## Reinforcement Learning China Summer School



# Game Theory Basics



Haifeng Zhang Institute of Automation, Chinese Academy of Sciences Aug 16, 2022

## **Outline**

- Motivation and Normal-form Game
- Extensive-form Game and Imperfect Information
- Bayesian Game and Incomplete Information
- Nash Equilibrium and Variants
- Theoretical Results of Nash Equilibrium
- Repeated Game and Learning Methods
- Evolutionary Game Theory and Coalitional Game Theory

Outline

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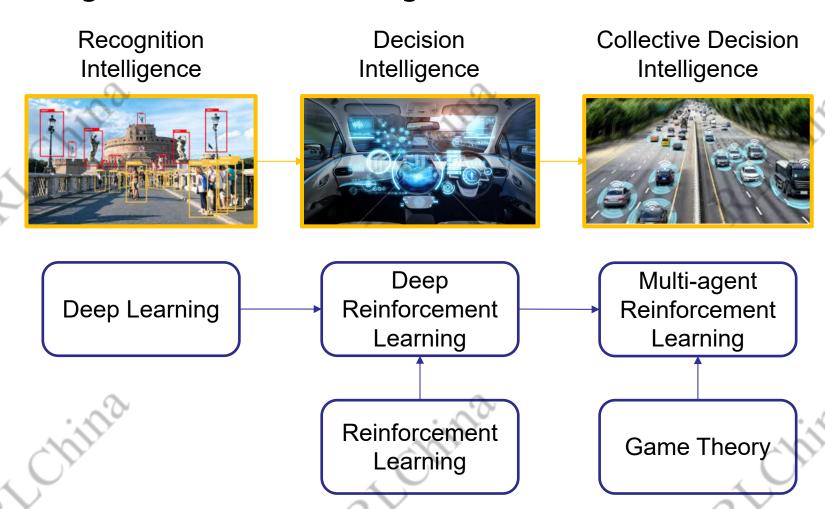
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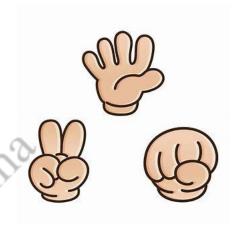
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## Collective Decision Intelligence

Progress of Artificial Intelligence



# Games in Reality



Rock, Scissors, Paper



Chess



Auction



Poker

## History of Game Theory

1934, Stackelberg,Stackelberg Equilibrium[1]1950, Nash,Mixed Nash Equilibrium[2]

**1967**, **Harsanyi**, Bayesian Nash Equilibrium in Bayesian game[5] **1994**, Papadimitriou, PPAD[8]

1951, Brown,Fictitious Play in Repeated game[3]1965, Selten,Subgame Perfect Equilibrium in Extensive-form Game[4]

1973, Smith & Price,Evolutional Game Theory[6]1974, Aumann,Correlated Equilibrium[7]

Till now, 18 game theorists received Nobel Prize in Economics!

# Elements of Game

- Players  $N = \{1, 2, ..., n\}$ 
  - $N = \{1,2\}$
- Strategies (actions)  $A = A_1 \times A_2 \times ... \times ...$ 
  - $A_1 = \{R, S, P\}$   $A_2 = \{R, S, P\}$



- Payoff (utility)  $u = (u_1, u_2, ..., u_n), u_i: A \to \mathbb{R}$ 
  - $u_1: A_1 \times A_2 \to \mathbb{R}$
  - $u_2: A_1 \times A_2 \to \mathbb{R}$

# Normal-form Game

## More than 2 players

| _                |         | $p_1$ | $p_2$          | $p_3$ |
|------------------|---------|-------|----------------|-------|
| loint            | R, R, R | 0     | -1             | 1     |
| Joint<br>Actions | R, R, S | 1 .   | <sup>ک</sup> 1 | 0     |
| Chir             |         |       |                |       |

# Rationality of Players

- Self-interested
  - Preference over game outcome
  - E.g. (paper, rock) is better than (rock, paper) for row player
- - Utility of (paper, rock) is 1 Utility of (rock, paper) is -1
- Objective
  - Act to maximize (expected) utility

## Common Knowledge

#### Definition

 All the players know p, they all know that they know p, they all know that they all know that they know p, and so on

#### Example

- ×
  - Alice knows 'the weather is good'
  - Bob Knows 'the weather is good'
- X
  - Alice and Bob knows 'the weather is good' respectively
  - Alice and Bob knows 'the opposite knows the weather is good'
- √
- Alice and Bob are told 'the weather is good' together

## Pure Strategy and Mixed Strategy

#### Pure Strategy

- $a_1 \in A_1 = \{Heads, Tails\}$
- $a_2 \in A_2 = \{Heads, Tails\}$

#### **Matching Pennies**

|       | Heads | Tails |
|-------|-------|-------|
| Heads | 1, -1 | -1, 1 |
| Tails | -1, 1 | 1, -1 |

- Mixed Strategy: Probability Distribution over Pure Strategy
  - $a_1 = (x_H, x_T), x_H \in [0,1], x_T \in [0,1], x_H + x_T = 1$
  - $a_2 = (y_H, y_T), y_H \in [0,1], y_T \in [0,1], y_H + y_T = 1$
- Expected Utility
  - $EU_1 = x_H y_H u_1(H, H) + x_H y_T u_1(H, T) + x_T y_H u_1(T, H) + x_T y_T u_1(T, T)$
  - $EU_2 = x_H y_H u_2(H, H) + x_H y_T u_2(H, T) + x_T y_H u_2(T, H) + x_T y_T u_2(T, T)$
- Example
  - $a_1 = (0.1, 0.9), a_2 = (0.3, 0.7)$
  - $EU_1 = 0.32, EU_2 = -0.32$

## Classic Games

- Zero-sum Game
  - $u_1(a) + u_2(a) = 0, \forall a \in A$
- Cooperative Game
  - $u_i(a) = u_j(a), \forall a \in A, i, j \in N$
- Coordination Game
  - Multiple Nash Equilibria Exist
- Social Dilemma [9]
  - Everyone suffers in an NE

#### **Matching Pennies**

|       | Heads | Tails |
|-------|-------|-------|
| Heads | 1, -1 | -1, 1 |
| Tails | -1, 1 | 1, -1 |

#### **Road Selection**

|       | Left | Right |
|-------|------|-------|
| Left  | 1, 1 | 0, 0  |
| Right | 0, 0 | 1, 1  |

#### Battle of Sex

|       | Party | Home  |
|-------|-------|-------|
| Party | 10, 5 | 0, 0  |
| Home  | 0, 0  | 5, 10 |

#### Prisoner's Dilemma

| V         | Cooperate | Defect |
|-----------|-----------|--------|
| Cooperate | 2, 2      | 0, 3   |
| Defect    | 3, 0      | 1, 1   |

