The University of Texas at Dallas CS 6364 Artificial Intelligence Fall 2020

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Homework 3: 200 points (30 points extra-credit)
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Submit only in eLearning

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PROBLEM 1: Inference in Propositional Logic (55 points)

a/ A propositional 2-CNF expression is a conjunction of clauses, each containing exactly two literals, e.g.:

$$(X \vee Y) \wedge (\neg X \vee Z) \wedge (\neg Y \vee W) \wedge (\neg Z \vee G) \wedge (\neg W \vee G)$$

(15 points) Prove using resolution that the above sentence entails G.

Solution:

Clauses:

S1: $(X \vee Y)$

S2: $(\neg X \lor Z)$

S3: $(\neg Y \vee W)$

S4: (¬Z∨ G)

S5: $(\neg W \vee G)$

 $\alpha = G$

 $\neg \alpha = \neg G$

Step1:

Applying resolution to S1 and S3 with Y, we get

 $X1: X \vee W$

Step2:

Pairing S1 and S2 with X, we get

X2: Y ∨ Z Step3:

Pairing X1 and S5, we get

X3: X ∨ G

Step4:

Pairing X2 and S3

 $X4: Z \vee W$

Step5:

Pairing X4 and S5

 $X5: Z \vee G$

Step6:

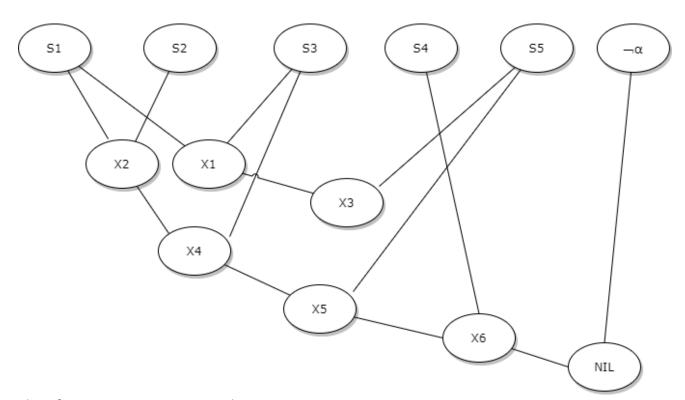
Pairing X5 and S4

X6: G

Step7:

Applying Resolution to X6 and $\neg \alpha$

X7: NIL



Therefore, given sentences entails G

b/ Given the following Knowledge Base:

$$1.~X \wedge B \wedge C \Longrightarrow A$$

$$2.\ A \wedge D \wedge E \Longrightarrow C$$

$$6.~G \land \neg F \Longrightarrow X$$

(15 points) Use backward-chaining inference to prove the query A.

Solution

Observation: Not all sentences in KB are in Horn Clauses!

S4: $E \land \neg F$

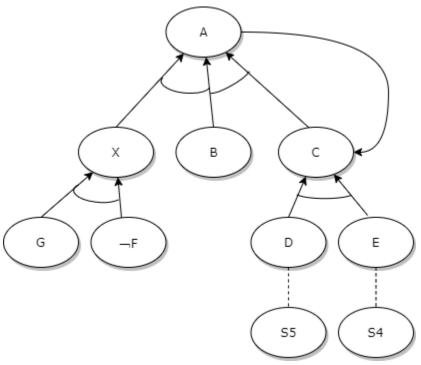
S5: $F \lor D$

We can build truth table for S4 and S5

Е	F	S4
F	H	F
F	Т	F
Т	F	Т
Т	Т	F

D	F	S5
F	F	F
F	Т	Т
Т	F	Т
Т	Т	Т

AND-OR Graph:



<u>Step 1:</u>

Suppose A is True

According to S1, X, B and C are True

Step 2:

B is True from S3 in KB.

<u>Step 3:</u>

Suppose X is True

¬F is True from S6 in KB and G is True from S8 in KB.

Step 4:

Suppose C is True

According to S2 in K: A, D and E are True.

Step 5:

If D = True regardless of F, then D shall be true given S5

Step 6:

If $\neg F$ is True and E = True, then E shall be true given S4

<u>Step 7:</u>

Since we assumed A is True, we get C is True.

