

In my example, I use the function y=+x+3 from [0,1] this means fraj=3 & fraj=2 50 Fraj.fraj 40 of 640 condition is not held for the bisection algo to guarantee a root. Also the bisection algorith checks to see if feat is the same sign as f(P), Because Frankland frank frank frank positive this statement will ALWAYS be called before the else Statement. Hence F(o;) is inching towards Overall the also will find the Closest number to zero in the interval, but not 0 if the fas. f(b) Lo condition is violated. Let f(x)=(x-1)10 p=1 & Pn=1+1. Show that | f(Pn) | 4 10-3 Whenever not but that IPn-P1/210-3 requires that no 1000. A represents the iteration number in the bisection algorithm meaning $n \in \mathbb{N}$ too. $|f(P_n)| \leq |o^{-3}| > |f(1+\frac{1}{n})| \leq |o^{-3}|$ => $f(1+\frac{1}{n}) = (1+\frac{1}{n}-1)^{10} = (\frac{1}{n})^{10} = (\frac{1}{n})^{10} = (\frac{1}{n}-1)^{10} = (\frac{1}{n}-1)^{10}$ => logic(n10) < logic(10-3) => logic(n10) < -3 => -logic(n10) <-3 Now lets see - logio (m+1)10) the next step, in this case (it nears - 10,00 (n+1)0) 4-10,10 (n10) 2-3. Lets try n=2. - legio (210) = -3.01 4-3 meaning by induction, this principal 1 holds Let 1P0-P1 < 10-3. \(1+\frac{1}{n}\)-1 \(\2\10^{-3}=\) \\\ \| \\ \| \| \2\10^{-3} => 121 < 13. Because A>1, A cannot be negative. 1 6 103 => 103.1 6 103.103 => 103 K1 => 103 KN 50 N21000. 12

4) If the bisection. Method is applied with starting interval

[2^m, 2^{m+1}], where m is a positive of negative integer; how

many steps should be taken to compute the root to full

Mochine precision on a 32-bit word-length computer?

[2ⁿ | 2^{m+1} | 2^{m+1} | 2^m | 2 | 2 | 2ⁿ⁺¹ | 2ⁿ⁺