图论

树的重心

```
int n, siz[N], wei[N];
std::vector< int > centroid, adj[N];
void GetCentroid(int x, int p) {
   siz[x] = 1;
   wei[x] = 0;
   for (const auto \&y : adj[x]) {
        if (y ^ p) {
            GetCentroid(y, x);
            siz[x] += siz[y];
            wei[x] = std::max(wei[x], siz[y]);
        }
    }
   wei[x] = std::max(wei[x], n - siz[x]);
    if (wei[x] \le n / 2) {
        centroid.push_back(x);
    }
}
```

最小树形图

有向图上的最小生成树(Directed Minimum Spanning Tree)称为最小树形图。

常用的算法是朱刘算法(也称 Edmonds 算法),可以在时间 $\Theta(nm)$ 内解决最小树形图问题。

Kruskal 重构树

原图中两个点之间的所有简单路径上最大边权的最小值 = 最小生成树上两个点之间的简单路径上的最大值 = $\mathrm{Kruskal}$ 重构树上两点之间的 LCA 的权值。

```
struct Edge {
   int u, v, w;
   bool operator < (const Edge &o) const {
      return w < o.w;
   }
};
struct DSU {
   std::vector< int > f;
   DSU(int n) : f(n + 1) {
      std::iota(begin(f), end(f), 0);
}
```

```
int find(int x) {
        return f[x] == x ? x : f[x] = find(f[x]);
   }
};
std::vector< int > adj[N];
struct Kruskal {
   std::vector< Edge > edges;
   int root;
   int val[N];
   void AddEdge(int u, int v, int w) {
        edges.push_back({u, v, w});
   void solve(int n) {
        std::sort(begin(edges), end(edges));
        root = n;
        DSU dsu(n * 2 - 1);
        for (auto &g : edges) {
            int fx = dsu.find(g.u), fy = dsu.find(g.v);
            if (fx != fy) {
                dsu.f[fx] = dsu.f[fy] = ++ root;
                val[root] = g.w;
                adj[root] = \{fx, fy\};
            }
       }
   }
};
```

LCA

```
template< typename T >
struct INFO {
   int v;
   Tw;
};
// index from 1 to n
template< typename T >
struct Tree {
   int n, cur;
   std::vector< int > dfn, pos, dep, fa;
   std::vector< T > dis;
   std::vector< std::vector< INFO< T > > adj;
   Tree(int _n) {
       n = _n;
        cur = 0;
        dfn = std::vector < int > (_n * 2);
        pos = dep = fa = std::vector< int > (_n + 1);
```

```
dis = std::vector < T > (n + 1);
        adj = std::vector< std::vector< INFO< T > > ( n + 1);
    }
   void AddEdge(int u, int v, T w = 1) {
        adj[u].push_back({v, w});
    }
   void dfs(int x, int p) {
        fa[x] = p;
        dep[x] = dep[p] + 1;
        dfn[++ cur] = x;
        pos[x] = cur;
        for (const auto \&[y, z] : adj[x]) {
            if (y ^ p) {
                dis[y] = dis[x] + z;
                dfs(y, x);
                dfn[++ cur] = x;
            }
       }
    }
    std::vector< int > LOG;
    std::vector< std::vector< int > > RMQ;
   void Build() {
        LOG = std::vector< int > (cur + 1);
        for (int i = 2; i <= cur; ++ i) {
            LOG[i] = LOG[i / 2] + 1;
        RMQ = std::vector< std::vector< int > > (cur + 1, std::vector< int > (LOG[cur]
+ 1));
        for (int i = 1; i <= cur; ++ i) {
            RMQ[i][0] = pos[dfn[i]];
        for (int j = 1; j <= LOG[cur]; ++ j) {</pre>
            for (int i = 1; i + (1 << j) - 1 <= cur; ++ i) {
                RMQ[i][j] = std::min(RMQ[i][j-1], RMQ[i+(1 << (j-1))][j-1]);
        }
    }
    int LCA(int x, int y) { // O(1) query lca
       x = pos[x];
        y = pos[y];
        if (x > y) {
            std::swap(x, y);
        int k = LOG[y - x + 1];
        return dfn[std::min(RMQ[x][k], RMQ[y - (1 << k) + 1][k])];
    }
```

```
int Distance(int x, int y) {
    return dis[x] + dis[y] - 2 * dis[LCA(x, y)];
}
```

判断负环

Spfa

时间复杂度 $\Theta(n imes m)$

```
bool vis[N];
int d[N], cnt[N];
vector< pair< int, int > > adj[N];
bool spfa(int n) {
    for (int i = 1; i <= n; ++ i) {
        adj[0].push_back({i, 0});
    }
    memset(d, -0x3f, sizeof d);
    d[0] = 0;
    vis[0] = true;
    queue< int > q;
    q.push(0);
    while (!q.empty()) {
       int u = q.front();
        q.pop();
        vis[u] = false;
        for (const auto \&[v, w] : adj[u]) {
            if (d[v] < d[u] + w) {
                d[v] = d[u] + w;
                if (!vis[v]) {
                    q.push(v);
                    if (++ cnt[v] >= n) {
                        return false;
                    }
                }
            }
        }
    }
   return true;
}
```

Bellman-Ford

时间复杂度 $\Theta(n^2 imes m)$

```
int d[N];
vector< pair< int, int > > adj[N];
bool Bellman_Ford(int n) {
   for (int t = 1; t < n; ++ t) {
        bool upd = false;
        for (int i = 1; i <= n; ++ i) {
            for (const auto \&[v, w] : adj[i]) {
                if (d[i] > d[v] + w) {
                    d[i] = d[v] + w;
                    upd = true;
                }
            }
        }
        if (!upd) {
            break;
        }
    for (int i = 1; i <= n; ++ i) {
        for (const auto \&[v, w] : adj[i]) {
            if (d[i] > d[v] + w) {
                return false;
        }
    }
   return true;
}
```

最小环

给出一个图, 问其中的有 n 个节点构成的边权和最小的环 $(n \geq 3)$ 是多大。

图的最小环也称围长。

Dijkstra

枚举所有边,每一次求删除一条边之后对这条边的起点跑一次最短路,ans $= \min(ans, dis(u, v) + w)$ 时间复杂度: $\Theta(m(m+n)\log n)$

Floyd

```
ans=\min(ans,w(i,k)+w(k,j)+dis(i,j)) dis(i,j)=\min(dis(i,j),dis(i,k)+dis(k,j)) 时间复杂度:\Theta(\mathrm{n}^3)
```

连通性相关

无向图

割边以及边双连通分量e-DCC的缩点

```
struct Edge {
   int from, to;
   Edge(int u, int v) : from(u), to(v) {}
};
struct e_DCC {
   int n, m, cnt, cur;
   std::vector< Edge > edges;
   std::vector< std::vector< int > > G;
   std::vector< int > dfn, low, col;
   std::vector< bool > bridge;
   e DCC(int n) : n(n), cnt(0), cur(0), G(n), dfn(n), low(n), col(n) {}
   void AddEdge(int u, int v) {
        edges.push_back(Edge(u, v));
        G[u].push_back((int)edges.size() - 1);
       m = (int)edges.size();
    }
   void tarjan(int x, int p) {
        dfn[x] = low[x] = ++ cur;
        for (const auto \&i : G[x]) {
            int y = edges[i].to;
            if (y == p) {
                continue;
            }
            if (!dfn[y]) {
                tarjan(y, x);
                low[x] = std::min(low[x], low[y]);
                if (low[y] > dfn[x]) {
                    bridge[i] = bridge[i ^ 1] = true;
                }
            } else {
```

```
low[x] = std::min(low[x], dfn[y]);
            }
       }
    }
   void dfs(int x) {
        col[x] = cnt;
        for (const auto &i : G[x]) {
            int y = edges[i].to;
            if (col[y] | bridge[i]) {
                continue;
            }
            dfs(y);
       }
    }
   void solve() {
        bridge.assign(m, 0);
        for (int i = 0; i < n; ++ i) {
            if (!dfn[i]) {
                tarjan(i, i);
        }
        for (int i = 0; i < n; ++ i) {
            if (!col[i]) {
                ++ cnt;
                dfs(i);
        }
        for (int i = 1; i < edges.size(); ++ i) {</pre>
            int x = edges[i ^ 1].to, y = edges[i].to;
            if (col[x] != col[y]) {
                // Add_New_Edge
            }
        }
   }
};
```

割点和点双连通分量v-DCC的缩点

```
struct v_DCC {
   int cur, cnt, root, top;
   std::vector< int > dfn, low, col, stk;
   std::vector< bool > cut;
   std::vector< std::vector< int > > adj;

   v_DCC(int n) : cur(0), cnt(0), top(0), dfn(n), low(n), col(n), stk(n), cut(n),
   adj(n) {}
```

```
void tarjan(int x) {
        dfn[x] = low[x] = ++ cur;
        stk[top ++] = x;
        if (x == root && adj[x].empty()) { // 孤立点
            ++ cnt;
            col[x] = cnt;
            return;
        int flag = 0;
        for (const auto &y : adj[x]) {
            if (!dfn[y]) {
               tarjan(y);
                low[x] = std::min(low[x], low[y]);
                if(low[y] >= dfn[x]) {
                    ++ flag;
                    if (x != root | flag > 1) {
                        cut[x] = true;
                    ++ cnt;
                    int z;
                    do {
                       z = stk[--top];
                        col[z] = cnt;
                    } while (z != y);
                    col[x] = cnt;
                }
            } else {
                low[x] = std::min(low[x], dfn[y]);
            }
       }
   }
};
```

有向图找强连通分量

```
struct SCC {
   int cur, cnt, top;
   std::vector< int > dfn, low, col, stk;
   std::vector< std::vector< int > > adj;

SCC(int n) : cur(0), cnt(0), top(0), dfn(n), low(n), col(n), stk(n), adj(n) {}

void tarjan(int x) {
   dfn[x] = low[x] = ++ cur;
   stk[top ++] = x;
   for (const auto &y : adj[x]) {
      if (!dfn[y]) {
```

```
tarjan(y);
                low[x] = std::min(low[x], low[y]);
            } else if (!col[y]) {
                low[x] = std::min(low[x], dfn[y]);
            }
        }
        if (low[x] == dfn[x]) {
            ++ cnt;
            int z;
            do {
                z = stk[--top];
                col[z] = cnt;
            } while (z != x);
       }
   }
};
```

欧拉图

```
struct Euler {
   int n, m, top, t, tot;
   std::vector< int > head, ver, next, ans;
   std::stack< int > stk;
   std::vector< bool > vis;
   Euler(int n, int m) : n(n), m(m), head(n), ver(m), next(n), ans(n), vis(m) \ \{\}
   void add_edge(int x, int y) {
        ver[++ tot] = y, next[y] = head[x], head[x] = tot;
    }
   void euler() {
        stk.push(1);
        while (top) {
            int x = stk.top(), i = head[x];
            while (i && vis[i]) i = next[i];
            if (i) {
                stk.push(ver[i]);
                vis[i] = vis[i ^ 1] = true;
                head[x] = next[i];
            } else {
                stk.pop();
                ans[++t] = x;
            }
        }
    }
   void show() {
        for (int i = t; i; -- i) {
            std::cout << ans[i] << '\n';
        }
```

```
};
```

网络流

最大流

Dinic

算法原理

```
对于二分图\Theta(m\sqrt{n})
```

最坏情况 $\Theta(n^2m)$

```
constexpr int INF = 0x3f3f3f3f;
template< typename T >
struct Edge {
   int from, to;
   T cap, flow;
   Edge(int u, int v, T c, T f) : from(u), to(v), cap(c), flow(f) {}
};
template< typename T >
struct Dinic {
   int n, m;
   int Source, Sink;
   std::vector< Edge< T > > edges;
   std::vector< std::vector< int > > G;
   std::vector< int > d, cur;
   std::vector< bool > vis;
    // number of vertices
   Dinic(int n) : n(n), d(n), cur(n), vis(n), G(n) \{ \}
   void AddEdge(int from, int to, T cap) {
        edges.push_back(Edge< T >(from, to, cap, 0));
        edges.push_back(Edge< T >(to, from, 0, 0));
        m = edges.size();
        G[from].push_back(m - 2);
        G[to].push_back(m - 1);
    }
   bool BFS() {
        vis.assign(n, 0);
        std::queue< int > Q;
        Q.push(Source);
        d[Sink] = 0;
        vis[Source] = true;
```

```
while (!Q.empty()) {
            int x = Q.front();
            Q.pop();
            for (int i = 0; i < G[x].size(); ++ i) {
                Edge< T > &e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = true;
                    d[e.to] = d[x] + 1;
                    Q.push(e.to);
               }
            }
       return vis[Sink];
   }
   T DFS(int x, T a) {
        if (x == Sink | | a == 0) {
           return a;
        }
        T flow = 0, f;
        for (int& i = cur[x]; i < G[x].size(); ++ i) {
            Edge< T > \&e = edges[G[x][i]];
            if (d[x] + 1 == d[e.to] & (f = DFS(e.to, std::min(a, e.cap - e.flow))) >
0) {
                e.flow += f;
                edges[G[x][i] ^ 1].flow -= f;
                flow += f;
                a -= f;
                if (a == 0) {
                    break;
                }
            }
       return flow;
    }
    T Maxflow(int s, int t) {
        this->Source = s;
        this->Sink = t;
        T flow = 0;
        while (BFS()) {
            cur.assign(n, 0);
            flow += DFS(Source, INF);
        }
       return flow;
   }
};
```

最小割

最大流最小割定理

```
f(s,t)_{max} = c(s,t)_{min}
```

输出方案

首先根据最小割最大流定理,我们跑一遍Dinic就可以求出最小割,这是参量网络中 s 和 t 已经不再联通了.我们可以从 s 开始跑一遍 dfs ,沿着所有还未满流的边搜索,所有能到达的节点就是和 s 再同一集合的节点.之后我们遍历每一条边,将边起点终点不在同一集合内的输出即可

输出方案即输出除源点外的S集合,这里有一个待研究的性质:**dinic最后一次bfs后,vis为1的节点集合即为最小割的S集合**.

