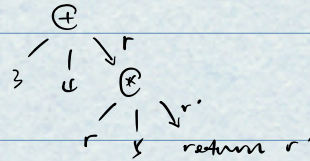
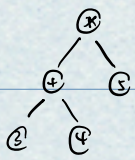


Continuation

$(3 + 4) * 5$

add 3 & 4 and pass the result r to the computation
mult r 5 and pass the result r' to the computation
return r'



① fun add $x\ y\ k = k(x + y)$
fun mult $x\ y\ k = k(x * y)$
↑
continuation

ex. add 3 4 (fun $r \Rightarrow$ mult $r\ 5$ (fun $r' \Rightarrow r'$))
 \Rightarrow (fun $r \Rightarrow \dots$) (3 + 4)
 \Rightarrow [7 / r] mult $r\ 5$ (fun $r' \Rightarrow r'$)
 \Rightarrow (fun $r' \Rightarrow r'$) (7 * 5)
 \Rightarrow [35 / r'] r' "tail recursive"
 \Rightarrow 35

add : $\text{int} \rightarrow \text{int} \rightarrow (\text{int} \rightarrow 'a) \rightarrow 'a$

② (* sum : $\text{int list} \rightarrow \text{int}$ *)
fun sum [] = 0
| sum ($x :: xs$) = $x + (\text{sum } xs)$

(* tsum : $\text{int list} \rightarrow \text{int}$ *)
ENS: $\text{tsum } (l, \text{acc}) = \text{sum } l + \text{acc}$
fun tsum ([], acc) = acc
| tsum ($x :: xs$, acc) = tsum (xs , $x + \text{acc}$)

no explicit accumulator

(* csum : $\text{int list} \rightarrow (\text{int} \rightarrow 'a) \rightarrow 'a$ *)
ENS: $\text{csum } l\ k \cong k(\text{sum } l)$

*)
fun csum [] $k = k(0)$
| csum ($x :: xs$) $k = \text{csum } xs$ (fun $s \Rightarrow k(x + s)$)

$S @ (x :: xs)$?

CPS means roughly call
functions tail-recursively with
functions as accumulator

initial continuation

csum [2, 3] Int.toString \rightarrow "5"

csum [2, 3] : ($\text{int} \rightarrow 'a$) $\rightarrow 'a$

csum [2, 3] Int.toString

\Rightarrow csum [3] (fun $s \Rightarrow \text{Int.toString } (2 + s)$) \cong "5"

\Rightarrow csum [1] (fun $s' \Rightarrow \text{Int.toString } (3 + s')$)

\Rightarrow (fun $s' \Rightarrow \text{Int.toString } (3 + s')$) 0

\Rightarrow

② datatype 'a tree = Empty | Node of 'a tree * 'a * 'a tree

(* inorder : 'a tree * 'a list → 'a list

ENS: $\text{inorder}(T, acc) \cong L @ acc$ where L consists of the elements of T as seen in an in order traversal

*)

fun inorder (Empty, acc) = acc

| inorder (Node(l, x, r), acc) = inorder(l, x :: inorder(r, acc))

(* same : int list * int list → bool

ENS: $\text{same}(X, Y) \Rightarrow \begin{cases} \text{true, if } X \cong Y \\ \text{false otherwise} \end{cases}$

*)

fun same ([], []) = true

| same (x :: xs, y :: ys) = $x = y$ and also same (xs, ys)

| same _ = false $\Rightarrow \text{len } X \neq \text{len } Y$

(* treematch : int tree * int list → bool

ENS: $\text{treematch}(T, L) \Rightarrow \begin{cases} \text{true, if the inorder traversal of } T \text{ is } L \\ \text{false, otherwise} \end{cases}$

fun treematch (T, L) = same (inorder(T, []), L)

! problem: if the first elements are different, still scan through the entire tree in inorder

(* prefix : int tree → int list → (int list → bool) → bool

REQ: k is total

ENS: $\text{prefix } T \ L \ k \Rightarrow \begin{cases} \text{true, if } \exists \text{ lists } L_1, L_2 \text{ s.t. } L \cong L_1 @ L_2, \text{ the inversed traversal of } T \text{ is } L_1, \\ \quad k(L_2) \cong \text{true}; \\ \text{false, otherwise} \end{cases}$

fun prefix Empty L k = $k(L) \cong L_1 = [] \Rightarrow L_2 = L$

| prefix (Node(left, x, right)) L k =

prefix left L (fun [] \Rightarrow false $\rightarrow \text{len}(T) > \text{len}(L)$

| (y :: ys) $\Rightarrow (x = y)$ and also (prefix right ys k))

fun treematch (T, L) = prefix T L List.null

'a list → bool (aka null) fun null [] = true

| null _ = false

@ success & failure continuation

(* search : ('a → bool) → 'a tree → ('a → 'b) → (unit → 'b) → 'b

success continuation failure continuation

REQ: p is total

ENS: $\text{search } p \ T \ sc \ fc \Rightarrow \begin{cases} sc(x) \text{ if } x \text{ is in the tree \& } p(x) \cong \text{true} \\ fc(), \text{ otherwise} \end{cases}$

*)

fun search p Empty sc fc = fc() cannot have x to call sc !

| search p (Node(l, x, r)) sc fc = if $p(x)$ then $sc(x)$

else search p l sc (fun () \Rightarrow search p r sc fc)

fun findEvens (T: int tree) : int option =

search (fun x => x mod 2 = 0) T SOME (fun () => NONE)
(fun x => SOME x)

Exception & handle

exception Silly

Silly : exn (exception or extensible)

raise Silly : 'a -> type? can accommodate! e.g. (if 3 = 4 then raise Silly else 0) : int
(if 3 = 4 then raise Silly else "0") : string

fun f(x) = f(x) f : 'a -> 'b

fun g(x) = raise Silly g : 'a -> 'b

(if 3 = 4 then g(0) else "0") : string g : int -> string

exception Rdiv of real constructors are also functions!

