

## Warmup 1

1.

$$\{t \mid \exists s \in \text{Tree}(s[t\_id] = t[t\_id] \wedge \\ \forall r \in \text{Tree}(r[\text{years\_to\_maturity}] \leq s[\text{years\_to\_maturity}]))\}$$

2.

$$\{n \mid \exists n_1 \in \text{Nursery}(n_1[n\_id] = n[n\_id] \wedge \\ \exists s \in \text{Sells}(n_1[n\_id] = s[n\_id] \wedge \\ \forall t \in \text{Tree}(t[t\_id] = s[t\_id] \wedge t[\text{height}] \geq 10))))\}$$

3.

$$\{n \mid \exists n_1 \in \text{Nursery}(n_1[n\_id] = n[n\_id] \wedge n_1[\text{state}] = \text{'California'} \wedge \\ \exists s \in \text{Sells}(n_1[n\_id] = s[n\_id] \wedge s[\text{price}] \geq 20 \wedge \\ \exists t \in \text{Tree}(t[t\_id] = s[t\_id] \wedge t[\text{height}] = 10))))\}$$

4.

$$\{n \mid \forall n_1 \in \text{Nursery}(n_1[n\_id] = n[n\_id1] \wedge \\ \forall n_2 \in \text{Nursery}(n_2[n\_id] = n[n\_id2] \wedge \\ \exists s_1 \in \text{Sell}(s_1[n\_id] = n_1[n\_id] \wedge \\ \exists s_2 \in \text{Sell}(s_2[n\_id] = n_2[n\_id] \wedge \\ n_1[\text{price}] > n_2[\text{price}]))))))\}$$

5.

- (a) Find id and name of trees that sold by the nursery named as "Johnny Appleseed"
- (b) Find id and name of all Nurseries that sold any kind of trees for highest price

## Warmup 2

1.

A query is safe when the result is a finite set of entities.

2.

$\{t \mid \neg(t \in Tress) \wedge t \in Shrubs\}$  is safe, since the result is finite.

$\{t \mid \neg(t \in Tress) \vee t \in Shrubs\}$  is not safe, because the result can be infinite.

## Warmup 3

1.

**Person** (p\_id, (first\_name, middle\_name, last\_name), age, (street\_name, state, zip\_code, country), {phone\_number})

**Gardener** (g\_id, years\_of\_experience)

**Tree** (species, max\_height, years\_to\_maturity)

2.

Job: relating gardener with person

Cultivate: relating gardener with tree

3.

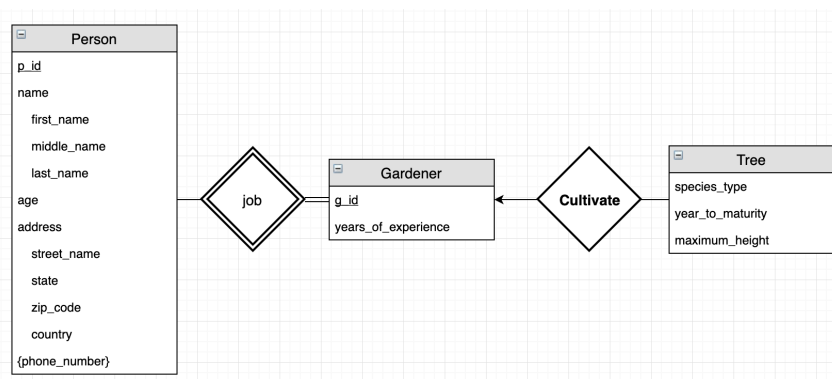


Figure 1: ER diagram

4.

