

ASP for Consistent Query Answering

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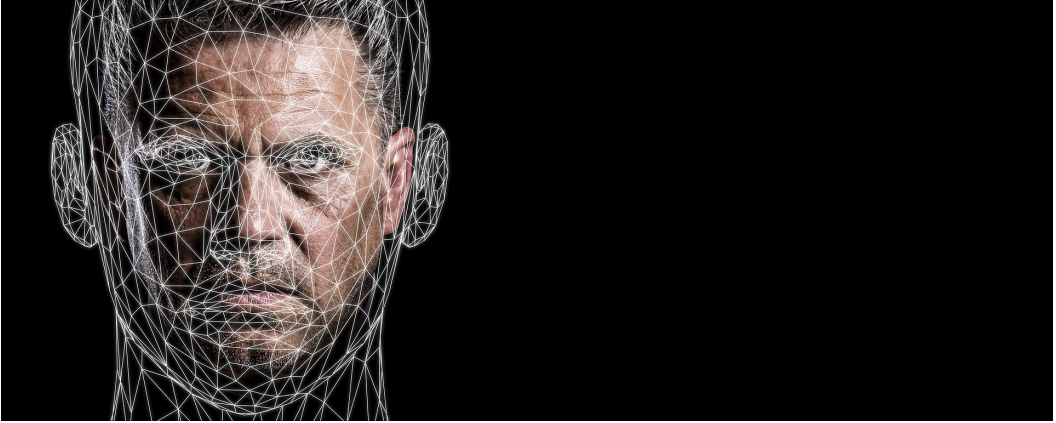


Fig. 1

Consistent query answering for inconsistent databases is a running problem...

CCS Concepts: • **Information systems** → **Database design and models**; **Database query processing**.

Additional Key Words and Phrases: Answer Set Programming, Consistent Query Answering

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1 INTRODUCTION

The aim of this article is to present a fair comparison between two methods for solving the problem of CERTAINTY(q). Considering an inconsistent database, a repair is a maximal set of tuples from this database that respects his constraints. The CERTAINTY(q) problem consists in answering the question of knowing if it exists a repair that falsifies the query. Depending on the query, the CERTAINTY(q) problem can have a first order complexity. For the queries that are in first order, we want to compare the efficiency of the generate-and-test method and of the first order rewriting method.

The comparison is realised here on few queries with ASP. For each query, we have measured the execution times of the two methods on databases of different sizes, while distinguishing the yes-instances (databases for which the CERTAINTY(q) problem is true) and the no-instances.

2 CHOSEN QUERIES

To make a one to one comparison with the results found by Akhil A.Dixit and Phokion G.Kolaitis in their "A SAT-Based System for Consistent Query Answering", we decided to reuse the same FO-rewritable queries they used to prove that the KW-fo rewriting can be more efficient by using ASP instead of SQL.

For an easiest implementation, we remove the free variables of the queries. At the end, here are the queries used for the tests we performed.

$$\begin{aligned} q_1 &:= \exists x, y, z, v, w (R_1(\underline{x}, y, z) \wedge R_2(\underline{y}, v, w)) \\ q_2 &:= \exists x, y, z, v, u, p (R_1(\underline{x}, y, z) \wedge R_3(\underline{y}, v) \wedge R_2(\underline{v}, u, p)) \\ q_3 &:= \exists x, y, z, v, u, (R_1(\underline{x}, y, z) \wedge R_2(\underline{y}, v, d)) \end{aligned}$$

Notice that d is a constant in q_3 .

3 QUERIES IMPLEMENTATION IN ASP

3.1 First query

FO Rewriting:

```
q1:-r1(X,Y,Z),not p1(X).
p1(X):-r1(X,Y,Z),not p2(Y).
p2(Y):-r2(Y,V,W).
```

```
certainty:-q1.
```

```
certainty:-not certainty.
```

```
#show certainty/0.
```

Generate and Test:

```
1{rr1(X,Y,Z):r1(X,Y,Z)}1:-r1(X,_,_).
1{rr2(X,Y,Z):r2(X,Y,Z)}1:-r2(X,_,_).
:-rr1(X,Y,Z),rr2(Y,V,W).
```

```
#show .
```

3.2 Second query

FO Rewriting:

```

q1 :- r1(X,_,_), not p1(X).
p1(X) :- r1(X,Y,_), not q3(Y).
q3(Y) :- r3(Y,_), not q2(Y).
q2(Y) :- r3(Y,V), not q1(V).
q1(V) :- r2(V,_,_).

