

1. (a) $T(n) = 3T(\frac{n}{2}) + n^2$

by Master Theorem

$c = \log_b a$ 其中 $a=3, b=2$, 故 $c = \log_2 3$

又 $f(n) = n^2$ 且 $f(n) \in \Omega(n^{\log_2 3 + \epsilon})$, $\epsilon > 0$

$\exists f(\frac{n}{2}) \leq d f(n)$ 其中 $d \in \mathbb{R}^+$ 且 $d < 1$

$T(n) = \Theta(n^2)$

(b) $T(n) = 4T(\frac{n}{2}) + n^2$

by Master Theorem, $c = \log_b a$

其中 $c = \log_2 4 = 2$

$f(n) = n^2$, $f(n) \in \Theta(n^{\log_2 4})$

故 $T(n) \in \Theta(n^2 \log n)$

(c) $T(n) = 16T(\frac{n}{4}) + n$

by Master Theorem, $c = \log_4 16 = 2$

$f(n) = n$, $f(n) \in \Theta(n^{\log_4 16 - \epsilon})$, $\epsilon > 0$

故 $T(n) = \Theta(n^2)$

(d) $T(n) = T(\frac{n}{2}) + 2^n$

by Master Theorem $c = \log_2 1 = 0$

$f(n) = 2^n$, $f(n) \in \Omega(n^{\log_2 1 + \epsilon})$, $a = n \log_2 n$

$\Rightarrow T(n) = \Theta(2^n)$

$\forall n > 0$ 且 n 充分大

$$2. (a) n^2 + 3n + 1 \in O(n^2)$$

在 $n \geq n_0$ 時
若存在 $c, n_0 > 0$, 且 $T(n) \leq c\phi(n)$

$$\Rightarrow n^2 + 3n + 1 \leq cn^2$$

$$\frac{n^2 + 3n + 1}{n^2} \leq c \Rightarrow 1 + \frac{3}{n} + \frac{1}{n^2} \leq c$$

$$c = 2 \text{ 且 } \forall n \geq 4 \text{ 均成立} \quad 1 + \frac{3}{n} + \frac{1}{n^2} \leq 2 \Rightarrow \frac{3}{n} + \frac{1}{n^2} \leq 1$$

$$(b) \log_b n \in O(n^k) \quad \forall b > 1, k > 0$$

if $\exists c > 0, n_0 > 0$ when $n \geq n_0, T(n) \leq c\phi(n)$

$$\Rightarrow \log_b n \leq cn^k \Rightarrow \frac{\log_b n}{n^k} \leq c$$

$$\lim_{n \rightarrow \infty} \frac{\log_b n}{n^k} = 0 \Rightarrow \lim_{n \rightarrow \infty} \frac{1}{\frac{n \log_b n}{k n^{k-1}}} = 0 \quad \text{由洛必达法则} \quad \text{故 } \log_b n \in O(n^k)$$

$$(c) n^k \in O(b^n)$$

if $\exists c > 0, n_0 > 0$ when $n \geq n_0, T(n) \leq c\phi(n)$

$$\Rightarrow n^k \leq cb^n \quad \frac{n^k}{cb^n} \leq 0$$

$$\lim_{n \rightarrow \infty} \frac{n^k}{cb^n} = 0 \Rightarrow \lim_{n \rightarrow \infty} \frac{k n^{k-1}}{c b^n \ln b} = 0 \quad \text{由洛必达法则} \quad \text{故 } n^k \in O(b^n)$$

$$3. \log(n!) \in \Theta(n \log n)$$

$$\log 1 + \dots + \log n \leq (\log n + \dots + \log n) \in O(n \log n)$$

$$\log 1 + \dots + \log n \geq \log \frac{n}{2} + \dots + \log n \quad \text{剩一半}$$

$$\log \frac{n}{2} + \dots + \log n \geq \log \frac{n}{2} + \dots + \log \frac{n}{2} \quad \text{把 } \log n \text{ 换成 } \log \frac{n}{2}$$

7.11)

$$\underbrace{\log \frac{n}{k}}_{\frac{n}{k}} + \underbrace{\log \frac{n}{k}}_{\frac{n}{k}} \Rightarrow \frac{n}{k} \log \frac{n}{k} = \frac{n}{k} (\log n - \log k) \in \Omega(n \log n)$$

$$\frac{n}{k} \log n \in \Theta(n \log n)$$

其中 $k > 1$ 是 constant