Contents	11 string 16 11.1KMP
1 basic 1.1 default 1.2 godcode 1.3 random 1.4 run.bat 1.5 run.sh	1 11.2minRotation 17 1 11.3PalindromeTree 17 1 11.4RollingHash 17 1 11.5SuffixArray 17 1 11.6trie 18 1 11.7Z-algorithm 18 11.8馬拉車 18
2 binarysearch 2.1 二分搜	1 12 tree 18 2 12.1DSUONTREE
3.2 fenwickTree	<pre>12.5HeavyLightDecomposition</pre>
3.8 2Dbit	<pre>#include <bits stdc++.h=""> using namespace std; #define masterspark ios::sync_with_stdio(0), cin.tie(0) ,cout.tie(0),cin.exceptions(cin.failbit);</bits></pre>
4.5 lowConvexHull	<pre>#define int long long #define pp pair<int, int=""> #define ff first #define ss second</int,></pre>
5.5 對偶建圖	<pre>#define forr(i,n) for(int i = 1; i <= n;++i) #define rep(i,j,n) for(int i = j; i < n;++i) #define PB push_back #define PF push_front #define EB emplace_back</pre>
6.1 Point 6.2 Line 6.3 Circle 6.4 圓多邊形面積 6.5 圓三角形面積 6.6 半平面交 6.7 圓線交 6.8 圓圓突 6.9 線線交 6.10ConvexHull 6.11Hulltrick 6.12點線距	<pre>#define all(v) (v).begin(), (v).end() #define FZ(x) memset(x, 0, sizeof(x)) //fill zero #define SZ(x) ((int)x.size()) bool chmin(auto &a, auto b) { return (b < a) and (a = b</pre>
6.13MEC 6.14MEC2 6.15旋轉卡尺 6.15旋轉卡尺 6.16Minkowski 6.17PointInPolygon 6.18UnionOfCircles 6.19UnionOfPolygons 6.20圓公切線 6.21點圓切線 6.21點圓切線	<pre>10</pre>
7 graph 7.1 BCC	// freopen("stdout","w",stdout); // cin >> t; while(t){ solve(); } return 0; }
8 math	13 1.2 godcode
8.1 DiscreteSqrt 8.2 excrt 8.3 exgcd 8.4 FFT 8.5 josephus 8.6 Theorem 8.7 Primes	#pragma GCC optimize("03,unroll-loops") #pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt") 編譯指令: g++ -std=c++20 -w -Wfatal-errors -Wall - Wshadow -fsanitize=undefined
8.8 millerrabin	15
8.14sieve	#define SECs ((double)clock() / CLOCKS_PER_SEC) struct KeyHasher { size_t operator()(const Key& k) const { return k.first + k.second * 100000; }
9.3 SmallestLexicographic	<pre>16 } }; 16 typedef unordered_map<key,int,keyhasher> map_t; 16 16 17 18 19 19 19 19 19 19 19</key,int,keyhasher></pre>

1.3 random

1.4 run.bat

```
@echo off
g++ ac.cpp -o ac.exe
g++ wa.cpp -o wa.exe
set /a num=1
:loop
    echo %num%
    python gen.py > input
    ac.exe < input > ac
    wa.exe < input > wa
    fc ac wa
    set /a num=num+1
if not errorlevel 1 goto loop
```

1.5 run.sh

```
set -e
for ((i=0;;i++))
do
    echo "$i"
    python gen.py > in
    ./ac < in > ac.out
    ./wa < in > wa.out
    diff ac.out wa.out || break
done
```

2 binarysearch

2.1 二分搜

```
int bsearch_1(int l, int r)
{
   while (l < r)
       int mid = l + r \gg 1;
       if (check(mid)) r = mid;
       else l = mid + 1;
   return 1;
// .....0000000000
int bsearch_2(int 1, int r)
   while (l < r)
   {
       int mid = l + r + 1 >> 1;
       if (check(mid)) l = mid;
       else r = mid - 1;
   return 1;
// 000000000.....
int m = *ranges::partition_point(views::iota(0LL,(int)1
    e9+9),[&](int a){
   return check(a) > k;
   });
//[begin,last)
//1111111000000000000
//搜左邊數過來第一個 ∅
//都是 1 會回傳 last
```

```
int partitionpoint(int L,int R,function<bool(int)> chk)
    {
    int l = L,r = R-1;
    while(r - l > 10){
        int m = l + (r-l)/2;
        if(chk(m)) l = m;
        else r = m;
    }
    int m = l;
    while(m <= r){
        if(!chk(m)) break;
        ++m;
    }
    if(!chk(m)) return m;
    else return R;
}

//手工
2.2 三分搜
int l = 1,r = 100;</pre>
```

```
int l = 1,r = 100;
while(l < r) {
    int lmid = l + (r - l) / 3; // l + 1/3区间大小
    int rmid = r - (r - l) / 3; // r - 1/3区间大小
    lans = cal(lmid),rans = cal(rmid);
    // 求凹函数的极小值
    if(lans <= rans) r = rmid - 1;
    else l = lmid + 1;
}</pre>
```

3 dataStructure

3.1 DSU

```
struct STRUCT_DSU {
     vector<int> f, sz;
STRUCT_DSU(i32 n) : f(n), sz(n) {
          for (int i = 0; i < n; i++) {
               f[i] = i;
               s\bar{z}[\bar{i}] = 1;
          }
     int find(int x) {
          if (x == f[x]) return x;
          f[x] = find(f[x]);
          return f[x];
     void merge(int x, int y) {
          x = find(x), y = find(y);
          if (x == y) return;
if (sz[x] < sz[y])</pre>
               swap(x, y);
          sz[x] += sz[y];
          f[y] = x;
     bool same(int a, int b) {
          return (find(a) == find(b));
};
```

3.2 fenwickTree

```
struct fenwick {
    // [0, n]
    #define lowbit(x) (x & -x)
    int n;
    vector<i64> v;
    fenwick(i32 _n) : n(_n + 1), v(_n + 2, 0) {}
    void _add(i32 x, i64 u){
        for(;x <= n; x += lowbit(x)) v[x] += u;
    }
    i64 _qry(i32 x){
        int ret = 0;
        for(; x ; x -= lowbit(x)) ret += v[x];
        return ret;
    }
    i32 _lowerbound(i64 k) {
        i64 sum = 0;
        i32 p = 0;
        for (i32 i = (1 << __lg(n)); i; i >>= 1) {
```

```
i32 nxt = p + i;
if (nxt <= n && sum + v[nxt] < k) {
                         sum += v[nxt];
                         p = nxt;
            }
            return p + 1;
   void add(i32 x, i64 v) { _add(x + 1, v); }
i64 qry(i32 x) { return _qry(x + 1); }
i64 qry(i32 l,i32 r) { return qry(r) - qry(l - 1); }
      i32 lower_bound(i64 k) { return _lowerbound(k) - 1;
};
```

```
3.3 segmentTree1
template<class Info>
struct SegmentTree {
    inline i32 cl(i32 x) { return x << 1; }</pre>
    inline i32 cr(i32 x) { return (x << 1) | 1; }
    i32 n;
    vector<Info> info;
SegmentTree() : n(0) {}
    SegmentTree(i32 n_, Info v_{-} = Info()) { init(n_, v_{-}
    ); }
template<class T>
    SegmentTree(vector<T> init_) { init(init_); }
    void init(i32 n_, Info v_ = Info()) { init(vector(
        n_, v_)); }
    template<class T>
    void init(vector<T> init_) {
        n = init_.size();
        info.assign(4 << __lg(n), Info());
function<void(i32, i32, i32)> build = [&](i32 p
            , i32 l, i32 r) {
if (r - l == 1) {
                 info[p] = init_[l];
                 return:
             i32 m = (l + r) >> 1;
             build(cl(p), l, m);
             build(cr(p), m, r);
            pull(p);
        build(1, 0, n);
    void pull(i32 p) { info[p] = merge(info[cl(p)],
        info[cr(p)]); }
    void modify(i32 p, i32 l, i32 r, i32 x, const Info
        &v) {
        if (r - l == 1) {
            info[p] = v;
             return;
        i32 m = (l + r) >> 1;
        if (x < m) modify(cl(p), l, m, x, v);</pre>
        else modify(cr(p), m, r, x, v);
        pull(p);
    void modify(i32 p, const Info &v) { modify(1, 0, n,
          p, v); }
    Info rangeQuery(i32 p, i32 l, i32 r, i32 x, i32 y)
        if (l >= y || r <= x) return Info();</pre>
        if (l >= x \&\& r <= y) return info[p];
        i32 m = (l + r) >> 1;
        return merge(rangeQuery(cl(p), l, m, x, y),
             rangeQuery(cr(p), m, r, x, y));
    Info rangeQuery(i32 l, i32 r) { return rangeQuery
        (1, 0, n, l, r); 
    template<class F>
    i32 findFirst(i32 p, i32 l, i32 r, i32 x, i32 y, F
        &&pred) {
        if (1 \ge y \mid | r \le x) return -1;
```

```
if (l >= x && r <= y && !pred(info[p])) return</pre>
              -1;
         if (r - l == 1) return l;
         i32 m = (l + r) >> 1;
         i32 res = findFirst(cl(p), l, m, x, y, pred);
         if (res == -1) res = findFirst(cr(p), m, r, x,
             y, pred);
         return res;
    template<class F>
    i32 findFirst(i32 l, i32 r, F &&pred) { return
         findFirst(1, 0, n, l, r, pred); }
    template<class F>
    i32 findLast(i32 p, i32 l, i32 r, i32 x, i32 y, F
         &&pred) {
         if (l >= y | | r <= x) return -1;
if (l >= x && r <= y && !pred(info[p])) return</pre>
             -1;
         if (r - l == 1) return l;
i32 m = (l + r) >> 1;
         i32 res = findLast(cr(p), m, r, x, y, pred);
         if (res == -1) res = findLast(cl(p), l, m, x, y
              , pred);
         return res;
    template<class F>
    i32 findLast(i32 l, i32 r, F &&pred) { return
         findLast(1, 0, n, l, r, pred); }
};
```

3.4 segmentTree2

```
// [l, r)
template<class Info, class Tag>
struct segTree {
     inline i32 cl(i32 x) { return x << 1; }</pre>
     inline i32 cr(i32 x) { return (x << 1) | 1; }
    i32 n:
    vector<Info> info;
    vector<Tag> tag;
    segTree(): n(0) {}
segTree(i32 n_, Info v_ = Info()) {
         init(n_{, v_{)};
    template<class T>
    segTree(vector<T> init_) {
         init(init_);
    void init(i32 n_, Info v_ = Info()) {
         init(vector(n_, v_));
     template<class T>
    void init(vector<T> init_) {
         n = init_.size();
         info.assign(4 << __lg(n), Info());
tag.assign(4 << __lg(n), Tag());</pre>
         function<void(i32, i32, i32)> build = [&](i32 p
              , i32 l, i32 r) {
if (r - l == 1) {
                  info[p] = init_[l];
                  return;
             i32 m = (l + r) >> 1;
             build(cl(p), l, m);
build(cr(p), m, r);
             pull(p, 1, r);
         build(1, 0, n);
     void pull(i32 p, i32 l, i32 r) {
         i32 m = (l + r) >> 1;
         push(cl(p), 1, m);
         push(cr(p), m, r)
         info[p] = merge(info[cl(p)], info[cr(p)]);
     void rangeModify(i32 p, i32 l, i32 r, i32 x, i32 y,
          const Tag &v) {
         push(p, 1, r);
         if (1 >= x \& r <= y) {
             tag[p] += v;
```

return;

```
i32 m = (l + r) >> 1;
          if (x < m) rangeModify(cl(p), l, m, x, y, v);
          if (y > m) rangeModify(cr(p), m, r, x, y, v);
          pull(p, l, r);
     Info rangeQuery(i32 p, i32 l, i32 r, i32 x, i32 y)
         push(p, l, r);
if (l >= y || r <= x) {
    return Info();</pre>
          if (1 >= x \& r <= y) {
               return info[p];
          i32 m = (l + r) >> 1;
          return merge(rangeQuery(cl(p), l, m, x, y),
               rangeQuery(cr(p), m, r, x, y));
     Info rangeQuery(i32 l, i32 r) { return rangeQuery
     (1, 0, n, l, r); }
void rangeModify(i32 l, i32 r, const Tag &v) {
     rangeModify(1, 0, n, l, r, v); } void push(i32 p, i32 l, i32 r) { // need compelete
          if (tag[p].add != 0) {
              info[p].v += tag[p].add * (r - l);
if (r - l != 1) {
                   tag[cl(p)].add += tag[p].add;
                   tag[cr(p)].add += tag[p].add;
               tag[p].add = 0;
          }
     }
};
```

