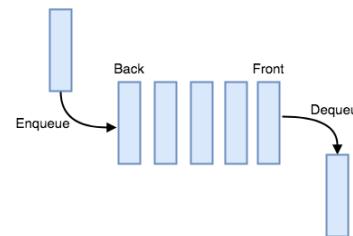


Data Structures and Algorithms

Queue



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Outlines

- ▶ **Introduction to Queue**
- ▶ **Queue ADT and its operations**
- ▶ **Queue Applications**
- ▶ **Implementation of Queue**
- ▶ **Array implementation of Queue**
- ▶ **Pseudocode for Array implementation of Queue**

What is Queue?



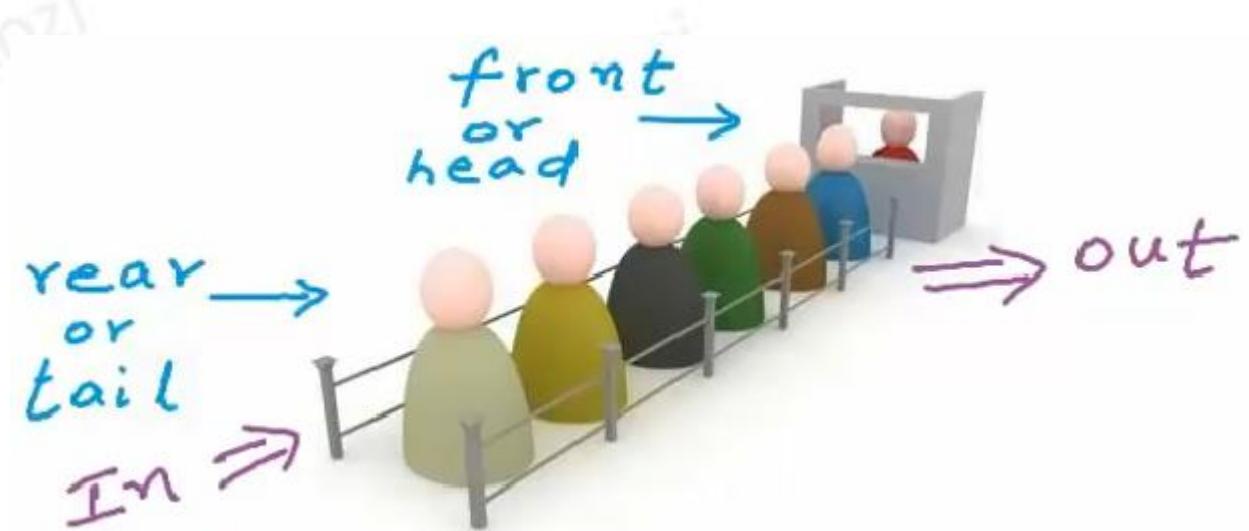
Introduction to Queue 1/2

- ▶ **Queue** data structure is exactly what we mean when we say **Queue** in real world.
- ▶ **Queue** is a data structure in which whatever goes in first, comes out first.
- ▶ First In First Out (FIFO).
 - ▶ Stack is LIFO.



Introduction to Queue 2/2

- ▶ Insertion must happen from one end called rear or tail.
- ▶ Removal must happen from the other end, which called front or head.



Queue Definition

- Formal definition of Queue ADT:

A list or collection with the restriction that insertion can be performed at one end (rear) and deletion can be performed at other end (front).