Outline

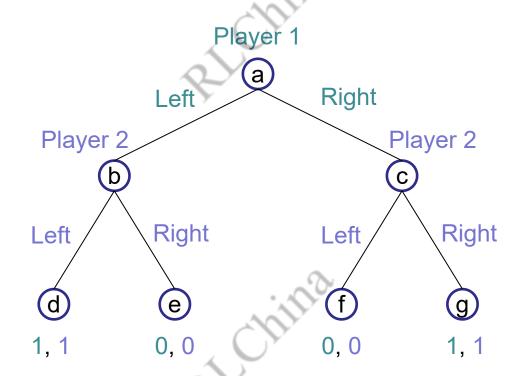


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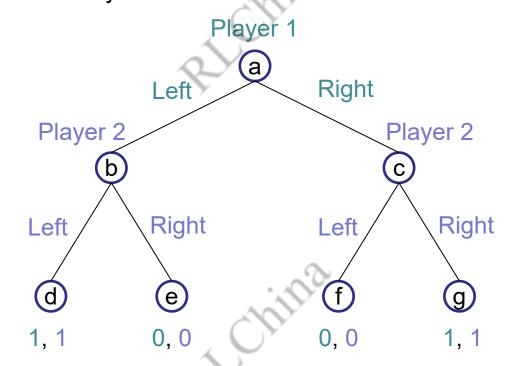
## Extensive-form Game

- Game Tree
  - Node: decision point for a specified player
  - Edge: action decided by the player
  - · Leaf: outcome of the game with payoff



## Strategies in Extensive-form Game

- Strategy Space
  - Player 1: {Left, Right}
  - Player 2: {(Left, Left), (Left, Right), (Right, Left), (Right, Right)}
     action in every node

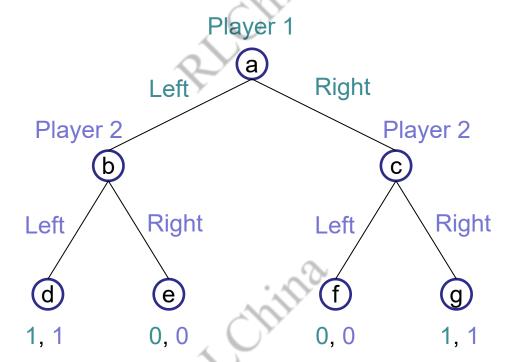


## Extensive-form vs. Normal-form

#### Equivalent Normal-form Game

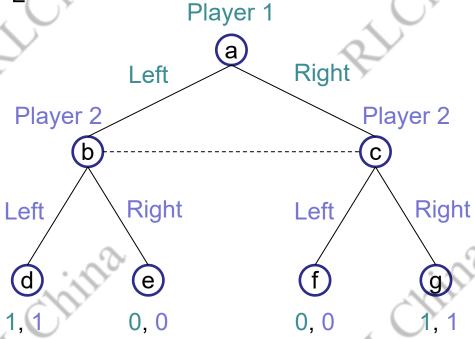
|       | (Left, Left) | (Left, Right) | (Right, Left) | (Right, Right) |
|-------|--------------|---------------|---------------|----------------|
| Left  | 1, 1         | 1, 1          | 0, 0          | 0, 0           |
| Right | 0, 0         | 1, 1          | 0, 0          | 1, 1           |

multiple step/state ↓ **dynamic** 



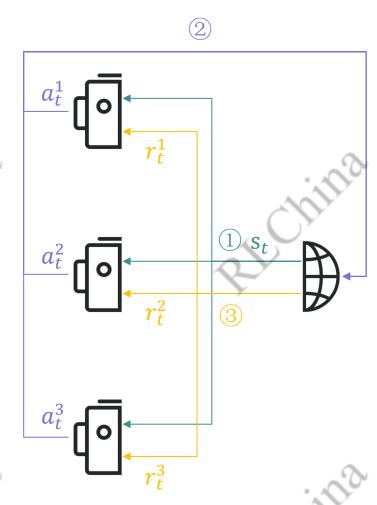
## Imperfect Information

- Imperfect Information Game
  - Some historical actions are invisible by other players
- Information Set
  - A set containing undistinguishable states, e.g. {b, c} is an information set for player 2
- Strategy Space
  - Player 1: {Left, Right}
  - Player 2: {Left, Right}
     action in every information set



## Markov Game (or Stochastic Game)

- Game Definition
  - State space S
  - Action space  $A = A_1 \times A_2 \times ... \times A_n$
  - Transition function  $p: S \times A \rightarrow S$
  - Reward function  $r: S \times A \to \mathbb{R}^n$
- Behavioral Strategy
  - Policy  $\pi_i: S \times A_i \rightarrow [0,1]$
- Properties
  - Simultaneous action (Normal-form)
  - Multiple step/state (Extensive-form)
  - · Immediate reward
  - Randomness
  - Cycle



Interaction at time-step t

# Summary of Strategy Representation

| -100                | Static Game<br>(Single Step/state) | Dynamic Game<br>(Multiple step/state)          |
|---------------------|------------------------------------|--|
| Pure Strategy       | $a_i \in A_i$                      | $\pi_i: S \to A_i \text{ or } \pi_i \in A_i^S$ |
| Mixed Strategy      | $a_i: A_i \to [0,1]$               | $\pi_i: A_i^S \to [0,1]$                       |
| Behavioral Strategy | $a_i: A_i \to [0,1]$               | $\pi_i: S \times A_i \to [0,1]$                |

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## **Example: Auction**

#### Game Definition

- Players has private value  $v_1$ ,  $v_2$
- Players decide biddings  $b_1$ ,  $b_2$
- Player i with higher bidding  $b_i$  has utility  $v_i b_i$
- The other player has utility 0
- Uncertainty of Private Value
  - $v_1 = 4, v_2 = 4$
  - $b_1 \in \{1,3\}, b_2 \in \{2,4\}$

| • | $v_1$ | = | 4, | $v_2$ | = | 5 |
|---|-------|---|----|-------|---|---|
|---|-------|---|----|-------|---|---|

• 
$$b_1 \in \{1,3\}, b_2 \in \{2,4\}$$

|           | $b_2 = 2$ | $b_2 = 4$ |
|-----------|-----------|-----------|
| $b_1 = 1$ | 0, 2      | 0, 0      |
| $b_1 = 3$ | 1, 0      | 0, 0      |

|           | $b_2 = 2$ | $b_2 = 4$ |
|-----------|-----------|-----------|
| $b_1 = 1$ | 0, 3      | 0, 1      |
| $b_1 = 3$ | 1, 0      | 0, 1      |

Players don't know the exact payoff matrix of the game!

