

# Language Models and Recurrent Neural Networks

COSC 7336: Advanced Natural Language Processing  
Fall 2017

# Announcements

- ★ Next Week no class (Mid-Semester Bash)
- ★ Assignment 2 is out
- ★ Mid term exam Oct. 27th
- ★ Paper presentation sign up coming up
- ★ Reminder to submit report, code and notebook for Assignment 1 today!

# Today's lecture

- ★ Short intro to language models
- ★ Recurrent Neural Networks
- ★ Demo: RNN for text

# Language Models

- ★ They assign a probability to a sequence of words:
  - Machine Translation:
    - $P(\text{high winds tonite}) > P(\text{large winds tonite})$
  - Spell Correction:
    - The office is about fifteen **minuets** from my house
    - $P(\text{about fifteen minutes from}) > P(\text{about fifteen minuets from})$
  - Speech Recognition
    - $P(\text{I saw a van}) \gg P(\text{eyes awe of an})$
  - Summarization, question-answering, OCR correction and many more!

# More Formally

- ★ Given a sequence of words predict the next one:
  - $P(w_5|w_1, w_2, w_3, w_4)$
- ★ Predict the likelihood of a sequence of words:
  - $P(W) = P(w_1, w_2, w_3, w_4, w_5 \dots w_n)$
- ★ How do we compute these?

# Chain Rule

- ★ Recall the definition of conditional probabilities:
  - $p(B|A) = P(A,B)/P(A)$  Rewriting:  $P(A,B) = P(A)P(B|A)$
- ★ More variables:
  - $P(A,B,C,D) = P(A)P(B|A)P(C|A,B)P(D|A,B,C)$
- ★ The Chain Rule in General
  - $P(x_1, x_2, x_3, \dots, x_n) = P(x_1)P(x_2|x_1)P(x_3|x_1, x_2) \dots P(x_n|x_1, \dots, x_{n-1})$

# Back to sequence of words

$P(\text{"I am the fire that burns against the cold"}) = P(I) \times P(\text{am}|I) \times P(\text{the}|I \text{ am}) \times P(\text{fire}|I \text{ am the}) \times P(\text{that}|I \text{ am the fire}) \times P(\text{burns}|I \text{ am the fire that}) \times P(\text{against}|I \text{ am the fire that burns}) \times P(\text{the}|I \text{ am the fire that burns against}) \times P(\text{cold}|I \text{ am the fire that burns against the})$

★ How do we estimate these probabilities?

$count(I \text{ am the fire that burns against the cold})$

$count(I \text{ am the fire that burns against the})$

# We shorten the context (history)

## Markov Assumption:

$P(\text{cold} \mid \text{I am the fire that burns against the}) \approx P(\text{cold} \mid \text{burns against the})$

This is:  $P(w_n \mid w_1^{n-1}) \approx P(w_n \mid w_{n-N+1}^{n-1})$

When  $N = 1$ , this is a unigram language model:

$$P(w_1 w_2 \dots w_n) \approx \prod_i P(w_i)$$

When  $k = 2$ , this is a bigram language model:

$$P(w_1^n) \approx \prod_{k=1}^n P(w_k \mid w_{k-1})$$



# Are all our problems solved?

# Zeros

Training set:

- ... denied the allegations
- ... denied the reports
- ... denied the claims
- ... denied the request

Test set:

- ... denied the offer
- ... denied the loan

$P(\text{"offer"} \mid \text{denied the}) = 0$

# Laplace Smoothing

- ★ Add one to all counts

- ★ Unigram counts:

$$P(w_i) = \frac{c_i}{N}$$

- ★ Laplace counts:

$$P_{\text{Laplace}}(w_i) = \frac{c_i + 1}{N + V}$$

- ★ Disadvantages:

  - Drastic change in probability mass

- ★ But for some tasks Laplace smoothing is a reasonable choice

# Leveraging Hierarchy of N-grams: **Backoff**

- ★ Adding one to all counts is too drastic
- ★ What about relying on shorter contexts?
  - Let's say we're trying to compute  $P(\text{against} \mid \text{that burns})$ , but  $\text{count}(\text{that burns against}) = 0$
  - We backoff to a shorter context:  $P(\text{against} \mid \text{that burns}) \approx P(\text{against} \mid \text{burns})$
- ★ We backoff to a lower n-gram
- ★ Katz Backoff (**discounted backoff**)

$$P_{\text{BO}}(w_n | w_{n-N+1}^{n-1}) = \begin{cases} P^*(w_n | w_{n-N+1}^{n-1}), & \text{if } C(w_{n-N+1}^n) > 0 \\ \alpha(w_{n-N+1}^{n-1}) P_{\text{BO}}(w_n | w_{n-N+2}^{n-1}), & \text{otherwise.} \end{cases}$$

# Leveraging Hierarchy of N-grams: Interpolation

- ★ Better idea: why not always rely on lower order N-grams?

$$\begin{aligned}\hat{P}(w_n|w_{n-2}w_{n-1}) = & \lambda_1 P(w_n|w_{n-2}w_{n-1}) \\ & + \lambda_2 P(w_n|w_{n-1}) \\ & + \lambda_3 P(w_n)\end{aligned}\quad \text{Subject to: } \sum_i \lambda_i = 1$$

- ★ Even better, condition on context:

$$\begin{aligned}\hat{P}(w_n|w_{n-2}w_{n-1}) = & \lambda_1(w_{n-2}^{n-1})P(w_n|w_{n-2}w_{n-1}) \\ & + \lambda_2(w_{n-2}^{n-1})P(w_n|w_{n-1}) \\ & + \lambda_3(w_{n-2}^{n-1})P(w_n)\end{aligned}$$

# Absolute Discounting

Bigram count in training	Bigram count in held out set
0	.0000270
1	0.448
2	1.25
3	2.24
4	3.23
5	4.21
6	5.23
7	6.21
8	7.21
9	8.26

For each bigram count in training data, what's the count in held out set?

Approx. a 0.75 difference!

