Principles of Programming Languages, 2024.02.02

Important notes

- Total available time: 2h (multichance students do not need to solve Exercise 1).
- You may use any written material you need, and write in English or in Italian.
- You cannot use electronic devices during the exam: every phone must be turned off and kept on your table.
- You cannot use library functions not covered in class in your code.

Exercise 1, Scheme (11 pts)

```
Consider the following data structure, written in Haskell:
data Expr a = Var a | Val Int | Op (Expr a) (Expr a)
instance Functor Expr where
 fmap (Val x) = Val x
 fmap g(Var x) = Var(g x)
 fmap g (Op a b) = Op (fmap g a) (fmap g b)
instance Applicative Expr where
 pure = Var
 <*> Val x =  Val x
 Val x <*> _ = Val x
 Var f < *> Var x = Var (f x)
 Var f < *> Op x y = Op (fmap f x) (fmap f y)
 Op f q < *> x = Op (f < *> x) (q < *> x)
instance Monad Expr where
 Val x >>= = Val x
 Var x >>= f = f x
```

Define an analogous in Scheme, with all the previous operations, where the data structures are encoded as lists – e.g. Op (Val 0) (Var 1) is represented in Scheme as '(Op (Val 0) (Var 1)).

Exercise 2, Haskell (11 pts)

Consider the following datatype definition.

Op a b >>= f = Op (a >>= f) (b >>= f)

```
data F b a = F (b \rightarrow b) a | Null
```

- 1) Make F an instance of Functor, Applicative, and Monad.
- 2) Using an example, show what >>= does in your implementation.

Exercise 3, Erlang (11 pts)

Consider a list L of tasks, where each task is encoded as a function having only one parameter: a PID P. When a task is called, it runs some operations and then sends back the results to P, in this form: $\{result, < Task_PID >, < Result_value >\}$; a task could also fail for some errors.

Define a server which takes L and runs in parallel all the tasks in it, returning the list of the results (the order is not important). Note: in case of failure, every task should be restarted only once; if it fails twice, its result should be represented with the atom bug.

Utilities: you can use from the standard libraries the following functions: *maps:from_list(<List of {Key, Value}>)*, which takes a list of pairs and builds the corresponding map, and *maps:values(<Map>)*, which is its inverse.

Solutions

```
Ex 1
```

```
;; Constructors
(define (var x) (list 'Var x))
(define (val x) (list 'Val x))
(define (op x y) (list 'Op x y))
;; predicates
(define (var? x) (eq? 'Var (car x)))
(define (val? x) (eq? 'Val (car x)))
(define (op? x) (eq? 'Op (car x)))
;; Functor
(define (fmap f e)
 (cond
  ((val? e) e)
  ((var? e)
          (let ((x (cadr e)))
           (var (f x))))
  ((op? e) (let ((x (cadr e))
              (y (caddr e)))
           (op (fmap f x) (fmap f y))))))
;; Applicative
(define pure var)
(define (<*> x y)
 (cond
  ((val? y) y)
((val? x) x)
  ((var? x)
          (let ((f (cadr x)))
           (if (var? y)
(var (f (cadr y)))
             (let ((a (cadr y))
              (b (caddr y)))
(op (fmap f a)(fmap f b))))))
  ((op? x) (let ((a (cadr x))
             (b (caddr x)))
          (op (<*> a y)(<*> b y))))))
;; Monad
(define (>>= x y)
 (cond
  ((val? x) x)
  ((var? x) (y (cadr x)))
  ((op? x) (let ((a (cadr x))
             (b (caddr x)))
          (op (>>= a y)(>>= b y))))))
```

Ex 2

Note: This data structure is basically a combination of State and Maybe, where the State pair is brought out of the function.

```
instance Functor (F x) where
 fmap f(Fgt) = Fg(ft)
 fmap Null = Null
instance Applicative (F x) where
 pure = F id
 Null <*> = Null

_ <*> Null = Null
 \overline{(F f x)} < *> (F g y) = F (f . g) (x y)
instance Monad (F x) where
 Null >>= _= Null
 Ffx >>= q = case q x of
            Null -> Null
            F f' x' -> F (f . f') x'
-- Example
runit (F f x) s = (f s, x)
ex = F (x -> 2*x) >>= x -> pure (x+1)
runit ex 1 -- result: (2,6)
```

Ex 3

```
getres(Fs) ->
   Self = self(),
   process_flag(trap_exit, true),
Pids = maps:from_list([{spawn_link(fun() -> F(Self) end), {F, wait}} || F <- Fs]),
   getres loop(Pids, length(Fs)).
getres_loop(Pids, 0) ->
   [R || {done, R} <- maps:values(Pids)];
getres_loop(Pids, Waiting) ->
   receive
      {result, Pid, R} ->
         getres loop(Pids#{Pid := {done, R}}, Waiting - 1);
      {'EXIT', Pid, Reason} when Reason /= normal ->
#{Pid := {F, Status}} = Pids,
         Self = self(),
         case Status of
            restart ->
               getres_loop(Pids#{Pid := {done, bug}}, Waiting - 1);
              NewPid = spawn_link(fun() -> F(Self) end),
               getres_loop(Pids#{NewPid => {F, restart}}, Waiting)
         end
   end.
```