







EGYPFWD Initiative

Advanced Embedded Systems Nanodegree, Real-Time Operating Systems Masterclass by SPRINTS Egypt.

Implementation of an EDF Scheduler based on freeRTOS

A Graduation Project submitted in partial Fulfillment of Real-Time Operating Systems Masterclass.

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Verifying the system implementation

Analytical Method

| ID | Task | Periodicity (ticks) | Execution Time(ms) | Calls in HyperPeriod |
|----|----------------------|---------------------|--------------------|-------------------------|
| 1 | Button_1_Monitor | 50 | 0.019 | 2 |
| 2 | Button_2_Monitor | 50 | 0.0193 | 2 |
| 3 | Periodic_Transmitter | 100 | 0.02465 | 1 |
| 4 | Uart_Receiver | 20 | 0.02352 | 5 |
| 5 | Load_1_Simulation | 10 | 5 | 10 |
| 6 | Load_2_Simulation | 100 | 12 | 1 |

• Calculation of HyperPeriod

$$HyperPeriod = LCM(Periodicities) = LCM(50,50,100,20,10,100)$$

 $HyperPeriod = 100$

• **CPU Load Calculations**

$$\mathbf{CPU \, LOAD} = \frac{\text{Total Time}}{\text{HyperPeriod}} *100$$

$$\textbf{\textit{Total Time}} = \sum_{i=1}^{6} \text{ExecutionTime}_{i} \text{*Num of Calls In HyperPeriod}_{i}$$

$$Total\ Time = 0.019 * 2 + 0.0193 * 2 + 0.02465 + 0.02352 * 5 + 5 * 10 + 12$$

$$U = CPU LOAD = \frac{62.21885}{100} *100 = 62.21885\%$$





• Schedulability Analysis using Rate Monotonic Utilization Bound

(Assuming the given set of tasks are scheduled using a fixed priority rate -monotonic scheduler)

if $U \leq Urm$ then the system is schedulable

$$Urm = n[2^{1/n} - 1]$$

$$Urm = 6*(2^{1/6}-1) = 0.7347723$$

The system is guaranteed schedulable since U < Urm.

• Schedulability Analysis using Time Demand Analysis

$$w_i(t) = e_i + \sum_{k=1}^{i-1} \left\lceil \frac{t}{p_k} \right\rceil e_k \quad \text{ for } 0 < t \leq p_i \quad \begin{array}{l} \text{W = Worst response time} \\ \text{E = Execution time} \\ \text{P = Periodicity} \\ \text{T = Time instance} \end{array}$$

Tasks are organized by lowest periodicity which is highest priority.

$$Task 5 \rightarrow Task 4 \rightarrow Task 1 \rightarrow Task 2 \rightarrow Task 3 \rightarrow Task 6$$

o Task 5

$$W(4) = 5 + 0 = 5ms$$

Since D = 10 ms and W(10) = 5ms
 $W(10) < D$

Task 5 is schedulable.

o Task 4

$$W(20) = 0.02352 + (20/10)*5 = 10.02352ms$$

Since D = 20 ms and W(20) = 10.02352ms
 $W(20) < D$

Task 4 is schedulable.

o Task 1

$$W(50) = 0.019 + (50/10)*5 + (50/20)*0.02352 = 25.607ms$$

Since D = 50 ms and W(50) = 25.607ms
 $W(50) < D$

Task 1 is schedulable.





o Task 2

$$W(50) = 0.0193 + (50/10)*5 + (50/20)*0.02352 + (50/50)*0.019$$

= 25.6263ms
Since D = 50 ms and W(50) = 25.6263ms
 $W(50) < D$

Task 2 is schedulable.

o Task 3

$$W(100) = 0.02465 + (100/50)*0.019 + (100/50)*0.0193 + (100/10)*5$$

 $+ (100/20)*0.02352 = 50.21885ms$
Since D = 100 ms and W(100) = 50.21885ms
 $W(100) < D$

Task 3 is schedulable.

o Task 6

$$W(100) = 12 + (100/100)*0.02465 + (100/50)*0.019 + (100/50)*0.0193 + (100/10)*5$$

 $+ (100/20)*0.02352 = 62.21885ms$
 Since D = 100 ms and W(100) = 62.21885ms
 $W(100) < D$

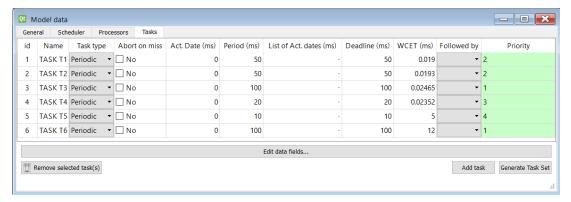
Task 6 is schedulable.

Therefore, the system is schedulable.

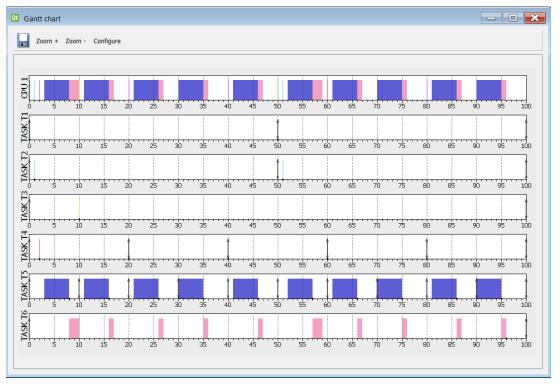




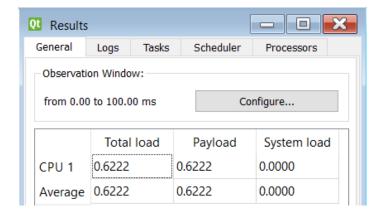
Simso Simulation Method



Task Creation



Gantt Chart of task activity

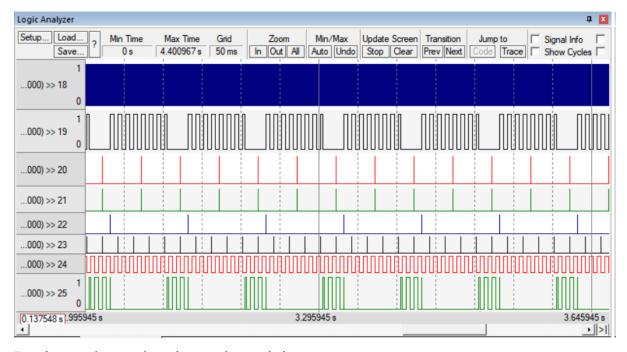


CPU Load which is the same as calculated.

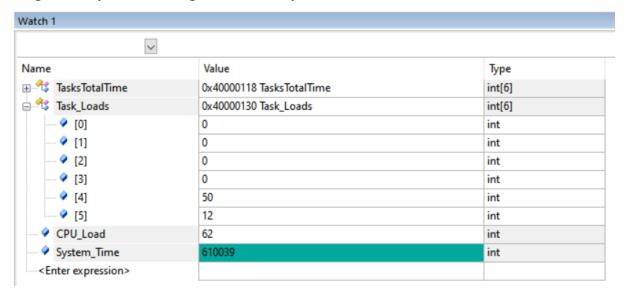




Keil Simulation Method



Logic Analyzer showing tasks activity



CPU Load which is the same as calculated as well as Simso simulation.

In conclusion, the CPU loads measured in the three methods are almost identical which proves a successful implementation of the EDF scheduler in the freeRTOS kernel.