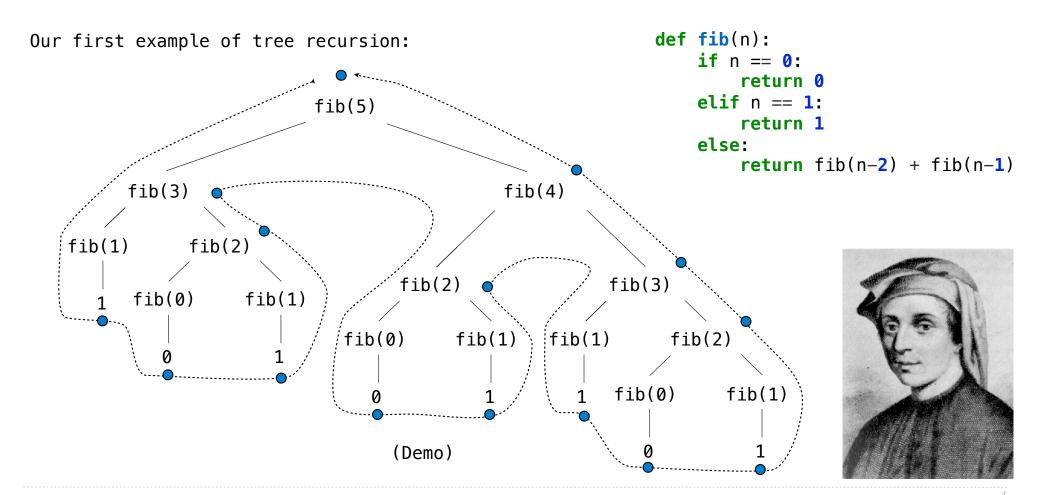
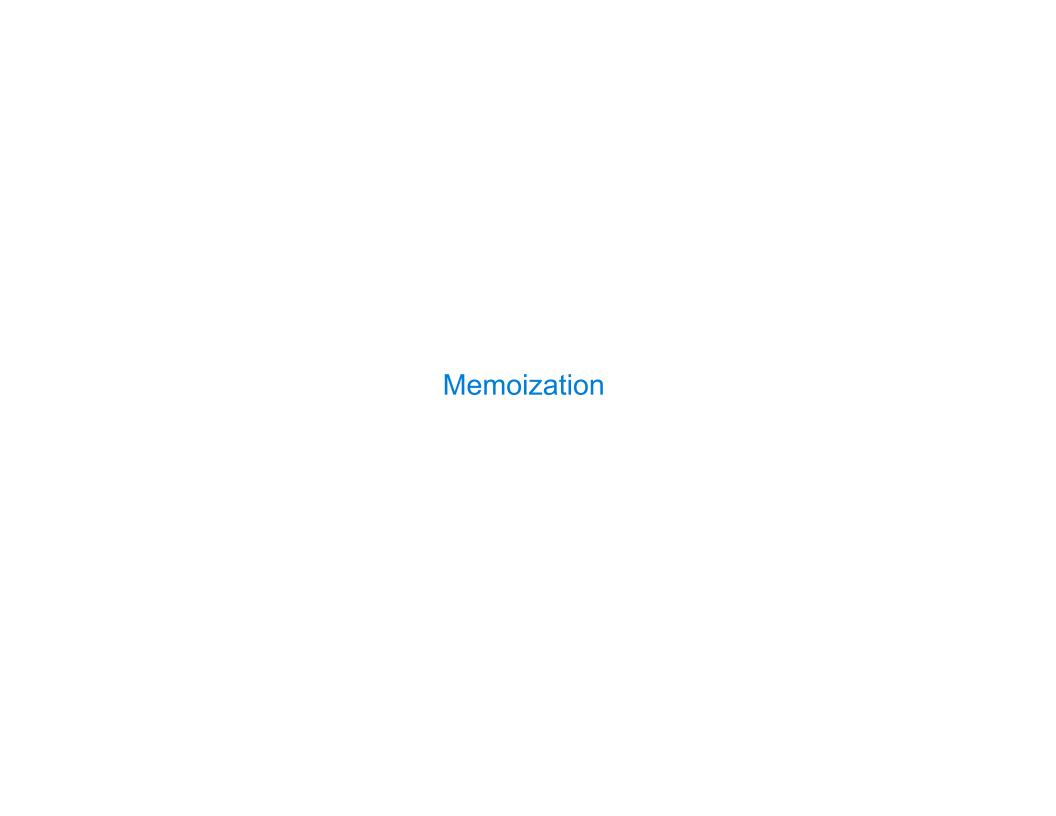


