

## Array

→ What is an array? → List → Collection of similar data type.

⊛ It can be defined as a data structure that is used to store a group or collection of items or elements sequentially inside a memory.

→ Syntax and declaration

int x [100];      X →

--	--	--	--	--	--	--	--	--	--

0   1   2   3   4   5   6   7   8   9 → indices

data-type      array-name      array-size

→  $x[3] = 5;$  → This will assign 5 value to the 3<sup>rd</sup> index of array x.

→  $x[3] = 8;$  → This will update the 3<sup>rd</sup> index of array x to 8.

⇒ Array elements can be accessed by using indices

⊛ eg →  $arr[4];$

→ Printing all the elements of an array -

```
int arr[7] = {2, 4, 5, 7, 8, 10, 13};
for (int i = 0; i <= 6; i++) {
    cout << arr[i] << " ";
}
```

Ans

Output → 2 4 5 7 8 10 13

→ Taking input (Using for loop)

```
for (int i = 0; i <= 6; i++) {
    cin >> arr[i];
}
```

Q) Given an array of marks of students, if the marks of any student is less than 35 then print its roll no. [roll no. = index]

→ int main() {

int marks[6];

// input

```
for (int i = 0; i <= 5; i++) {
    cin >> marks[i];
}
```

```
for (int i = 0; i <= 5; i++) {
```

```
    if (marks[i] < 35) cout << i << " ";
}
```

}

→ In this question we can also take the no. of students manually as n and can run the for loop from i = 0 to i <= n-1 to take input and print output.

## → Types of arrays

- 1) 1 dimensional array
- 2) 2 dimensional array  $\rightarrow$  matrix

## → Size and ~~size~~ of size of operator

$\Rightarrow$  `int arr[] = {2, 3, 4, 1, 2, 9, 10, 3, 14, 15, 20};`

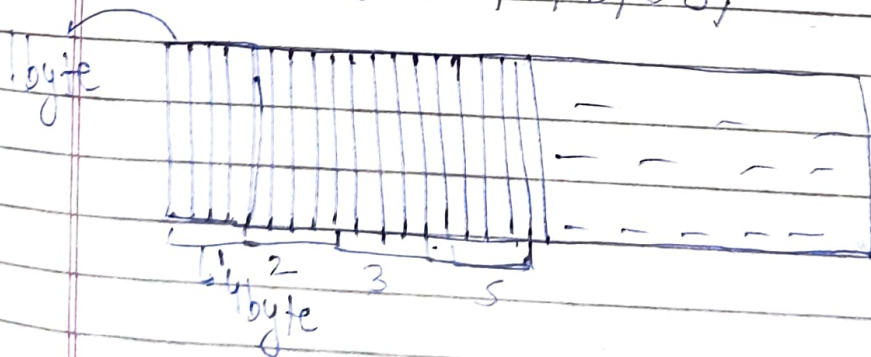
`int n = size of (arr) / size of (arr[0]);`

$\hookrightarrow$  This will be the length of the array

## $\Rightarrow$ Memory Allocation in arrays

→ Continuous memory allocation

→ `int arr[4] = {1, 2, 3, 5};`



→ `&arr` or `&arr[0]` or `&arr[0]`

$\hookrightarrow$  This will give us the address of first byte of first element of the array



→ Calculate the sum of all elements in the given array.

→ Approach

```
int sum = 0;
```

```
for (int i = 0; i <= size - 1; i++) {
```

```
    sum = sum + arr[i];
```

```
}
```

```
cout << sum;
```

⇒ Linear Search

→ Find the element  $x$  in the array. Take array and  $x$  as input.