最小费用最大流

Dinic

```
constexpr int INF = 0x3f3f3f3f;
template< typename T >
struct Edge {
   int from, to;
   T flow, cost;
   Edge(int from, int to, T flow, T cost) : from(from), to(to), flow(flow), cost(cost)
{}
};
template< typename T >
struct Dinic {
   int n, m;
   int Source, Sink;
   std::vector< Edge< T > > edges;
   std::vector< std::vector< int > > G;
   std::vector< T > incf, d;
   std::vector< int > pre;
   std::vector< bool > vis;
    // number of vertices
   Dinic(int n) : n(n), G(n), incf(n), d(n), pre(n), vis(n) \{ \}
   void AddEdge(int from, int to, T flow, T cost) {
        edges.push_back(Edge< T >(from, to, flow, cost));
        edges.push back(Edge< T >(to, from, 0, -cost));
        m = edges.size();
        G[from].push_back(m - 2);
        G[to].push_back(m - 1);
    }
   bool SPFA() {
```

```
d.assign(n, INF);
        incf.assign(n, 0);
        vis.assign(n, false);
        std::queue< int > Q;
        Q.push(Source);
        vis[Source] = true;
        incf[Source] = INF;
        d[Source] = 0;
        while (!Q.empty()) {
            int x = Q.front();
            Q.pop();
            vis[x] = false;
            for (int i = 0; i < G[x].size(); ++ i) {
                Edge< T > &e = edges[G[x][i]];
                if (d[x] + e.cost < d[e.to] && e.flow) {
                    d[e.to] = d[x] + e.cost;
                    pre[e.to] = G[x][i];
                    incf[e.to] = std::min(incf[x], e.flow);
                    if (!vis[e.to]) {
                        vis[e.to] = true;
                        Q.push(e.to);
                }
            }
        }
        return d[Sink] != INF;
    }
    std::pair< T, T > MCMF(int s, int t) {
        this->Source = s;
        this->Sink = t;
        T maxflow = 0, mincost = 0;
        while (SPFA()) {
            maxflow += incf[Sink];
            mincost += incf[Sink] * d[Sink];
            for (int i = Sink; i != Source; i = edges[pre[i] ^ 1].to) {
                edges[pre[i]].flow -= incf[Sink];
                edges[pre[i] ^ 1].flow += incf[Sink];
            }
        }
        return {maxflow, mincost};
    }
};
```

重链剖分

```
struct Tree {
  int dfn[N], siz[N], dep[N], son[N], f[N], top[N];
```

```
vector< int > adj[N];
   void dfs1(int x, int p) {
        dep[x] = dep[p] + 1;
        f[x] = p;
        siz[x] = 1;
        for (auto &y : adj[x]) {
            if (y ^ p) {
                dfs1(y, x);
                siz[x] += siz[y];
                if (siz[y] > siz[son[x]]) {
                    son[x] = y;
            }
        }
   void dfs2(int x, int tp) {
        top[x] = tp;
        dfn[x] = ++ dfn[0];
        if (!son[x]) {
            return;
        dfs2(son[x], tp);
        for (auto &y : adj[x]) {
            if (y != son[x] && y != f[x]) {
                dfs2(y, y);
            }
        }
    int lca(int x, int y) {
        while (top[x] != top[y]) {
            if (dep[top[x]] < dep[top[y]]) {</pre>
                swap(x, y);
            x = f[top[x]];
       return dfn[x] < dfn[y] ? x : y;</pre>
   }
};
```

长链剖分

还没学,鸽...

虚树

```
template< typename T >
struct rmq {
   vector< T > v;
```

```
int n;
          static const int b = 30;
          vector<int> mask, t;
          int op(int x, int y) {
                    return v[x] < v[y] ? x : y;
          }
          int msb(int x) { return __builtin_clz(1) - __builtin_clz(x); }
          rmq() {}
          rmq(const \ vector < T > \&v_) : v(v_), n(v.size()), mask(n), t(n) {
                     for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
                                at = (at << 1) & ((1 << b) - 1);
                               while (at and op(i, i - msb(at \& -at)) == i) at ^= at & -at;
                     for (int i = 0; i < n / b; i++) t[i] = b * i + b - 1 - msb(mask[b * i + b - 1 - msb(mask
1]);
                    for (int j = 1; (1 << j) <= n / b; j++)
                                for (int i = 0; i + (1 << j) <= n / b; <math>i++)
                                          t[n / b * j + i] = op(t[n / b * (j - 1) + i], t[n / b * (j - 1) + i + i]
(1 << (j - 1))]);
          }
          int small(int r, int sz = b) { return r - msb(mask[r] \& ((1 << sz) - 1)); }
          T query(int 1, int r) {
                     if (r - 1 + 1 \le b) return small(r, r - 1 + 1);
                     int ans = op(small(1 + b - 1), small(r));
                     int x = 1 / b + 1, y = r / b - 1;
                     if (x \le y) {
                               int j = msb(y - x + 1);
                                ans = op(ans, op(t[n / b * j + x], t[n / b * j + y - (1 << j) + 1]));
                    return ans;
        }
};
namespace lca {
          vector<int> g[N];
          int v[2 * N], pos[N], dep[2 * N];
          int t;
          rmq<int> RMQ;
          void dfs(int i, int d = 0, int p = -1) {
                     v[t] = i, pos[i] = t, dep[t++] = d;
                     for (int j : g[i])
                               if (j != p) {
```

```
dfs(j, d + 1, i);
                v[t] = i, dep[t++] = d;
            }
    }
   void build(int n, int root) {
        t = 0;
        dfs(root);
        RMQ = rmq < int > (vector < int > (dep, dep + 2 * n - 1));
    }
    int lca(int a, int b) {
        a = pos[a], b = pos[b];
        return v[RMQ.query(min(a, b), max(a, b))];
    }
    int dist(int a, int b) {
        return dep[pos[a]] + dep[pos[b]] - 2 * dep[pos[lca(a, b)]];
    }
}
vector<int> virt[N];
int build virt(vector<int> v) {
    auto cmp = [&](int i, int j) { return lca::pos[i] < lca::pos[j]; };</pre>
   sort(v.begin(), v.end(), cmp);
   for (int i = v.size() - 1; i; i--) v.push_back(lca::lca(v[i], v[i - 1]));
   sort(v.begin(), v.end(), cmp);
   v.erase(unique(v.begin(), v.end()), v.end());
   for (int i : v) virt[i].clear();
   for (int i = 1; i < v.size(); i++) {
        virt[lca::lca(v[i - 1], v[i])].push_back(v[i]);
   return v[0];
}
```

2-SAT

```
struct TWO_SAT {
   int head[N], ver[N], Next[N];
   int dfn[N], low[N], col[N];
   int n, m, tot, num, cnt;
   stack< int > st;
   bool vis[N];
   void tarjan(int x) {
      dfn[x] = low[x] = ++ num;
      st.push(x);
```

```
for (int i = head[x]; i; i = Next[i]) {
            int y = ver[i];
            if (!dfn[y]) {
                tarjan(y);
                low[x] = min(low[x], low[y]);
            } else if (!col[y]) {
                low[x] = min(low[x], dfn[y]);
            }
        if (low[x] == dfn[x]) {
            col[x] = ++ cnt;
            while (st.top() != x) {
                col[st.top()] = cnt;
                st.pop();
            }
            st.pop();
        }
    }
   bool check() {
        for (int i = 1; i \le n; ++ i) {
            if (col[i] == col[i + n]) return false;
        return true;
    }
   void find() {
        vector< int > ret;
        for (int i = 1; i <= n; ++ i) {
            int x = col[i], y = col[i + n];
            if (vis[x]) {
                ret.push_back(1); continue;
            if (vis[y]) {
                ret.push_back(0); continue;
            if (x < y) {
                ret.push_back(1); vis[x] = true;
            } else {
                ret.push_back(0); vis[y] = true;
            }
       }
    }
};
```

图的匹配

二分图最大匹配

找增广路 $\Theta(nm)$

```
struct augment path {
   std::vector< std::vector< int > > g;
   std::vector< int > pa; // 匹配
   std::vector< int > pb;
   std::vector< int > vis; // 访问
               // 两个点集中的顶点数量
   int n, m;
                    // 时间戳记
   int dfn;
   int res;
                    // 匹配数
   augment_path(int_n, int_m) : n(_n), m(_m) {
       assert(0 <= n && 0 <= m);
       pa = std::vector < int > (n, -1);
       pb = std::vector< int >(m, -1);
       vis = std::vector< int >(n);
       g.resize(n);
       res = 0;
       dfn = 0;
    }
   void AddEdge(int from, int to) {
       assert(0 <= from && from < n && 0 <= to && to < m);
       g[from].push back(to);
   }
   bool dfs(int v) {
       vis[v] = dfn;
       for (int u : g[v]) {
           if (pb[u] == -1) {
               pb[u] = v;
              pa[v] = u;
               return true;
           }
       for (int u : g[v]) {
           if (vis[pb[u]] != dfn && dfs(pb[u])) {
               pa[v] = u;
               pb[u] = v;
               return true;
           }
       }
       return false;
    }
   std::pair< int, std::vector< int > > find_max_unweighted_matching() {
       while (true) {
```

```
## dfn;
int cnt = 0;
for (int i = 0; i < n; ++ i) {
    if (pa[i] == -1 && dfs(i)) {
        ++ cnt;
    }
}
if (cnt == 0) {
    break;
}
res += cnt;
}
return {res, pa};
}
</pre>
```

二分图最大权匹配

Hungarian Algorithm (Kuhn-Munkres Algorithm)

 $\Theta(n^3)$

```
template< typename T >
struct hungarian { // km
   int n;
   std::vector< int > matchx; // 左集合对应的匹配点
   std::vector< int > matchy; // 右集合对应的匹配点
   std::vector< int > pre; // 连接右集合的左点
   std::vector< bool > visx; // 拜访数组 左
   std::vector< bool > visy; // 拜访数组 右
   std::vector< T > lx;
   std::vector< T > ly;
   std::vector< std::vector< T > > g;
   std::vector< T > slack;
   T inf;
   T res;
   std::queue< int > q;
   int org_n;
   int org m;
   hungarian(int _n, int _m) {
       org_n = _n;
       org_m = _m;
       n = std::max(_n, _m);
       inf = std::numeric limits< T >::max();
       res = 0;
       g = std::vector< std::vector< T > >(n, std::vector< T >(n));
       matchx = std::vector< int >(n, -1);
       matchy = std::vector< int >(n, -1);
```

```
pre = std::vector< int >(n);
       visx = std::vector< bool >(n);
       visy = std::vector< bool >(n);
       lx = std::vector < T > (n, -inf);
       ly = std::vector< T >(n);
       slack = std::vector< T >(n);
   }
   void AddEdge(int u, int v, T w) {
       g[u][v] = std::max(w, 0); // 负值还不如不匹配 因此设为0不影响
   }
   bool check(int v) {
       visy[v] = true;
       if (matchy[v] != -1) {
           q.push(matchy[v]);
           visx[matchy[v]] = true; // in S
           return false;
       }
       // 找到新的未匹配点 更新匹配点 pre 数组记录着"非匹配边"上与之相连的点
       while (v != -1) {
           matchy[v] = pre[v];
           std::swap(v, matchx[pre[v]]);
       return true;
   }
   void bfs(int i) {
       while (!q.empty()) {
           q.pop();
       }
       q.push(i);
       visx[i] = true;
       while (true) {
           while (!q.empty()) {
               int u = q.front();
               q.pop();
               for (int v = 0; v < n; ++ v) {
                   if (!visy[v]) {
                       T delta = lx[u] + ly[v] - g[u][v];
                       if (slack[v] >= delta) {
                           pre[v] = u;
                           if (delta) {
                               slack[v] = delta;
                           } else if (check(v)) { // delta=0 代表有机会加入相等子图 找增广
路
                               // 找到就return 重建交错树
                               return;
                           }
```

```
}
            }
        }
        // 没有增广路 修改顶标
        T a = inf;
        for (int j = 0; j < n; ++ j) {
            if (!visy[j]) {
                a = std::min(a, slack[j]);
            }
        }
        for (int j = 0; j < n; ++ j) {
            if (visx[j]) { // S
                lx[j] = a;
            if (visy[j]) { // T
                ly[j] += a;
            } else { // T'
                slack[j] -= a;
            }
        for (int j = 0; j < n; ++ j) {
            if (!visy[j] && slack[j] == 0 && check(j)) {
                return;
            }
        }
   }
}
std::pair< T, std::vector< int > > find max weighted matching() {
    // 初始顶标
    for (int i = 0; i < n; ++ i) {
        for (int j = 0; j < n; ++ j) {
            lx[i] = std::max(lx[i], g[i][j]);
        }
    }
    for (int i = 0; i < n; ++ i) {
        fill(slack.begin(), slack.end(), inf);
        fill(visx.begin(), visx.end(), false);
        fill(visy.begin(), visy.end(), false);
       bfs(i);
    }
   // custom
    for (int i = 0; i < n; ++ i) {
        if (g[i][matchx[i]] > 0) {
            res += g[i][matchx[i]];
        } else {
```

```
matchx[i] = -1;
}

/*

std::cout << res << "\n";

for (int i = 0; i < org_n; i++) {
    std::cout << matchx[i] + 1 << " \n"[i + 1 == org_n];
}

*/

return {res, matchx};
}
};</pre>
```