Therefore, our assumption that A is True was correct.

c/ Use propositional logic inference rules to decide which of the following sentences are entailed by the Sentence 1: $(X \lor Y) \land (\neg Z \lor \neg W \lor Q)$

Sentence 2: $(X \vee Y)$

Sentence 3: $(X \lor Y \lor Z) \land (Y \land Z \land W \Rightarrow Q)$

Sentence 4: $(X \vee Y) \wedge (\neg W \vee Q)$

To get full credit you need to write if:

i. S1 Entails S2 or not, and if so, which propositional inference rules you have applied to reach the conclusion. Detail the results of each inference rule. (5 points)

Solution:

S1: $(X \vee Y) \wedge (\neg Z \vee \neg W \vee Q)$

S2: $(X \vee Y)$

By using AND elimination on S1, we get S2.

Hence, we can say that S1 entails S2.

ii. S1 Entails S3 or not, and if so, which propositional inference rules you have applied to reach the conclusion. Detail the results of each inference rule. (5 points)

Solution:

S1: $(X \vee Y) \wedge (\neg Z \vee \neg W \vee Q)$

S3: $(X \lor Y \lor Z) \land (Y \land Z \land W \Rightarrow Q)$

By using OR introduction in S1, we get

S1: $(X \lor Y \lor Z) \land (\neg Y \lor \neg Z \lor \neg W \lor Q)$

By using de Morgan, we get

S1: $(X \lor Y \lor Z) \land (\neg (Y \land Z \land W) \lor Q)$

By introducing implication in the above, we get

S1: $(X \lor Y \lor Z) \land (Y \land Z \land W \Rightarrow Q)$

We can say that S1 entails S3.

iii. S1 Entails S4 or not, and if so, which propositional inference rules you have applied to reach the conclusion. Detail the results of each inference rule. (5 points)

Solution:

S1: $(X \vee Y) \wedge (\neg Z \vee \neg W \vee Q)$

S4: $(X \vee Y) \wedge (\neg W \vee Q)$

In S1 and S4, entailment depends on $(\neg Z \lor \neg W \lor Q)$, since we cannot eliminate Z. We can say that S1 does not entail S4.

d/ Demonstrate whether the following sentences are valid, satisfiable, or neither. Motivate and detail your demonstrations.

<u>Sentence 1:</u> ((Smart \lor Beautiful) \Rightarrow (Interesting \lor Boring)) \Leftrightarrow ((Smart \Rightarrow Interesting) \lor (Beautiful \Rightarrow boring)) (5 points)

Solution

We can use Truth table to check for validity and satisfiability

Smar t	Beautif ul	Interestin g	Borin g	Smart ∨ Beautif ul	Interestin g ∨ Boring	((Smart ∨ Beautiful) ⇒ (Interestin g ∨ Boring)	Smart ⇒ Interestin g	Beautif ul ⇒ boring	(Smart ⇒ Interestin g) ∨ (Beautiful ⇒ boring)	Sentenc e 1
Т	Т	Т	Т	Т	Т	Т	Т	Т	T	Т
Т	Т	Т	F	Т	Т	Т	Т	F	T	Т
Т	Т	F	Т	Т	Т	Т	F	Т	T	Т
Т	Т	F	F	Т	F	F	F	F	F	Т
Т	F	Т	Т	Т	Т	Т	Т	Т	Т	Т
Т	F	Т	F	Т	Т	Т	Т	Т	Т	Т
Т	F	F	Т	Т	Т	Т	F	Т	T	Т
Т	F	F	F	Т	F	F	F	Т	T	F
F	Т	Т	Т	Т	Т	Т	Т	Т	T	Т
F	Т	Т	F	Т	Т	Т	Т	F	T	Т
F	Т	F	Т	Т	Т	Т	Т	Т	T	Т
F	Т	F	F	Т	F	F	Т	F	Т	F
F	F	Т	Т	F	Т	Т	Т	Т	Т	Т
F	F	Т	F	F	Т	Т	Т	Т	Т	Т
F	F	F	Т	F	Т	Т	Т	Т	Т	Т
F	F	F	F	F	F	Т	Т	Т	F	F

Since we are getting False for 3 models, the **sentence is not Valid**. Since we are getting True for other models, **the sentence is satisfiable**.

