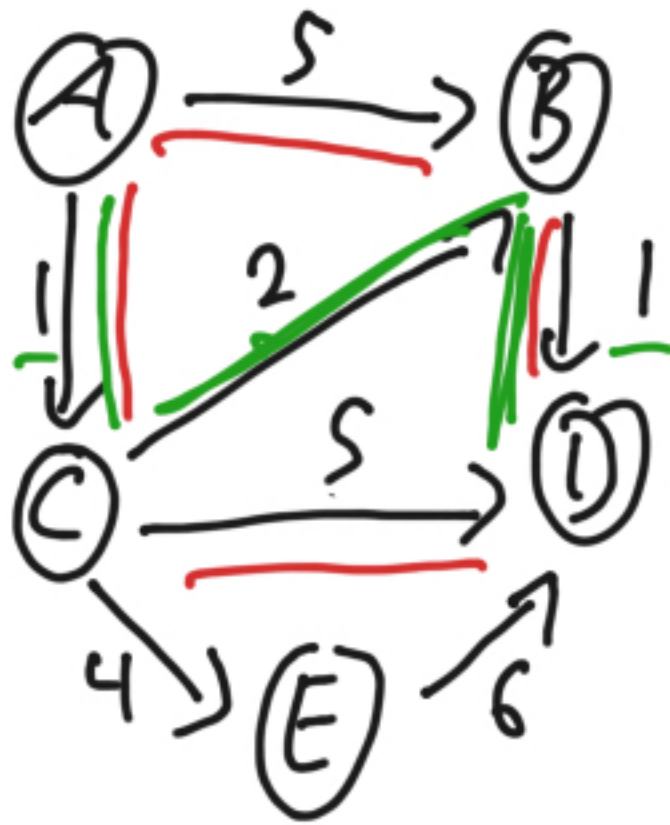


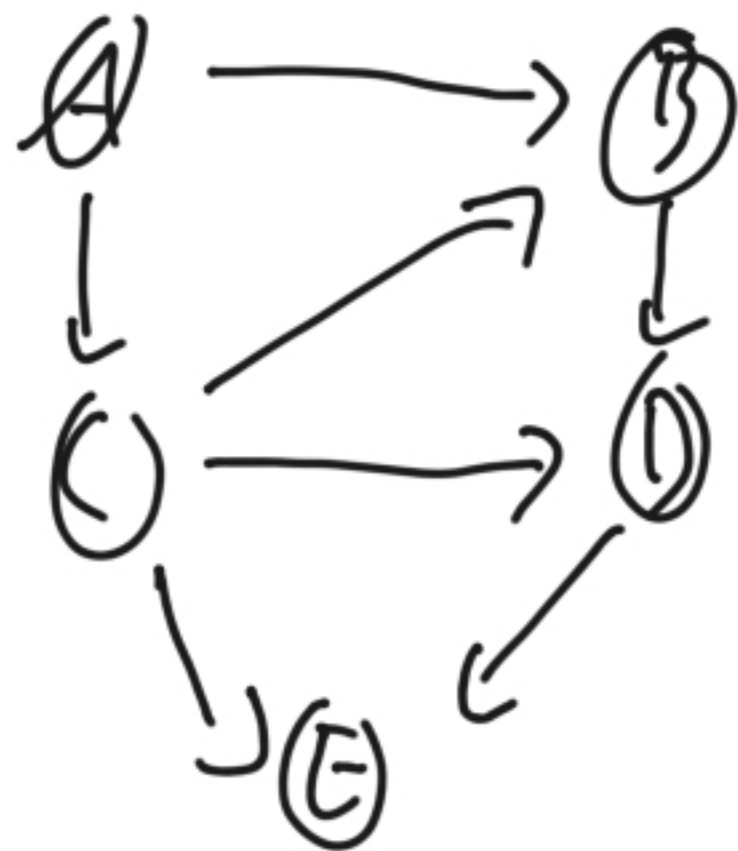
Graphs



vertices $V = \{A, B, C, D, E\}$

Edges $E = \{(A, B), (A, C), (B, D), (C, D), (C, B), (C, E), (E, D)\}$

directed v. undirected graphs
weighted v. unweighted graphs



Path: Sequence of vertices
between vertices s and t

$\langle s, i, j, \dots, t \rangle$