Queue Operations

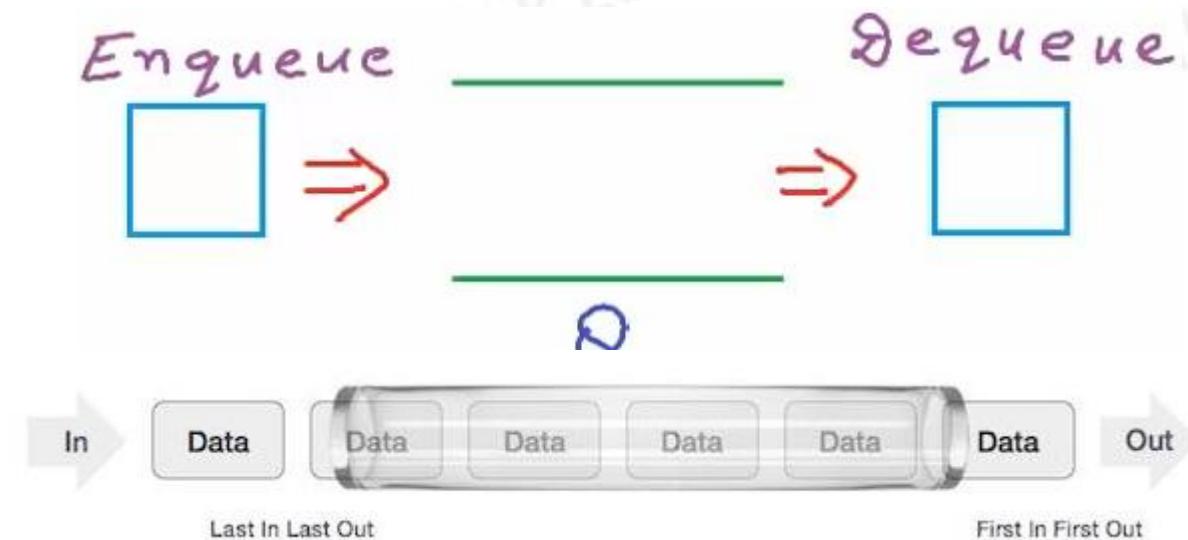
Operations available with Queue:

- The main two operations are:
 - Insertion an element, called **EnQueue**.
 - Removing an element, called **DeQueue**.
- Returns the element at the top of Queue, **Front**.
 - Should not remove any element from the Queue.
- Check if Queue is empty or not, **IsEmpty**.
- **EnQueue** and **DeQueue** one element at a time.
- All operations can be performed in constant time -> $O(1)$.

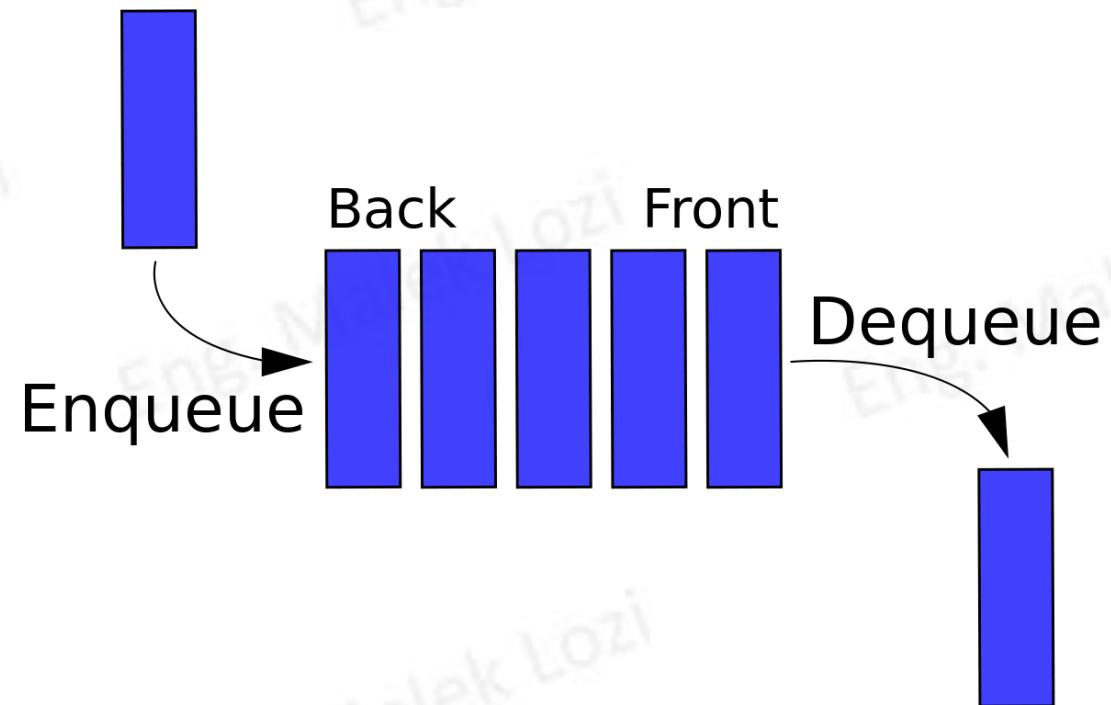
- (1) *EnQueue (x)*
- (2) *DeQueue ()*
- (3) *front ()*
- (4) *IsEmpty ()*

