}, \{c_1\}\}$
a_1	9	$\{\{a_1\}, \{b_1\}, \{b_0, c_1\}, \{c_1, d_1\}\}$
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Rank: topological numbering of the nodes in SLCG, nodes with higher rank cannot be successor of nodes with lower rank.

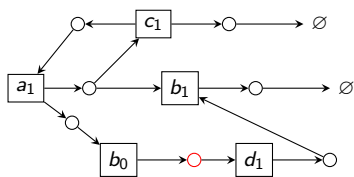
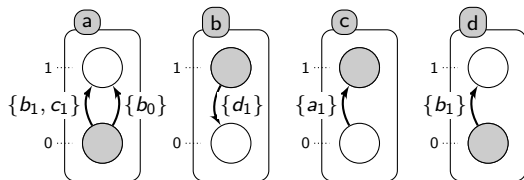
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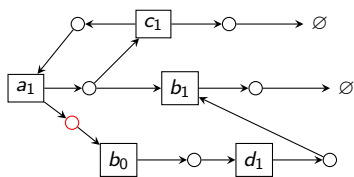
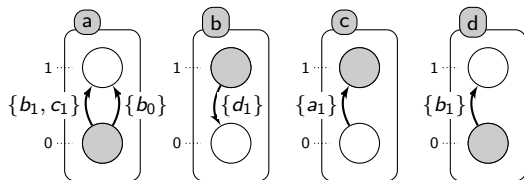
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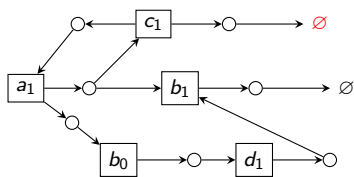
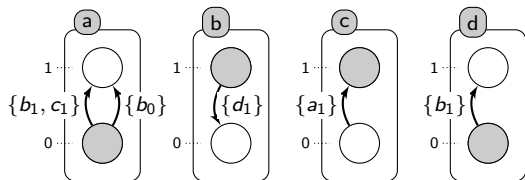
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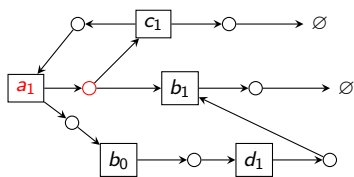
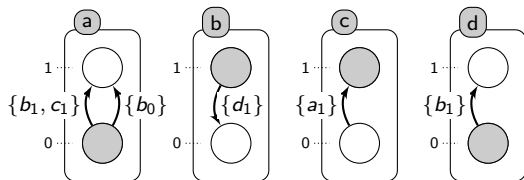
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$\{d_1\} \rightarrow b_0$	5	$\{\{b_1\}, \{d_1\}\}$
b_0	6	$\{\{b_0\}, \{b_1\}, \{d_1\}\}$
$\{b_0\} \rightarrow a_1$	7	$\{\{b_0\}, \{b_1\}, \{d_1\}\}$
\emptyset (of c_1)	8	\emptyset
c_1	9	$\{\{c_1\}\}$
$\{b_1, c_1\} \rightarrow a_1$	9	$\{\{b_1\}, \{c_1\}\}$
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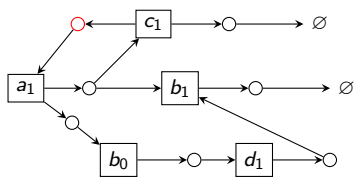
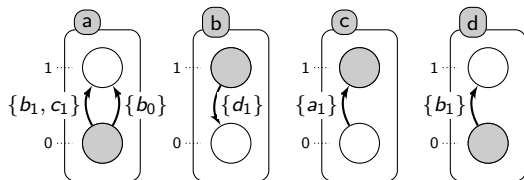
Cut Set