such that there is an
edge between each pair of
vertices,

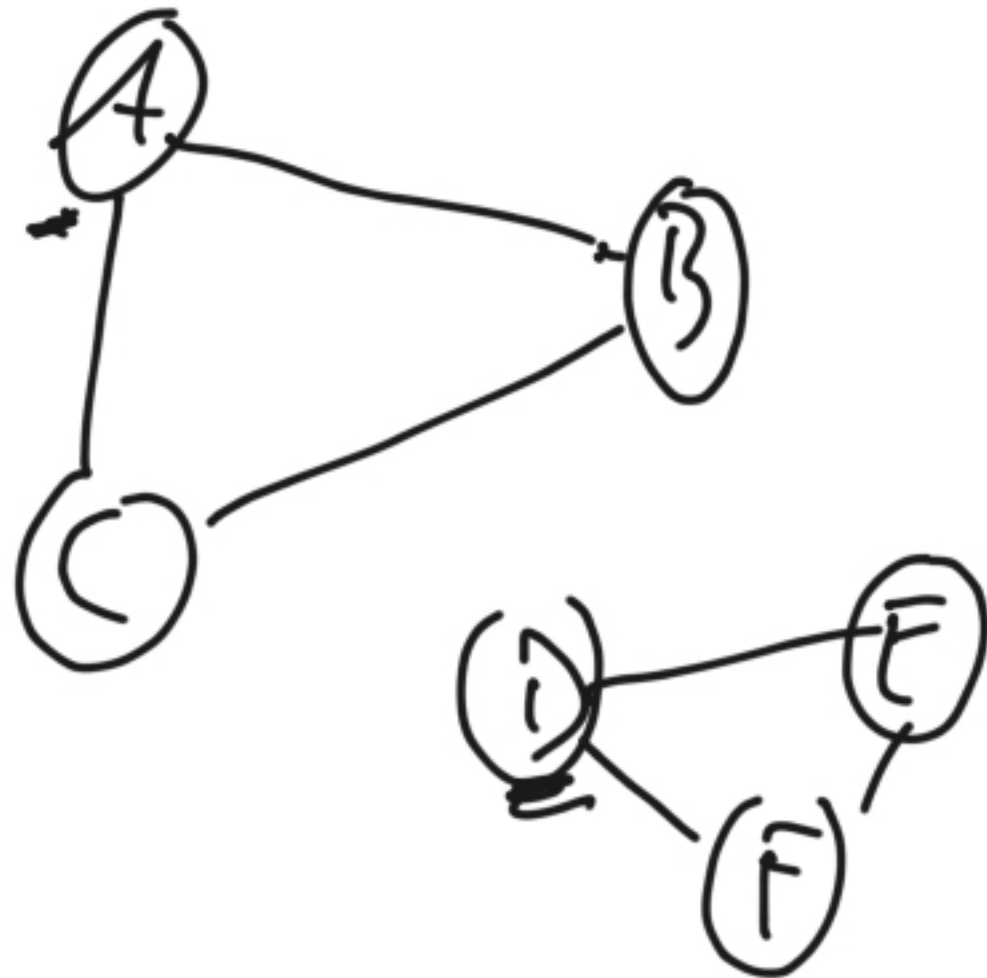
$(s, i) \in E, (i, j) \in E, \dots$

Length of path: # edges
 $= \# \text{ vertices} - 1$

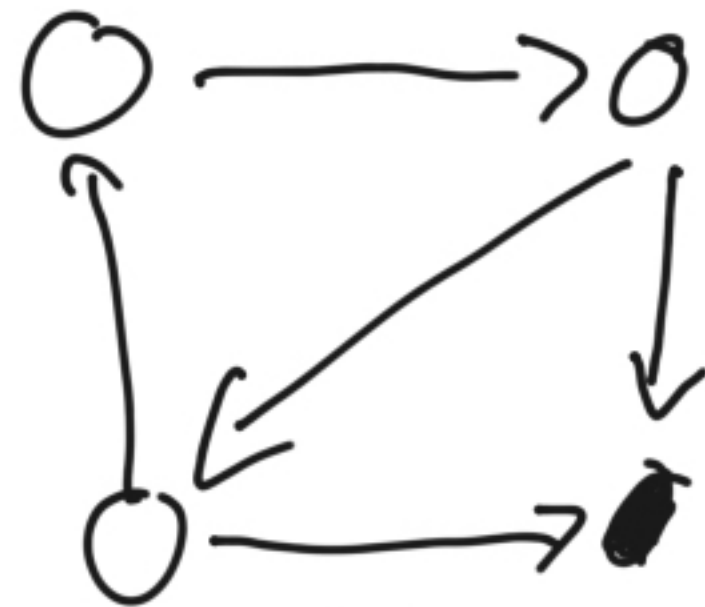
$\langle A, B, D \rangle$ has length 2

Connectedness

a graph is connected if there is some path between any pair of vertices.



