

# Constraint Satisfaction Problems & solutions

CS 220 Data, analytics, TA session

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# Recap: CSP definition

## • Components:

- A set of variables:  $x_1, x_2, \dots, x_n$
- A set of constraints:  $C_1, C_2, \dots, C_m$
- Each variable  $x_i$  has a non-empty domain  $D_i$  of possible values.
- Each constraint  $C_i$ :
  - Involves a subset of variables.
  - Specifies the allowable combinations of values for that subset.

## • State:

- Defined by an assignment of values to some or all variables.
- An assignment that does not violate any constraints is called a *consistent* or *legal* assignment.
- A *complete assignment* is one in which every variable is assigned a value.

## • Solution:

- A complete assignment that satisfies all constraints.

# Recap: CSP in searching

- **Initial State:**

- The empty assignment, in which all variables are unassigned.

- **Successor Function:**

- A value can be assigned to any unassigned variable.
- The assignment must not conflict with previously assigned variables.

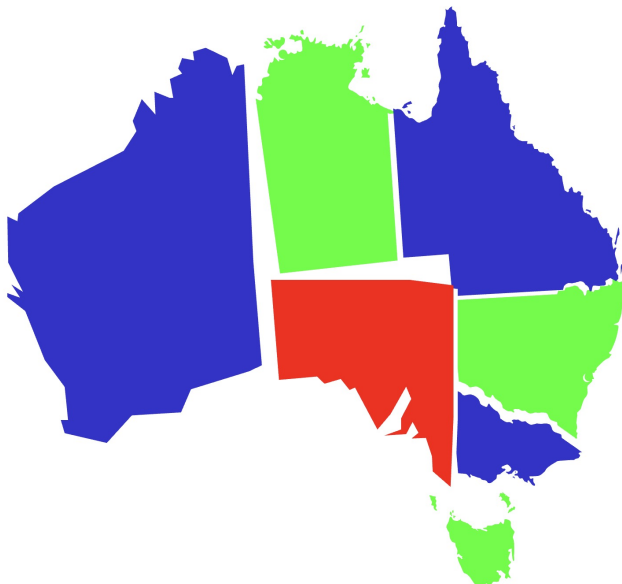
- **Goal Test:**

- The current assignment is complete (all variables are assigned).

- **Path Cost:**

- A constant cost (e.g., 1) is incurred for each step.

# Map Coloring Problem as a CSP



# Map Coloring Problem as a CSP

- **Problem Description:**

- Given a map consisting of multiple regions, some of which are adjacent.
- Each region must be assigned a color.
- Adjacent regions cannot share the same color.

- **CSP Representation:**

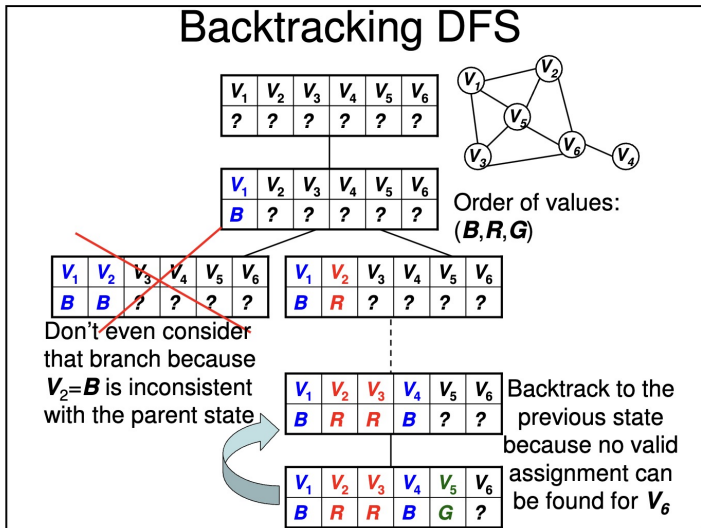
- **Variables:** Each region is a variable  $x_i$ , representing the color of the region.
- **Domains:** Each variable  $x_i$  has a domain  $D_i$  of possible colors (e.g., {red, blue, green}).
- **Constraints:** For any two adjacent regions  $x_i$  and  $x_j$ ,  $x_i \neq x_j$ .

- **Solution:**

- A valid assignment of colors to all regions such that no adjacent regions share the same color.

- **Plain Search:** DFS or BFS
  - Use a brute-force approach to evaluate each potential solution.
  - $n!d^n$  possible nodes, but  $n^d$  assignments.
- **Backtracking Search:** Systematically explore possible assignments and backtrack when constraints are violated.
  - Basic strategy for solving CSPs.
  - Prunes the search space by abandoning paths that lead to invalid solutions.
- **Heuristic Search:** Improve efficiency by prioritizing the search order.
  - Use heuristics such as Minimum Remaining Values (MRV) or Least Constraining Value (LCV).

## Backtracking DFS



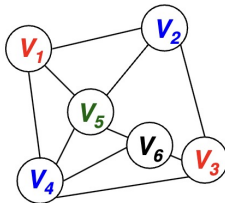
# Backtracking

- Uninformed search.
- Can be improved by forward checking or consistency.

## Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

|     | $V_1$ | $V_2$ | $V_3$ | $V_4$ | $V_5$ | $V_6$ |
|-----|-------|-------|-------|-------|-------|-------|
| $R$ | O     |       | O     |       |       | X     |
| $B$ |       | O     |       | O     |       | X     |
| $G$ |       |       |       |       | O     | X     |



There are no valid assignments left for  $V_6$  we need to backtrack

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# Arc consistency (AC-3)

**Goal:** Ensure each value of one variable has a consistent value in related variables.

## Steps:

- ① **Initialize:** Add all arcs  $(X, Y)$  to a queue  $Q$ .
- ② **Process Arcs:**
  - While  $Q$  is not empty:
    - ① Remove an arc  $(X, Y)$  from  $Q$ .
    - ② Check if any value in  $D(X)$  violates the constraint with  $D(Y)$ :  
$$\text{Remove } x \in D(X) \text{ if no } y \in D(Y) \text{ satisfies } C(X, Y).$$
  - ③ If  $D(X)$  changes, re-add all related arcs  $(Z, X)$  to  $Q$ .
- ③ **Terminate:** When  $Q$  is empty, domains are arc-consistent. If a domain  $D(X)$  is empty, this searching has no solutions.

# Example with AC-3 Algorithm

## Problem:

- **Variables:**  $X, Y, Z$
- **Domains:**  $D(X) = \{1, 2, 3\}, D(Y) = \{2, 3\}, D(Z) = \{1, 2, 3\}$
- **Constraints:**  $X < Y, Y \neq Z$

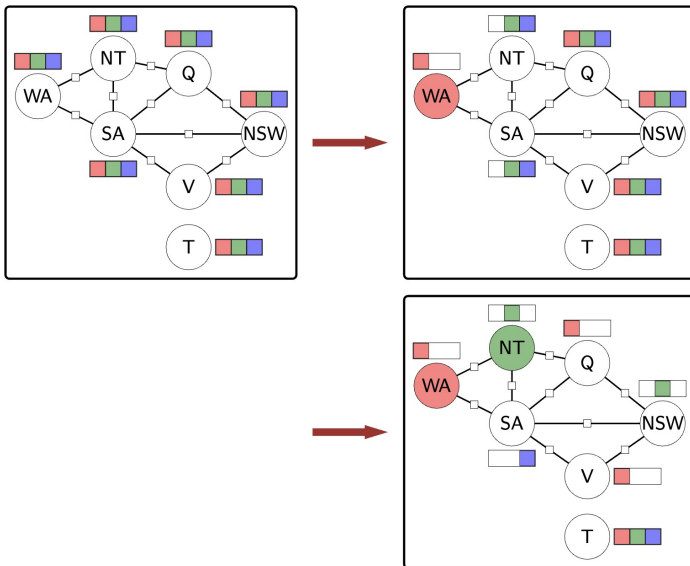
## AC-3 Steps:

- 1 Initialize  $Q = \{(X, Y), (Y, X), (Y, Z), (Z, Y), (X, Z), (Z, X)\}$
- 2 Process  $(X, Y)$ : Remove 3 from  $D(X)$  ( $X < Y$  fails for 3).
- 3 Update  $D(X) = \{1, 2\}$ , recheck related arcs.
- 4 Continue until  $Q = \emptyset$ .

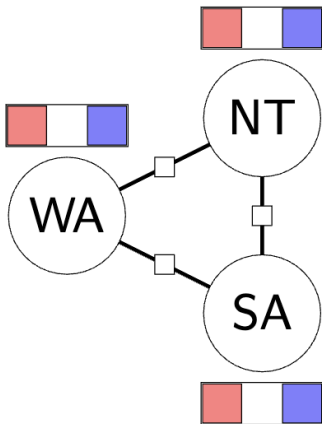
## Result:

$$D(X) = \{1, 2\}, \quad D(Y) = \{2, 3\}, \quad D(Z) = \{1, 2, 3\}$$

# AC-3 example: earlier than forward checking



## AC-3 example



# K-consistency: general formulation

- **Node consistency:** single value constraints.
- **Arch consistency:** 2-consistency.
- **Path consistency:** compatibility of three-variables.
- **K-consistency:** high-order consistency.

## Variable Selection Heuristics:

- **Minimum Remaining Value (MRV):**

- *Idea:* Choose the variable with the fewest legal values remaining.
- *Reason:* Reduces conflict potential and reveals unsolvable branches earlier.

- **Degree Heuristic:**

- *Idea:* When MRV ties, choose the variable with the most constraints on other unassigned variables.
- *Reason:* Solves critical variables earlier, reducing complexity for others.

## Value Selection Heuristics:

- **Least Constraining Value (LCV):**

- *Idea:* Select the value that imposes the fewest restrictions on the legal values of other variables.
- *Reason:* Minimizes future conflicts, increasing the likelihood of successful assignment.

# Variable selection heuristics

- Minimum Remaining Values (MRV)

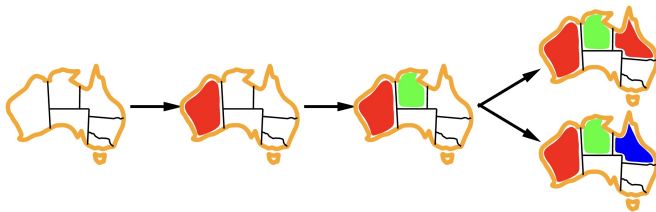


- Degree Heuristic



# Value selection heuristics

- Least Constraining Value (LCV)





# References

- CS 220 slides, Xin Gao, KAUST
- CS 188 slides, Stuart Russell, UC Berkeley
- 15-381 slides, Martial Hebert and Mike Lewicki, CMU
- CS 221 slides and an online lecture, Percy Liang, Stanford