

INT102

Algorithmic Foundations And Problem Solving Graph Theory

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Learning outcomes

- ✓ Able to tell what is an undirected graph and what is a directed graph
- ✓ Know how to represent an undirected graph using matrix and list
- ✓ Understand what Euler path / circuit and able to determine whether such path / circuit exists in an undirected graph
- ✓ Able to apply BFS and DFS to traverse a graph
- Know how to represent a graph using matrix and list
- Able to tell what a tree is

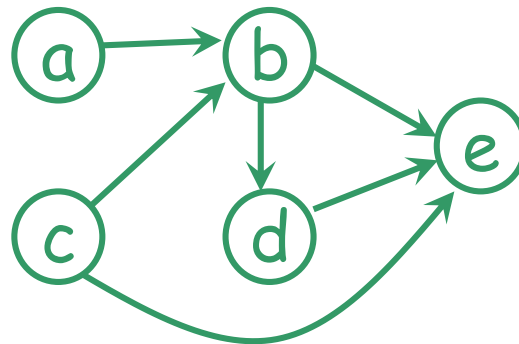
Directed graph ...

Directed graph

Given a directed graph G , a vertex a is said to be **connected to** a vertex b if there is a path from a to b .

E.g., G represents the routes provided by a certain airline. That means, a vertex represents a city and an edge represents a flight from a city to another city. Then we may ask question like: Can we fly from one city to another?

Reminder: A directed graph $G=(V,E)$ consists of a set of vertices V and a set of edges E . Each edge is an ordered pair of vertices.



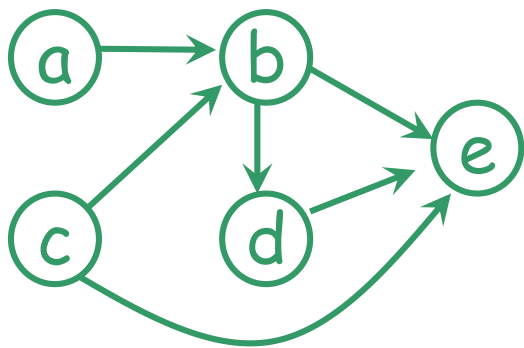
$E = \{ (a,b), (b,d), (b,e), (c,b), (c,e), (d,e) \}$

N.B. (a,b) is in E , but (b,a) is NOT

In/Out degree (in directed graphs)

The in-degree of a vertex v is the number of edges *leading to* the vertex v .

The out-degree of a vertex v is the number of edges *leading away* from the vertex v .



	<u>in-deg(v)</u>	<u>out-deg(v)</u>
a	0	1
b	2	2
c	0	2
d	1	1
e	3	0

sum: 6

6

Always equal?

Representation (of directed graphs)

Similar to undirected graph, a directed graph can be represented by adjacency matrix, adjacency list, incidence matrix or incidence list.

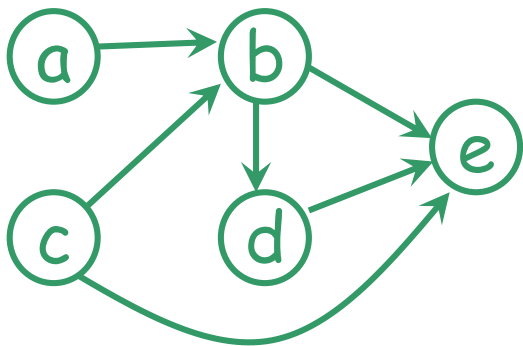
Adjacency matrix / list

Adjacency matrix M for a **directed** graph with n vertices:

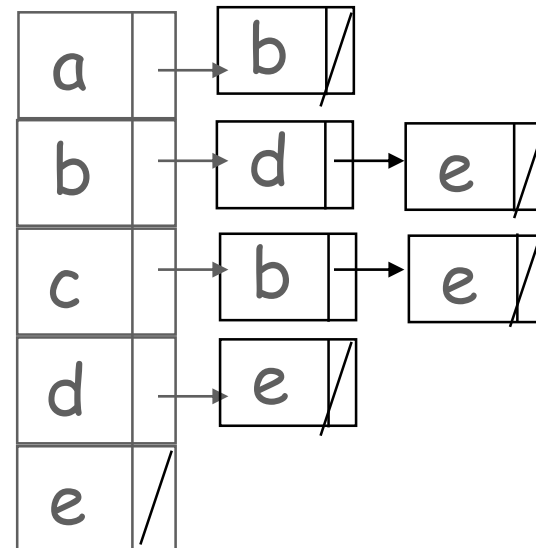
- M is an $n \times n$ matrix
- $M(i, j) = 1$ if (i, j) is an edge

Adjacency list:

