

Constraint Satisfaction Problems & solutions

CS 220 Data, analytics, TA session

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Recap: CSP definition

- **Components:**

- A set of variables: x_1, x_2, \dots, x_n
- A set of constraints: C_1, C_2, \dots, C_m
- Each variable x_i has a non-empty domain D_i of possible values.
- Each constraint C_i :
 - Involves a subset of variables.
 - Specifies the allowable combinations of values for that subset.

- **State:**

- Defined by an assignment of values to some or all variables.
- An assignment that does not violate any constraints is called a *consistent* or *legal* assignment.
- A *complete assignment* is one in which every variable is assigned a value.

- **Solution:**

- A complete assignment that satisfies all constraints.

Recap: CSP in searching

- **Initial State:**

- The empty assignment, in which all variables are unassigned.

- **Successor Function:**

- A value can be assigned to any unassigned variable.
 - The assignment must not conflict with previously assigned variables.

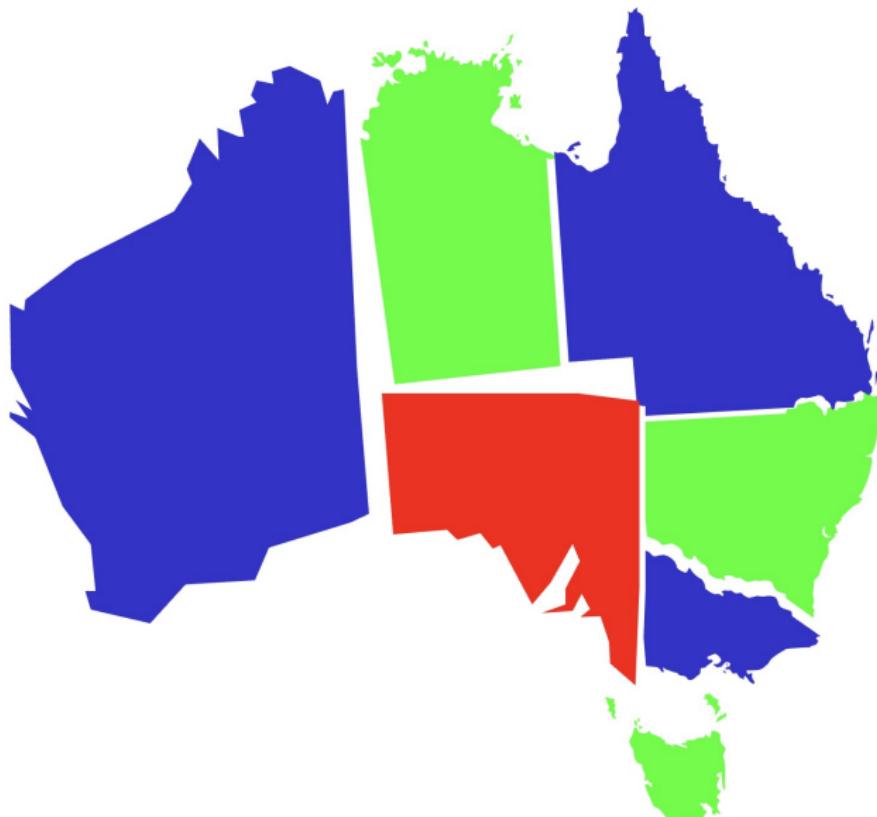
- **Goal Test:**

- The current assignment is complete (all variables are assigned).

- **Path Cost:**

- A constant cost (e.g., 1) is incurred for each step.

Map Coloring Problem as a CSP



Map Coloring Problem as a CSP

- **Problem Description:**

- Given a map consisting of multiple regions, some of which are adjacent.
- Each region must be assigned a color.
- Adjacent regions cannot share the same color.

- **CSP Representation:**

- **Variables:** Each region is a variable x_i , representing the color of the region.
- **Domains:** Each variable x_i has a domain D_i of possible colors (e.g., {red, blue, green}).
- **Constraints:** For any two adjacent regions x_i and x_j , $x_i \neq x_j$.

- **Solution:**

- A valid assignment of colors to all regions such that no adjacent regions share the same color.