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$\{b_1\} \rightarrow d_1$	3	$\{\{b_1\}\}$
d_1	4	$\{\{b_1\}, \{d_1\}\}$
$\{d_1\} \rightarrow b_0$	5	$\{\{b_1\}, \{d_1\}\}$
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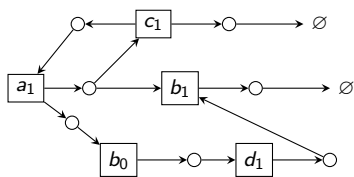
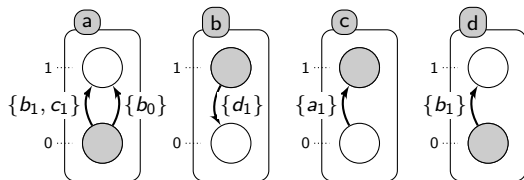
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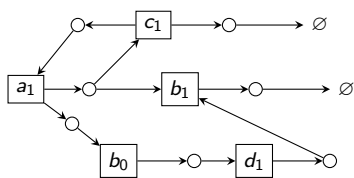
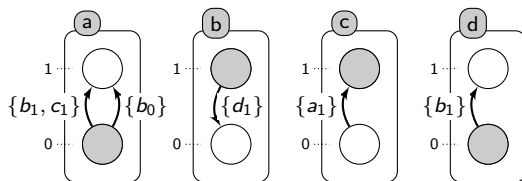
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By inhibiting the reachability of one of the sets in cut set

$\{\{a_1\}, \{b_1\}, \{b_0, c_1\}, \{c_1, d_1\}\}$, a_1 shall not be reachable.

Corresponding transition sets can be deduced.

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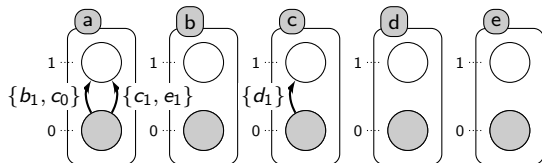
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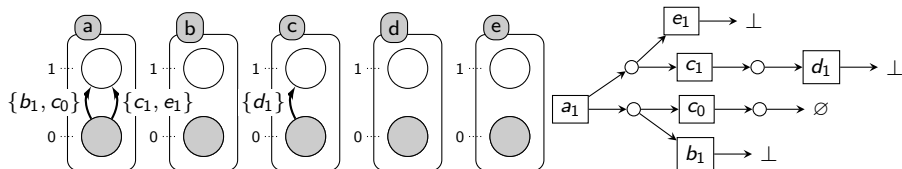
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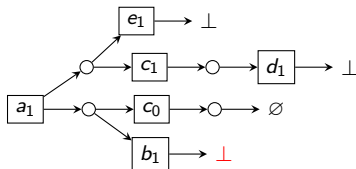
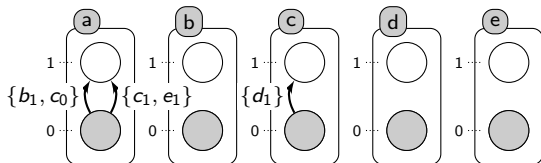
Completion Set



Completion Set



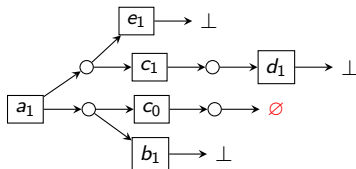
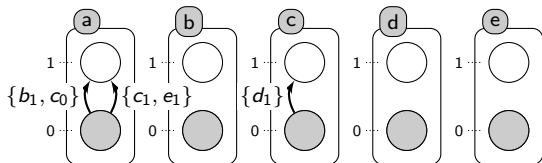
Completion Set