一般图最大匹配

开花算法(Blossom Algorithm,也被称做带花树)

```
// graph
template< typename T >
class graph {
public:
   struct edge {
        int from;
        int to;
        T cost;
   };
   std::vector< edge > edges;
   std::vector< std::vector< int > > g;
   int n;
   graph(int _n) : n(_n) {
       g.resize(n);
   }
   virtual int add(int from, int to, T cost) = 0;
};
// undirectedgraph
template< typename T >
class undirectedgraph : public graph< T > {
public:
   using graph< T >::edges;
   using graph< T >::g;
   using graph< T >::n;
   undirectedgraph(int _n) : graph< T >(_n) {}
```

```
int AddEdge(int from, int to, T cost = 1) {
       assert(0 <= from && from < n && 0 <= to && to < n);
       int id = (int) edges.size();
       g[from].push_back(id);
       g[to].push_back(id);
       edges.push_back({from, to, cost});
       return id;
   }
};
// blossom / find max unweighted matching
template< typename T >
std::vector< int > find_max_unweighted_matching(const undirectedgraph< T > &g) {
   std::mt19937 rng(std::chrono::steady clock::now().time since epoch().count());
   std::vector< int > match(g.n, -1); // 匹配
   std::vector< int > aux(g.n, -1);
                                      // 时间戳记
                                      // "o" or "i"
   std::vector< int > label(g.n);
                                      // 花根
   std::vector< int > orig(g.n);
   std::vector< int > parent(g.n, -1); // 父节点
   std::queue< int > q;
   int aux time = -1;
   auto lca = [\&](int v, int u) {
       aux time++;
       while (true) {
           if (v != -1) {
               if (aux[v] == aux_time) { // 找到拜访过的点 也就是LCA
                   return v;
               aux[v] = aux_time;
               if (match[v] == -1) {
                   v = -1;
               } else {
                   v = orig[parent[match[v]]]; // 以匹配点的父节点继续寻找
               }
           }
           std::swap(v, u);
       }
   }; // lca
   auto blossom = [\&](int v, int u, int a) {
       while (orig[v] != a) {
           parent[v] = u;
           u = match[v];
           if (label[u] == 1) { // 初始点设为"o" 找增广路
               label[u] = 0;
               q.push(u);
           }
```

```
orig[v] = orig[u] = a; // 缩花
       v = parent[u];
   }
}; // blossom
auto augment = [\&](int v) {
   while (v != -1) {
       int pv = parent[v];
       int next v = match[pv];
       match[v] = pv;
       match[pv] = v;
       v = next_v;
   }
}; // augment
auto bfs = [&](int root) {
   fill(label.begin(), label.end(), -1);
   iota(orig.begin(), orig.end(), 0);
   while (!q.empty()) {
       q.pop();
   q.push(root);
   // 初始点设为 "o", 这里以"0"代替"o", "1"代替"i"
   label[root] = 0;
   while (!q.empty()) {
       int v = q.front();
       q.pop();
       for (int id : g.g[v]) {
           auto &e = g.edges[id];
           int u = e.from ^ e.to ^ v;
           if (label[u] == -1) { // 找到未拜访点
               label[u] = 1;  // 标记 "i"
               parent[u] = v;
               if (match[u] == -1) { // 找到未匹配点
                                      // 寻找增广路径
                   augment(u);
                  return true;
               }
               // 找到已匹配点 将与她匹配的点丢入queue 延伸交错树
               label[match[u]] = 0;
               q.push(match[u]);
               continue;
           } else if (label[u] == 0 && orig[v] != orig[u]) {
               // 找到已拜访点 且标记同为"o" 代表找到"花"
               int a = lca(orig[v], orig[u]);
               // 找LCA 然后缩花
               blossom(u, v, a);
               blossom(v, u, a);
           }
       }
```

```
return false;
   }; // bfs
   auto greedy = [\&]() {
       std::vector< int > order(g.n);
       // 随机打乱 order
       iota(order.begin(), order.end(), 0);
       shuffle(order.begin(), order.end(), rng);
       // 将可以匹配的点匹配
       for (int i : order) {
           if (match[i] == -1) {
               for (auto id : g.g[i]) {
                   auto &e = g.edges[id];
                   int to = e.from ^ e.to ^ i;
                   if (match[to] == -1) {
                       match[i] = to;
                       match[to] = i;
                       break;
                   }
               }
           }
   }; // greedy
   // 一开始先随机匹配
   greedy();
   // 对未匹配点找增广路
   for (int i = 0; i < g.n; i++) {
       if (match[i] == -1) {
           bfs(i);
       }
   }
   return match;
}
```

数据结构

并查集

维护联通性

```
struct UnionFind {
   std::vector< int > f, siz;
   UnionFind(int n) : f(n), siz(n, 1) {
        iota(begin(f), end(f), 0);
    int find(int x) {
        return f[x] == x ? x : f[x] = find(f[x]);
   bool Merge(int x, int y) { // 按秩合并
        x = find(x), y = find(y);
       if (x != y) {
            if (siz[x] < siz[y]) {
                std::swap(x, y);
            siz[x] += siz[y];
            f[y] = x;
        return false;
    }
    int operator [] (int x) {
       return find(x);
    }
};
```

带权并查集

```
int find(int x) {
    if (f[x] == x) {
        return x;
    }
    int rt = find(f[x]);
    val[x] += val[f[x]];
    return f[x] = rt;
}
```

可持久化并查集

给定 n 个集合,第 i 个集合内初始状态下只有一个数,为 i。

有m次操作。操作分为3种:

- 1 a b 合并 a, b 所在集合;
- $2 ext{ k}$ 回到第 k 次操作(执行三种操作中的任意一种都记为一次操作)之后的状态;

• 3 a b 询问 a, b 是否属于同一集合,如果是则输出 1,否则输出 0。

```
#include<bits/stdc++.h>
using namespace std;
const int N = 2e5 + 5;
int n, m;
struct node {
   int val, rnk, ls, rs; // 父节点、树的深度、左右儿子编号
} pt[N << 5];</pre>
int a[N], rt[N], tot;
int build(int 1, int r) {
   int now = ++ tot;
   if (1 == r) {
        pt[now].val = a[1];
        pt[now].rnk = 1;
       return now;
    }
   pt[now].ls = build(1, mid);
   pt[now].rs = build(mid + 1, r);
   return now;
int update(int x, int 1, int r, int tarjet, int newroot) {
   int now = ++ tot;
   pt[now] = pt[x];
   if (1 == r) {
        pt[now].val = newroot;
        pt[newroot].rnk = max(pt[newroot].rnk, pt[now].rnk + 1);
       return now;
    }
    if (tarjet <= mid) pt[now].ls = update(pt[x].ls, 1, mid, tarjet, newroot);</pre>
   else pt[now].rs = update(pt[x].rs, mid + 1, r, tarjet, newroot);
   return now;
int query(int x, int 1, int r, int tarjet) {
   if (l == r) return pt[x].val;
   if (tarjet <= mid) return query(pt[x].ls, l, mid, tarjet);</pre>
   else return query(pt[x].rs, mid + 1, r, tarjet);
}
int getroot(int x, int 1, int r, int tarjet) {
    int ans = query(x, 1, r, tarjet);
   if (ans == tarjet) return ans;
   else return getroot(x, 1, r, ans);
}
signed main() {
   cin.tie(nullptr)->sync_with_stdio(false);
   cin >> n >> m;
   for (int i = 1; i \le n; ++ i) a[i] = i;
```

```
rt[0] = build(1, n);
    int op, a, b, k;
   for (int i = 1; i <= m; ++ i) {
        cin >> op;
        if (op == 1) {
            cin >> a >> b;
            int root_a = getroot(rt[i - 1], 1, n, a);
            int root_b = getroot(rt[i - 1], 1, n, b);
            if (pt[root a].rnk < pt[root b].rnk) {</pre>
                rt[i] = update(rt[i - 1], 1, n, root_a, root_b);
            } else {
                rt[i] = update(rt[i - 1], 1, n, root_b, root_a);
        } else if (op == 2) {
            cin >> k;
            rt[i] = rt[k];
        } else {
            cin >> a >> b;
            rt[i] = rt[i - 1];
            cout << (getroot(rt[i], 1, n, a) == getroot(rt[i], 1, n, b)) << '\n';</pre>
        }
    }
   return 0;
}
```

单调栈

第i个元素之后第一个大于 a_i 的元素的**下标**

```
#include<bits/stdc++.h>
using LL = long long;

signed main() {
    std::cin.tie(nullptr)->sync_with_stdio(false);

    int n;
    std::cin >> n;
    std::vector< int > a(n), ans(n);
    for (auto &x : a) {
        std::cin >> x;
    }
    std::stack< int > st;
    for (int i = n - 1; i >= 0; -- i) {
        while (!st.empty() && a[st.top()] <= a[i]) {
            st.pop();
        }
}</pre>
```

```
ans[i] = st.empty() ? -1 : st.top();
    st.push(i);
}
for (int i = 0; i < n; ++ i) {
    std::cout << ans[i] + 1 << " \n"[i + 1 == n];
}
return 0;
}</pre>
```

单调队列

区间长度为k的最大最小值

```
#include<bits/stdc++.h>
using LL = long long;
signed main() {
    std::cin.tie(nullptr)->sync_with_stdio(false);
   int n, k;
   std::cin >> n >> k;
    std::vector< int > a(n), Max(n), Min(n);
    for (auto \&x : a) {
        std::cin >> x;
   std::deque< int > p, q;
    for (int i = 0; i < n; ++ i) {
        while (!q.empty() && a[q.back()] >= a[i]) {
            q.pop_back();
        q.push_back(i);
        if (q.back() - q.front() \ge k) {
            q.pop_front();
        if (i >= k - 1) {
            Min[i] = a[q.front()];
        }
        while (!p.empty() && a[p.back()] < a[i]) {</pre>
            p.pop_back();
        p.push_back(i);
        if (p.back() - p.front() >= k) {
            p.pop_front();
        if (i >= k - 1) {
            Max[i] = a[p.front()];
```

```
}
}
for (int i = k - 1; i < n; ++ i) {
    std::cout << Min[i] << " \n"[i + 1 == n];
}
for (int i = k - 1; i < n; ++ i) {
    std::cout << Max[i] << " \n"[i + 1 == n];
}
return 0;
}</pre>
```

ST表

```
int lg[N];
struct Sparse_Table {
   vector< vector< int > > f;
   Sparse\_Table(int n) : f(n + 1, vector < int > (lg[n] + 1, 0)) \{ \}
   void Build(int n, int *a) {
        for (int i = 2; i <= n; ++ i) {
            lg[i] = lg[i >> 1] + 1;
        for (int i = 1; i <= n; ++ i) {
            f[i][0] = a[i];
        for (int j = 1; j \le lg[n]; ++ j) {
            for (int i = 1; i + (1 << j) - 1 <= n; ++ i) {
                f[i][j] = max(f[i][j-1], f[i+(1 << (j-1))][j-1]);
            }
        }
   }
    int Query(int 1, int r) {
        int k = \lg[r - 1 + 1];
        return \max(f[1][k], f[r - (1 << k) + 1][k]);
    }
};
```

树状数组

一维树状数组

单点修改,区间查询 OR 区间修改,单点查询

```
template< typename T >
struct FenWick {
```

```
int n;
   std::vector< T > c;
   FenWick(int n) : n(n), c(n + 1) \{ \}
   void add(int i, T d) {
        for (; i \le n; i += i \& -i) {
           c[i] += d;
       }
   }
   void add(int 1, int r, T d) {
       add(1, d);
        add(r + 1, -d);
   T get(int i) {
        T sum = 0;
        for (; i; i -= i & -i) {
           sum += c[i];
        }
       return sum;
   }
    T get(int 1, int r) {
       return get(r) - get(l - 1);
   }
};
```

区间修改,区间查询

```
template< typename T >
struct Fenwick {
   int n;
   std::vector< T > c1, c2;
   Fenwick(int n) : n(n), c1(n + 1), c2(c1) {}
   void add(int i, T d) {
       for (int j = i; i \le n; i += i \& -i) {
            c1[i] += d;
            c2[i] += j * d;
        }
    }
   void add(int 1, int r, T d) {
       add(1, d);
       add(r + 1, -d);
    }
    T get(int i) {
       T sum = 0;
        for (int j = i; i; i = i & -i) {
            sum += (j + 1) * c1[i] - c2[i];
        }
        return sum;
    T get(int 1, int r) {
```

```
return get(r) - get(l - 1);
}
};
```

二维树状数组

单点修改,区间查询OR区间修改,单点查询

```
template< typename T >
struct Fenwick {
   int n, m;
   std::vector< vector< T > > c;
   Fenwick(int n, int m) : n(n), m(m), c(n + 1), vector< T > (m + 1)) {}
   void add(int x, int y, T d) {
        for (; x \le n; x += x & -x) {
            for (; y \le m; y += y & -y) {
                c[x][y] += d;
        }
   }
    void add(int x1, int y1, int x2, int y2, T d) {
        add(x1, y1, d);
        add(x1, y2 + 1, -d);
        add(x2 + 1, y1, -d);
        add(x2 + 1, y2 + 1, d);
    }
    T get(int x, int y, T sum = 0) {
        for (; x; x = x \& -x) {
            for (; y; y -= y & -y) {
                sum += c[x][y];
        }
        return sum;
    T get(int x1, int y1, int x2, int y2) {
       return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1) + get(x1 - 1, y1 - 1);
};
```

区间修改,区间查询(左上角为(x1,y1),右下角为(x2,y2))

```
template< typename T >
struct Fenwick {
   int n, m;
   std::vector< vector< T > > c1, c2, c3, c4;
   Fenwick(int n, int m) : n(n), m(m), c1(n + 1, vector< T > (m + 1)), c2(c1), c3(c1),
c4(c1) {}
   void add(int x, int y, T d) {
```

```
for (int i = x; i \le n; i += i \& -i) {
            for (int j = y; j \le m; j += j \& -j) {
                c1[i][j] += d;
                c2[i][j] += x * d;
                c3[i][j] += y * d;
                c4[i][j] += x * y * d;
            }
        }
    void add(int x1, int y1, int x2, int y2, T d) {
        add(x1, y1, d);
        add(x1, y2 + 1, -d);
        add(x2 + 1, y1, -d);
        add(x2 + 1, y2 + 1, d);
    T get(int x, int y, T sum = 0) {
        for (int i = x; i; i = i & -i) {
            for (int j = y; j; j == j \& -j) {
                sum += (x + 1) * (y + 1) * c1[i][j];
                sum = (y + 1) * c2[i][j];
                sum = (x + 1) * c3[i][j];
                sum += c4[i][j];
            }
        }
       return sum;
   }
    T get(int x1, int y1, int x2, int y2) {
        return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1) + get(x1 - 1, y1 - 1);
    }
};
```