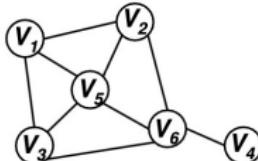
Solutions for CSP

- **Plain Search:** DFS or BFS
 - Use a brute-force approach to evaluate each potential solution.
 - $n!d^n$ possible nodes, but n^d assignments.
- **Backtracking Search:** Systematically explore possible assignments and backtrack when constraints are violated.
 - Basic strategy for solving CSPs.
 - Prunes the search space by abandoning paths that lead to invalid solutions.
- **Heuristic Search:** Improve efficiency by prioritizing the search order.
 - Use heuristics such as Minimum Remaining Values (MRV) or Least Constraining Value (LCV).

Backtracking

Backtracking DFS

V_1	V_2	V_3	V_4	V_5	V_6
?	?	?	?	?	?



V_1	V_2	V_3	V_4	V_5	V_6
B	?	?	?	?	?

Order of values:
 (B, R, G)

V_1	V_2	V_3	V_4	V_5	V_6
B	B	?	?	?	?

V_1	V_2	V_3	V_4	V_5	V_6
B	R	?	?	?	?

Don't even consider
that branch because
 $V_2 = B$ is inconsistent
with the parent state



V_1	V_2	V_3	V_4	V_5	V_6
B	R	R	B	?	?

Backtrack to the
previous state
because no valid
assignment can
be found for V_6

V_1	V_2	V_3	V_4	V_5	V_6
B	R	R	B	G	?

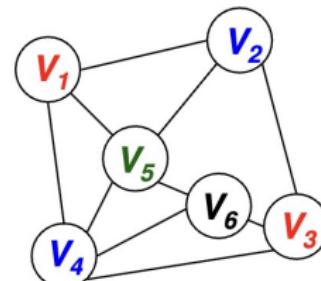
Backtracking

- Uninformed search.
- Can be improved by forward checking or consistency.

Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

	V_1	V_2	V_3	V_4	V_5	V_6
R	O		O			X
B		O		O		X
G				O	X	



There are no valid assignments left
for V_6 we need to backtrack

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Arc consistency (AC-3)

Goal: Ensure each value of one variable has a consistent value in related variables.

Steps:

- ① **Initialize:** Add all arcs (X, Y) to a queue Q .
- ② **Process Arcs:**
 - While Q is not empty:
 - ① Remove an arc (X, Y) from Q .
 - ② Check if any value in $D(X)$ violates the constraint with $D(Y)$:
Remove $x \in D(X)$ if no $y \in D(Y)$ satisfies $C(X, Y)$.
 - ③ If $D(X)$ changes, re-add all related arcs (Z, X) to Q .
 - ③ **Terminate:** When Q is empty, domains are arc-consistent. If a domain $D(X)$ is empty, this searching has no solutions.

Example with AC-3 Algorithm

Problem:

- **Variables:** X, Y, Z
- **Domains:** $D(X) = \{1, 2, 3\}, D(Y) = \{2, 3\}, D(Z) = \{1, 2, 3\}$
- **Constraints:** $X < Y, Y \neq Z$

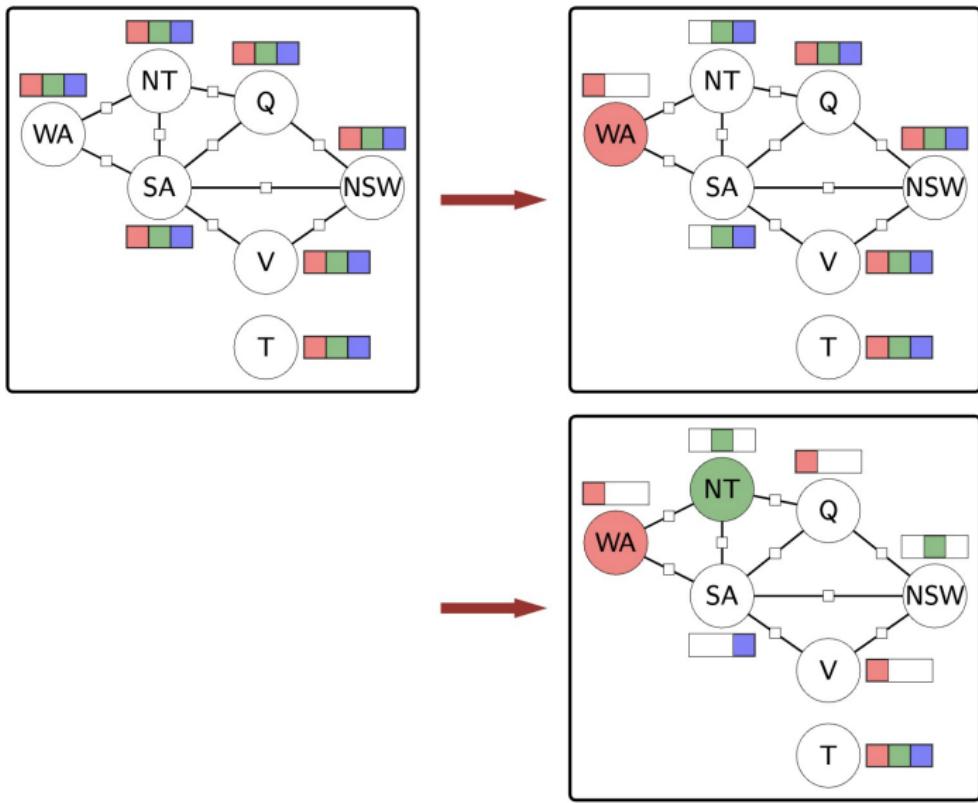
AC-3 Steps:

