MODERN OPERATING SYSTEMS

LECTURE 6

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AGENDA

- ✓ FILE SYSTEMS
- ✓ DEFINITION AND CONCEPTS
- ✓ PECULIARITIES OF VARIOUS FILE SYSTEMS
- ✓ FAT VERSIONS
- ✓ NTFS
- ✓ NTFS OPERATING PECULIARITIES

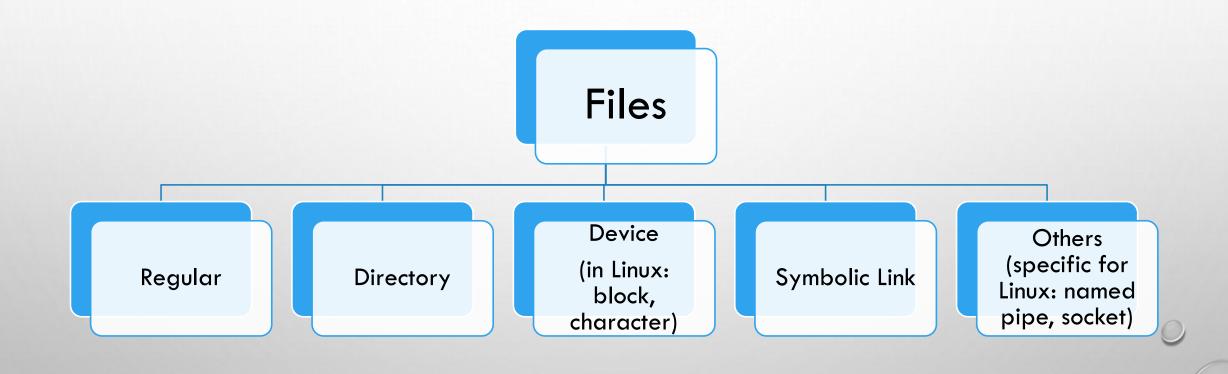


DEFINITIONS & MAIN CONCEPTS

- FILE IS A LINEAR ARRAY OF BYTES, STORED PERSISTENTLY AND HAVING A NAME
- > SEVERAL USERS CAN SHARE INFORMATION STORED IN A FILE
- IN EVERY FILE SYSTEM THERE IS A UNIQUE STRUCTURE (RECORD IN MFT (NTFS), INDEXED NODE, INODE (EXT2FS, EXT3FS)) WHICH IDENTIFIES A FILE
- > THERE ARE FILES OF VARIOUS TYPES IN A FILE SYSTEM



TYPES OF FILES



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FILE SYSTEM

FILE SYSTEM IS A SYSTEM WHICH MAIN GOAL IS TO SUPPORT FILE ORGANIZATION, AND WHICH IS COMPRISED OF 3 MAIN COMPONENTS:

- > A SET OF VARIOUS TYPE FILES
- > SPECIAL DATA STRUCTURES (OFTEN IN THE FORM OF TABLES WHICH HELP TO FIND THE NECESSARY DATA IN THE DISK, I.E. SUPPORT THE PROCESS OF CONVERTING SYMBOLIC FILE NAMES INTO PHYSICAL ADDRESSES)
- > PROGRAMS WHICH REALIZES ALL THE OPERATIONS WITH FILES (E.G. READING, WRITING, DELETING, AND ETC.)



FILE BLOCK

- PFILE SYSTEM BLOCK (CLUSTER IN NTFS & FAT) IS A GROUP OF CONSECUTIVE SECTORS (THE NUMBER OF SECTORS IS ANY DEGREE OF 2 IT COULD BE 1, 2, 4, 8, ... SECTORS) THAT OPTIMIZE STORAGE ADDRESSING.
- > IT IS THE MINIMAL SIZE OF A FILE IN A FILE SYSTEM (MINIMAL ADDRESSING VOLUME)
- MODERN FILE SYSTEMS GENERALLY USE BLOCK SIZES FROM 1 TO 128 SECTORS (512-64K BYTES).