<u>Sentence 2:</u> (Tall \vee Gorgeous) $\vee \neg$ (Tall \Rightarrow Gorgeous) (5 points)

Solution

Gorgeous	Tall	$((Tall \lor Gorgeous) \lor \neg (Tall \Rightarrow Gorgeous))$
F	F	F
F	Т	Т
Т	F	Т
Т	Т	Т

Since we are getting False for 1 model, the **sentence is not Valid**. Since we are getting True for other models, **the sentence is satisfiable**. <u>PROBLEM 2:</u> (25 points) Logic Representations a/ According to political pundits, a person who is a radical (R) is electable (E) if he/she is conservative (C), but otherwise is not electable. Which of the following are correct representations in propositional logic of this assertion? Explain why. (15 points)

i.
$$(R \wedge E) \Leftrightarrow C$$

Solution:

It means that a person is both radical and electable can be conservative and a person who is conservative is both radical and electable. It is not a correct representation of the given assertion.

ii.
$$R \Rightarrow (E \Leftrightarrow C)$$

Solution:

It means that a person who is radial, is electable if and only if the person is conservative. It is a correct representation of the assertion

iii.
$$R \Rightarrow ((C \Rightarrow E) \lor \neg E)$$

Solution:

By applying implication elimination, we get $R \Rightarrow (C \lor E \lor \neg E)$. $E \lor \neg E$ is always True, which means that a person is electable regardless if that person is conservative. Hence it is an incorrect representation.

b/ Unification: For each pair of literals, find the Most General Unifier and the Most General Common Substitution Instance:

(2 points)
$$\{P(x), P(A)\}$$

Solution:

Most General Unifier: $\Theta = \{x/A\}$

Most General Common Substitution Instance: x= A

(4 points)
$$\{P[f(x), y, g(y)], P[f(x), z, g(x)]\}$$

Solution:

Most General Unifier: $\Theta = \{y/x\}$

Most General Common Substitution Instance: y=x

(4 points) $\{P[f(x, g(A,y)), g(A,y)], P[f(x,z),z]\}$

Solution:

Most General Unifier: $\Theta = \{g(A, y)/z\}$

Most General Common Substitution Instance: g(A, y) = z

PROBLEM 3: First-Order Logic (FOL) representations (40 points)

Write in FOL the following statements by defining first your vocabulary (i.e. predicates, constants, variables, functions, etc):

1. (2 points) Some leaves turn red each Fall.

Leaf(x): x is a leaf Red(x): x is red

Season(x): x is season

 $\exists x [\text{Leaf}(x) \land \text{Season}(\text{Fall}) \land \text{Red}(x))]$

2. (3 points) Some trees lose all their leaves when winter comes.

Tree(x): x is a tree

Leaf (x, y): y is a leaf of tree x

Lose (x, y):x is a tree that loses the leaf y

Season (x): x is a season

 $\exists x \ \forall y [Tree(x) \land Leaf(x, y) \land Season(Winter) => Lose(x, y)]$

3. (2 points) Flowers are always nice and they smell lovely.

Flower(x): x is a flower

Nice(x): x is nice

Smell_Lovely(x): x smells lovely

 \forall x[flower(x)=>Nice(x) \land Smell_Lovely(x)]

4. (2 points) One flower does not bring Spring.

Flower(x): x is a flower Season(x): x is a season

 $\exists x \neg [(Flower(x) => Season(Spring)]$

5. (2 points) Every flower fades at some point.

Flower(x): x is a flower

Fade(x): x is something that fades

 \forall x [Flower(x)=> Fade(x)]

6. (2 points) Only one flower is left in the vase.

