

Dynamic Pricing

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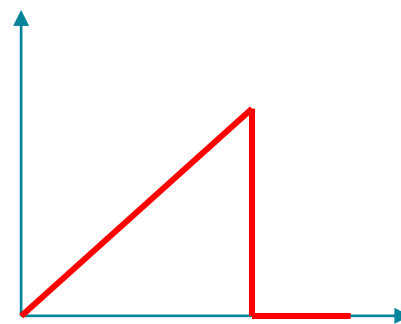
Generic Setting

For $t = 1, 2, \dots, T$:

- Customers generate a **valuation** y_t **secretly**;
- We propose a **price** p_t ;
- If $p_t \leq y_t$, then we get a **reward** $r_t = p_t$;
- Else $p_t > y_t$, then we get a **reward** $r_t = 0$.

3 Properties:

- **Boolean-censored** valuation
- **Bandit** feedback
- **Direction sensitive** reward



Identical or Various Items

- Identical items: no feature
 - Aim at a best **fixed price**.
 - Here ``dynamics'' come from explorations.
- Various items: with features
 - For each time, a vector $x_t \in \mathbb{R}^d$ describes the **feature** of current item, and helps our pricing.
 - Aim at a best ``feature \rightarrow price'' **policy**.
 - Here ``dynamics'' come from different items, and explorations.

Non-feature dynamic pricing

History



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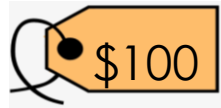


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Deal



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No Deal

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Deal w/ highest price

Non-feature setting

There are mainly three categories:

1. Fixed valuation

y_t is fixed but secret

2. Stochastic valuation

- ① y_t is drawn i.i.d. from a fixed **unknown** distribution.
- ② y_t is drawn from adversarial **unknown** distribution \mathbb{D}_t 's.

3. Adversarial valuation

y_t is obviously arbitrarily chosen.

Non-feature setting -- Regret

Regret is defined as the difference between:

1. Max (expected) reward of a best **fixed** price, i.e.

$$\max_p \sum_{t=1}^n p \cdot 1(p \leq y_t)$$

2. (Expected) reward of our algorithm, i.e.