- > EFFICIENCY (I.E. THE RATE OF ACCESS TO FILES)
- ELEVEL OF FRAGMENTATION (WHETHER FILE DATA ARE TO BE PLACED IN THE NEIGHBORING BLOCKS OR SCATTERED ON THE DISK SURFACE)
- > MAXIMUM VOLUME SUPPORTED BY THIS FS (MAXIMAL FILE SIZE)
- > CHANCE TO FILE SYSTEM RECOVERY



FILE SYSTEMS (FS) IN WINDOWS

- FAT (FAT12, FAT16, FAT32) FILE ALLOCATION TABLE THE MAIN FS STRUCTURE
- > NTFS NEW TECHNOLOGY FS
- > REFS (RESILIENT FILE SYSTEM)
 - ✓ IS THE LATEST PRODUCT OF MICROSOFT INTRODUCED WITH WINDOWS 8 AND NOW AVAILABLE FOR WINDOWS 10.
 - ✓ THE FILE SYSTEM ARCHITECTURE ABSOLUTELY DIFFERS FROM OTHER WINDOWS FILE SYSTEMS AND IS MAINLY ORGANIZED IN A FORM OF THE B+-TREE.
 - ✓ REFS HAS HIGH TOLERANCE TO FAILURES DUE TO NEW FEATURES INCLUDED INTO THE SYSTEM



- FAT (FILE ALLOCATION TABLE) IS ONE OF THE SIMPLEST FILE SYSTEM TYPES, WHICH HAS BEEN AROUND SINCE THE 1980S
- > ITS VERSIONS: FAT12, FAT16, FAT32.
- > 12, 16, 32 STAND FOR THE NUMBER OF BITS USED TO ENUMERATE FILE SYSTEM BLOCKS
- FAT CONSISTS OF:
 - > THE FILE SYSTEM DESCRIPTOR SECTOR (BOOT SECTOR OR SUPERBLOCK),
 - > THE FILE SYSTEM BLOCK ALLOCATION TABLE (REFERRED AS THE FILE ALLOCATION TABLE)
 - > PLAIN STORAGE SPACE FOR STORING FILES AND FOLDERS.
- FILES IN FAT ARE STORED IN DIRECTORIES. EACH DIRECTORY IS AN ARRAY OF 32-BYTE RECORDS, EACH DEFINING A FILE OR EXTENDED ATTRIBUTES OF A FILE (E.G. A LONG FILE NAME). A FILE RECORD ATTRIBUTES THE FIRST BLOCK OF A FILE.
- > ANY NEXT BLOCK CAN BE FOUND THROUGH THE BLOCK ALLOCATION TABLE BY USING IT AS A LINKED LIST.

EXAMPLE OF A FILE ALLOCATION TABLE

2	3	4	5	6	7	8	9	10	11	12
0	O	5	6	Е	9	0		Е	0	0
				O			О	O		
				F			F	F		



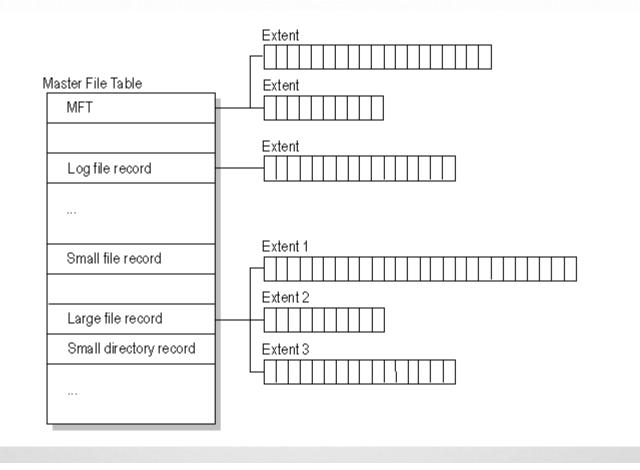
- THE MAIN SYSTEM STRUCTURE OF NTFS IS "MFT MASTER FILE TABLE". IN NTFS, DISK IS DIVIDED INTO MFT SPACE AND SPACE FOR FILE STORAGE. THE MFT PART OCCUPIES AROUND 12% OF THE DISK TO STORE THE MFT METAFILE AND REST 88% SPACE IS USED FOR DATA STORAGE.
- EACH FILE ON AN NTFS VOLUME IS REPRESENTED BY A RECORD (OR SEVERAL RECORDS) IN THE MASTER FILE TABLE (MFT).
- > NTFS RESERVES THE FIRST 16 RECORDS OF THE TABLE FOR SPECIAL INFORMATION (META FILES)
- THE FIRST RECORD OF THIS TABLE DESCRIBES THE MASTER FILE TABLE ITSELF, FOLLOWED BY A MFT MIRROR RECORD.
- > THE THIRD RECORD OF THE MFT IS THE LOG FILE, USED FOR FILE RECOVERY



