Assignment 5, Optimization methods (Maxim Kaledin)

Problem 3

Apply Sylvester's criterium to ensure that matrix is positive semidefinite. All principal minors should be non-negative, these are

$$-25y_1^2 \ge 0,$$

$$(4y_1 + 2)(-25y_1^2) \ge 0,$$

$$3y_2(4y_1 + 2) \ge 0,$$

$$3y_2 \ge 0,$$

$$4y_1 + 2 \ge 0.$$

The first inequality gives $y_1=0$, and as result constraints are $D=\{y:y_1=0,y_2\geq 0\}$. Obviously, all feasible solutions are optimal and optimal goal is 0.

Let F_1, F_2, G are such symmetric matrices that the initial matrix $A = y_1F_1 + y_2F_2 + G$:

$$F_1 = \begin{bmatrix} 0 & 5 & 0 \\ 5 & 0 & 0 \\ 0 & 0 & 4 \end{bmatrix}, F_2 = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 0 \end{bmatrix}, G = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 2 \end{bmatrix}.$$

Then we can write down the dual problem in the following way (denoting $Z \succeq 0$ as symmetric matrix of dual variables and scalar product as $(A, B) = \operatorname{tr}(A^T B)$):

$$\label{eq:state_equation} \begin{split} \max_{Z\succeq 0} -(Z,G), \quad s.t. \\ (Z,F_1) = 2, (Z,F_2) = 0. \end{split}$$

Rewrite it in the coordinate form:

$$\max_{Z \succeq 0} -2z_{33}, \quad s.t.$$

$$10z_{12} + 4z_{33} = 2, \qquad z_{22} = 0.$$

Now add Sylvester's criterium inequalities

$$z_{11} \ge 0, \ z_{22} \ge 0, \ z_{33} \ge 0,$$

 $z_{11}z_{22} - z_{12}^2 \ge 0,$
 $z_{11}z_{33} - z_{13}^2 \ge 0,$
 $z_{22}z_{33} - z_{23}^2 \ge 0$
 $det(Z) > 0.$

This conditions can be significantly simplified if we add constraints from the dual problem. Finally, we obtain

$$\max_{Z} -2z_{33}, \quad s.t.$$

$$10z_{12} + 4z_{33} = 2,$$

$$z_{22} = 0, z_{12} = 0, z_{23} = 0$$

$$z_{1} \ge 0, z_{33} \ge 0,$$

$$z_{11}z_{33} - z_{13}^{2} \ge 0.$$

Since $z_{12} = 0$ the entry $z_{33} = 1/2$ and we can somehow assign other variables to satisfy feasibility. Thus the optimal dual goal is -1 which is not equal to optimal primal goal 0.

Problem 4

1

The problem can be reformulated as

$$\begin{aligned} & \min_{x,t} t \quad s.t. \\ & t \geq a_i^T x + b_i \quad \forall i \in \{1,..,m\}. \end{aligned}$$

Corresponding Lagrange dual problem is

$$\max_{\lambda} \inf_{x,t} \left[t + \sum_{i=1}^{m} \lambda_i (a_i^T x + b_i - t) \right],$$
s.t. $\lambda \ge 0$,

where $\lambda_i \geq 0$ are dual variables. We can easily simplify this because it is a linear program:

$$\max_{\lambda} b^{T} \lambda, \quad s.t.$$

$$\begin{bmatrix} A \\ -1...-1 \end{bmatrix} \lambda = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ -1 \end{bmatrix}$$

2

Let

$$\hat{f}_2(y) = \frac{1}{\alpha} \log \left(\sum_{i=1}^m \exp \alpha y_i \right),$$

and $y_i = a_i^T x + b_i$ be new constraints. First, formulate the primal problem:

$$\min_{x,y} \hat{f}_2(y)$$
s.t. $y_i = a_i^T x + b_i \quad \forall i \in \{1,..,m\}.$

The dual one is defined as

$$\begin{aligned} & \max_{\lambda} \inf_{x,y} \left[\frac{1}{\alpha} \log \left(\sum_{i=1}^{m} \exp \alpha y_i \right) + \sum_{i=1}^{m} \lambda_i (a_i^T x + b_i - y_i) \right] = \\ & = \inf_{y} \left[\frac{1}{\alpha} \log \left(\sum_{i=1}^{m} \exp \alpha y_i \right) - \sum_{i=1}^{m} \lambda_i y_i \right] + \inf_{x,y} \left[\sum_{i=1}^{m} \lambda_i (a_i^T x + b_i) \right]. \end{aligned}$$