二叉搜索树&平衡树

旋转Treap

```
#include <bits/stdc++.h>
using namespace std;
struct Node {
    Node *ch[2];
    int val, rank;
    int rep_cnt;
    int siz;
    Node(int val) : val(val), rep_cnt(1), siz(1) {
        ch[0] = ch[1] = nullptr;
        rank = rand();
    }
    void upd_siz() {
        siz = rep_cnt;
    }
}
```

```
if (ch[0] != nullptr) siz += ch[0]->siz;
        if (ch[1] != nullptr) siz += ch[1]->siz;
   }
};
class Treap {
private:
   Node *root;
   enum rot_type {
       LF = 1, RT = 0
   };
   int q_prev_tmp = 0, q_nex_tmp = 0;
   void _rotate(Node *&cur, rot_type dir) { // 0为右旋, 1为左旋
        Node *tmp = cur->ch[dir];
        cur->ch[dir] = tmp->ch[!dir];
        tmp->ch[!dir] = cur;
        tmp->upd_siz(), cur->upd_siz();
        cur = tmp;
   void _insert(Node *&cur, int val) {
        if (cur == nullptr) {
            cur = new Node(val);
            return;
        } else if (val == cur->val) {
            cur->rep cnt++;
            cur->siz++;
        } else if (val < cur->val) {
            _insert(cur->ch[0], val);
            if (cur->ch[0]->rank < cur->rank) {
                _rotate(cur, RT);
            cur->upd_siz();
        } else {
            _insert(cur->ch[1], val);
            if (cur->ch[1]->rank < cur->rank) {
                _rotate(cur, LF);
            cur->upd_siz();
        }
    }
   void _del(Node *&cur, int val) {
        if (val > cur->val) {
            _del(cur->ch[1], val);
            cur->upd siz();
        } else if (val < cur->val) {
            _del(cur->ch[0], val);
            cur->upd siz();
        } else {
            if (cur->rep_cnt > 1) {
                cur->rep_cnt--, cur->siz--;
```

```
return;
        }
        uint8 t state = 0;
        state |= (cur->ch[0] != nullptr);
        state = ((cur->ch[1] != nullptr) << 1);
        // 00都无,01有左无右,10,无左有右,11都有
        Node *tmp = cur;
        switch (state) {
            case 0:
                delete cur;
                cur = nullptr;
                break;
            case 1: // 有左无右
                cur = tmp->ch[0];
                delete tmp;
               break;
            case 2: // 有右无左
                cur = tmp->ch[1];
                delete tmp;
                break;
            case 3:
                rot type dir = cur->ch[0]->rank < cur->ch[1]->rank ? RT : LF;
                _rotate(cur, dir);
                del(cur->ch[!dir], val);
                cur->upd_siz();
                break;
        }
    }
}
int query rank(Node *cur, int val) {
    int less_siz = cur->ch[0] == nullptr ? 0 : cur->ch[0]->siz;
    if (val == cur->val)
       return less_siz + 1;
    else if (val < cur->val) {
        if (cur->ch[0] != nullptr)
            return _query_rank(cur->ch[0], val);
        else
           return 1;
    } else {
        if (cur->ch[1] != nullptr)
            return less_siz + cur->rep_cnt + _query_rank(cur->ch[1], val);
        else
           return cur->siz + 1;
    }
}
int query val(Node *cur, int rank) {
    int less_siz = cur->ch[0] == nullptr ? 0 : cur->ch[0]->siz;
    if (rank <= less_siz)</pre>
       return _query_val(cur->ch[0], rank);
```

```
else if (rank <= less siz + cur->rep cnt)
            return cur->val;
        else
            return _query_val(cur->ch[1], rank - less_siz - cur->rep_cnt);
    int _query_prev(Node *cur, int val) {
        if (val <= cur->val) {
            if (cur->ch[0] != nullptr) return _query_prev(cur->ch[0], val);
        } else {
            q_prev_tmp = cur->val;
            if (cur->ch[1] != nullptr) _query_prev(cur->ch[1], val);
            return q_prev_tmp;
        }
        return -1145;
    int _query_nex(Node *cur, int val) {
        if (val >= cur->val) {
            if (cur->ch[1] != nullptr) return _query_nex(cur->ch[1], val);
        } else {
            q_nex_tmp = cur->val;
            if (cur->ch[0] != nullptr) _query_nex(cur->ch[0], val);
            return q nex tmp;
        }
        return -1145;
    }
public:
   void insert(int val) { _insert(root, val); }
   void del(int val) { _del(root, val); }
   int query_rank(int val) { return _query_rank(root, val); }
   int query val(int rank) { return query val(root, rank); }
    int query_prev(int val) { return _query_prev(root, val); }
    int query_nex(int val) { return _query_nex(root, val); }
};
Treap tr;
int main() {
   srand(0);
   int t;
   scanf("%d", &t);
   while (t--) {
        int mode;
        int num;
        scanf("%d%d", &mode, &num);
        switch (mode) {
            case 1: tr.insert(num); break;
            case 2: tr.del(num); break;
            case 3: printf("%d\n", tr.query_rank(num)); break;
            case 4: printf("%d\n", tr.query_val(num)); break;
            case 5: printf("%d\n", tr.query_prev(num)); break;
```

```
case 6: printf("%d\n", tr.query_nex(num)); break;
}
}
```

无旋Treap

区间翻转

```
#include <bits/stdc++.h>
using namespace std;
struct Node {
   Node *ch[2];
   int val, prio;
   int cnt;
   int siz;
   bool to_rev = false;
   Node(int _val) : val(_val), cnt(1), siz(1) {
        ch[0] = ch[1] = nullptr;
        prio = rand();
    inline int upd_siz() {
        siz = cnt;
        if (ch[0] != nullptr) {
            siz += ch[0]->siz;
        if (ch[1] != nullptr) {
            siz += ch[1]->siz;
        return siz;
   inline void pushdown() {
        swap(ch[0], ch[1]);
        if (ch[0] != nullptr) {
            ch[0]->to_rev ^= 1;
        if (ch[1] != nullptr) {
            ch[1]->to_rev ^= 1;
        to_rev = false;
    inline void check_tag() {
        if (to_rev) {
            pushdown();
        }
    }
};
struct Seg_treap {
```

```
Node *root;
#define siz(_) (_ == nullptr ? 0 : _->siz)
array< Node*, 2 > split(Node *cur, int sz) {
    if (cur == nullptr) return {nullptr, nullptr};
    cur->check_tag();
    if (sz <= siz(cur->ch[0])) {
        auto temp = split(cur->ch[0], sz);
        cur->ch[0] = temp[1];
        cur->upd siz();
        return {temp[0], cur};
    } else {
        auto temp = split(cur->ch[1], sz - siz(cur->ch[0]) - 1);
        cur->ch[1] = temp[0];
        cur->upd_siz();
        return {cur, temp[1]};
    }
}
Node *merge(Node *sm, Node *bg) {
    if (sm == nullptr && bg == nullptr) {
        return nullptr;
    }
    if (sm != nullptr && bg == nullptr) {
        return sm;
    }
    if (sm == nullptr && bg != nullptr) {
        return bg;
    sm->check_tag(), bg->check_tag();
    if (sm->prio < bg->prio) {
        sm->ch[1] = merge(sm->ch[1], bg);
        sm->upd_siz();
        return sm;
    } else {
        bg \rightarrow ch[0] = merge(sm, bg \rightarrow ch[0]);
        bg->upd_siz();
        return bg;
    }
}
void insert(int val) {
    auto temp = split(root, val);
    auto l_tr = split(temp[0], val - 1);
    Node *new_node;
    if (l tr[1] == nullptr) new node = new Node(val);
    Node *l_tr_combined = merge(l_tr[0], l_tr[1] == nullptr ? new_node : l_tr[1]);
    root = merge(1 tr combined, temp[1]);
void seg_rev(int 1, int r) {
    auto less = split(root, 1 - 1);
    auto more = split(less[1], r - l + 1);
```

```
more[0]->to rev = true;
        root = merge(less[0], merge(more[0], more[1]));
    }
   void print(Node *cur) {
        if (cur == nullptr) return;
        cur->check_tag();
        print(cur->ch[0]);
        cout << cur->val << " ";</pre>
        print(cur->ch[1]);
   }
};
Seg_treap tr;
int main() {
   srand(time(NULL));
   int n, m;
   cin >> n >> m;
   for (int i = 1; i <= n; i++) {
        tr.insert(i);
    }
   while (m--) {
        int 1, r;
        cin >> 1 >> r;
        tr.seg_rev(1, r);
   tr.print(tr.root);
}
```

线段树

拓展性极强

好像也没什么板子, 纯手敲...

关于区间 \gcd ,可以维护差分数组d,然后查询 $\gcd(a[L],\gcd(d[L+1],\cdots,d[R]))$

若有区间修改,则d[L] + x, d[R+1] - x, 再维护一个区间加和单点查询

```
struct Info {
   int 1, r;
   LL sum, lz;
   friend Info operator + (const Info &1, const Info &r) {
        Info rt;
        rt.l = l.l;
        rt.r = r.r;
        rt.sum = l.sum + r.sum;
        rt.lz = 0;
```

```
return rt;
    }
};
namespace SegmentTree {
    Info inf[N << 2];</pre>
    void Push_up(int x) {
        \inf[x] = \inf[x << 1] + \inf[x << 1 | 1];
    }
    void Build(int 1, int r, int x = 1) {
        if (1 == r) {
            inf[x] = \{1, 1, 0, 0\};
            return;
        }
        int mid = (1 + r) >> 1;
        Build(1, mid, x \ll 1);
        Build(mid + 1, r, x << 1 | 1);
        Push up(x);
    void Set_lz(int x, LL v) {
        \inf[x].lz += \inf[x].lz;
        \inf[x].sum += (\inf[x].r - \inf[x].l + 1) * \inf[x].lz;
    void Push_down(int x) {
        if (!inf[x].lz) {
            return;
        }
        Set_lz(x << 1, inf[x].lz);
        Set_lz(x \ll 1 \mid 1, inf[x].lz);
        inf[x].lz = 0;
    void Update(int ql, int qr, LL val, int x = 1) {
        if (ql > qr \mid | ql > inf[x].r \mid | qr < inf[x].l) {
            return;
        }
        if (ql \le inf[x].l \&\& inf[x].r \le qr) {
            Set_lz(x, val);
            return;
        Push_down(x);
        Update(ql, qr, val, x \ll 1);
        Update(ql, qr, val, x \ll 1 \mid 1);
        Push_up(x);
    Info Query(int ql, int qr, int x = 1) {
        if (ql \le \inf[x].l \&\& \inf[x].r \le qr) {
            return inf[x];
        }
        Push_down(x);
        int mid = (\inf[x].l + \inf[x].r) >> 1;
```

```
if (qr <= mid) {
    return Query(ql, qr, x << 1);
} else if (ql > mid) {
    return Query(ql, qr, x << 1 | 1);
} else {
    return Query(ql, qr, x << 1) + Query(ql, qr, x << 1 | 1);
}
};</pre>
```

线段树合并&分裂

P5494 【模板】线段树分裂

```
struct Info {
   int cnt;
   Info() : cnt(0) {}
   friend Info operator + (const Info &1, const Info &r) {
        rt.cnt = 1.cnt + r.cnt;
        return rt;
   }
};
int n, q, a[N], col[N], rt[N];
set< int > s;
set< int >::iterator it;
struct SegmentTree {
   int tot;
   int col[N];
   int ls[N \ll 5], rs[N \ll 5], rub[N \ll 5], Bottom;
   Info inf[N << 5];</pre>
   SegmentTree() : tot(0), Bottom(0) {}
    // 新开节点
   int New() {
        return Bottom ? rub[Bottom --] : ++ tot;
    }
    // 删除节点
   void Delete(int &x) {
        ls[x] = rs[x] = 0;
       inf[x] = Info();
       rub[++ Bottom] = x;
        x = 0;
   void Push_Up(int x) {
        inf[x] = inf[ls[x]] + inf[rs[x]];
    // 使树x里权值为pos的个数增加k
   void Modify(int &x, int 1, int r, int pos, LL k) {
```

```
if (!x) {
        x = New();
    }
    if (1 == r) {
        inf[x].cnt += k;
        return;
    }
    int mid = (1 + r) >> 1;
    if (pos <= mid) {
        Modify(ls[x], l, mid, pos, k);
    } else {
        Modify(rs[x], mid + 1, r, pos, k);
    }
    Push_Up(x);
// 合并xy
int Merge(int x, int y, int 1, int r) {
    if (!x || !y) {
       return x | y;
    }
    int z = New();
    if (1 == r) {
        inf[z] = inf[x] + inf[y];
    } else {
        int mid = (1 + r) >> 1;
        ls[z] = Merge(ls[x], ls[y], l, mid);
        rs[z] = Merge(rs[x], rs[y], mid + 1, r);
        Push_Up(z);
        Delete(x); // 删除,看情况
        Delete(y);
    }
    return z;
}
// 将树x的[ql,qr]里的信息分裂到树y里
void Split(int &x, int &y, int 1, int r, int q1, int qr) {
    if (ql > qr || ql > r || qr < 1) {
       return;
    if (ql <= l && r <= qr) {
       y = x;
        x = 0;
       return;
    }
    y = ++ tot;
    int mid = (1 + r) >> 1;
    Split(ls[x], ls[y], l, mid, ql, qr);
    Split(rs[x], rs[y], mid + 1, r, ql, qr);
    Push_Up(x);
    Push_Up(y);
```

```
// 前k个给x,其余给y
void split(int p, int &x, int &y, int l, int r, int k) {
    if (!p) {
       x = y = 0;
       return;
    }
    if (1 == r) { // 断边
       if (k) {
           x = p;
           y = 0;
        } else {
           y = p;
           x = 0;
       return;
    }
    int mid = (1 + r) >> 1;
    if (inf[ls[p]].cnt >= k) { // 右儿子给y, 递归左儿子
       y = p;
       x = New();
        split(ls[p], ls[x], ls[y], l, mid, k);
    } else { // 左儿子给x, 递归右儿子
       x = p;
       y = New();
        split(rs[p], rs[x], rs[y], mid + 1, r, k - inf[ls[p]].cnt);
    Push_Up(x);
    Push_Up(y);
}
// 在树x里查询第k大
int Query_Kth(int x, int 1, int r, int k) {
   if (1 == r) {
       return 1;
    }
    int mid = (1 + r) >> 1;
    if (\inf[ls[x]].cnt \ge k) {
       return Query Kth(ls[x], 1, mid, k);
    return Query_Kth(rs[x], mid + 1, r, k - inf[ls[x]].cnt);
// 在树x里查询[ql,qr]的个数和
LL Query Cnt(int x, int 1, int r, int ql, int qr) {
    if (ql > qr || ql > r || qr < 1) {
       return 0;
    if (ql <= l && r <= qr) {
       return inf[x].cnt;
    }
```

```
int mid = (1 + r) >> 1;
       return Query_Cnt(ls[x], 1, mid, ql, qr) + Query_Cnt(rs[x], mid + 1, r, ql, qr);
   }
   int query(int x, int 1, int r) {
        if (1 == r) {
           return 1;
       int mid = (1 + r) >> 1;
        if (inf[ls[x]].cnt) {
           return query(ls[x], l, mid);
       return query(rs[x], mid + 1, r);
   }
   // 将区间分裂成[1,x]和(x,r]
   void split(int p) {
        it = s.lower_bound(p);
        if (*it == p) {
           return;
        }
       int r = *it, l = *prev(it) + 1;
       if (col[r]) {
            split(rt[r], rt[r], rt[p], 1, n, r - p);
        } else {
           split(rt[r], rt[p], rt[r], 1, n, p - 1 + 1);
       col[p] = col[r];
        s.insert(p);
    // 区间排序,0升序,1降序
   void Sort(int 1, int r, int op) {
        split(1 - 1);
        split(r);
       int x = 0;
        for (it = s.lower_bound(1); *it <= r;) {</pre>
           x = Merge(x, rt[*it], 1, n);
           rt[*it] = 0;
           set< int >::iterator IT = it;
           ++ it;
           s.erase(IT);
        }
        rt[r] = x;
       col[r] = op;
        s.insert(r);
} st;
```

