

1. (5 pts each, 30 pts total) Mark the following reasoning patterns as S (= sound, carries true premises to true conclusions) or U (= unsound, may carry true premises to false conclusions). Premises are shown above the line, conclusions below the line. Here, “ \Rightarrow ” means “implies” and “ \neg ” means “not.” The first one is done for you as an example.

a. S
$$\begin{array}{l} P \Rightarrow Q \\ \hline P \\ Q \end{array}$$

b. U
$$\begin{array}{l} P \Rightarrow Q \\ \hline Q \\ P \end{array}$$

c. U
$$\begin{array}{l} P \Rightarrow Q \\ \hline P \text{ or } \neg Q \end{array}$$

d. S
$$\begin{array}{l} P \Rightarrow \neg Q \\ \hline Q \\ \neg P \end{array}$$

e. S
$$\begin{array}{l} P \Rightarrow Q \\ \hline \neg Q \\ \neg P \end{array}$$

f. S
$$\begin{array}{l} P \Rightarrow Q \\ \hline \neg P \text{ or } Q \end{array}$$

g. U
$$\begin{array}{l} \neg P \Rightarrow Q \\ \hline P \\ \neg Q \end{array}$$

2. (5 pts each, 40 pts total) In each of the following, KB is a set of sentences, $\{\}$ is the empty set of sentences, and S is a single sentence. Recall \models means “entails” and \vdash means “derives,” where \vdash_i means “inference procedure i derives.” Use these keys:

Snd = Sound.

Unsnd = Unsound.

C = Complete.

I = Incomplete.

V = Valid.

Sat = Satisfiable.

Unsat = Unsatisfiable.

N = None of the above.

For each blank below, write in the key above that best corresponds to the correct term.

(a) Suppose some inference procedure i has the property, that for some KB and some S , $KB \models S$ but not $KB \vdash_i S$. Then the inference procedure i is I.

(b) Let S be given in advance. Suppose that for some KB_1 , $KB_1 \models S$; but that for some other KB_2 , $KB_2 \models \neg S$. Then S is Sat.

(c) Suppose some inference procedure i has the property, that for any KB and any S , whenever $KB \models S$ then $KB \vdash_i S$. Then the inference procedure i is C.

(d) Suppose inference procedure i has the property, that for some KB and some S , $KB \vdash_i S$ but not $KB \models S$. Then the inference procedure i is Unsnd.

(e) Let S be given in advance. Suppose that $\{\} \models S$. Then S is V.

(f) Suppose some inference procedure i has the property, that for any KB and any S , whenever $KB \vdash_i S$ then $KB \models S$. Then the inference procedure i is Snd.

(g) Suppose that $KB \models S$, then the sentence $(KB \Rightarrow S)$ is V.

(h) Suppose that $KB \models S$, then the sentence $(KB \text{ and } \neg S)$ is Unsat.

3. Consider the KB shown below.

a. (5 pts each, 15 pts total) Translate the following **KB** into Conjunctive Normal Form. The first one is done for you as an example (it was already in Conjunctive Normal Form ;-)).

A. $P \vee R$. $P \vee R$

B. $Q \Rightarrow S$. $\neg Q \vee S$

C. $P \Rightarrow Q$. $\neg P \vee Q$

D. $R \Rightarrow S$. $\neg R \vee S$

b. (15 pts total, -5 for each wrong step, but not negative. The order may vary, if proof is correct.) Write a complete resolution proof that $KB \models S$. Show the two clauses that you resolve in front of the symbol \vdash , and the resulting clause after \vdash . You may not require all of the lines provided. The sentence labeled "E." adds the negated goal. The first one is done for you as an example.

E. $\neg S$

(a) $\neg S$, $\neg Q \vee S$, \vdash $\neg Q$.

(b) $\neg Q$, $\neg P \vee Q$, \vdash $\neg P$.

(c) $\neg P$, $P \vee R$, \vdash R .

(d) R , $\neg R \vee S$, \vdash S .

(e) S , $\neg S$, \vdash \square .

Other proofs are fine if correct. For example, at step (d) above you could have resolved with $\neg S$:

(d) $\neg S$, $\neg R \vee S$, \vdash $\neg R$.

(e) $\neg R$, R , \vdash \square .