# Absolute Discounting

$$P_{\text{AbsoluteDiscounting}}(w_i|w_{i-1}) = \frac{C(w_{i-1}w_i) - d}{\sum_v C(w_{i-1}v)} + \lambda(w_{i-1})P(w_i)$$

# How much do we want to trust unigrams?

- ★ Instead of  $P(w)$ : “How likely is  $w$ ”
- ★  $P_{\text{continuation}}(w)$ : “How likely is  $w$  to appear as a novel continuation?”
- ★ For each word, count the number of bigram types it completes

$$P_{\text{CONTINUATION}}(w) = \frac{|\{v : C(vw) > 0\}|}{|\{(u', w') : C(u'w') > 0\}|}$$



# Interpolated Kneser-Ney

★ **Intuition:** Use estimate of  $P_{\text{CONTINUATION}}(w_i)$

$$P_{\text{KN}}(w_i | w_{i-n+1}^{i-1}) = \frac{\max(c_{\text{KN}}(w_{i-n+1}^i) - d, 0)}{\sum_v c_{\text{KN}}(w_{i-n+1}^{i-1} v)} + \lambda(w_{i-n+1}^{i-1}) P_{\text{KN}}(w_i | w_{i-n+2}^{i-1})$$

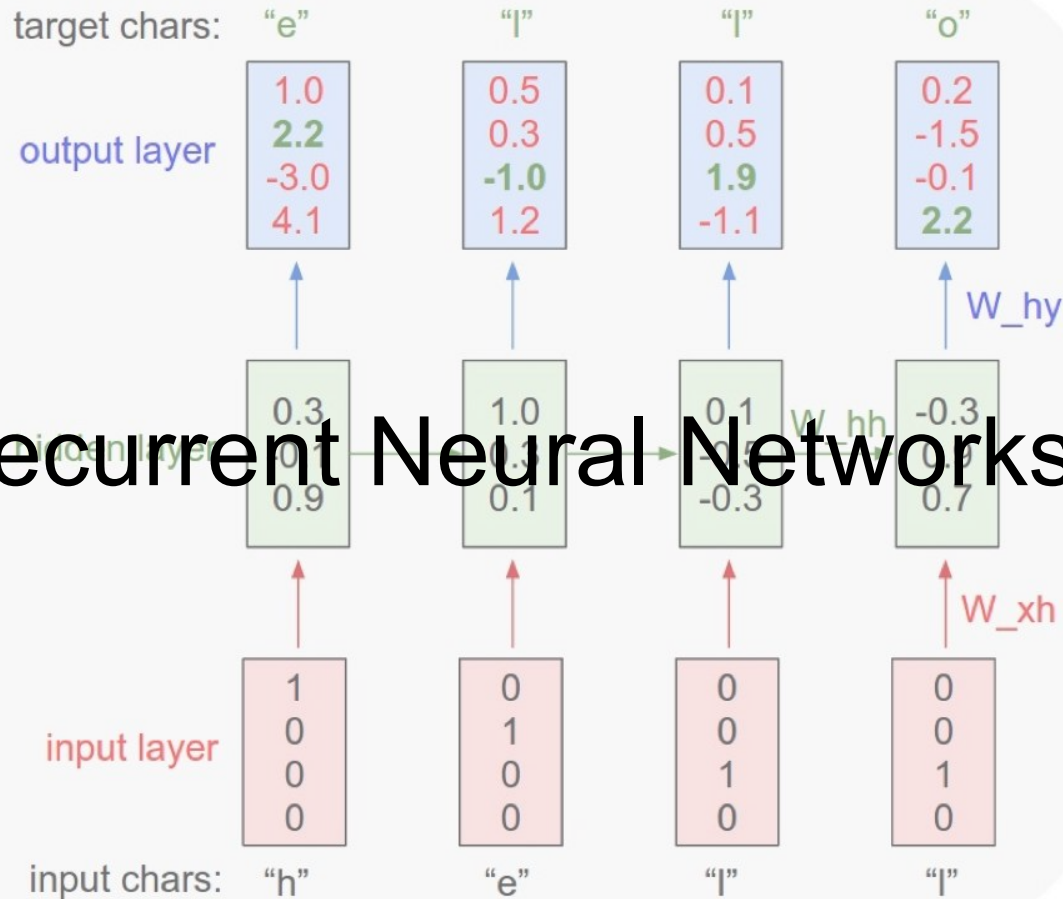
# Evaluating Language Models

- ★ Ideal: Evaluate on end task (extrinsic)
- ★ Intrinsic evaluation: use perplexity:

$$\text{PP}(W) = \sqrt[N]{\prod_{i=1}^N \frac{1}{P(w_i | w_1 \dots w_{i-1})}}$$

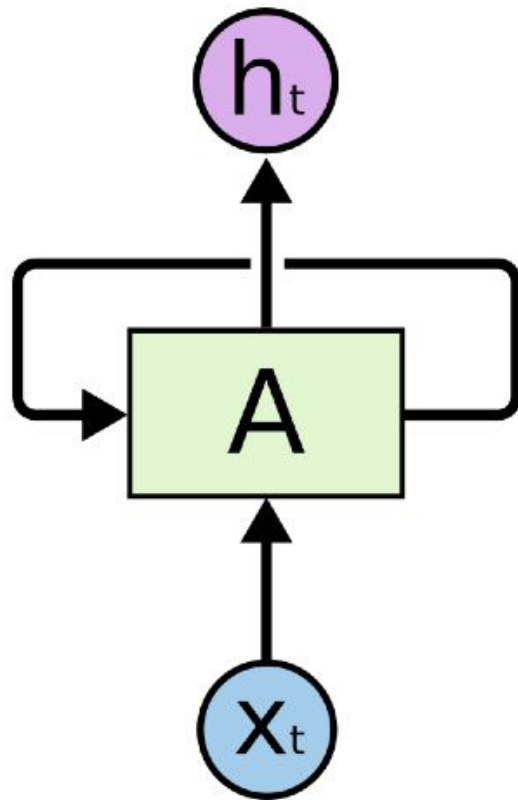
- ★ Perplexity is inversely proportional to the probability of W