## Incomplete Information

- Recall the Elements of a Game
  - Players  $N = \{1, 2, ..., n\}$
  - Action space  $A = A_1 \times A_2 \times ... \times A_n$
  - Payoff functions  $u = (u_1, u_2, ..., u_n), u_i: A \to \mathbb{R}$
- Incomplete Information Game
  - Players know: N and A
  - Players don't completely know: u
  - Criteria: whether players have private information when game starts
- Example
  - Auction
  - Mahjong
  - Werewolves of Miller's Hollow

## **Bayesian Game**

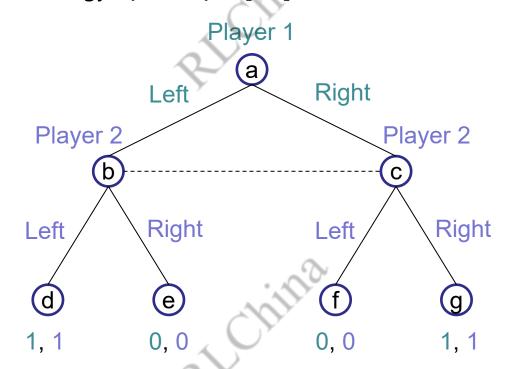
- Basic Idea
  - Payoff function  $p_i$  is unknown, but the distribution of  $p_i$  is known
- Elements of Bayesian Game
  - Players  $N = \{1, 2, ..., n\}$ , action space  $A = A_1 \times A_2 \times ... \times A_n$
  - Player type space  $\Theta = \Theta_1 \times \Theta_2 \times ... \times \Theta_n$
  - Distribution over types  $d: \Theta \rightarrow [0,1]$
  - Payoff functions  $u = (u_1, u_2, ..., u_n), u_i: \Theta \times A \rightarrow \mathbb{R}$
- Strategy
  - Pure strategy  $\pi_i: \Theta_i \to A_i$
  - Mixed strategy  $\pi_i: \Theta_i \times A_i \rightarrow [0,1]$
- Example
  - $\Theta_1 = \{4\}, \Theta_2 = \{4,5\}$
  - d(4,4) = 0.3, d(4,5) = 0.7

|               |             | $b_2 = 2$ | $b_2 = 4$ |
|---------------|-------------|-----------|-----------|
| $v_2 = 4$ 0.3 | $b_1 = 1$   | 0, 3      | 0, 1      |
| 3.8           | $h_{*} = 3$ | 1 0       | 0 1       |

| _         |           | $b_2 = 2$ | $b_2 = 4$ |
|-----------|-----------|-----------|-----------|
| $v_2 = 5$ | $b_1 = 1$ | 0, 2      | 0, 0      |
| 0.7       | $b_1 = 3$ | 1, 0      | 0, 0      |

## Dynamic Bayesian Game

- Belief System in Imperfect Information Extensive-form Game
  - Distribution over the states in an information set  $b_i: S \to [0,1]$
- Strategy
  - Pure strategy  $\pi_i: S \to A_i$
  - Behavioral strategy  $\pi_i: S \times A_i \rightarrow [0,1]$



## **Summary of Game Representation**

|         |           | Complete                                     | Incomplete                     |
|---------|-----------|--|--------------------------------|
| Static  |           | Normal-form Game,<br>e.g. Prisoner's Dilemma | Bayesian Game,<br>e.g. Auction |
| Dynamic | Perfect   | Extensive-form Game,<br>e.g. Chess           | Texas Hold'em Poker            |
| O       | Imperfect | StarCraft                                    | Mahjong                        |

Dynamic Bayesian game

- Harsanyi Transformation
  - Incomplete Information → Imperfect Information
  - Introduce a nature player who decides the type of each player

Outline

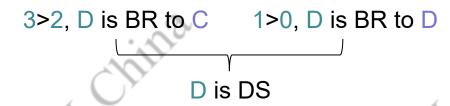
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## Game Solution Reasoning

- Best Response (BR)
  - Given  $a_{-i} \in A_1 \times ... \times A_{i-1} \times A_{i+1} \times ... \times A_n$
  - $a_i$  is best response to  $a_{-i} \Leftrightarrow u_i(a_i, a_{-i}) \ge u_i(a_i', a_{-i})$ ,  $\forall a_i' \in A_i$
- Dominant Strategy (DS)
  - $a_i$  is dominant strategy  $\Leftrightarrow$  Given any  $a_{-i}$ ,  $a_i$  is best response
- Example

#### Prisoner's Dilemma

|               | Cooperate (C) | Defect (D) |
|---------------|---------------|------------|
| Cooperate (C) | 2, 2          | 0, 3       |
| Defect (D)    | 3, 0          | 1, 1       |



# Game Solution Concept: Nash Equilibrium

#### Definition

• A joint strategy (or strategy profile)  $a \in A$  is a Nash Equilibrium  $\Leftrightarrow a_i$  is best response to  $a_{-i}$  holds for every player i

#### Example

#### **Matching Pennies**

|       | Heads | Tails |
|-------|-------|-------|
| Heads | 1, -1 | -1, 1 |
| Tails | -1, 1 | 1, -1 |

#### **Road Selection**

|       | Left | Right |
|-------|------|-------|
| Left  | 1, 1 | 0, 0  |
| Right | 0, 0 | 1, 1  |

#### Battle of Sex

|       | Party | Home  |
|-------|-------|-------|
| Party | 10, 5 | 0, 0  |
| Home  | 0, 0  | 5, 10 |

#### Prisoner's Dilemma

|           | Cooperate | Defect |
|-----------|-----------|--------|
| Cooperate | 2, 2      | 0, 3   |
| Defect    | 3, 0      | 1, 1   |

## Pareto Optimality vs. Nash Equilibrium

- Pareto Optimality (PO)
  - A joint strategy (or strategy profile)  $a \in A$  achieves Pareto optimality  $\Leftrightarrow \nexists a' \in A$  s. t.  $\textcircled{1} \forall i, u_i(a') \geq u_i(a), \textcircled{2} \exists i, u_i(a') > u_i(a)$
  - A Pareto optimality is not necessarily a Nash equilibrium
  - A Nash equilibrium is not necessarily a Pareto optimality

| Chicken |      | ,    | Stay Hunt |      |      | Filsoner's Dilemina |      |      |
|---------|------|------|-----------|------|------|---------------------|------|------|
|         | С    | D    |           | С    | D    |                     | С    | D    |
| С       | 3, 3 | 1, 4 | С         | 3, 3 | 0, 2 | С                   | 2, 2 | 0, 3 |
| D       | 4, 1 | 0, 0 | D         | 2, 0 | 1, 1 | D                   | 3, 0 | 1, 1 |

Stan Hunt

C-D is PO and NE

Chicken

D-D is NE but not PO

C-C is PO but not NE

Drisonar's Dilamma

## Mixed-Strategy Nash Equilibrium

#### Definition

- A mixed-strategy profile  $(a_1, a_2, ..., a_n)$ ,  $a_i \in PD(A_i)$  is a Nash Equilibrium  $\Leftrightarrow a_i$  is best response to  $a_{-i}$  holds for every player i
- Example (Rock-Scissors-Paper)

