# Gajendra Processor

Lab Report

submitted by

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Under the supervision

 $\mathbf{of}$ 

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# 1 General CPU Register

## 1.1 Diagram

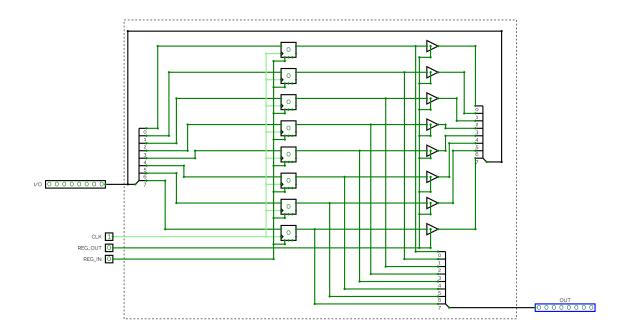


Figure 1: 8-bit Register

# 1.2 I/O Control

It takes 8-bit input from the common bus. Which is stored in it. The 8-bit value stored in it is given as output. REG\_in and REGout are used as the control inputs for controlling the input and output.

# 1.3 Working

If REG\_in is 1 the input is read from the common bus and when REG\_out is 1 the output is written to the common bus.

# 2 Program Counter

# 2.1 Diagram

The ROM consists of both the instructions and data for the program that has to be run , The program counter points to the address of the instruction in ROM , that has to be processed

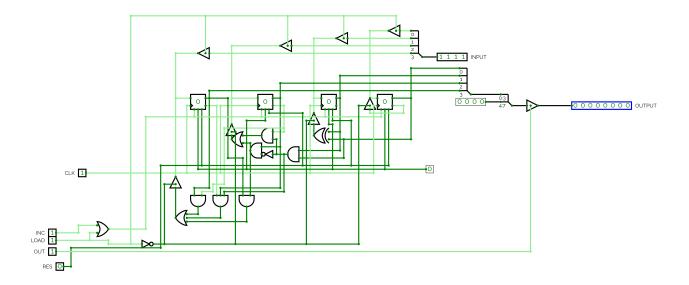


Figure 2: Program Counter

next . At the starting of the program the program counter is reset to 0000 , It then sends the address that is stored in it (i.e 0000 in this case) to the Memory Address Register , The program counter is then incremented by 1 (i.e to 0001 in this case) , After the first instruction is executed then the program counter sends the current address stored in it (i.e 0001 in this case) , this cycle keeps on repeating until the program ends.

The following are the functionalities that are present in the Program Counter.

## 2.2 PC\_INC

If  $PC \perp INC = 1$  then the value of the address stored in the Program Counter is incremented by 1 in each positive clock triggering.

#### 2.3 PC\_OUT

If  $PC\_OUT = 1$  then the Program Counter sends the address stored in it into the Least Significant Bits (0-3) of the **common bus** of the **Gajendra Processor** through **output**.

## 2.4 PC LOAD

If PC\_LOAD = 1 then the Program Counter takes the Least Significant Bits (0-3) of the **common bus** into it through the **input** and stores it as the next instruction.

# 3 MEMORY

## 3.1 Diagram

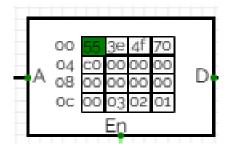


Figure 3: Memory

## 3.2 Diagram

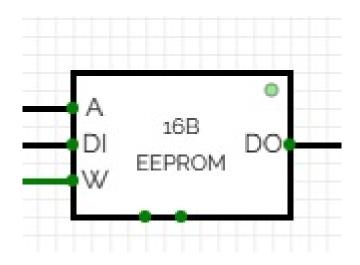


Figure 4: EEPROM

Memory is a ROM which contains both instruction and data required for a program. The following are the functionalities that are present in the Memory. We are using EEPROM here for STA instruction .

#### 3.3 MEM\_OUT

If MEM\_OUT = 1 then the memory releases a 8-bit output into the common bus.

# 4 MEMORY ADDRESS REGISTER

During a computer run , the 4-bit address in the Programming Counter is latched into the MAR . Then the MAR applies this 4-bit address to the Memory where a read operation is performed. The following are the functionalities that are present in the MAR .

