

Topics:

- 1) Linked lists
- 2) Stack
- 3) Queue

Linked lists

Qn: You are given a sequence of elements given some queries (deletion of element). For each query In what time complexity can you do it in terms of size of array n ?

Trivial idea: I store them in an array then



Takes in worst case (when we delete first element) $O(n)$ time.

With linked list this could be done in $O(1)$

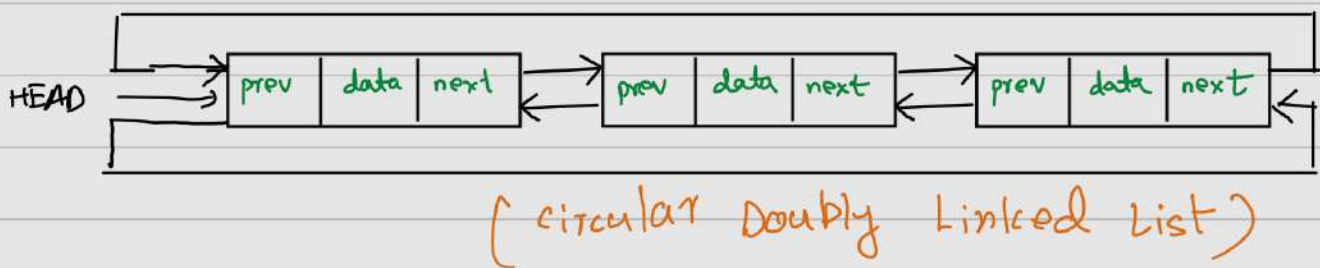
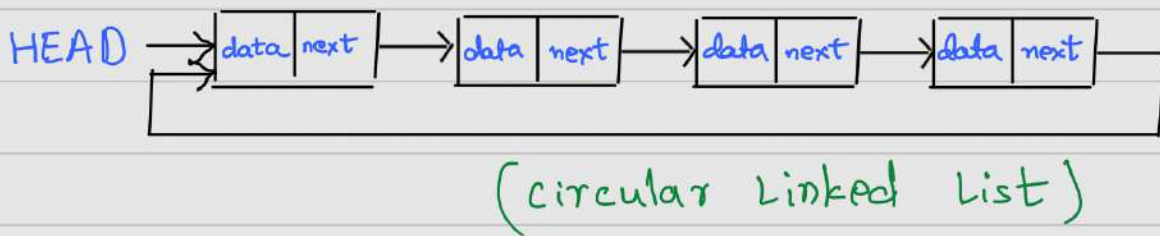
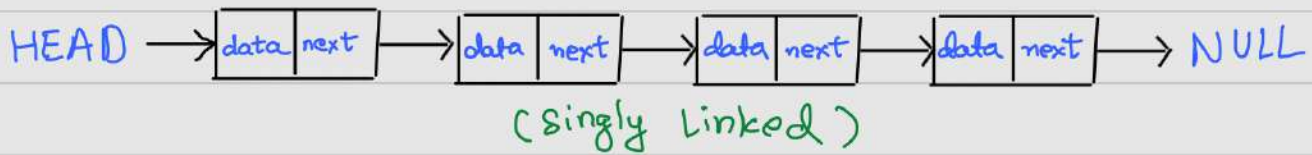
Similarly insertion in a sequence can be done in $O(1)$.
(Think how to do this in array?)

What is Linked list?

A linear list consisting of a series of connected nodes.

Node = $\begin{cases} \text{Data} \\ \text{Address of next node (pointer)} \end{cases}$

Types: Singly Linked, Doubly Linked, Circular Linked list



Representation

```
class Node {  
    public:  
        int data;  
        Node* next;  
}; // Singly Linked
```

```
class Node {  
    public:  
        int data;  
        Node* next;  
        Node* prev;  
}; // Doubly Linked
```

★ Main power: Power to break and rejoin the chain

Linked list Operations: (Singly Linked lists)

★ Traversal

★ Insertion

★ Deletion

TRAVERSAL:

- * HEAD points to 1st node.
- * NEXT pointer of last node is NULL.

If you have the HEAD you can traverse whole list.