3.5 persistantSegTree

```
struct pSeg{
    struct node{
        int v;
node *1,*r;
    int n;
    vector<node*> ver;
    node* build(int l,int r){
        node* x = new node();
        if(l == r){
            x \rightarrow v = 0;
            return x;
        int m = (l+r)/2;
        x->l = build(l,m);
        x->r = build(m+1,r);
        x->v = x->l->v + x->r->v;
        return x;
    void init(int _n){
        n = _n+2;
        ver.PB(build(0,n-1));
    int qry(node* now,int l,int r,int ql,int qr){
        if(ql \ll l \& r \ll qr){
            return now->v;
        int m = (1+r)/2, ret = 0;
        if(ql <= m)ret += qry(now->1,1,m,ql,qr);
        if(qr > m )ret += qry(now->r,m+1,r,ql,qr);
        return ret;
    node* upd(node* prv,int l,int r,int p,int v){
        node* x = new node();
        if(l == r){
            return x;
        int m = (l+r)/2;
        if(p \ll m) {
            x\rightarrow l = upd(prv\rightarrow l, l, m, p, v);
            x->r = prv->r;
        }else{
            x->1 = prv->1;
            x->r = upd(prv->r,m+1,r,p,v);
```

3.6 countMinimumSeg

```
//count zeros on segmentTree
struct segTree{
     #define cl (i<<1)
     #define cr ((i << 1)+1)
     pp_seg[MXN*4]
     int tag[MXN*4];
     pp comb(pp a,pp b){
          if(a.ff < b.ff) return a;
          if(a.ff > b.ff) return b;
          return pp{a.ff,a.ss+b.ss};
     void push(int i,int l,int r){
          if(tag[i]){
              seg[i].ff += tag[i];
if(r - l > 1){
    tag[cl] += tag[i];
                   tag[cr] += tag[i];
               tag[i] = 0;
          }
     void pull(int i,int l,int r){
   int m = (r-l)/2 + l;
          push(cl,l,m);
          push(cr,m,r);
          seg[i] = comb(seg[cl],seg[cr]);
     void build(int i,int l,int r){
    if(r - l <= 1){</pre>
              seg[i] = pp{0,1};
              return:
          int m = (r-1)/2 + 1;
          build(cl,1,m);
          build(cr,m,r);
          pull(i,l,r);
     void upd(int i,int l,int r,int ql,int qr,int x){
          push(i,1,r);
if(q1 <= 1 && r <= qr){</pre>
              tag[i] += x;
              return:
          int m = (r-1)/2 + 1;
          if(ql < m) upd(cl,l,m,ql,qr,x);</pre>
          if(qr > m) upd(cr,m,r,ql,qr,x);
          pull(i,l,r);
     int qry(){
          //count zero
          if(seg[1].ff == 0) return seg[1].ss;
          return 0;
     void upd(int l,int r,int x){
          upd(1,0,MXN,l,r,x);
}st;
```

3.7 LiChaoSegTree

```
const int inf = numeric_limits<i64>::max()/2;
```

```
struct Line {
    // y = ax + b
                                                                 void add(int l,int u,int r,int d,int x){
    i64 a{0}, b{-inf};
                                                                     ++r,++d;
    i64 operator()(i64 x) {
    return a * x + b;
                                                                     add(1,u,x);
                                                                     add(1,d,-x);
                                                                     add(r,u,-x);
};
                                                                     add(r,d,x);
struct Seg{
                                                                 int qry(int l,int u,int r,int d){
    int l, r;
Seg *ls{},*rs{};
                                                                     --l,--u;
                                                                     return qry(r,d) - qry(r,u) - qry(l,d) + qry(l,u)
    Line f{};
    Seg(int l, int r) : l(l), r(r) {}
                                                                 }
    void add(Line g){
                                                            };
        int m = (1+r)/2
        if (g(m) > f(m)) swap(g, f);
                                                             4
                                                                 dp
        if(g.b == -inf | | r - l == 1) return;
        if(g.a < f.a){
                                                             4.1 digit
            if(!ls) ls = new Seg(l,m);
            ls->add(g);
                                                            ll dp[MXN_BIT][PRE_NUM][LIMIT][F0];//字串位置, 根據題目
        }else{
                                                                 的值,是否上界,前導0
                                                             ll dfs(int i,int pre, bool lim, bool f0, const string&
            if(!rs) rs = new Seg(m,r);
            rs->add(g);
        }
                                                                 if(v[i][pre][f0][lim]) return dp[i][pre][f0][lim];
                                                                 v[i][pre][f0][lim] = true;
    i64 qry(i64 x){
        int m = (l+r) / 2;
                                                                 if(i == str.size())
        i64 y = f(x);
                                                                     return dp[i][pre][f0][lim] = 1;
        if(x < m \&\& ls) y = max({y,ls->qry(x)});
        if(x >= m \&\& rs) y = max({y,rs->qry(x)});
                                                                 ll ret = 0, h = \lim ? str[i] : '9';
        return y;
                                                                 for(int j='0'; j<=h; j++){
   if(abs(j-pre)>=2 || f0){
    }
auto add = [&](Line g,int ql,int qr){ //新增線段 [ql,qr
                                                                         ret += dfs(i+1, j, j==h && lim, f0 && j=='0
                                                                              ', str);
    auto find = [&](auto &&self,Seg * now,int l,int r)
         -> void {
                                                                 return dp[i][pre][f0][lim] = ret;
        if(ql <= l \& r <= qr){
                                                            }
            now->add(g);
            return;
                                                             4.2 p_median
        int m = (l+r) / 2;
                                                             void p_Median(){
        if(ql < m) {
                                                                 for (int i=1; i<=N; ++i)
                                                                     for (int j=i; j<=N; ++j){
    m = (i+j)/2,d[i][j] = 0;
            if(!now->ls) now->ls = new Seg(l,m);
            self(self,now->ls,l,m);
                                                                                                           // m是中位
                                                                         數, d[i][j]為距離的總和
for (int k=i; k<=j; ++k) d[i][j] += abs(arr
            if(!now->rs) now->rs = new Seg(m,r);
                                                                              [k] - arr[m]);
            self(self,now->rs,m,r);
                                                                 for (int p=1; p<=P; ++p)</pre>
                                                                     for (int n=1; n<=N; ++n){
    find(find,st,-ninf,ninf);
                                                                         dp[p][n] = 1e9;
                                                                         for (int k=p; k<=n; ++k)
   if (dp[p-1][k-1] + d[k][n] < dp[p][n]){</pre>
//Seg *st = new Seg(-ninf,ninf); // [l,r)
                                                                                  dp[p][n] = dp[p-1][k-1] + d[k][n];
3.8 2Dbit
                                                                                  r[p][n] = k;
                                                                                                  // 從第k個位置往右
                                                                                      到第 j個 位置
struct fenwick{
                                                                             }
    #define lowbit(x) (x&-x)
                                                                     }
    int n,m;
                                                            }
    vector<vector<int>> v;
    4.3 sosdp
    void add(int x,int y,int u){
                                                             // 求子集和 或超集和 -> !(mask & (1 << i))
        ++x,++y;
                                                             for(int i = 0; i<(1<<N); ++i) F[i] = A[i]; //預處理 狀
        for(;x < n; x += lowbit(x)){
                                                                 態權重
            for(int j = y; j < m; j += lowbit(j)) v[x][j
                                                             for(int i = 0; i < N; ++i)
                 ] += u;
                                                             for (int s = 0; s < (1 << N); ++s)
                                                               if (s & (1 << i))
    int qry(int x,int y){
                                                                 F[s] += F[s \land (1 << i)];
        ++x,++y;
        int ret = 0;
                                                             //窮舉子集合
        for(; x ; x -= lowbit(x)){
                                                            for(int s = mask; s ; s = (s-1)\&mask;)
             for(int j = y; j; j -= lowbit(j)) ret += v[
                                                             4.4 MinimumSteinerTree
                 x][j];
                                                            int dp[MXN][(1<<11)],vis[MXN];
//dp[i][S] -> 選了前K個點 以第i個點為第K+1個點的 生成
        return ret;
    //(l,u) <= (r,d)
                                                                 (1..K+1)的最小生成樹
                                                            rep(s,0,(1<<K)) forr(i,N) dp[i][s] = INF;
rep(j,0,K) dp[j+1][(1<<j)] = 0;</pre>
    //d -
```

//u +

```
rep(s,0,(1<< K)){
    forr(i,N){
      for(int a = s; a; a=(a-1)&s)
      dp[i][s] = min(dp[i][s], dp[i][s^a] + dp[i][a]);
           // node
    FZ(vis);
    priority_queue<pp,vector<pp>,greater<pp>> Q;
    forr(i,N) Q.emplace(dp[i][s],i);
    while(Q.size()){
      auto [d,u] = Q.top();Q.pop();
      if(vis[u]) continue;
      vis[u] = 1;
      for(auto [v,w]:E[u]){
   if(dp[u][s]+w < dp[v][s]) {
          dp[v][s] = dp[u][s]+w;
          Q.emplace(dp[v][s],v);
   }
rep(i,K+1,N+1) cout << dp[i][(1<<K)-1] <<'\n';
```

lowConvexHull

```
struct Line {
  mutable ll m, b, p;
   bool operator<(const Line& o) const { return m < o.m;</pre>
  bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
   const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
  return a / b - ((a ^ b) < 0 && a % b); }</pre>
  bool isect(iterator x, iterator y) {
     if (y == end()) { x->p = inf; return false; }
if (x->m == y->m) x->p = x->b > y->b ? inf : -inf;
     else x->p = div(y->b - x->b, x->m - y->m);
     return x->p >= y->p;
  void insert_line(ll m, ll b) {
     auto z = insert({m, b, 0}), y = z++, x = y;
while (isect(y, z)) z = erase(z);
if (x != begin() && isect(--x, y)) isect(x, y =
          erase(y));
     while ((y = x) != begin() \&\& (--x)->p >= y->p)
       isect(x, erase(y));
   ll eval(ll x) {
     assert(!empty());
     auto l = *lower_bound(x);
     return l.m * x + l.b;
};
```

flow 5

5.1 Dinic

```
struct Dinic{
  struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
    n = _n; s = _s; t = _t;
for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
    E[u].PB({v,f,SZ(E[v])});
    E[v].PB(\{u,0,SZ(E[u])-1\});
  bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;</pre>
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()){
       int u = que.front(); que.pop();
       for (auto it : E[u]){
```

```
if (it.f > 0 && level[it.v] == -1){
  level[it.v] = level[u]+1;
           que.push(it.v);
    } } }
    return level[t] != -1;
  int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
      if (it.f > 0 && level[it.v] == level[u]+1){
         int tf = DFS(it.v, min(nf,it.f));
         res += tf; nf -= tf; it.f -= tf;
         E[it.v][it.re].f += tf;
         if (nf == 0) return res;
    } }
    if (!res) level[u] = -1;
    return res;
  int flow(int res=0){
    while ( BFS() )
      res += DFS(s,2147483647);
    return res;
} }flow;
```

