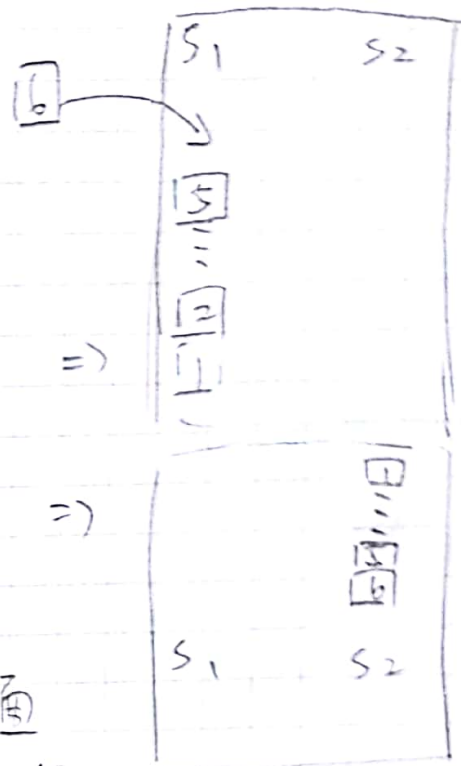


1. (1) Stack is LIFO

用 stack 來作出 enqueue

令 stack1, stack 為 s_1, s_2 .

enqueue: $\begin{cases} s_1 \cdot \text{push}(1, 2, 3 \dots) \\ s_2 \cdot \text{push}(s_1 \cdot \text{pop}()) \end{cases} \Rightarrow$