Rdiv : real -> exn Rdiv (3.14) : exn

raise Rdiv (3.14) : 'a
a keyword, not function

fun rdivide (r, r2) =

if Real.abs (r2) < 0.0001 then raise Rdiv (r) rdivide : real * real -> real
else r / r2

rdivide (3.14, 0.0) => --- raise Rdiv (3.14) not a value! error message at top level

(

rdivide (3.14, 0.0)
Block of code ---

) handle Rdiv (r) => "oops, you tried to divide" ^ (Real.toString (r)) ^ " by 0"
↑
Suppose this has type [3.14 / r] need return a string

(

Block of
code
[expression]

) handle Rdiv (r) => 9.99999 * r
| Silly => 1.0
| _ => 0.0

(do handle P_i => e_i

| P₂ => e₂

⋮

| P_n => e_n) : t

e₀, e₁, ..., e_n must all have the same type t,
which is then the type of the expression & the
patterns p₁, ..., p_n must match exceptions.

If e₀ evaluates to v then that is the value of the overall expression.

n - Queens

local decisions are sometimes hard => backtracking!

exception Conflict (i, j) represent positions on $n \times n$ board
 $\downarrow \downarrow$
 col row $1 \leq i, j \leq n$

$(\ast \text{Threat} : \text{int} \times \text{int} \rightarrow \text{int} \times \text{int} \rightarrow \text{bool} \ast)$

fun Threat $(x, y) (a, b) = x=a \text{ or else } y=b \text{ or else } x+y=a+b \text{ or else } x-y=a-b$

$(\ast \text{conflict} : \text{int} \times \text{int} \rightarrow (\text{int} \times \text{int}) \text{ list} \rightarrow \text{bool} \ast)$

fun conflict $p = \text{List.exists} (\text{threat } p)$ $\text{threat } p : \text{int} \times \text{int} \rightarrow \text{bool}$

$(\ast a \rightarrow \text{bool}) \rightarrow \ast a \text{ list} \rightarrow \text{bool}$

\downarrow
 $\ast a : \text{int} \times \text{int}$

returns $\text{int} \times \text{int} \text{ list} \rightarrow \text{bool}$

1st attempt raise exception

$(\ast \text{addqueen} : \text{int} \times \text{int} \times (\text{int} \times \text{int}) \text{ list} \rightarrow (\text{int} \times \text{int}) \text{ list} \ast)$

col n existing queen all queen
 placements placements

& raise Conflict if it is not possible to place the remaining queens

queens : $\text{int} \rightarrow (\text{int} \times \text{int}) \text{ list}$ or raise Conflict

\ast

fun addqueen $(i, n, Q) =$

let
 fun try $j =$

$(\text{if conflict } (i, j) Q \text{ then raise Conflict}$

check if can move queen up

else if $i=n$ then $(i, j) :: Q$

else addqueen $(i+1, n, (i, j) :: Q)$ handle Conflict \Rightarrow if $j=n$ then raise Conflict else try $(j+1)$

in
 try 1

end

fun queens $n = \text{addqueen} (1, n, [])$ alternative : fun queens $n = \text{SOME} (\text{addqueen} (1, n, []))$

eg. queens 4 $\Rightarrow [(4, 3), (3, 1), (2, 4), (1, 2)]$ in reverse order!

handle Conflict $\Rightarrow \text{NONE}$

queens 1 $\Rightarrow [(1, 1)]$

queens 2 \Rightarrow raise an unhandled conflict

2nd attempt: success / failure confirmation

$(\ast \text{addqueen} : \text{int} \times \text{int} \times (\text{int} \times \text{int}) \text{ list} \rightarrow ((\text{int} \times \text{int}) \text{ list} \rightarrow \ast)$

fun addqueen (i, n, Q) so $fc =$

let

fun try $j =$ let
 fun fnew $() = \text{if } j=n \text{ then } fc() \text{ else try } (j+1)$

in

if conflict $(i, j) Q$ then fnew $()$

else if $i=n$ then $(i, j) :: Q$

else addqueen $(i+1, n, (i, j) :: Q)$ so fnew

end

in

try 1

end

fun queens $n = \text{addqueen} (1, n, []) \text{ SOME } (fn () \Rightarrow \text{NONE})$

3rd attempt: use options

fun try =

case (if conflict $(i, j) Q$ then NONE

else if $i=n$ then $(i, j) :: Q$

else addqueen $(i+1, n, (i, j) :: Q)$ of


```
if result = NONE then try(j+1, a) else  
NONE => if j = n then NONE else try(j+1)  
| result => result
```

Multi-purpose Continuation

```
(* find : int tree -> (int -> (unit -> 'a) -> 'a) -> (unit -> 'a) -> 'a *)
```

```
fun find Empty se fc = fc()
```

```
| find (Node(l, x, r)) se fc =
```

```
let
```

```
fun fnew() = find l se (fun () => find r se fc)
```

```
in
```

```
if x mod 2 = 0 then se x fnew else fnew()
```

```
end
```

```
fun findfirst T = find T (fun x => fun fc => SOME(x)) (fun () => NONE)
```

```
fun findall T = find T (fun x => fun fc => x::(fc())) (fun () => [])
```

```
fun count T = find T (fun x => fun fc => 1+fc()) (fun () => 0)
```