5.2 isap

```
struct Maxflow {
   static const int MAXV = 20010;
   static const int INF = 1000000;
   struct Edge {
     int v, c, r;
Edge(int _v, int _c, int _r):
   v(_v), c(_c), r(_r) {}
   int s, t;
   vector<Edge> G[MAXV*2];
int iter[MAXV*2], d[MAXV*2], gap[MAXV*2], tot;
   void init(int x) {
     tot = x+2;
     s = x+1, t = x+2;
     for(int i = 0; i <= tot; i++) {</pre>
       G[i].clear();
       iter[i] = d[i] = gap[i] = 0;
   } }
   void addEdge(int u, int v, int c) {
     G[u].push_back(Edge(v, c, SZ(G[v])));
     G[v].push\_back(Edge(u, 0, SZ(G[u]) - 1));
   int dfs(int p, int flow) {
     if(p == t) return flow;
     for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
       Edge &e = G[p][i];
       if(e.c > 0 \& d[p] == d[e.v]+1) {
          int f = dfs(e.v, min(flow, e.c));
          if(f) {
            e.c -= f:
            G[e.v][e.r].c += f;
            return f;
     } } }
     if((--gap[d[p]]) == 0) d[s] = tot;
     else {
       d[p]++:
       iter[p] = 0;
       ++gap[d[p]];
     return 0;
   int solve() {
     int res = 0;
     gap[0] = tot;
     for(res = 0; d[s] < tot; res += dfs(s, INF));</pre>
     return res;
   void reset() {
     for(int i=0;i<=tot;i++) {</pre>
       iter[i]=d[i]=gap[i]=0;
5.3
       KM
```

struct KM{ // max weight, for min negate the weights

```
int n, mx[MXN], my[MXN], pa[MXN];
ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
                                                                            横槓(n*(m-1)); B: 直槓((n-1)*m); C: 斜槓((n-1)
                                                                     *(m-1));
//n 列 m 行平面圖 (1-base) S起點 (左上) T 終點 (右下)
  bool vx[MXN], vy[MXN];
void init(int _n) { // 1-based, N個節點
                                                                     forr(s,(n-1)){
                                                                          int M = (m-1)*2;
     for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);</pre>
                                                                          forr(i,M){
                                                                               int id = i + (s-1)*M;
   void addEdge(int x, int y, ll w) \{g[x][y] = w;\} // \pm
                                                                               if(i&1){
                                                                                   int u = (s < n-1) ? ((i+1) + s*M) : T;
int e = (i > 1) ? id - 1 : T;
add(id,e,B[s-1][(i-1)/2]);
        邊的集合節點x連邊右邊集合節點y權重為w
  void augment(int y) {
     for(int x, z; y; y = z)
  x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
                                                                                   add(id,u,A[s][(i-1)/2]);
  void bfs(int st) {
                                                                                    if(i == M) add(id,S,B[s-1][m-1])
                                                                                    if(s == 1) add(id,S,A[s-1][i/2-1]);
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
                                                                                    int w = C[s-1][i/2-1];
     queue<int> q; q.push(st);
     for(;;) {
  while(q.size()) {
                                                                                    add(id,id-1,w);
                                                                              }
          int x=q.front(); q.pop(); vx[x]=1;
                                                                          }
          for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
                                                                     }
            ll t = lx[x]+ly[y]-g[x][y];
                                                                             最小花費最大流 dijkstra 不能負值
            if(t==0){
              pa[y]=x
              if(!my[y]){augment(y);return;}
                                                                     struct MinCostMaxFlow{
            vy[y]=1, q.push(my[y]);
}else if(sy[y]>t) pa[y]=x,sy[y]=t;
                                                                     typedef int Tcost;
                                                                        static const int MAXV = 20010;
       } }
                                                                        static const int INFf = 1000000;
                                                                        static const Tcost INFc = 1e9;
       11 cut = INF;
       for(int y=1; y<=n; ++y)</pre>
                                                                        struct Edge{
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
                                                                          int v, cap;
       for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;</pre>
                                                                          Tcost w;
                                                                          int rev;
         if(vy[j]) ly[j] += cut;
else sy[j] -= cut;
                                                                          Edge(){}
                                                                          Edge(int t2, int t3, Tcost t4, int t5)
                                                                          : v(t2), cap(t3), w(t4), rev(t5) {}
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
  if(!my[y]){augment(y); return;}</pre>
                                                                        int V, s, t;
         vy[y]=1, q.push(my[y]);
                                                                        vector<Edge> g[MAXV];
                                                                       void init(int n, int _s, int _t){
   V = n; s = _s; t = _t;
   for(int i = 0; i <= V; i++) g[i].clear();</pre>
    } }
  11 solve(){ // 回傳值為完美匹配下的最大總權重
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0); fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y) //
                                                                        void addEdge(int a, int b, int cap, Tcost w){
                                                                          g[a].push_back(Edge(b, cap, w, (int)g[b].size()));
g[b].push_back(Edge(a, 0, -w, (int)g[a].size()-1));
          1-base
       lx[x] = max(lx[x], g[x][y])
     for(int x=1; x<=n; ++x) bfs(x);</pre>
                                                                        Tcost d[MAXV];
     11 \text{ ans} = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
                                                                        int id[MAXV], mom[MAXV];
                                                                        bool inqu[MĀXV];
     return ans:
} }graph;
                                                                        queue<int> q;
                                                                        pair<int,Tcost> solve(){
                                                                          int mxf = 0; Tcost mnc = 0;
5.4 匈牙利
                                                                          while(1){
bool dfs(int u){
                                                                            fill(d, d+1+V, INFc);
     for(int i : edge[u]){
                                                                            fill(inqu, inqu+1+V, 0);
          fill(mom, mom+1+V, -1);
                                                                            mom[s] = s;
                                                                            d[s] = 0;
                   match[i] = u; match[u] = i; // 紀錄匹配
                                                                            q.push(s); inqu[s] = 1;
                   return true;
                                                                            while(q.size()){
              }
                                                                               int u = q.front(); q.pop();
                                                                               inqu[u] = 0;
for(int i = 0; i < (int) g[u].size(); i++){</pre>
         }
     return false;
                                                                                 Edge &e = g[u][i];
                                                                                 int v = e.v
                                                                                 if(e.cap > 0 \& d[v] > d[u]+e.w){
int hungarian(){
     int ans = 0;
                                                                                   d[v] = d[u] + e.w;
    memset(match, -1, sizeof(match));
for(int i = 1;i <= lhs; i++){
    // 記得每次使用需清空vis陣列
                                                                                   mom[v] = u;
                                                                                   id[v] = i;
                                                                                    if(!inqu[v]) q.push(v), inqu[v] = 1;
         memset(vis, 0, sizeof(vis));
                                                                            } } }
         if(dfs(i)) ans++;
                                                                            if(mom[t] == -1) break ;
                                                                            int df = INFf;
     return ans;
                                                                            for(int u = t; u != s; u = mom[u])
                                                                            df = min(df, g[mom[u]][id[u]].cap);
for(int u = t; u != s; u = mom[u]){
  Edge &e = g[mom[u]][id[u]];
5.5 對偶建圖
                                                                               e.cap
auto add = [&](int u,int v,int w){
                                                                              g[e.v][e.rev].cap += df;
     E[u].EB(v,w);
     E[v].EB(u,w);
                                                                            mxf += df;
                                                                            mnc += df*d[t];
};
```

8

```
NTOU Miaotomata
                                                                  friend T ori(pt a, pt b, pt c) { return (b - a) ^ (c
                                                                       - a); }
    return {mxf,mnc};
} }flow;
                                                                  friend T abs2(pt a) { return a * a; }
                                                                };
                                                                using numbers::pi; // c++20
5.7 最小花費最大流 SPFA
                                                                const ld pi = acos(-1);
                                                                const ld eps = 1e-8L;
struct zkwflow{
                                                                using Pt = pt<ld>;
  static const int maxN=10000;
                                                                int sgn(ld x) { return (x > -eps) - (x < eps); } //
  struct Edge{ int v,f,re; ll'w;};
int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
                                                                    dcmp == sgn
                                                                ld abs(Pt a) { return sqrt(abs2(a)); }
ld arg(Pt x) { return atan2(x.y, x.x); }
  vector<Edge> E[maxN];
  void init(int _n,int _s,int _t){
                                                                bool argcmp(Pt a, Pt b) { // arg(a) < arg(b)</pre>
    n=_n, s=_s, t=_t;
                                                                    int f = (Pt\{a.y, -a.x\} > Pt\{\}? 1 : -1) * (a != Pt
    for(int i=0;i<n;i++) E[i].clear();</pre>
                                                                         {});
                                                                    int g = (Pt\{b.y, -b.x\} > Pt\{\} ? 1 : -1) * (b != Pt
  void addEdge(int u,int v,int f,ll w){
    E[u].push_back({v,f,(int)E[v].size(),w});
                                                                    {});
return f == g ? (a ^ b) > 0 : f < g;
    E[v].push\_back({u,0,(int)}E[u].size()-1,-w});
                                                                Pt unit(Pt x) { return x / abs(x); }
Pt rotate(Pt u) { // pi / 2
  bool SPFA(){
    fill_n(dis,n,LLONG_MAX); fill_n(vis,n,false);
                                                                    return {-u.y, u.x};
    queue<int> q; q.push(s); dis[s]=0;
    while (!q.empty()){
                                                                Pt rotate(Pt u, ld a) {
      int u=q.front(); q.pop(); vis[u]=false;
                                                                    Pt v{sin(a), cos(a)};
return {u ^ v, u * v};
      for(auto &it:E[u]){
         if(it.f>0&&dis[it.v]>dis[u]+it.w){
           dis[it.v]=dis[u]+it.w;
           if(!vis[it.v]){
             vis[it.v]=true; q.push(it.v);
                                                                istream & operator>>(istream &s, Pt &a) { return s >> a.
    } } } }
                                                                    x \gg a.y;
    return dis[t]!=LLONG_MAX;
                                                                ostream &operator<<(ostream &s, Pt &a) { return s << "(
                                                                     " << a.x << ",
                                                                                     " << a.y << ")";}
  int DFS(int u,int nf){
    if(u==t) return nf;
                                                                bool collinearity(Pt a, Pt b, Pt c) { // 三點共線
    int res=0; vis[u]=true;
for(int &i=ptr[u];i<(int)E[u].size();i++){</pre>
                                                                    return ((b - a) \wedge (c - a)) == 0;
      auto &it=E[u][i];
      if(it.f>0&&dis[it.v]==dis[u]+it.w&&!vis[it.v]){
                                                                6.2 Line
         int tf=DFS(it.v,min(nf,it.f));
         res+=tf,nf-=tf,it.f-=tf;
                                                                struct Line {
         E[it.v][it.re].f+=tf;
                                                                    Pt a, b;
         if(nf==0){ vis[u]=false; break; }
                                                                    Pt dir() const { return b - a; }
      }
                                                                int PtSide(Pt p, Line L) {
    return res;
                                                                    // return sgn(ori(L.a, L.b, p) / abs(L.a - L.b));
                                                                    return sgn(ori(L.a, L.b, p));
  pair<int,ll> flow(){
    int flow=0; ll cost=0;
                                                                bool PtOnSeq(Pt p, Line L) {
    while (SPFA()){
                                                                    return PtSide(p, L) = 0 and sgn((p - L.a) * (p - L
      fill_n(ptr,n,0);
                                                                         .b)) <= 0;
       int f=DFS(s,INT_MAX)
      flow+=f; cost+=dis[t]*f;
                                                                Pt proj(Pt p, Line l) {
                                                                    Pt dir = unit(l.b - l.a);
    return{ flow,cost };
                                                                    return l.a + dir * (dir * (p - l.a));
    // reset: do nothing
} flow;
                                                                6.3 Circle
6
     geometry
                                                                struct Cir {
6.1 Point
                                                                    Pt o;
                                                                    ld r;
using ld = long double;
                                                                bool disjunct(const Cir &a, const Cir &b) {
template<class T>
struct pt{
                                                                    return sgn(abs(a.o - b.o) - a.r - b.r) >= 0;
  T x, y;
  pt(T_x,T_y):x(_x),y(_y){}
                                                                bool contain(const Cir &a, const Cir &b) {
  pt():x(0),y(0){}
                                                                    return sgn(a.r - b.r - abs(a.o - b.o)) >= 0;
                                                                }
  pt operator * (T c){ return pt(x*c,y*c);}
pt operator / (T c){ return pt(x/c,y/c);}
                                                                       圓多邊形面積
                                                                6.4
  pt operator + (pt a){ return pt(x+a.x,y+a.y);}
  pt operator - (pt a){ return pt(x-a.x,y-a.y);}
                                                                double CirclePoly(Cir C, const vector<Pt> &P) {
                                                                    auto arg = [&](Pt p, Pt q) { return atan2(p ^ q, p * q); };

double r2 = C.r * C.r / 2;
     operator * (pt a){ return x*a.x + y*a.y;}
     operator ^ (pt a){ return x*a.y - y*a.x;}
```

bool operator== (pt a) const { return x == a.x and y

== a.x && y < a.y);;

== a.y;;

auto tri = [&](Pt p, Pt q) {

auto det = a * a - b;

* C.r)/ abs2(d);