Logical representation of Queue

- Logical representation of a Queue:
 - Container opened from Two sides.
 - An element can be inserted from one side.
 - An element can be removed from the other side.

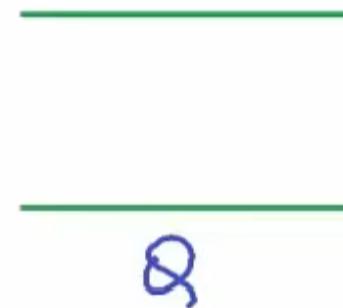


Insertion & Removing an element (EnQueue) & (DeQueue)

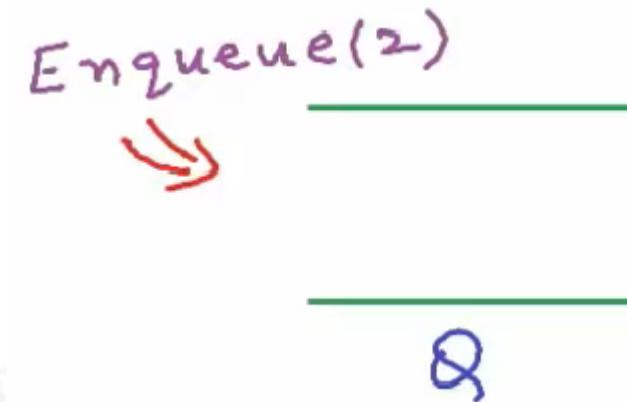


(EnQueue & DeQueue) (1/3)

- We have an empty Queue called Q:



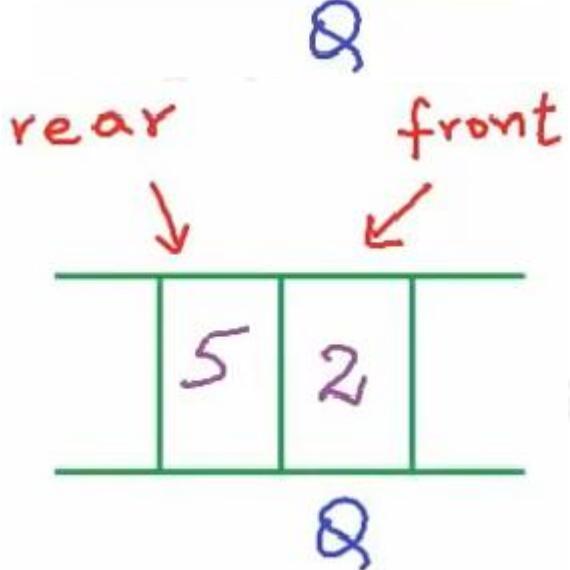
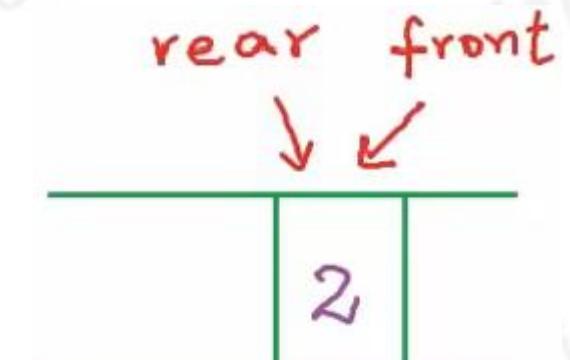
- Inserting number 2.