## 4.1 MAR JN

If MAR  $_{\rm IN}=1$ , then the MAR takes the 4-bit address (Least Significant in common bus) from the Program Counter through common bus.

# 5 INSTRUCTION REGISTER

## 5.1 Diagram

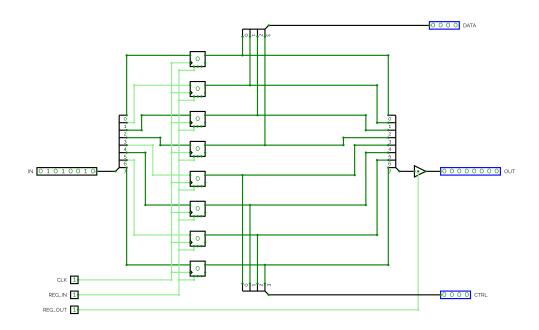


Figure 5: Instruction Register

To fetch an instruction from the memory the computer does a memory read operation . This places the contents of the addressed memory location on the common bus. The instruction register loads this content into it on the next positive clock edge .

The contents of the instruction register are split into two nibbles of 4-bits each (least 4 significant bits , most 4-significant bits), The most 4-significant bits are sent into **Controller** and it will also send the 8-bit content into the **Instruction decoder**.

The following are the functionalities that are present in the Instruction register.

#### 5.2 INS\_REG\_IN

If INS\_REG \_IN = 1 , then the INS register takes the 8-bit content as input from the common bus .

#### 5.3 INS\_REG\_OUT

If INS\_REG  $\_OUT = 1$ , then INS register loads its 8-bit content into Instruction decoder which further enter into the common bus .

## 5.4 CTRL

The Most Significant 4-bits(4-7) automatically goes to the **Controller** and **INS Decoder** as instruction from INS Register.

# 6 INSTRUCTION DECODER

## 6.1 Diagram

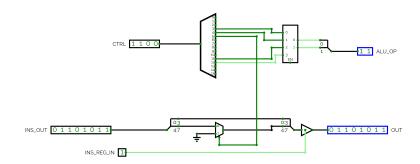


Figure 6: Instruction Decoder

It helps the ALU by giving the instruction whether to add or subtract, by using the most significant 4-bits (i.e CTRL). It also helps us with the instruction **LDI** by making the 4 - most significant bits into 0000, so that the remaining 4-bits gets stored in our **accumulator** (i.e register A).

The following are the functionalities that are present in the Instruction Decoder .

#### 6.2 INS DEC OUT

If INS  $_{-}DEC _{-}OUT = 1$ , then it takes the 8-bit content from OUT of INS register.

#### 6.3 ALU\_OUT

The INS\_DEC generates a 2-bit input for the **ALU**.

It generates the following outputs for the corresponding instructions .

$$00 \implies Addition$$

$$01 \implies Subtraction$$

$$10 \implies LShift$$

$$11 \implies RShift$$

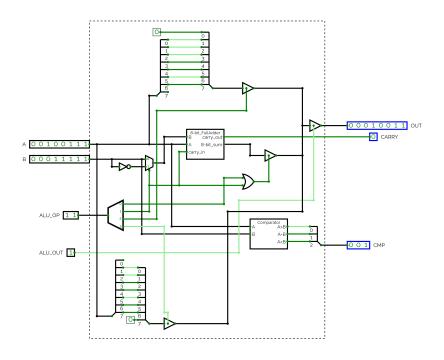


Figure 7: Gajendra ALU

# 7 Arithmetic and Logic Unit(ALU)

The ALU performs the arithmetic operations that are executed by the computer. The operations that are supported are Addition(ADD) , Subtraction(SUB) and SHIFT . It stores the final value in it and it also returns the final 8-bit into Accumulator or REG  $\_A$  .

#### 7.1 ADD

The ADD operation is used to add two 8-bit numbers using a 8-bit full adder which adds without carry .

#### 7.2 SUB

The SUB operation is used to subtract two 8-bit numbers using a 8-bit full adder with an input from the INS\_DEC which triggers sub operation.