Pt d = q - p; auto a = (d * p) / abs2(d), b = (abs2(p) - C.r)

if (det <= 0) return arg(p, q) * r2;</pre>

```
-a - sqrt(det), t = min(1., -
          auto s = max(0.
                                                                               // Counterclockwise
                a + sqrt(det));
          if (t < 0 \text{ or } 1 <= s) return arg(p, q) * r2;
Pt u = p + d * s, v = p + d * t;
return arg(p, u) * r2 + (u \land v) / 2 + arg(v, q)
                                                                                  圓圓交
                                                                          6.8
                                                                          vector<Pt> CircleInter(Cir a, Cir b) {
                                                                               double d2 = abs2(a.o - b.o), d = sqrt(d2);
if (d < max(a.r, b.r) - min(a.r, b.r) || d > a.r +
     double sum = 0.0;
                                                                                    b.r) return {};
     for (int i = 0; i < P.size(); i++)</pre>
                                                                               Pt u = (a.0 + b.0) / 2 + (a.0 - b.0) * ((b.r * b.r - a.r * a.r) / (2 * d2));
     sum += tri(P[i] - C.o, P[(i + 1) \% P.size()] - C.o)
                                                                               return sum;
}
6.5 圓三角形面積
                                                                               if (sgn(v.x) == 0 \text{ and } sgn(v.y) == 0) \text{ return } \{u\};
                                                                               return {u - v, u + v}; // counter clockwise of a
double CircleTriangle(Pt a, Pt b, double r) {
   if (sgn(abs(a) - r) <= 0 and sgn(abs(b) - r) <= 0)</pre>
                                                                                  線線交
                                                                          6.9
          return abs(a ^ b) / 2;
                                                                          bool isInter(Line 1, Line m) {
     if (abs(a) > abs(b)) swap(a, b);
                                                                               if (PtOnSeg(m.a, 1) or PtOnSeg(m.b, 1) or
     auto I = CircleLineInter({{{}}, r{{}}, {a, b{}});
                                                                                    PtOnSeg(l.a, m) or PtOnSeg(l.b, m))
                                                                                    return true;
     erase_if(I, [&](Pt x) { return !Pt0nSeg(x, {a, b});
                                                                               return PtSide(m.a, 1) * PtSide(m.b, 1) < 0 and</pre>
            });
                                                                                        PtSide(l.a, m) * PtSide(l.b, m) < 0;
     if (I.size() == 1) return abs(a ^ I[0]) / 2 +
    SectorArea(I[0], b, r);
     if (I.size() == \overline{2})^{-2}
                                                                          Pt LineInter(Line l, Line m) {
           return SectorArea(a, I[0], r) + SectorArea(I
                                                                               double s = ori(m.a, m.b, l.a), t = ori(m.a, m.b, l.
                [1], b, r) + abs(I[0] \wedge I[1]) / 2;
                                                                               return (l.b * s - l.a * t) / (s - t);
     return SectorArea(a, b, r);
                                                                          6.10 ConvexHull
6.6 半平面交
                                                                          vector<Pt> Hull(vector<Pt> P) {
bool cover(Line L, Line P, Line Q) {
                                                                               sort(all(P));
     // PtSide(LineInter(P, Q), L) <= 0 or P, Q parallel
i128 u = (Q.a - P.a) ^ Q.dir();
i128 v = P.dir() ^ Q.dir();
                                                                               P.erase(unique(all(P)), P.end());
                                                                               P.insert(P.end(), P.rbegin() + 1, P.rend());
                                                                               vector<Pt> stk;
     i128 x = P.dir().x * u + (P.a - L.a).x * v;
i128 y = P.dir().y * u + (P.a - L.a).y * v;
return sgn(x * L.dir().y - y * L.dir().x) * sgn(v)
                                                                               for (auto p : P) {
                                                                                    auto it = stk.rbegin();
                                                                                    while (stk.rend() - it >= 2 and \
    ori(*next(it), *it, p) <= 0 and \
    (*next(it) < *it) == (*it < p)) {</pre>
vector<Line> HPI(vector<Line> P) {
                                                                                         it++;
     // line P.a -> P.b 的逆時針是半平面
     sort(all(P), [&](Line l, Line m) {
                                                                                    stk.resize(stk.rend() - it);
          if (argcmp(l.dir(), m.dir())) return true;
if (argcmp(m.dir(), l.dir())) return false;
                                                                                    stk.push_back(p);
          return ori(m.a, m.b, 1.a) > 0;
                                                                               stk.pop_back();
                                                                               return stk;
     });
     int n = P.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {</pre>
          if (i and !argcmp(P[i - 1].dir(), P[i].dir()))
                                                                          6.11 Hulltrick
                continue;
          while (l < r and cover(P[i], P[r - 1], P[r])) r</pre>
                                                                          struct Convex {
                                                                               int n;
          while (l < r \text{ and } cover(P[i], P[l], P[l + 1])) l
                                                                               vector<Pt> A, V, L, U;
                                                                               Convex(const vector<Pt> &_A) : A(_A), n(_A.size())
          P[++r] = P[i];
                                                                                     \{ // n >= 3
                                                                                    auto it = max_element(all(A));
     while (l < r and cover(P[l], P[r - 1], P[r])) r--;
while (l < r and cover(P[r], P[l], P[l + 1])) l++;
if (r - l <= 1 or !argcmp(P[l].dir(), P[r].dir()))</pre>
                                                                                    L.assign(A.begin(), it + 1);
                                                                                    U.assign(it, A.end()), U.push_back(A[0]);
                                                                                    for (int i = 0; i < n; i++) {
    V.push_back(A[(i + 1) % n] - A[i]);
           return {}; // empty
     if (cover(P[l + 1], P[l], P[r]))
  return {}; // infinity
return vector(P.begin() + l, P.begin() + r + 1);
                                                                               int inside(Pt p, const vector<Pt> &h, auto f) {
                                                                                    auto it = lower_bound(all(h), p, f);
| }
                                                                                    if (it == h.end()) return 0;
6.7 圓線交
                                                                                    if (it == h.begin()) return p == *it;
                                                                                    return 1 - sgn(ori(*prev(it), p, *it));
vector<Pt> CircleLineInter(Cir c, Line 1) {
     Pt H = proj(c.o, 1);
                                                                               // 0: out, 1: on, 2: in
     Pt dir = unit(l.b - l.a);
double h = abs(H - c.o);
                                                                               int inside(Pt p) {
                                                                                    return min(inside(p, L, less{}), inside(p, U,
     if (sgn(h - c.r) > 0) return {};
double d = sqrt(max((double)0., c.r * c.r - h * h))
                                                                                         greater{}));
```

if (sgn(d) == 0) return {H};
return {H - dir *d, H + dir * d};

static bool cmp(Pt a, Pt b) { return sgn(a ^ b) >

// A[i] is a far/closer tangent point

```
for (int k = 0; k < j; k++) {
   if (C.inside(P[k])) continue;
   C.o = Center(P[i], P[j], P[k]);</pre>
     int tangent(Pt v, bool close = true) {
   assert(v != Pt{});
         auto l = V.begin(), r = V.begin() + L.size() -
                                                                                   C.r = abs(C.o - P[i]);
         if (v < Pt{}) l = r, r = V.end();
         if (close) return (lower_bound(l, r, v, cmp) -
                                                                          }
              V.begin()) % n;
         return (upper_bound(l, r, v, cmp) - V.begin())
                                                                      return C;
    }
// closer tangent point array[0] -> array[1] 順時針
                                                                 6.14 MEC2
     array<int, 2> tangent2(Pt p) {
         array<int, 2> t{-1, -1};
if (inside(p) == 2) return t;
                                                                 PT arr[MXN];
                                                                 int n = 10;
         if (auto it = lower_bound(all(L), p); it != L.
  end() and p == *it) {
                                                                 double checky(double x, double y) {
                                                                      double cmax = 0;
for (int i = 0; i < n; i++) { // 過程中回傳距離^2
              int s = it - L.begin();
              return {(s + 1) % n, (s - 1 + n) % n};
                                                                           避免不必要的根號運算
                                                                          cmax = max(cmax, (arr[i].x - x) * (arr[i].x - x)
         if (auto it = lower_bound(all(U), p, greater{})
                                                                               ) + (arr[i].y - y) * (arr[i].y - y));
              ; it != U.end() and p == *it) {
              int s = it - U.begin() + L.size() - 1;
                                                                      return cmax;
              return \{(s + 1) \% n, (s - 1 + n) \% n\};
                                                                 double checkx(double x) {
                                                                      double yl = -1e9, yr = 1e9;
while (yr - yl > EPS) {
         for (int i = 0; i != t[0]; i = tangent((A[t[0]
             = i] - p), 0));
(int i = 0; i != t[1]; i = tangent((p - A[t
                                                                          double ml = (yl + yl + yr) / 3, mr = (yl + yr +
              [1] = i]), 1));
                                                                                yr) / 3;
         return t;
                                                                          if (checky(x, ml) < checky(x, mr))</pre>
                                                                               yr = mr;
     int find(int l, int r, Line L) {
                                                                          else
         if (r < 1) r += n;
                                                                               yl = ml;
         int s = PtSide(A[1 % n], L);
                                                                      }
         return *ranges::partition_point(views::iota(l,
                                                                 signed main() {
                                                                      double xl = -1e9, xr = 1e9; while (xr - xl > EPS) {
              [&](int m) {
                  return PtSide(A[m % n], L) == s;
                                                                          double ml = (xl + xl + xr) / 3, mr = (xl + xr + xr) / 3
                                                                                xr) / 3
    };
// Line A_x A_x+1 interset with L
                                                                          if (checkx(ml) < checkx(mr))</pre>
    vector<int> intersect(Line L) {
                                                                               xr = mr;
         int l = tangent(L.a - L.b), r = tangent(L.b - L.b)
                                                                          else
                                                                               xl = ml;
              .a);
         if (PtSide(A[l], L) * PtSide(A[r], L) >= 0)
                                                                      }
              return {};
         return {find(l, r, L) % n, find(r, l, L) % n};
                                                                 6.15
                                                                          旋轉卡尺
    }
};
                                                                 auto RotatingCalipers(const vector<Pt> &hull) { // 最遠
         點線距
                                                                      點對 回傳距離平方
6.12
                                                                      int n = hull.size();
double PtSegDist(Pt p, Line 1) {
                                                                      auto ret = abs2(hull[0]);
     double ans = min(abs(p - 1.a), abs(p - 1.b));
                                                                      ret = 0;
     if (sgn(abs(1.a - 1.b)) == 0) return ans;
                                                                      if (hull.size() <= 2) return abs2(hull[0] - hull</pre>
    if (sgn((1.a - 1.b) * (p - 1.b)) < 0) return ans; if (sgn((1.b - 1.a) * (p - 1.a)) < 0) return ans;
                                                                           [1]);
                                                                      for (int i = 0, j = 2; i < n; i++) {
    Pt a = hull[i], b = hull[(i + 1) % n];
    while(ori(hull[j], a, b) <
     return min(ans, abs(ori(p, l.a, l.b)) / abs(l.a - l
         .b));
                                                                                (ori(hull[(j + 1) % n], a, b)))
double SegDist(Line 1, Line m) {
                                                                               j = (j + 1) \% n;
                                                                          chmax(ret, abs2(a - hull[j]));
     return PtSegDist({0, 0}, {1.a - m.a, 1.b - m.b});
                                                                          chmax(ret, abs2(b - hull[j]));
6.13 MEC
                                                                      return ret;
                                                                 }
Pt Center(Pt a, Pt b, Pt c) {
    Pt x = (a + b) / 2;
                                                                 6.16 Minkowski
    Pt y = (b + c) / 2
     return LineInter(\{x, x + rotate(b - a)\}, \{y, y + a\}
                                                                 // P, Q, R(return) are counterclockwise order convex
         rotate(c - b)});
                                                                      polygon
                                                                 vector<Pt> Minkowski(vector<Pt> P, vector<Pt> Q) {
Cir MEC(vector<Pt> P)
                                                                      auto cmp = [\&](Pt a, Pt b) {
    mt19937 rng(time(0));
                                                                          return Pt{a.y, a.x} < Pt{b.y, b.x};
    auto reorder = [&](auto &R) {
                                                                          rotate(R.begin(), min_element(all(R), cmp), R.
                                                                               end());
         C = {P[i], 0};
for (int j = 0; j < i; j++) {
                                                                          R.push\_back(R[0]), R.push\_back(R[1]);
              if (C.inside(P[j])) continue;
                                                                      const int n = P.size(), m = Q.size();
```

reorder(P), reorder(Q);

vector<Pt> R;