→ Approach

```
int n;
```

```
cin >> n;
```

```
int arr[n];
```

// Take inputs of the array elements

```
int x;
```

```
cin >> x;
```

// Loop through array elements  
if (arr[i] == x) cout << "Element is present";

→ Find maximum value  $\text{arr}$  of all the elements in an array

# Approach

arr	2	3	5	10	13	15
	0	1	2	3	4	5

int max = arr[0];

// loop from  $i=1$  to  $i \leq n-1$

if (arr[i] > max)

max = arr[i]

}

→ cout << max;

→ Find second largest element

# Approach

① Find Max element

② Traverse through the array,

if (smax < arr[i] && arr[i] != max)

smax = arr[i];

⇒ Pointers in array

h int arr[] = {4, 2, 3, 4}

int\* ptr = arr; → store the address of first element of an array.

int\* ptr1 = &arr;

→ Now, we can access the array elements using pointer.

Eg. `ptr[0] = ;` // It will give us the element at 0<sup>th</sup> index.

⇒ Note → Whenever we pass an array to a function, we pass the address of the first element of the array.

→ Printing the array using pointer

```
int arr[] = {2, 4, 6, 8}
int* ptr = arr;
```

```
for (int i = 0; i <= 4arr; i++) {
    cout << *ptr << " ";
    ptr++;
}
```

⇒ Vectors in C++  
 ( — x — x — )  
 → Dynamic array

→ Syntax

declaration → `vector <data-type> name;`

## → Basic operations on vector

→ `vector<int> v;`

① `v.push_back(6);` → `v[6]`

`v.push_back(1);` → `v[6][1]` (Size increased by  $\times 2$ )

`v.push_back(9);` → `v[6][9]` ( " )

`v.push_back(0);` → `v[6][9][0]`

→ `v.size` → It will give the no. of elements in the vector

→ `v.capacity` → It gives us the total capacity of the vector (It increase by  $\times 2$ )

② `v.pop_back();` → Remove the last element of the vector

→ To print vector elements

③ `for (int i=0; i<v.size(); i++) {`  
`cout << v[i] << " ";`  
`}`

// `v.pop_back();` → updates the size but not the capacity.



→ We can ~~also~~ also utilize vector with size

(\*) `vector<int> v(5);`

↳ Size of vector is 5

(\*) `vector<int> v(5, 7);`

↳ all the places in vector will be assigned 7.

→ Taking input in a vector

(\*) `vector<int> v(5);`

(\*) `for (int i = 0; i < 5; i++) {`  
`cin >> v[i];`  
`}` // Only possible if we have declared the size of the vector.

# If we have not declared the size

(\*) `for (int i = 0; i < n; i++) {`  
`int x;`  
`cin >> x;`  
`v.push_back(x);`  
`}` ↳ Can be asked to user

(3) `v.at(index)` → it gives the element present at a that index



④ `sort(v.begin(), v.end());`

This will sort all the elements in the ascending order.

⇒ Passing Vectors to function

→ Normally vectors are passed by ~~change~~ values in a function unlike array that gets passed by reference.

→ But we can also pass it by reference using (& → ampersand)

⇒ `void change (Vector<int> a) {`  
`a[0]=100;`  
`}`

`int main () {`

`Vector<int> v;`

`v.push-back (10);`

`change (v); // Pass by value`

`}`

⇒ `void change (Vector<int> &a) {`  
`a[0]=100`  
`}`

`int main () {`

`Vector<int> v;`

`v.push-back (10);`

`change (v); // Pass by reference`

`}`

~~WAP to reverse the array without using any extra array.~~

→ WAP to reverse the array without using any extra array.

Approach

	0	1	2	3	4	5	6	7
v	6	1	2	4	3	4	3	1
	↑ i							↑ j

→ swap(v[i], v[j])

→ i++, j--

	0	1	2	3	4	5	6	7
v	1	1	2	4	3	4	3	6
		↑ i					↑ j	

→ Again repeat the process.

⊛ This is known as 2 pointers approach.