#### **7.3** SHIFT

The LSHIFT operation performs a bitwise left shift on the value in register A whereas the RSHIFT operation performs a bitwise right shift on the value in register A.

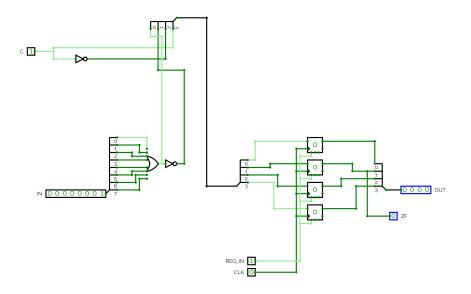


Figure 8: Status Register

# 8 Status Register

# 8.1 Diagram

# 8.2 I/O

It takes 2 inputs ALU \_OUT and Carry from ALU and gives a 4-bit output in which the **the** most significant bit is C(1 if carry = 1) and the next bit gives  $NC(\overline{C})$  and the next bit gives  $Z(1 \text{ if}(ALU \_OUT = 1))$  and the least significant gives  $Z(\overline{Z})$ .

## 8.3 Uses

It is used to implement JUMP and JUMP if zero conditions .

## 9 Control Processor

## 9.1 Diagram

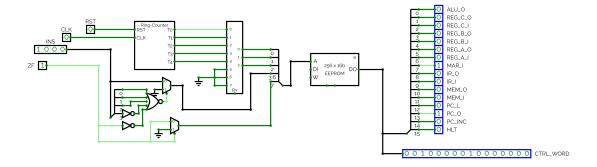


Figure 9: Control Processor

# 9.2 I/O

The control processor takes 4 bit instruction as input and we use a ring counter of length 5 indicating the 5 T-states. The EEPROM in the control processor takes 8-bit address in which the first bit is inherently zero and next 4 bits are the instruction and the last 3 bits are the T-state that is to be implemented. The EEPROM gives 16-bit output which is the control word which is also the output of the Control Processor.

# 9.3 Working

This control word controls all the other elements like REG\_in,REG\_out,PC\_load, PC\_out, PC\_inc, flag, MEM\_in, MEM\_out, MAR\_in, INS\_out, INS\_out, ALU\_out The control processor is the most important component of our processor as it controls all the other components . When Ins is  $1000 \, (JNZ)$  and ZF = 1, this instruction changes to NOP (0000) and the MSB of address entering into EEPROM will also be 0 (i.e the ZF will enter A as 0). In this way other half of the EEPROM need not be wasted and optimized.

# 9.4 Ring Counter

Every instruction has some micro instructions called T-states that are to be executed in order one at a time , these control words are generated by accessing the memory locations in the EEPROM , the T-states are governed by the ring counter .

## 9.5 EEPROM Programming

The EEPROM is loaded in such a way that every instruction has 8 memory locations in the EEPROM for its micro instructions. The each memory locations contains the control word of that particular micro instruction.

# 10 Microinstructions and Controller Logic Design

## 10.1 NOP:

```
NOP: 1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

0 T2

0 T3

0 T4
```

Figure 10: Micro Instruction for NOP

#### 10.2 LDA:

```
LDA :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<MAR_IN T2

1<<MEM_OUT|1<<REGA_IN T3

0 T4
```

Figure 11: Micro instruction for LDA

#### 10.3 ADD:

```
ADD :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<MAR_IN T2

1<<MEM_IN|1<<REGA_OUT T3

0 T4
```

Figure 12: Micro instruction for ADD

# 10.4 SUB:

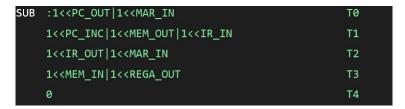


Figure 13: Micro instruction for SUB

# 10.5 LDI:

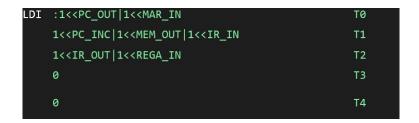


Figure 14: Micro instruction for LDI

## 10.6 JMP:



Figure 15: Micro instruction for JMP

## 10.7 SWAP:

```
      SWAP :1<<PC_OUT|1<<MAR_IN</td>
      T0

      1<<PC_INC|1<<MEM_OUT|1<<IR_IN</td>
      T1

      1<<REGA_OUT|1<<REGB_IN</td>
      T2

      1<<REGC_OUT|1<<REGA_IN</td>
      T3

      1<<REGB_OUT|1<<REGC_IN</td>
      T4
```