• 
$$a_1 = (1/3, 1/3, 1/3), a_2 = (1/3, 1/3, 1/3)$$

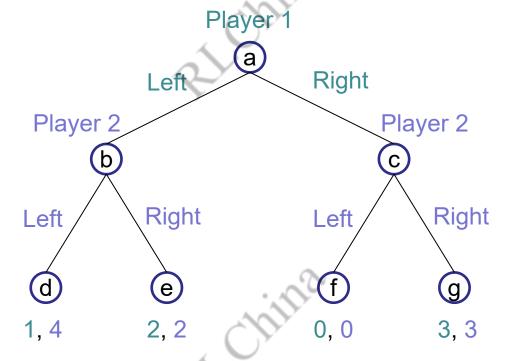
- $EU_1(a_1, a_2) = 0 \ge EU_1(a_1', a_2) = 0$ ,  $\forall a_1' \in A_1$
- $EU_2(a_1, a_2) = 0 \ge EU_2(a_1, a_2') = 0$ ,  $\forall a_2' \in A_2$

|     |     | 1/3   | 1/3   | 1/3   |
|-----|-----|-------|-------|-------|
| 1.  |     | R     | S     | Р     |
| 1/3 | R   | 0, 0  | 1, -1 | -1, 1 |
| 1/3 | ⊘ S | -1, 1 | 0, 0  | 1, -1 |
| 1/3 | Р   | 1, -1 | -1, 1 | 0, 0  |

## Nash Equilibrium in Extensive-form Game

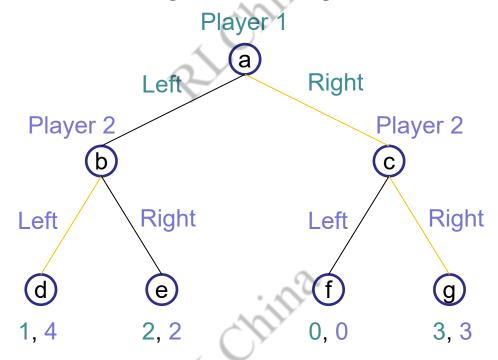
#### Incredible Threat

|       | (Left, Left) | (Left, Right) | (Right, Left) | (Right, Right) |
|-------|--------------|---------------|---------------|----------------|
| Left  | 1, 4 ?       | 1, 4          | 2, 2          | 2, 2           |
| Right | 0, 0         | 3, 3          | 0, 0          | 3, 3 ?         |



# Subgame Perfect Nash Equilibrium (SPNE)

- Definition
  - An NE is SPNE ⇔ the NE holds in every subgame
- Solution
  - Backward induction: Right (Left, Right)



## Bayesian Nash Equilibrium

- Recall Bayesian Game
  - Player type space  $\Theta = \Theta_1 \times \Theta_2 \times ... \times \Theta_n$
  - Distribution over types  $d: \Theta \rightarrow [0,1]$
  - Payoff functions  $u = (u_1, u_2, ..., u_n), u_i: \Theta \times A \rightarrow \mathbb{R}$
- Strategy in Bayesian Game
  - Pure strategy  $\pi_i: \Theta_i \to A_i$
  - Mixed strategy  $\pi_i: \Theta_i \times A_i \rightarrow [0,1]$
- Bayesian Nash Equilibrium (BNE)
  - Assume each player i knows her own type  $\theta_i \in \Theta_i$
  - Set expected utility  $\mathbb{E}[u_i|\pi,\theta] = \sum_{a \in A} (\prod_{j \in N} \pi_j(\theta_j, a_j) u_i(\theta, a))$
  - $\pi$  is BNE  $\Leftrightarrow \pi_i \in \operatorname{argmax}_{\pi_i}$ ,  $\sum_{\theta_{-i} \in \Theta_{-i}} d(\theta_i, \theta_{-i}) \mathbb{E}[u_i | \pi_i', \pi_{-i}, \theta_i, \theta_{-i}]$  holds for each player i with her own type  $\theta_i$

## Bayesian Nash Equilibrium: Example

#### Auction

• 
$$A_1 = \{1,3\}, A_2 = \{2,4\}, \Theta_1 = \{4\}, \Theta_2 = \{4,5\}, d(4,4) = 0.3, d(4,5) = 0.7$$

#### Strategy

• 
$$\pi_1(4,1) = x, \pi_1(4,3) = 1 - x$$

• 
$$\pi_2(4,2) = y_1, \pi_2(4,4) = 1 - y_1$$

• 
$$\pi_2(5,2) = y_2, \pi_2(5,4) = 1 - y_2$$

| 0     |          |   |
|-------|----------|---|
| $v_2$ | <u> </u> | 4 |
|       |          | _ |
| · ·   | 0.3      | 3 |

|           | $b_2 = 2$ | $b_2 = 4$ |
|-----------|-----------|-----------|
| $b_1 = 1$ | 0, 3      | 0, 1      |
| $b_1 = 3$ | 1, 0      | 0, 1      |

#### Equilibrium

• 
$$\mathbb{E}[u_1|\pi_1,\pi_2,4,4] = (1-x)y_1$$

• 
$$\mathbb{E}[u_1|\pi_1,\pi_2,4,5] = (1-x)y_2$$

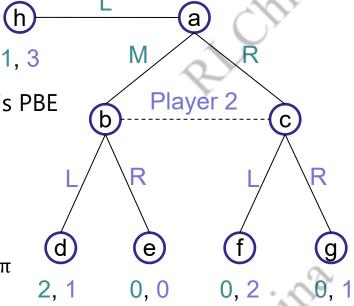
$$v_2 = 5$$
 0.7

|           | $b_2 = 2$ | $b_2 = 4$ |
|-----------|-----------|-----------|
| $b_1 = 1$ | 0, 2      | 0, 0      |
| $b_1 = 3$ | 1, 0      | 0, 0      |

- $\mathbb{E}[u_2|\pi_1,\pi_2,4,4] = 3xy_1 + x(1-y_1) + (1-x)(1-y_1)$
- $\mathbb{E}[u_2|\pi_1,\pi_2,4,5] = 2xy_2$
- $(x, y_1, y_2)$  satisfies  $x = \operatorname{argmax}_{x} 0.3(1 x)y_1 + 0.7(1 x)y_2$  and  $y_1 = \operatorname{argmax}_{y_1} 3xy_1 + x(1 y_1) + (1 x)(1 y_1)$  and  $y_2 = \operatorname{argmax}_{y_2} 2xy_2$

