

Homework 1

ECE 590

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Question 1: Prove that:

$$\forall P, Q \in \mathbb{B}. P \rightarrow Q \leftrightarrow \neg P \vee Q$$

Proof is here:

Definiation: None

Lemmas and Axioms:

1. $A_{LEM} \forall P. \neg P \vee P$
2. *Demorgan's Laws* :
 - 2.1 $P \vee Q = \neg(\neg P \wedge \neg Q)$
 - 2.2 $P \wedge Q = \neg(\neg P \vee \neg Q)$
3. Contrapositive: $P \Rightarrow Q = \neg Q \Rightarrow \neg P$

Proof Goal:

$$\forall P, Q \in \mathbb{B}. P \rightarrow Q \leftrightarrow \neg P \vee Q$$

Step: 1. $\forall P, Q \in \mathbb{B}. P \rightarrow Q \leftrightarrow \neg P \vee Q$

Pick any P and Q, and...

1.1 $\forall P, Q \in \mathbb{B}. (P \rightarrow Q) \rightarrow \neg P \vee Q$

Assume $P \rightarrow Q$ and...

1.1.1 $\neg P \vee P$

By U_{spec} of A_{LEM} .

1.1.2 $\neg P \vee Q$

By cases on $P \vee \neg P$:

case 1: P :

1.1.2.1 Q

By MP on assumption 1.1.

1.1.2.2 $\neg P \vee Q$

By or-intro on case 1 and 1.2.1.1.

case 2: $\neg P$:

1.1.2.2 $\neg P \vee Q$

By or-intro on assumption 1.1.

1.2 $\forall P, Q \in \mathbb{B}. \neg P \vee Q \rightarrow (P \rightarrow Q)$

Assume $(\neg P \vee Q)$ and...

1.2.1 $P \rightarrow Q$

a. Assume P , and...

b. By contradiction, assume $\neg Q$ and...

1.2.1.1 $P \wedge \neg Q$

By and-intro on assumption 1.2.1.a and 1.2.1.b.

1.2.1.2 $\neg(\neg P \vee Q)$

By Demorgan's law on 1.2.1.1..

1.2.1.3 *False* (\perp)

Contradiction between 1.2.1.2 and assumption $(\neg P \vee Q)$.

1.3 $\forall P, Q \in \mathbb{B}. (P \rightarrow Q) \leftrightarrow \neg P \vee Q$

Iff-intro on 1.1 and 1.2.

Question 2: Prove that:

$$(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x)) \leftrightarrow \neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$$

Proof is here:

Defination: None

Lemmas and Axioms:

1. Demorgan's Laws:

$$1.1 \quad P \vee Q = \neg(\neg P \wedge \neg Q)$$

$$1.2 \quad P \wedge Q = \neg(\neg P \vee \neg Q)$$

Proof Goal:

$$(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x)) \leftrightarrow \neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$$

Step: 1. $(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x)) \leftrightarrow \neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$

Pick any $x \in \mathbb{Z}$, and...

1.1 $(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x)) \rightarrow \neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$

Assume $(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x))$ and...

1.1.1 $\neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$

By contradiction, assume $\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x))$

1.1.1.1 $\neg(f(a) \rightarrow g(a))$

By exists-elim on assumption 1.1.1.

1.1.1.1.1 $(f(a) \rightarrow g(a))$

By forall-elim on assumption 1.1.

1.1.1.1.2 *False* (\perp)

Contradiction between 1.1.1.1 and 1.1.1.1.1.

1.2 $\neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x))) \rightarrow (\forall x \in \mathbb{Z}. f(x) \rightarrow g(x))$

Assume $\neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$ and...

Pick any x , and...

1.2.1 $(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x))$

Assume $f(x)$, and...

By contradicton, assume $\neg g(x)$

1.2.1.1 $\neg(\exists x \in \mathbb{Z}. \neg(\neg f(x) \vee g(x)))$

Using meaning of implication on assumption 1.2.

1.2.1.2 $\neg(\exists x \in \mathbb{Z}. f(x) \wedge \neg g(x))$

Using Demorgan's Laws Based on 1.2.1.1

1.2.1.3 $f(x) \wedge \neg g(x)$

Using and-intro based on assumption 1.2.1

1.2.1.4 $\exists x \in \mathbb{Z}. f(x) \wedge \neg g(x)$

By exists intro on 1.2.1.3 with witness $x=x$.