Figure 16: Micro instruction for SWAP

## 10.8 JNZ:

```
JNZ :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<PC_LOAD T2

0 T3

0 T4
```

Figure 17: Micro instruction for JNZ

## 10.9 MOVAB:

```
MOVAB:1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<REGA_OUT|1<<REGB_IN T2

0 T3

0 T4
```

Figure 18: Micro instruction for MOVAB

## 10.10 MOVBA:

```
MOVBA:1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<REGB_OUT|1<<REGA_IN T2

0 T3

0 T4
```

Figure 19: Micro instruction for MOVBA

## 10.11 LSHIFT:

```
      SHIFT:1<<PC_OUT|1<<MAR_IN</td>
      T0

      1<<PC_INC|1<<MEM_OUT|1<<IR_IN</td>
      T1

      1<<ALU_OUT|1<<REGA_IN</td>
      T2

      0
      T3

      0
      T4
```

Figure 20: Micro instruction for LSHIFT

# 10.12 RSHIFT:

```
      SHIFT:1<<PC_OUT|1<<MAR_IN</td>
      T0

      1<<PC_INC|1<<MEM_OUT|1<<IR_IN</td>
      T1

      1<<ALU_OUT|1<<REGA_IN</td>
      T2

      0
      T3

      0
      T4
```

Figure 21: Micro instruction for RSHIFT

## 10.13 OUT:

```
OUT :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<REGA_OUT|1<<REGC_IN T2

0 T3

0 T4
```

Figure 22: Micro instruction for OUT  $\,$ 

# 10.14 HLT:

HLT	:1< <pc_out 1<<mar_in< th=""><th>Т0</th></pc_out 1<<mar_in<>	Т0
	1< <pc_inc 1<<mem_out 1<<ir_in< th=""><th>T1</th></pc_inc 1<<mem_out 1<<ir_in<>	T1
	1< <hlt< th=""><th>T2</th></hlt<>	T2
1	0	Т3
1	0	T4

Figure 23: Micro instruction for HLT

# 11 Instruction Set

## 11.1 NOP:

## Description

Loads the Instruction Register and increments the Program Counter.

Operation:

(i) (i)None

Syntax: Operands: Program Counter:

(i) NOP X None  $PC \leftarrow PC + 1$ 

8-bit op-code:

0000	VVVV
0000	^^^

Figure 24: Instruction Set for NOP

#### 11.2 LDA:

# Description

Loads the value stored in the address into the accumulator.

Operation:

(i) (i) $R_A \leftarrow *X$ 

Syntax: Operands: Program Counter:

(i) LDA X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

0001	
0001	XXXX

Figure 25: Instruction Set for LDA

## 11.3 ADD:

## Description

Stores the value in the given address in register B and adds that value to the accumulator.

Operation:

- (i) (i)  $R_B \leftarrow *X$ 
  - (ii)  $R_A \leftarrow R_A + R_B$

Syntax: Operands: Program Counter:

(i) ADD X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

0011 xxxx

Figure 26: Instruction Set for ADD

## 11.4 SUB:

## Description

Stores the value in the given address in register B and subtracts it from the value in the accumulator.

Operation:

- (i) (i) $R_B \leftarrow *X$ 
  - (ii)  $R_A \leftarrow R_A R_B$

Syntax: Operands: Program Counter:

(i) SUB X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

0100	
0100	XXXX

Figure 27: Instruction Set for SUB

## 11.5 LDI:

# Description

Loads the given value into the accumulator.

Operation:

(i) (i) $R_A \leftarrow X$ 

Syntax: Operands: Program Counter: LDI X  $0 \le x \le 15$   $PC \leftarrow PC + 1$ 

8-bit op-code:

(i)

0101 xxxx

Figure 28: Instruction Set for LDI

#### 11.6 JMP:

## Description

Jumps to the given address unconditionally.

