Recitation for indexing

Part1 Btree

Storage Types

- → Cache
 - Fastest/most costly; volatile;
- → Main Memory
 - ◆ Fast access; too small for entire db; volatile
- → Disk
 - Long-term storage of data; random access; non-volatile
- → Flash Memory
 - No seeks (cheap reads);

How Disk Works

- → Surface divided into tracks -> sectors (smallest unit of data that can be R/W)
 - ◆ Disk arm swings to position head on right track
 - Platter spins continually as data is R/W from sector
- → Measuring Disk Speed
 - Access Time:
 - Seek Time: time to find right track
 - Latency Time: time for sector to appear under head
 - Data Transfer Rate:
 - The rate at which data can be retrieved
- → Sequential I/O < Random I/O</p>

Problem Statement

- → Many queries reference only a fraction of records
- → It is inefficient for the system to read every single tuple
- → A better design is for the system to locate reference directly

Index

→ Indexes

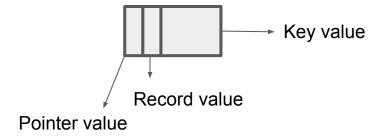
- Auxiliary data structures over relations that can improve the search time
- ◆ Think about index in a textbook
- We search for an index, find corresponding pages, and then read information to find what we are looking for
- The index is much smaller than the book and has words sorted in alphabetical order

B-trees

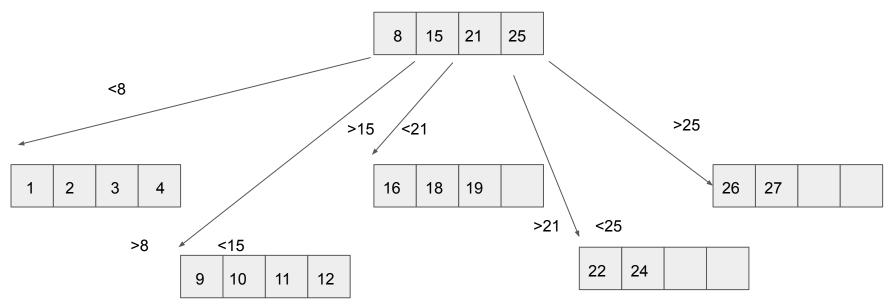
- → Most successful family of index schemes (B-trees, B+-trees,...)
- → Can be used for primary/secondary, clustering/non-clustering index
- → Balanced "n-way" search trees
 - ♦ Helpful to think of it as a tree structure with n-pointers in each node

B-tree nodes

- → Key values are ordered
- → MAXIMUM: n pointers
- → MINIMUM: [n/2] pointers
 - ◆ Exception: root's minimum = 2
- → Internal nodes must have n-1 key values where n is the number of pointers

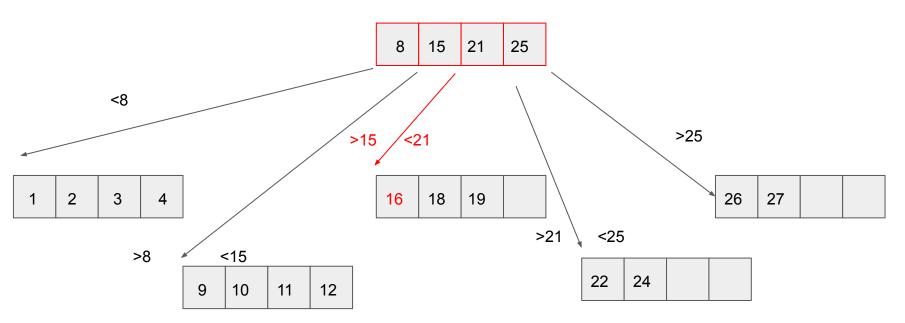


B-tree of order 5

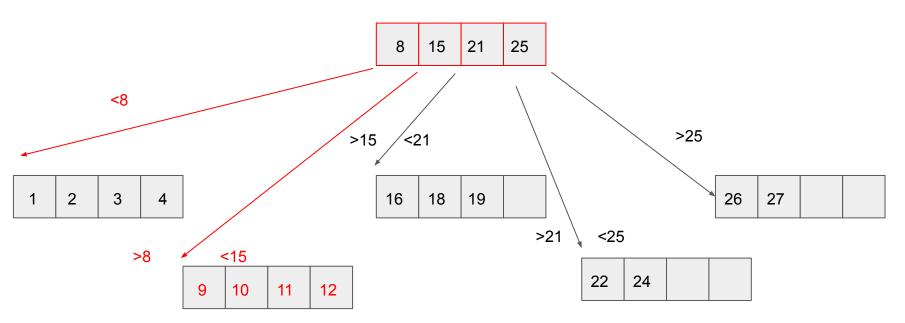


- → Key values appear once
- → Record pointers accompany keys

- → Let's run through an exact match query
 - ◆ Let's find the value 16?