directed graph

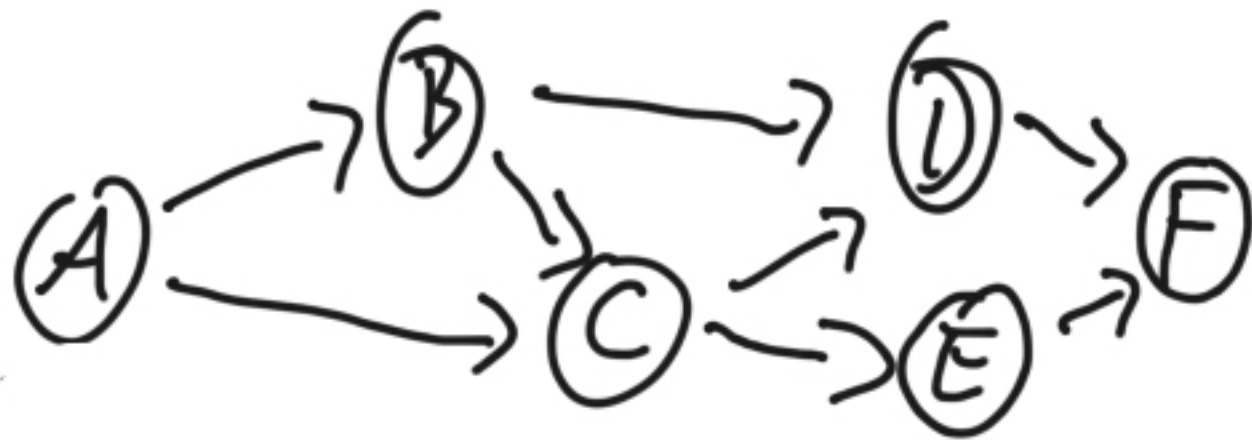


weakly connected

not strongly connected

A directed graph contains a cycle if there is some path $\langle A \dots A \rangle$ of length ≥ 2 .

A directed graph without cycles is called a Directed Acyclic Graph (DAG)



Topological Sort on a DAG



Queue:

A B C D E F

Step 1:

Compute in-degree of each vertex. If there is a path between s and t in the graph,

Step 2:

use a queue and enqueue all vertices with in-degree 0.

while queue not empty:

$s \leftarrow \text{queue.dequeue()}$

add s to output

for all neighbors t of s :

$t.\text{indegree} - 1$

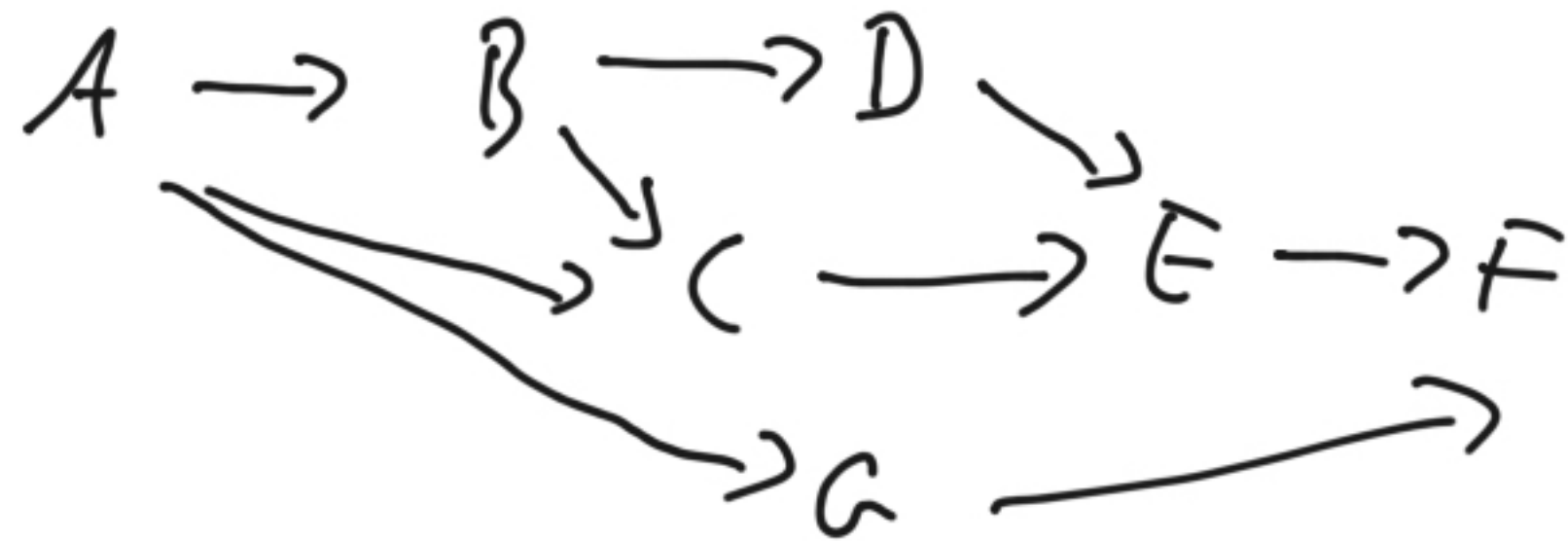
if $t.\text{indegree} == 0$:

$\text{queue.enqueue}(t)$

then s must appear before t in the top sort.

A B C D E F

A B C E D -



BFS

Breadth first search

cost = { 'A' : 0 }



use a queue q
Set discovered
 $q.enqueue(start)$

while not found and q not empty.

if s not in discovered

visit s

$s = q.dequeue()$

for each neighbor t of s :

if not t in discovered:

discovered.add(t)

if t is goal:

stop

enqueue(t)

cost(t) = cost(s) + 1

DISCOVER t