META FILES

- THE META FILE NAMES ARE PREFIXED BY A '\$'SIGN FOR EXAMPLE: \$MFT, \$BOOT (BOOT SECTOR), \$LOGFILE (JOURNALING SUPPORT FILE), \$.(ROOT DIRECTORY) ETC.
- > \$VOLUME CONTAINS INFORMATION ABOUT THE VOLUME, NAMELY THE VOLUME OBJECT IDENTIFIER, VOLUME LABEL, FILE SYSTEM VERSION, AND VOLUME FLAGS
- > \$ATTRDEF A TABLE OF MFT ATTRIBUTES THAT ASSOCIATES NUMERIC IDENTIFIERS WITH NAMES
- > . ROOT DIRECTORY
- > \$BITMAP AN ARRAY OF BIT ENTRIES: EACH BIT INDICATES WHETHER ITS CORRESPONDING CLUSTER IS USED (ALLOCATED) OR FREE (AVAILABLE FOR ALLOCATION)
- > \$BOOT VOLUME BOOT RECORD. THIS FILE IS ALWAYS LOCATED AT THE FIRST CLUSTERS ON THE VOLUME. IT CONTAINS BOOTSTRAP CODE
- > \$BADCLUS A FILE THAT CONTAINS ALL THE CLUSTERS MARKED AS HAVING BAD SECTORS. THIS FILE SIMPLIFIES CLUSTER MANAGEMENT BY THE CHKDSK UTILITY





MFT Record for a Small File or Directory:

Standard information Recurity descriptor Security Data or index

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LINUX FILE SYSTEMS OF LINUX

- > EXT2, EXT3, EXT4 A "NATIVE" LINUX FILE SYSTEM. THIS FILE SYSTEM FALLS UNDER ACTIVE DEVELOPMENTS AND IMPROVEMENTS.
 - > EXT3 FILE SYSTEM IS JUST AN EXTENSION OF EXT2 THAT USES TRANSACTIONAL FILE WRITING OPERATIONS WITH A JOURNAL.
 - EXT4 IS A FURTHER DEVELOPMENT OF EXT3, EXTENDED WITH THE SUPPORT OF OPTIMIZED FILE ALLOCATION INFORMATION (EXTENTS) AND EXTENDED FILE ATTRIBUTES. THIS FILE SYSTEM IS FREQUENTLY USED AS A "ROOT" FILE SYSTEM FOR MOST LINUX INSTALLATIONS.
- > REISERFS AN ALTERNATIVE LINUX FILE SYSTEM OPTIMIZED FOR STORING OF SMALL FILES. IT HAS GOOD CAPABILITY OF FILES SEARCH AND ENABLES COMPACT FILES ALLOCATION. HOWEVER, THIS FILE SYSTEM NO LONGER RECEIVES ACTIVE SUPPORT.
- > XFS A FILE SYSTEM DERIVED FROM SGI COMPANY AND WAS INITIALLY USED FOR COMPANY'S IRIX SERVERS. NOW XFS

 SPECIFICATIONS ARE IMPLEMENTED IN LINUX. XFS FILE SYSTEM HAS GREAT PERFORMANCE AND IS WIDELY USED TO STORE FILES.

 OPTIMIZED FOR STORING BIG FILES
- > JFS A FILE SYSTEM DEVELOPED BY IBM FOR THE COMPANY'S POWERFUL COMPUTING SYSTEMS. JFS1 USUALLY STANDS
 FOR JFS, JFS2 IS THE SECOND RELEASE. CURRENTLY, THIS FILE SYSTEM IS OPEN-SOURCE AND IMPLEMENTED IN MOST MODERN LINUX VERSIONS.
- ► BTRFS A FILE SYSTEM DESIGNED BY ORACLE SUPPORTED BY THE MAINLINE LINUX KERNEL SINCE 2009. THE FILE SYSTEM IS AIMED AT BETTER RELIABILITY AND SCALABILITY, OFFERING HIGHER FAULT TOLERANCE, EASIER ADMINISTRATION, ETC. TOGETHER WITH A NUMBER OF ADVANCED FEATURES, BUT STILL CANNOT BE CONSIDERED FULLY STABLE.



FILE SYSTEMS OF MAC OS

- > HFS+, AN EXTENSION TO THEIR OWN HFS FILE SYSTEM
 - ✓ USES B-TREES FOR PLACING AND LOCATING FILES.
 - ✓ THE INFORMATION CONCERNING FREE AND USED FILE BLOCKS IS KEPT IN THE ALLOCATION FILE
 - ✓ ALL BLOCKS ASSIGNED TO EACH FILE AS EXTENDS ARE RECORDED IN THE EXTENDS OVERFLOW FILE
 - ✓ ALL FILE ATTRIBUTES ARE LISTED IN THE ATTRIBUTES FILE
 - ✓ DATA RELIABILITY IS IMPROVED THROUGH JOURNALING