## Recursive Computation of the Fibonacci Sequence





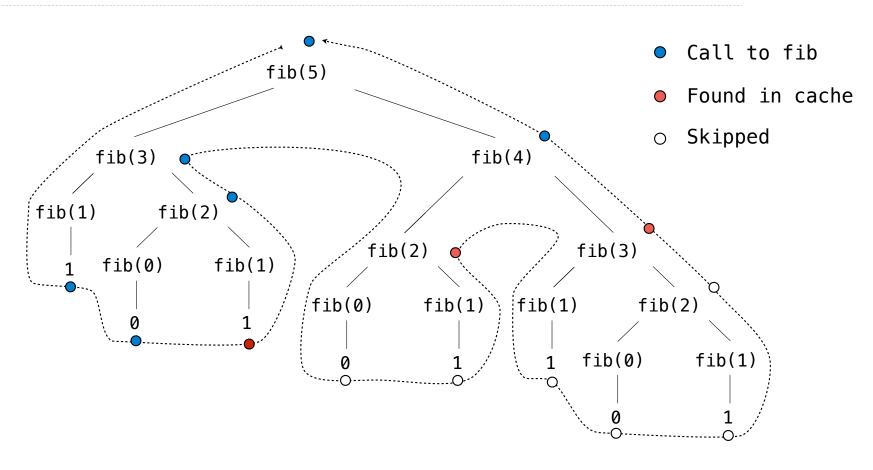
## Memoization

**Idea:** Remember the results that have been computed before 保存已经计算过的函数值

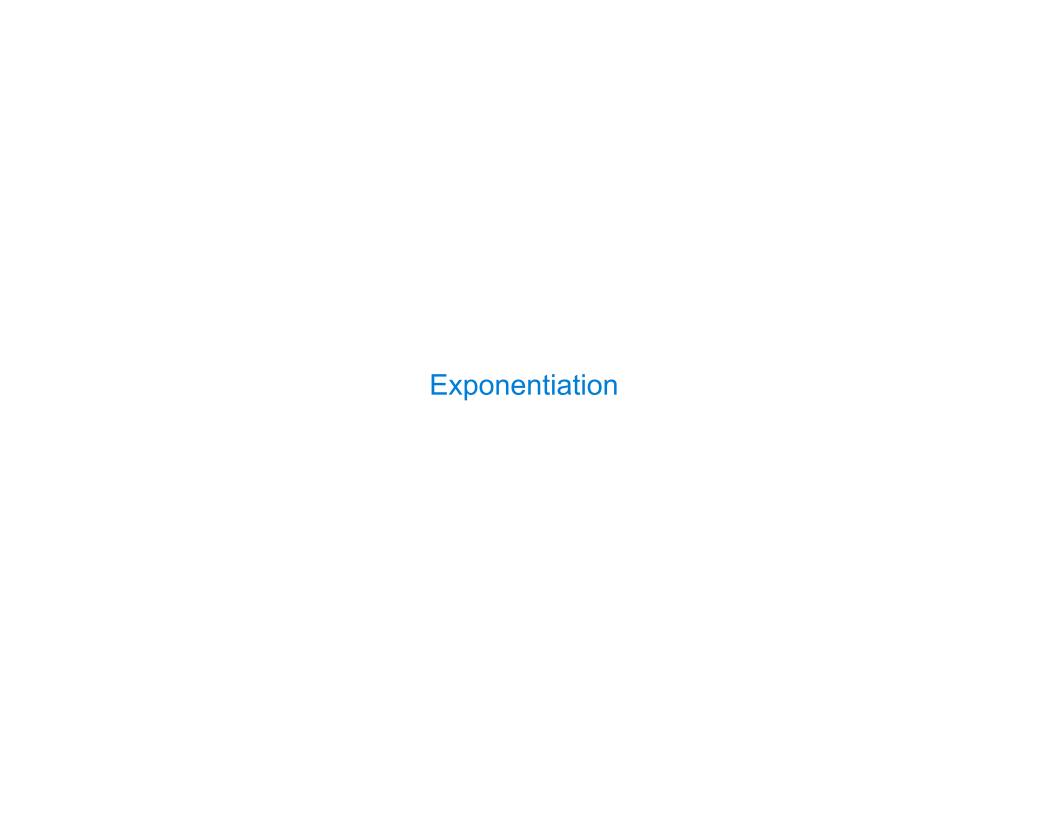
(Demo)

