

Distributed Model Predictive Control for Nonlinear Networked Systems with Asynchronous Communication

ZHOU Yuanqiang LI Dewei
XI Yugeng CEN Lihui

Shanghai Jiao Tong University
Key Laboratory of System Control and Information Processing, Ministry of
Education
Central South University

July 28, 2017

Introduction

Distributed model predictive control for networked nonlinear systems:

- ▶ The large-scale systems with decoupled dynamics
- ▶ The large-scale systems with coupled dynamics

This presentation considers the distributed networked nonlinear systems subject to asynchronous communication.

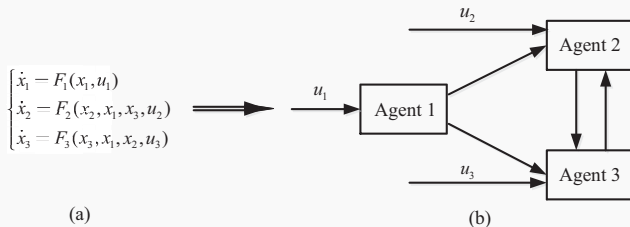


Figure: The concerned control structure

System setup

Nonlinear dynamics of each agent \mathcal{A}_i , $i \in \mathcal{I} := I[1, N]$:

$$\dot{x}_i(t) = f_i(x_i(t), u_i(t)) + g_i\left(\{x_j(t)\}_{j \in \mathcal{N}_i}\right), \quad t \geq t_0 \quad (1)$$

where $x_i \in \mathcal{X}_i \subseteq \mathbb{R}^n$, $u_i \in \mathcal{U}_i \subseteq \mathbb{R}^m$.

Define a graph: $\mathcal{G} = (\mathcal{V}, \epsilon)$ with $\mathcal{V} = \{1, \dots, N\}$,
 $\epsilon = \{(i, j) \in \mathcal{V} \times \mathcal{V}\}$.

- ▶ By the the connected graph of (7), we get the upstream neighbor set: $\mathcal{N}_i \subseteq \mathcal{I} \setminus \{i\}$;
- ▶ Through the communication of agent \mathcal{A}_i , we obtain the downstream neighbor set: $\mathcal{D}_i \subseteq \mathcal{I} \setminus \{i\}$;

Then, we have:

- ▶ $\mathcal{N}_i := \{j \in \mathcal{V} | (j, i) \in \epsilon\}$, $\mathcal{D}_i := \{j \in \mathcal{V} | (i, j) \in \epsilon\}$;
- ▶ $\mathcal{N}_i \cup \mathcal{D}_i \subset \mathcal{I}$, $i \notin \mathcal{N}_i \cap \mathcal{D}_i$;
- ▶ $j \in \mathcal{N}_i$ if and only if $i \in \mathcal{D}_j$.

System setup

Let $M = |\epsilon|$ and $d_i = |\mathcal{N}_i|$. There is at least one agent in \mathcal{I} satisfying $\mathcal{N}_i \neq \emptyset$.

$$0 < \max_{i \in \mathcal{I}} d_i \leq M \quad (2)$$

If $d_i \neq 0$, then let $\mathcal{N}_i = \{i_1, i_2, \dots, i_{d_i}\}$. Denote the sequences of state trajectories for all neighbors in \mathcal{N}_i as $x_{-i}(t)$, i.e.,

$$x_{-i}(t) = \text{vect}\{x_{i_1}(t), x_{i_2}(t), \dots, x_{i_{d_i}}(t)\} \quad (3)$$

the whole system can be described as

$$\dot{x} = f(x, u) + g(x) := F(x, u) \quad (4)$$

where $f(x, u) = \text{vect}(f_1(x_1, u_1), f_2(x_2, u_2), \dots, f_N(x_N, u_N))$ and $g(x) = \text{diag}(g_1(x_{-1}), g_2(x_{-2}), \dots, g_N(x_{-N}))$.

Preliminary results

linearized (1) around the origin, we get

$$\dot{x}_i(t) = A_i x_i(t) + B_i u_i(t) + \sum_{j \in \mathcal{N}_i} A_{ij} x_j(t), \quad t \geq t_0 \quad (5)$$

where $A_i = \partial f_i / \partial x_i|_{(0,0)}$, $B_i = \partial f_i / \partial u_i|_{(0,0)}$, $A_{ij} = \partial g_i / \partial x_j|_{(0,0)}$.