可持久化权值线段树

```
struct President Tree {
   int tot;
   int rt[N], ls[N \ll 5], rs[N \ll 5], sum[N \ll 5], cnt[N \ll 5];
   int insert(int x, int 1, int r, int pos, int k) {
        int now = ++ tot;
        cnt[now] = cnt[x] + k;
        sum[now] = sum[x] + pos * k; // 若需离散化,则为初始值
        ls[now] = ls[x], rs[now] = rs[x];
        if (1 == r) {
           return now;
        }
        int mid = (1 + r) >> 1;
        if (pos <= mid) {</pre>
            ls[now] = insert(ls[x], l, mid, pos, k);
        } else {
            rs[now] = insert(rs[x], mid + 1, r, pos, k);
        }
        return now;
    }
    // 第k小
    int query kth(int x, int y, int l, int r, int kth) {
        if (1 == r) {
           return 1;
        }
        int lcnt = cnt[ls[x]] - cnt[ls[y]];
        int mid = (1 + r) >> 1;
       if (lcnt >= kth) {
            return query_kth(ls[x], ls[y], l, mid, kth);
        return query kth(rs[x], rs[y], mid + 1, r, kth - lcnt);
    }
    int query_cnt(int x, int y, int l, int r, int ql, int qr) {
        if (ql > qr || ql > r || qr < 1) {
            return 0;
        if (ql <= l && r <= qr) {
           return cnt[x] - cnt[y];
        int mid = (1 + r) >> 1;
        return query_cnt(ls[x], ls[y], l, mid, ql, qr) + query_cnt(rs[x], rs[y], mid +
1, r, ql, qr);
    int query sum(int x, int y, int l, int r, int ql, int qr) {
        if (ql > qr || ql > r || qr < 1) {
            return 0;
        }
```

```
if (ql <= l && r <= qr) {
        return sum[x] - sum[y];
}
int mid = (l + r) >> 1;
    return query_sum(ls[x], ls[y], l, mid, ql, qr) + query_sum(rs[x], rs[y], mid +
1, r, ql, qr);
};
```

树上主席树

```
\operatorname{rt}[\mathbf{x}] = \operatorname{insert}(\operatorname{rt}[\mathbf{y}], 1, \mathbf{m}, \operatorname{val}[\mathbf{y}]), \mathbf{y} \in \operatorname{son}[\mathbf{x}] 查询第k小时需要对x, y, \operatorname{lca}(\mathbf{x}, \mathbf{y}), \operatorname{fa}[\operatorname{lca}(\mathbf{x}, \mathbf{y})]一并查询 \operatorname{lcnt} = \operatorname{ls}[\mathbf{x}] + \operatorname{ls}[\mathbf{y}] - \operatorname{ls}[\operatorname{lca}(\mathbf{x}, \mathbf{y})] - \operatorname{ls}[\operatorname{fa}[\operatorname{lca}(\mathbf{x}, \mathbf{y})]]
```

带修主席树

```
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 2e5 + 5, int INF = 0x3f3f3f3f, int mod = 1e9 + 7;
int n, m;
int a[N];
int b[N], num;
int lnum, rnum;
int root[N], L[N], R[N], tot;
struct Seg { int 1, r, cnt;} tree[N * 200];
#define l(x) tree[x].l
#define r(x) tree[x].r
#define cnt(x) tree[x].cnt
int x[N], y[N], z[N];
bool vis[N];
inline int lowbit(int x) { return x & -x;}
int Get_Pos(int x) {
   return lower_bound(b + 1, b + num + 1, x) - b;
int update(int now, int 1, int r, int pos, int val) {
   int x = ++ tot;
   tree[x] = Seg\{l(now), r(now), cnt(now) + val\};
   if(l == r) return x;
   int mid = (1 + r) / 2;
    if(pos \le mid) l(x) = update(l(now), l, mid, pos, val);
   else r(x) = update(r(now), mid + 1, r, pos, val);
   return x;
}
```

```
void add(int x, int pos, int val) {
    for(int i = x; i <= n; i += lowbit(i)) {</pre>
        root[i] = update(root[i], 1, num, pos, val);
   }
}
int query(int 1, int r, int k) {
   if(l == r) return l;
   int lcnt = 0;
   for(int i = 1; i <= lnum; ++ i) lcnt -= cnt(l(L[i]));
   for(int i = 1; i <= rnum; ++ i) lcnt += cnt(l(R[i]));
   int mid = (1 + r) / 2;
   if(lcnt >= k) {
        for(int i = 1; i <= lnum; ++ i) L[i] = l(L[i]);
        for(int i = 1; i <= rnum; ++ i) R[i] = l(R[i]);
        return query(l, mid, k);
   } else {
        for(int i = 1; i \le lnum; ++ i) L[i] = r(L[i]);
        for(int i = 1; i \le rnum; ++ i) R[i] = r(R[i]);
        return query(mid + 1, r, k - lcnt);
   }
}
signed main() {
   cin.tie(nullptr)->sync_with_stdio(false);
   cin >> n >> m;
    for(int i = 1; i <= n; ++ i) {
        cin >> a[i];
        b[++ num] = a[i];
    }
    for(int i = 1; i <= m; ++ i) {
        string s;
        cin >> s;
        if(s == "C") {
            cin >> x[i] >> y[i];
            vis[i] = true;
            b[++ num] = y[i];
        } else {
            cin >> x[i] >> y[i] >> z[i];
        }
    }
    sort(b + 1, b + num + 1);
   num = unique(b + 1, b + num + 1) - (b + 1);
    for(int i = 1; i \le n; ++ i) add(i, Get Pos(a[i]), 1);
    for(int i = 1; i <= m; ++ i) {
        if(vis[i]) {
            add(x[i], Get Pos(a[x[i]]), -1);
            a[x[i]] = y[i];
            add(x[i], Get_Pos(a[x[i]]), 1);
        } else {
```

```
lnum = rnum = 0;
    for(int j = x[i] - 1; j; j -= lowbit(j)) L[++ lnum] = root[j];
    for(int j = y[i]; j; j -= lowbit(j)) R[++ rnum] = root[j];
    int pos = query(1, num, z[i]);
    cout << b[pos] << '\n';
}
return 0;
}</pre>
```

珂朵莉树

珂朵莉树 - OI Wiki

核心思想:把值相同的区间合并成几个节点保存在set里。

时间复杂度: $\Theta(q \times \log n \times \log n)$

```
struct Node {
   int 1, r;
   mutable int v; // 永远都可以被修改
   bool operator < (const Node &o) const {
       return 1 < o.1;
   }
};
set< Node > odt;
set< Node >::iterator Split(int x) {
   auto it = prev(odt.upper_bound({x, 0, 0}));
   if (it->1 == x) {
       return it;
   }
   auto [1, r, v] = *it;
   odt.erase(it);
   odt.insert(\{1, x - 1, v\});
   return odt.insert({x, r, v}).first;
}
void Assign(int 1, int r, int v) {
   auto R = Split(r + 1), L = Split(1);
    // TODO
   odt.erase(L, R);
   odt.insert({l, r, v});
}
void Maintain(int 1, int r) {
   auto R = Split(r + 1), L = Split(1);
   for (auto x = L; x != R; ++ x) {
       // TODO
    }
}
```

可持久化01Trie

```
constexpr int ALPHA SIZE = 30, N = 2e5 + 5;
struct Trie {
   int cur, rt[N], ch[N * ALPHA_SIZE][2], val[N * ALPHA_SIZE];
   void insert(int o, int lst, int v) {
        for (int i = ALPHA_SIZE; i >= 0; -- i) {
            val[o] = val[lst] + 1; // 在原版本的基础上更新
            if ((v \& (1 << i)) == 0) {
               if (!ch[o][0]) {
                   ch[0][0] = ++ cur;
               ch[0][1] = ch[lst][1];
               o = ch[o][0];
               lst = ch[lst][0];
            } else {
               if (!ch[o][1]) {
                   ch[0][1] = ++ cur;
               ch[0][0] = ch[lst][0];
               o = ch[o][1];
               lst = ch[lst][1];
            }
        }
       val[o] = val[lst] + 1;
        // printf("%d\n",0);
   }
    // 查询[l, r]内max{a[i]^v}
    // cout << query(st.rt[r], st.rt[l - 1], v) << '\n';
   int query(int o1, int o2, int v) {
        int ret = 0;
        for (int i = ALPHA_SIZE; i \ge 0; -- i) {
            // printf("%d %d %d\n",o1,o2,val[o1]-val[o2]);
            int t = ((v \& (1 << i)) ? 1 : 0);
            if (val[ch[o1][!t]] - val[ch[o2][!t]]) {
               ret += (1 << i);
               o1 = ch[o1][!t];
               o2 = ch[o2][!t]; // 尽量向不同的地方跳
            } else {
               o1 = ch[o1][t];
               o2 = ch[o2][t];
            }
        }
```

```
return ret;
}
```

例题:Five Day Couple

```
#include<bits/stdc++.h>
using namespace std;
using LL = long long;
const int N = 5e6 + 5, MOD = 998244353, INF = 0x3f3f3f3f;
struct Trie {
   int cnt, rt[N], ch[N * 33][2], val[N * 33];
   void insert(int o, int lst, int v) {
        for (int i = 30; i >= 0; i--) {
            val[o] = val[lst] + 1; // 在原版本的基础上更新
            if ((v \& (1 << i)) == 0) {
                if (!ch[o][0]) ch[o][0] = ++cnt;
                ch[0][1] = ch[lst][1];
                o = ch[o][0];
                lst = ch[lst][0];
            } else {
                if (!ch[o][1]) ch[o][1] = ++cnt;
                ch[0][0] = ch[1st][0];
                o = ch[o][1];
                lst = ch[lst][1];
            }
        }
       val[o] = val[lst] + 1;
   int query(int o1, int o2, int v) {
        int ret = 0;
        for (int i = 30; i >= 0; i--) {
            int t = ((v & (1 << i)) ? 1 : 0);
            if (val[ch[o1][!t]] - val[ch[o2][!t]])
                ret += (1 << i), o1 = ch[o1][!t],
                o2 = ch[o2][!t]; // 尽量向不同的地方跳
            else
                o1 = ch[o1][t], o2 = ch[o2][t];
        }
       return ret;
   }
} st;
int n, m, a[N];
signed main() {
   cin.tie(nullptr)->sync with stdio(false);
   cin >> n;
```

```
for (int i = 1; i <= n; ++ i) {
    cin >> a[i];
    st.rt[i] = ++ st.cnt;
    st.insert(st.rt[i], st.rt[i - 1], a[i]);
}
cin >> m;
while (m -- ) {
    int c, 1, r;
    cin >> c >> 1 >> r;
    cout << st.query(st.rt[r], st.rt[max(l - 1, 0)], c) << '\n';
}
return 0;
}</pre>
```

根号数据结构

分块

思想: 将长度为n的序列均分为 \sqrt{n} 个长度为 \sqrt{n} 的连续子序列,并记录每个'块'的编号、始末下标,然后进行操作. 每次操作时间复杂度为 θ $\left(\sqrt{n}\right)$,总的时间复杂度为 Θ $\left(\mathbf{n}\times\sqrt{\mathbf{n}}\right)$

初始化

 $\Theta(n)$

```
int num, st[N], ed[N], belong[N], siz[N];
void Pre_work(int n) {
    num = sqrt(n);
    for (int i = 1; i <= num; ++ i) {
        st[i] = n / num * (i - 1) + 1;
        ed[i] = n / num * i;
    }
    ed[num] = n;
    for (int i = 1; i <= num; ++ i) {
        for (int j = st[i]; j <= ed[i]; ++ j) {
            belong[j] = i;
        }
        siz[i] = ed[i] - st[i] + 1;
    }
}</pre>
```