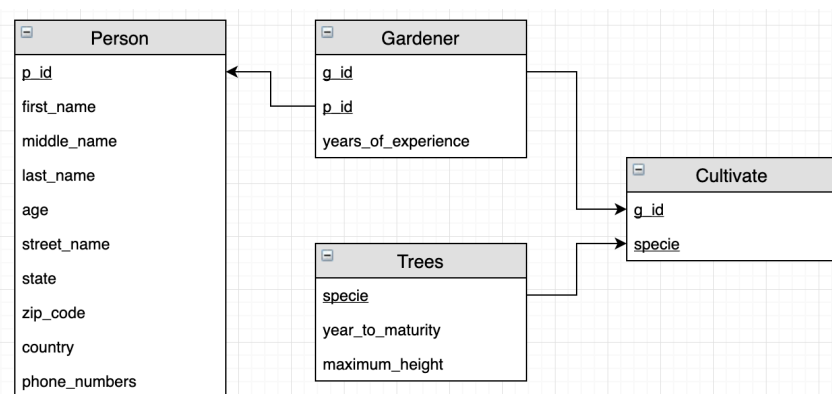


Figure 2: Database schema

## Problem 4

1.

**User** (user\_id, (first\_name, last\_name), email, minor, (guardian\_first\_name, guardian\_last\_name))

**Explanation** : The customer and passenger can be the same person, and the information for customer overlapped with the information for passenger. Therefore, I create a new entity called user to combine customer and passenger.

**Route**(route\_id, stop\_number, origin, destination, stop1, stop2)

**Explanation** : The entity route indicates all possible routes. A route\_id is used to identify a route. The stop\_number is derived attribute to show the number of stops in the current route.

**Flight**(flight\_number, departure\_time, route\_id, arrive\_time, aircraft\_id, first\_section\_available\_number, {first\_section\_available\_seats}, business\_section\_available\_number,{business\_section\_available\_seats}, economy\_section\_available\_number,{economy\_section\_available\_seats})

**Explanation** : The entity flight saves the information of a flight. The reason why I didn't create a entity for aircraft is that the same aircraft may serve different flight and it's impossible to record available seats for different flight in one entity.

2.

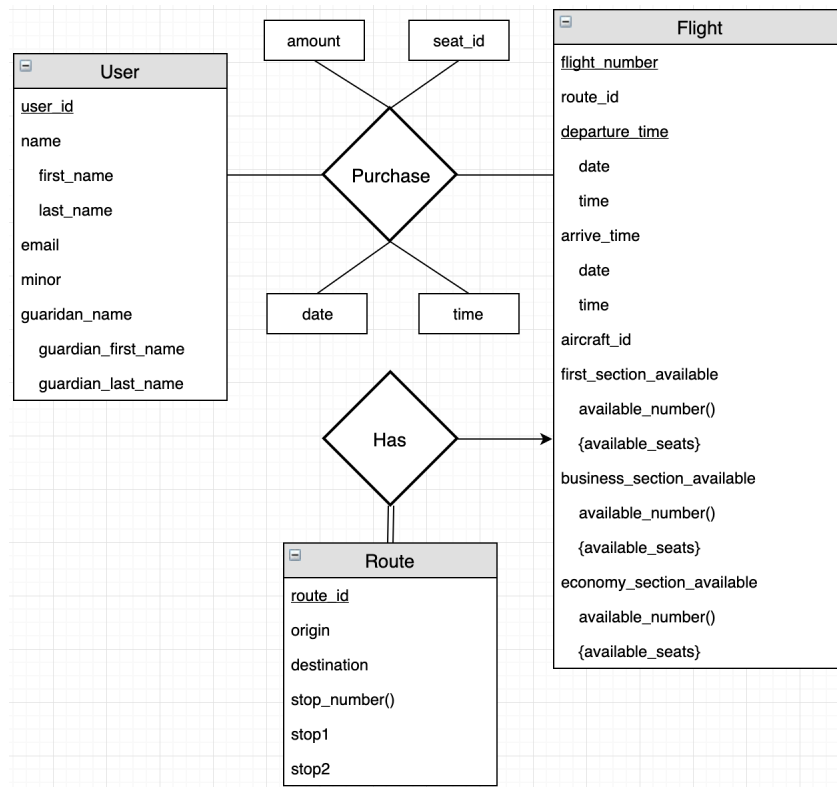


Figure 3: ER diagram

3.