 $C = \{(P[i] + P[j]) / 2, abs(P[i] - P[j]) / \}$

11

```
NTOU Miaotomata
         (int i = 0, j = 0, s; i < n or j < m; ) {
R.push_back(P[i] + Q[j]);
                                                                   vector<double> PolyUnion(const vector<vector<Pt>>> &P) {
                                                                        const int n = P.size();
         s = sgn((P[i + 1] - P[i]) \wedge (Q[j + 1] - Q[j]));
                                                                        vector<double> Area(n + 1);
         if (s >= 0) i++;
                                                                        vector<Line> Ls;
                                                                        for (int i = 0; i < n; i++)
  for (int j = 0; j < P[i].size(); j++)</pre>
         if (s <= 0) j++;
                                                                                  Ls.push_back(\{P[i][j], P[i][(j+1) \% P[i].
    return R;
}
                                                                                      size()]})
6.17 PointInPolygon
int inPoly(Pt p, const vector<Pt> &P) {
     const int n = P.size();
     int cnt = 0;
                                                                        };
    for (int i = 0; i < n; i++) {
         Pt a = P[i], b = P[(i + 1) \% n];
         if (PtOnSeg(p, {a, b})) return 1; // on edge
if ((sgn(a.y - p.y) == 1) ^ (sgn(b.y - p.y) ==
                                                                             Line L = Ls[l];
              cnt += sgn(ori(a, b, p));
    return cnt == 0 ? 0 : 2; // out, in
}
6.18
        UnionOfCircles
// Area[i] : area covered by at least i circle
// TODO:!!!aaa!!!
vector<double> CircleUnion(const vector<Cir> &C) {
    const int n = C.size();
vector<double> Area(n + 1)
    auto check = [&](int i, int j) {
   if (!contain(C[i], C[j]))
              return fals
         return sgn(C[i].r - C[j].r) > 0 or (sgn(C[i].r).r)
              - C[j].r) == 0 and i < j);
                                                                             });
                                                                             Pt lst{0, 0};
    struct Teve {
         double ang; int add; Pt p;
         bool operator<(const Teve &b) { return ang < b.</pre>
```

```
ana: }
    auto ang = [\&](Pt p) \{ return atan2(p.y, p.x); \};
    for (int i = 0; i < n; i++) {
         int cov = 1;
         vector<Teve> event;
         for (int j = 0; j < n; j++) if (i != j) {
   if (check(j, i)) cov++;</pre>
              else if (!check(i, j) and !disjunct(C[i], C
                   [j])) {
                   auto I = CircleInter(C[i], C[j]);
                   assert(I.size() == 2);
double a1 = ang(I[0] - C[i].o), a2 =
                       ang(I[1] - C[i].o);
                   event.push_back({a1, 1, I[0]});
event.push_back({a2, -1, I[1]});
                   if (a1 > a2) cov++;
         if (event.empty()) {
    Area[cov] += pi * C[i].r * C[i].r;
              continue;
         sort(all(event));
         event.push_back(event[0]);
         for (int j = 0; j + 1 < event.size(); j++) {
              cov += event[j].add;
              Area[cov] += (event[j].p \land event[j + 1].p)
              double theta = event[j + 1].ang - event[j].
                   ang;
              if (theta < 0) theta += 2 * pi;
              Area[cov] += (theta - sin(theta)) * C[i].r
                   * C[i].r / 2.;
         }
    return Area;
}
```

6.19 UnionOfPolygons

```
auto cmp = [&](Line &l, Line &r) {
   Pt u = l.b - l.a, v = r.b - r.a;
   if (argcmp(u, v)) return true;
   if (argcmp(v, u)) return false;
            return PtSide(l.a, r) < 0;</pre>
      sort(all(Ls), cmp);
for (int l = 0, r = 0; l < Ls.size(); l = r) {</pre>
           while (r < Ls.size() and !cmp(Ls[l], Ls[r])) r</pre>
           vector<pair<Pt, int>> event;
for (auto [c, d] : Ls) {
                 if (sgn((L.a - L.b) \land (c - d)) != 0) {
                      int s1 = PtSide(c, L) == 1;
int s2 = PtSide(d, L) == 1;
                      if (s1 ^ s2) event.emplace_back(
                LineInter(L, {c, d}), s1 ? 1 : -1);
} else if (PtSide(c, L) == 0 and sgn((L.a -
L.b) * (c - d)) > 0) {
                      event.emplace_back(c, 2)
                      event.emplace_back(d, -2);
           int cov = 0, tag = 0;
           for (auto [p, s] : event) {
                 if (cov >= tag) {
                     Area[cov] += lst ^ p;
Area[cov - tag] -= lst ^ p;
                 if (abs(s) == 1) cov += s;
                 else tag += s / 2;
                 lst = p;
           }
      for (int i = n - 1; i >= 0; i--) Area[i] += Area[i
            + 1];
      for (int i = 1; i <= n; i++) Area[i] /= 2;
      return Area;
};
           圓公切線
 6.20
 vector<Line> CircleTangent(Cir c1, Cir c2, int sign1) {
      // sign1 = 1 for outer tang, -1 for inter tang
      vector<Line> ret;
      ld d_sq = abs2(c1.o - c2.o);
      if (sgn(d_sq) == 0) return ret;
      ld d = sqrt(d_sq);
      Pt v = (c2.0 - c1.0) / d;
ld c = (c1.r - sign1 * c2.r) / d;
      if (c * c > 1) return ret;
ld h = sqrt(max(0.0, 1.0 - c * c));
      for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
   Pt n = Pt(v.x * c - sign2 * h * v.y, v.y * c +
        sign2 * h * v.x);
           Pt p1 = c1.0 + n * c1.r;
            Pt p2 = c2.0 + n * (c2.r * sign1);
            if (sgn(p1.x - p2.x) == 0 \& sgn(p1.y - p2.y)
                 p2 = p1 + rotate(c2.o - c1.o);
           ret.push_back({p1, p2});
    return ret;
 }
           點圓切線
 6.21
vector<Line> CircleTangent(Cir c, Pt p) {
```

```
vector<Line> z
                                                                                dfn[u] = low[u] = step++;
     double d = abs(p - c.o);
                                                                                stk[top++] = u;
     if (sgn(d - c.r) == 0) {
                                                                                for (auto v : E[u]) {
                                                                                    if (v == f) continue;
          Pt i = rotate(p - c.o);
                                                                                     if (dfn[v] == -1) {
          z.push_back({p, p + i});
                                                                                         DFS(v, u);
low[u] = min(low[u], low[v]);
     } else if (d > c.r) {
          double o = acos(c.r / d);
         Pt i = unit(p - c.o);
Pt j = rotate(i, o) * c.r;
                                                                                         if (low[v] >= dfn[u]) {
                                                                                              int z
          Pt k = rotate(i, -o) * c.r;
                                                                                              sccv[nScc].clear();
         z.push_back(\{\hat{c}.\hat{o} + \hat{j}, p\});
z.push_back(\{c.o + k, p\});
                                                                                              do {
                                                                                                   z = stk[--top];
                                                                                                   sccv[nScc].PB(z);
     return z;
                                                                                              } while (z != v);
}
                                                                                              sccv[nScc++].PB(u);
6.22 最近點對
                                                                                    } else
                                                                                         low[u] = min(low[u], dfn[v]);
pair<ld, pair<i32, i32>> ClosestPair(vector<Pt> &P) {
    // ans = dis * dis !!注意ans overflow問題
                                                                               }
                                                                          }
     if (P.size() == 1) { return {1e200L, {0, 0}}; }
                                                                           vector<vector<int>> solve() {
     pair<i32, i32> ansi;
                                                                               vector<vector<int>> res;
     auto ans = abs2(P[0] - P[1]);
                                                                                for (int i = 0; i < n; i++) dfn[i] = low[i] =
     ansi = \{0, 1\};
                                                                                     -1;
                                                                                for (int i = 0; i < n; i++)
if (dfn[i] == -1) {
     auto upd = [&](const Pt &a, const Pt &b) {
          auto dis = abs2(a - b);
          if (dis < ans) ans = dis, ansi.FF = a.id, ansi.</pre>
                                                                                         top = 0;
                                                                                         DFS(i, i);
               SS = b.id:
     auto cmpy = [](const Pt &a, const Pt &b) { return a
                                                                                REP(i, nScc) res.PB(sccv[i]);
          .y < b.y; };
                                                                                return res;
     vector<Pt> t(P.size() + 1);
                                                                      } graph;
     function<void(i32, i32)> rec = [&](i32 l, i32 r) {
          if (r - l <= 3) {
    for (i32 i = l; i <= r; i++)
                                                                      7.2 SCC
                   for (i32 j = i + 1; j <= r; j++) upd(P[i], P[j]);
                                                                      struct Scc{
                                                                        int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
              sort(P.begin() + l, P.begin() + r + 1, cmpy
                                                                        void init(int _n){
              return;
                                                                          n = n:
                                                                           for (int i=0; i<= n; i++)</pre>
                                                                             E[i].clear(), rE[i].clear();
          i32 m = (l + r) >> 1;
          auto midx = P[m].x;
                                                                        void addEdge(int u, int v){
          rec(l, m), rec(m + 1, r);
                                                                          E[u].PB(v); rE[v].PB(u);
          i32 tsz = 0;
          inplace_merge(P.begin() + l, P.begin() + m + 1,
                                                                        void DFS(int u){
          P.begin() + r + 1, cmpy);
for (i32 i = l; i <= r; i++) {
                                                                           vst[u]=1;
                                                                           for (auto v : E[u]) if (!vst[v]) DFS(v);
              if (abs(P[i].x - midx) * abs(P[i].x - midx)
                                                                           vec.PB(u);
                     >= ans) continue;
              for (i32 j = tsz - 1; j >= 0 && (P[i].y - t
    [j].y) * (P[i].y - t[j].y) < ans; j--)
    upd(P[i], t[j]);
t[tsz++] = P[i];</pre>
                                                                        void rDFS(int u){
                                                                           vst[u] = 1; bln[u] = nScc;
                                                                           for (auto v : rE[u]) if (!vst[v]) rDFS(v);
         }
                                                                        void solve(){
    };
                                                                          nScc = 0;
     sort(all(P));
                                                                           vec.clear();
     rec(0, P.size() - 1);
                                                                           fill(vst, vst+n+1, 0);
     return make_pair(sqrt(ans), ansi);
                                                                           for (int i=0; i<=n; i++)
}
                                                                             if (!vst[i]) DFS(i);
                                                                           reverse(vec.begin(),vec.end());
                                                                           fill(vst, vst+n+1, 0);
      graph
                                                                           for (auto v : vec)
7.1 BCC
                                                                             if (!vst[v]){
                                                                               rDFS(v); nScc++;
#define REP(i, n) for (int i = 0; i < n; i++)
struct BccVertex {
     int n, nScc, step, dfn[MXN], low[MXN];
vector<int> E[MXN], sccv[MXN];
                                                                     };
                                                                             支配樹
                                                                      7.3
     int top, stk[MXN];
     void init(int _n) {
                                                                      #define REP(i, s, e) for (int i = (s); i \leftarrow (e); i \leftarrow (e); i \leftarrow (e); i \leftarrow (e)
          n = _n;
                                                                      #define REPD(i, s, e) for (int i = (s); i >= (e); i--)
struct DominatorTree { // O(N) 1-base
          nScc = step = 0;
          for (int i = 0; i < n; i++) E[i].clear();</pre>
                                                                           int n, s;
                                                                          vector<int> g[MAXN], pred[MAXN];
vector<int> cov[MAXN];
     void addEdge(int u, int v) {
          E[u].PB(v);
          E[v].PB(u);
                                                                           int dfn[MAXN], nfd[MAXN], ts;
                                                                          int par[MAXN]; // idom[u] s到u的最後一個必經點int sdom[MAXN], idom[MAXN];
```