## Perfect Bayesian (Nash) Equilibrium

- Motivation
  - SPNE is not enough for some imperfect information game
  - Example: (L,R) is an SPNE but is incredible
- Recall Dynamic Bayesian Game
  - Belief function  $b_i: S \to [0,1]$
  - Behavioral strategy  $\pi_i: S \times A_i \rightarrow [0,1]$
- Perfect Bayesian Equilibrium (PBE) [10]
  - A strategy profile  $\pi$  with a belief system b is PBE
  - Sequential rationality
    - Each player has best expected utility in each information set following b and  $\pi$
  - Consistency of beliefs with Strategies
    - Beliefs b are correct according to strategies  $\pi$



Player 1

# Summary of Nash Equilibrium

|         |           | Complete                            | Incomplete                           |
|---------|-----------|-------------------------------------|--------------------------------------|
| Static  |           | Nash Equilibrium                    | Bayesian Nash<br>Equilibrium         |
| Dynamic | Perfect   | Subgame Perfect<br>Nash Equilibrium | Perfect Bayesian<br>Nash Equilibrium |
|         | Imperfect |                                     |                                      |

- Harsanyi Transformation
  - Incomplete Information → Imperfect Information
  - Introduce a nature player who decides the type of each player

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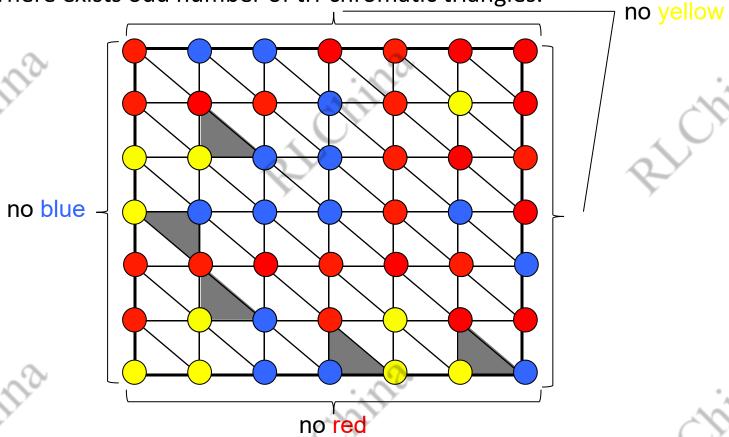
## Existence of Nash Equilibrium

- Nash's Theorem [2]
  - Every finite game has a mixed-strategy Nash equilibrium.
  - Proof: apply Brouwer's fixed point theorem.
- Brouwer's Fixed Point Theorem
  - Let D be a convex, compact subset of the Euclidean space. If  $f: D \to D$  is continuous, then there exists  $x \in D$  such that f(x) = x.
  - Proof: apply Sperner's lemma.
- Sperner's Lemma
  - Color the boundary using three colors in a legal way. No matter how the internal nodes are colored, there exists a tri-chromatic triangle. In fact an odd number of those.

# Sperner's Lemma (2-D)

#### Lemma

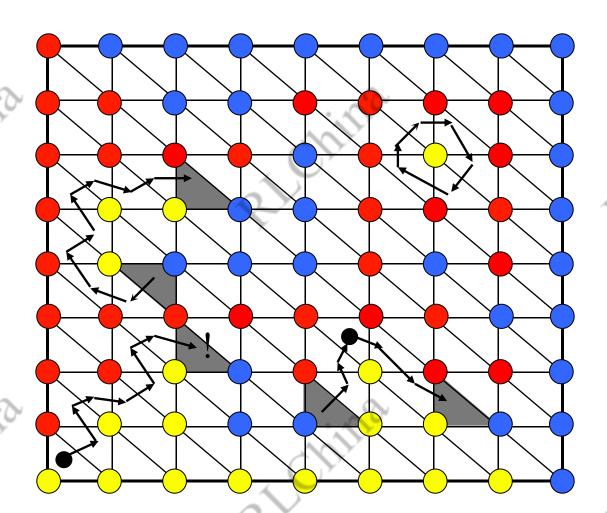
There exists odd number of tri-chromatic triangles.



