

Logic Review

CS 536: Science of Programming, Fall 2019

Due Wed Sep 4, 11:59 pm

A. Why?

- We use propositions and predicates to write program specifications.
- Propositions and predicates can be related or manipulated syntactically or semantically.
- States describe memory; an expression has a value relative to a state.

B. Objectives

At the end of this homework, you should be able to

- Describe the relationship between syntactic equality and semantic equality.
- Translate expressions, propositions, and predicates to and from English.
- Design predicate functions for simple properties on values and arrays.

C. Formatting and Submitting Your Work

- You don't have to use a word processor to write out your answers: Feel free to convert logical symbols into ASCII text: For \wedge , \vee , \rightarrow , \neg , \forall , \exists , write `and`, `or`, `->`, `!`, `all`, and `exist`. For \Rightarrow , \Leftrightarrow , \equiv , and \neq , write `=>`, `<=>`, `==`, and `!=`.

D. Problems [50 points total]

Quantified variables range over \mathbb{Z} unless otherwise specified.

1. [6 = 2 * 3 points] What is the full parenthesization of
 - a. $p \wedge \neg r \wedge s \rightarrow \neg q \vee r \rightarrow \neg p \leftrightarrow \neg s \rightarrow t$?
 - b. $\exists m. 0 \leq m < n \wedge \forall j. 0 \leq j < m \rightarrow b[0] \leq b[j] \leq b[m]^*$
2. [6 = 3 * 2 points] Give the minimal parenthesization of each of the following by showing what remains after removing all redundant parentheses. Hint: To avoid getting confused about which parentheses match each other, try rewriting the given parentheses with subscripts: $(_1 \dots)_1$ versus $(_2 \text{ and })_2$ and so on.
 - a. $((\neg(p \vee q) \vee r) \rightarrow (((\neg q) \vee r) \rightarrow ((p \vee (\neg r)) \vee (q \wedge s))))$
 - b. $(\exists i. (((0 \leq i) \wedge (i < m)) \wedge (\forall j. (((m \leq j) \wedge (j < n)) \rightarrow (b[i] = b[j])))))$. (This predicate asks “Is there a value in $b[0..m-1]$ > every value in $b[m..n-1]$?”)
 - c. $\forall x. ((\exists y. (p \rightarrow q)) \rightarrow (\forall z. (q \vee (r \wedge s))))$

* Leave $(0 \leq j < m)$ as is; don't expand it to $((0 \leq j) \wedge (j < m))$. Don't forget to parenthesize $(b[0])$, e.g.

3. [4 = 2 * 2 points] Say whether the given propositions or predicates are \equiv or \neq . Briefly justify your answer.
- Is $p \wedge q \vee \neg r \rightarrow \neg p \rightarrow q \equiv ((p \wedge q) \vee ((\neg r \rightarrow ((\neg p) \rightarrow q))))$?
 - Is $\forall x . p \rightarrow \exists y . q \rightarrow r \equiv ((\forall x . p) \rightarrow (\exists y . q)) \rightarrow r$?
4. [6 = 2 * 3 points] Say whether each of the following is a tautology, contradiction, or contingency. If it's a contingency, show an instance when the proposition is true and show an instance where it's false.
- $(p \rightarrow (q \rightarrow r)) \leftrightarrow ((p \rightarrow q) \rightarrow r)$
 - $(\forall x \in \mathbb{Z} . \forall y \in \mathbb{Z} . f(x, y) > 0) \rightarrow (\exists x \in \mathbb{Z} . \exists y \in \mathbb{Z} . f(x, y) > 0)$. Rely on the idea that for $(\forall u . \phi)$ to be false, we need some value for u for which ϕ is false. I.e., we need $(\exists u . \neg \phi)$. Similarly, for $(\exists v . \psi)$ to be false, we need ψ to be false for every value of v . I.e., we need $(\forall v . \neg \psi)$.
5. [2 points] Which of the following mean(s) $p \rightarrow q$ and which mean $q \rightarrow p$?
- p is sufficient for q
 - p only if q
6. [6 = 3 * 2 points] Let e_1 and e_2 be expressions.
- In general, does $e_1 \neq e_2$ imply $e_1 \neq e_2$? If yes, briefly justify (a sentence or two is fine); if no, give a counterexample (specific values for e_1 and e_2 that show that this implication does not always hold).
 - In general, does $e_1 = e_2$ imply $e_1 \equiv e_2$? Again give a brief justification or counterexample.
7. [6 points] The goal is to show that $p \wedge \neg(q \wedge r) \rightarrow q \wedge r \rightarrow \neg p$ is a tautology by proving it is $\Leftrightarrow T$. To do this, complete the proof of equivalence below using (only) the propositional logic rules (from Lecture 2). Be sure to include the names of the rules. There's more than one correct answer [just give one of them].
- $$\begin{array}{ll}
 p \wedge \neg(q \wedge r) \rightarrow q \wedge r \rightarrow \neg p & \\
 [you \text{ fill in}] & \text{Defn } \rightarrow \\
 [you \text{ fill in}] & \text{Defn } \rightarrow \\
 [and \text{ so on}] &
 \end{array}$$
8. [6 points] Simplify $\neg(\forall x . (\exists y . x \leq y) \vee \exists z . x \geq z)$ to a predicate that has no uses of \neg . Present a proof of equivalence. You'll need DeMorgan's Laws. Also use rules like " $\neg(e_1 \leq e_2) \Leftrightarrow e_1 > e_2$ by negation of comparison".
9. [6 points] Write the definition of a predicate function $GT(b, x, m, k)$ that yields true iff $x > b[m]$, ... $b[m+k-1]$. E.g., in the state $\{b = (1, 3, -2, 8, 5)\}$, $GT(b, 4, 0, 3)$ is true; $GT(b, 0, 1, 2)$ is false. You can assume without testing that the indexes $m, \dots, m+k-1$ are all in range. If $k \leq 0$, the sequence $b[m], b[m+1], \dots, b[m+k-1]$ is empty and $GT(b, x, m, k)$ is true. (It's straightforward to write GT so that this is not a special case.) Remember, this has to be a **predicate function**, not a program that calculates a boolean value.
- Hint: Check the discussion in Lecture 2 about trying to translate

10. [4 = 2 * 2 points] For each of the following, there are three possibilities:

- The state is *ill-formed* (doesn't meet the criteria for being a state).
- The state is *well-formed but not proper* for the given predicate
- The state is legal, proper, and satisfies the predicate.

Notation: x, y, \dots, i, \dots are variables of type integer. b is a variable of type array of integer. $\alpha, \beta, \gamma, \dots$ are semantic values; so are items spelled out in English like *two plus two*.

- $\{x = \text{ten}, y = \text{eight plus one}\} \models x + y > 0$.
- $\{c = \alpha, d = 2\alpha, e = 3\alpha\}$ (for some α) and $d/c + (0 * z) \leq 0$.