void DFS(int u, int f) {

```
int mom[MAXN], mn[MAXN];
                                                                     int popcount(const Int& val) { return val.count();
    inline bool cmp(int u, int v) { return dfn[u] < dfn</pre>
         [v]; }
                                                                     int lowbit(const Int& val) { return val._Find_first
    int eval(int u) {
                                                                         (); }
         if (mom[u] == u) return u;
                                                                     int ans, stk[MXN];
                                                                     int id[MXN], di[MXN], deg[MXN];
         int res = eval(mom[u]);
         if (cmp(sdom[mn[mom[u]]], sdom[mn[u]])) mn[u] =
                                                                     Int cans:
              mn[mom[u]];
                                                                     void maxclique(int elem_num, Int candi) {
         return mom[u] = res;
                                                                         if (elem_num > ans) {
                                                                             ans = elem_num;
                                                                             cans.reset();
for (int i = 0; i < elem_num; i++) cans[id[</pre>
     void init(int _n, int _s) {
         ts = 0;
         n = _n;
                                                                                  stk[i]]] = 1;
         s = _s;
         REP(i, 1, n) g[i].clear(), pred[i].clear();
                                                                         int potential = elem_num + popcount(candi);
                                                                         if (potential <= ans) return;</pre>
    void addEdge(int u, int v) {
   g[u].push_back(v);
                                                                         int pivot = lowbit(candi);
                                                                         Int smaller_candi = candi & (~linkto[pivot]);
         pred[v].push_back(u);
                                                                         while (smaller_candi.count() && potential > ans
    void dfs(int u) {
                                                                              int next = lowbit(smaller_candi);
         ts++;
                                                                              candi[next] = !candi[next];
         dfn[u] = ts;
                                                                              smaller_candi[next] = !smaller_candi[next];
         nfd[ts] = u;
                                                                              potential--;
         for (int v : g[u])
   if (dfn[v] == 0) {
      par[v] = u;
                                                                              if (next == pivot || (smaller_candi &
                                                                                  linkto[next]).count()) {
                                                                                  stk[elem_num] = next;
                 dfs(v);
                                                                                  maxclique(elem_num + 1, candi & linkto[
             }
                                                                                      next]);
                                                                             }
    void build() {
                                                                         }
         REP(i, 1, n) {
                                                                     int solve() {
    for (int i = 0; i < n; i++) {</pre>
             idom[i] = par[i] = dfn[i] = nfd[i] = 0;
             cov[i].clear();
                                                                              id[i] = i;
             mom[i] = mn[i] = sdom[i] = i;
                                                                              deg[i] = v[i].count();
         dfs(s);
         REPD(i, n, 2) {
    int u = nfd[i];
                                                                         sort(id, id + n, [&](int id1, int id2) { return
                                                                         deg[id1] > deg[id2]; });
for (int i = 0; i < n; i++) di[id[i]] = i;</pre>
             if (u == 0) continue;
for (int v : pred[u])
                                                                         for (int i = 0; i < n; i++)
                                                                             for (int j = 0; j < n; j++)
    if (v[i][j]) linkto[di[i]][di[j]] = 1;</pre>
                 if (dfn[v]) {
                      eval(v)
                      if (cmp(sdom[mn[v]], sdom[u])) sdom
                                                                         Int cand;
                                                                         cand.reset();
                           [u] = sdom[mn[v]];
                                                                         for (int i = 0; i < n; i++) cand[i] = 1;
             cov[sdom[u]].push_back(u);
                                                                         ans = 1;
             mom[u] = par[u];
                                                                         cans.reset();
             for (int w : cov[par[u]]) {
                                                                         cans[0] = 1;
                                                                         maxclique(0, cand);
                  eval(w)
                  if (cmp(sdom[mn[w]], par[u]))
                                                                         return ans;
                      idom[w] = mn[w];
                                                                } solver;
                      idom[w] = par[u];
                                                                7.5 最小圈
             cov[par[u]].clear();
                                                                /* minimum mean cycle O(VE) */
         REP(i, 2, n) {
                                                                struct MMC{
             int u = nfd[i];
                                                                #define E 101010
             if (u == 0) continue;
                                                                #define V 1021
             if (idom[u] != sdom[u]) idom[u] = idom[idom
                                                                #define inf 1e9
                                                                #define eps 1e-6
                                                                  struct Edge { int v,u; double c; };
         }
                                                                  int n, m, prv[V][V], prve[V][V], vst[V];
                                                                  Edge e[E];
} domT;
                                                                  vector<int> edgeID, cycle, rho;
                                                                  double d[V][V];
7.4 最大團
                                                                  void init( int _n )
                                                                  { n = _n; m = 0; }
// WARNING: TYPE matters
struct MaxClique { // 0-base
    typedef bitset<MXN> Int;
                                                                  void addEdge( int vi , int ui , double ci )
                                                                  { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
    Int linkto[MXN], v[MXN];
    int n;
                                                                    void init(int _n) {
```

for(int j=0; j<m; j++) {
 int v = e[j].v, u = e[j].u;
 if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {

d[i+1][u] = d[i][v]+e[j].c;

prv[i+1][u] = v;

prve[i+1][u] = j;

for (int i = 0; i < n; i++) {

void addEdge(int a, int b) { v[a][b] = v[b][a] = 1;

linkto[i].reset(); v[i].reset();

```
if (t >= 2) {
                                                                                        do \{b = rand() \% (p - 2) + 2;
   double solve(){
                                                                                        } while (mypow(b, p / 2, p) != p - 1);
     // returns inf if no cycle, mmc otherwise
                                                                                     pb = mypow(b, h, p);
} int s = mypow(a, h / 2, p);
     double mmc=inf;
     int st = -1;
                                                                                     for (int step = 2; step <= t; step++) {
  int ss = (((i64)(s * s) % p) * a) % p;
     bellman_ford();
     for(int i=0; i<n; i++) {
        double avg=-inf;
                                                                                        for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);</pre>
                                                                                     if (ss + 1 == p) s = (s * pb) % p;

pb = ((i64)pb * pb) % p;

} x = ((i64)s * a) % p; y = p - x;
        for(int k=0; k<n; k++) {</pre>
           if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i</pre>
                 1)/(n-k)):
           else avg=max(avg,inf);
                                                                                  } return true:
                                                                               }
        if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                               8.2 excrt
     fill(vst,0); edgeID.clear(); cycle.clear(); rho.
                                                                               i128 exgcd(i128 a, i128 b, i128 &x, i128 &y){
    if (b == 0) return x=1, y=0, a;
           clear():
     for (int i=n; !vst[st]; st=prv[i--][st]) {
        vst[st]++;
                                                                                     int d = exgcd(b, a \% b, y, x);
        edgeID.PB(prve[i][st]);
                                                                                     y -= a / b * x;
                                                                                     return d;
        rho.PB(st);
                                                                               // as -> 算式答案
     while (vst[st] != 2) {
        if(rho.empty()) return inf;
                                                                                // ns -> 模數 MOD
                                                                               i128 CRT(vector<i64> as, vector<i64> ns) {
        int v = rho.back(); rho.pop_back();
        cycle.PB(v);
                                                                                     i32 n = as.size();
i128 a1, a2, n1, n2;
        vst[v]++;
                                                                                     bool flag = false
                                                                                     auto china = [&]() {
     reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
                                                                                           i128 d = a2 - a1;
                                                                                           i128 x, y;
     return mmc;
                                                                                           i128 g = exgcd(n1, n2, x, y);
                                                                                           if (d % g == 0) {
    x = ((x * d / g) % (n2 / g) + (n2 / g)) % (
7.6 kShortestPath
                                                                                                     n2 / g);
while(Q.size()){
                                                                                                a1 = x * n1 + a1;
     auto [dx,x] = Q.top();Q.pop();
                                                                                                n1 = (n1 * n2) / g;
     if(dis[x].size() >= k) continue;
                                                                                           } else {
     dis[x].PB(dx)
                                                                                                flag = true;
     for(auto [v,w]:E[x]) Q.emplace(w+dx,v);
                                                                                     };
7.7 結論
                                                                                     a1 = as[0], n1 = ns[0];
                                                                                     for (i32 i = 1; i < n; i++) {
   • 2-SAT :
     (a_i \lor a_j) = true \ \forall (i,j)
對於任意限制 (x \lor y)
建兩條有向邊 (要多編號 \neg x)
                                                                                           a2 = as[i], n2 = ns[i];
                                                                                           china();
                                                                                           if (flag) return -1;
     x \rightarrow \neg y and y \rightarrow \neg x
     跑 scc
                                                                                     return a1;
     \operatorname{scc.bln}[x] < \operatorname{scc.bln}[\neg x] \Leftrightarrow x \text{ is true} \\ \operatorname{scc.bln}[\neg x] < \operatorname{scc.bln}[x] \Leftrightarrow x \text{ is false}
                                                                               }
     \exists x \text{ which scc.bln}[x] == \text{scc.bln}[\neg x] \Leftrightarrow \# M
                                                                               8.3 exgcd
   • 差分約束:
                                                                               int exgcd(int a,int b,int&x,int&y){
      n 個變數及 m 個約束條件
     求滿足所有 x_j - x_i \le b_k (i, j \in [1, n], k \in [1, m]) 的一組 x_1 ... x_n 可轉成 x_j - x_i \le b_k \to x_j \le x_i + b_k 結論就是 使得所有 x_j 變小以滿足上式
                                                                                     if(b==0)return x=1,y=0,a;
                                                                                     int d = exgcd(b,a\%b,y,x);
                                                                                     y=a/b*x;
                                                                                     return d;
      建邊跑 SPFA/Bellman
     要多建起點 s 連到所有 i 且邊權 0, \mathrm{dis}[s] = 0 有負環則無解,否則起點到所有 i 的距離為一組解
                                                                               }
     x_j - x_i \le k \Rightarrow \mathsf{addEdge}\ i \xrightarrow{k} j
                                                                                8.4 FFT
     x_j - x_i \ge k \Rightarrow \mathsf{addEdge} \ j \overset{-k}{\longrightarrow} i
                                                                               const int MAXN = 262144;
     x_j = x_i \Rightarrow \mathsf{addEdge}\ i \stackrel{0}{\longrightarrow} j \mathsf{and}\ j \stackrel{0}{\longrightarrow} i
                                                                                // (must be 2^k)
                                                                               // before any usage, run pre_fft() first
typedef long double ld;
8 math
                                                                                typedef complex<ld> cplx; //real() ,imag()
                                                                               const ld PI = acosl(-1);
const cplx I(0, 1);
8.1 DiscreteSqrt
void calcH(i64 &t, i64 &h, const i64 p) {
                                                                                cplx omega[MAXN+1];
  i64 tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
                                                                                void pre_fft(){
                                                                                  for(int i=0; i<=MAXN; i++)
  omega[i] = exp(i * 2 * PI / MAXN * I);</pre>
// solve equation x^2 \mod p = a
// !!!! (a != 0) !!!!!!
                                                                                // n must be 2^k
bool solve(i64 a, i64 p, i64 &x, i64 &y) {
  if(p == 2) { x = y = 1; return true; }
int p2 = p / 2, tmp = mypow(a, p2, p);
if (tmp == p - 1) return false;
                                                                               void fft(int n, cplx a[], bool inv=false){
                                                                                  int basic = MAXN / n;
                                                                                   int theta = basic;
  if ((p + 1) \% 4 == 0) {
                                                                                  for (int m = n; m >= 2; m >>= 1) {
```

x=mypow(a,(p+1)/4,p); y=p-x; return true;

i64 t, h, b, pb; calcH(t, h, p);