(EnQueue & DeQueue) (2/3)

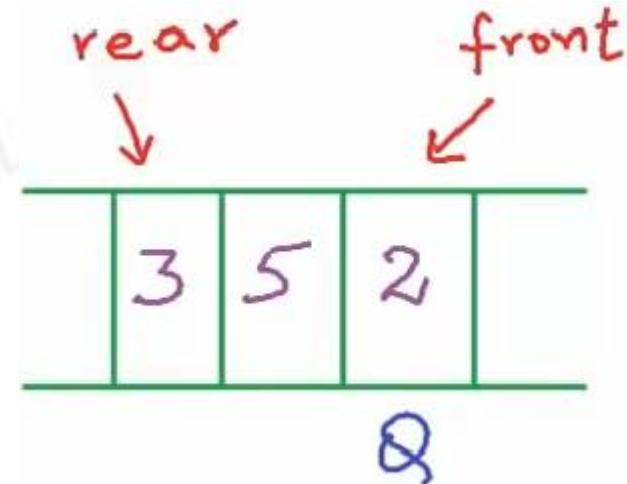
- After the EnQueue, the Queue will look like:
 - It contains one integer.

- Inserting number 5.
 - It will be at rear of the Queue.

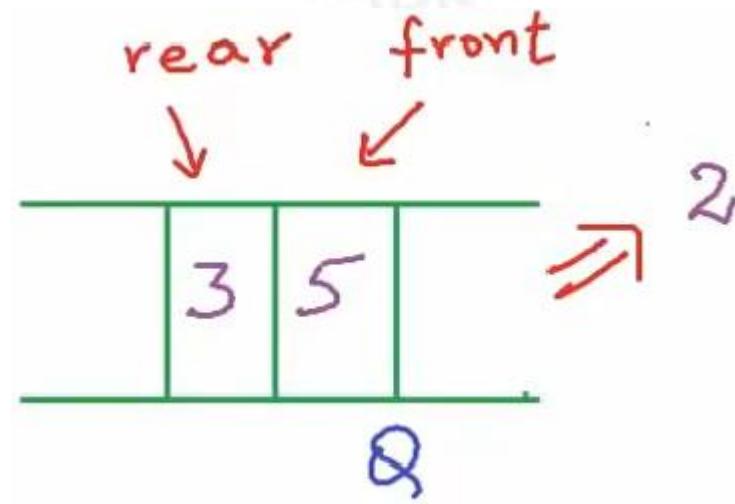


(EnQueue & DeQueue) (3/3)

- ▶ Inserting number 3:



- ▶ Removing an element.
- ▶ We will get number 2.



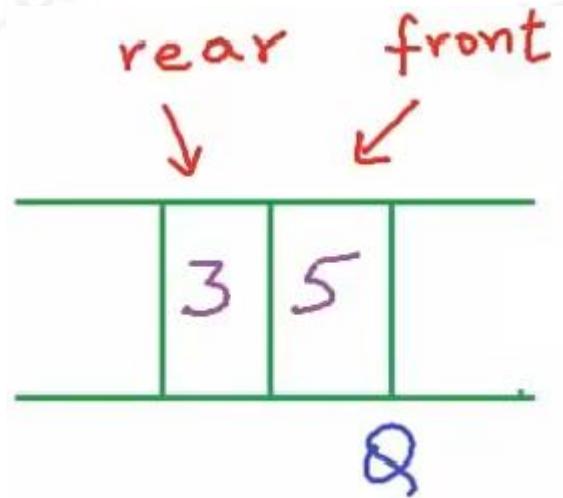
Calling Front () & IsEmpty() in Queue

► Calling Front at this stage:

► Will return number 5.

► Calling IsEmpty at this stage:

► Will return False.



Queue Applications

- Real applications of Queue.
 - Shared resource that can handle only one request at a time.
 - Queuing the requests.
 - The request that comes first, get served first.
 - Process scheduling.
 - Processor on computer is a shared resource.



Implementation of Queue

- Referring to the formal definition, A Queue is a special kind of **List**, in which an element can be inserted at One end and an element can be removed from the other end.