Assumption 2.

For each agent \mathcal{A}_i in (7), there exists a decoupled static feedback matrix K_i such that $A_{di} := A_i + B_i K_i$ is Hurwitz.

Lemma 1[24]. For each \mathcal{A}_i in (1) with $R_i > 0$, $Q_i > 0$, there exist $P_i > 0$ and $\varepsilon_i > 0$, satisfying: (I) $P_i A_{di} + A_{di}^T P_i = -\bar{Q}_i$; (II) $PA_o + A_o^T P \leq \bar{Q}/2$, where $\bar{Q}_i = Q_i + K_i^T R_i K_i$, $A_o = A_c - A_d$, $\Xi = \text{diag}(\Xi_1, \dots, \Xi_N)$ for any set $\Xi_i \in \{P_i, K_i, \bar{Q}_i, A_{di}\}$, such that the set $\Omega_i(\varepsilon_i) \triangleq \{x_i(t) : \|x_i(t)\|_{P_i}^2 \leq \varepsilon_i^2\}$ is a positively invariant region for $\dot{x}_i(t) = f_i(x_i(t), u_i(t)) + g_i(x_{-i}(t))$. Additionally, $x_i \in \mathcal{X}_i$, $K_i x_i \in \mathcal{U}_i$ for all $x_i \in \Omega_i(\varepsilon_i)$.

Problem description

With the different control and communication frequency, define the update time sequence $\{t_k^i, k \in \mathbb{N}\}$. At each t_k^i ,

- (1) Measure $x_i(t_k^i)$ for \mathcal{A}_i by the sampler;
- (2) Obtain $\hat{u}_i^*(\cdot; t_k^i)$ by the designed optimization problem P_i and apply to \mathcal{A}_i ;
- (3) Send out states and controls to all neighbors in \mathcal{D}_i ;

Since for agents \mathcal{A}_i and \mathcal{A}_l , $i, l \in \mathcal{I}$, $l \in \mathcal{N}_i$, $\{t_k^i\}_{k=1}^\infty \not\equiv \{t_k^l\}_{k=1}^\infty$.

The control target:

To design an effective control optimization problem P_i for agent \mathcal{A}_i which also coordinates the asynchronous communication for all its neighboring agents.

Develop an DMPC to steer the state to $\Omega_i(\varepsilon_i)$

To-be-minimized cost function $J_i := J_i(x_i(t_k^i), u_i(s; t_k^i))$ associated with (7) is defined as

$$J_i = \int_{t_k^i}^{t_k^i + T_i} \|x_i(s; t_k^i)\|_{Q_i}^2 + \|u_i(s; t_k^i)\|_{R_i}^2 ds + \|x_i(t_k^i + T_i; t_k^i)\|_{P_i}^2 \quad (6)$$

For \mathcal{A}_i at t_k^i , design an *OCP* (P_i) as

$$\hat{u}_i^*(s; t_k^i) = \arg \min_{\hat{u}_i(s; t_k^i)} J_i(x_i(t_k^i), \hat{u}_i(s; t_k^i))$$

$$s.t. \quad \dot{\hat{x}}_i(s; t_k^i) = f_i(\hat{x}_i(s; t_k^i), \hat{u}_i(s; t_k^i)) + g_i(\tilde{x}_{-i}(s; t_k^i)) \quad (7a)$$

$$\|\hat{x}_i(s; t_k^i)\|_{P_i} \leq \frac{T}{s - t_k^i} \frac{\alpha_i}{(q+1)} \varepsilon_i \quad (7b)$$

$$\hat{u}_i(s; t_k^i) \in \mathcal{U}_i \quad (7c)$$

Virtual input design with asynchronous communication

Define

$$[k]_i = \max\{l : t_l^j \leq t_k^i, j \in \mathcal{N}_i, l = 0, 1, 2, \dots\} \quad (8)$$