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## Memoized Tree Recursion



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## Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
                                                                                 b^n = \begin{cases} 1 & \text{if } n = 0\\ b \cdot b^{n-1} & \text{otherwise} \end{cases}
       if n == 0:
              return 1
       else:
              return b * exp(b, n-1)
def exp_fast(b, n):
       if n == 0:
                                                                                                  快速幂算法
              return 1
       elif n % 2 == 0:
                                                                                  b^{n} = \begin{cases} 1 & \text{if } n = 0\\ (b^{\frac{1}{2}n})^{2} & \text{if } n \text{ is even}\\ b \cdot b^{n-1} & \text{if } n \text{ is odd} \end{cases}
              return square(exp_fast(b, n//2))
       else:
              return b * exp_fast(b, n-1)
def square(x):
       return x * x
```

(Demo)

## Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
    if n == 0:
        return 1
    else:
        return b * exp(b, n-1)

def exp_fast(b, n):
    if n == 0:
        return 1
    elif n % 2 == 0:
        return square(exp_fast(b, n//2))
    else:
        return b * exp_fast(b, n-1)

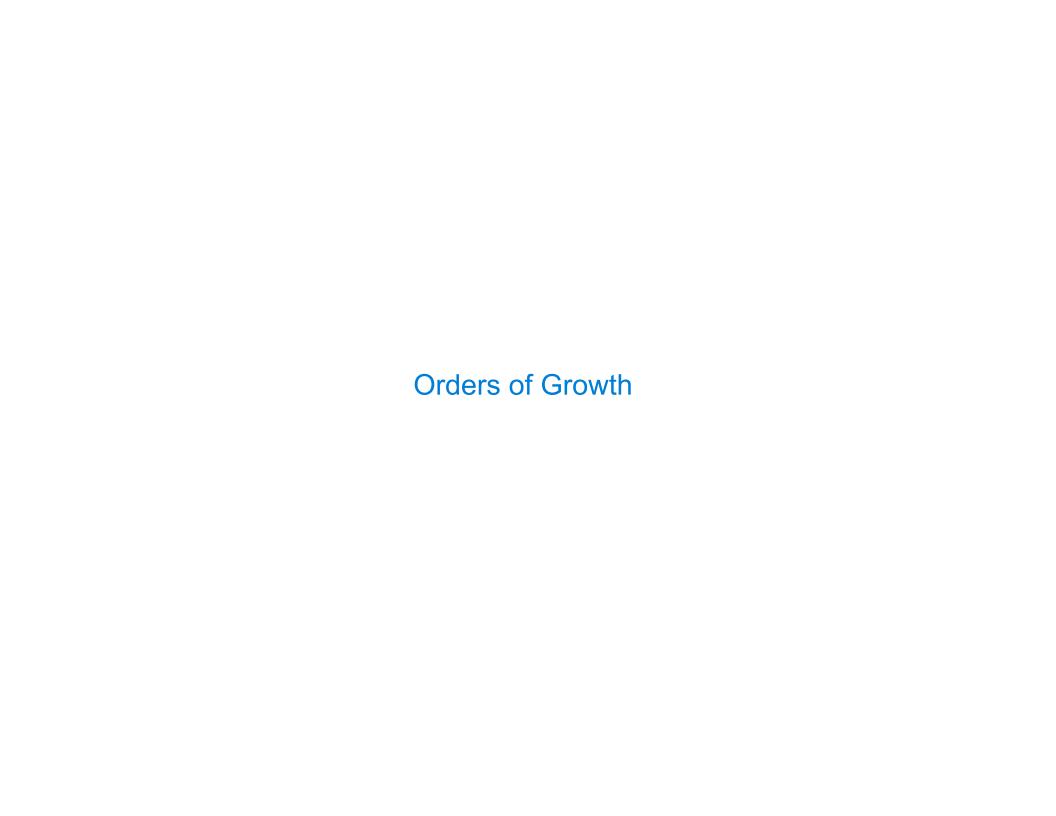
def square(x):
    return x * x
```

#### Linear time:

- Doubling the input doubles the time
- 1024x the input takes 1024x as much time

#### Logarithmic time:

- Doubling the input increases the time by a constant C
- 1024x the input increases the time by only 10 times C