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$\{d_1\} \rightarrow c_1$	8	$\{\{d_1\}\}$
c_1	9	$\{\{c_1, d_1\}\}$
\perp (of e_1)	10	\emptyset
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$\{c_1, e_1\} \rightarrow a_1$	12	$\{\{c_1, e_1\}, \{d_1, e_1\}\}$
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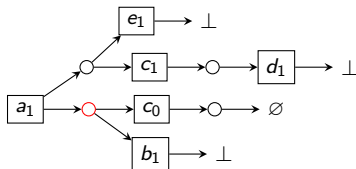
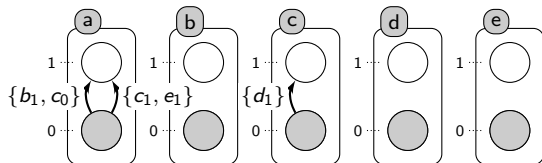
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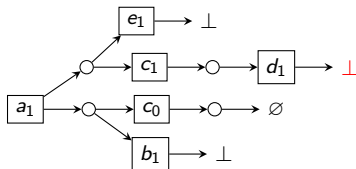
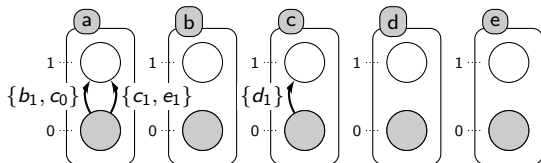
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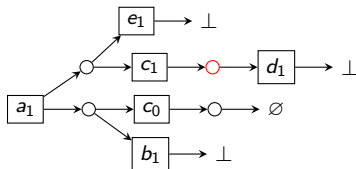
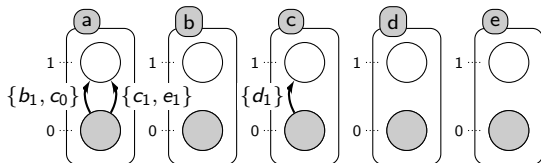
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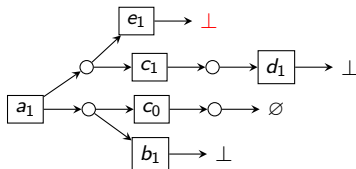
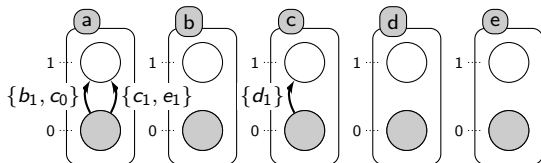
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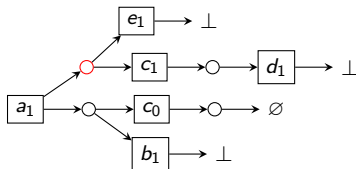
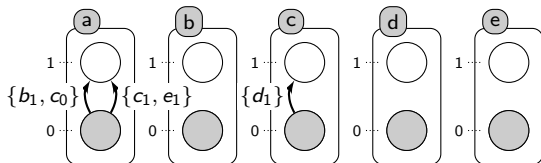
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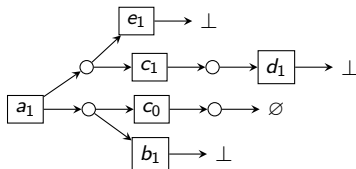
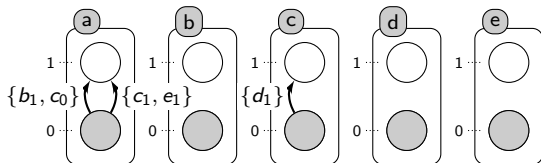
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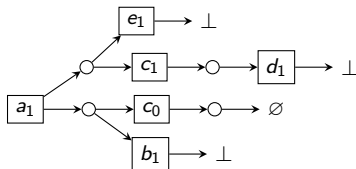
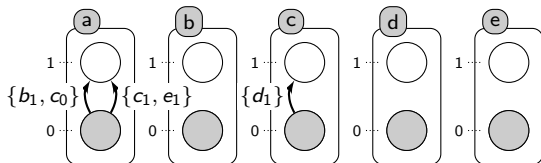


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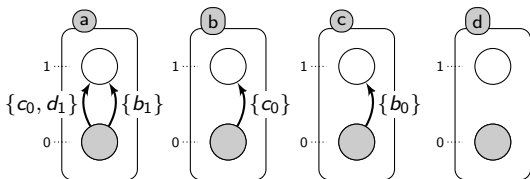
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 - Compute with the SLCG involving all the wanted states in *a priori* knowledge