- each vertex u has a list of vertices pointed to by an edge leading away from u



	a	b	c	d	e
a	0	1	0	0	0
b	0	0	0	1	1
c	0	1	0	0	1
d	0	0	0	0	1
e	0	0	0	0	0

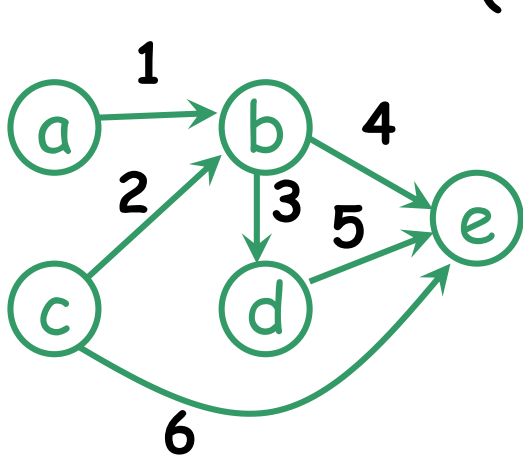


Incidence matrix / list

Incidence matrix M for a directed graph with n vertices and m edges is an $m \times n$ matrix

- $M(i, j) = 1$ if edge i is leading away from vertex j
- $M(i, j) = -1$ if edge i is leading to vertex j

Incidence list: each edge has a list of two vertices (leading away is 1st and leading to is 2nd)

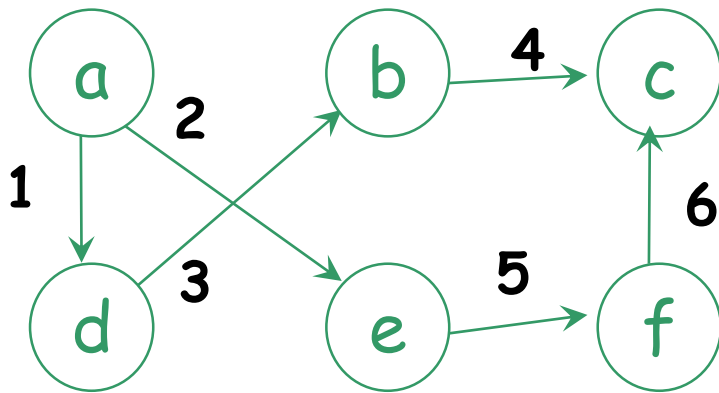


	a	b	c	d	e
1	1	-1	0	0	0
2	0	-1	1	0	0
3	0	1	0	-1	0
4	0	1	0	0	-1
5	0	0	0	1	-1
6	0	0	1	0	-1

1	→	a	→	b	/
2	→	c	→	b	/
3	→	b	→	d	/
4	→	b	→	e	/
5	→	d	→	e	/
6	→	c	→	e	/

Exercise

Give the adjacency matrix and incidence matrix of the following graph



labels of edge
are edge
number

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- ✓ Able to apply BFS and DFS to traverse a graph
- **Able to tell what a tree is**

Tree ...

Trees

An undirected graph $G=(V,E)$ is a tree if G is connected and acyclic (i.e., contains no cycles)

Other equivalent statements:

1. There is exactly one path between any two vertices in G
2. G is connected and removal of one edge disconnects G
3. G is acyclic and adding one edge creates a cycle
4. G is connected and $m=n-1$ (where $|V|=n$, $|E|=m$)

Trees

An undirected graph $G=(V,E)$ is a tree if G is connected and acyclic (i.e., contains no cycles)

Other equivalent statements:

1. There is exactly one path between any two vertices in G
(coz G is connected and acyclic)
2. G is connected and removal of one edge disconnects G
(removal of an edge $\{u,v\}$ disconnects at least u and v because of [1])
3. G is acyclic and adding one edge creates a cycle
(adding an edge $\{u,v\}$ creates one more path between u and v , a cycle is formed)
4. G is connected and $m=n-1$ (where $|V|=n$, $|E|=m$)

Lemma: $P(n)$: If a tree T has n vertices and m edges, then $m=n-1$.

optional, self-study

Proof: By induction on the number of vertices.

Basic step: A tree with single vertex does not have an edge.

Induction step: $P(n-1) \Rightarrow P(n)$ for $n > 1$?

Remove an edge from the tree T . By [2], T becomes disconnected. Two connected components T_1 and T_2 are obtained, neither contains a cycle (the cycle is also present in T otherwise).

