

**CSCI 335**  
**Software Design and Analysis III**  
**Lecture 11 Part 2: Hashing**

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## Agenda

- Hash tables
- Hash tables without Linked Lists
  - Linear Probing
  - Quadratic Probing
- Hash tables with Linked Lists
- Separate Chaining

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## Hash Applications

- Symbol tables (compilers)
- Graphs where nodes are strings (e.g. names of cities)
- Spell-checkers
- Password checking
- • ...

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## Hashing

- Several methods of implementing hash table.
- Compare these methods analytically.
- Show numerous applications of hashing.
- Compare hash tables with binary trees.

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## Hashing

- Hash Table ADT

- An array of some fixed size containing the items.
- Implementation is called hashing.
- Used for insertions, deletions, and finds in constant average time.
- Ordering operations are not supported efficiently.
  - findMin, findMax and printing entire table in linear time not supported.

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- Find Ada: Linear search
- Find Ada : When index known
  - myData = Array(8)

Ada

Jan	Tim	Mia	Sam	Leo	Ted	Bea	Lou	Ada	Max	Zoe
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
-----	---	----	---	-----	---	----	-----	---

0	1	2	3	4	5	6	7	8	9	10
---	---	---	---	---	---	---	---	---	---	----

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
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0	1	2	3	4	5	6	7	8	9	10
---	---	---	---	---	---	---	---	---	---	----

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1

Mia										
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1

Tim Mia										
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1

Bea  
Zoe  
Jan  
Ada  
Leo  
Sam  
Lou  
Max  
Ted

<https://www.asciitable.xyz/>

Tim Mia										
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2
Sam	S	83	a	97	m	109	289	3
Lou	L	76	o	111	u	117	304	7
Max	M	77	a	97	x	120	294	8
Ted	T	84	e	101	d	100	285	10

Tim Mia										
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2
Sam	S	83	a	97	m	109	289	3
Lou	L	76	o	111	u	117	304	7
Max	M	77	a	97	x	120	294	8
Ted	T	84	e	101	d	100	285	10

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Index number = sum\_ASCII\_codes Mod array\_size

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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Find Ada 262 Mod 11 = 9

myData = Array(9)

Ada
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Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Acknowledgement for example: Kevin Drumm, Computer Science series

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## Hashing

- Data is stored in a table (array/vector).
  - Let us call size of table T and number of elements stored in table N.
  - Load factor  $\lambda = N / T$  (helps to analyze performance).
- Ideally  $O(1)$  find/insert/delete.
  - But no ability for sorting-based features.
- Hash function gives index in constant time:
  - Ideally simple to compute and ensure 2 different keys get different cells.
  - In practice, should distribute keys evenly among cells
  - Hash Function: Key  $\rightarrow \{0, 1, 2, \dots, T-1\}$ .
  - Key: Integer/String/etc. – depending on application.
  - Mapping from keys to indices is not unique  $\Rightarrow$ 
    - Collision resolution is required.

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## Hashing Algorithm

- Calculation applied to a key to transform it into an address.
- For numeric keys, divide the key by the number of available addresses,  $n$ , and take the remainder  

$$\text{address} = \text{key} \bmod n$$
- For alphanumeric keys, divide the sum of ASCII codes in a key by the number of available addresses,  $n$  and take their remainder.
- Folding method divides key into equal parts then add the parts together
  - Phone number 1-212-650-5031 becomes  $01+21+26+50+50+31=179$
  - Can divide by some constant and take the remainder

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## Collision Resolution

- If when an element is inserted, it hashes to the same value as an already inserted element
- Collision needs to be resolved.
- Simple methods:
  - Open addressing
  - Closed addressing
- Alternatives:
  - Cuckoo Hashing
  - Hopskotch Hashing

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## Collisions

- Open addressing
  - Linear probing
  - Plus 3 rehash
  - Quadratic probing (failed attempts)<sup>2</sup>
  - Double hashing
- Closed addressing
  - Separate chaining

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## Open Addressing with Linear probing

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## Open Addressing with Linear probing

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4

Bea	Tim			Mia	Zoe						
0	1	2	3	4	5	6	7	8	9	10	

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## Open Addressing with Linear probing

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1
Moe	M	83	o	111	e	101	289	3
Lou	L	76	o	111	u	117	304	7
Rae	R	82	a	97	e	101	280	5
Max	M	77	a	97	x	120	294	8
Tod	T	84	o	111	d	100	295	9

Bea	Tim			Mia	Zoe	Sue					
0	1	2	3	4	5	6	7	8	9	10	

Acknowledgement for example: Kevin Drumm, Computer Science series

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## Open Addressing with Linear probing

Mia	M	77	i	105	a	97	279	4
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Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
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Rae	R	82	a	97	e	101	280	5
Max	M	77	a	97	x	120	294	8
Tod	T	84	o	111	d	100	295	9

Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod	
0	1	2	3	4	5	6	7	8	9	10	

Acknowledgement for example: Kevin Drumm, Computer Science series

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## Open Addressing with Linear probing

Find Rae  $280 \text{ Mod } 11 = 5$ 

myData = Array(5)

Rae

$$\text{Load Factor} = \frac{\text{total \# of items}}{\text{size of array}}$$

Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod	
0	1	2	3	4	5	6	7	8	9	10	

Acknowledgement for example: Kevin Drumm, Computer Science series

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## Summary

- More about other collision resolution methods next lecture
- Used to index large amounts of data
- Address of each key calculated using the key itself.
- Collisions resolved when open or closed addressing
- Hashing is widely used in database indexing, compilers, caching, password authentication and more
- Insertion, deletion and retrieval occur in constant time.

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