## **Quadratic Time**

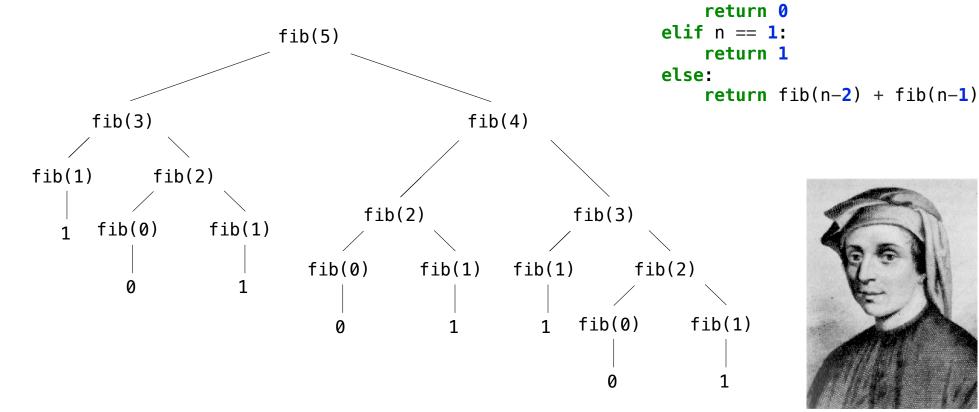
Functions that process all pairs of values in a sequence of length n take quadratic time

```
3
                                                                    7
                                                                          6
def overlap(a, b):
    count = 0
                                                         0
                                                               0
                                                   4
    for item in a:
        for other in b:
                                                   5
            if item == other:
                 count += 1
    return count
                                                               0
                                                         0
                                                   6
overlap([3, 5, 7, 6], [4, 5, 6, 5])
                                                         0
                                                               1
                                                                    0
                                                                          0
                                                   5
```

(Demo)

## **Exponential Time**

Tree-recursive functions can take exponential time





def fib(n):