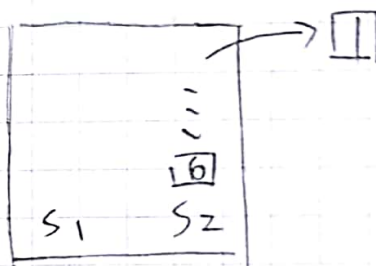


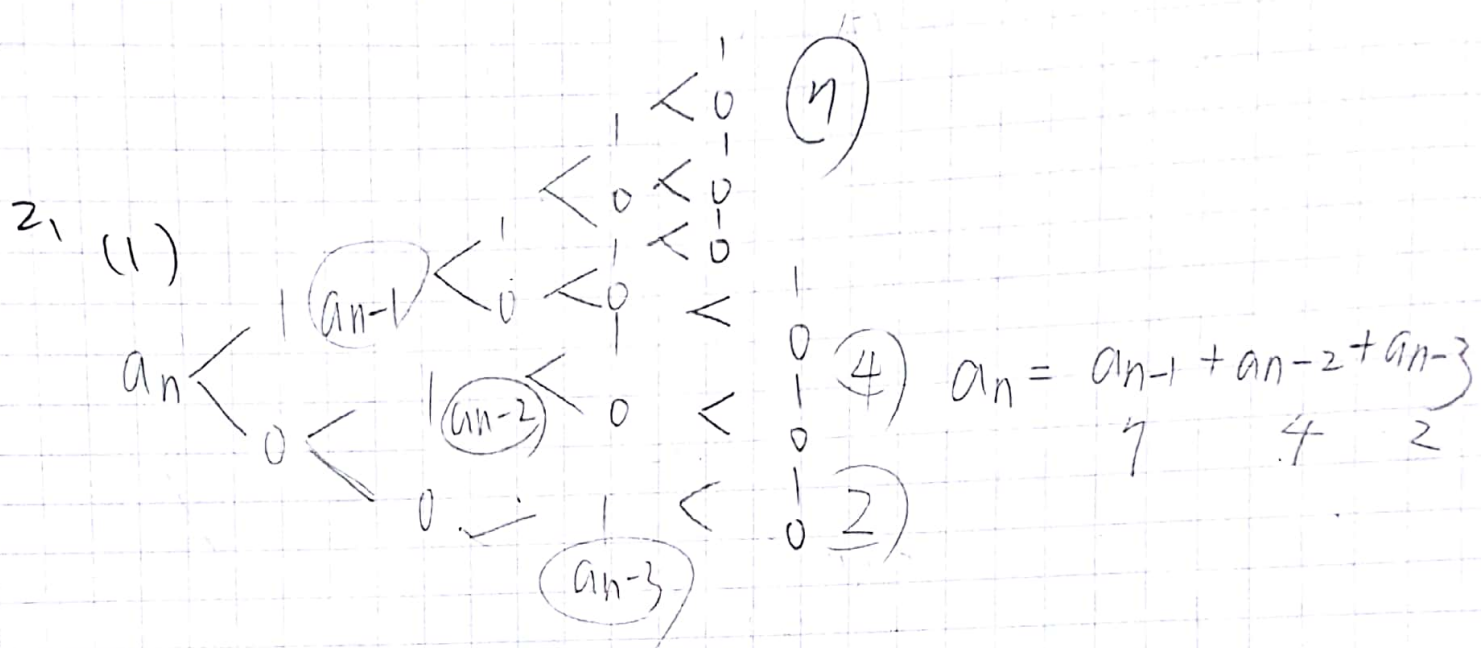
queue 之性質為 FIFO, 舊的在前面
故符合

(2) $s_2 \cdot \text{pop}()$

\Rightarrow

將 stack2 中之值從高
至低 pop 出去。





(2) 由(1)可看出 $a_1 = 2, a_2 = 4, a_3 = 7$

(3)

a_1	2
a_2	4
a_3	7
a_4	13
a_5	24
a_6	44
a_7	81
a_8	149
a_9	274
a_{10}	504

3.

pop 和 push 之量需要相等 balance

排列出之組合有 $C^{pop, push}(n, n)$ 種, 但因

有 stack 之限制, 不會每種路線均可行.

此問題可以轉化為走格子問題, 而

3.1.2, 並不少考量入此轉化問題。

When element 1 $C_1 = 1$

element 1, 2 $C_2 = 2$

element 1, 2, 3 $C_3 = 5$

element 1, 2, 3, 4 $C_4 = 14$

符合

Catalan Numbers

可以得出 $C_n = C_n^{2n} - C_{n-1}^{2n}$ ← 不合法之方式

$$= \frac{(2n)!}{n!n!} - \frac{(2n)!}{(n-1)!(n+1)!}$$

$$= \frac{(2n)!}{n!(n-1)!} \left(\frac{1}{n} - \frac{1}{n+1} \right)$$

$$= \frac{(2n)!}{n!(n-1)!} \times \frac{1}{n(n+1)}$$

$$= \frac{1}{n+1} C_n^{2n} \neq$$

4.

寫出 Fibonacci

0 1 1 2 3 5 8 13 21 ...

換用數學歸納法

$$F_1 \times F_3 - (F_2)^2 = -1 \quad \text{成立}$$

$$F_2 \times F_4 - (F_3)^2 = (-1)^2 \quad \text{成立}$$

若存在 $n \leq k$ 時亦成立則 $n = k+1$ 時：

$$F_k \times F_{k+2} - (F_{k+1})^2$$

$$\Rightarrow F_k (F_k + F_{k+1}) - (F_k + F_{k-1})^2$$

$$\Rightarrow F_k^2 + F_k F_{k+1} - F_k^2 - 2F_k F_{k-1} - F_{k-1}^2$$

$$\Rightarrow [F_k^2 + F_{k+1} F_k - F_{k-1} (2F_k + F_{k-1})] - F_{k-1}^2$$

$$\Rightarrow [F_k^2 + F_{k+1} F_k - F_{k-1} (F_k + F_{k+1})] - F_{k-1}^2$$

$$\Rightarrow [F_k^2 + F_{k+1} F_k - F_{k-1} F_k - F_{k-1} F_{k+1}] - F_{k-1}^2$$

$$\Rightarrow F_k (F_k + F_{k+1} - F_{k-1}) - F_{k-1} F_{k+1} - F_{k-1}^2$$

$$\Rightarrow F_k (F_k + \underbrace{F_{k+1} - F_{k-1}}_{=F_k}) - F_{k-1} F_{k+1} - F_{k-1}^2$$

$$\Rightarrow 2F_k^2 - F_{k-1} F_{k+1} - F_{k-1}^2 = (-1) [F_{k-1} F_{k+1} - F_k^2] \quad \text{X}$$