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Flower(x): x is a flower
```

Vase(x): x is left inside vase

 $\exists x \forall y (Flower(x) => Vase(x)) \land (Flower(y)=> Vase(y) => x=y$

7. (2 points) Every person that buys flowers is sensitive.

Person(x): x is a person Flower(y): y is a flower

Buys (x, y): x buys y

Sensitive(x): x is sensitive

 $\forall x \exists y [Person(x) \land Flower(y) \land Buys (x, y) => Sensitive(x)]$

8. (3 points) Poets are sensitive but they do not buy flowers, they write beautiful poems.

Poet(x): x is a poet

Flower(y): y is a flower Buys (x, y): x buys y

Sensitive(x): x is sensitive

Poem(x): x is a poem

Beautiful(x): x is beautiful

Writes(x,y): x writes y

 $\forall x \text{ Poet}(x) => \text{Sensitive}(x) \land (\forall y \text{ [Flower}(y) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y)]) \land (\forall z \text{ [Poem}(z) \land \text{Writes}(x, z) => \neg \text{Buys}(x, y) => \neg \text{Writes}(x, z$

Beautiful(z)])

9. (2 points) No poet will kill an animal.

Poet(x): x is a poet

Animal(x): x is animal

Kill (x, y): x kills y

 $\forall x,y [Poet(x) \land Animal(y) => \neg Kill (x, y)]$

10. (**3 points**) There is an agent who sells policies only to people who are not insured.

Agent(x): x is agent

Insured(x): x is insured

Sell policy(x, y): x sells policies to y

 $\exists x \ Agent(x) \Rightarrow \forall y \ [\neg \ Insured(y) \Rightarrow Sell_policy(x, y)]$

11. (2 points) There is a barber who shaves all men in town who do not shave themselves.

Barber(x): x is a barber

Man(y): y is a man Shave (x, y): x shaves y $\forall x \forall y [Barber(x) \land Man(y) \land \neg Shave (x, y) => Shave (x, y)]$

12. (**10 points**) A person born outside the US, one of who has at least one parent who is a US citizen by birth is a US citizen by descent.

Person(x): x is a person
Born (x, y): x is born in y
Parent (x, y): x is parent of y
Citizen (x, y): x is citizen of y $\forall x \exists y [Person(x) \land \neg Born (x, US) \land Parent (y, x) \land Citizen (y, US) => Citizen (x, US)]$

13. (**2 points**) There is a flower that smells nice in the house.

Flower(x): x is a flower
House(y): y is a house
Inside (x, y): x is inside y

Smell_Nice(x): x is a flower that smells nice

 $\exists x,y \; Flower(x) \land House(y) \land Inside(x, y) \land Smell-Nice(x)$

14. (**3 points**) John bought only two flowers.

Flower(x): x is a flower Bought(x, y): x bought y $\exists x,y \ [Flower(x) \land Bought(John, x) \land Bought(John, y) \land Flower(y)] \land \forall z \ [Bought(John, z) \land Flower(z)] => (x=z \lor y=z)]$

PROBLEM 4: Refutation in First-Order Logic (80 points)

The purpose of this assignment is to give you experience in proving facts with the resolution method and in exposing you to Prover9, an automatic theorem prover that can help you devise your refutations.

Consider the following helpful pointers for using Prover9:

Installation (https://www.cs.unm.edu/~mccune/mace4/gui/v05.html) For linux users, install python-wxtools also.

Help Manual (https://www.cs.unm.edu/~mccune/prover9/manual/2009-02A/)

Simple tutorial (www.cs.utsa.edu/~bylander/cs5233/prover9-intro.pdf)

You are asked to solve the following puzzle.

- 1. Anyone who rides a Harley is a rough character.
- 2. Every biker rides [something that is] either a Harley or a BMW.
- 3. Anyone who rides any BMW is a yuppie.
- 4. Every yuppie is a lawyer.
- 5. Any nice girl does not date anyone who is a rough character.
- 6. Mary is a nice girl, and John is a biker.
- 7. (Conclusion) If John is not a lawyer, then Mary does not date John.
- i. (14 points) Represent these clauses in first order logic, using only these predicates:

Harley(x), Rides(x,y), Rough(x), Biker(x), BMW(x), Yuppie(x), Lawyer(x), Nice(x), Date(x,y)

Solution:

- 1. $\forall x \ \forall y \ [Rides (x, y) \land Harley(y) => Rough(x)]$
- 2. $\forall x [Biker(x) => \exists y Rides (x, y) \land (BMW(y)V Harley(y))]$
- 3. $\forall x \ \forall y \ [Rides (x, y) \land BMW(y) => Yuppie(x)]$
- 4. $\forall x [Yuppie(x) => Lawyer(x)]$
- 5. $\forall x \forall y [Nice(x)=> \neg (Date(x, y) \land Rough(y))]$
- 6. Nice (Mary) ∧Biker (John)
- 7. ¬Lawyer(John)=> ¬Date(Mary, John)
- ii. (14 points) Convert the logic sentences to clause form, skolemizing as necessary.
 - 1. Removing ∀ and Applying implication elimination and de Morgan.
 - S1: \neg Ride (x, y) $\vee \neg$ Harley(y) \vee Rough(x)
 - 2. Eliminate \exists by skolemizing using y = f(x)
 - S2: \neg Biker(x) \vee Rides(x, f(x))
 - S3: $\neg Biker(x) \lor Harley(y) \lor BMW(y)$
 - 3. Applying implication elimination and de Morgan
 - S4: $\neg Rides(x, y) \lor \neg BMW(y) \lor Yuppie(x)$
 - 4. Applying implication elimination and de Morgan
 - S5: \neg Yuppie(x) \vee Lawyer(x)
 - 5. Applying implication elimination and de Morgan
 - S6: $\neg Nice(x) \lor \neg Date(x, y) \lor \neg Rough(y)$
 - 6. S7: Nice (Mary)
 - 7. S8: Biker (John)

iii. (42 points) Prove by hand whether the conclusion is true by using resolution refutation (i.e. negate the conclusion and show its unsatisfiability with the rest of the knowledge base). Make sure to document the substitutions you use.