> RECENTLY RELEASED APFS

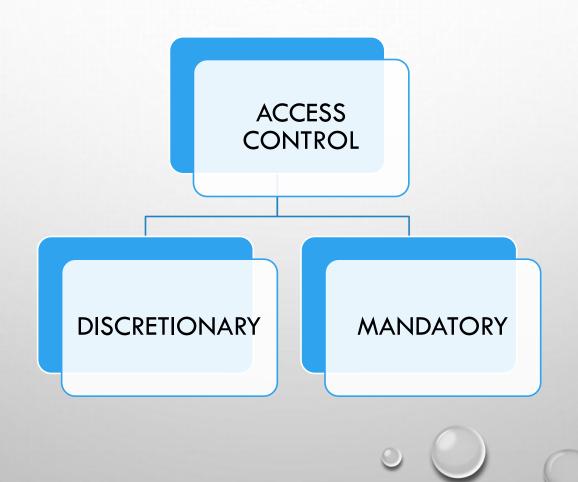
- ✓ IS AIMED TO ADDRESS FUNDAMENTAL ISSUES PRESENT IN ITS PREDECESSOR AND WAS DEVELOPED TO EFFICIENTLY WORK WITH MODERN FLASH STORAGES AND SOLID-STATE DRIVES
- ✓ ALL THE FILE CONTENTS AND METADATA ABOUT FILES, FOLDERS ALONG WITH OTHER APFS STRUCTURES ARE KEPT IN THE APFS CONTAINER.
- ✓ THE CONTAINER SUPERBLOCK STORES INFORMATION ABOUT THE NUMBER OF BLOCKS IN THE CONTAINER, THE BLOCK SIZE, ETC.
- ✓ INFORMATION ABOUT ALL ALLOCATED AND FREE BLOCKS OF THE CONTAINER IS MANAGED WITH THE HELP OF BITMAP STRUCTURES.



WORKING IN NTFS



ACCESS CONTROL TYPES



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DISCRETIONARY ACCESS CONTROL (DAC)

- IS A TYPE OF ACCESS CONTROL DEFINED AS A MEANS OF RESTRICTING ACCESS TO OBJECTS
 BASED ON IDENTITY OF SUBJECTS AND/OR GROUPS TO WHICH THEY BELONGS;
- IS "DISCRETIONARY IN THE SENSE THAT A SUBJECT WITH A CERTAIN ACCESS PERMISSION IS CAPABLE OF PASSING THAT PERMISSION ON TO ANY OTHER SUBJECT

(OBJECTS - FILES AND FOLDERS, SUBJECTS - USERS OR THEIR APPLICATIONS)

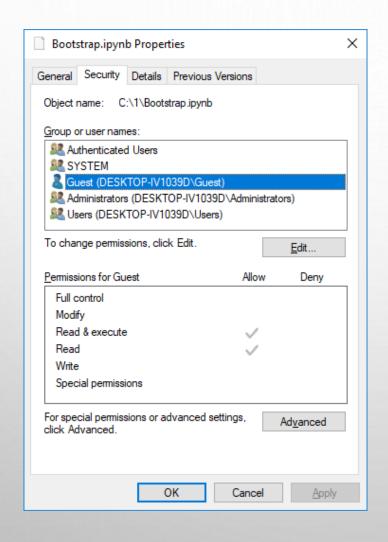


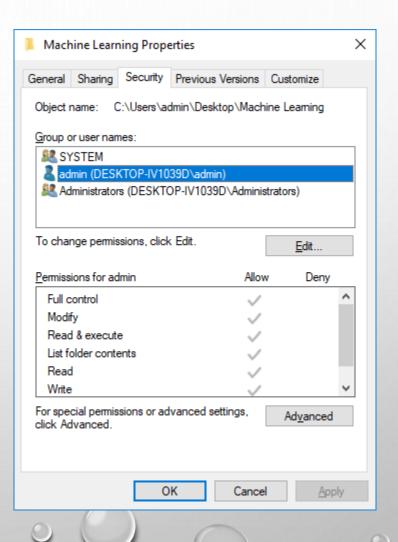
- EVERY FILE MUST HAVE AN OWNER (THERE IS NO "NO ONE'S" FILES OR FOLDERS). THE USER WHO CREATES THE FILE (FOLDER) HAS BECOME ITS OWNER
- THE OWNER SET PERMISSIONS FOR WORING WITH THE FILE (FOLDER) FOR ALL OTHER USERS, AND THESE PERMISSIONS COULD BE INTERPRETED IN THE ONLY ONE WAY
- THERE IS A SUPER USER IN THE SYSTEM WHO IS ABLE TO BECOME THE OWNER OF ANY FILE (FOLDER) IN THE SYSTEM

