

# MAIN MEMORY

PAGING - IMPORTANT FORMULAS

## POWER OF 2

### FOR MAIN MEMORY

Physical Address Space = Size of main memory ( $2^M$ )

**Size of main memory** = Total number of frames  $\times$  Frame size

Frame size = Page size ( $2^F = 2^P$ )

**Total number of frames** = Size of main memory / Frame size

If number of frames in main memory =  $2^M / 2^F = 2^{M-F} = 2^X$ , then number of bits in frame number = **X** bits

If Page size =  $2^P$ , then number of bits in page = **P** bits

If size of main memory =  $2^M$ , then number of bits in physical address = **M** bits

page number	page offset
p	d
m - n	n

### FOR PROCESS

Virtual Address Space = Size of process

Number of pages the process is divided = Process size / Page size

If process size =  $2^X$ , then number of bits in virtual address space = **X** bits

### FOR PAGE TABLE

Size of page table = Number of entries in page table  $\times$  Page table entry size

Number of entries in pages table = Number of pages the process is divided

Page table entry size = Number of bits in frame number + Number of bits used for optional fields, if any.

In general, if the given address consists of 'n' bits, then using 'n' bits,  $2^n$  locations are possible.

- Then, size of memory =  $2^n \times$  Size of one location.

If the memory is **byte-addressable**, then size of one location = 1 byte.

- Thus, size of memory =  $2^n$  bytes.

If the memory is **word-addressable** where 1 word = m bytes, then size of one location = m bytes.

- Thus, size of memory =  $2^n \times m$  bytes.