**if** n == **0**:

Time for n+n

Time for input n+1

Time for input n

#### Common Orders of Growth

**Exponential growth.** E.g., recursive fib Incrementing *n* multiplies *time* by a constant

$$a \cdot b^{n+1} = (a \cdot b^n) \cdot b$$

Quadratic growth. E.g., overlap

Incrementing n increases time by n times a constant

$$a \cdot (n+1)^2 = (a \cdot n^2) + a \cdot (2n+1)$$

Linear growth. E.g., slow exp

Incrementing n increases time by a constant

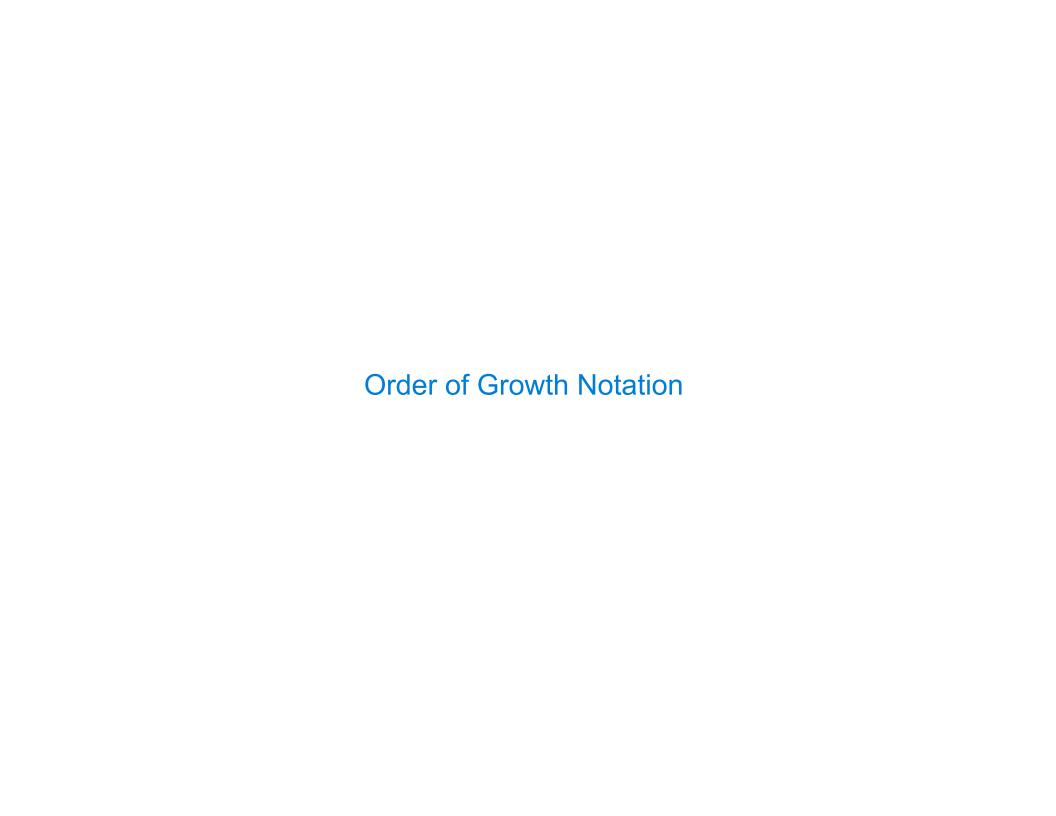
$$a \cdot (n+1) = (a \cdot n) + a$$

Logarithmic growth. E.g., exp\_fast

Doubling n only increments time by a constant

$$a \cdot \ln(2 \cdot n) = (a \cdot \ln n) + a \cdot \ln 2$$

Constant growth. Increasing n doesn't affect time



	Lower bound and Upper bound	Upper bound
Big Theta and Big O Notation for Orders of Growth		
<b>Exponential growth.</b> E.g., recursive fib Incrementing $n$ multiplies $time$ by a constant	$\Theta(b^n)$	$O(b^n)$
Quadratic growth. E.g., overlap Incrementing $n$ increases $time$ by $n$ times a constant	$\Theta(n^2)$ tant	$O(n^2)$
<b>Linear growth.</b> E.g., slow exp Incrementing $n$ increases $time$ by a constant	$\Theta(n)$	O(n)
<b>Logarithmic growth.</b> E.g., $exp_fast$ Doubling $n$ only increments $time$ by a constant	$\Theta(\log n)$	$O(\log n)$
Constant growth. Increasing $n$ doesn't affect times	ne $\Theta(1)$	O(1)



## **Space and Environments**

Which environment frames do we need to keep during evaluation?

At any moment there is a set of active environments

Values and frames in active environments consume memory

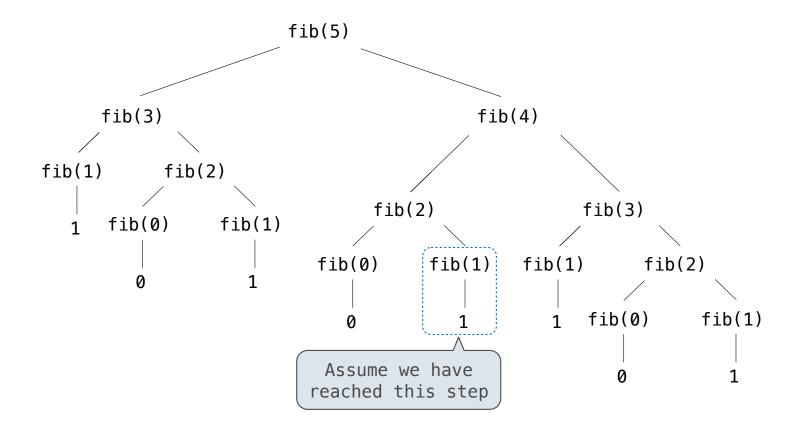
Memory that is used for other values and frames can be recycled

#### **Active environments:**

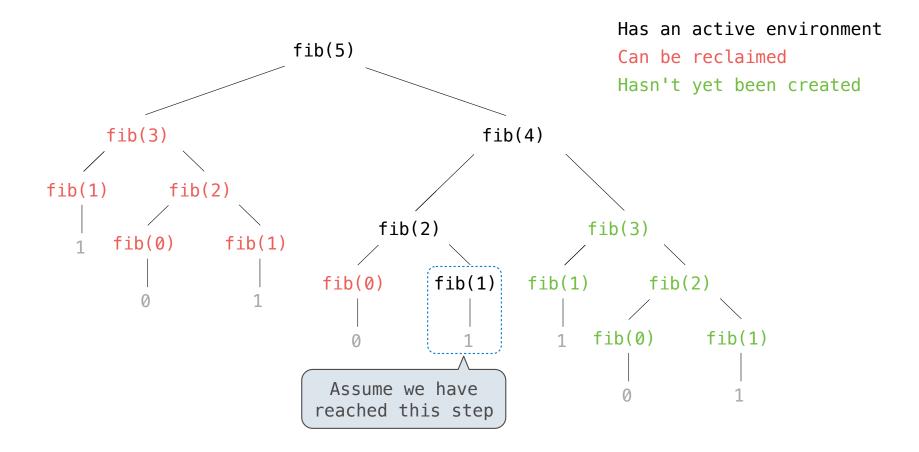
- Environments for any function calls currently being evaluated
- Parent environments of functions named in active environments

(Demo)

# Fibonacci Space Consumption



# Fibonacci Space Consumption



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