MAIN ISSUES OF ACCESS CONTROL IN NTFS

- DAC IS REALIZED FOR THIS FILE SYSTEM, THUS ROLE OF THE "OWNER" IS IMMENSELY IMPORTANT IN NTFS
- THE OWNER CAN "ALLOW" OR "DENY" SOME OPERATIONS FOR VARIOUS USERS AND GROUPS
- > YOU MUST TAKE IN ACCOUNT INTERACTIONS BETWEEN THESE "ALLOW" AND "DENY", AS IF THEY CAN BE SET FOR THE USERS INDIVIDUALLY, AND FOR DIFFERENT GROUPS WHERE THESE USERS TOOK PART AS WELL
- > IT IS NECESSARY TO PAY ATTENTION TO PERMISSION INHERITANCE

"SECURITY" TAB IN FILE/FOLDER "PROPERTIES" WINDOW





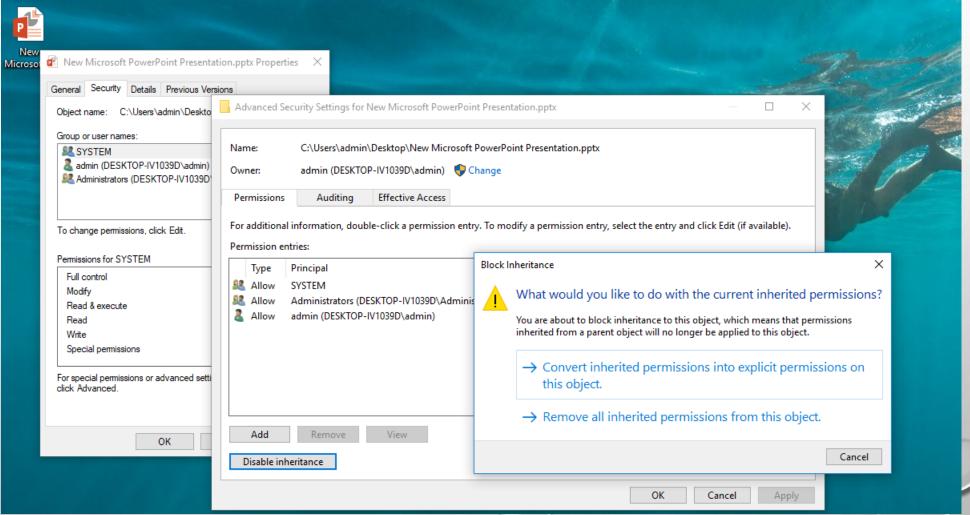


RULES

- > ALL PERMISSIONS (INDICATED AS "ALLOW"") FOLLOW THE RULE OF ADDITION (THEY ARE SUMMARIZED)
- EXAMPLE: IF THE USER NAMED "USER1" HAS ALLOWANCE TO READ PERSONALLY, AND READING IS DENIED FOR THEM AS A MEMBER OF THE GROUP "USERS", "USER1" WILL NOT BE ABLE TO READ THIS FILE)