例题:Yuno loves sqrt technology III

询问区间众数出现次数

```
#include<bits/stdc++.h>
```

```
using namespace std;
using LL = long long;
const int N = 5e5 + 5, MOD = 998244353, INF = 0x3f3f3f3f;
int n, m, a[N], b[N], tot;
vector< int > pos[N];
int st[N], ed[N], belong[N], siz[N], num;
int f[720][720], cnt[N], idx[N];
int Query(int 1, int r) {
   int x = belong[1], y = belong[r], ans = 0;
   if (x == y) {
        for (int i = 1; i \le r; ++ i)
            cnt[a[i]] = 0;
        for (int i = 1; i <= r; ++ i)
            ans = max(ans, ++ cnt[a[i]]);
       return ans;
   ans = f[x + 1][y - 1];
    for (int i = 1; i \le ed[x]; ++ i) {
        while (idx[i] + ans < pos[a[i]].size() && pos[a[i]][idx[i] + ans] <= r)
            ++ ans;
    for (int i = st[y]; i \le r; ++ i) {
        while (idx[i] - ans \ge 0 \&\& pos[a[i]][idx[i] - ans] \ge 1)
   return ans;
}
signed main() {
    cin.tie(nullptr)->sync with stdio(false);
   cin >> n >> m;
    for (int i = 1; i \le n; ++ i) {
        cin >> a[i];
        b[i] = a[i];
    }
   sort(b + 1, b + n + 1);
    tot = unique(b + 1, b + n + 1) - (b + 1);
    for (int i = 1; i <= n; ++ i) {
        a[i] = lower_bound(b + 1, b + tot + 1, a[i]) - b;
        pos[a[i]].push_back(i);
        idx[i] = (int)pos[a[i]].size() - 1;
   num = sqrt(n);
    for (int i = 1; i <= num; ++ i) {
        st[i] = n / num * (i - 1) + 1;
        ed[i] = n / num * i;
    ed[num] = n;
```

```
for (int i = 1; i <= num; ++ i) {
        for (int j = st[i]; j \le ed[i]; ++ j) {
            belong[j] = i;
        siz[i] = ed[i] - st[i] + 1;
        memset(cnt, 0, sizeof cnt);
        for (int j = i; j <= num; ++ j) {
            f[i][j] = f[i][j - 1];
            for (int k = st[j]; k \le ed[j]; ++ k) {
                f[i][j] = max(f[i][j], ++ cnt[a[k]]);
            }
        }
    }
   int lastans = 0;
   while (m --) {
        int 1, r;
        cin >> 1 >> r;
        1 ^= lastans;
        r ^= lastans;
        cout << (lastans = Query(1, r)) << '\n';</pre>
   return 0;
}
```

莫队

奇偶化排序,减少左右指针移动次数

```
bool operator < (const Query &o) const {
        return 1 / block == o.1 / block ? r != o.r && ((1 / block) & 1) ^ (r < o.r) : 1
< o.1;
    }</pre>
```

普通莫队套路都一样

```
int L = 1, R = 0;
for (int i = 1; i <= q; ++ i) {
    auto &[l, r, id] = Q[i];
    while (L > l) {
        add(-- L);
    }
    while (R < r) {
        add(++ R);
    }
    while (L < l) {
        del(L ++ );
    }
}</pre>
```

```
while (R > r) {
    del(R -- );
}
ans[id] = res;
}
```

树上莫队通过欧拉序将查询路径线段化,特判lca

例题:COT2 - Count on a tree II

求u ⇒ v路径上不同颜色个数

```
#include<bits/stdc++.h>
using namespace std;
using LL = long long;
const int N = 5e5 + 5;
int n, q, a[N];
int b[N], m;
vector< int > adj[N];
int seg[N], fir[N], las[N], num;
int dep[N], f[N][20], t;
void dfs(int x, int p) {
   seg[++ num] = x; fir[x] = num;
   for (int i = 1; i <= t; ++ i) {
        f[x][i] = f[f[x][i-1]][i-1];
    for (auto &y : adj[x]) {
        if (y != p) {
            dep[y] = dep[x] + 1;
            f[y][0] = x;
            dfs(y, x);
    }
   seg[++ num] = x; las[x] = num;
}
int lca(int x, int y) {
    if (dep[x] > dep[y]) swap(x, y);
    for (int i = t; ~i; -- i) {
       if (dep[f[y][i]] >= dep[x]) y = f[y][i];
    if (x == y) return x;
    for (int i = t; ~i; -- i) {
        if (f[x][i] != f[y][i]) {
            x = f[x][i];
            y = f[y][i];
        }
   return f[x][0];
}
```

```
int block;
struct Query {
    int 1, r, id, p;
   bool operator < (const Query &x) const {</pre>
        return 1 / block != x.1 / block ? 1 / block < x.1 / block : r < x.r;
    }
} Q[N];
bool vis[N];
int ans[N], cnt[N], res;
void add(int x) {
   if (++ cnt[a[x]] == 1) ++ res;
}
void del(int x) {
    if (-- cnt[a[x]] == 0) -- res;
void option(int x) {
   vis[x] ^= 1;
   if (!vis[x]) del(x);
    else add(x);
}
signed main() {
    cin.tie(nullptr)->sync with stdio(false);
    cin >> n >> q;
    for (int i = 1; i <= n; ++ i) {
       cin >> a[i];
        b[i] = a[i];
    sort(b + 1, b + n + 1);
    m = unique(b + 1, b + n + 1) - (b + 1);
    for (int i = 1; i <= n; ++ i) {
        a[i] = lower_bound(b + 1, b + m + 1, a[i]) - b;
    for (int i = 1; i < n; ++ i) {
       int u, v;
       cin >> u >> v;
        adj[u].push_back(v);
        adj[v].push_back(u);
    }
    t = ceil(log2(n));
    dep[1] = 1;
    dfs(1, 0);
    block = sqrt(n * 2);
    for (int i = 1; i \le q; ++ i) {
        int x, y;
        cin >> x >> y;
        if (fir[x] > fir[y]) swap(x, y);
        int p = lca(x, y);
        if (x == p) Q[i] = \{fir[x], fir[y], i, 0\};
```

```
else Q[i] = \{las[x], fir[y], i, p\};
    sort(Q + 1, Q + q + 1);
    int L = 1, R = 0;
    for (int i = 1; i \le q; ++ i) {
        auto \&[l, r, id, p] = Q[i];
        while (L > 1) option(seg[-- L]);
        while (R < r) option(seg[++ R]);</pre>
        while (L < 1) option(seg[L ++ ]);</pre>
        while (R > r) option(seg[R -- ]);
        if (p) option(p);
        ans[id] = res;
        if (p) option(p);
    }
    for (int i = 1; i <= q; ++ i) {
        cout << ans[i] << '\n';
    return 0;
}
```

树上启发式合并

子树信息的查询

- 遍历所有轻儿子, 递归结束时消除它们的贡献
- 遍历所有重儿子,保留它的贡献
- 再计算当前子树中所有轻子树的贡献
- 更新答案
- 如果当前点是轻儿子,消除当前子树的贡献

```
/* 第一个dfs计算每个节点的重儿子son */
void calc(int x, int p, bool flag) {
   if (flag) {
       /* 相当于add操作 */
   } else {
       /* 相当于delete操作 */
   for (int i = head[x]; i; i = Next[i]) {
       int y = ver[i];
       if (!vis[y] && y != p) {
           calc(y, x, flag);
       }
   }
void dfs2(int x, int p, bool keep) {
   for (int i = head[x]; i; i = Next[i]) {
       int y = ver[i];
       if (y != son[x] \&\& y != p) {
```

```
dfs2(y, x, false);
}

if (son[x]) {
    dfs2(son[x], x, true);
    vis[son[x]] = true;
}

calc(x, p, true), vis[son[x]] = false;
/* 统计答案 */
if (!keep) {
    calc(x, p, false); // 消除轻子树的贡献
}
```

由于每个节点到跟节点最多经过 $\log n$ 条轻边和 $\log n$ 条重边,所以总的时间复杂度是优秀的 $O(n\log n)$!

路径信息的查询

计算几何

还没学.先放个板子

```
const double eps = 1e-8;
int sign(double x) {
   if (fabs(x) < eps) {
        return 0;
   }
   if (x < 0) {
       return -1;
    }
   return 1;
}
int cmp(double x, double y) {
   if (fabs(x - y) < eps) {
       return 0;
    if (x < y) {
       return -1;
   return 1;
struct Point {
   double x, y;
   Point(double x = 0, double y = 0) : x(x), y(y) {}
   bool operator<(const Point& B) const {</pre>
        return x == B.x ? y < B.y : x < B.x;
    }
   bool operator==(const Point& B) const {
```

```
return !sign(x - B.x) && !sign(y - B.y);
   Point operator+(const Point& B) const {
       return Point(x + B.x, y + B.y);
    }
   Point operator-(const Point& B) const {
       return Point(x - B.x, y - B.y);
    }
   Point operator*(const double a) const {
       return Point(x * a, y * a);
    }
   Point operator/(const double a) const {
       return Point(x / a, y / a);
    }
   double operator*(const Point& B) const {
       return x * B.x + y * B.y;
   double operator^(const Point& B) const {
       return x * B.y - y * B.x;
    }
   double length() {
       return sqrt(x * x + y * y);
   }
   Point trunc(double r) { // 化为长度为r的向量
       double 1 = length();
       if (!sign(l)) {
           return *this;
        }
       r /= 1;
       return Point(x * r, y * r);
   }
   Point toleft() {
       return Point(-y, x);
   }
    friend int relation(Point a, Point b, Point c) {
        int flag = sign((b - a) ^ (c - a));
        if (flag > 0) {
           return 1;
        } else if (flag < 0) {</pre>
           return -1;
        }
       return 0;
   }
};
using Vector = Point;
struct Line {
   Point p;
   Vector v;
   double rad;
```

```
Line() {}
   Line(Point P, Vector V) : p(P), v(V) {
        rad = atan2(v.y, v.x);
    }
   Point get_point_in_line(double t) {
       return p + v * t;
    }
   bool operator<(const Line& L) const {</pre>
        if (!cmp(rad, L.rad)) {
           return L.onLeft(p) > 0;
        }
       return rad < L.rad;
   }
   int onLeft(const Point& a) const { // 点a是否在直线的左边(>0:左 <0:右)
        return relation(p, p + v, a);
   }
   friend Point getIntersection(Line a, Line b) {
       Vector u = a.p - b.p;
       double t = (b.v ^ u) / (a.v ^ b.v);
       return a.get_point_in_line(t);
   friend bool on right(const Line& a, const Line& b, const Line& c) { // b,c的交点是否
在直线a的右边
       Point o = getIntersection(b, c);
       return a.onLeft(o) <= 0;</pre>
    }
   friend Line getTranslation(Line a, double dist) { // 将直线a向左(逆时针)平移dist
       Vector u = a.v.toleft().trunc(dist);
       return Line(a.p + u, a.v);
   };
};
vector<Point> half_plane(vector<Line> line) {
   sort(line.begin(), line.end());
   int hh = 0, tt = -1, num = line.size();
   vector<int> q(num + 10);
   for (int i = 0; i < num; ++ i) {
        if (i && !cmp(line[i].rad, line[i - 1].rad)) {
           continue;
        while (hh + 1 \le tt \&\& on_right(line[i], line[q[tt - 1]], line[q[tt]])) {
           -- tt;
        while (hh + 1 \le tt \&\& on right(line[i], line[q[hh]], line[q[hh + 1]])) {
           ++ hh;
        q[++ tt] = i;
   while (hh + 1 <= tt && on_right(line[q[hh]], line[q[tt - 1]], line[q[tt]])) {
        -- tt;
```

```
while (hh + 1 <= tt && on_right(line[q[tt]], line[q[hh]], line[q[hh + 1]])) {
       ++ hh;
    }
    if (tt - hh + 1 <= 2) {
       return {};
    }
   q[++ tt] = q[hh];
    vector<Point> res;
   for (int i = hh; i < tt; ++ i) {
        res.push_back(getIntersection(line[q[i]], line[q[i + 1]]));
   return res;
}
double getArea(vector<Point>& p) {
   int n = p.size();
   double ans = 0;
   for (int i = 0; i < n; ++ i) {
        ans += (p[i] ^ p[(i + 1) % n]);
    }
   return ans / 2;
}
```

字符串

字符串Hash

 $Hash[S+T] = Hash[S] imes P^{length[T]} + Hash[T]$ 其中,P代表将字符串以P进制存储

```
using ULL = unsigned long long;
struct Hash {
    ULL pw[N];
    const ULL base = 131ULL;
    void Calc_Power() {
        pw[0] = 1;
        for (int i = 1; i < N; ++ i) {
            pw[i] = pw[i - 1] * base;
        }
    }
    vector< ULL > Get_hash(const string &str) {
        vector< ULL > hs(str.size() + 1);
        for (int i = 0; i < str.size(); ++ i) {
            hs[i + 1] = hs[i] * base + str[i];
        }
        return hs;</pre>
```

```
}
ULL Get_val(const vector< ULL > &hs, int 1, int r) {
    return hs[r] - hs[l - 1] * pw[r - l + 1];
}
};
```

KMP

next[i]代表S中以i结尾的非前缀子串与S的前缀能够匹配的最长长度

```
// index from 1 to n
vector< int > calc next(char *s, int n) {
   vector< int > next(n + 1);
   for (int i = 2, j = 0; i \le n; ++ i) {
        while (j > 0 \&\& s[i] != s[j + 1]) {
            j = next[j];
        if (s[i] == s[j + 1]) {
           ++ j;
       next[i] = j;
   }
   return next;
}
// if i % (i - next[i]) == 0
// s[1 ~ i]的最小循环节长度为i - next[i],循环次数为i / (i - next[i])
vector< int > match(char *s, int n, char *t, int m) {
   vector< int > f(m + 1);
   vector< int > next = calc_next(s, n);
   for (int i = 1, j = 0; i \le m; ++ i) {
        while (j > 0 \&\& (j == n \mid | t[i] != s[j + 1])) {
            j = next[j];
        if (t[i] == s[j + 1]) {
           ++ j;
        f[i] = j; // if (f[i] == n) 此时就是A在B中的某一次出现
    }
   return f;
}
```

exKMP

定义 $\mathbf{z}[i]$ 代表 \mathbf{S} 的前缀与 \mathbf{S} 中以 \mathbf{i} 开头的后缀能够匹配的最长长度

```
vector<int> z_function(string s) {
```

```
int n = (int) s.length();
   vector<int> z(n);
   for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (i \le r \&\& z[i-1] \le r-i+1) {
            z[i] = z[i - 1];
        } else {
            z[i] = max(0, r - i + 1);
            while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) {
               ++z[i];
           }
       }
        if (i + z[i] - 1 > r) {
           1 = i;
           r = i + z[i] - 1;
        }
   }
   return z;
}
```

Manacher

```
char s[N], str[N * 2];
int Len[N * 2];
int len;
void init() {
   int k = 0;
    str[k++] = '@';
    for (int i = 0; i < len; ++i) {
        str[k++] = '#';
        str[k++] = s[i];
    }
    str[k++] = '#';
    len = k;
    str[k] = '?';
}
int manacher() {
    int mx = 0, id = 0;
    int ans = 0;
    for (int i = 1; i \le len - 1; ++ i) {
        if (mx > i) {
            Len[i] = min(Len[id * 2 - i], mx - i);
        } else {
            Len[i] = 1;
```