$$\sum_{t=1}^n p_t \cdot 1(p_t \leq y_t)$$

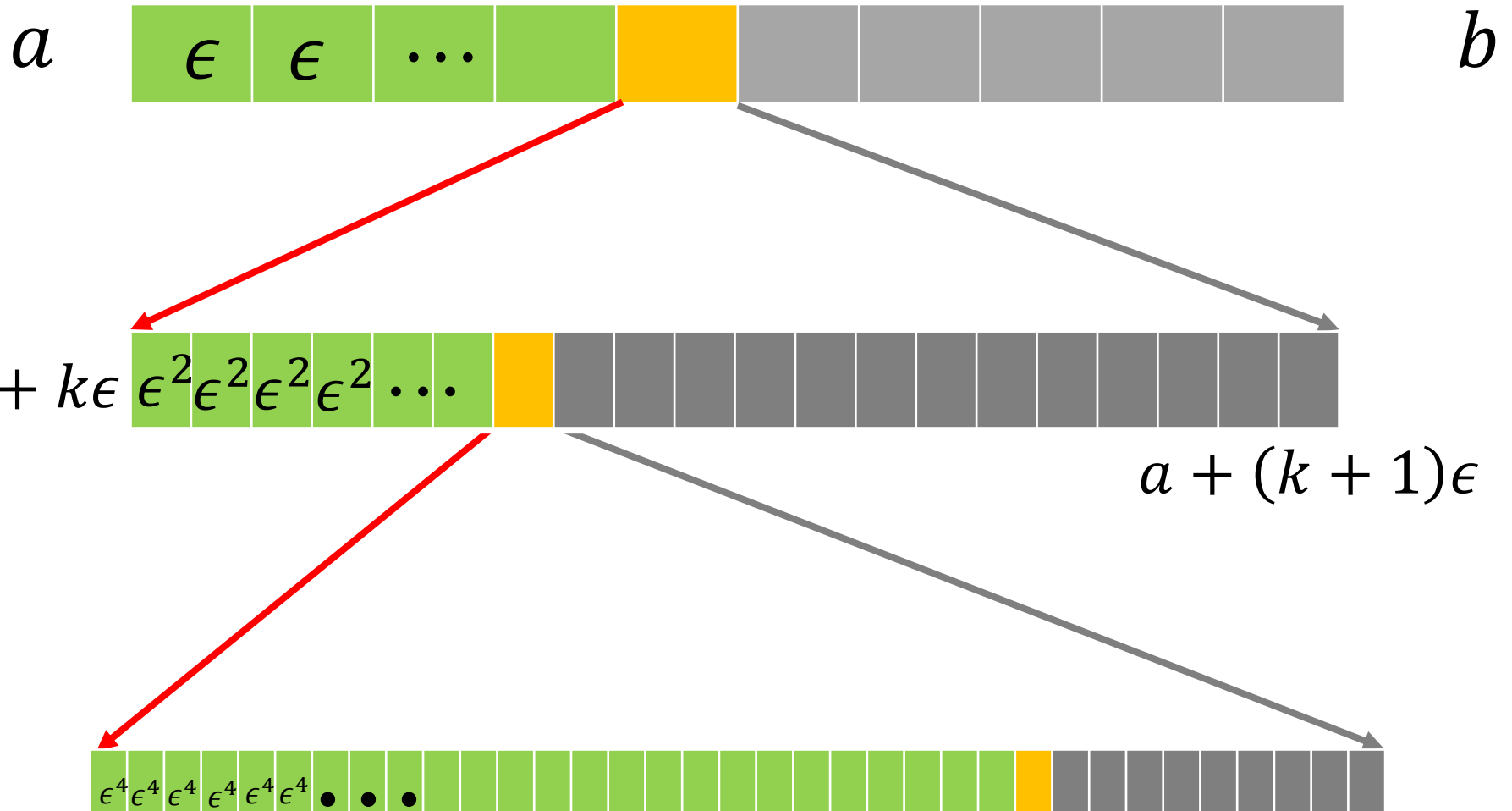
Non-feature setting -- Regrets

	Fixed valuation	Stochastic valuation	Arbitrary valuation
Upper Bound	$O(\log \log T)$	$O(T^{2/3})$ in general $O(\sqrt{T})$ for \mathbb{C}^2 demand curve	$O(T^{2/3})$
	[KL03]	[KL03]	[KL03]
Lower Bound	$\Omega(\log \log T)$	$\Omega(T^{2/3})$ in general, and for Lipschitz demand curve	$\Omega(T^{2/3})$
	[KL03]		[KL03]

Fixed valuation: searching algorithm

- Intuitively, we may apply binary search.
- But binary search can be as expensive as $O(\log T)$.
 - Indeed, a binary search will approach y_t in the fastest rate: at most $O(\log T)$ times of trials.
 - But it might fail in all $O(\log T)$ trials.
- To match the $\Omega(\log \log T)$ lower bound, we instead applies a *squaring search* method.

Squaring search [KL03]



Stochastic/Adversarial valuation: bandit algorithm

- Discretize the price space as: $\{T^{-\alpha}, 2T^{-\alpha}, \dots, 1 - T^{-\alpha}, 1\}$, and play multi-armed bandits.
- Error caused by discretization intervals: $O(T^{1-\alpha})$;
- Error caused by multi-armed bandits: $O(\sqrt{TK}) = O(T^{\frac{1+\alpha}{2}})$
- Minimax regret = $O\left(\min_{\alpha} \max\{T^{1-\alpha}, T^{\frac{1+\alpha}{2}}\}\right) = O(T^{\frac{2}{3}})$.

Feature-based dynamic pricing



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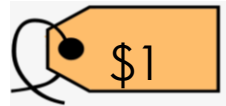


