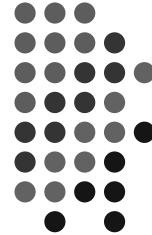


Trees

6B

Heaps & Other Trees



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1

Heap

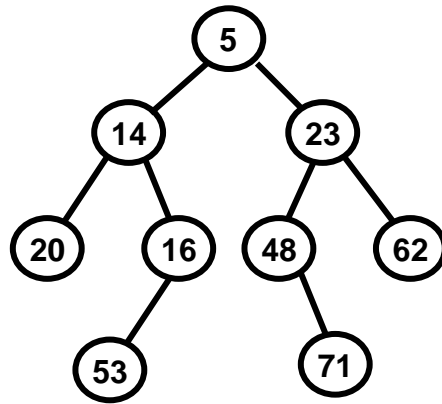


- A min-heap is a binary tree such that
 - the data contained in each node is less than (or equal to) the data in that node's children.
 - the binary tree is complete
- A max-heap is a binary tree such that
 - the data contained in each node is greater than (or equal to) the data in that node's children.
 - the binary tree is complete

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2

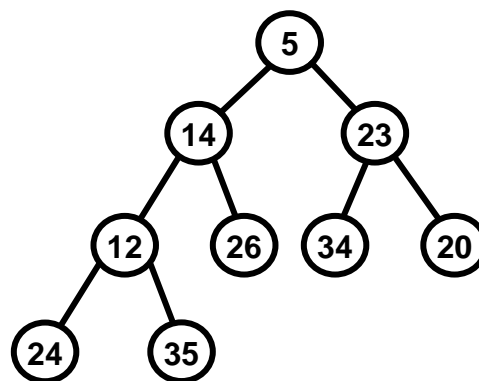
Is it a min-heap?



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3

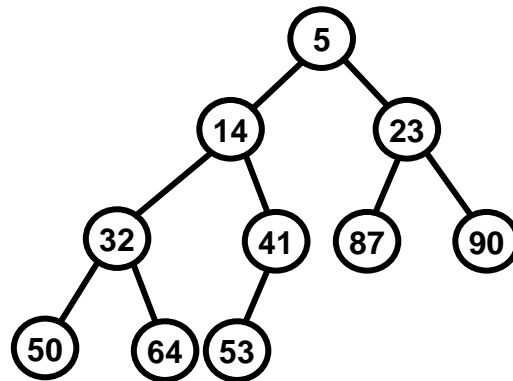
Is it a min-heap?



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4

Is it a min-heap?



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5

Using heaps



What are min-heaps good for?
(What operation is extremely fast when
using a min-heap?)

The difference in level between any two leaves
in a heap is at most what?

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6



Storage of a heap

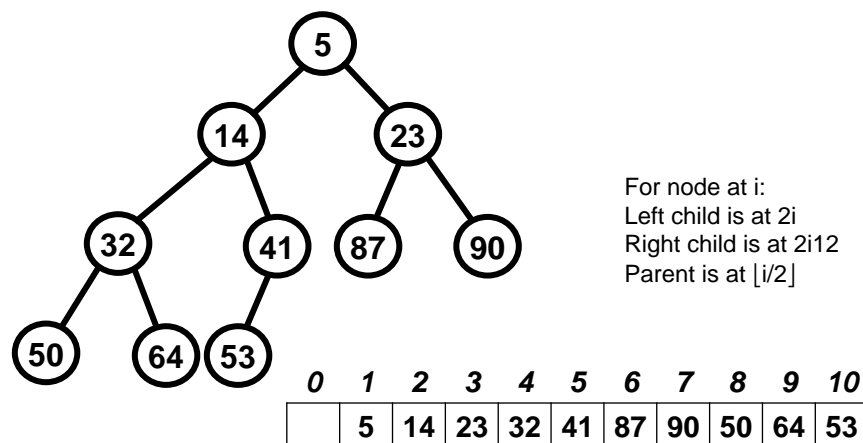
- Use an array to hold the data.
- Store the root in position 1.
 - We won't use index 0 for this implementation.
- For any node in position i ,
 - its left child (if any) is in position $2i$
 - its right child (if any) is in position $2i + 1$
 - its parent (if any) is in position $i/2$
(use integer division)