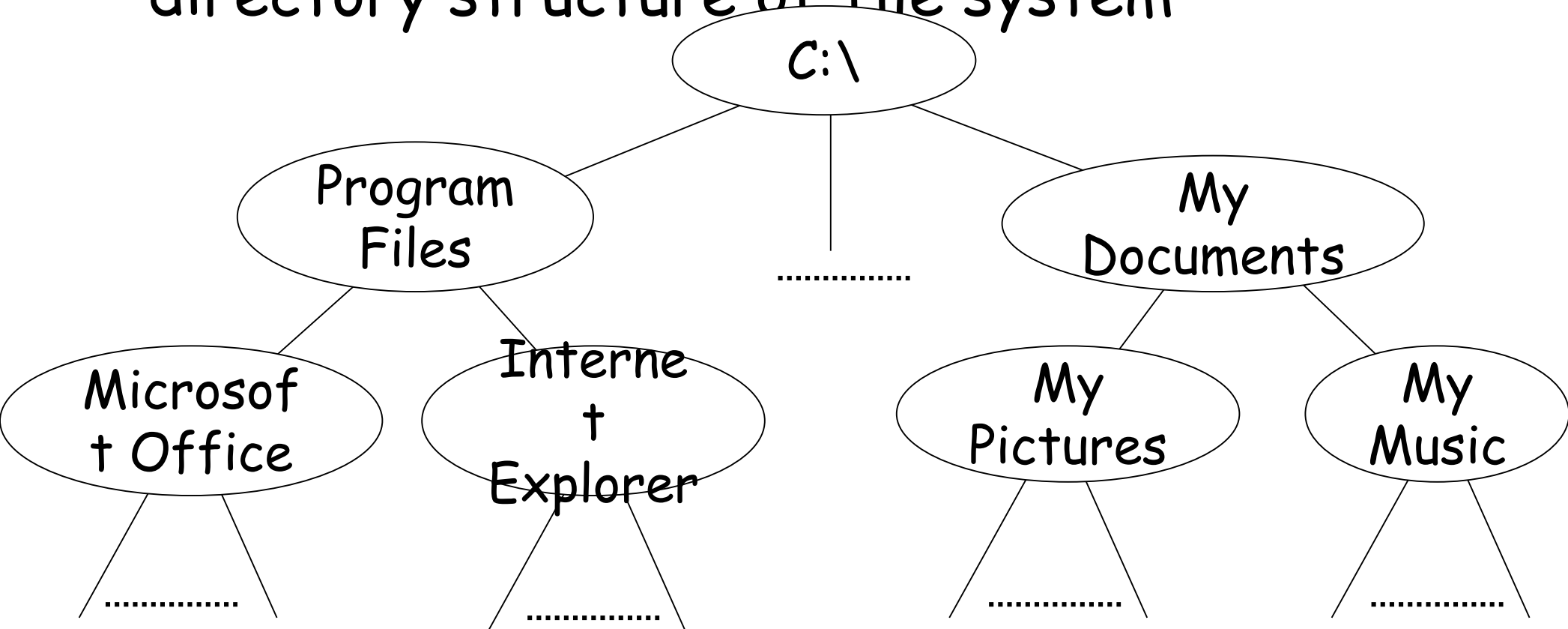
Therefore, both T_1 and T_2 are trees. Let n_1 and n_2 be the number of vertices in T_1 and T_2 . [$n_1+n_2 = n$]

By the induction hypothesis, T_1 and T_2 contains n_1-1 and n_2-1 edges.

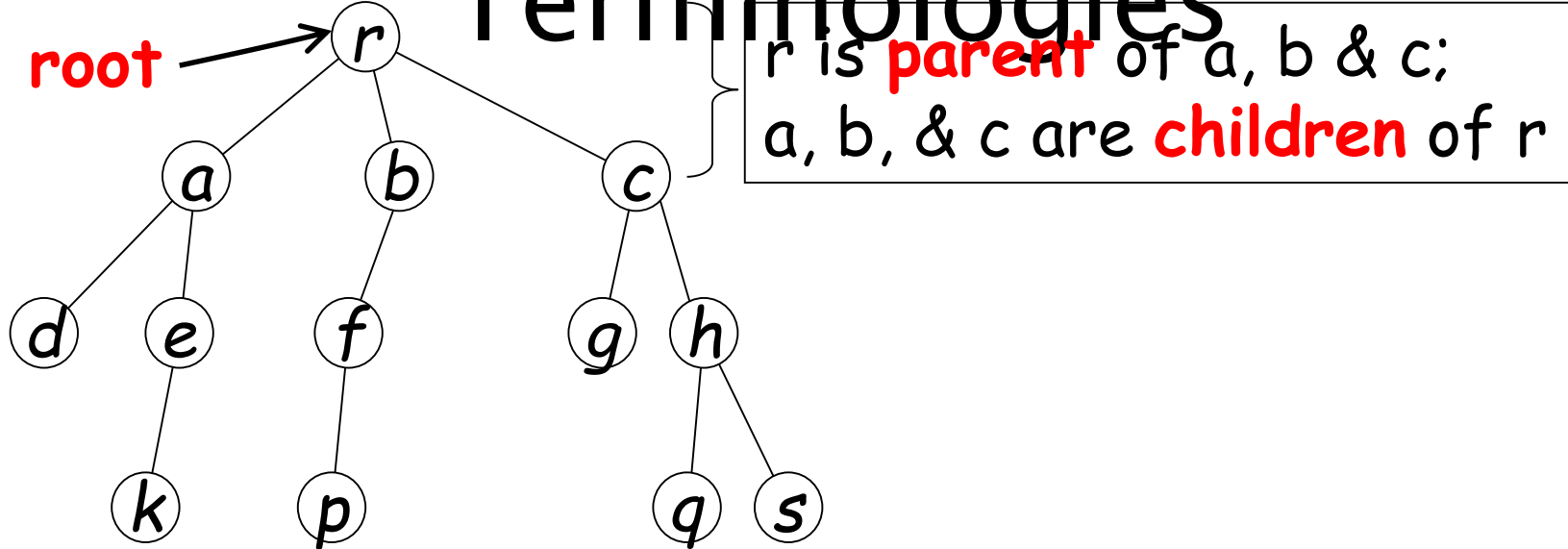
Hence, T contains $(n_1-1) + (n_2-1) + 1 = n-1$ edges.