⇒ Reverse part of array

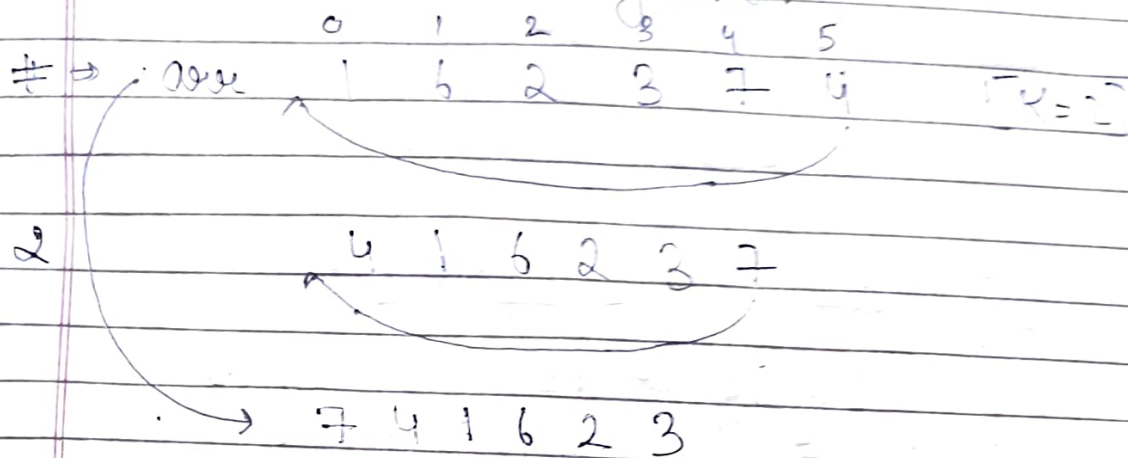
arr → 1 6 2 3 7 4

reverse → 4 7 3 2 6 1

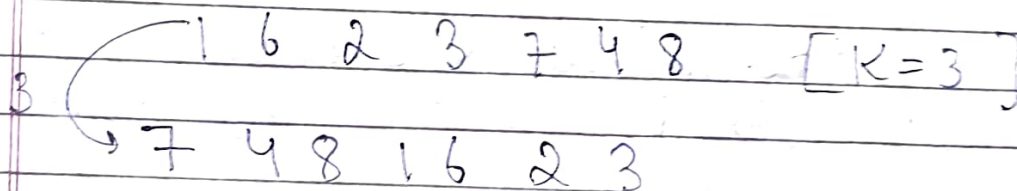
reverse (1, 4) → 1 7 3 2 6 4

reverse (0, 3) → 3 2 6 1 7 4

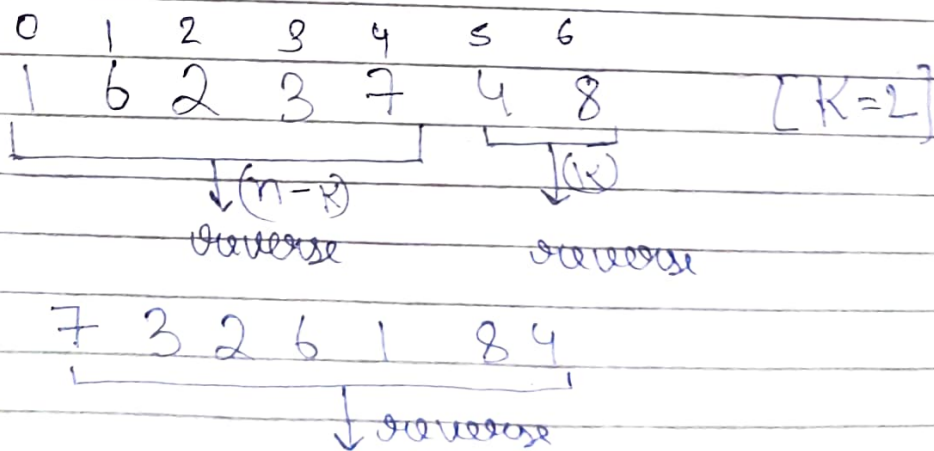
⇒ Rotate the given array  $a$  by  $k$  steps, where  $k$  is non-negative.