- ① Initialize $Q = \{(X, Y), (Y, X), (Y, Z), (Z, Y), (X, Z), (Z, X)\}$
- ② Process (X, Y) : Remove 3 from $D(X)$ ($X < Y$ fails for 3).
- ③ Update $D(X) = \{1, 2\}$, recheck related arcs.
- ④ Continue until $Q = \emptyset$.

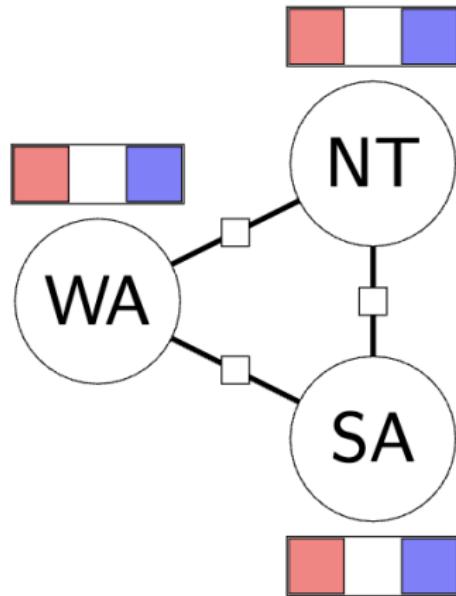
Result:

$$D(X) = \{1, 2\}, \quad D(Y) = \{2, 3\}, \quad D(Z) = \{1, 2, 3\}$$

AC-3 example: earlier than forward checking



AC-3 example



K-consistency: general formulation

- **Node consistency:** single value constraints.
- **Arch consistency:** 2-consistency.
- **Path consistency:** compatibility of three-variables.
- **K-consistency:** high-order consistency.

Variable Selection Heuristics:

- **Minimum Remaining Value (MRV):**

- *Idea:* Choose the variable with the fewest legal values remaining.
 - *Reason:* Reduces conflict potential and reveals unsolvable branches earlier.

- **Degree Heuristic:**

- *Idea:* When MRV ties, choose the variable with the most constraints on other unassigned variables.
 - *Reason:* Solves critical variables earlier, reducing complexity for others.

Value Selection Heuristics:

- **Least Constraining Value (LCV):**

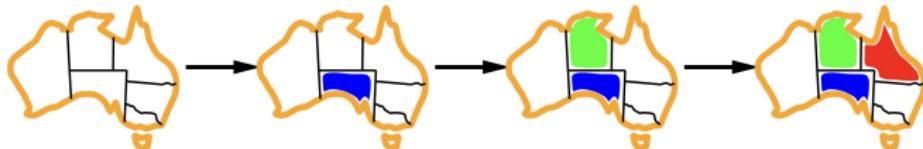
- *Idea:* Select the value that imposes the fewest restrictions on the legal values of other variables.
 - *Reason:* Minimizes future conflicts, increasing the likelihood of successful assignment.

Variable selection heuristics

- Minimum Remaining Values (MRV)



- Degree Heuristic



Value selection heuristics

- Least Constraining Value (LCV)



References

- CS 220 slides, Xin Gao, KAUST
- CS 188 slides, Stuart Russell, UC Berkeley
- 15-381 slides, Martial Hebert and Mike Lewicki, CMU
- CS 221 slides and an online lecture, Percy Liang, Stanford