AC自动机

fail 指针指向所有模式串的前缀中匹配当前状态的最长后缀

fail指针相当于S的后缀集合

P5357 【模板】AC 自动机 (二次加强版)

求出每个模式串 T_i 在文本串S中出现次数

```
struct Aho_Corasick_Automaton {
   static constexpr int ALPFABET_SIZE = 26, N = 1e5;
   int trie[N][ALPFABET SIZE], fail[N], cnt[N], tot;
   void Insert(const std::string &str) {
        int rt = 0;
        for (const auto &x : str) {
            if (!trie[rt][x - 'a']) {
                trie[rt][x - 'a'] = ++ tot;
            rt = trie[rt][x - 'a'];
        ++ cnt[rt];
    }
   void Build() {
        std::queue< int > q;
        for (int i = 0; i < ALPFABET_SIZE; ++ i) {</pre>
            if (trie[0][i]) {
                q.push(trie[0][i]);
            }
        }
        while (!q.empty()) {
            int x = q.front();
            q.pop();
            for (int i = 0; i < ALPFABET_SIZE; ++ i) {</pre>
                if (trie[x][i]) {
```

后缀自动机

len代表每个endpos等价类串的最长长度,len[fa[a]] + 1 = minlen(a)

【模板】后缀自动机 (SAM)

请你求出8的所有出现次数不为1的子串的出现次数乘上该子串长度的最大值

```
struct Suffix_Automaton {
   static constexpr int ALPHBET_SIZE = 26, N = 1e6;
   int last = 1, cntNodes = 1;
   int ch[N * 2][ALPHBET_SIZE], fa[N * 2], len[N * 2];
    int Extend(int c) {
        int p = last, np = last = ++ cntNodes;
        len[np] = len[p] + 1;
        while (p && !ch[p][c]) {
            ch[p][c] = np;
           p = fa[p];
        }
        if (!p) {
            fa[np] = 1;
        } else {
            int q = ch[p][c];
            if (len[p] + 1 == len[q]) {
                fa[np] = q;
            } else {
                int nq = ++ cntNodes;
                len[nq] = len[p] + 1;
                std::copy(ch[q], ch[q] + ALPHBET_SIZE, ch[nq]);
                fa[nq] = fa[q];
                fa[q] = fa[np] = nq;
                while (p \&\& ch[p][c] == q) {
                    ch[p][c] = nq;
                    p = fa[p];
```

```
}
return np;
}
```

回文自动机

每个节点代表了一个回文串, $\mathrm{ch}[\mathrm{S}][\mathrm{c}]$ 代表在状态S的左右各添加一个字符c

那么从根节点出发沿着 $\mathrm{ch}[S][c]$ 即可组成 cSc

 $\operatorname{fail}[S]$ 代表回文串S的最长回文后缀, last 代表上一个字符结尾的最长回文子串

P5496 【模板】回文自动机 (PAM)

求以s[i]结尾的回文串个数

```
struct Palindromic_Automaton {
   int s[N], top; // 原串
   int ch[N][26], fail[N], len[N], tot, last;
   Palindromic_Automaton() {
        s[0] = -114514;
       tot = -1;
       New(0); // 偶根
       New(-1); // 奇根
       fail[0] = 1;
       last = 0;
    }
   int New(int length) {
       len[++ tot] = length;
       return tot;
    }
   int Get_Fail(int x) { // 找最长的回文后缀
        while (s[top - len[x] - 1] != s[top]) {
           x = fail[x];
        }
        return x;
   void Extend(int c) {
        s[++ top] = c;
        int now = Get_Fail(last);
        if (!ch[now][c]) {
           int x = New(len[now] + 2);
            fail[x] = ch[Get_Fail(fail[now])][c];
            ch[now][c] = x;
```

```
}
last = ch[now][c];
}
```

动态规划

数位dp

例题 windy数

不含前导零且相邻两个数字之差至少为 2 的正整数被称为 windy 数。windy 想知道,在 a 和 b 之间,包括 a 和 b ,总共有多少个 windy 数?

对于全部的测试点,保证 $1 \le a \le b \le 2 \times 10^9$ 。

```
#include<bits/stdc++.h>
using namespace std;
using LL = long long;
const int N = 205;
int a[N], cnt;
int dp[N][N][2][2];
int dfs(int pos, int pre, bool lead, bool limit) {
   if (pos > cnt) return 1;
   if (dp[pos][pre][lead][limit]) return dp[pos][pre][lead][limit];
   int res = limit ? a[pos] : 9;
   int ans = 0;
   for (int i = 0; i <= res; ++ i) {
        if (abs(i - pre) < 2) continue;
        if (lead \&\& i == 0) ans += dfs(pos + 1, pre, lead, limit \&\& i == a[pos]);
        else ans += dfs(pos + 1, i, 0, limit && i == a[pos]);
   return dp[pos][pre][lead][limit] = ans;
}
int calc(int x) {
   cnt = 0;
   memset(dp, 0, sizeof dp);
   for (cnt = 0; x; x \neq 10) a[++ cnt] = x % 10;
   reverse(a + 1, a + cnt + 1);
   return dfs(1, -2, 1, 1);
signed main() {
    cin.tie(nullptr)->sync_with_stdio(false);
```

```
int l, r;
cin >> l >> r;
cout << calc(r) - calc(l - 1) << '\n';
return 0;
}</pre>
```

最长上升子序列计数

 $\Theta(n^2)$

```
pair< int, int > dp[N];
cin >> n;
for (int i = 1; i <= n; ++ i) {
    cin >> a[i];
}
int max_len = 1;
for (int i = 1; i \le n; ++ i) {
    dp[i] = \{1, 1\};
    for (int j = 1; j < i; ++ j) {
        if (a[j] \le a[i]) \{
            if (dp[j].first + 1 > dp[i].first) {
                dp[i] = {dp[j].first + 1, dp[j].second};
            } else if (dp[j].first + 1 == dp[i].first) {
                dp[i].second += dp[j].second;
            }
        }
    }
    max_len = max(max_len, dp[i].first);
}
int ans = 0;
for (int i = 1; i <= n; ++ i) {
    if (dp[i].first == max_len)
        ans += dp[i].second;
cout << ans << '\n';
```

 $\Theta(n \log n)$

```
#include<bits/stdc++.h>
using namespace std;
using LL = long long;
const int N = 4e5 + 5, MOD = 1e9 + 7;

int n, a[N];
struct SegTree {
   int len, cnt;
}
```

```
} st[N << 2];</pre>
SegTree operator + (const SegTree &1, const SegTree &r) {
    if (!1.len | r.len > 1.len)
       return r;
   if (!r.len | | 1.len > r.len)
        return 1;
   return {l.len, (l.cnt + r.cnt) % MOD};
}
void build(int x, int l, int r) {
   if (1 == r) {
        st[x] = \{0, 1\};
        return;
   }
   int mid = (1 + r) >> 1;
   build(x << 1, 1, mid);
   build(x << 1 | 1, mid + 1, r);
   st[x] = st[x << 1] + st[x << 1 | 1];
void update(int x, int l, int r, int pos, const SegTree &val) {
   if (1 == r) {
        st[x] = val;
        return;
    }
   int mid = (1 + r) >> 1;
    if (pos <= mid)
        update(x << 1, 1, mid, pos, val);
   else
        update(x \ll 1 \mid 1, mid + 1, r, pos, val);
   st[x] = st[x << 1] + st[x << 1 | 1];
SegTree query(int x, int l, int r, int ql, int qr) {
   if (ql > qr \mid | ql > r \mid | qr < l)
       return {0, 1};
   if (ql <= l && r <= qr)
        return st[x];
   int mid = (1 + r) >> 1;
   return query(x \ll 1, 1, mid, ql, qr) + query(x \ll 1 | 1, mid + 1, r, ql, qr);
int b[N], m;
signed main() {
    cin.tie(nullptr)->sync_with_stdio(false);
   cin >> n;
   for (int i = 1; i <= n; ++ i) {
        cin >> a[i];
        b[i] = a[i];
   sort(b + 1, b + n + 1);
   m = unique(b + 1, b + n + 1) - (b + 1);
```

```
build(1, 1, m);
   int max len = 0;
   for (int i = 1; i \le n; ++ i) {
        a[i] = lower_bound(b + 1, b + m + 1, a[i]) - b;
        auto x = query(1, 1, m, 1, a[i] - 1);
       ++ x.len;
       max_len = max(max_len, x.len);
        auto y = query(1, 1, m, a[i], a[i]);
        // 当且仅当f[j]+1=f[i]才能将其合并
       update(1, 1, m, a[i], x + y);
        // auto z = query(1, 1, m, a[i], a[i]);
        // cout << z.len << ' ' << z.cnt << '\n';
   }
   int ans = 0;
   for (int i = 1; i <= m; ++ i) {
        auto x = query(1, 1, m, i, i);
       if (x.len == max len)
           ans = (ans + x.cnt) % MOD;
    }
   cout << ans << '\n';
   return 0;
}
```

数学

数论分块

给定n和k,求 $\sum_{i=1}^n k\%i$, $1 \le n, k \le 10^9$

```
LL ans = 0;
for (int x = 1, gx; x <= n; x = gx + 1) {
    gx = k / x ? min(k / (k / x), n) : n;
    ans -= 1LL * (k / x) * (x + gx) * (gx - x + 1) / 2;
}</pre>
```

扩展欧几里得算法

对于任意整数a, b,存在一对整数x, y,满足 $a \cdot x + b \cdot y = \gcd(a, b)$

```
void exGCD(int a, int b, int &x, int &y) {
  if (b == 0) { x = 1, y = 0; return a;}
  int d = exGCD(b, a % b, x, y);
  int z = x; x = y; y = z - y * (a / b);
  return d;
}
```

矩阵快速幂

求Fibonacci第n项

```
(0 \le n \le 2 	imes 10^9)
```

```
void mul(int f[2], int a[2][2]) {
   int c[2];
   memset(c, 0, sizeof c);
    for (int j = 0; j < 2; ++ j) {
        for (int k = 0; k < 2; ++ k) {
            c[j] = (c[j] + (LL)f[k] * a[k][j]) % MOD;
        }
    }
    memcpy(f, c, sizeof c);
}
void mulself(int a[2][2]) {
   int c[2][2];
   memset(c, 0, sizeof c);
    for (int i = 0; i < 2; ++ i) {
        for (int j = 0; j < 2; ++ j) {
            for (int k = 0; k < 2; ++ k) {
                c[i][j] = (c[i][j] + (LL)a[i][k] * a[k][j]) % MOD;
        }
   memcpy(a, c, sizeof c);
void solve() {
   cin >> n;
   int f[2] = \{0, 1\};
   int a[2][2] = \{\{0, 1\}, \{1, 1\}\};
    for (; n; n >>= 1) {
        if (n & 1) mul(f, a);
        mulself(a);
   cout << f[0] << '\n';
}
```

高斯消元

```
double a[20][20], b[20], c[20][20];
int n;
void solve() {
   cin >> n;
   for (int i = 1; i \le n + 1; ++ i) {
       for (int j = 1; j \le n; ++ j) {
           cin >> a[i][j];
       }
    }
   // c:系数矩阵,b:常数,二者一起构成增广矩阵
   for (int i = 1; i <= n; ++ i) {
       for (int j = 1; j \le n; ++ j) {
           c[i][j] = 2 * (a[i][j] - a[i + 1][j]);
           b[i] += a[i][j] * a[i][j] - a[i + 1][j] * a[i + 1][j];
       }
    }
    // 高斯消元(数据保证一定有解)
   for (int i = 1; i \le n; ++ i) {
       // 找到x[i]的系数不为0的一个方程
       for (int j = i; j \le n; ++ j) {
            if (fabs(c[j][i]) > 1e-8) {
               for (int k = 1; k \le n; ++ k) {
                   swap(c[i][k], c[j][k]);
               swap(b[i], b[j]);
           }
       }
       // 消去其他方程的x[i]的系数
       for (int j = 1; j \le n; ++ j) {
           if (i == j) {
               continue;
            double rate = c[j][i] / c[i][i];
            for (int k = i; k \le n; ++ k) {
               c[j][k] -= c[i][k] * rate;
           b[j] -= b[i] * rate;
       }
    }
    for (int i = 1; i \le n; ++ i) {
       printf("%.3f%c", b[i] / c[i][i], i == n ? '\n' : ' ');
    }
}
```

组合数取模相关

```
Mint inv[N], fac[N], finv[N];
void Pre_Work() {
    fac[0] = fac[1] = 1;
    inv[1] = 1;
    finv[0] = finv[1] = 1;
    for(int i = 2; i < N; ++ i) {
        fac[i] = fac[i-1] * i;
        inv[i] = P - P / i * inv[P * i];
        finv[i] = finv[i - 1] * inv[i];
    }
}
Mint C(int x, int y) {
    return fac[x] * finv[x - y] * finv[y];
}</pre>
```

线性筛

最小质因子

```
constexpr int NUMBER_SIZE = 2e5 + 5;
int v[NUMBER_SIZE], pri[NUMBER_SIZE], Pcnt;
void Pre_Work() {
    Pcnt = 0;
    for (int i = 2; i < NUMBER_SIZE; ++ i) {
        if (!v[i]) {
            v[i] = i;
            pri[tot ++ ] = i;
        }
        for (int j = 0; j < Pcnt && 1LL * pri[j] * i < NUMBER_SIZE; ++ j) {
            if (pri[j] > v[i]) {
                 break;
            }
            v[pri[j] * i] = pri[j];
        }
    }
}
```

欧拉函数

```
constexpr int NUMBER_SIZE = 2e5 + 5;
int phi[NUMBER_SIZE], pri[NUMBER_SIZE], Pcnt;
bool is_prime[NUMBER_SIZE];
void Pre_Work() {
   std::fill(is_prime + 1, is_prime + NUMBER_SIZE, 1);
   Pcnt = 0;
```

```
is prime[1] = 0;
    phi[1] = 1;
    for (int i = 2; i < NUMBER SIZE; ++ i) {</pre>
        if (is_prime[i]) {
            pri[++ Pcnt] = i;
            phi[i] = i - 1;
        }
        for (int j = 1; j <= Pcnt && 1LL * i * pri[j] < NUMBER_SIZE; ++ j) {</pre>
            is_prime[i * pri[j]] = 0;
            if (i % pri[j]) {
                phi[i * pri[j]] = phi[i] * phi[pri[j]];
                phi[i * pri[j]] = phi[i] * pri[j];
                break;
            }
       }
  }
}
```