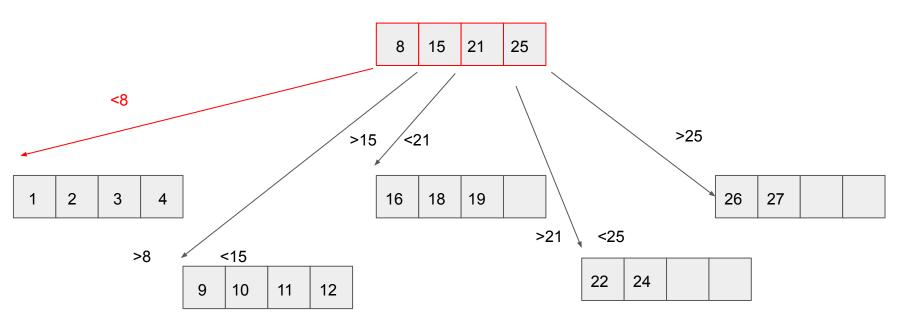
- → Let's run through a range query
 - ♦ Let's find the values between 6 and 13?



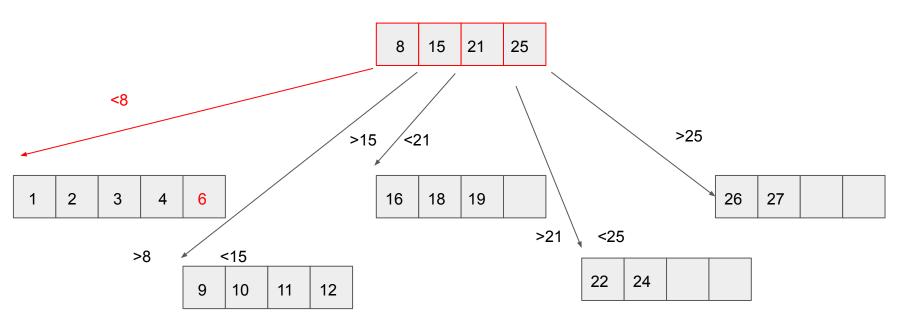
How To Maintain B-trees?

- → B-tree rules must be obeyed on every insert + delete
- → Rules:
 - ♦ Insert in leaf, if room exists
 - On overflow (no more room)
 - Split: create a new internal node
 - Redistribute keys
 - So that it preserves B-tree properties
 - Push middle key up (recursively)

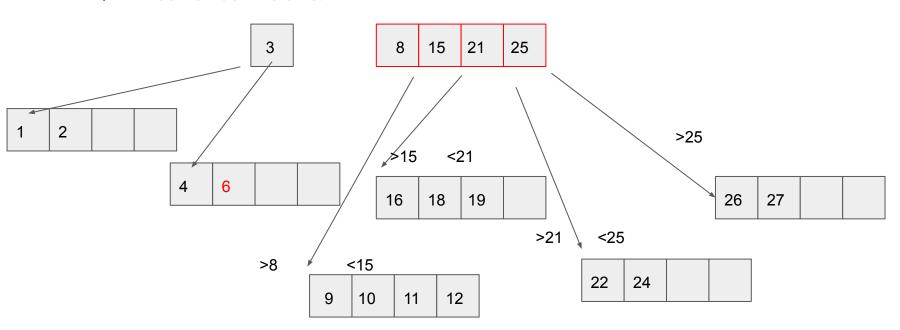
- → Let's run through an overflow insertion
 - ◆ Let's insert the value 6:



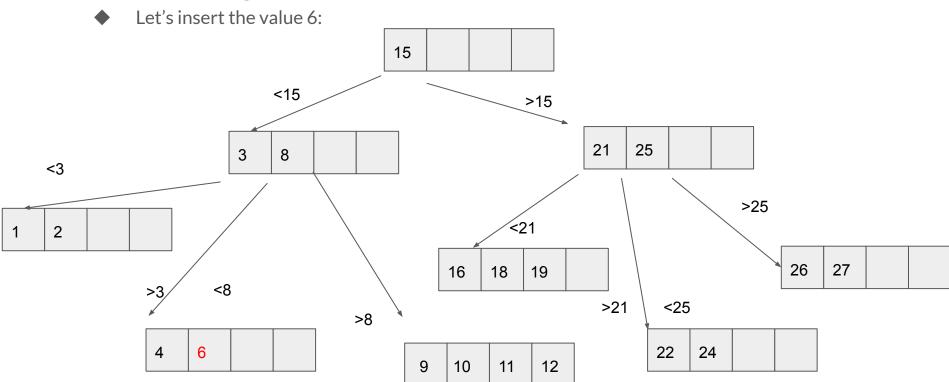
- → Let's run through an overflow insertion
 - ◆ Let's insert the value 6:



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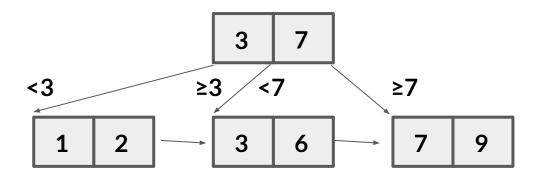
→ Let's run through an overflow insertion



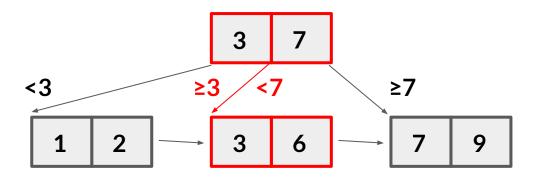
B+-trees

- → Allow sequential operations
 - String all leaf nodes together
 - Every key appears at leaf level (some keys show up more than once)
- → This mean non-leaf nodes do not contain data for the key, only the pointer value and key value

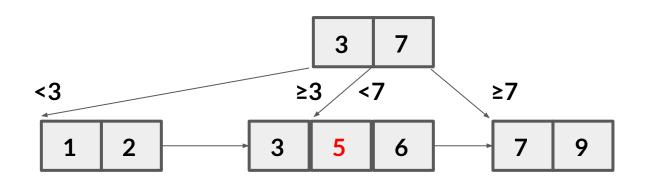
→ Insert 5 into B-tree of order 3



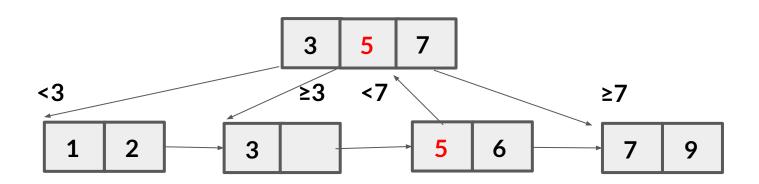
→ Insert 5

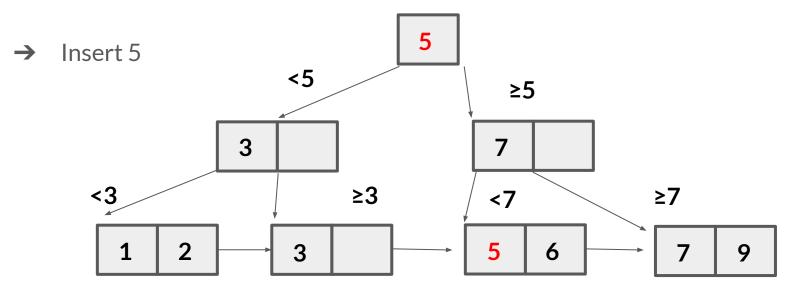


→ Insert 5



→ Insert 5





Recitation for indexing

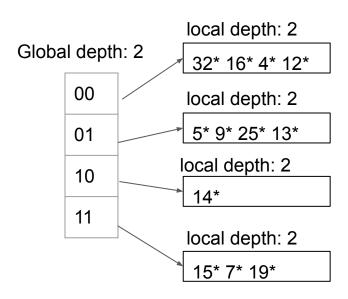
Part-2 hashing

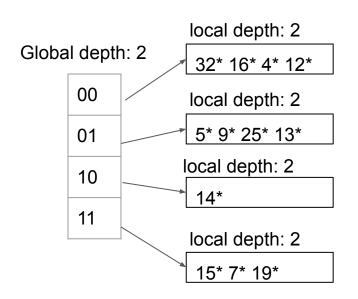
When do we need hashing?

Hash-based indexes are best for equality selections. Cannot support range searches.