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Storage of a heap



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8

Inserting into a min-heap



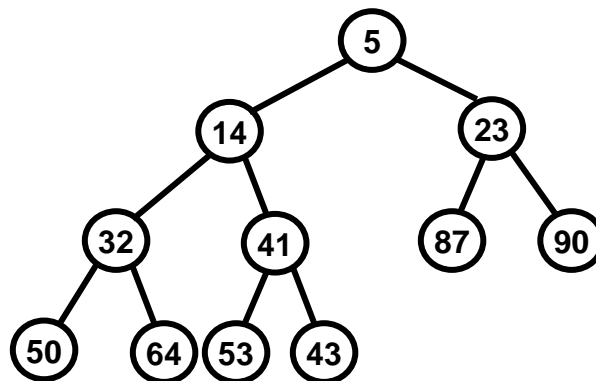
- Place the new element in the next available position in the array.
- Compare the new element with its parent. If the new element is smaller, then swap it with its parent.
- Continue this process until either
 - the new element's parent is smaller than or equal to the new element, or
 - the new element reaches the root (index 0 of the array)

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9

Inserting into a min-heap

Insert 43

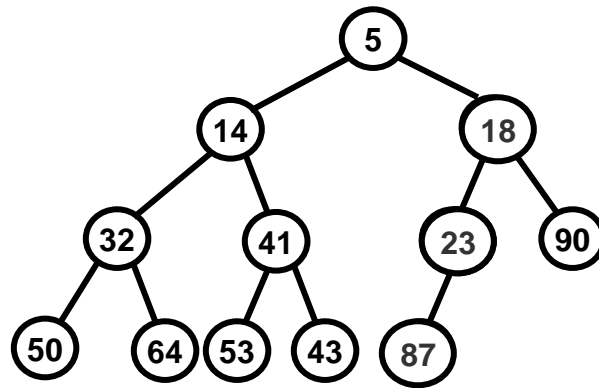


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Inserting into a min-heap

Insert 18

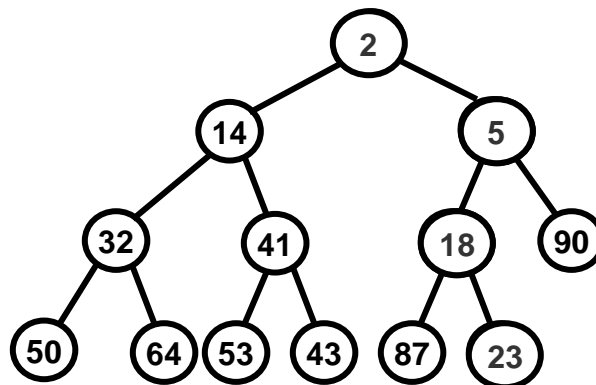


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11

Inserting into a min-heap

Insert 2



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12

Removing from a heap



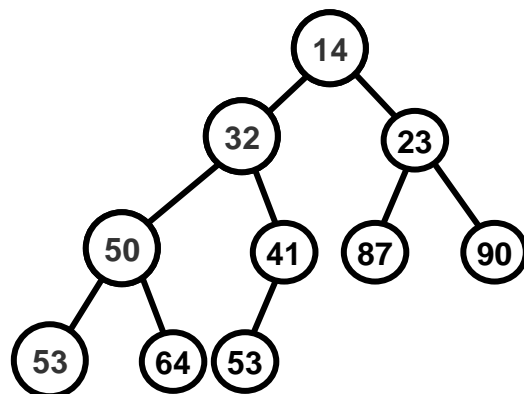
- Place the root element in a variable to return later.
- Remove the last element in the deepest level and move it to the root.
- While the moved element has a value greater than at least one of its children, swap this value with the smaller-valued child.
- Return the original root that was saved.

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13

Removing from a min-heap

Remove min



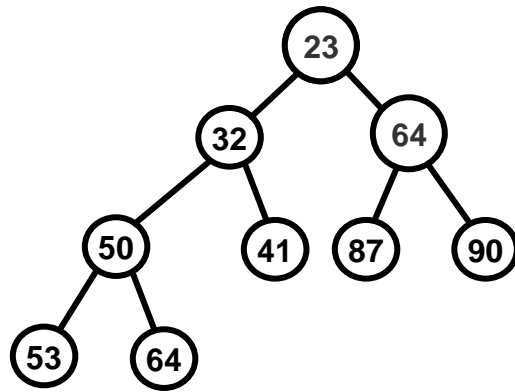
returnValue 5

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14

Removing from a min-heap

Remove min



returnValue 14

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15

Efficiency of heaps



Assume the heap has N nodes.

Then the heap has $\lceil \log_2(N+1) \rceil$ levels.

- Insert

Since the insert swaps at most once per level, the order of complexity of insert is $O(\log N)$

- Remove

Since the remove swaps at most once per level, the order of complexity of remove is also $O(\log N)$

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