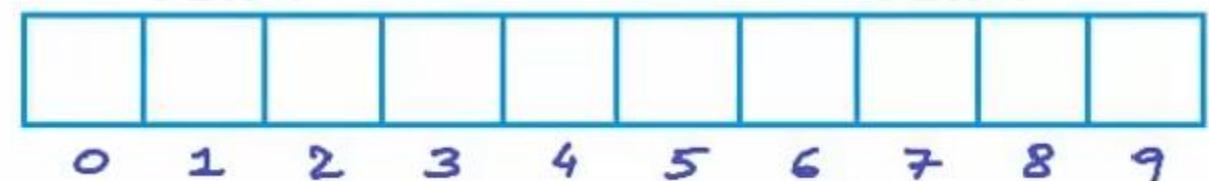
We can implement Queues using:

- 1) Arrays
- 2) Linked Lists

Array implementation of Queue (1/9)

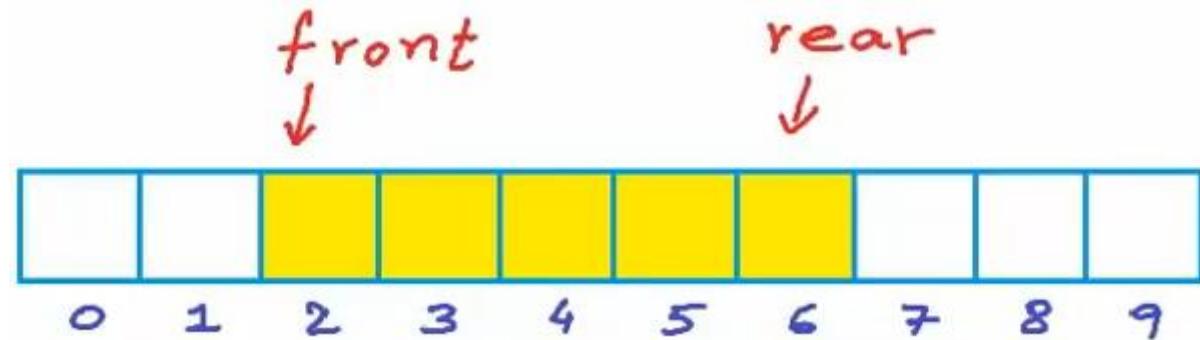
- Let us create a Queue of Integers.
- First, we create an array of integers.
 - Array of Ten integers.

`int A[10]`



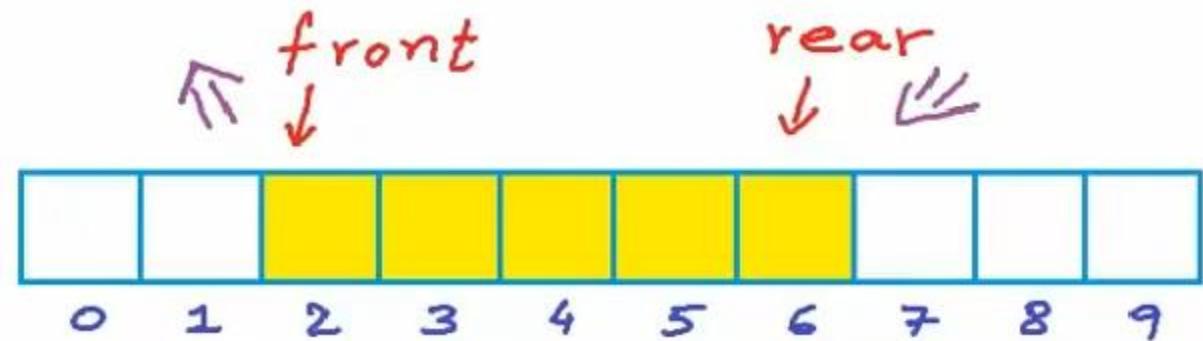
Array implementation of Queue (2/9)

- At any point, some part after array starting at index marked as front to index marked as rear will be our Queue.
- Rest of the Array position are free space that can be used to expand the Queue.