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Deal



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No Deal



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Deal with a high price

History



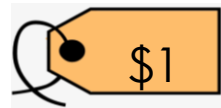
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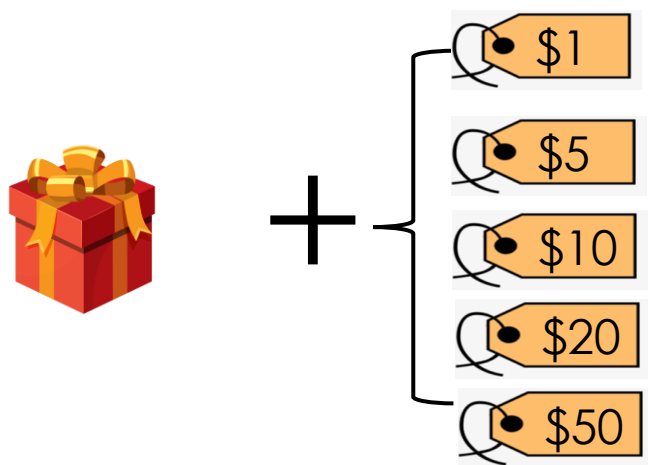
No Deal

⋮



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 \Rightarrow **Deal** w/ highest price



Feature-based setting

We assume a *linear model* for the feature-valuation relationships: $y_t = x_t^\top \theta^* + N_t$, where θ^* is fixed but unknown over time, and N_t can be:

1. $N_t \equiv 0, \forall t = 1, 2, \dots, T$.
2. $N_t \sim_{i.i.d.} \mathbb{D}$, where \mathbb{D} is fixed and could be:
 - ① Totally known (e.g. Standard Gaussian)
 - ② Parametric (e.g. Gaussian, or Laplacian)
 - ③ Totally unknown beside bounds

Feature-based setting -- Regret

In this setting, a regret is defined as the difference between:

1. Max expected reward of an **omniscient** seller, i.e.

$$\sum_{t=1}^n \max_{p_t^*} \mathbb{E}_{N_t \sim \mathbb{D}} [p_t^* \cdot 1(p_t^* \leq x_t^\top \theta^* + N_t) | x_t, \theta^*, \mathbb{D}]$$

2. Expected reward of our algorithm, i.e.

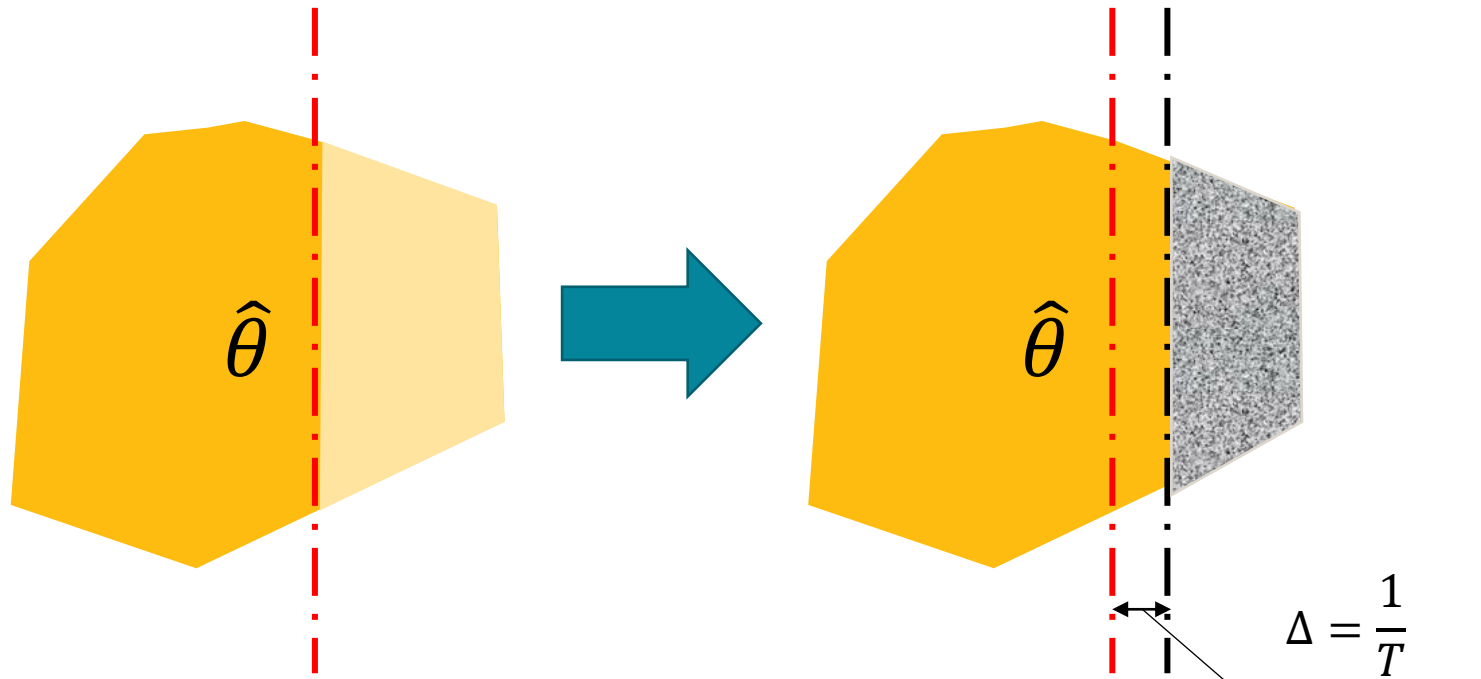
$$\sum_{t=1}^n \mathbb{E}_{N_t \sim \mathbb{D}} [p_t \cdot 1(p_t \leq x_t^\top \theta^* + N_t)]$$