Rooted trees

Tree with hierarchical structure, e.g.,
directory structure of file system

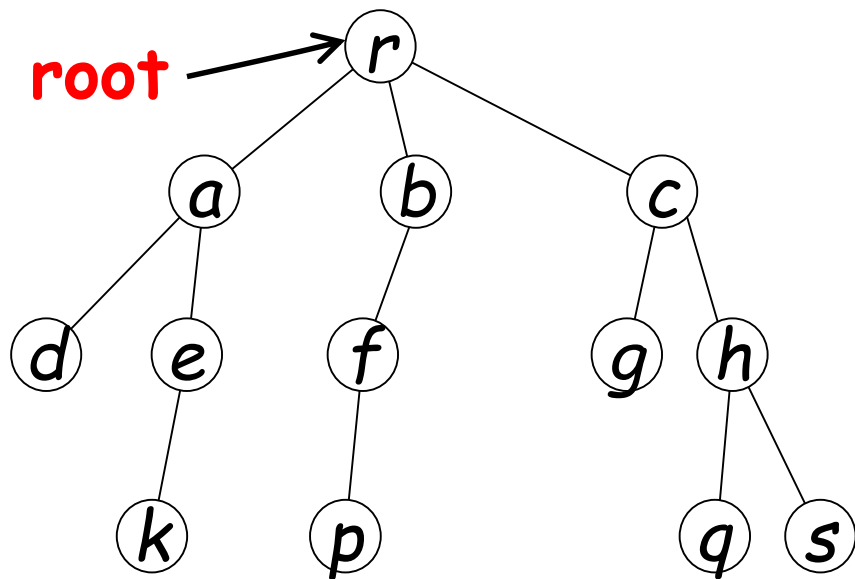


Terminologies



- Topmost vertex is called the root.
- A vertex **u** may have some children directly below it, **u** is called the parent of its children.
- Degree of a **vertex** is the no. of children it has. (N.B. it is different from the degree in an unrooted tree.)
- Degree of a **tree** is the max. degree of all vertices.
- A vertex with no child (degree-0) is called a leaf. All others are called internal vertices.