```
Negating the conclusion:
        S9: ¬Lawyer(John)
        S10: Date(Mary, John)
Resolving S2 and S8, we get:
X1: Rides (John, f(John))
\Theta = \{x/ John\}
Applying resolution to S3 and S8, we get:
X2: Harley(f(John)) \vee BMW(f(John))
\Theta = \{x/ John\}
Applying resolution to S1 and X1, we get:
X3: \negHarley(f(John)) \vee Rough (John)
\Theta = \{x/John, y/f(John)\}
Resolving S4 and X1:
X4: BMW(f(John)) \times Yuppie (John)
\Theta = \{x/John, y/f(john)\}
Applying resolution to S6 and S7, we get:
X5: \negDate (Mary, y) \vee \negRough(y)
\Theta = \{x/Mary\}
Resolving S10 and X5, we get:
X6: ¬Rough (John)
\Theta = \{y/John\}
Resolving S5 and S9:
X7: ¬Yuppie (John)
\Theta = \{x/John\}
Resolving X4 and X7:
X8: \negBMW(f(John))
Resolving X3 and X6:
X9: ¬Harley (f(John))
```

```
Resolving X2 and X8: X10: Harley (f (John))
```

Pairing X10 and X9: NIL

Therefore, if John is not a lawyer, Mary does not date him

iv.(**20 points)** Use Prover9 to perform automatically the refutation. Submit a report with three parts:

- I. Assumptions and goal;
- II. The input and output of prover9 (The input of prover 9 should be in plain text)
- III. Conclusion

Solution:

```
Assumptions:
```

```
\label{eq:continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous
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Goal:

```
-lawyer(John) -> -date(Mary, John).
```