```

```

certainty:-q1.
certainty:-not certainty.

```

```
#show certainty / 0.
```

Generate and Test:

```

1{ rr1(X,Y,Z): r1(X,Y,Z)}1:- r1(X,_,_).
1{ rr2(X,Y,Z): r2(X,Y,Z)}1:- r2(X,_,_).
1{ rr3(X,Y): r3(X,Y)}1:- r3(X,_).
:- rr1(X,Y,Z), rr3(Y,V), rr2(V,U,D).

```

```
#show .
```

3.3 Third query

FO Rewriting:

```

q1:-r1(X,Y,Z), not p1(X).
p1(X):-r1(X,Y,Z), not p2(Y).
p2(Y):-r4(Y,V,W),W=w.

```

```

certainty:-q1.
certainty:-not certainty.

```

```
#show certainty / 0.
```

Generate and Test:

```

1{ rr1(X,Y,Z): r1(X,Y,Z)}1:- r1(X,_,_).
1{ rr4(X,Y,Z): r4(X,Y,Z)}1:- r4(X,_,_).
:- rr1(X,Y,Z), rr4(Y,V,w).

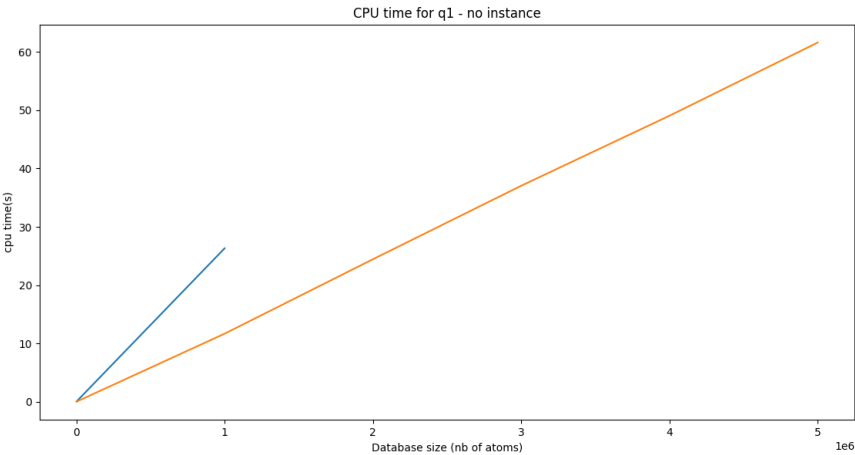
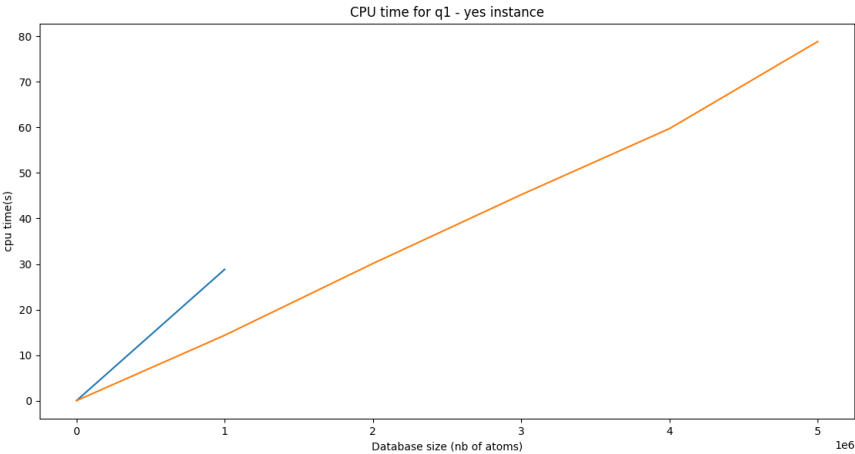
```

```
#show .
```

4 RESULTS

For each graph, the blue line corresponds to the time taken by the generate-and-test method and the orange line to the time taken by the FO method. When a result is not in the graph, that means that the execution of the program was interrupted for insufficient memory. For example, the times

for the generate-and-test for the yes-instance for q_1 for the db sizes granter than 1 million are absent.



We see that the fo rewriting leads to better results, in terms of cpu time, that the generate-and-test method.

5 CONCLUSION

This project was very fun

ACKNOWLEDGMENTS