# Recurrent Neural Networks

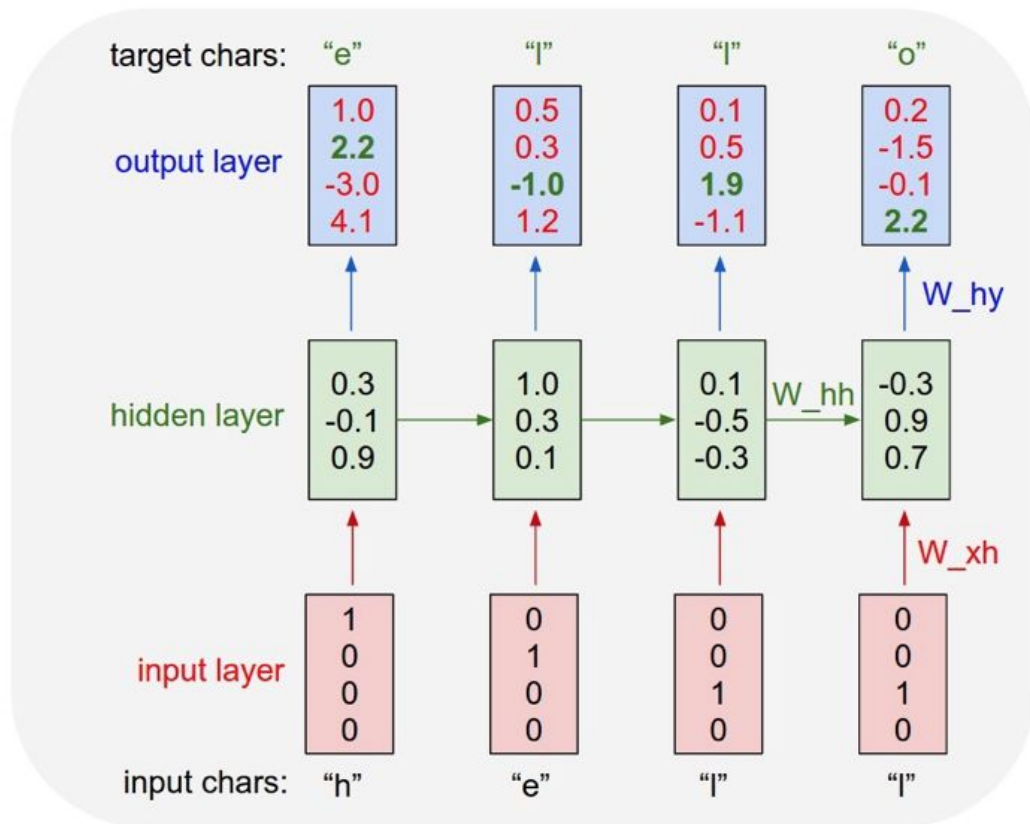


# Recurrent NNs

- ★ Neural networks with memory
- ★ Feed-forward NN: output exclusively depends on the current input
- ★ Recurrent NN: output depends on current and previous states
- ★ This is accomplished through lateral/backward connections which carry information while processing a sequence of inputs



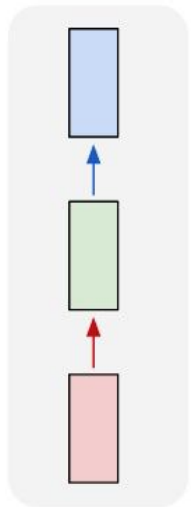
# Character-level language model



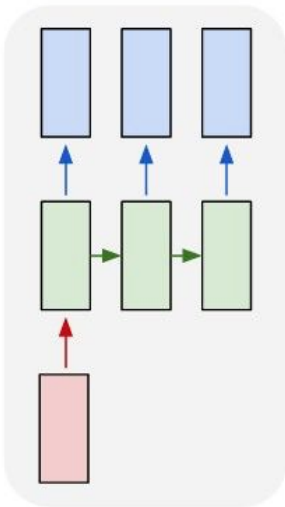
(source:  
<http://karpathy.github.io/2015/05/21/rnn-effectiveness/>)