Ref: http://people.csail.mit.edu/costis/6896sp10/

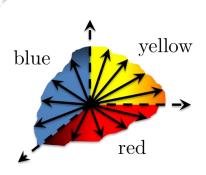
# Proof of Sperner's Lemma

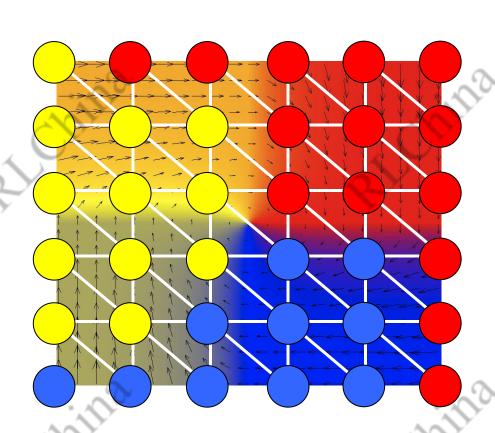
Proof Sketch



# Sperner's Lemma to Brouwer's Theorem

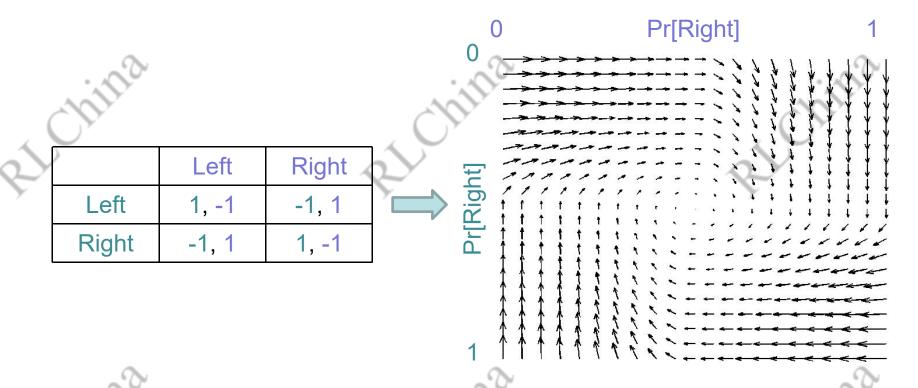
Proof Sketch



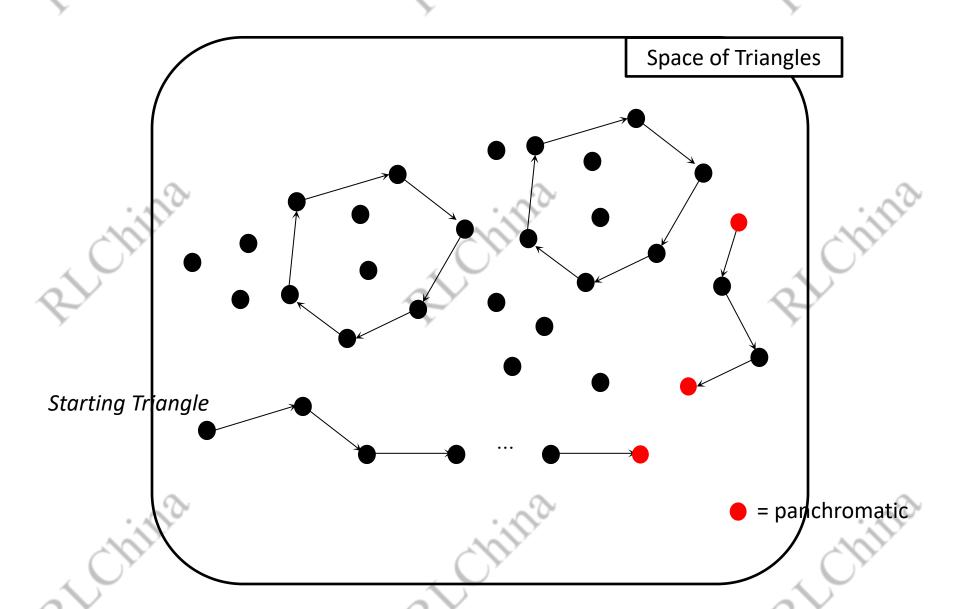


# Brouwer's Theorem to Nash's Theorem

Proof Sketch



# End-of-the-Line to Sperner's Lemma

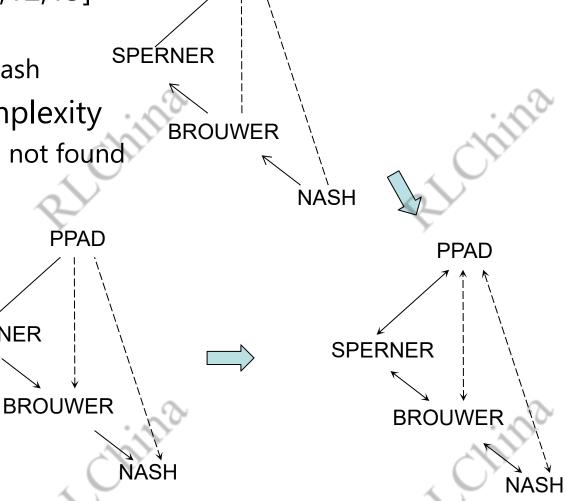


## **PPAD Complexity Class**

- PPAD-complete [11,12,13]
  - End-of-the-line
  - Sperner, Brouwer, Nash
- **Computational Complexity** 
  - Poly-time algorithm not found

**SPERNER** 

**PPAD** 



**PPAD** 

Outline

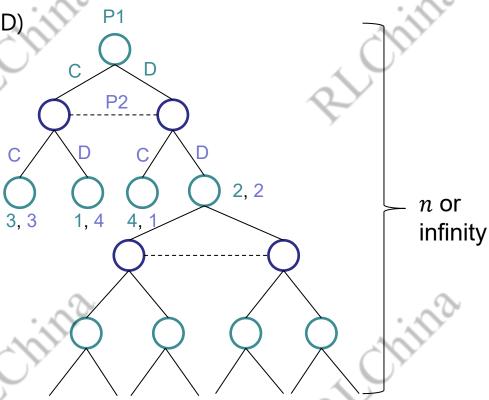
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PICHIL

- Motivation and Normal-form Game
- Extensive-form Game and Imperfect Information
- Bayesian Game and Incomplete Information
- Nash Equilibrium and Variants
- Theoretical Results of Nash Equilibrium
- Repeated Game and Learning Methods
- Evolutionary Game Theory and Coalitional Game Theory

#### Repeated Game

- Definition
  - A normal-form game is played over and again by the same players
  - The game repeated in each period is referred to as the stage game
- Example
  - Iterated Prisoners' Dilemma (IPD)
  - Reward: average or discounted
  - Memory: perfect recall



#### Memory

- Historical Behavior
  - At stage t, the action profile is  $a_t$
  - Each player remembers the action profiles at last k stages
  - We say the players have *k*-memory
- Relation to Markov Game
  - Memory is regarded as state

| 1-memory |   |   |          |  |   |  |  |
|----------|---|---|----------|--|---|--|--|
|          | 1 | 2 | 3        |  | m |  |  |
| P1       | C | C | D        |  | D |  |  |
| P2       | D | D | С        |  | C |  |  |
|          |   |   |          |  |   |  |  |
| Pn       | D | С | D        |  | С |  |  |
| 7        |   | П | <b>☆</b> |  |   |  |  |

as state

1-memory

#### k-memory

|    | 1 | 2 | 3 | <br>m |
|----|---|---|---|-------|
| P1 | С | С | D | <br>D |
| P2 | D | D | С | <br>С |
|    |   |   |   |       |
| Pn | D | C | D | <br>C |



#### Tit-for-tat

- Idea [11]
  - The Tit-for-tat strategy copies what the other player previously choose.
  - Nice: start by cooperating.
  - Clear: be easy to understand and adapt to.
  - Provocable: retaliate against anti-social behavior.
  - Forgiving: cooperate when faced with pro-social play.

|   | С    | D    |
|---|------|------|
| С | 2, 2 | 0, 3 |
| D | 3, 0 | 1, 1 |

Prisoner's Dilemma

|    | 1 | 2 | 3 | 4 |  |
|----|---|---|---|---|--|
| P1 | С | С | С | С |  |
| P2 | С | С | С | С |  |

|    | 1 | 2 | 3 | 4 |  |
|----|---|---|---|---|--|
| P1 | Ó | Ω | O | D |  |
| P2 | ۵ | O | D | O |  |

cooperate by playing strategy (C,C)

payoff = 
$$2 + 2\gamma + 2\gamma^2 + 2\gamma^3 + \dots = 2\frac{\gamma^n - 1}{\gamma - 1} = \frac{2}{1 - \gamma}$$

a player deviates to defecting (D)

payoff = 
$$3 + 0\gamma + 3\gamma^2 + 0\gamma^3 + ... = \frac{3}{1 - \gamma^2}$$

# Win-stay, lose-shift

|   | С    | D    |
|---|------|------|
| С | 2, 2 | 0, 3 |
| D | 3, 0 | 1, 1 |

- Idea [12]
  - Repeat if it was rewarded by 2 or 3

Prisoner's Dilemma

- Shift if it was punished by 0 or 1
- Advantage: tolerant, one round of mutual defection followed by a return to cooperation
- Disadvantage: fares poorly against inveterate defectors

| >  | 1 | 2 | 3 | 4 |  |
|----|---|---|---|---|--|
| P1 | O | D | O | O |  |
| P2 | D | D | O | C |  |

P2 deviates to defecting(D) initially

Payoff 
$$> \frac{3}{1 - \gamma^2}$$

 1
 2
 3
 4
 ...