1.2.1.5 *False* (\perp)

Contradiction between assumption 1.2.1.3. and 1.2.1.4.

1.3 $(\forall x \in \mathbb{Z}. f(x) \rightarrow g(x)) \leftrightarrow \neg(\exists x \in \mathbb{Z}. \neg(f(x) \rightarrow g(x)))$

Iff-intro on 1.1 and 1.2.

Question 3: Prove that:

$$\forall n \in \mathbb{N}. \sum_{i=0}^n i^3 = \frac{1}{4}n^2(n+1)^2$$

Proof is here:

Defination: None

Lemmas and Axioms: None

Proof Goal:

$$\forall n \in \mathbb{N}. \sum_{i=0}^n i^3 = \frac{1}{4}n^2(n+1)^2$$

Step: 1. $\forall n \in \mathbb{N}. \sum_{i=0}^n i^3 = \frac{1}{4}n^2(n+1)^2$

By weak induction on n . Cases:

Base: 0

Goal: $\sum_{i=0}^0 i^3 = \frac{1}{4}0^2(0+1)^2$

Proof: $0 = 0$ (trivial)

Base: 1.

Goal: $\sum_{i=0}^1 i^3 = \frac{1}{4}1^2(1+1)^2$

Proof: $1 = 1$ (trivial)

Ind.

IH: $\sum_{i=0}^n i^3 = \frac{1}{4}n^2(n+1)^2$

Goal: $\sum_{i=0}^{n+1} i^3 = \frac{1}{4}(n+1)^2((n+1)+1)^2$

Proof:

1.3.1

$$\sum_{i=0}^{n+1} i^3 = \sum_{i=0}^n i^3 + (n+1)^3$$

Definition of summation.

1.3.2

$$\sum_{i=0}^{n+1} i^3 = \frac{1}{4}n^2(n+1)^2 + (n+1)^3$$

Substitution of IH on 1.3.1.

1.3.3

$$\frac{1}{4}n^2(n+1)^2 + (n+1)^3 = \frac{1}{4}(n+1)^2((n+1)+1)^2$$

Because $\frac{1}{4}n^2(n+1)^2 + (n+1)^3$

$$= (n+1)^2 \left[\frac{1}{4}n^2 + (n+1) \right]$$

$$= \frac{1}{4}(n+1)^2(n^2 + 4n + 4)$$

$$= \frac{1}{4} (n+1)^2 ((n+1)+1)^2$$

Question 4: Prove that:

$$\forall A, B, C \in P(\mathbb{Z}). A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$$

Proof is here:

Definiation:

1. $\forall A, B \in P(\mathbb{Z}). A \subseteq B$ is equivalent to $\forall x. (x \in A \rightarrow x \in B)$

Lemmas and Axioms: None

Proof Goal:

$$\forall A, B, C \in P(\mathbb{Z}). A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$$

Step: 1. $\forall A, B, C \in P(\mathbb{Z}). A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$

Pick any A, B and C, and...

1.1 $A \subseteq C$

Assume $A \subseteq B \wedge B \subseteq C$ and...

1.1.1 $\forall x. (x \in A \rightarrow x \in C)$

Pick any x and...

1.1.1.1 $x \in A \rightarrow x \in C$

Using Definition 1 and universal-elim on 1.1.1.

Assume $x \in A$.

1.1.1.1.1 $A \subseteq B$

Using and-elim on assumption 1.1.

1.1.1.1.2 $x \in A \rightarrow x \in B$

Using Definition 1 and universal-elim on 1.1.1.1.1.

1.1.1.1.3 $x \in B$

By implication-elim with assumption 1.1.1.1 on 1.1.1.1.2.

1.1.1.1.4 $B \subseteq C$

Using and-elim on assumption 1.1.

1.1.1.1.5 $x \in B \rightarrow x \in C$

Using Definition 1 and universal-elim on 1.1.1.1.4.

1.1.1.1.6 $x \in C$

By implication-elim with 1.1.1.1.3 on 1.1.1.1.5.

Question 5:

Definiation:

1. D_i : Dragons on the island.

2. $ged(D_i)$: If dragon D_i is green eyed, $ged(D_i)$ is true, otherwise is false.