Pseudocode:



```
Node* temp = HEAD;
while (temp != NULL)
{
    cout << temp->data << " ";
    temp = temp->next;
}
```

* INSERTION

- allocate memory for new node
- store data
- store appropriate location in the next field } VVI

Insert at beginning

Pseudocode:

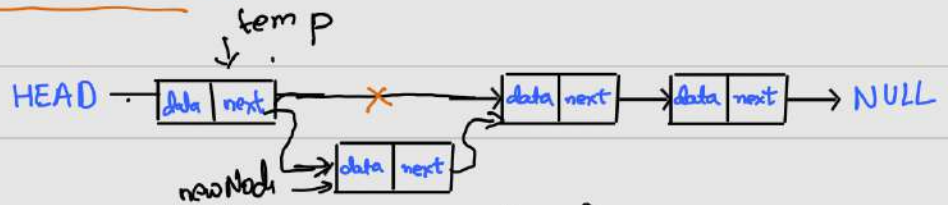


```
Node* newNode = malloc(sizeof(Node)); // new Node;
newNode->data = 4;
newNode->next = HEAD;
HEAD = newNode;
```

Similarly we can do for last Node; (Always make rough sketches of which pointer has to go where)

Insert anywhere in middle

Pseudocode:



```
Node* newNode = malloc(sizeof(Node));
newNode->data = 4;
```

```
Node* temp = head;
```

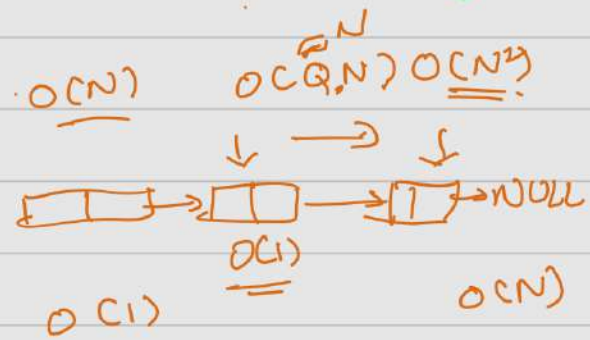
```
for (int i = 1; i < position; i++)
```

```
{ if (temp->next != NULL)
```

```
{ temp = temp->next; }
```

```
}
```

// position is acc. to 1-based indexing



```
newNode->next = temp->next;
temp->next = newNode;
```

As you can feel this code does insertion in $O(N)$ time. Then how is linked list useful?

When we talk about insertion & deletion then we assume that we have a direct access to position like in array we have index & in linked list we have ptr of element where insertion takes place & thereafter we calculate the time taken.

E.g.- Deletion of duplicate elements of a sequence

DELETION

- Just set next field of previous node to the node which is present just after the node to be deleted.

delete from anywhere in middle

- traverse to its before and do as said above.

Pseudocode:



```
Node* temp = HEAD;
```

```
for( int i=1 ; i < pos ; i++ )
```

```
{ if (temp->next != NULL)
```

```
{ temp = temp->next; }
```

```
}
```

```
temp->next = temp->next->next;
```

Doubly Linked List

Node {
 *prev
 data
 *next

Representation :

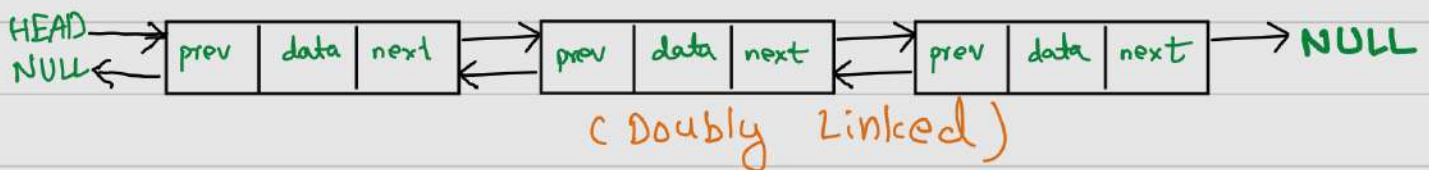
```
class Node {
```

```
public :
```

```
int data;
```

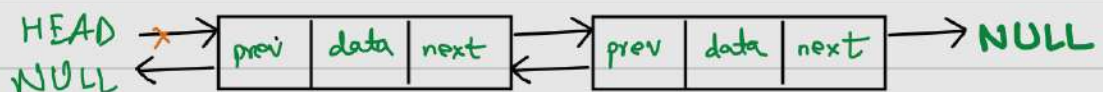
```
Node* next;
```

```
Node* prev };
```



INSERTION :

Insertion at beginning:



```
HEAD->prev = new Node
```



```
newNode->next = HEAD;
```

- Create a New node , Set its pointers accordingly

Pseudocode:

```
Node* newNode = malloc (sizeof(Node));
```

```
newNode → data = 4;
```

```
newNode → prev = NULL;
```