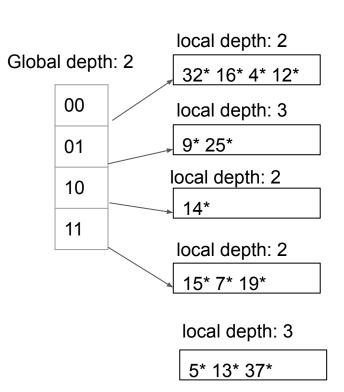
Two kinds of hashing:

- Static hashing: will degrade performance with long overflow chains
- Dynamic hashing: Extendible hashing and linear hashing



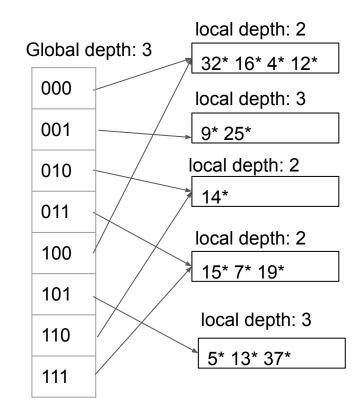


Insert 37

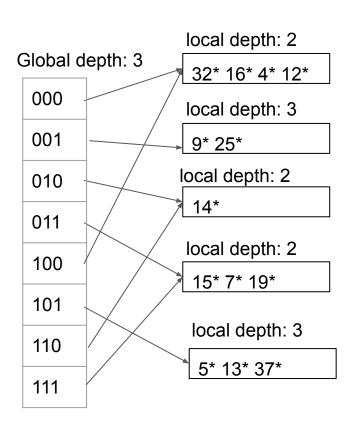


local depth: 2 Global depth: 2 32* 16* 4* 12* 00 local depth: 3 9* 25* 01 local depth: 2 10 14* 11 local depth: 2 15* 7* 19* local depth: 3 5* 13* 37*

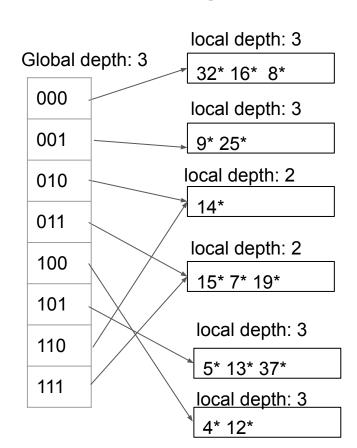
Insert 37



What will happen when Inserting 8*?

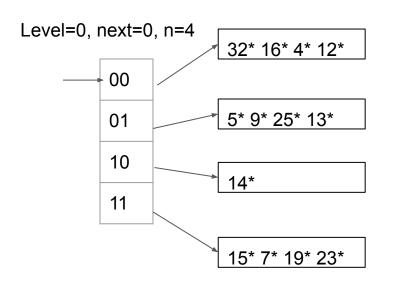


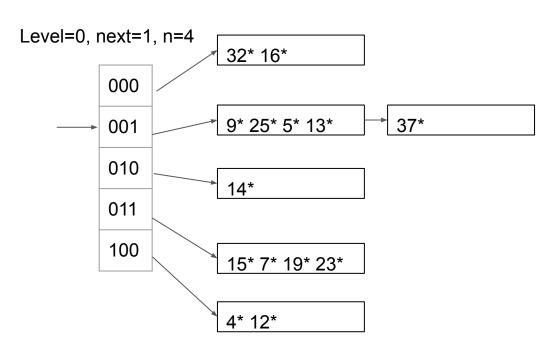
What will happen when Inserting 8*?



Linear Hashing

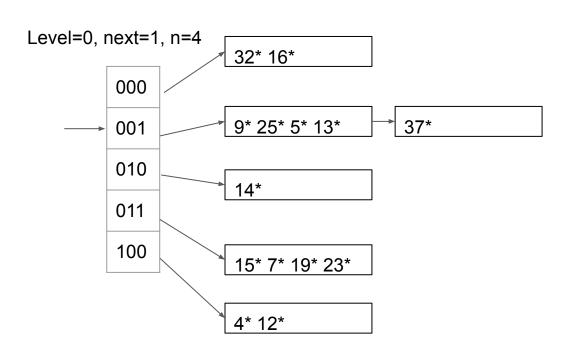
Insert 37





Linear Hashing

What will happen when inserting 27?



Linear Hashing

What will happen when inserting 27?