} else {

int mh = m >> 1;
for (int i = 0; i < mh; i++) {</pre>

cplx w = omega[inv ? MAXN-(i*theta%MAXN)]

```
: i*theta%MAXN];
                                                                                                • Bell 數 (有 n 個人, 把他們拆組的方法總數):
          for (int j = i; j < n; j += m) {</pre>
                                                                                                  B_0 = 1
B_n = \sum_{k=0}^{n} s(n, k) \quad (second - stirling)
B_{n+1} = \sum_{k=0}^{n} {n \choose k} B_k
            int k = j + mh;
            cplx x = a[j] - a[k];
            a[j] += a[k];
                                                                                                • Wilson's theorem :
            a[k] = w * x;
                                                                                                  (p-1)! \equiv -1 \pmod{p}
        }
      theta = (theta * 2) % MAXN;
                                                                                                • Fermat's little theorem :
                                                                                                  a^p \equiv a (mod \ p)
   int i = 0;
for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
                                                                                                • Euler's totient function:
                                                                                                   A^{B^C} mod \ p = pow(A, pow(B, C, p-1)) mod \ p
      if (j < i) swap(a[i], a[j]);</pre>
                                                                                                • 歐拉函數降冪公式: A^B \mod C = A^B \mod \phi(c) + \phi(c) \mod C
   if(inv) for (i = 0; i < n; i++) a[i] /= n;
                                                                                                • 6 的倍數:
                                                                                                   (a-1)^3 + (a+1)^3 + (-a)^3 + (-a)^3 = 6a
cplx arr[MAXN+1];
inline void mul(int _n,i64 a[],int _m,i64 b[],i64 ans
                                                                                                • Standard young tableau (標準楊表): \lambda = (\lambda_1 \geq \cdots \geq \lambda_k), \sum \lambda_i = n \text{ denoted by } \lambda \vdash n \\ \lambda \vdash n \text{ 意思為 } \lambda \text{ 整數拆分 } n \text{ eg. } n = 10, \lambda = (6,4) \text{ 此拆分可表示一種楊表}
      ]([]
   int n=1, sum=_n+_m-1;
   while(n<sum)</pre>
                                                                                                   形狀。
                                                                                                  楊表: 第 1 列 \lambda_1 行 \cdots 第 k 列 \lambda_k 行的方格圖。標準楊表: 每列從左到右遞增,每行從上到下遞增。
Let T 為某一 Permutation 跑 RSK 後的標準楊表,則此 Permutation 的 LDS、LIS 長度分別為 T 的列、行數。
     n << =1;
   for(int i=0;i<n;i++) {
    double x=(i<_n?a[i]:0),y=(i<_m?b[i]:0);
      arr[i]=complex<double>(x+y,x-y);
                                                                                                • RSK Correspondence:
   fft(n,arr);
                                                                                                  A permutation is bijective to (P,Q) 一對標準楊表 P: Permutation 跑 RSK 算法的結果,可為半標準楊表。
   for(int i=0;i<n;i++)</pre>
      arr[i]=arr[i]*arr[i];
                                                                                                   Q : 可用來還原 Permutation (像排列矩陣)。
   fft(n,arr,true);
for(int i=0;i<sum;i++)</pre>
                                                                                                • Hook length formula (形狀為 \lambda 的標準楊表個數):
                                                                                                  \begin{array}{l} f^{\lambda} = \frac{n!}{\prod h_{\lambda}(i,j)} \\ h_{\lambda}(i,j) = \text{number of pair } (x,y) \text{ where } (x=i \vee y=j) \wedge (x,y) \geq (i,j) \end{array}
      ans[i]=(i64)(arr[i].real()/4+0.5);
                                                                                                   且 (x,y) 落在形狀為 \lambda 的表上。
                                                                                                  Recursion: (i) f^{(0,\dots,0)} = 1
8.5 josephus
                                                                                                   (ii) f^{(\lambda_1,\dots,\lambda_m)} = \sum_{k=1}^m f^{(\lambda_1,\dots,\lambda_{k-1},\lambda_k-1,\lambda_{k+1},\dots,\lambda_m)}
int josephus(int n, int m){ //n人每m次
       int ans = 0;
      for (int i=1; i<=n; ++i)
                                                                                            8.7 Primes
            ans = (ans + m) \% i;
                                                                                                 Prime
                                                                                                                 Root
                                                                                                                          Prime
                                                                                                                                           Root
      return ans;
                                                                                                  7681
                                                                                                                 17
                                                                                                                          167772161
}
                                                                                                 12289
                                                                                                                 11
                                                                                                                          104857601
                                                                                                 40961
                                                                                                                          985661441
8.6
        Theorem
                                                                                                 65537
                                                                                                                          998244353
                                                                                                  786433
                                                                                                                10
                                                                                                                          1107296257
                                                                                                                                          10
    • Lucas's Theorem :
                                                                                                 5767169
                                                                                                                          2013265921
                                                                                                                                           31
      For n,m\in\mathbb{Z}^* and prime P, C(m,n) mod P=\Pi(C(m_i,n_i)) where
                                                                                                  7340033
                                                                                                                          2810183681
                                                                                                                                          11
       m_i is the i-th digit of m in base P.
                                                                                                  23068673
                                                                                                                          2885681153
                                                                                                 469762049
                                                                                                                3
                                                                                                                          605028353
   • Stirling approximation :
                                                                                            8.8 millerrabin
      n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\frac{1}{12n}}
   • Stirling Numbers(permutation |P|=n with k cycles): S(n,k)= coefficient of x^k in \Pi_{i=0}^{n-1}(x+i)
                                                                                            // n < 4,759,123,141
                                                                                                                                        3: 2, 7, 61
                                                                                                                                                2, 13, 23, 1662803
6: pirmes <= 13
                                                                                            // n < 1,122,004,669,633
                                                                                                                                        4:
                                                                                            // n < 3,474,749,660,383
   - Stirling Numbers(Partition \boldsymbol{n} elements into \boldsymbol{k} non-empty set):
                                                                                            // n < 2^{64}
      S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^{n}
                                                                                            // 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                                            // Make sure testing integer is in range [2, n-2] if
                                                                                            // you want to use magic.
   - Pick's Theorem : A=i+b/2-1 \\ A: Area, i: grid number in the inner, b: grid number on the side
                                                                                            bool witness(i64 a,i64 n,i64 u,int t){
   • Catalan number : C_n = \binom{2n}{n}/(n+1) C_n^{n+m} - C_{n+1}^{n+m} = (m+n)! \frac{n-m+1}{n+1} for n \geq m C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)!n!}
                                                                                               if(!a) return 0;
                                                                                               i64 x=mypow(a,u,n);
                                                                                               for(int i=0;i<t;i++) {</pre>
                                                                                                  i64 nx=mul(x,x,n);
      \begin{array}{lll} C_0 = 1 & and & C_{n+1} = 2(\frac{2n+1}{n+2})C_n \\ C_0 = 1 & and & C_{n+1} = \sum_{i=0}^n C_i C_{n-i} & for & n \geq 0 \end{array}
                                                                                                  if(nx==1\&&x!=1\&&x!=n-1) return 1;
                                                                                                  x=nx:
                                                                                               }
    • Euler Characteristic:
                                                                                               return x!=1;
      planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2
                                                                                            bool mii64er_rabin(i64 n) {
       V,E,F,C: number of vertices, edges, faces(regions), and compo-
      nents
                                                                                               // iterate s times of witness on n
    • Kirchhoff's theorem :
      A_{ii}=deg(i), A_{ij}=(i,j)\in E\ ?-1:0 , Deleting any one row, one column, and call the det(A)
                                                                                               if(n<2) return 0;</pre>
                                                                                               if(!(n\&1)) return n == 2;
                                                                                               i64 u=n-1; int t=0;
    \bullet Polya' theorem ( c is number of color, m is the number of cycle
                                                                                               // n-1 = u*2^t
       size):
                                                                                               while(!(u&1)) u>>=1, t++;
      (\sum_{i=1}^{\stackrel{\centerdot}{m}}c^{\gcd(i,m)})/m
                                                                                               while(s--){
   • Burnside lemma: |X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|
                                                                                                  i64 a=magic[s]%n;
                                                                                                  if(witness(a,n,u,t)) return 0;
                                                                                               return 1;
    • 錯排公式: (n 個人中, 每個人皆不再原來位置的組合數):
```

}

 $dp[0] = 1; \dot{dp}[1] = 0;$

dp[i] = (i-1) * (dp[i-1] + dp[i-2]);

8.9 phi

8.10 pollardrho

```
// does not work when n is prime O(n^(1/4))
i64 f(i64 x, i64 c, i64 mod){ return add(mul(x,x,mod),c
    ,mod); }
i64 poi64ard_rho(i64 n) {
    i64 c = 1, x = 0, y = 0, p = 2, q, t = 0;
    while (t++ % 128 or gcd(p, n) == 1) {
        if (x == y) c++, y = f(x = 2, c, n);
        if (q = mul(p, abs(x-y), n)) p = q;
        x = f(x, c, n); y = f(f(y, c, n), c, n);
    }
    return gcd(p, n);
}
```

8.11 primes

```
* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231
* 999888733, 98789101, 987777733, 999991921, 1010101333
* 1010102101, 1000000000039, 100000000000037
* 2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[ N ] , p_tbl[ N ];
vector<int> primes;
void sieve() {
  mu[1] = p_tbl[1] = 1;
  primes.push_back( i );
       mu[ i ] = -1;
    for( int p : primes ){
       int x = i *
                    р;
       if( x >= M ) break;
      p_tbl[ x ] = p;
mu[ x ] = -mu[ i ];
if( i % p == 0 ){
         mu[x] = 0;
         break;
vector<int> factor( int x ){
  vector<int> fac{ 1 };
  while (x > 1)
    int fn = SZ(fac), p = p_tbl[ x ], pos = 0;
    while( x \% p == 0){
      x /= p;

for( int i = 0 ; i < fn ; i ++ )

fac.PB( fac[ pos ++ ] * p );
  return fac;
```

8.12 Euler

```
int Euler(int n){
  int now = n;
  for (int i = 2; i * i <= n; i++)
    if (n % i == 0){
      now = now - now / i;
      while (n % i == 0) n = n / i;
      }
  if (n > 1) now = now - now / n;
  return now;
}
```

8.13 quickeuler

```
vector<int> pri;
bool not_prime[MXN + 10];
int phi [MXN + \bar{1}0];
void quick_euler(int n) {
    phi[1] = 1;
     for (int i = 2; i <= n; i++) {
         if (!not_prime[i]) {
              pri.push_back(i);
              phi[i] = i - 1;
         for (int pri_j : pri) {
    if (i * pri_j > n)
                   break;
              not_prime[i * pri_j] = true;
if (i % pri_j == 0) {
                   phi[i * pri_j] = phi[i] * pri_j;
                   break;
              phi[i * pri_j] = phi[i] * phi[pri_j];
         }
    }
}
```

8.14 sieve

9 other

9.1 cda

```
// 三維偏序 (求 arr[j] < arr[i] (每一維嚴格小於), i!=j
    j 的個數)
// 先照 x 排序 merge sort排y 最後BIT動態求z的順序個數
// 左區間的 x < 右區間的
void cdq(int ll,int rr){
    if(ll == rr) return;
    int m = (ll+rr)/2;
    cdq(ll,m),cdq(m+1,rr);
    int i = ll,j = m+1,t = 0;
auto work = [&](){
        ans += BIT.qry(arr[j].z); //計數
        temp[t++] = arr[j++];
    while(i <= m && j <= rr){</pre>
        if(arr[i].y <= arr[j].y){</pre>
             BIT.add(arr[i].z,1); //二維偏序求法
             temp[t++] = arr[i++];
        else work();
    while(i <= m) temp[t++] = arr[i++];</pre>
    while(j <= rr) work();
BIT.reset(); //操作復原
    rep(k,0,t) arr[k+ll] = temp[k];
//[l,r)
auto cdq = [&](auto&& self,auto l,auto r){
    if((r - l) <= 1) return;
auto m = (r - l) / 2 + l;
    self(self,l,m);
    self(self,m,r);
    auto i = l, j = m;
    auto work = [\&](){}
    while(i != m && j != r){
        if(arr[*i][1] <= arr[*j][1]) {
             ++i;
        }else work();
```

```
}
while(j != r) work();
clear();
inplace_merge(l,m,r,[&](auto a,auto b){
    return arr[a][1] < arr[b][1];
});
};
cdq(cdq,all(ord));//排ord</pre>
```

9.2 DeBruijnSequence

```
//求由所有 N 長度bitstring作為substring 最短的字串 B(2,
   N) //B(k,N): 以k個字元作為N長度字串節點
//00110 -> 00 01 11 10
//建圖 : 點為substrings 邊用 0 1 連接
//走訪: 000 -1-> 001
  解為 Hamiltonian 路徑 (剛好所有節點走過一遍)
// 可同構到 N-1 圖上的Eulerian Circuit (每條邊 N-1 圖上
    的邊 代表 N 圖上的一個點)
vector<int> edges[1<<(N-1)];</pre>
vector<int> ans;
void dfs(int x){ // Eulerian Circuit
   while(edges[x].size()){
       int u = edges[x].back();
       edges[x].pop_back();
       ans.push_back(u&1);
       dfs(u);
   }
void solve(int n){
   if(n == 1) {
       ans = \{1,0\};
       return:
    for(int i = 0; i < (1<<(n-1)); ++i){
       edges[i].push_back((i<<1)&((1<<(n-1))-1)); // 0
       edges[i].push_back(((i<<1)+1)&((1<<(n-1))-1));
   for(int i = 0; i < n-1; ++i) ans.push_back(0); //初
       始狀態
   dfs(0);
}
```

9.3 SmallestLexicographic

```
//對於可化作DAG的回朔問題求最小字典序的選擇
//建反圖 (反著做回來) (把以 i 結尾變成 以 i 開頭)
//結論 : i <- j (i < j) 取最小的 a[j]
for(int j = N; j; --j) {
    for(auto i:E[j])
    dp[i] = min(dp[i],dp[j]);
}</pre>
```

