

Ref

t ref $(t: \text{any ML type})$

$\boxed{v} \rightarrow \text{a cell}$ e.g. $\boxed{7}$ a value of type int ref containing the value 7 of type int .
 \downarrow
 ref 7

ref e If $e \rightarrow v$, then $\text{ref } e \rightarrow \boxed{v}$

$!e$ If $e \rightarrow \boxed{v}$, then $!e \rightarrow v$

$e_1 := e_2$ If $e_1 \rightarrow \boxed{w}$ and if $e_2 \rightarrow v$, then replace w with v in the cell and return $()$

e.g. $\text{val } c = \text{ref } 7$ $\left\{ \begin{array}{l} \boxed{7} / c \\ \text{val } () = c := 4 \\ \text{val } v = !c \end{array} \right\}$
 $\left\{ \begin{array}{l} 4 / c \\ 4 / v \end{array} \right\}$

type rules $\left\{ \begin{array}{l} \text{ref } e: t \text{ ref if } e: t \\ !e: t \text{ if } e: t \text{ ref} \\ e_1 := e_2: \text{unit if } e_1: t \text{ ref and } e_2: t \end{array} \right.$

ref is a constructor, $'a \rightarrow 'a \text{ ref}$
 can pattern matching

aliasing

$\text{val } c = \text{ref } 10$
 $\text{val } w = !c$
 $\text{val } d = c$
 $\text{val } () = d := 42$
 $\text{val } v = !c \rightarrow 10 / w \ 42 / v$



c & d are aliases

equality

$\text{val } c = \text{ref } 10$
 $\text{val } e = \text{ref } 10$
 $\text{val } d = c$
 $\Rightarrow !c = !e$
 $c = d \rightarrow \text{true}$
 $c = e \rightarrow \text{false}$

equal iff bound to the same cell!

sequential expression

$(e_1; e_2; \dots; e_n) : t_n$ if $\exists t_i$ s.t. $e_i : t_i, i = 1, \dots, n$

evaluate left \rightarrow right

if e_i loops forever / raises exc, overall expression loops forever / raises exc

e.g. $(\text{print}(\dots); \text{ref } 10) \rightarrow \text{ref } 10$

Store: the set of accessible reference cells & their contents

$\{e; s\} \Rightarrow \{e'; s'\}$ with e, e', s, s' stores

$e \approx e'$ independent of store: $\{e; s\} \Rightarrow \{v; s'\}$ and $\{e'; s\} \Rightarrow \{v; s'\}$ with v a value & initial stores s

Race condition

fun deposit $a \ n = a := !a + n : \text{int ref} \rightarrow \text{int} \rightarrow \text{int ref}$

fun withdraw $a \ n = a := !a - n$

val $chk = \text{ref } 100$

val $- = (\text{deposit } chk \ 50; \text{withdraw } chk \ 50)$

$!chk \rightarrow 70$

val $- = (\text{deposit } chk \ 50, \text{withdraw } chk \ 50)$

no definitive $!chk$

	Persistent	Ephemeral
Sequential	Functional programming	FP is fine, reasoning about
Parallel	is gradual	concurrency??


```
(* reach : graph → int * int → bool *)
```

```
fun reach (g: graph) (x, y) =
```

```
  let
```

```
    fun dfs n = (n = y) or else (list.exists dfs (g n))
```

```
  in
```

```
    dfs x
```

```
  end
```

check if y is reachable from any of n's neighbors

⚠ possible infinite loop

```
(* mem : graph → int * int → bool *)
```

```
fun mem (n: int) = list.exists (fn x => n = x)
```

```
fun reachable (g: graph) (x, y) =
```

```
  let
```

```
    val visited = ref []
```

```
    fun dfs n = (n = y) or else (not (mem n (!visited)) andalso (visited := n :: (!visited);
```

```
  in
```

```
    dfs x
```

```
  end
```

```
list.exists dfs (g n)))
```

```
signature RANDOM =
```

```
sig
```

```
  type gen (* abstract *)
```

```
  val init : int → gen
```

```
  val random : gen →
```

```
end
```

```
structure R :> RANDOM =
```

```
struct
```

```
  type gen = real ref
```

```
  val a = 16817.0
```

```
  val m = 2147483647.0
```

```
  fun next r = a * r - m * real (floor (a * r / m))
```

```
  val init = ref 0.0
```

```
  fun random y b = (g := next (!g); floor ((!g / m) * (real b)))
```

```
end
```

Stream Memorization

```
fun delay d =
```

```
  let
```

```
    val cell = ref d
```

```
    fun memoFn () =
```

```
      let
```

```
        val r = d()
```

```
      in
```

```
        (cell := (fn () => r); r)
```

```
      end handle E => (cell := (fn () => raises E); raise E)
```

```
    val _ = cell := memoFn
```

```
  in
```

```
    Stream (fn () => !cell())
```

```
  end
```