Input and Output:

```
yuppie(x) \rightarrow lawyer(x).
nice(x) \rightarrow -(date(x,y) \& rough(y)).
nice (Mary) & biker (John).
end of list.
formulas (goals).
-lawyer(John) -> -date(Mary, John).
end of list.
======= end of input
______
% Enabling option dependencies (ignore applies only on input).
_____
% Formulas that are not ordinary clauses:
1 rides(x,y) & harley(y) \rightarrow rough(x) # label(non clause).
[assumption].
2 biker(x) \rightarrow (exists y rides(x,y)) & (bmw(y) | harley(y)) #
label (non clause). [assumption].
3 rides(x,y) & bmw(y) \rightarrow yuppie(x) # label(non_clause).
[assumption].
4 yuppie(x) \rightarrow lawyer(x) # label(non clause). [assumption].
5 nice(x) \rightarrow -(date(x,y) & rough(y)) # label(non clause).
[assumption].
6 nice(Mary) & biker(John) # label(non clause). [assumption].
7 -lawyer(John) -> -date(Mary, John) # label(non clause) #
label(goal). [goal].
==========
% Clauses before input processing:
formulas (usable).
end of list.
formulas(sos).
-rides(x,y) | -harley(y) | rough(x). [clausify(1)].
-biker(x) | rides(x, f1(x, y)). [clausify(2)].
-biker(x) | bmw(y) | harley(y). [clausify(2)].
-rides(x,y) | -bmw(y) | yuppie(x). [clausify(3)].
-yuppie(x) | lawyer(x). [clausify(4)].
```

```
-\text{nice}(x) \mid -\text{date}(x,y) \mid -\text{rough}(y). [clausify(5)].
nice(Mary). [clausify(6)].
biker(John). [clausify(6)].
-lawyer (John). [deny(7)].
date(Mary, John). [deny(7)].
end of list.
formulas (demodulators).
end of list.
===========
Eliminating rides/2
8 -biker(x) | rides(x, f1(x, y)). [clausify(2)].
9 -rides(x,y) | -harley(y) | rough(x). [clausify(1)].
Derived: -biker(x) \mid -harley(fl(x,y)) \mid rough(x).
[resolve(8,b,9,a)].
10 -rides(x,y) | -bmw(y) | yuppie(x). [clausify(3)].
Derived: -bmw(f1(x,y)) \mid yuppie(x) \mid -biker(x).
[resolve(10, a, 8, b)].
Eliminating biker/1
11 biker(John). [clausify(6)].
12 -biker(x) | bmw(y) | harley(y). [clausify(2)].
Derived: bmw(x) \mid harley(x). [resolve(11, a, 12, a)].
13 -biker(x) | -harley(f1(x,y)) | rough(x). [resolve(8,b,9,a)].
Derived: -harley(f1(John,x)) \mid rough(John). [resolve(13,a,11,a)].
14 -bmw(f1(x,y)) | yuppie(x) | -biker(x). [resolve(10,a,8,b)].
Derived: -bmw(f1(John,x)) \mid yuppie(John). [resolve(14,c,11,a)].
Eliminating yuppie/1
15 -bmw(f1(John,x)) \mid yuppie(John). [resolve(14,c,11,a)].
16 -yuppie(x) | lawyer(x). [clausify(4)].
Derived: -bmw(f1(John,x)) \mid lawyer(John). [resolve(15,b,16,a)].
Eliminating nice/1
17 nice (Mary). [clausify(6)].
18 -\text{nice}(x) \mid -\text{date}(x, y) \mid -\text{rough}(y). [clausify(5)].
Derived: -date(Mary,x) \mid -rough(x). [resolve(17,a,18,a)].
Eliminating lawyer/1
19 -bmw(f1(John,x)) | lawyer(John). [resolve(15,b,16,a)].
20 -lawyer(John). [deny(7)].
Derived: -bmw(f1(John,x)). [resolve(19,b,20,a)].
Eliminating date/2
21 -date(Mary,x) | -rough(x). [resolve(17,a,18,a)].
```

```
22 date (Mary, John). [deny(7)].
Derived: -rough(John). [resolve(21,a,22,a)].
Eliminating bmw/1
23 -bmw(f1(John,x)). [resolve(19,b,20,a)].
24 bmw(x) | harley(x). [resolve(11,a,12,a)].
Derived: harley(f1(John,x)). [resolve(23,a,24,a)].
Eliminating harley/1
25 harley(f1(John,x)). [resolve(23,a,24,a)].
26 -harley(fl(John,x)) \mid rough(John). [resolve(13,a,11,a)].
Derived: rough (John). [resolve (25, a, 26, a)].
Eliminating rough/1
27 rough (John). [resolve (25, a, 26, a)].
28 -rough (John). [resolve (21, a, 22, a)].
Derived: $F. [resolve(27, a, 28, a)].
======================= end predicate elimination
Auto denials: (no changes).
Term ordering decisions:
Predicate symbol precedence: predicate order([ ]).
Function symbol precedence: function order([]).