Heuristic Revision

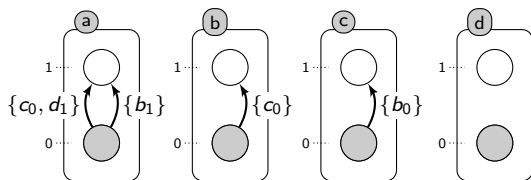


$$\alpha = \langle a_0, b_0, c_0, d_0 \rangle, \omega = a_1$$

$$U_K = \{(\alpha, b_1), (\alpha, d_1)\}$$

$$R_K = \{(\alpha, a_1)\}$$

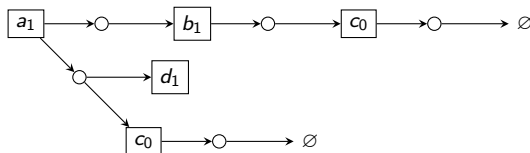
Heuristic Revision



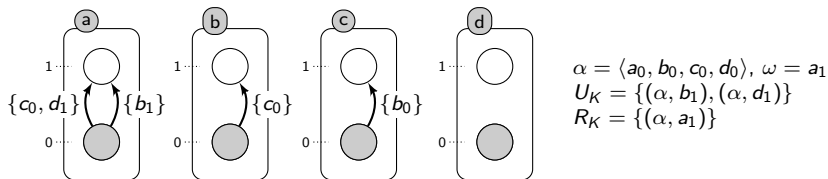
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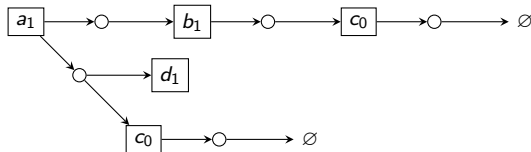
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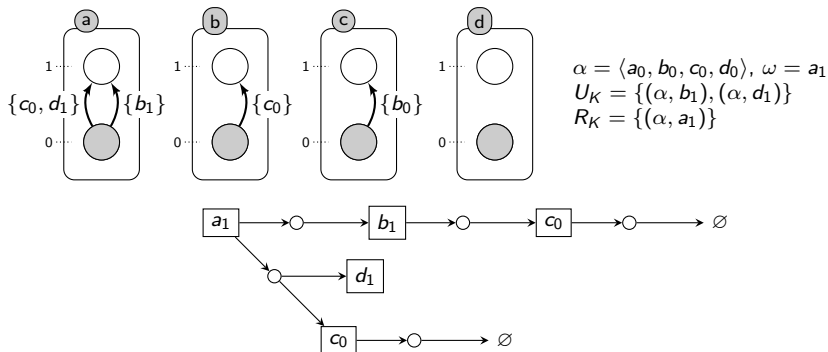
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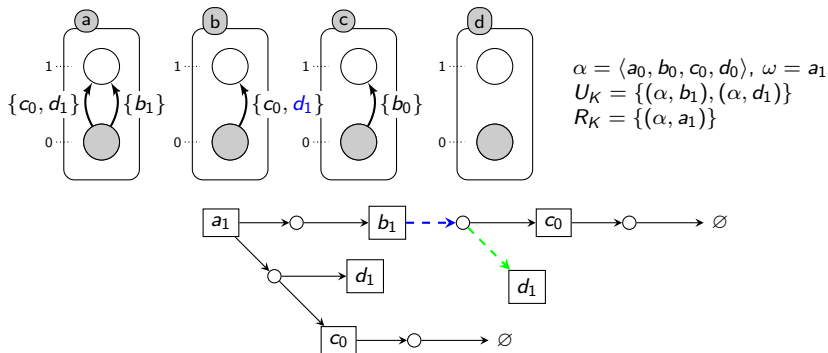
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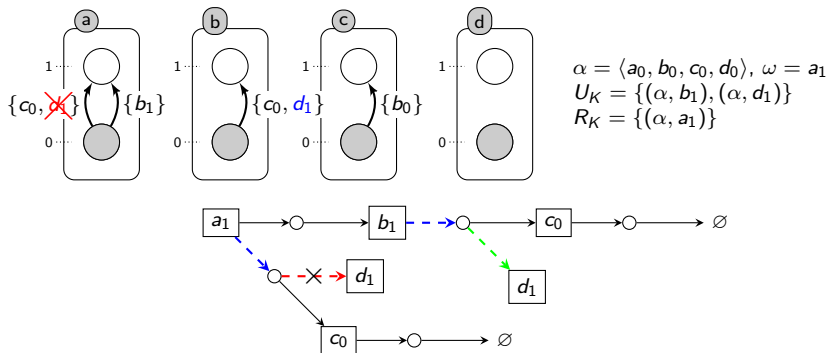
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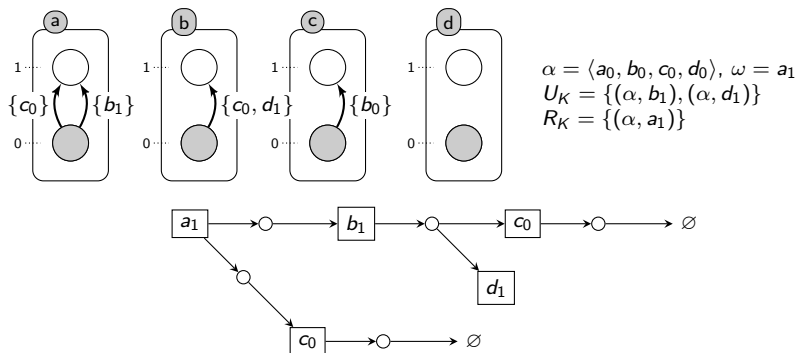
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- Check the reachability of (α, a_1) : reachable, finish