# Sequence learning alternatives

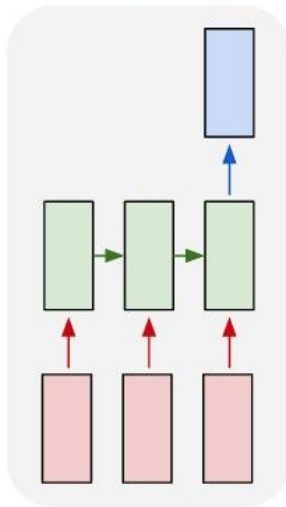
one to one



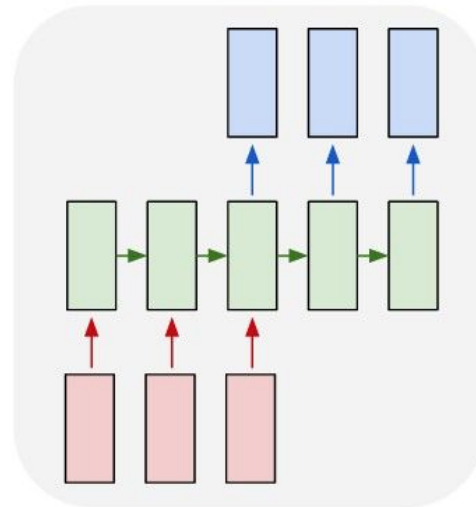
one to many



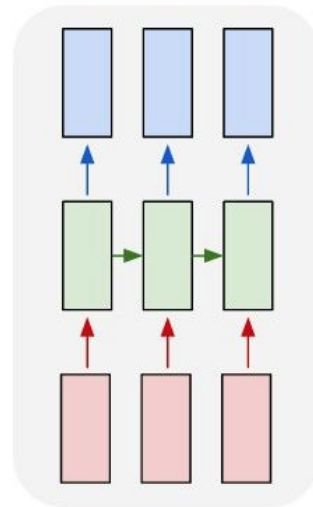
many to one



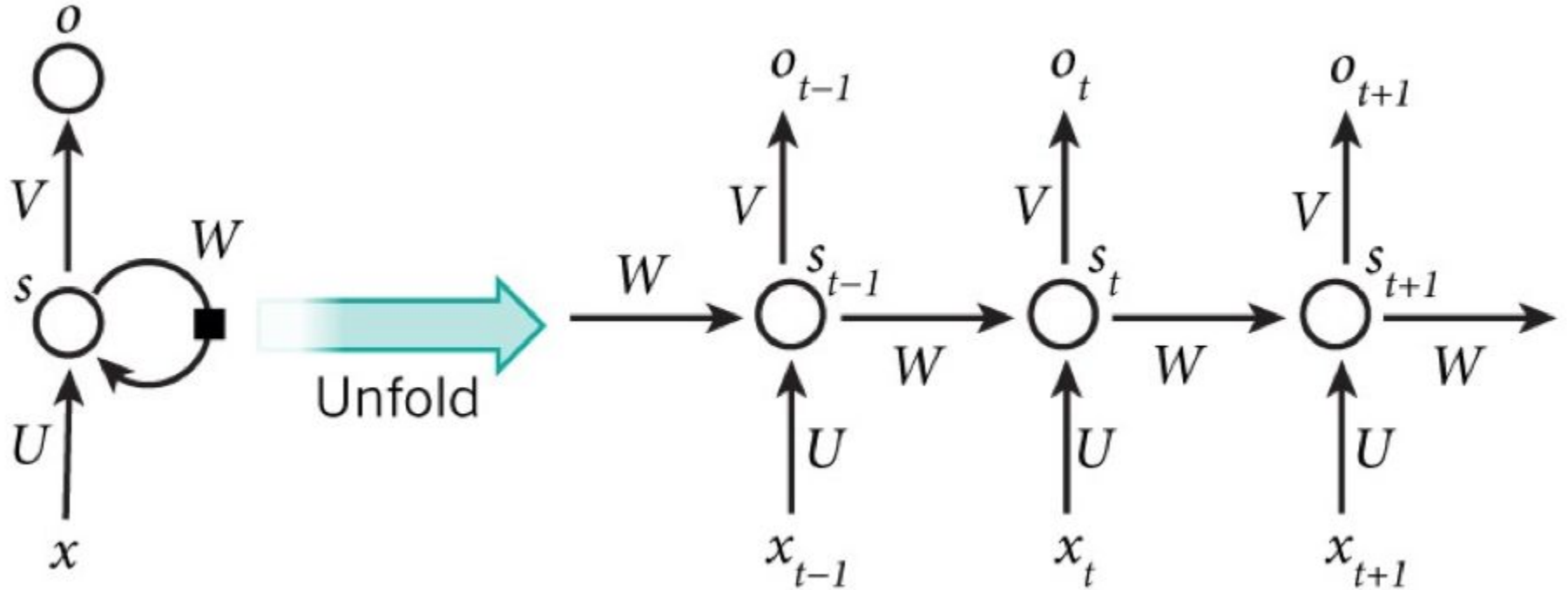
many to many



many to many

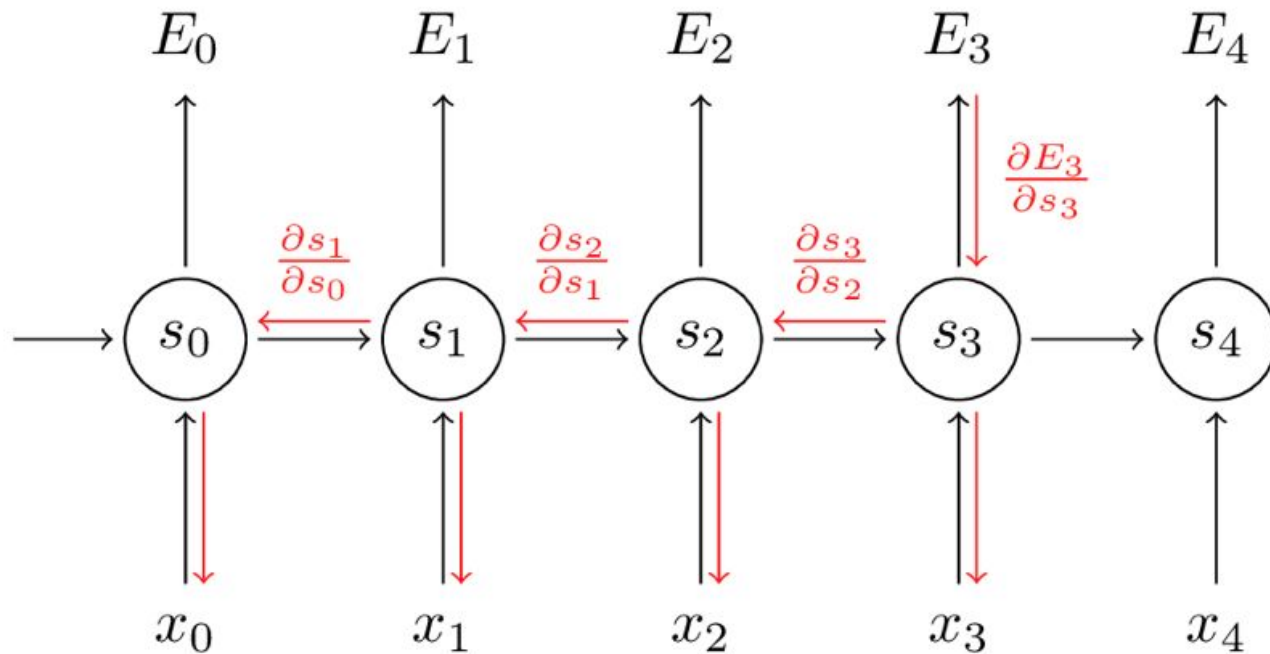


# Network unrolling



(source: <http://www.wildml.com/2015/09/recurrent-neural-networks-tutorial-part-1-introduction-to-rnns/>)