After inverse order: (no changes).
Unfolding symbols: (none).
Auto inference settings:
  % set(neg binary resolution). % (HNE depth diff=0)
  % clear(ordered res). % (HNE depth diff=0)
  % set(ur resolution). % (HNE depth diff=0)
   % set(ur resolution) -> set(pos ur resolution).
    % set(ur resolution) -> set(neg ur resolution).
Auto process settings: (no changes).
======== PROOF
% Proof 1 at 0.01 (+ 0.01) seconds.
% Length of proof is 29.
% Level of proof is 8.
% Maximum clause weight is 0.
% Given clauses 0.
```

```
1 rides(x,y) & harley(y) \rightarrow rough(x) # label(non clause).
[assumption].
2 biker(x) \rightarrow (exists y rides(x,y)) & (bmw(y) | harley(y)) #
label(non clause). [assumption].
3 rides(x,y) & bmw(y) \rightarrow yuppie(x) # label(non clause).
[assumption].
4 yuppie(x) -> lawyer(x) # label(non clause). [assumption].
5 nice(x) \rightarrow -(date(x,y) & rough(y)) # label(non clause).
[assumption].
6 nice(Mary) & biker(John) # label(non clause). [assumption].
7 -lawyer(John) -> -date(Mary, John) # \overline{label}(non clause) #
label(goal). [goal].
8 -biker(x) | rides(x, f1(x, y)). [clausify(2)].
9 -rides(x,y) | -harley(y) | rough(x). [clausify(1)].
10 -rides(x,y) | -bmw(y) | yuppie(x). [clausify(3)].
11 biker(John). [clausify(6)].
12 -biker(x) | bmw(y) | harley(y). [clausify(2)].
13 -biker(x) | -harley(f1(x,y)) | rough(x). [resolve(8,b,9,a)].
14 -bmw(f1(x,y)) | yuppie(x) | -biker(x). [resolve(10,a,8,b)].
15 -bmw(f1(John,x)) \mid yuppie(John). [resolve(14,c,11,a)].
16 -yuppie(x) | lawyer(x). [clausify(4)].
17 nice (Mary). [clausify(6)].
18 -\text{nice}(x) \mid -\text{date}(x,y) \mid -\text{rough}(y). [clausify(5)].
19 -bmw(f1(John,x)) \mid lawyer(John). [resolve(15,b,16,a)].
20 -lawyer(John). [deny(7)].
21 -date (Mary, x) | -rough(x). [resolve (17, a, 18, a)].
22 date(Mary, John). [deny(7)].
23 -bmw(f1(John,x)). [resolve(19,b,20,a)].
24 bmw(x) | harley(x). [resolve(11, a, 12, a)].
25 harley(f1(John,x)). [resolve(23,a,24,a)].
26 -harley(f1(John,x)) \mid rough(John). [resolve(13,a,11,a)].
27 rough (John). [resolve(25, a, 26, a)].
28 -rough (John). [resolve (21,a,22,a)].
29 $F. [resolve(27,a,28,a)].
====== end of proof
_____
Given=0. Generated=1. Kept=0. proofs=1.
Usable=0. Sos=0. Demods=0. Limbo=0, Disabled=22. Hints=0.
Weight deleted=0. Literals deleted=0.
Forward subsumed=0. Back subsumed=0.
Sos limit deleted=0. Sos displaced=0. Sos removed=0.
New demodulators=0 (0 lex), Back demodulated=0. Back unit deleted=0.
Demod attempts=0. Demod rewrites=0.
```

Conclusion:

From lines 27,28, we can conclude that, if John is not a lawyer, Mary does not date him.

Extra-credit: (30 points) Use Prover9 to automatically perform the refutation of the following:

The Pigs and Balloons Puzzle

- (1) All, who neither dance on tight ropes nor eat penny-buns, are old.
- (2) Pigs, that are liable to giddiness, are treated with respect.
- (3) A wise balloonist takes an umbrella with him.
- (4) No one ought to lunch in public who looks ridiculous and eats pennybuns.
- (5) Young creatures, who go up in balloons, are liable to giddiness.
- (6) Fat creatures, who look ridiculous, may lunch in public, provided that they do not dance on tight ropes.
- (7) No wise creatures dance on tight ropes, if liable to giddiness.
- (8) A pig looks ridiculous, carrying an umbrella.
- (9) All, who do not dance on tight ropes, and who are treated with respect are fat.

Show that no wise young pigs go up in balloons.

-Lewis Carroll, Symbolic Logic,

Submit a report with three parts:

Assumptions and goal;

II.	The input and output of prover9 (The input of prover 9 should be in plain text) III. Conclusion