→ Algorithm



→ For  $K=2$



① →

reverse Part  $(0, n-k-1, v)$   $n = v.size()$

reverse Part  $(n-k, n-1, v)$

reverse Part  $(0, n-1, v)$



→ Sort the arrays of 0's and 1's.

# int n;

0 1 0 0 1 1 0 1

⇒ Method-I (Two Pass Method)

int noZ = 0; → Count of zeros

int noO = 0; → Count of ones

// noZ = 4

// noO = 4

① → Count the number of zeros and no. of ones in one pass (i.e. in one iteration)

② → In second pass fill the 0th to  $(noZ-1)^{th}$  index with 0 and rest with 1.

→ T.C →  $O(2n)$

⇒ Method-II (Two Pointers Method)

0 1 2 3 4 5 6 7

→ 1 1 0 0 1 1 0 1

① Make 2 pointer variables

int i = 0;

int j = n-1; // n → no. of elements

② while ( $i < j$ ) {

```

    if (arr[j] == 1) j--;
    if (arr[i] == 0) i++;
    if (arr[i] == 1 && arr[j] == 0) swap(i, j);
    i++;
    j--;
}

```

⇒ Sort the array of 0's, 1's and 2's

(Leetcode - 75)

→ Method 1 (Two Pass Method)

void sortColors (vector<int> & nums) {

int n = nums.size();

int no2 = 0;

int no0 = 0;

int no1 = 0;

for (int i = 0; i < n; i++) {

if (nums[i] == 0) no2++;

else if (nums[i] == 1) no0++;

else no1++;

}

// Fill

for (int i = 0; i < n; i++) {

if (i < no2) nums[i] = 0;

else if (i < (no2 + no0)) nums[i] = 1;

else nums[i] = 2;

}

return;

⇒ Method-2 (3-Pointer Algorithm) (Dutch flag algorithm)

→ Approach →  
`int lo = 0;`  
`int mid = 0;`  
`int hi = n-1;` //  $n \rightarrow$  size of array

→ ~~Overall~~

`0 - lo-1` → all the 0's

`hi+1 - n-1` → all the 2's

⇒ `while (mid <= hi) {`

① `if (nums[mid] == 2) {`  
`swap (mid, hi);`  
`hi--;`  
`}`

② `if (nums[mid] == 0) {`  
`swap (lo, mid);`  
`mid++;`  
`lo++;`

③ `if (nums[mid] == 1) {`  
`mid++;`

`}`



⇒ Merge Two sorted arrays (Leetcode - 88)

arr1 → 1 4 5 8

arr2 → 2 3 6 7 10

→ (Sorted)

→ Merge these 2 sorted array by creating another array.

⇒ Algorithm

→ arr1 

0	1	2	3
1	4	5	8

 → Size → n  
i

arr2 

0	1	2	3	4
2	3	6	7	10

 Size → m  
j

arr3 

0	1	2	3	4	5	6	7	8
1	2	3	4	5	6	7	8	10

 → Size → n+m  
k

```

1. int i=0, j=0, k=0;
   while (i < n & j < m) {
       if (arr1[i] < arr2[j]) {
           arr3[k] = arr1[i];
           i++;
           k++;
       }
       else {
           arr3[k] = arr2[j];
           j++;
           k++;
       }
   }

```

```

if (i == n) {
    arr1 pooria with gya
    while (j < m)
        arr3
        arr3[k] = arr2[j];
        j++;
        k++;
    }
}

```

```

if (j == m) { // arr2 pooria with gya
    while (i < n) {
        arr3
        arr3[k] = arr1[i];
        k++;
        i++;
    }
}
}

```

⇒ Find the next permutation of array  
(Leetcode - 31)

→ Algorithm

① Find Pivot index

```

int idx = -1;
for (int i = n-2; i >= 0; i--) {
    if (arr[i] < arr[i+1]) {
        idx = i;
    }
}

```

- ② if ( $idx == -1$ )  
    reverse / sort the whole array.
- ③ Find the element just greater than the  $(idx)^{th}$  element from  $idx+1$  to end.
- ④ Swap that element with  $arr[idx]$ ;