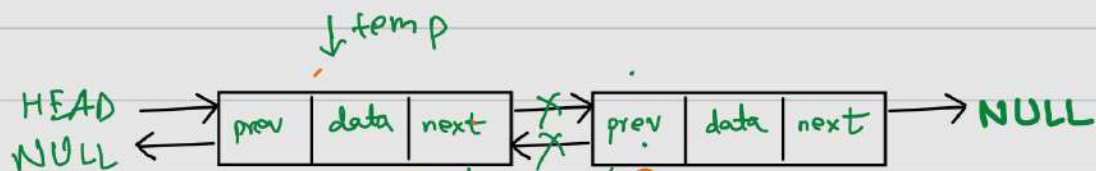
```
newNode → next = HEAD;
```

```
HEAD → prev = newNode
```

```
HEAD = newNode
```



Insertion anywhere in middle



$newNode \rightarrow next = temp \rightarrow next$

$newNode \rightarrow prev = temp;$

$temp \rightarrow next \rightarrow prev = newNode;$

$temp \rightarrow next = newNode;$

Pseudocode:

```
Node* newNode = malloc (sizeof(Node));
```

```
newNode → val = 4;
```

```
Node* temp = HEAD;
```

```
for(int i=1; i<pos; i++)
```

```
{ if (temp → next != NULL)
```

```
{ temp = temp → next; }
```

```
}
```



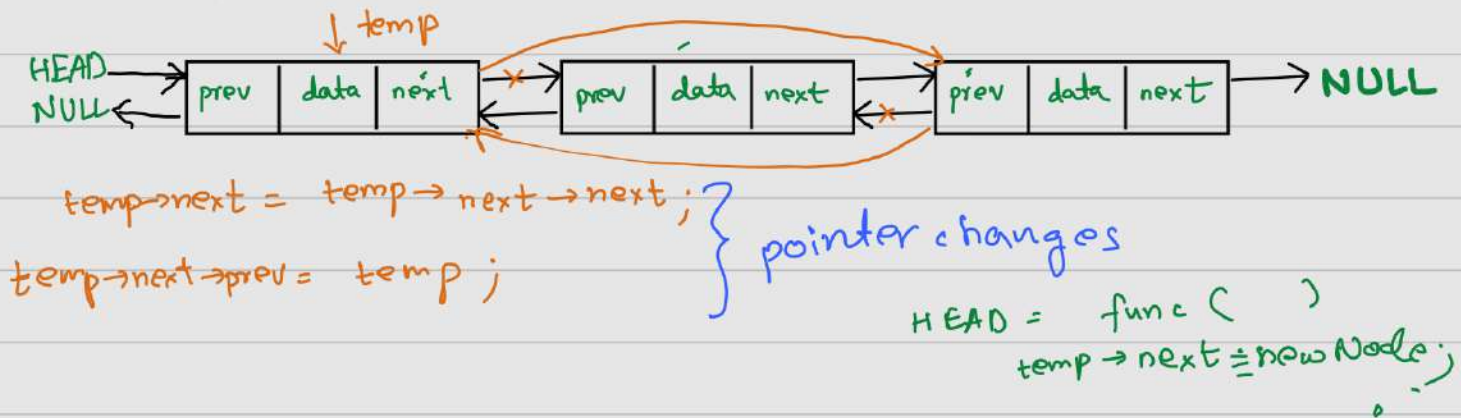
```

newNode → next = temp → next;
newNode → prev = temp;
temp → next → prev = newNode;
temp → next = newNode;

```

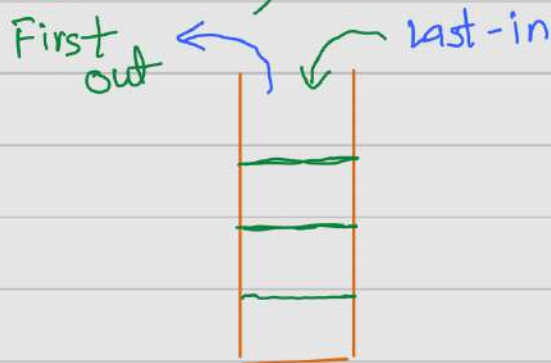
DELETION

Deletion of inner node



STACK

Linear data structure following principle of (LIFO)
(last-in First-out)



Putting an item on top is called "push" & removing is called "pop"

Main Operations:

★ Push - adds element on top