# Backpropagation through time (BPTT)

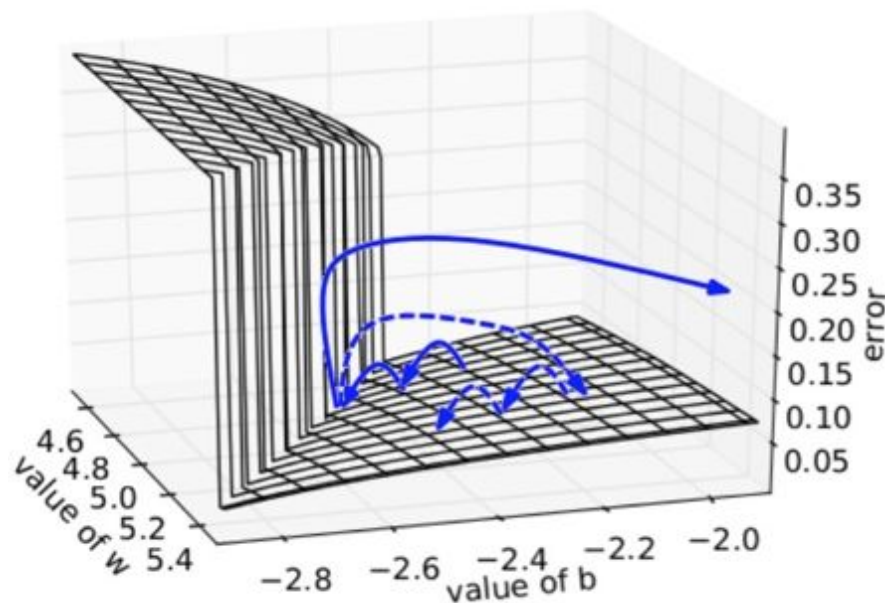


(source: <http://www.wildml.com/2015/10/recurrent-neural-networks-tutorial-part-3-backpropagation-through-time-and-vanishing-gradients/>)



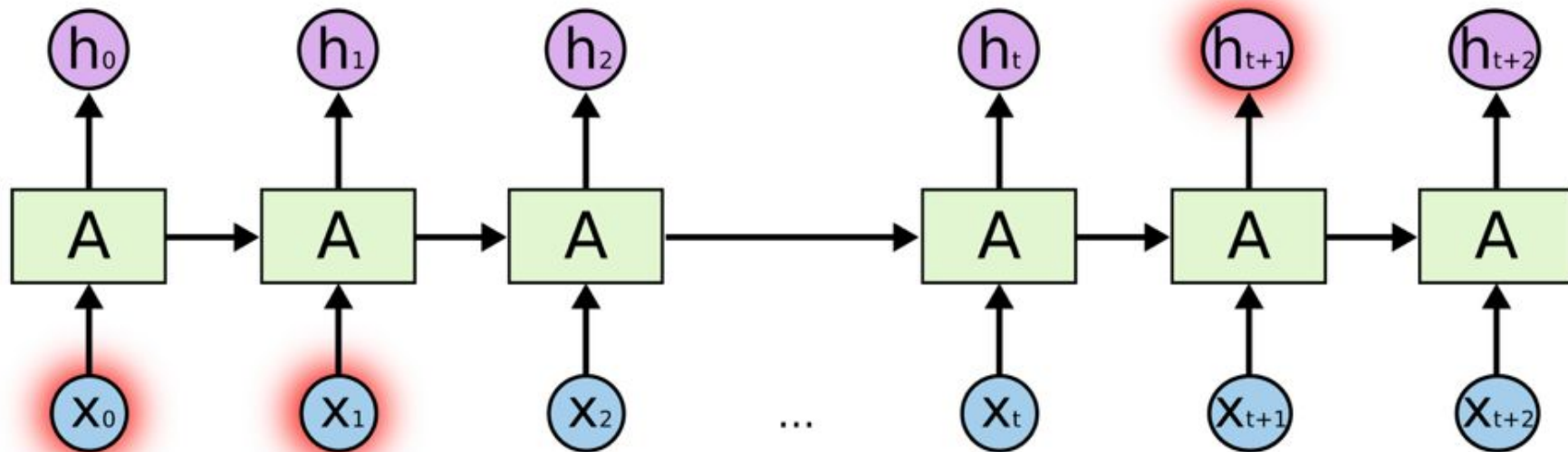
# BPTT is hard

- ★ The vanishing and the exploding gradient problem
- ★ Gradients could vanish (or explode) when propagated several steps back
- ★ This makes difficult to learn long-term dependencies.



Razvan Pascanu, Tomas Mikolov, and Yoshua Bengio. 2013. On the difficulty of training Recurrent Neural Networks. Proc. of ICML, abs/1211.5063.