For $j \in \mathcal{N}_i := \{i_{d_1}, i_{d_2}, \dots, i_{d_i}\}$, $\tilde{x}_j(s; t_k^i)$ for \mathcal{A}_j is generated by following the rules:

r1 : IF $t_{[k]_i}^j + T_i \leq t_k^i$, THEN take

$$\tilde{x}_j(s; t_k^i) = x_j^L(s), \quad s \in [t_{[k]_i}^j + T_i, t_k^i + T_i] \quad (9)$$

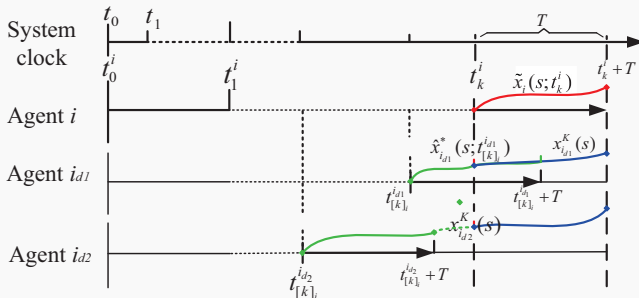
r2 : IF $t_{[k]_i}^j + T_i > t_k^i$, THEN take

$$\tilde{x}_j(s; t_k^i) = \begin{cases} \hat{x}_j^*(s; t_{[k]_i}^j), & s \in [t_k^i, t_{[k]_i}^j + T_i) \\ x_j^L(s), & s \in [t_{[k]_i}^j + T_i, t_k^i + T_i] \end{cases} \quad (10)$$

where $x_j^L(s)$ follows the solution of linearization dynamics

$$\begin{cases} \dot{x}_j^L(s) = A_{dj}x_j^L(s) \\ x_j^L(t_{[k]_i}^j + T_i) = \hat{x}_j^*(t_{[k]_i}^j + T_i; t_{[k]_i}^j) \end{cases} \quad (11)$$

Virtual input design with asynchronous communication



Consider $\mathcal{N}_i = \{i_{d1}, i_{d2}\}$,

- ▶ When $t_{[k]_i}^j + T_i \leq t_k^i$, then $\tilde{x}_j(s; t_k^i)$ is taken over by the closed-loop linearization response which ignores coupling;
- ▶ For $t_{[k]_i}^j + T_i > t_k^i$, $\tilde{x}_j(s; t_k^i)$ is constructed by the remainder of the previously predicted trajectory $\hat{x}_j^*(s; t_{[k]_i}^j)$, concatenated with $x_j^L(s)$ in (11)

Asynchronous DMPC with dual-mode strategy

The developed asynchronous DMPC approach will be improved by the dual-mode strategy.

The terminal region for \mathcal{A}_i : $\mathcal{X}_i^f := \{x_i : W_i(\hat{x}_i(t_k^i + T_i; t_k^i)) \leq \varepsilon_i^2\}$.
For agent \mathcal{A}_i at time t_k^i ,

- ▶ when $x_i(t_k^i) \notin \Omega_i(\alpha_i \varepsilon_i)$, then the optimization problem of DMPC P_i is solved and the optimized control input $u_i(s; t_k^i) = \hat{u}_i^*(s; t_k^i)$ is applied during $[t_k^i, t_{k+1}^i)$;
- ▶ when $x_i(t_k^i) \in \Omega_i(\alpha_i \varepsilon_i)$, then the control input is switched into the static feedback control law as

$$\hat{u}_i(s; t_k^i) = K_i x_i(s; t_k^i) \quad (12)$$

Algorithm 1:

- 0) *Initialization*: $x_i(t_0)$, $\hat{u}_i^*(s; t_0)$ (or $u_i(s; t_0)$) at time $t_0^i = t_0$, $\tilde{x}_i(s; t_0^i)$, $s \in [t_0^i, t_0^i + T_i]$, and $T_i \in (0, \infty)$, $q \in \{1, 2, 3, \dots\}$.
- 1) Apply $\hat{u}_i^*(s; t_k^i)$ for $s \in [t_k^i, t_k^i + T_i]$ and measure $x_i(s; t_k^i)$, if $x_i(t) \in \Omega_i(\alpha_i \varepsilon_i)$, then go to step 4).
- 2) Obtain $(t_{[k+1]i}^j, x_j(t_{[k+1]i}^j), \hat{x}_j^*(s; t_{[k+1]i}^j))$ for all \mathcal{A}_j where $j \in \mathcal{N}_i$, generate $\tilde{x}_j(s; t_{k+1}^i)$ according to (9) or (10) and stack them together into $\tilde{u}_{-i}(s; t_k^i)$.
- 3) Solve P_i in (7) to obtain $\hat{u}_i^*(s; t_{k+1}^i)$, $s \in [t_{k+1}^i, t_{k+1}^i + T_i]$, compute $\hat{x}_i^*(s; t_{k+1}^i)$, $s \in [t_{k+1}^i, t_{k+1}^i + T_i]$ and send it to all neighbors \mathcal{A}_j in \mathcal{D}_i , then go to step 1).
- 4) Generate the terminal control policy as (12) at instant time t_k^i and apply it to agent \mathcal{A}_i for all $[t_k^i, t_{k+1}^i)$.

Theorem 1 (*Boundedness*). For each agent \mathcal{A}_i with dynamics described in (7), at update time instant t_k^i , if the trajectory $\tilde{x}_j(s; t_k^i)$ is generated by following the rules (9) or (10), then we have

$$\max_{j \in \mathcal{N}_i} \|\hat{x}_j^*(s; t_k^i) - \tilde{x}_j(s; t_k^i)\|_{P_i} \leq \frac{2}{q+1} \bar{\alpha}_i \bar{\varepsilon}_i \quad (13)$$

where $s \in [t_k^i, t_k^i + T_i]$ and $\bar{\alpha}_i \bar{\varepsilon}_i = \max_{j \in \mathcal{N}_i \cup \{i\}} \{\alpha_j \varepsilon_j\}$.

Theorem 2 (*Feasibility*). For each agent \mathcal{A}_i in (7), suppose that Assumption 1-2 hold true and $x_i(t_0) \in \mathcal{X}_i^0$. If the prediction horizon T_i , the minimum update period $\Delta t_i = \min_{k \in \mathbb{N}} \{t_{k+1}^i - t_k^i\}$ and the global constant q are designed to satisfy the following inequalities

$$e^{L_{f_i} T_i} T_i \leq \frac{\alpha_i \varepsilon_i}{4d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i} \quad (14a)$$

$$\Delta t_i \geq \frac{2d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i e^{L_{f_i} T_i} T_i^2}{\alpha_i \varepsilon_i - 2d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i e^{L_{f_i} T_i} T_i} \quad (14b)$$

$$q \geq \frac{(\alpha_i - 1)L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i + (4 - 2\alpha_i)d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i e^{L_{f_i} T_i} T_i}{\alpha_i \varepsilon_i - 4d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i e^{L_{f_i} T_i} T_i} \quad (14c)$$

then the designed Algorithm 1 for \mathcal{A}_i is recursively feasible to steer the actual trajectory to the positively invariant set $\Omega_i(\varepsilon_i)$.

Remark 4: Theorem 2 indicates that the update time interval sequence $\{t_{k+1}^i - t_k^i, k \in \mathbb{N}\}$ for each \mathcal{A}_i is required to satisfy the allowable lower bound. Since

$T_i \geq t_{k+1}^i - t_k^i \geq \Delta t_i = \min_{k \in \mathbb{N}} \{t_{k+1}^i - t_k^i\}$, then the predictive horizon T is bounded as

$$e^{L_{f_i} \Delta t_i} \Delta t_i \leq e^{L_{f_i} T_i} T_i \leq \frac{\alpha_i \varepsilon_i}{4d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i} \quad (15)$$

Since q plays an important role in guaranteing the feasibility of the proposed algorithm, α_i can be used to adjust the contracting effect of the constraint for the predicted state variables.