莫比乌斯函数

```
constexpr int NUMBER SIZE = 2e5 + 5;
int mu[NUMBER_SIZE], pri[NUMBER_SIZE], v[NUMBER_SIZE], Pcnt;
void Pre_Work() {
   Pcnt = 0;
   mu[1] = 1;
   for (int i = 2; i < NUMBER_SIZE; ++ i) {</pre>
        if (!v[i]) {
            mu[i] = -1;
            pri[Pcnt ++ ] = i;
        for (int j = 0; j < Pcnt && 1LL * i * pri[j] < NUMBER_SIZE; ++ j) {</pre>
            v[i * pri[j]] = 1;
            if (i % pri[j] == 0) {
                mu[i * pri[j]] = 0;
                break;
            mu[i * pri[j]] = -mu[i];
        }
   }
}
```

约数个数

```
constexpr int NUMBER SIZE = 2e5 + 5;
int d[NUMBER SIZE], pri[NUMBER SIZE], num[NUMBER SIZE], Pcnt;
bool v[NUMBER_SIZE];
void Pre_Work() {
   Pcnt = 0;
   d[1] = 1;
   for (int i = 2; i < NUMBER_SIZE; ++ i) {</pre>
        if (!v[i]) {
            v[i] = 1;
            pri[Pcnt ++ ] = i;
            d[i] = 2;
            num[i] = 1;
        for (int j = 0; j < Pcnt && 1LL * i * pri[j] < NUMBER SIZE; ++ j) {
            v[pri[j] * i] = 1;
            if (i % pri[j] == 0) {
                num[i * pri[j]] = num[i] + 1;
                d[i * pri[j]] = d[i] / num[i * pri[j]] * (num[i * pri[j]] + 1);
                break;
            } else {
                num[i * pri[j]] = 1;
                d[i * pri[j]] = d[i] * 2;
        }
   }
}
```

约数和

```
constexpr int NUMBER SIZE = 2e5 + 5;
int g[NUMBER_SIZE], f[NUMBER_SIZE], pri[NUMBER_SIZE], Pcnt;
bool v[NUMBER_SIZE];
void Pre_Work() {
    Pcnt = 0;
    g[1] = f[1] = 1;
    for (int i = 2; i < NUMBER_SIZE; ++ i) {</pre>
        if (!v[i]) {
            v[i] = 1;
            pri[Pcnt ++ ] = i;
            g[i] = i + 1;
            f[i] = i + 1;
        for (int j = 0; j < Pcnt && 1LL * i * pri[j] < NUMBER_SIZE; ++ j) {</pre>
            v[pri[j] * i] = 1;
            if (i % pri[j] == 0) {
                g[i * pri[j]] = g[i] * pri[j] + 1;
```

杂项

快读快写

```
template < typename T > void read(T &x) {
   x = 0;
   T f = 1;
   char ch = getchar();
   while (ch < '0' | | ch > '9') {
        if (ch == '-') {
            f = -1;
        }
       ch = getchar();
   }
   while (ch >= '0' && ch <= '9') {
        x = (x << 1) + (x << 3) + (ch^48);
       ch = getchar();
   x = x * f;
template < typename T > void write(T x) {
   if (x < 0) {
       putchar('-');
       x = -x;
    }
    if (x < 10) {
        putchar(x + '0');
    } else {
       write(x / 10);
        putchar(x % 10 + '0');
}
template < typename T, typename ...Arg > void read(T &x, Arg &...arg) {
   read(x);
```

```
read(arg...);
}
template < typename T, typename ...Arg > void write(T &x,Arg &...arg) {
    write(x);
    putchar(' ');
    write(arg...);
}
```

吸氧优化

```
#pragma GCC optimize(3, "Ofast")
#pragma GCC optimize(2, "Ofast")
```

取模运算

```
constexpr int P = 998244353;
// assume -P <= x < 2P
int norm(int x) {
   if (x < 0) {
       x += P;
   }
   if (x \ge P) {
       x = P;
    }
   return x;
}
template<class T>
T power(T a, LL b) {
   T res = 1;
   for (; b; b /= 2, a *= a) {
        if (b % 2) {
           res *= a;
        }
    }
   return res;
}
struct Mint {
   int x;
   Mint(int x = 0) : x(norm(x)) \{ \}
   Mint(LL x) : x(norm(x % P)) {}
   int val() const {
        return x;
   Mint operator-() const {
        return Mint(norm(P - x));
    }
   Mint inv() const {
```

```
assert(x != 0);
        return power(*this, P - 2);
    }
   Mint &operator*=(const Mint &rhs) {
        x = LL(x) * rhs.x % P;
        return *this;
    }
   Mint &operator+=(const Mint &rhs) {
        x = norm(x + rhs.x);
       return *this;
    }
   Mint &operator == (const Mint &rhs) {
       x = norm(x - rhs.x);
       return *this;
   Mint &operator/=(const Mint &rhs) {
        return *this *= rhs.inv();
    friend Mint operator*(const Mint &lhs, const Mint &rhs) {
       Mint res = lhs;
       res *= rhs;
        return res;
    }
   friend Mint operator+(const Mint &lhs, const Mint &rhs) {
       Mint res = lhs;
       res += rhs;
       return res;
    friend Mint operator-(const Mint &lhs, const Mint &rhs) {
       Mint res = lhs;
       res -= rhs;
        return res;
    friend Mint operator/(const Mint &lhs, const Mint &rhs) {
       Mint res = lhs;
       res /= rhs;
       return res;
    friend std::istream &operator>>(std::istream &is, Mint &a) {
       LL v;
       is >> v;
        a = Mint(v);
        return is;
    }
    friend std::ostream &operator<<(std::ostream &os, const Mint &a) {
        return os << a.val();
    }
};
```

随机化

随机排序

```
std::mt19937 rng(std::chrono::system_clock::now().time_since_epoch().count());
std::shuffle(begin(a), end(a), rng);
```

程序运行时间

```
std::cout << (double)clock() / CLOCKS_PER_SEC * 1000 << " ms\n";</pre>
```

高精度运算

```
class bign {
public:
   int len, s[N];//数的长度, 记录数组
//构造函数
   bign();
   bign(const char *);
   bign(int);
   bool sign;//符号 1正数 0负数
   string toStr() const;//转化为字符串,主要是便于输出
   friend istream &operator>>(istream &, bign &);//重载输入流
   friend ostream & operator << (ostream &, bign &);//重载输出流
//重载复制
   bign operator=(const char *);
   bign operator=(int);
   bign operator=(const string);
//重载各种比较
   bool operator>(const bign &) const;
   bool operator>=(const bign &) const;
   bool operator<(const bign &) const;</pre>
```

```
bool operator<=(const bign &) const;</pre>
   bool operator==(const bign &) const;
   bool operator!=(const bign &) const;
//重载四则运算
   bign operator+(const bign &) const;
   bign operator++();
   bign operator++(int);
   bign operator+=(const bign &);
   bign operator-(const bign &) const;
   bign operator--();
   bign operator--(int);
   bign operator==(const bign &);
   bign operator*(const bign &) const;
   bign operator*(const int num) const;
   bign operator*=(const bign &);
   bign operator/(const bign &) const;
   bign operator/=(const bign &);
//四则运算的衍生运算
   bign operator%(const bign &) const;//取模(余数)
   bign factorial() const;//阶乘
   bign Sqrt() const;//整数开根(向下取整)
   bign pow(const bign &) const;//次方
//一些乱乱的函数
   void clean();
   ~bign();
};
bign::bign() {
   memset(s, 0, sizeof(s));
   len = 1;
   sign = 1;
```

```
bign::bign(const char *num) {
   *this = num;
}
bign::bign(int num) {
   *this = num;
}
string bign::toStr() const {
    string res;
    res = "";
    for (int i = 0; i < len; i++) {
        res = (char) (s[i] + '0') + res;
    }
    if (res == "") {
       res = "0";
    }
    if (!sign && res != "0") {
        res = "-" + res;
    return res;
}
istream &operator>>(istream &in, bign &num) {
    string str;
    in >> str;
    num = str;
   return in;
}
ostream &operator << (ostream &out, bign &num) {
    out << num.toStr();</pre>
   return out;
}
bign bign::operator=(const char *num) {
    memset(s, 0, sizeof(s));
    char a[N] = "";
    if (num[0] != '-') {
        strcpy(a, num);
    } else {
        for (int i = 1; i < strlen(num); ++ i) {</pre>
            a[i - 1] = num[i];
    }
    sign = !(num[0] == '-');
    len = strlen(a);
```

```
for (int i = 0; i < strlen(a); ++ i) {</pre>
        s[i] = a[len - i - 1] - 48;
    }
    return *this;
}
bign bign::operator=(int num) {
    char temp[N];
    sprintf(temp, "%d", num);
    *this = temp;
   return *this;
}
bign bign::operator=(const string num) {
    const char *tmp;
    tmp = num.c_str();
    *this = tmp;
    return *this;
}
bool bign::operator<(const bign &num) const {</pre>
    if (sign ^ num.sign) {
        return num.sign;
    }
    if (len != num.len) {
        return len < num.len;</pre>
    for (int i = len - 1; i >= 0; -- i) {
        if (s[i] != num.s[i]) {
            return sign ? (s[i] < num.s[i]) : (!(s[i] < num.s[i]));
        }
    }
   return !sign;
}
bool bign::operator>(const bign &num) const {
   return num < *this;
}
bool bign::operator<=(const bign &num) const {</pre>
    return !(*this > num);
}
bool bign::operator>=(const bign &num) const {
   return !(*this < num);</pre>
bool bign::operator!=(const bign &num) const {
    return *this > num || *this < num;
```

```
bool bign::operator==(const bign &num) const {
    return !(num != *this);
}
bign bign::operator+(const bign &num) const {
    if (sign ^ num.sign) {
        bign tmp = sign ? num : *this;
        tmp.sign = 1;
        return sign ? *this - tmp : num - tmp;
    bign result;
    result.len = 0;
    int temp = 0;
    for (int i = 0; temp | i < (max(len, num.len)); ++ i) {</pre>
        int t = s[i] + num.s[i] + temp;
        result.s[result.len ++ ] = t % 10;
        temp = t / 10;
    }
    result.sign = sign;
    return result;
}
bign bign::operator++() {
    *this = *this + 1;
    return *this;
}
bign bign::operator++(int) {
    bign old = *this;
    ++(*this);
   return old;
}
bign bign::operator+=(const bign &num) {
    *this = *this + num;
    return *this;
}
bign bign::operator-(const bign &num) const {
    bign b = num, a = *this;
    if (!num.sign && !sign) {
        b.sign = 1;
        a.sign = 1;
        return b - a;
    }
    if (!b.sign) {
        b.sign = 1;
```

```
return a + b;
    }
    if (!a.sign) {
        a.sign = 1;
        b = bign(0) - (a + b);
        return b;
    }
    if (a < b) {
       bign c = (b - a);
       c.sign = false;
       return c;
   bign result;
   result.len = 0;
    for (int i = 0, g = 0; i < a.len; ++ i) {
        int x = a.s[i] - g;
        if (i < b.len) {
            x = b.s[i];
        }
        if (x >= 0) {
            g = 0;
        }
        else {
            g = 1;
            x += 10;
        result.s[result.len ++ ] = x;
   result.clean();
   return result;
}
bign bign::operator*(const bign &num) const {
   bign result;
   result.len = len + num.len;
   for (int i = 0; i < len; ++ i) {
        for (int j = 0; j < num.len; ++ j) {
            result.s[i + j] += s[i] * num.s[j];
       }
    }
   for (int i = 0; i < result.len; ++ i) {</pre>
        result.s[i + 1] += result.s[i] / 10;
       result.s[i] %= 10;
   result.clean();
   result.sign = !(sign ^ num.sign);
    return result;
```

```
bign bign::operator*(const int num) const {
    bign x = num;
    bign z = *this;
    return x * z;
}
bign bign::operator*=(const bign &num) {
    *this = *this * num;
   return *this;
}
bign bign::operator/(const bign &num) const {
    bign ans;
    ans.len = len - num.len + 1;
    if (ans.len < 0) {
        ans.len = 1;
       return ans;
    }
    bign divisor = *this, divid = num;
    divisor.sign = divid.sign = 1;
    int k = ans.len - 1;
    int j = len - 1;
    while (k \ge 0) {
        while (divisor.s[j] == 0) {
            -- j;
        }
        if (k > j) k = j;
        char z[N];
        memset(z, 0, sizeof(z));
        for (int i = j; i >= k; -- i) {
            z[j - i] = divisor.s[i] + '0';
        }
        bign dividend = z;
        if (dividend < divid) {</pre>
            -- k;
            continue;
        }
        int key = 0;
        while (divid * key <= dividend) {
            ++ key;
        }
        -- key;
        ans.s[k] = key;
        bign temp = divid * key;
        for (int i = 0; i < k; ++ i) {
            temp = temp * 10;
```

```
divisor = divisor - temp;
        -- k;
    }
    ans.clean();
    ans.sign = !(sign ^ num.sign);
    return ans;
}
bign bign::operator/=(const bign &num) {
    *this = *this / num;
    return *this;
}
bign bign::operator%(const bign &num) const {
    bign a = *this, b = num;
    a.sign = b.sign = 1;
    bign result, temp = a / b * b;
    result = a - temp;
    result.sign = sign;
   return result;
}
bign bign::pow(const bign &num) const {
    bign result = 1;
    for (bign i = 0; i < num; ++ i) {
        result = result * (*this);
    return result;
}
bign bign::factorial() const {
    bign result = 1;
    for (bign i = 1; i <= *this; ++ i) {
        result *= i;
    return result;
}
void bign::clean() {
    if (len == 0) {
       ++ len;
    while (len > 1 && s[len - 1] == '\0') {
       -- len;
    }
}
bign bign::Sqrt() const {
```

```
if (*this < 0)return -1;
if (*this <= 1)return *this;
bign l = 0, r = *this, mid;
while (r - 1 > 1) {
    mid = (1 + r) / 2;
    if (mid * mid > *this) {
        r = mid;
    } else {
        l = mid;
    }
}
return l;
}
```