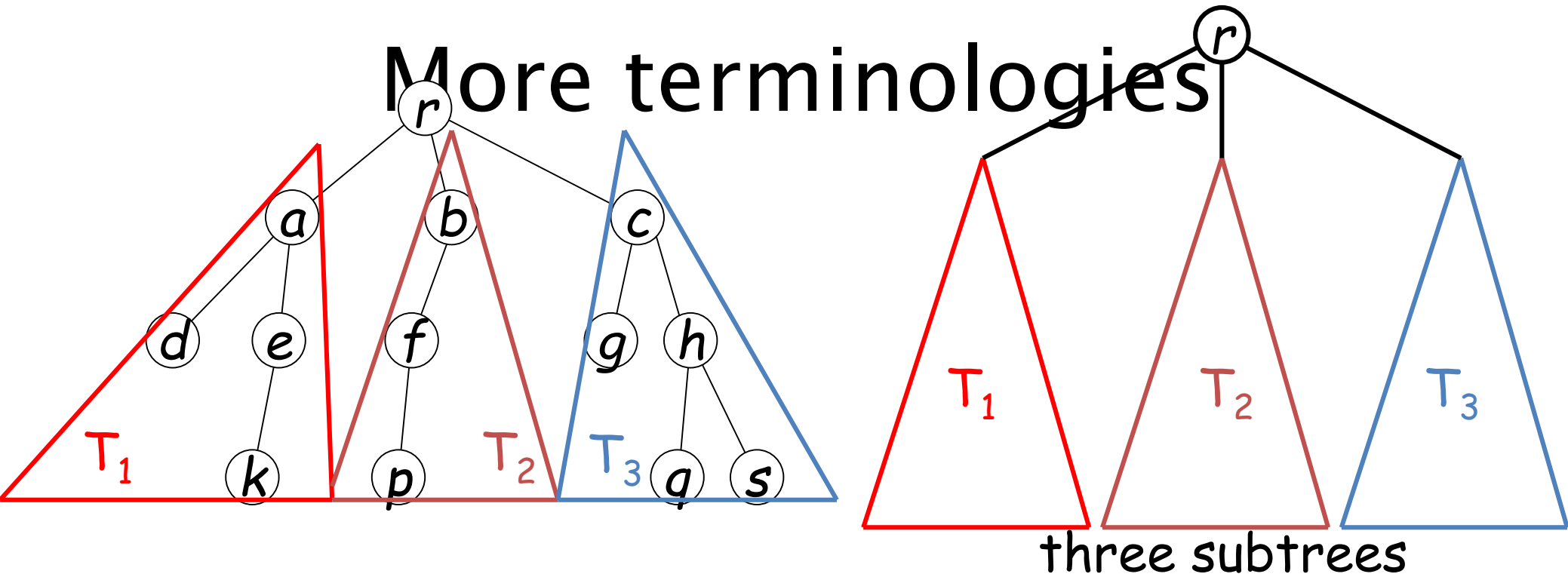
Terminologies



deg-0:
deg-1:
deg-2:
deg-3:

What is the
degree of
this tree?

More terminologies



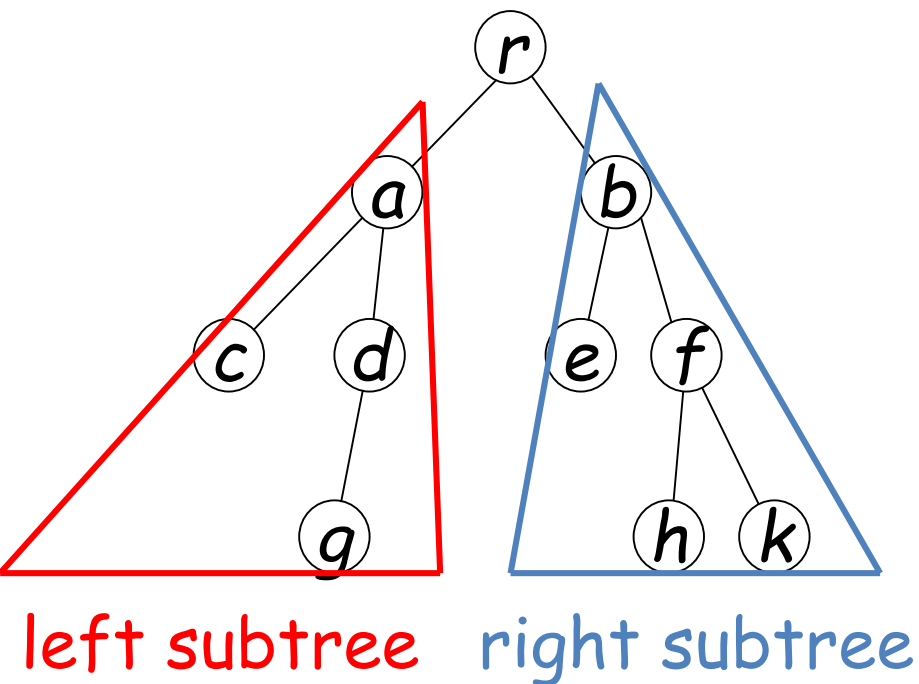
➤ We can define a tree **recursively**

- A single vertex is a tree.
- If T_1, T_2, \dots, T_k are **disjoint** trees with roots r_1, r_2, \dots, r_k , the graph obtained by attaching a *new vertex* r to each of r_1, r_2, \dots, r_k with a single edge forms a tree T with root r .
- T_1, T_2, \dots, T_k are called **subtrees** of T .

which are the roots of the subtrees?