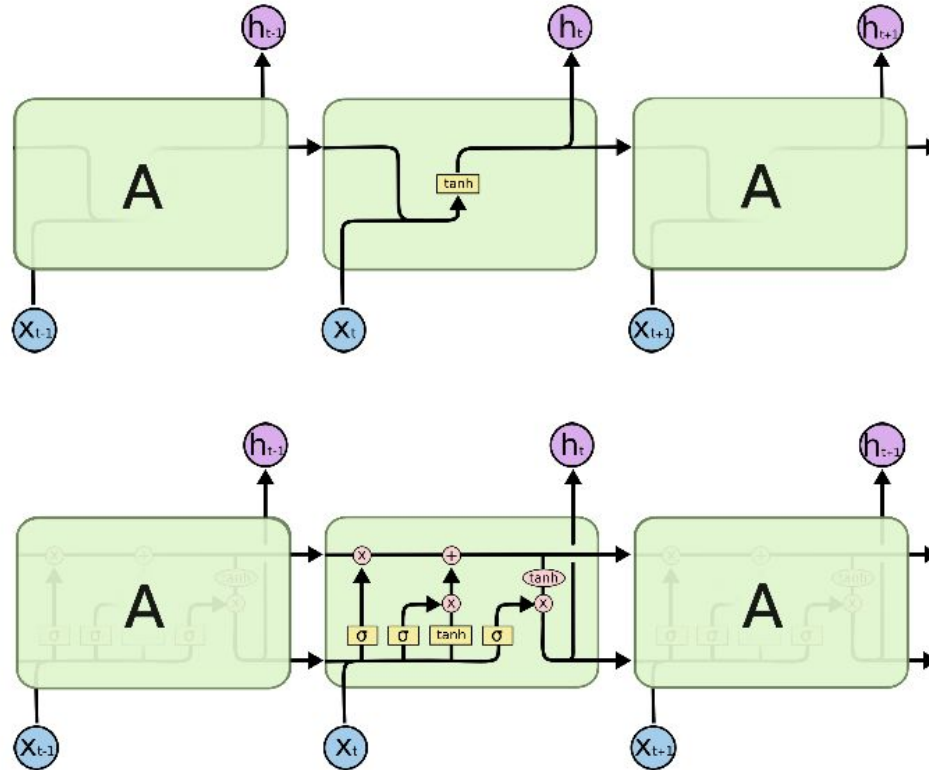
# Long term dependencies



# Long short-term memory (LSTM)

- ★ LSTM networks solve the problem of long-term dependency problem.
- ★ They use gates that allow to keep memory through long sequences and be updated only when required.
- ★ Hochreiter, Sepp, and Jürgen Schmidhuber. "Long short-term memory." Neural computation 9, no. 8 (1997): 1735-1780.

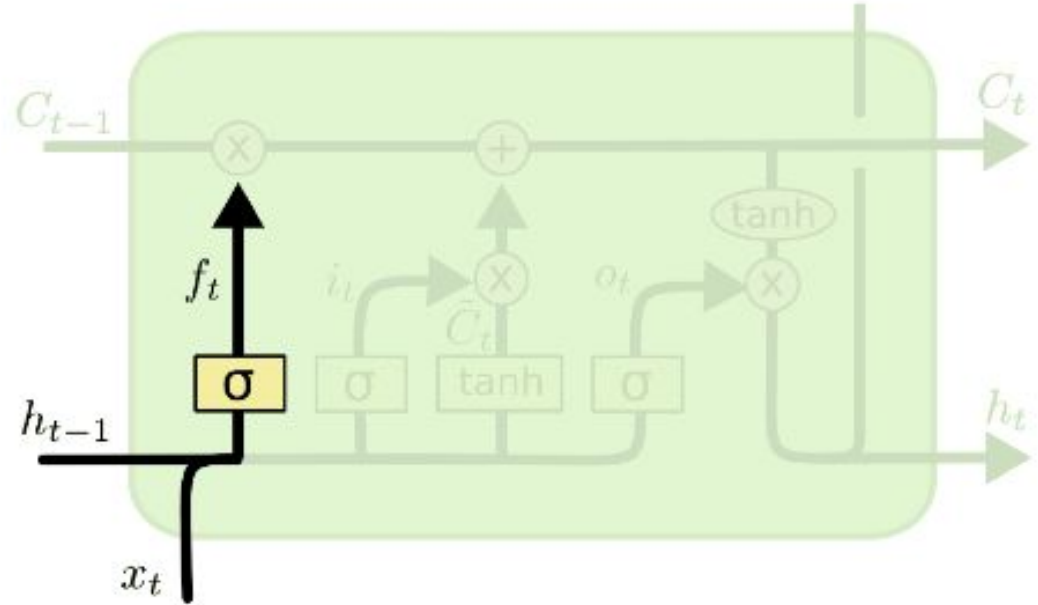
# Conventional RNN vs LSTM



(source: <http://colah.github.io/posts/2015-08-Understanding-LSTMs/>)

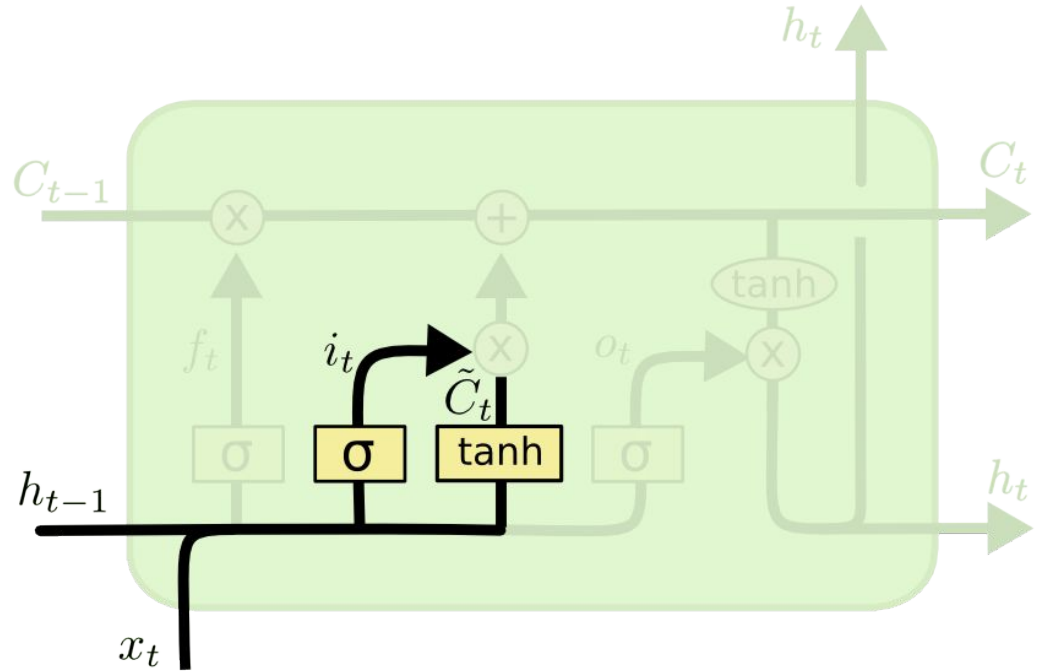
# Forget Gate

- ★ Controls the flow of the previous internal state  $C_{t-1}$
- ★  $f_t=1 \Rightarrow$  keep previous state
- ★  $f_t=0 \Rightarrow$  forget previous state