10 random

10.1 XORShift

```
const i64 mask = std::chrono::steady_clock::now().
    time_since_epoch().count();
//13 17 5
//13 17 7
i64 shift(i64 x) { // XOR shift (1-1 func)
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    x ^= mask;
    return x;
}</pre>
```

11 string

11.1 KMP

```
//pi[i] = 最大的 k 使得 s[0...(k-1)] = s[i-(k-1)...i]
vector<int> prefunc(const string& s){
  int n = s.size();
  vector<int> pi(n);
  for(int i=1,j=0;i<n;++i){
    j = pi[i-1];
    while(j && s[j] != s[i]) j = pi[j-1]; //取次小LCP</pre>
```

```
if(s[j] == s[i]) ++j;
  pi[i] = j;
}
return pi;
}
//找 s 在 str 中出現的所有位子
vector<int> kmp(string str, string s) {
  vector(int> nxt = prefunc(s);
  vector(int> ans;
  for (int i = 0, j = 0; i < SZ(str); i++) {
    while (j && str[i] != s[j]) j = nxt[j - 1];
    if (str[i] == s[j]) j++;
    if (j == SZ(s)) {
       ans.push_back(i - SZ(s) + 1);
       j = nxt[j - 1];
    }
  }
  return ans;
}</pre>
```

11.2 minRotation

```
// rotate(begin(s),begin(s)+minRotation(s),end(s))
#define rep(i, s, e) for (int i = (s); i < (e); i++)
int minRotation(string s) {
   int a = 0, N = s.size();
   s += s;
   rep(b, 0, N) rep(k, 0, N) {
      if (a + k == b || s[a + k] < s[b + k]) {
        b += max(0LL, k - 1);
        break;
      }
      if (s[a + k] > s[b + k]) {
        a = b;
        break;
      }
   return a;
}
```

11.3 PalindromeTree

```
// len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文子字串在整個字串中的出現次數
  fail[s]是他長度次長的回文後綴, aba的fail是a
// fail[s] -> s 建邊是顆樹
const int MXN = 1000010;
struct PalT{
  int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
  int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN] = \{-1\};
  int newNode(int 1,int f){
    len[tot]=l,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++;
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
```

```
tot=lst=n=0;
newNode(0,1),newNode(-1,1);
for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}
}palt;
```

11.4 RollingHash

```
struct RollingHash{
#define psz 2
     vector<ll> primes={17, 75577};
     vector<ll> MOD={998244353, 10000000007};
    vector<array<ll, psz>> hash, base;
void init(const string &s){
         hash.clear(); hash.resize(s.size());
base.clear(); base.resize(s.size());
for(int i=0;i<psz;i++){</pre>
              hash[0][i] = s[0];
              base[0][i] = 1;
          for(int i=1;i<s.size();i++){</pre>
              base[i][j] = base[i-1][j] * primes[j] %
              }
         }
    array<ll, psz> getHash(int_l,int r){
          if(l == 0) return hash[r];
         array<ll, psz> ret = hash[r];
for(int i=0;i<psz;i++){</pre>
              ret[i] -= hash[l-1][i] * base[r-l+1][i] %
                    MOD[i];
              if(ret[i]<0) ret[i]+=MOD[i];</pre>
         return ret;
}Hash;
```

11.5 SuffixArray

```
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
        hei[N], r[N];
  int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
    memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
    mkhei(n);
  void mkhei(int n){
    REP(i,n) r[\_sa[i]] = i;
    hei[0] = 0;
    REP(i,n) if(r[i]) {
  int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
       hei[r[i]] = ans;
    }
  void sais(int *s, int *sa, int *p, int *q, bool *t,
    int *c, int n, int z){
bool uniq = t[n-1] = true, neq;
    int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
          lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
    memcpy(x, c, sizeof(int) * z); \
    \label{eq:memcpy} \begin{array}{ll} \text{memcpy}(x + 1, \ c, \ sizeof(int) * (z - 1)); \\ \text{REP}(i,n) \ if(sa[i] \&\& \ !t[sa[i]-1]) \ sa[x[s[sa[i]-1]]) \end{array}
    ]-1]]++] = sa[i]-1; \
memcpy(x, c, sizeof(int) * z); \
     for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i
          MSO(c, z);
```

```
REP(i,n) uniq &= ++c[s[i]] < 2;
REP(i,z-1) c[i+1] += c[i];</pre>
                    if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
                   for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i
+1] ? t[i+1] : s[i]<s[i+1]);
                   MAGIC(\vec{R}EP1(i,1,\vec{n}-1) \ \vec{if(t[i]} \ \&\& \ !t[i-1]) \ sa[--x[s[i]] \ \vec{a} 
                                       ]]]=p[q[i]=nn++]=i);
                    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
                            neq=lst<0|lmemcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa]
                                                 [i])*sizeof(int));
                            ns[q[lst=sa[i]]]=nmxz+=neq;
                    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
                                             + 1);
                   MAGIC(for(int i = nn - 1; i \ge 0; i--) sa[--x[s[p[
                                       nsa[i]]]] = p[nsa[i]];
         }
}sa;
// H [i] 第 i 跟前面的最大共同前綴
// SA[i] 第 i 小是從第幾個字元開始
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
          // should padding a zero in the back
// ip is int array, len is array length
          // ip[0..n-1] != 0, and ip[len] = 0
          ip[len++] = 0;
          sa.build(ip, len, 128); // 注意字元個數 for (int i=0; i<len; i++) {
                   H[i] = sa.hei[i + 1];
                   SA[i] = sa.\_sa[i + 1];
           // resulting height, sa array \in [0,len)
```

11.6 trie

```
//01 bitwise trie
struct trie{
    trie *nxt[2]; // 差別
int cnt; //紀錄有多少個數字以此節點結尾
    int cnt;
               //有多少數字的前綴包括此節點
    int sz;
    trie():cnt(0),sz(0){
       memset(nxt,0,sizeof(nxt));
};
//創建新的字典樹
trie *root;
void insert(int x){
    trie *now = root; // 每次從根節點開始
    for(int i=22;i>=0;i--){ // 從最高位元開始往低位元走
       now->sz++;
       //cout<<(x>>i&1)<<endl;
       if(now->nxt[x>>i&1] == NULL){ //判斷當前第 i 個
           位元是 0 還是 1
           now->nxt[x>>i&1] = new trie();
       now = now->nxt[x>>i&1]; //走到下一個位元
   now->cnt++:
    now->sz++;
```

11.7 Z-algorithm

11.8 馬拉車

12 tree

12.1 DSUONTREE

```
int ans[MXN], color[MXN], son[MXN];
map<int, int> mp[MXN];
void dfs(int_x, int f){
    if(son[x]){
        dfs(son[x], x);
        swap(mp[x], mp[son[x]]);
        ans[x] = ans[son[x]];
    mp[x][color[x]]++;
    ans[x] = max(ans[x], mp[x][color[x]]);
    for(int i : edge[x]){
         if(i == f | i == son[x])
                                        continue;
        dfs(i, x);
        for(auto j : mp[i]){
             mp[x][j.first] += j.second;
             ans[x] = max(ans[x], mp[x][j.first]);
    }
}
```

12.2 EularTour

```
int timing=0;
int in[N],out[N];
void dfs(int u){
    in[u] = ++timing;//這時進入u
    for(int nxt : g[u]){//跑過所有孩子
        dfs(nxt);
    }
    out[u] = timing;//這時離開u 不會++
```

12.3 LCA

```
int n, q;
int anc[MAXN][25], in[MAXN], out[MAXN];
vector<int> edge[MAXN];
int timing = 1;
void dfs(int cur, int fa) {
    anc[cur][0] = fa;
    in[cur] = timing++;
for (int nex : edge[cur]) {
         if (nex == fa) continue;
         dfs(nex, cur);
    out[cur] = timing++;
void init() {
    dfs(1, 0);
    for (int i = 1; i < 25; i++) {
         for (int cur = 1; cur <= n; cur++) {</pre>
             anc[cur][i] = anc[anc[cur][i - 1]][i - 1];
    }
bool isanc(int u, int v) { return (in[u] <= in[v] &&</pre>
    out[v] <= out[u]); }
int lca(int a, int b) {
```

```
if (isanc(a, b)) return a;
     if (isanc(b, a)) return b;
     for (int i = 24; i >= 0; i--) {
          if (anc[a][i] == 0) continue;
          if (!isanc(anc[a][i], b)) a = anc[a][i];
     return anc[a][0];
}
int t = 0,tt = 0;
vector<int> dfn(n),in(n),out(n),dep(n);
vector anc(n,vector<int>(20));
auto pdfs = [&](auto &&self,int x,int f,int d = 0) ->
     void {
     in[x] = ++t;
     anc[\bar{x}][0] = f;
     dep[x] = d;
dfn[x] = ++tt;
     for(auto u:E[x]){
          if(u == f) continue;
          self(self,u,x,d+1);
     out[x] = ++t;
pdfs(pdfs,0,0);
for(int k = 1; k < 20;++k){
  for(int i = 0; i < n;++i){
    anc[i][k] = anc[anc[i][k-1]][k-1];</pre>
auto isanc = [&](int u,int v){
     return in[u] <= in[v] && out[v] <= out[u];</pre>
}:
auto lca = [\&](int x, int y){
     if(isanc(x,y)) return x;
     if(isanc(y,x)) return y;
for(int i = 19; i >= 0; --i){
          if(!isanc(anc[x][i],y)) x = anc[x][i];
     return anc[x][0];
};
```

12.4 treehash

```
map<vector<int>,int> id; //rooted
int dfs(int x,int f){
    vector<int> s;
    for(int u:E[x]){
         if(u == f) continue;
         s.PB(dfs(u,x));
    sort(all(s));
if(!id.count(s)) id[s] = id.size();
    return id[s];
}
const i64 mask = std::chrono::steady_clock::now().
    time_since_epoch().count();
//13 17 5
//13 17 7
i64 shift(i64 x) { // XOR shift (1-1 func)
  x ^= mask;
  x ^= x << 13;
  x \wedge = x \gg 7;
  x ^= x << 17;
  x \wedge = mask;
  return x:
}
int dfs(int x,int f){
   int ret = 1; // 需要常數
     for(int u:E[x]){
         if(u == f) continue;
         ret += shift(dfs(u,x));
    // ret ^= rand_mask //如果xor hash被卡
    return ret;
```

HeavyLightDecomposition

```
int t = 0;
vector\langle int \rangle dep(n+1),p(n+1),sz(n+1),dfn(n+1),son(n+1);
auto dfs = [&](auto &&self,int x,int f,int d = 0) ->
    void {
    ++sz[x],dep[x] = d,p[x] = f;
    for(auto u:E[x]){
        if(u == f) continue;
        self(self,u,x,d+1);
        sz[x] += sz[u];
        if(!son[x] | | sz[u] > sz[son[x]]) son[x] = u;
   }
vector<int> top(n+1);
auto dfsa = [&](auto &&self,int x,int f,int now) ->
    void {
    dfn[x] = ++t;
    top[x] = now;
    if(son[x]) self(self,son[x],x,now);
    for(auto u:E[x]){
   if(u == f || u == son[x]) continue;
        self(self,u,x,u);
   }
dfs(dfs,1,1);
dfsa(dfsa,1,1,1);
auto lca = [\&](int x, int y){
    while(top[x] != top[y]){
        if(dep[top[x]] < dep[top[y]]) swap(x,y);</pre>
        x = p[top[x]];
    return dep[x] < dep[y] ? x : y ;</pre>
// 如果要開線段樹 要每個鏈都開一顆 (比較快)
12.6 VirtualTree
//求關鍵點的虛樹
//thm1: 照dfn (dfs序) 排序後的 "相鄰點" 求lca可求出全
    點對的lca
auto virTree = [&](vector<int> key){
    auto cmp = [&](int a,int b){return dfn[a] < dfn[b</pre>
        ];};
    sort(all(key), cmp);
    auto res = vector<int>(all(key));
    for(int i = 1; i < key.size();++i){</pre>
        res.PB(lca(key[i-1],key[i]));
    sort(all(res), cmp);
    res.erase(unique(all(res)),res.end());
    return res; // res: 全點對lca集 + 關鍵點集
```

for(int i = 1; i < ret.size(); ++i){</pre>

virTree的邊 //query: 路徑詢問 //且會全部算到

int LCA = lca(ret[i-1],ret[i]); query(LCA,ret[i]); // 2. LCA -> ret[i] 是一條