Operation:

(i) (i)PC  $\leftarrow$  k

Syntax: Operands: Program Counter:

(i) JMP k  $0 \le k \le 15$  PC  $\leftarrow$  k

8-bit op-code:

0110 kkkk

Figure 29: Instruction Set for JMP

## 11.7 SWAP:

## Description

Swaps the values in registers A and C using the register B.

Operation:

- (i)  $(i)R_B \leftarrow R_A$ 
  - (ii)  $R_A \leftarrow R_C$
  - (iii)  $R_C \leftarrow R_B$

Syntax: Operands: Program Counter:

(i) SWAP X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

0111 xxxx

Figure 30: Instruction for SWAP

#### 11.8 JNZ:

## Description

Jumps to the given address only if the NZ flag is 0 else it acts as NOP.

Operation:

(i) (i)PC  $\leftarrow$  k / PC  $\leftarrow$  PC + 1

Syntax: Operands: Program Counter:  $0 \le x \le 15 \qquad PC \leftarrow k / PC \leftarrow PC + 1$ 

8-bit op-code:

1000 xxxx

Figure 31: Instruction Set for JNZ

#### 11.9 MOVAB:

## Description

Copies the value in register B to register A and retains the value in register B.

Operation:

(i) (i) $R_A \leftarrow R_B$ 

Syntax: Operands: Program Counter:

(i) MOVAB X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

1001 xxxx

Figure 32: Instruction Set for MOVAB

## 11.10 MOVBA:

# Description

Copies the value in register B to register A and retains the value in register B.

Operation:

(i) (i) $R_B \leftarrow R_A$ 

Syntax: Operands: Program Counter:

(i) MOVBA X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

1010	xxxx
------	------

Figure 33: Instruction Set for MOVBA

## 11.11 LSHIFT:

## Description

Bitwise shifts the value in register A by one bit to the left, discarding the leftmost bit and introducing a zero at the rightmost bit.

Operation:

(i) (i)  $R_A \ll 1$ 

Syntax: Operands: Program Counter:

(i) LSHIFT X

 $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

1011	xxxx
------	------

Figure 34: Instruction Set for LSHIFT

## 11.12 RSHIFT:

## Description

Bitwise shifts the value in register A by one bit to the right, discarding the rightmost bit and introducing a zero at the leftmost bit.

Operation:

(i) (i)  $R_A >> 1$ 

Syntax: Operands: Program Counter:

(i) RSHIFT X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

⊥ 11∩∩	VVVV
1100	XXXX
1100	

Figure 35: Instruction Set for RSHIFT

## 11.13 OUT:

## Description

Copies the value in Register A to Register C and the value is displayed on the screen.

Operation:

(i) (i) $R_A \leftarrow X$ 

Syntax: Operands: Program Counter:

(i) OUT X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

1101 xxxx

Figure 36: Instruction Set for OUT

#### 11.14 HLT:

## Description

Terminates the program i.e, no further instruction is executed unless the system is restarted.

Operation:

(i) (i) $R_A \leftarrow X$ 

Syntax: Operands: Program Counter:

(i) HLT X  $0 \le x \le 15$  PC  $\leftarrow$  PC + 1

8-bit op-code:

1110 xxxx

Figure 37: Instruction Set for HLT

# 12 Sample programs

#### 12.1 For Addition and Subtraction

LOAD c ADD f SUB d OUT HALT

0xd : 2 0xc : 30xf : 5

#### 12.2 For L-shift

LDI 6 LSHIFT OUT HALT

## 12.3 For Multiplication

This example is for 3\*5 . (CUSTOM values can be given in the mentioned places). Load for EEPROM for the multiplication of 3 and 5 :

(0x1a,0x3c,0x2e,0x1d,0x4b,0x2d,0x1e,0x81,0xd0,0xe0,0x00,0x01,0x03,0x05,0x00,0x00)0x03,0x05 (i.e c, d can be changed)

LDA a
ADD c
STA e
LDA d
SUB b
STA d
LDA e
JNZ 1 (To 2nd step(i.e ADD c))
OUT
HALT

0xA: 0 0xb: 1 0xc: 3

0xd: updated value of the cycle i=5.

0xe: value of final value after n cycles in loop (i.e 3\*n) is stored.