 P1
 C
 C
 D
 C
 ...

 P2
 C
 D
 D
 D
 ...

P2 is an inveterate defectors

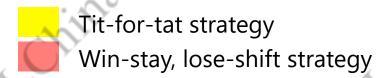
payoff<sub>1</sub> = 2 + 0
$$\gamma$$
 + 1 $\gamma$ <sup>2</sup> + 0 $\gamma$ <sup>3</sup> + ... =  $\frac{1\gamma}{1 - \gamma^2}$  + 1  
payoff<sub>2</sub> = 2 + 3 $\gamma$  + 1 $\gamma$ <sup>2</sup> + 3 $\gamma$ <sup>3</sup> + ... =  $\frac{1}{1 - \gamma^2}$  + 1

# Strategies in Iterated Prisoner's Dilemma

|   | С    | D    |
|---|------|------|
| С | 2, 2 | 0, 3 |
| D | 3, 0 | 1, 1 |

Prisoner's Dilemma

| Action profile at time t | <b>:</b> | l | Play | er 1 | stra | ateg | ies | at tii | me i | t + 1 | l wit | :h 1- | meı | mor | y | ) |
|--------------------------|----------|---|------|------|------|------|-----|--------|------|-------|-------|-------|-----|-----|---|---|
| $(a_1, a_2) = (C, C)$    | ) C      | С | С    | С    | С    | С    | С   | С      | D    | D     | D     | D     | D   | D   | D | D |
| $(a_1, a_2) = (C, D)$    | ) C      | С | С    | С    | D    | D    | D   | D      | С    | С     | С     | С     | D   | D   | D | D |
| $(a_1, a_2) = (D, C)$    | ) C      | С | D    | D    | С    | С    | D   | D      | С    | С     | D     | D     | С   | С   | D | D |
| $(a_1, a_2) = (D, D)$    | ) C      | D | С    | D    | С    | D    | С   | D      | С    | D     | С     | D     | С   | D   | С | D |



#### Folk Theorem

- Game Setting
  - n-player infinitely-repeated game G = (N, A, u) with average reward
- Enforceable
  - A payoff profile r is **enforceable** if  $r_i \ge v_i$ ,  $v_i = \min_{s_{-i} \in S_{-i}} \max_{s_i \in S_i} u_i(s_{-i}, s_i)$
- Feasible
  - A payoff profile r is **feasible** if there exist rational, non-negative values  $\alpha_a$  such that for all i, we can express  $r_i$  as  $\sum_{a \in A} \alpha_a u_i(a)$ , with  $\sum_{a \in A} \alpha_a = 1$
- Folk Theorem
  - r is **feasible** and **enforceable**  $\Rightarrow$  r is the payoff in some Nash equilibrium

|   | C    | D    |
|---|------|------|
| С | 2, 2 | 0, 3 |
| D | 3, 0 | 1, 1 |

Prisoner's Dilemma

$$(v_1, v_2) = (1, 1)$$
  
(-1, -1) is not enforceable, not feasible  
(0.5, 2) is not enforceable, **feasible**  
(5, 5) is **enforceable**, not feasible  
(2, 2) is **enforceable**, **feasible**

# Fictitious Play

#### Definition

 Each player plays a best response to assessed strategy of the opponent and observe the opponent's actual play and update beliefs.

**Matching Pennies** 

|       | Heads | Tails |
|-------|-------|-------|
| Heads | 1, -1 | -1, 1 |
| Tails | -1, 1 | 1, -1 |

|       |            | Attitude . |                  |                  |
|-------|------------|------------|------------------|------------------|
| Round | 1's action | 2's action | 1's beliefs      | 2's beliefs      |
| 0     |            |            | (1.5, 2)         | (2, 1.5)         |
| 1     | Т          | Т          | (1.5, 3)         | (2, <b>2.5</b> ) |
| 2     | Т          | Н          | (2.5, 3)         | (2, 3.5)         |
| 3     | Т          | Н          | <b>(3.5</b> , 3) | (2, 4.5)         |
| 4     | Н          | H 10       | (4.5, 3)         | (3, 4.5)         |
|       |            |            |                  |                  |

# Convergence of Fictitious Play

- Fictitious Play → Convergence
  - Each of the following are a sufficient conditions for the empirical frequencies of play to converge in fictitious play:
    - The game is zero sum;
    - The game is solvable by iterated elimination of strictly dominated strategies;
    - The game is a potential game;
    - The game is 2 n and has generic payoffs.
- Convergence → Nash Equilibrium
  - If the empirical distribution of each player's strategies converges in fictitious play, then it converges to a Nash equilibrium.
- Results in Extensive-form Game with Imperfect Information
  - Fictitious self-play converges to approximate Nash equilibrium [13]
  - AlphaStar for StarCraft [14]

# No-regret Learning

- Regret
  - Let  $a^t$  be the action profile played at time t
  - Regret of player i for not playing action  $a'_i$  at time t is  $R^t(a'_i) = u_i(a'_i, a^t_{-i}) u_i(a^t)$
  - Regret cumulated from time 1 to T is  $CR^{T}(a_{i}') = \sum_{t=1}^{T} R^{t}(a_{i}')$
- Regret Matching
  - At each time step, each action is chosen with probability proportional to its cumulated regret:  $\sigma_i^{t+1}(a_i) = \frac{CR^t(a_i)}{\sum_{a_i' \in A_i} CR^t(a_i')}$
  - Converge to correlated equilibrium
- No-regret learning in Extensive-form Game
  - Counterfactual Regret Minimization (CFR)
  - DeepStack for Texas Hold'em poker [15]

Outline

- A. Chili.
- Motivation and Normal-form Game
- Extensive-form Game and Imperfect Information
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# **Evolutional Game Theory**

- Motivation
  - Nash equilibrium is static, the dynamic of strategy is not described
  - Players are not fully rational
- Basic Idea
  - Strategy is inherent and player can not select strategy by herself
  - Player with high payoff is has more chance to be reproduced
- Evolutionary Stable Strategy (ESS)
  - If almost every member of the population follows a strategy, no mutant (that is, an individual who adopts a novel strategy) can successfully invade.