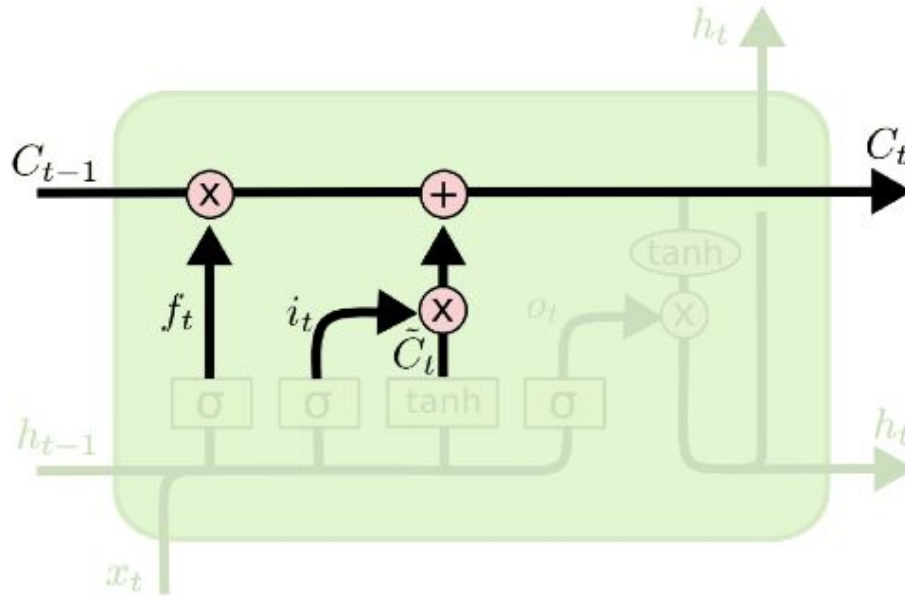


# Input Gate

- ★ Controls the flow of the input state  $x_t$
- ★  $i_t=1 \Rightarrow$  take input into account
- ★  $i_t=0 \Rightarrow$  ignore input



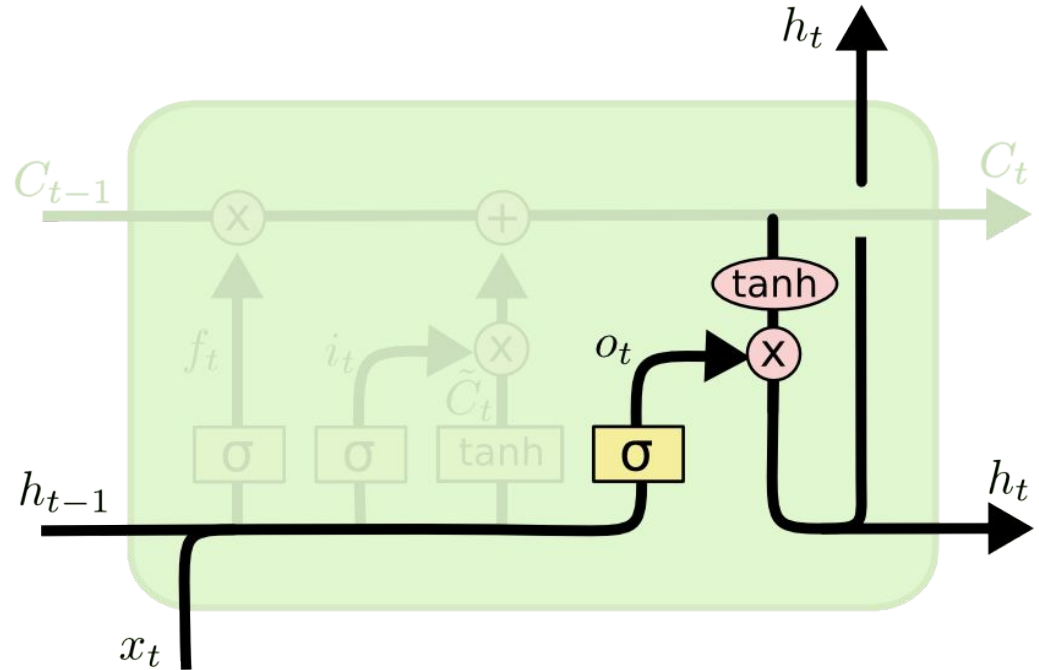
# Current state calculation



$$C_t = f_t * C_{t-1} + i_t * \tilde{C}_t$$

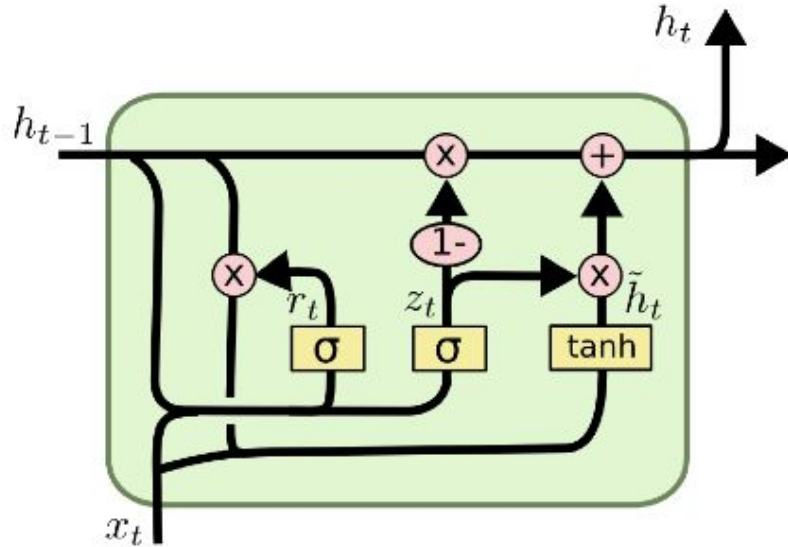
# Output Gate

- ★ Controls the flow of information from the internal state  $x_t$  to the outside  $h_t$
- ★  $o_t=1 \Rightarrow$  allows internal state out
- ★  $o_t=0 \Rightarrow$  doesn't allow internal state out





# Gated Recurrent Unit (GRU)



$$z_t = \sigma(W_z \cdot [h_{t-1}, x_t])$$

$$r_t = \sigma(W_r \cdot [h_{t-1}, x_t])$$

$$\tilde{h}_t = \tanh(W \cdot [r_t * h_{t-1}, x_t])$$

$$h_t = (1 - z_t) * h_{t-1} + z_t * \tilde{h}_t$$

Cho, K., Van Merriënboer, B., Gulcehre, C., Bahdanau, D., Bougares, F., Schwenk, H., & Bengio, Y. (2014). Learning phrase representations using rnn encoder-decoder for statistical machine translation. arXiv preprint arXiv:1406.1078.