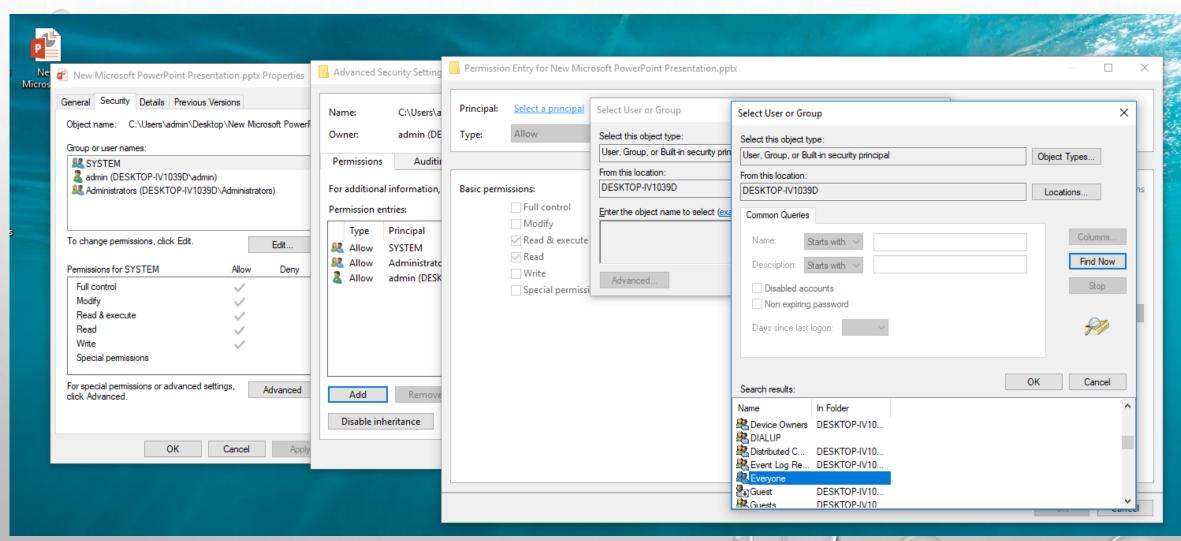
INHERITANCE

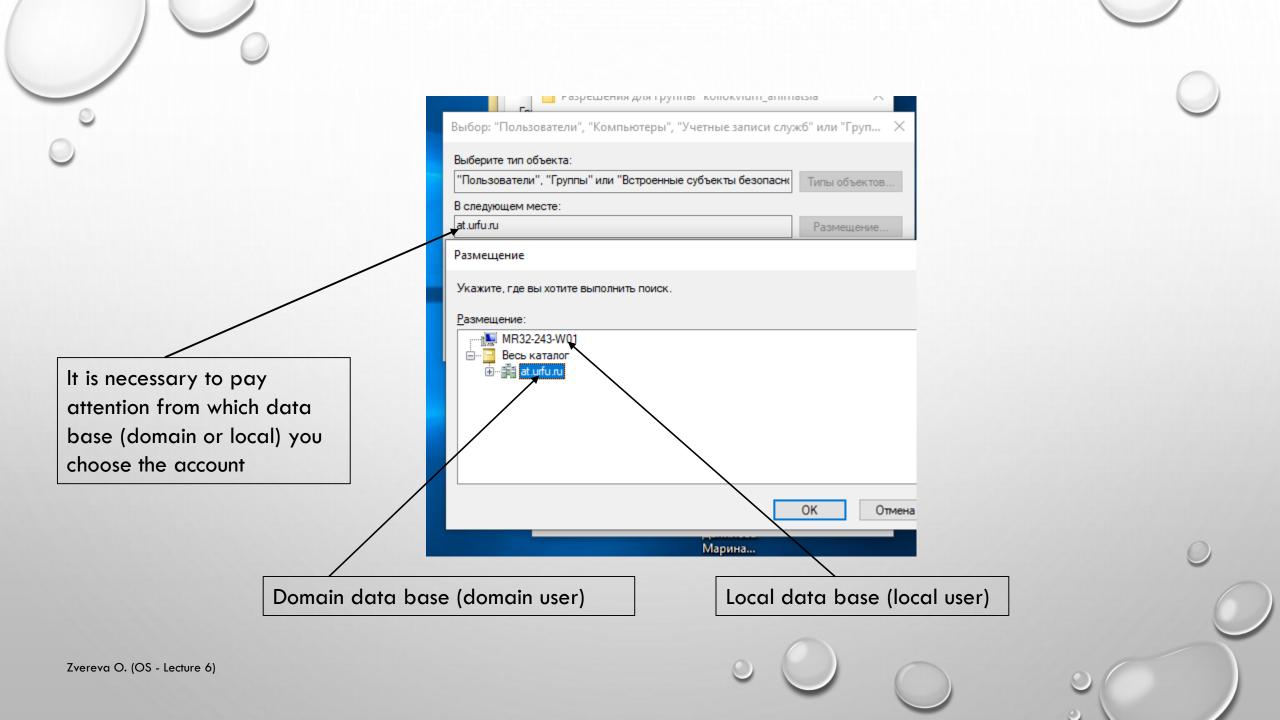


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SECURITY SETTINGS





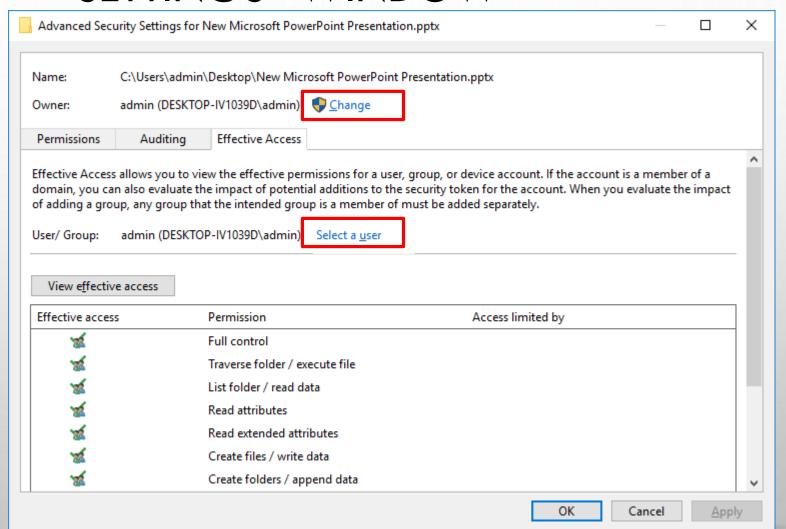


DIFFERENT TYPES OF PERMISSIONS



			_
	Permission Entry for New Microsoft PowerPoint Presentation.pp	otx	– 🗆 X
Permission Entry for New Microsoft PowerPoint Presentation.pptx Principal: BATCH Select a principal Type: Allow	Principal: BATCH Select a principal Type: Allow	•	
Basic permissions: Full control Modify Read & execute Read Write Special permissions	Advanced permissions: Full control Traverse folder / execute file List folder / read data Read attributes Read extended attributes Create files / write data Create folders / append data	 Write attributes Delete Read permissions Change permissions Take ownership 	Show basic permissions Clear all

"EFFECTIVE ACCESS" TAB IN "ADVANCED SECURITY SETTINGS" WINDOW



- ➤ GRAPHICAL
 INTERFACE
 (PREVIOUS SLIDE)
- "TAKEOWN"
 COMMAND
 (FROM THE
 COMMAND LINE)

TO CHANGE THE OWNER

Administrator: Co	ommand Prompt		_		×
C:\Windows\system32>takeown /?					