(a) The oracle's proposition: 'There is at least one dragon on this island with green eyes ' means that:

$$\exists D_i. ged(D_i)$$

(b) This particular dragon(D_0)'s proof is here:

Axioms:

1. at least one dragon is green-eyed:

$$\exists D_i. ged(D_i)$$

2. all dragons except D_0 are not green-eyed:

$$\forall D_i. (\neg (D_i = D_0) \rightarrow \neg ged(D_i))$$

3. $A \vee (B \wedge \neg B) \leftrightarrow A$

Proof Goal:

$$ged(D_0)$$

Step: 1. $ged(D_0)$

1.1 $ged(D_x)$

By exists-elim on assumption axiom 1.

1.2 $\forall D_i. (D_i = D_0 \vee \neg ged(D_i))$

By meaning of implication on axiom 2.

1.3 $D_x = D_0 \vee \neg ged(D_x)$

By universal-elim on 1.2.

$$1.4 \quad D_x = D_0 \vee \neg ged(D_x) \wedge ged(D_x)$$

By and-intro on 1.1 and 1.3

$$1.5 \quad D_x = D_0$$

Using axiom 3 on 1.4.

$$1.6 \quad ged(D_0)$$

By MP on 1.1 and 1.5.

(c) Under the scenario above (only D_0 has green eyes), what happens to the red eyed dragons?:

At 1st midnight, only one dragon turn into butterfly, so all red eyed dragons know that that dragon is the only green-eyed one, so they all turn into butterflies at next midnight.

(d)

- i All green-eyed dragons will turn into butterflies on the Nth day, and all red-eyed dragons will turn into butterflies on the (N+1)thday.

$$\text{ii } \forall n \in \mathbb{N}. |\{d | ged(d)\}| = n \rightarrow (\forall D_i. (i < n \rightarrow bfly(D_i) = n) \wedge (i \geq n \rightarrow bfly(D_i) = n + 1))$$

iii Proof:

Axioms:

- 1. at least one dragon is green-eyed:

$$\exists D_i. ged(D_i)$$

- 2. all dragons except n dragons ($n > 1$) are not green-eyed:

$$\forall D_i. (\neg((D_i = D_0) \vee (D_i = D_1) \vee \dots \vee (D_i = D_{n-1})) \rightarrow \neg ged(D_i))$$

and it is equivalent to

$$\forall D_i. ((D_i = D_n) \vee (D_i = D_{n+1}) \vee \dots \vee (D_i = D_{99}) \rightarrow \neg ged(D_i))$$

- 3. $\neg(A \vee B) \rightarrow C \leftrightarrow (\neg A \rightarrow C) \vee (\neg B \rightarrow C)$

- 4. If there are n green-eyed dragon, then there won't be any dragon turning into butterfly before nth midnight:

$$\forall n \in \mathbb{N}. |\{d | ged(d)\}| = n \rightarrow \forall D_i. \neg(bfly(D_i) = n - 1)$$

Step: 1. $\forall n \in \mathbb{N}. |\{d|ged(d)\}| = n \rightarrow$
 $(\forall D_i. (i < n \rightarrow bfly(D_i) = n) \wedge (i \geq n \rightarrow bfly(D_i) = n + 1))$
 Assume $\forall n \in \mathbb{N}. |\{d|ged(d)\}| = n.$
 By weak induction on n . Cases:

Base: 2

Goal: $(\forall D_i. (i < 2 \rightarrow bfly(D_i) = 2) \wedge (i \geq 2 \rightarrow bfly(D_i) = 3))$

Proof:

1.2.1 $\forall D_i. (i < 2 \rightarrow bfly(D_i) = 2)$

Assume $i < 2$, and...

1.2.1.1 $\forall D_i. (\neg((D_i = D_0) \vee (D_i = D_1)) \rightarrow \neg ged(D_i))$

By axiom 2, and...

1.2.1.2 $\forall D_i. ((D_i = D_0) \vee (D_i = D_1) \vee \neg ged(D_i))$

By meaning of implication on 1.2.1.1.

Assume $\forall D_i. (D_i = D_0) \rightarrow ged(D_i).$

1.2.1.3 $\neg(bfly(D_0) = 1)$

By axiom 4.

