

⇒

Space Complexity

Order of growth of memory (or RAM) space in terms of input size.

```
→ int getSum (int n)
{
    return n * (n+1) / 2;
}
```

$[O(1)] \leq$

or

$[O(1)] \leq$

```
→ int getSum2 (int n)
{
    int sum = 0;
    for (int i = 1; i <= n; i++)
        sum = sum + i;
    return sum;
}
```

→ $[O(1)] \leq$ as we need only 3 variables
i.e. n , sum & i

```
→ int arrSum (int arr[], int n)
{
```

```
    int sum = 0;
```

```
    for (int i = 0; i < n; i++)
```

```
        sum = sum + arr[i];
```

```
    return sum;
}
```

$O(n)$ as the space is growing linearly
 $[n \times arr[i]] \leq$

⇒

Auxiliary Space → Order of growth of extra space or temporary space in terms of input size.

```

→ int arrSum (int arr[], int n)
{

```

```

    int sum = 0;
    for (int i = 0; i < n; i++)
        sum = sum + arr[i];
    return sum;
}

```

Space complexity = $O(n)$

Auxiliary space = $O(1)$ as no other extra space is used.

```

→ int fun (int n)
{

```

```

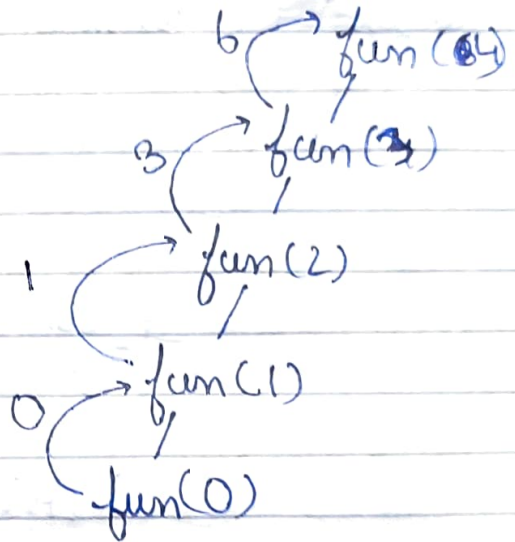
    if (n <= 0)
        return 0;

```

```

    return n + fun(n-1);
}

```



= $[O(n)]$

```

→ int fib (int n)
{

```

```

    if (n == 0 || n == 1)
        return n;

```

```

    return fib(n-1) + fib(n-2);
}

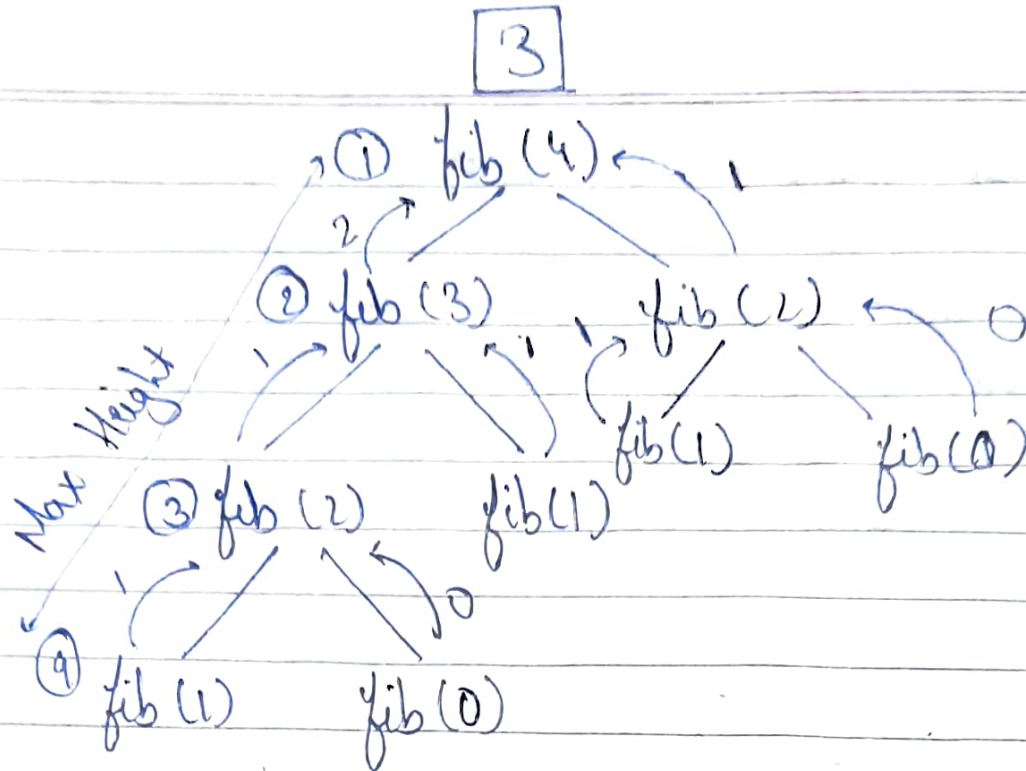
```

0, 1, 1, 2, 3, 5, ...

↓ ↓ ↓

n=0 n=1 n=2

Fibonacci No.'s



→ Recursion formula

(*) Auxiliary space → Max length of stack to leave path.

∴ It is growing linearly
 $[O(n)] \leq$