3

Let us show the right side of inequality, in order to do that we can exploit monotonicity of exponent:

$$\frac{1}{\alpha} \log \sum_{i=1}^{m} \exp\left(\alpha(a_i^T x + b_i)\right) - \max_i(a_i^T x + b_i) \le$$

$$\le \frac{1}{\alpha} \log\left(m \exp\left(\alpha \max_i(a_i^T x + b_i)\right)\right) - \max_i(a_i^T x + b_i) =$$

$$= \frac{\log(m)}{\alpha}.$$

Problem 5

Define Lagrange function of a problem $L(x,\lambda) = ||Ax - b||_2^2 + \sum_{i=1}^m \lambda_i (c_i \cdot x - d_i)$. Since there are no inequality constraints, KKT conditions are exactly

$$\frac{\partial L}{\partial x} = 0$$
$$\frac{\partial L}{\partial \lambda} = 0,$$

where derivatives denote the gradients truncated to needed set of variables. These sets of constraints are exactly

$$\frac{\partial L}{\partial x_j} = 2\sum_{i=1}^m a_{ij} \sum_{k=1}^n (a_{ik}x_k - b_i) + \sum_{i=1}^p \lambda_i c_{ij},$$

$$Cx = d.$$

Let us write the first expression shorter:

$$\nabla_x L = 2A^T (Ax - b) + C^T \lambda.$$

Since the goal function is convex, solution to system

$$\begin{cases} 2A^{T}(Ax - b) + C^{T}\lambda = 0, \\ Cx = d \end{cases}$$

exists and is an optimal solution of primal problem. With pseudo-inverse (it exists since C has full row rank):

$$x = (C^T C)^{-1} C^T d,$$

$$\lambda : C^T \lambda = -2A^T (Ax - b).$$

Now derive the dual problem:

$$\max_{\lambda} \inf_{x} L(x, \lambda).$$

Lagrangian L remains convex for all λ , so the solution to inner problem is given by the system of normal equations

$$2A^T(Ax - b) + C^T\lambda = 0$$

and can be found based on any λ since matrix A^TA has full rank. We could write the solution explicitly:

$$x = (A^T A)^{-1} \left(-\frac{1}{2} C^T \lambda + A^T b \right).$$

Problem 6

(1)

Consider problem (4) firstly:

$$\min_{x} c^{T} x$$
s.t. $Ax = b, x \ge 0$.

Its dual problem is

$$\max_{\lambda,\mu} \inf_{x} \left(c^T x + \sum_{i=1}^m \lambda_i (a_i \cdot x - b_i) - \sum_{j=1}^n \mu_j x_j \right),$$
s.t. $\mu \ge 0$.

It can be simplified to

$$\max_{\lambda} b^T \lambda,$$
s.t. $A^T y < c$.

KKT-conditions for primal problem are

- $c + A^T \lambda \mu = 0$ (stationarity, $\nabla_x L(x, \lambda, \mu) = 0$),
- feasibility conditions (as in problem statement),
- Ax = b (from $\nabla_{\lambda} L(x, \lambda, \mu) = 0$),
- $\mu_j x_j = 0 \quad \forall j$ (complementary slackness).

(2)

The idea is the same. Consider problem (5):

$$\min_{x} c^{T} x - \tau \sum_{i=1}^{n} \log(x_{i})$$
s.t. $Ax = b$.

Its dual problem is (let (1/x) denotes coordinate-wise division)

$$\max_{\lambda} \inf_{x} \left(c^{T} x - \tau \sum_{i=1}^{n} \log(x_i) + \sum_{i=1}^{m} \lambda_i (a_i \cdot x - b_i) \right).$$

and KKT-conditions are

- ullet $c- au(1/x)+A^T\lambda=0$ (stationarity, $abla_x L(x,\lambda,\mu)=0$),
- feasibility conditions (as in problem statement),
- Ax = b (from $\nabla_{\lambda} L(x, \lambda, \mu) = 0$).

Since the Hessian matrix of inf sub-problem is positive definite (it is even diagonal), infimum is achieved when neccessary conditions are satisfied (i.e. stationarity):

$$(1/x) = \frac{1}{\tau}(c + A^T \lambda) = y(\lambda).$$

So, simplified dual is

$$\max_{\lambda} \left(c^T(1/y(\lambda)) + \tau \sum_{i=1}^n \log(y(\lambda)_i) + \sum_{i=1}^m \lambda_i (a_i \cdot (1/y(\lambda)) - b_i) \right).$$

After some manipulation we can obtain

$$x_i = \frac{\tau}{c_i + (a_{\cdot i})^T \lambda}$$

and simplify the dual even more

$$\max_{\lambda} \left(\tau \sum_{i=1}^{n} \frac{c_i}{c_i + (a_{\cdot i})^T \lambda} - n\tau \log \tau + \tau \sum_{i=1}^{n} \log(c_i + a_i^T \lambda) + \sum_{i=1}^{m} \lambda_i (a_{i \cdot}(1/y(\lambda)) - b_i) \right).$$

Consider closely the sum

$$\sum_{i=1}^{m} \lambda_{i} a_{i} \cdot (1/y(\lambda)) = \sum_{i=1}^{m} \lambda_{i} \sum_{j=1}^{n} \frac{a_{ij} \tau}{c_{j} + (a_{\cdot j})^{T} \lambda} =$$

$$= \sum_{i=1}^{n} \sum_{j=1}^{m} \frac{\lambda_{i} a_{ij} \tau}{c_{j} + (a_{\cdot j})^{T} \lambda} = \sum_{j=1}^{n} \frac{(a_{\cdot j})^{T} \lambda \tau}{c_{j} + (a_{\cdot j})^{T} \lambda}.$$

It cancels with the first sum and results in $m\tau$. Finally, the dual problem is

$$\max_{\lambda} \left(\tau \sum_{j=1}^{n} \log(c_j + (a_{\cdot j})^T \lambda) - b^T \lambda \right).$$