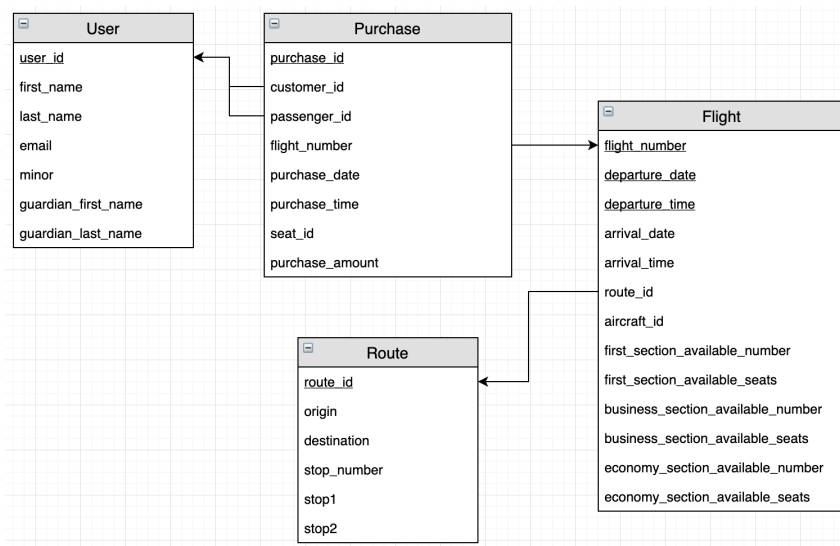


Figure 4: ER diagram

## Problem 5

1.

### Relational algebra

# select the route\_id from the Table Route where origin is "PVD" and the destination is "LAX" and the number of stops is no more than 1

$$Temp1 \leftarrow \pi_{route\_id}(\sigma_{origin='PVD' \wedge destination='LAX' \wedge stop\_number \leq 1}(Route))$$

# natural join Temp1 and Flight and then select the flight that the departure date is September 30, 2019 and has business seats available

$$Result \leftarrow \sigma_{departure\_date='September30,2019' \wedge business\_section\_available\_number > 0}(Flight \bowtie Temp1)$$

### Tuple relational calculus

$$\{f \mid f \in Flight \wedge f[departure\_date] = "Sep30, 2019" \wedge f[business\_section\_available\_number] > 0 \wedge \\ \exists r \in Route(r[route\_id] = f[route\_id] \wedge r[origin] = "PVD" \wedge r[destination] = "LAX" \wedge \\ r[stop\_number] \leq 1)\}$$

2.

Let "destination" be the group's destination.

Let "location" be their current airport location.

Let "date" be the current date.

Let "time" be the current time.

### Relational algebra

# select the route\_id from the Table Route where origin is current location and destination is their original destination

$$Temp1 \leftarrow \pi_{route\_id}(\sigma_{origin='location' \wedge destination='destination' \wedge stop\_number \leq 1}(Route))$$

# natural join Temp1 and Flight and then select the flight that the departure date is later than current time

$$Temp2 \leftarrow \sigma_{departure\_date \leq 'date' \wedge departure\_time > 'time' > 0}(Flight \bowtie Temp1)$$

# select the flights that have more than 3 available seats

$$Result \leftarrow \sigma_{first\_section\_available\_number + business\_section\_available\_number + economy\_section\_available\_number \geq 3}(Temp2)$$

### Tuple relational calculus

$$\{f \mid f \in Flight \wedge f[departure\_time] > "time" \wedge f[departure\_date] \geq "date" \wedge \\ f[first\_section\_available\_number] + f[business\_section\_available\_number] + \\ f[economy\_section\_available\_number] \geq 3 \wedge \\ \exists r \in Route(r[route\_id] = f[route\_id] \wedge r[origin] = "location" \wedge \\ r[destination] = "destination")\}$$

## Problem 6

1.

- prompt response: TAs give prompt response on Piazza.
- great recitation: Recitation is even better than lecture.
- too much homework: I spent more than 10 hours on both HW1 and HW2.

2. "Big Data" processing

3. Capture for recitation.