TAKEOWN [/S system [/U username [/P [password]]]] /F filename [/A] [/R [/D prompt]]					
	allows an adminis by re-assigning	trator to recover access to a file that file ownership.			
Parameter List: /S	system	Specifies the remote system to connect to.			
/υ	[domain\]user	Specifies the user context under which the command should execute.			
/P	[password]	Specifies the password for the given user context. Prompts for input if omitted.			
/F	filename	Specifies the filename or directory name pattern. Wildcard "*" can be used to specify the pattern. Allows sharename\filename.			
/A		Gives ownership to the administrators group instead of the current user.			
/R		Recurse: instructs tool to operate on files in specified directory and all			~

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TAKEOWN /F C:\WINDOWS\SYSTEM32\ACME.EXE

TAKEOWN /F %WINDIR%*.TXT

(IF YOU DO NOT USE KEY "/A" THEN THE CURRENT USER BECOMES FILE/FOLDER OWNER")



- A FILE ACCESS AUDIT IS

 MONITORING AND

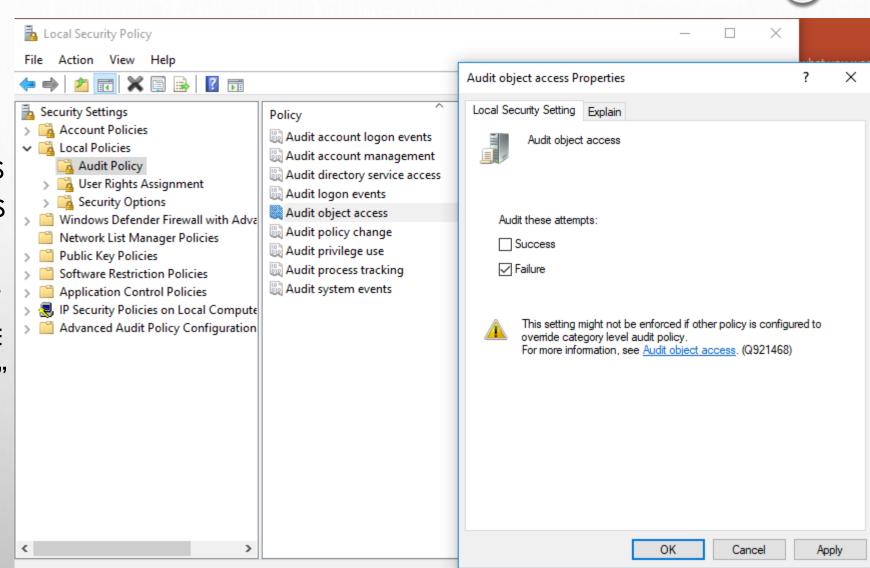
 RECORDING THE ACTIONS

 OF VARIOUS TYPE ACCESS

 TO FILES/FOLDERS
- TO SET THIS TYPE AUDIT IT
 IS NECESSARY TO USE THE
 "LOCAL SECURITY POLICY"
 UTILITY (1ST STAGE)