Theorem 3 (*Stability*). For agent \mathcal{A}_i in (7), suppose that Assumption 1-2 hold, $x_i(t_0) \in \mathcal{X}_i^0$. If the conditions in Theorem 1 are satisfied and the prediction horizon T_i , the constant integer q , the shrinkage rate α_i and the minimum update time interval Δt_i are designed to make the following hold:

$$\begin{aligned} & \lambda_M \frac{4d_i L_{g_i} T_i^2 \bar{\alpha}_i \bar{\varepsilon}_i}{(q+1)^2} \left\{ d_i L_{g_i} \bar{\alpha}_i \bar{\varepsilon}_i \frac{e^{2L_{f_i}(T_i - \Delta t_i)} - 1}{2L_{f_i}} \right. \\ & \left. + \alpha_i \varepsilon_i \left[\frac{e^{2L_{f_i}(T_i - \Delta t_i)} - 1}{4L_{f_i}} + \frac{(T_i - \Delta t_i)^2}{2T_i \Delta t_i} - \frac{2}{T_i - \Delta t_i} \right] \right\} \\ & < [\lambda_m \Delta t_i + 1 - \alpha_i^2] \varepsilon_i^2 \end{aligned} \quad (16)$$

where $\lambda_M = \lambda_{\max}(P_i^{-1/2} Q_i P_i^{-1/2})$, $\lambda_m = \lambda_{\min}(P_i^{-1/2} Q_i P_i^{-1/2})$. Then by application of the proposed algorithm for each \mathcal{A}_i , the closed-loop networked system in (4) is asymptotically stable.

Remark 5: Based on Theorem 2, Theorem 3 and their proofs, we have:

Principles of choosing parameters

- S1** Calculate $\mathcal{N}_i, d_i, \mathcal{D}_i, \varepsilon_i, P_i, K_i, L_{f_i}, L_{g_i}$ for each agent \mathcal{A}_i ;
- S2** Choose $\alpha_i, \{t_k^i\}, k \in \mathbb{N}$ and T_i , calculate $\bar{\alpha}_i \bar{\varepsilon}_i$, to satisfy the condition in Theorem 1;
- S3** Choose q to make the inequality in Theorem 2 hold;
- S4** Examine whether the condition in Theorem 3 could be satisfied. Otherwise, choose a smaller α_i and go back to Step S2, until (16) is satisfied.

Application to Van der Pol oscillators

Consider a walking bipedal locomotor:

$$\left\{ \begin{array}{l} \ddot{\theta}_1(t) = 0.1[1 - 5.25\theta_1^2(t)]\dot{\theta}_1(t) - \theta_1(t) + u_1(t) \\ \ddot{\theta}_2(t) = 0.01[1 - 6070(\theta_2(t) - \theta_{2e})^2]\dot{\theta}_2(t) - 4(\theta_2(t) - \theta_{2e}) \\ \quad + 0.057\theta_1(t)\dot{\theta}_1(t) + 0.1(\dot{\theta}_2(t) - \dot{\theta}_3(t)) + u_2(t) \\ \ddot{\theta}_3(t) = 0.01[1 - 192(\theta_3(t) - \theta_{3e})^2]\dot{\theta}_3(t) - 4(\theta_3(t) - \theta_{3e}) \\ \quad + 0.057\theta_1(t)\dot{\theta}_1(t) + 0.1(\dot{\theta}_3(t) - \dot{\theta}_2(t)) + u_3(t) \end{array} \right.$$

with The constraints for $\theta_i(t)$ and $u_i(t)$ are

$$-\frac{\pi}{2} \leq \theta_i(t) - \theta_{ie} \leq \frac{\pi}{2}, \quad |u_i(t)| \leq 1, \quad \forall t \geq 0$$

To perform the control strategy, we define the states

$x_i = [x_{i1} \ x_{i2}]^T$ for dynamics in (7), where $x_{i2} = \dot{\theta}_i$, $x_{i1} = \theta_i - \theta_{ie}$ with θ_{ie} the desired constant angle for $i \in \{1, 2, 3\}$. Then, the system can be converted to

