

EXAM IN CATEGORY THEORY

Let \mathcal{C} be a category. A *retract* of an object $X \in \mathcal{C}$ is a commutative diagram in \mathcal{C} :

$$\begin{array}{ccc} & X & \\ i \nearrow & & \searrow r \\ Y & \xrightarrow{\text{Id}_Y} & Y \end{array}$$

By abusing the terminology, we also say that Y is a retract of X .

Note that if we let $e = i \circ r$ then the morphism $e : X \rightarrow X$ is idempotent, i.e. $e^2 = e$. We say that \mathcal{C} is *idempotent complete* if every idempotent morphism in \mathcal{C} comes from a retract.

Problem 1. Show that the following conditions are equivalent for an idempotent morphism $e : X \rightarrow X$:

- (1) e comes from a retract of X .
- (2) The diagram $X \begin{smallmatrix} \xrightarrow{e} \\ \text{Id}_X \end{smallmatrix} X$ admits an equalizer.
- (3) The diagram $X \begin{smallmatrix} \xrightarrow{e} \\ \text{Id}_X \end{smallmatrix} X$ admits a coequalizer.

Problem 2. Show that if an idempotent morphism $e : X \rightarrow X$ comes from a retract Y of X then Y is unique up to unique isomorphism.

Problem 3. Show that for any idempotent complete category \mathcal{D} , $\text{Fun}(\mathcal{C}, \mathcal{D})$ is also idempotent complete.

We construct a category $\hat{\mathcal{C}}$ as follows. An object of $\hat{\mathcal{C}}$ is a pair (X, e) where $X \in \mathcal{C}$ and $e : X \rightarrow X$ is an idempotent morphism in \mathcal{C} . A morphism $(X, e) \rightarrow (Y, d)$ in $\hat{\mathcal{C}}$ is a morphism $f : X \rightarrow Y$ in \mathcal{C} such that $d \circ f = f = f \circ e$. The composition law of $\hat{\mathcal{C}}$ is induced from that of \mathcal{C} .

Problem 4. Show that the category $\hat{\mathcal{C}}$ is well-defined.

Problem 5. Show that $\hat{\mathcal{C}}$ is idempotent complete.

Problem 6. Show that the assignment $X \mapsto (X, \text{Id}_X)$, $f \mapsto f$ defines a fully faithful functor $I : \mathcal{C} \rightarrow \hat{\mathcal{C}}$.

Let \mathcal{E} be an idempotent complete category. We say that a functor $I : \mathcal{C} \rightarrow \mathcal{E}$ exhibits \mathcal{E} as *idempotent completion* of \mathcal{C} if for any idempotent complete category \mathcal{D} , composition with I induces an equivalence $\text{Fun}(\mathcal{E}, \mathcal{D}) \simeq \text{Fun}(\mathcal{C}, \mathcal{D})$.

Problem 7. Show that the embedding $I : \mathcal{C} \rightarrow \hat{\mathcal{C}}$ exhibits $\hat{\mathcal{C}}$ as idempotent completion of \mathcal{C} .

Problem 8. Use Problem 7 to show that \mathcal{C} is idempotent complete if and only if $\mathcal{C} \simeq \hat{\mathcal{C}}$.



Let \mathcal{S} be a locally small category which admits all small colimits. We say that an object $X \in \mathcal{S}$ is *completely compact* if the representable functor $\text{Hom}_{\mathcal{S}}(X, -) : \mathcal{S} \rightarrow \text{Set}$ preserves all small colimits.

Problem 9. Show that if $X \in \mathcal{S}$ is completely compact then any retract of X is completely compact.

In what follows, we assume \mathcal{C} is small and let $j : \mathcal{C} \rightarrow \mathcal{P}(\mathcal{C})$ be the Yoneda embedding.

Problem 10. Show that $\mathcal{P}(\mathcal{C}) \simeq \mathcal{P}(\hat{\mathcal{C}})$. [Hint: Use Problem 7.]

Problem 11. Show that $j(X) \in \mathcal{P}(\mathcal{C})$ is completely compact for any $X \in \mathcal{C}$.

Problem 12. Show that if $S \in \mathcal{P}(\mathcal{C})$ is completely compact, then there exists $X \in \mathcal{C}$ such that S is a retract of $j(X)$. [Hint: Use the fact that every object of $\mathcal{P}(\mathcal{C})$ is a small colimit of objects in $j(\mathcal{C})$.]

Problem 13. Let $\tilde{\mathcal{C}} \subset \mathcal{P}(\mathcal{C})$ be the full subcategory consisting of the completely compact objects. Show that the embedding $j : \mathcal{C} \rightarrow \tilde{\mathcal{C}}$ exhibits $\tilde{\mathcal{C}}$ as idempotent completion of \mathcal{C} . [Hint: Combine the above four problems.]