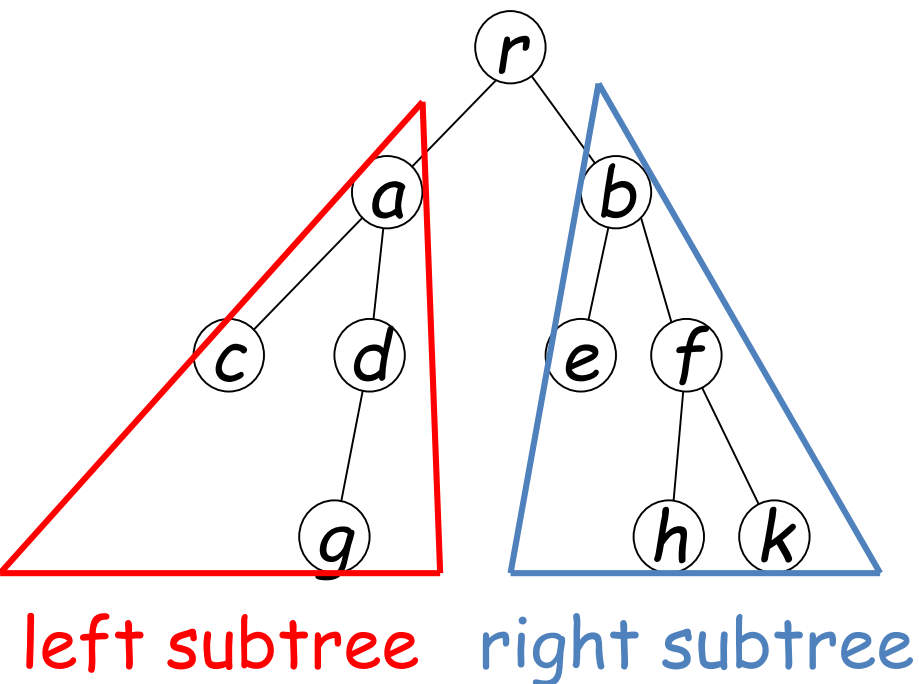
Binary tree

- a tree of degree at most TWO
- the two subtrees are called left subtree and right subtree (may be empty)



Binary tree

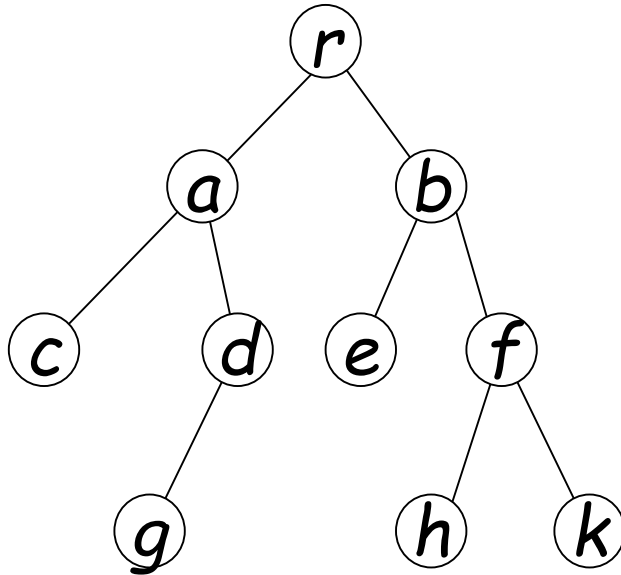
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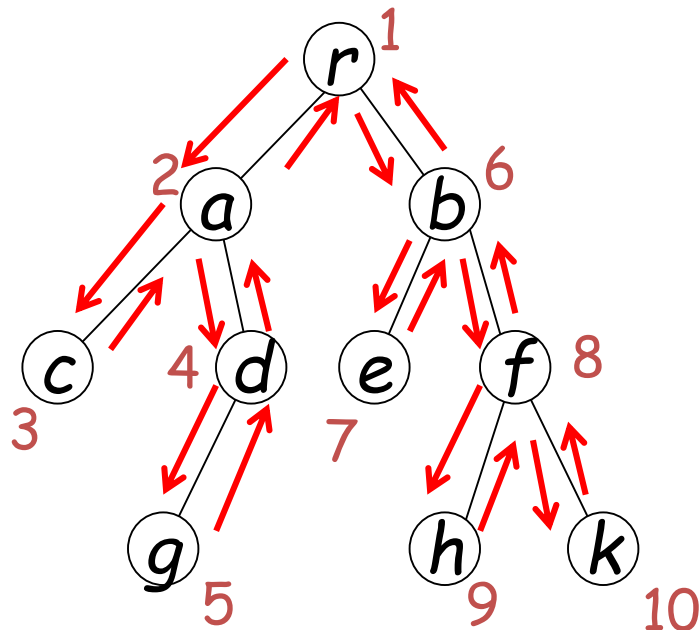
There are *three* common ways to traverse a binary tree:

- **preorder** traversal - vertex, left subtree, right subtree
- **inorder** traversal - left subtree, vertex, right subtree
- **postorder** traversal - left subtree, right subtree, vertex