$$\begin{aligned} \mathcal{A}_1 : & \begin{cases} \dot{x}_{11}(t) = x_{12}(t) \\ \dot{x}_{12}(t) = 0.1 [1 - 5.25x_{11}^2(t)] x_{12}(t) - x_{11}(t) + u_1(t) \end{cases} \\ \mathcal{A}_2 : & \begin{cases} \dot{x}_{21}(t) = x_{22}(t) \\ \dot{x}_{22}(t) = 0.01 [1 - 6070x_{21}^2(t)] x_{22}(t) - 4x_{21}(t) + u_2(t) \\ \quad + 0.057x_{11}(t)x_{12}(t) + 0.1(x_{22}(t) - x_{32}(t)) \end{cases} \\ \mathcal{A}_3 : & \begin{cases} \dot{x}_{31}(t) = x_{32}(t) \\ \dot{x}_{32}(t) = 0.01 [1 - 192x_{31}^2(t)] x_{32}(t) - 4x_{31}(t) + u_3(t) \\ \quad + 0.057x_{11}(t)x_{12}(t) + 0.1(x_{32}(t) - x_{22}(t)) \end{cases} \end{aligned}$$

$$K_1 = \begin{bmatrix} 0.25 \\ -2.1 \end{bmatrix}, \quad K_i = \begin{bmatrix} 1.5 \\ -3.61 \end{bmatrix}, \quad i = 2, 3$$

$$A_c = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 \\ -1 & 0.1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & -4 & 0.11 & 0 & -0.1 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & -0.1 & -4 & 0.11 \end{bmatrix}$$

$$J_i(x_i(t), u_i(t)) = \int_t^{t+T_i} (30x_{i1}^2 + 30x_{i2}^2 + 0.1u_i^2) d\tau + \|x_i(t+T_i)\|_{P_i}^2$$

$$P_1 = \begin{bmatrix} 53.3 & 20 \\ 20 & 17.6 \end{bmatrix}, \quad P_2 = P_3 = \begin{bmatrix} 37.2 & 6.0 \\ 6.0 & 6.2 \end{bmatrix}$$

$$\begin{cases} L_{f_1} = 4.1 \\ L_{g_1} = 0 \end{cases} ; \quad \begin{cases} L_{f_i} = 4 \\ L_{g_i} = 0.1 \end{cases}, \quad i = 2, 3$$

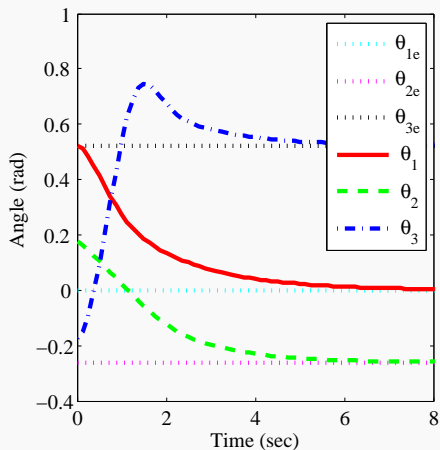


Figure: The closed-loop response of state one

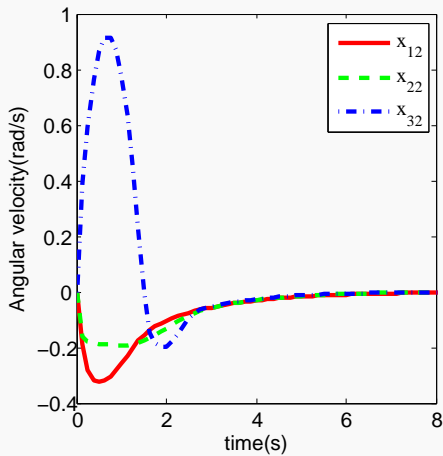


Figure: The closed-loop response of state two

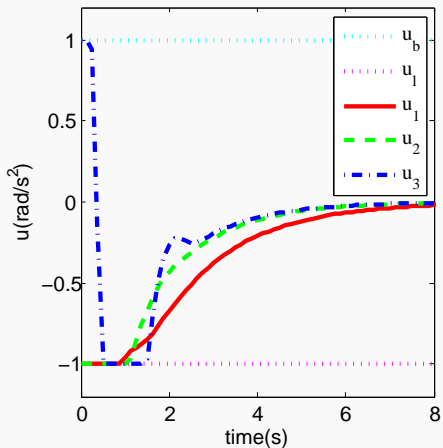


Figure: The inputs of closed-loop system

Conclusion

The basic idea behind this algorithm is to utilize the information associated with the interconnected upstream neighbors of complex networked systems to a distributed design.

Questions ?

Thank you!