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AUDIT

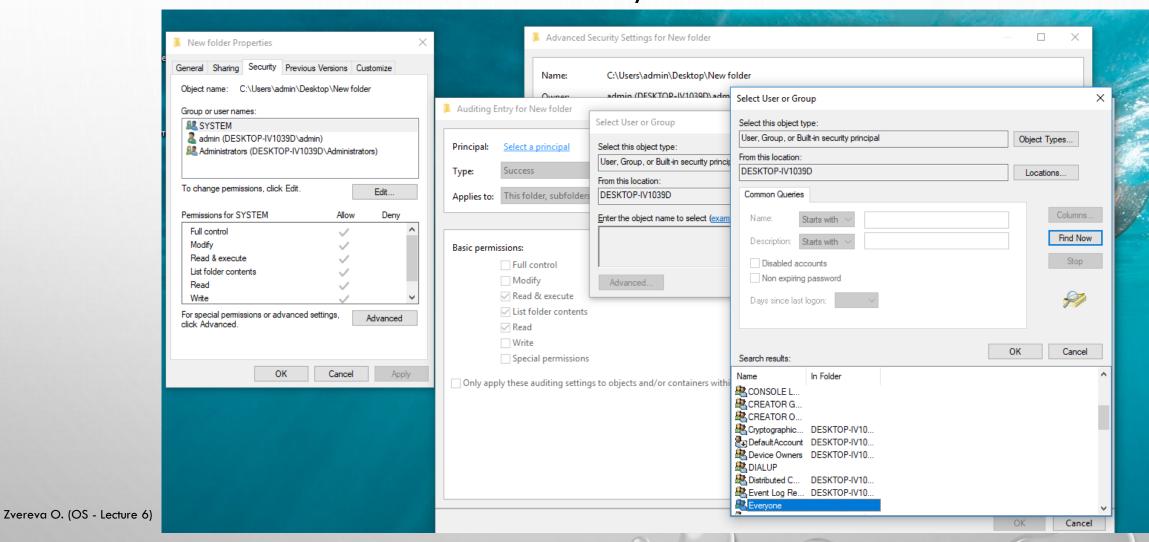




TYPES OF AUDIT

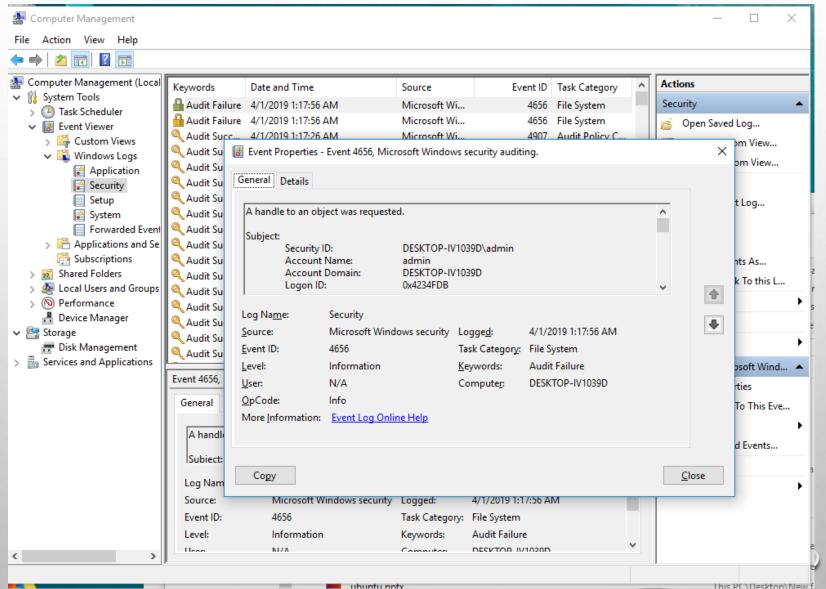
- > "SUCCESS": IF AN ACTION HAD BEEN ALLOWED TO THE USER (OR USERS) AND THIS USER WAS A SUCCESS IN REALIZING THIS ACTION
- FAILURE": IF AN ACTION HAD BEEN DENIED, BUT THE USER TRIED TO REALIZE THIS ACTION, THEY WAS NOT PERMITTED AND FAILED
- > BOTH TYPE ATTEMPTS COULD BE FIXED IN THE SPECIAL SYSTEM LOG (JOURNAL)

2ND STAGE OF AUDIT: IDENTIFYING THE OBJECT, ACTION AND USER/USERS





3RD STAGE: AUDIT RESULTS





TO STOP FILE ACCESS AUDITING

- FOR A FILE/FOLDER (PROPERTIES | SECURITY | ADVANCED SECURITY SETTINGS | AUDITING -> DELETE UNNECESSARY AUDITING)
- FOR THE SYSTEM IN A WHOLE (IMPLEMENTING THE "LOCAL SECURITY POLICY" UTILITY)