1.2.1.4 $ged(D_1)$

Based on 1.2.1.3

1.2.1.5 $bfly(D_0) = 2 \wedge bfly(D_1) = 2$

Based on assumption 1.2.1.2 and 1.2.1.4.

1.2.2 $\forall D_i. (i \geq 2 \rightarrow bfly(D_i) = 3)$

Assume $i \geq 2$, and...

1.2.2.1 $\forall D_i. ((D_i = D_n) \vee (D_i = D_{n+1}) \vee \dots \vee (D_i = D_{99}) \rightarrow \neg ged(D_i))$

By axiom 2, and...

1.2.2.2 $\forall D_i. ((D_i = D_n) \rightarrow \neg ged(D_i)) \wedge \forall D_i. ((D_i = D_{n+1}) \rightarrow \neg ged(D_i))$

$\wedge \dots \wedge \forall D_i. ((D_i = D_{99}) \rightarrow \neg ged(D_i))$

By axiom 3.

1.2.2.3 $\neg ged(D_n) \wedge \neg ged(D_{n+1}) \wedge \dots \wedge \neg ged(D_{99})$

Based on the proof in (b).

1.2.2.4 $bfly(D_n) = 3 \wedge bfly(D_{n+1}) = 3 \wedge \dots \wedge bfly(D_{99}) = 3$

Based on 1.2.1.5 and 1.2.2.3.

Ind.

IH: $\forall n \in \mathbb{N}. |\{d|ged(d)\}| = n \rightarrow$

$(\forall D_i. (i < n \rightarrow bfly(D_i) = n) \wedge (i \geq n \rightarrow bfly(D_i) = n + 1))$

Proof: $\forall n \in \mathbb{N}. |\{d|ged(d)\}| = n + 1 \rightarrow$

$$(\forall D_i. (i < n + 1 \rightarrow bfly(D_i) = n + 1) \wedge (i \geq n + 1 \rightarrow bfly(D_i) = n + 1 + 1))$$

1.3.1 $\forall D_i. (i < n + 1 \rightarrow bfly(D_i) = n + 1)$

Assume $i < n + 1$, and...

1.3.1.1 $\forall D_i. (\neg((D_i = D_0) \vee (D_i = D_1) \vee \dots \vee (D_i = D_n)) \rightarrow \neg ged(D_i))$

By axiom 2, and...

1.3.1.2 $\forall D_i. ((D_i = D_0) \vee (D_i = D_1) \vee \dots \vee (D_i = D_n) \vee \neg ged(D_i))$

By meaning of implication on 1.3.1.1.

Assume $\forall D_i. (D_i = D_0) \vee (D_i = D_1) \vee \dots \vee (D_i = D_{n-1}) \rightarrow ged(D_i)$.

1.3.1.3 $\neg(bfly(D_0) = n)$

By axiom 4.

1.3.1.4 $ged(D_n)$ Based on 1.2.1.3

1.3.1.5 $bfly(D_0) = n + 1 \wedge bfly(D_1) = n + 1 \wedge \dots \wedge bfly(D_n) = n + 1$

1.3.2 $\forall D_i. (i \geq n \rightarrow bfly(D_i) = n + 1)$

Assume $i \geq n$, and...

1.3.2.1 $\forall D_i. ((D_i = D_n) \vee (D_i = D_{n+1}) \vee \dots \vee (D_i = D_{99}) \rightarrow \neg ged(D_i))$

By axiom 2, and...

1.3.2.2 $\forall D_i. ((D_i = D_n) \rightarrow \neg ged(D_i)) \wedge \forall D_i. ((D_i = D_{n+1}) \rightarrow \neg ged(D_i)) \wedge \dots \wedge \forall D_i. ((D_i = D_{99}) \rightarrow \neg ged(D_i))$

By axiom 3.

1.3.2.3 $\neg ged(D_n) \wedge \neg ged(D_{n+1}) \wedge \dots \wedge \neg ged(D_{99})$

Based on the proof in (b).

1.3.2.4 $bfly(D_n) = n + 1 \wedge bfly(D_{n+1}) = n + 1 \wedge \dots \wedge bfly(D_{99}) = n + 1$

Based on 1.3.1.5 and 1.3.2.3.