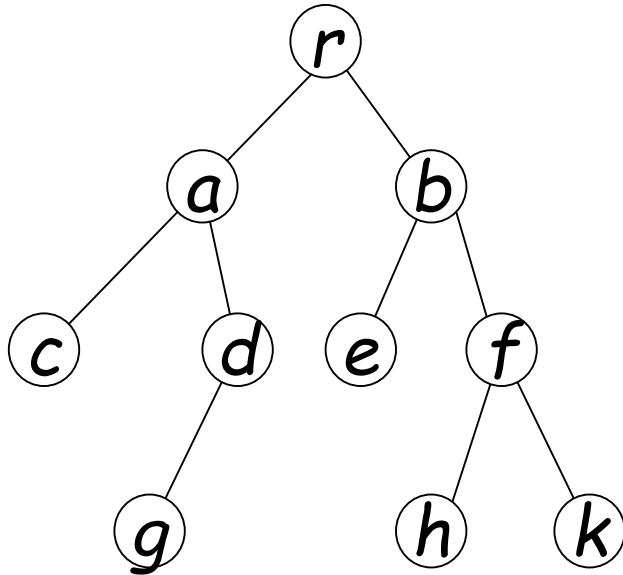
Traversing a binary tree



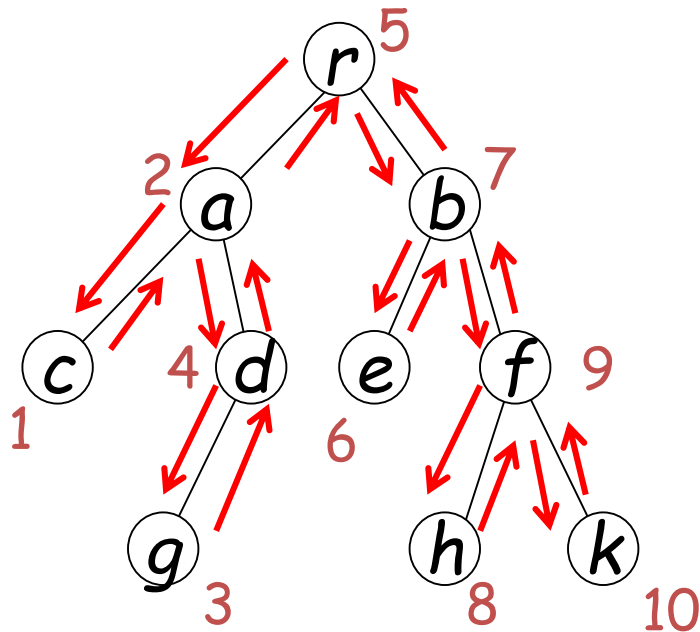
preorder traversal
- vertex, left subtree, right subtree



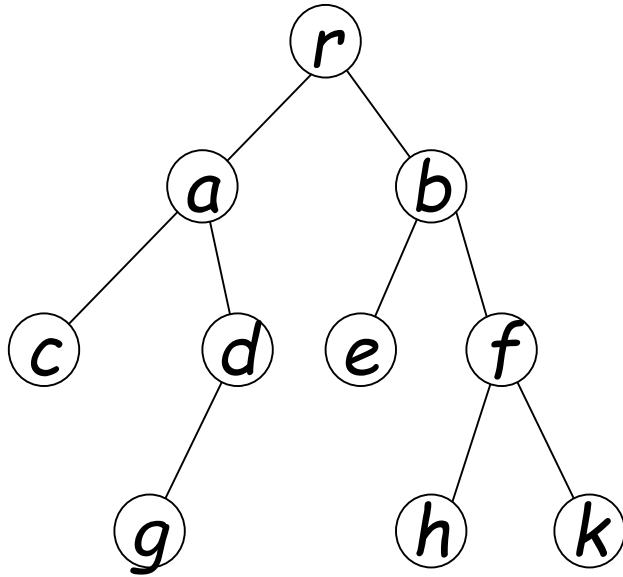
Traversing a binary tree



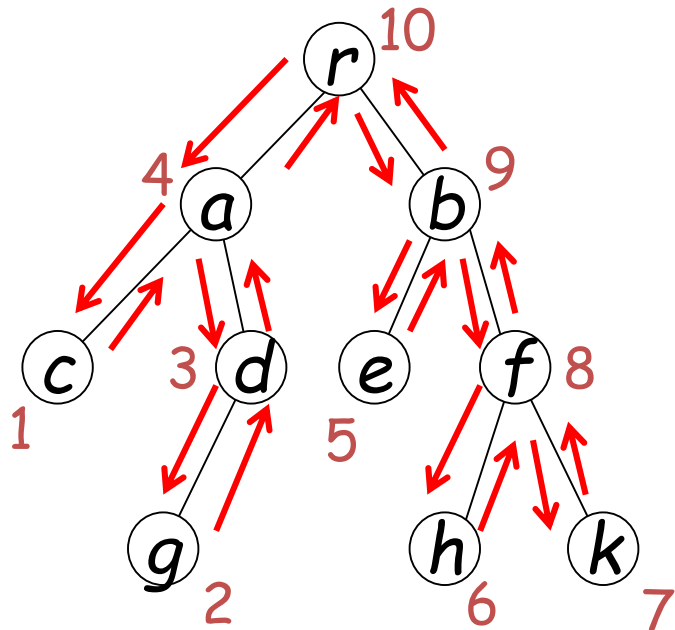
inorder traversal
- left subtree, vertex, right subtree