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 - By picking only reachability as revision criteria, inferred models are bisimilar to original systems in the sense of reachability

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Reachability Analysis and Revision of Dynamics of Biological Regulatory Networks

Xinwei Chai

École Centrale de Nantes
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Co-encadrant de thèse : Morgan MAGNIN, Professeur des universités, École Centrale de Nantes

May 24, 2019

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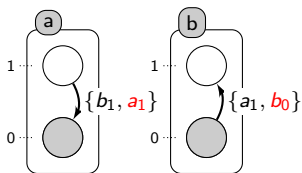
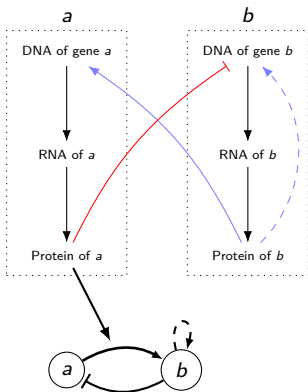


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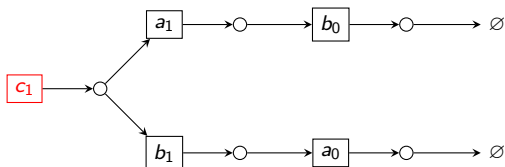
Self Regulation



- The self regulation of b can contain two elements: self activation/inhibition and intrinsic degradation.
- However they are not necessary in ABAN case. As the transitions are between Boolean states, the state changes necessarily from one to another \rightarrow there is no need to add self regulation related conditions in the transitions.

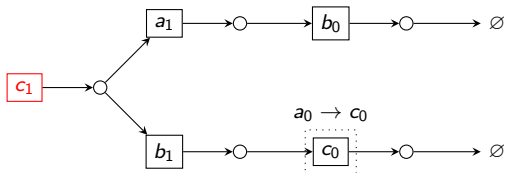
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- It returns inconclusive instances in the counterexample given in the presentation and also in λ -phage model



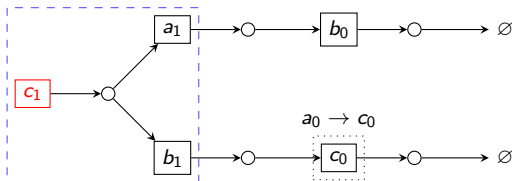
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