Feature-based setting -- Regrets

	Deterministic	Noisy		
		Known distribution	Parametric distribution	Non-parametric distribution
Upper Bound	$O(\log \log T)$	$O(\log T)$	$O(\sqrt{T})$ for sub-Gaussian noise	$O(T^{\frac{3}{4}})$
	[PLS 18]	[CLPL16] [JN19] $[XW21]$	[JN19]	
Lower Bound	$\Omega(\log \log T)$	$\Omega(\log T)$	$\Omega(\sqrt{T})$ for Gaussian noise	$\Omega(T^{\frac{2}{3}})$
	[KL 03]	[JN19]	[BR12] $[XW21]$	[KL03]

Shallow Pricing: for small-variance noise [CLPL16]

- For noises with small ($O\left(\frac{1}{T \log T}\right)$) variance, we may directly apply binary search with compromission.



EMLP: for Known \mathbb{D} and Stochastic x_t

- Since we **know** the noise distribution, we can use **Max Likelihood Estimate** to achieve a $\hat{\theta}$.
 - $\|\hat{\theta}_t - \theta^*\| = o\left(\frac{1}{\sqrt{t}}\right)$
- We upper bound the regret with $C \cdot \|\hat{\theta}_t - \theta^*\|^2$.
- The regret bound is then $= o\left(\sum_{t=1}^T \frac{1}{t}\right) = o(\log T)$

ONSP: for Known \mathbb{D} and Adversarial x_t

- Since x_t is adversarial, for MLE we do not necessarily have $\hat{\theta} \rightarrow \theta^*$.
- However, we can make use of likelihood functions as *surrogation loss*, i.e. $\text{Regret}(\hat{\theta}) \leq C' \cdot (L(\theta^*) - L(\hat{\theta}))$
 - On the one hand, $L(\theta)$ is exp-concave (strongly convex).
 - On the other hand, we upper bound the regret with $C \cdot \|\hat{\theta}_t - \theta^*\|^2$.
- We apply Online Newton Steps (ONS) to optimize $L(\theta)$.

For unknown noise distribution: why *not* a contextual bandits model?

- Action space
 - In contextual bandits, actions are **discrete**.
 - In dynamic pricings, actions are **continuous**. Extra discretization suffers an error-action trade-off.
- Policy set
 - In contextual bandits, we have a **finite** set of policies.
 - In feature-based dynamic pricing, we **do not** have a finite set of policies.

For unknown noise distribution: our approach

- Continuous action space
 - We discretize the action space at a rational scale.
- Unknown/infinite policy set
 - Define a policy as θ and \mathbb{D} .
 - Use covering number to finitely cover all policies.
 - Use Catalan number to count the finite subset.
 - Apply contextual bandits.
- There is still a huge gap.
 - Upper bound: $O(T^{\frac{3}{4}})$; Lower bound: $\Omega(T^{\frac{2}{3}})$

Reference

- [**KL03**] Kleinberg, R., & Leighton, T. (2003, October). The value of knowing a demand curve: Bounds on regret for online posted-price auctions. In *44th FOCS, 2003. Proceedings.* (pp. 594-605). IEEE.
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