Traversing a binary tree

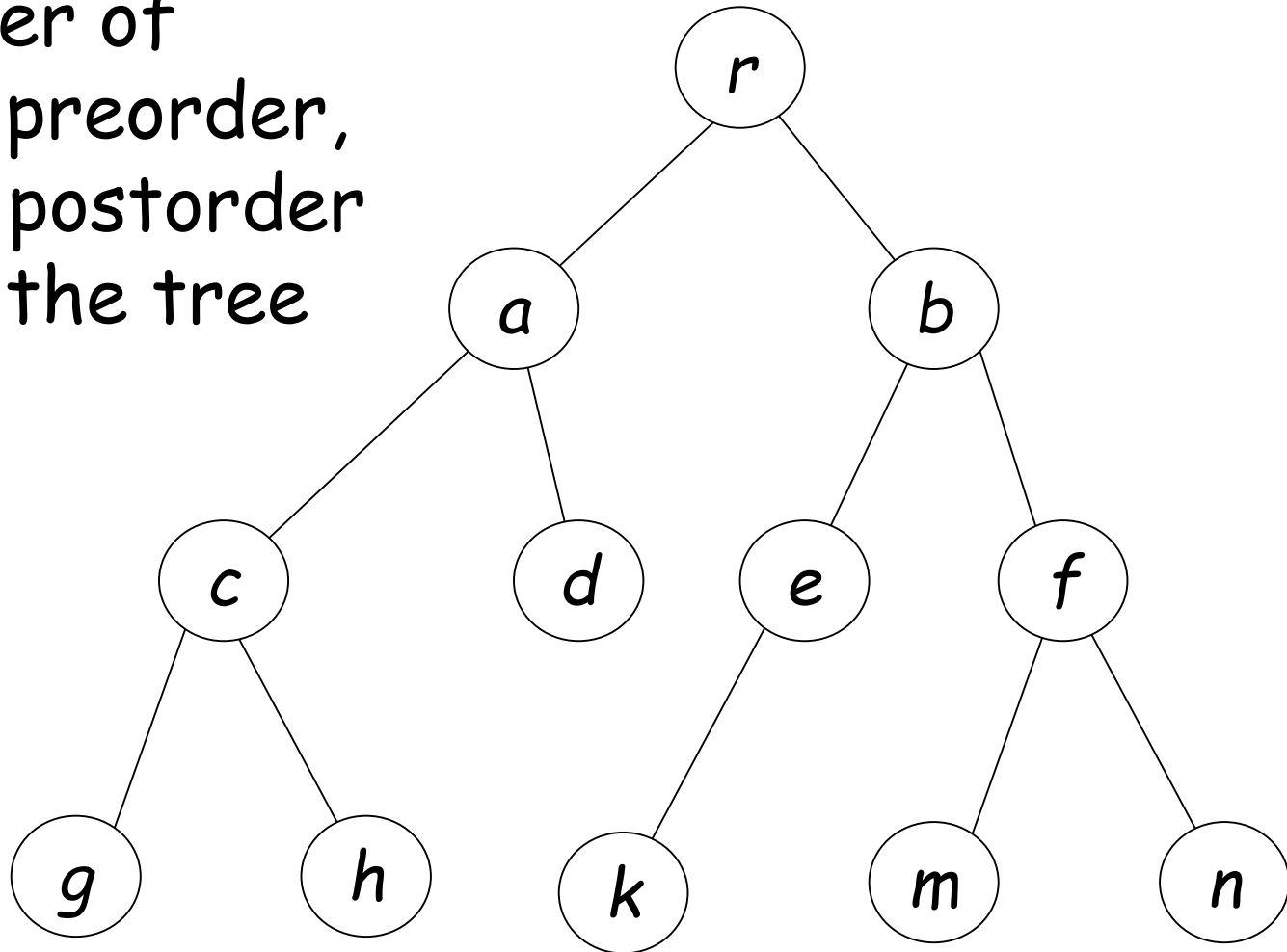


postorder traversal
- left subtree, right subtree, vertex



Example

Give the order of traversal of preorder, inorder, and postorder traversal of the tree



preorder:
inorder:
postorder:

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