Ref: https://plato.stanford.edu/entries/game-evolutionary/

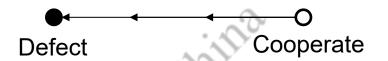
# Replicator Dynamics

#### Definition

- $\dot{x}_i = x_i [f_i(x) \phi(x)], \phi(x) = \sum_{j=1}^n x_j f_j(x)$
- x is distribution of types(strategies) over the population
- $f_i(x)$  is the fitness for type i in population x
- $\varphi(x)$  is the average fitness of the population

| 7 | С    | D    |
|---|------|------|
| С | 2, 2 | 0, 3 |
| D | 3, 0 | 1, 1 |

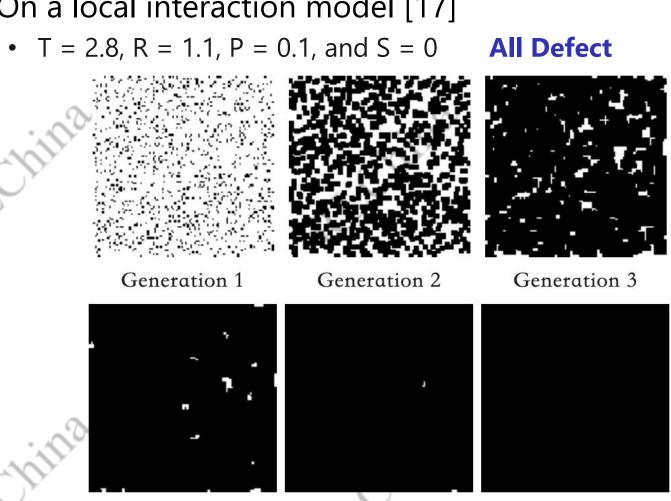
Prisoner's Dilemma



## Replicator Dynamics: Experiment

• On a local interaction model [17]

Generation 4

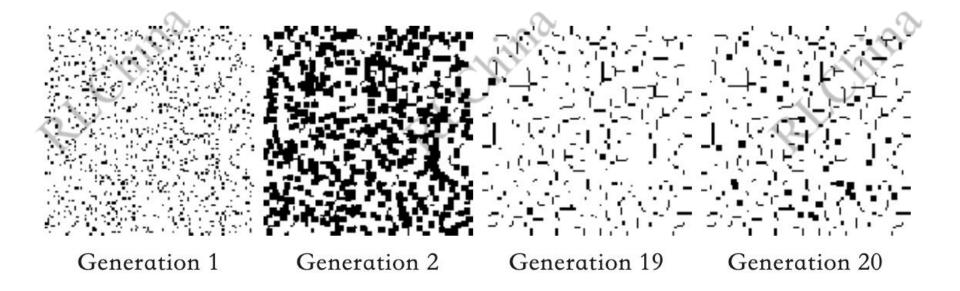


Generation 5

Generation 6

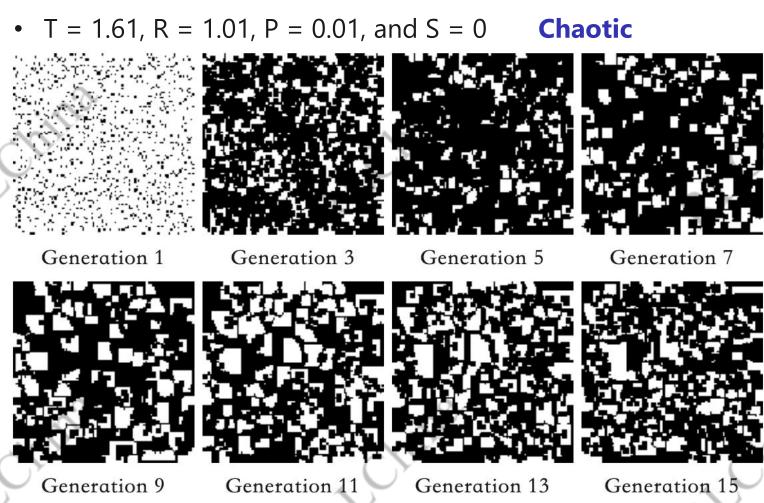
### Replicator Dynamics: Experiment

- On a local interaction model
  - T = 1.2, R = 1.1, P = 0.1, and S = 0 **Cooperate**



# Replicator Dynamics: Experiment

On a local interaction model



# **Coalitional Game Theory**

#### Definition

- N is a set of players indexed by i
- $v: 2^N \to \mathbb{R}$  associates with each coalition  $S \subseteq N$  a payoff v(S),  $v(\emptyset) = 0$

#### Question

- Which coalition will form?
- How to allocate payoff among coalition members?

# Shapley Value

- Definition
  - $\varphi_i(N, v) = \frac{1}{N!} \sum_{S \subseteq N \setminus \{i\}} |S|! (|N| |S| 1)! [v(S \cup \{i\}) v(S)]$
- Property
  - Symmetry
    - $\varphi_i(N,v) = \varphi_j(N,v)$  if  $v(S \cup \{i\}) = v(S \cup \{j\})$  for  $\forall S, i, j \notin S$
  - Dummy Player
    - $\varphi_i(N, v) = 0$  if  $v(S \cup \{i\}) = v(S)$  for  $\forall S$
  - Additivity
    - $\varphi_i(N,v_1+v_2)=\varphi_i(N,v_1)+\varphi_i(N,v_2)$  for  $\forall i$ , where  $(v_1+v_2)(S)=v_1(S)+v_2(S)$  for  $\forall S$

#### Core

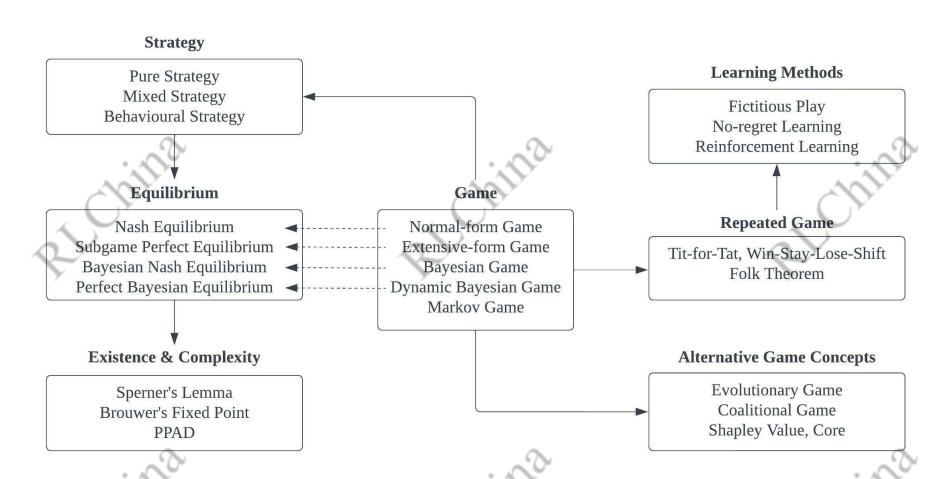
#### Definition

• A payoff vector x is in the **core** of a coalition game (N, v), iff  $\forall S \subseteq N, \sum_{i \in S} x_i \ge v(S)$ 

#### Property

- A payoff vector in the core is a stable distribution of the grand coalition
- In some sense, core is a more stable solution concept than Nash equilibrium because every group of the players, rather than every single player, has no intention to deviate

#### Summary



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