(source: <http://colah.github.io/posts/2015-08-Understanding-LSTMs/>)

# The Unreasonable Effectiveness of Recurrent Neural Networks

- ★ Famous [blog](#) entry from Andrej Karpathy (UofS)
- ★ Character-level language models based on multi-layer LSTMs.
- ★ Data:
  - Shakspeare plays
  - Wikipedia
  - LaTeX
  - Linux source code

# Algebraic geometry book in LaTeX

*Proof.* Omitted. □

**Lemma 0.1.** *Let  $\mathcal{C}$  be a set of the construction.*

*Let  $\mathcal{C}$  be a gerber covering. Let  $\mathcal{F}$  be a quasi-coherent sheaves of  $\mathcal{O}$ -modules. We have to show that*

$$\mathcal{O}_{\mathcal{O}_X} = \mathcal{O}_X(\mathcal{L})$$

*Proof.* This is an algebraic space with the composition of sheaves  $\mathcal{F}$  on  $X_{\text{étale}}$  we have

$$\mathcal{O}_X(\mathcal{F}) = \{\text{morph}_1 \times_{\mathcal{O}_X} (\mathcal{G}, \mathcal{F})\}$$

where  $\mathcal{G}$  defines an isomorphism  $\mathcal{F} \rightarrow \mathcal{F}$  of  $\mathcal{O}$ -modules. □

**Lemma 0.2.** *This is an integer  $Z$  is injective.*

*Proof.* See Spaces, Lemma ?? □

**Lemma 0.3.** *Let  $S$  be a scheme. Let  $X$  be a scheme and  $X$  is an affine open covering. Let  $\mathcal{U} \subset X$  be a canonical and locally of finite type. Let  $X$  be a scheme. Let  $X$  be a scheme which is equal to the formal complex.*

*The following to the construction of the lemma follows.*

*Let  $X$  be a scheme. Let  $X$  be a scheme covering. Let*

$$b : X \rightarrow Y' \rightarrow Y \rightarrow Y \rightarrow Y' \times_X Y \rightarrow X.$$

*be a morphism of algebraic spaces over  $S$  and  $Y$ .*

*Proof.* Let  $X$  be a nonzero scheme of  $X$ . Let  $X$  be an algebraic space. Let  $\mathcal{F}$  be a quasi-coherent sheaf of  $\mathcal{O}_X$ -modules. The following are equivalent

- (1)  $\mathcal{F}$  is an algebraic space over  $S$ .
- (2) If  $X$  is an affine open covering.

Consider a common structure on  $X$  and  $X$  the functor  $\mathcal{O}_X(U)$  which is locally of finite type. □

This since  $\mathcal{F} \in \mathcal{F}$  and  $x \in \mathcal{G}$  the diagram

$$\begin{array}{ccc} S & \xrightarrow{\quad} & \\ \downarrow & & \\ \xi & \xrightarrow{\quad} & \mathcal{O}_{X'} \\ \text{gor}_x & \uparrow & \searrow \\ & \alpha' & \\ & \downarrow & \\ & \alpha' & \xrightarrow{\quad} \alpha \end{array} \quad \begin{array}{c} X \\ \downarrow \\ \text{Mor}_{\text{Sets}} \\ \text{d}(\mathcal{O}_{X_{X/\mathcal{F}}}, \mathcal{G}) \end{array}$$

is a limit. Then  $\mathcal{G}$  is a finite type and assume  $S$  is a flat and  $\mathcal{F}$  and  $\mathcal{G}$  is a finite type  $\mathcal{F}$ . This is of finite type diagrams, and

- the composition of  $\mathcal{G}$  is a regular sequence,
  - $\mathcal{O}_{X'}$  is a sheaf of rings.
- 

*Proof.* We have see that  $X = \text{Spec}(R)$  and  $\mathcal{F}$  is a finite type representable by algebraic space. The property  $\mathcal{F}$  is a finite morphism of algebraic stacks. Then the cohomology of  $X$  is an open neighbourhood of  $U$ . □

*Proof.* This is clear that  $\mathcal{G}$  is a finite presentation, see Lemmas ??.

A reduced above we conclude that  $U$  is an open covering of  $\mathcal{C}$ . The functor  $\mathcal{F}$  is a “field

$$\mathcal{O}_{X,x} \longrightarrow \mathcal{F}_x \rightarrow \mathcal{O}_{X_{\text{étale}}}^{-1} \longrightarrow \mathcal{O}_{X_{\lambda}}^{-1} \longrightarrow \mathcal{O}_{X_{\lambda}}^{\vee}$$

is an isomorphism of covering of  $\mathcal{O}_{X_{\lambda}}$ . If  $\mathcal{F}$  is the unique element of  $\mathcal{F}$  such that  $X$  is an isomorphism.

The property  $\mathcal{F}$  is a disjoint union of Proposition ?? and we can filtered set of presentations of a scheme  $\mathcal{O}_X$ -algebra with  $\mathcal{F}$  are opens of finite type over  $S$ .

If  $\mathcal{F}$  is a scheme theoretic image points. □

If  $\mathcal{F}$  is a finite direct sum  $\mathcal{O}_{X_{\lambda}}$  is a closed immersion, see Lemma ?? . This is a sequence of  $\mathcal{F}$  is a similar morphism.

# Linux source code

```
/*
 * Increment the size file of the new incorrect UI_FILTER group information
 * of the size generatively.
 */
static int indicate_policy(void)
{
    int error;
    if (fd == MARN_EPT) {
        /*
         * The kernel blank will coeld it to userspace.
         */
        if (ss->segment < mem_total)
            unblock_graph_and_set_blocked();
        else
            ret = 1;
        goto bail;
    }
    segaddr = in_SB(in.addr);
    selector = seg / 16;
    setup_works = true;
    for (i = 0; i < blocks; i++) {
        seq = buf[i++];
        bpf = bd->bd.next + i * search;
        if (fd) {
            current = blocked;
```

# LSTM language model demo