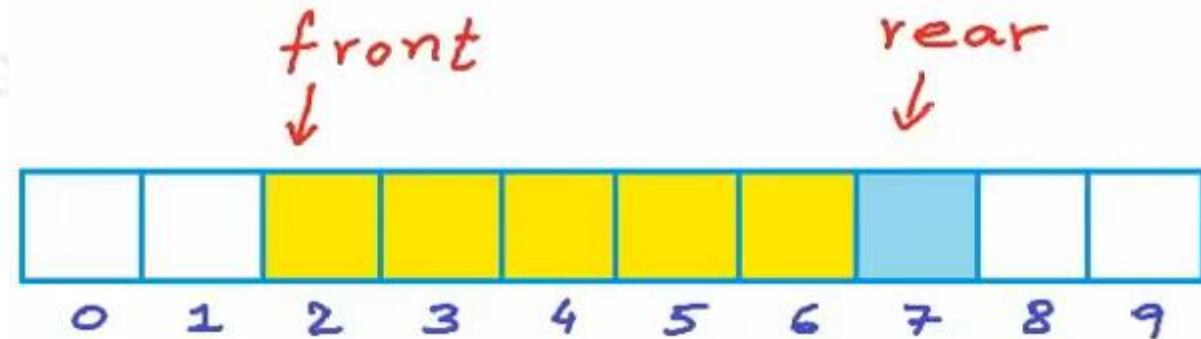
Array implementation of Queue (3/9)

- An element must always be added from rear side and must always be removed from front side.



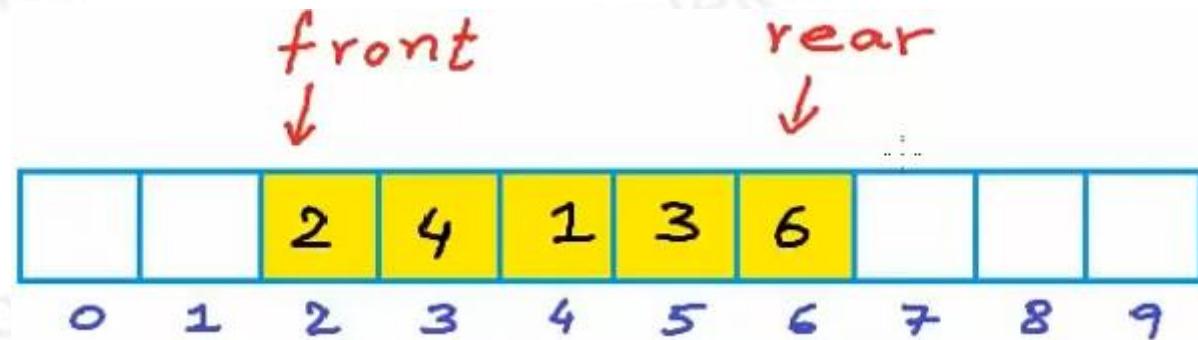
Array implementation of Queue (4/9)

- To insert an element (EnQueue), we can increment rear.
- In this position we can add the new element.

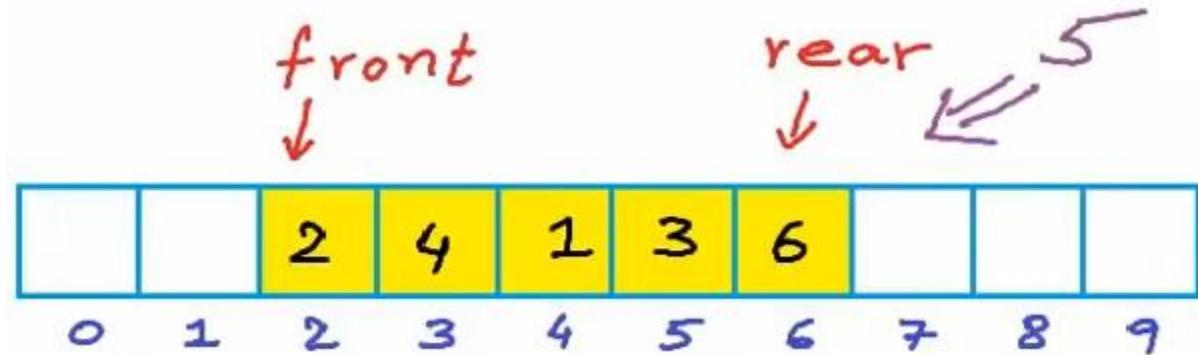


Array implementation of Queue (5/9)

- Let us fill some values in the Queue.

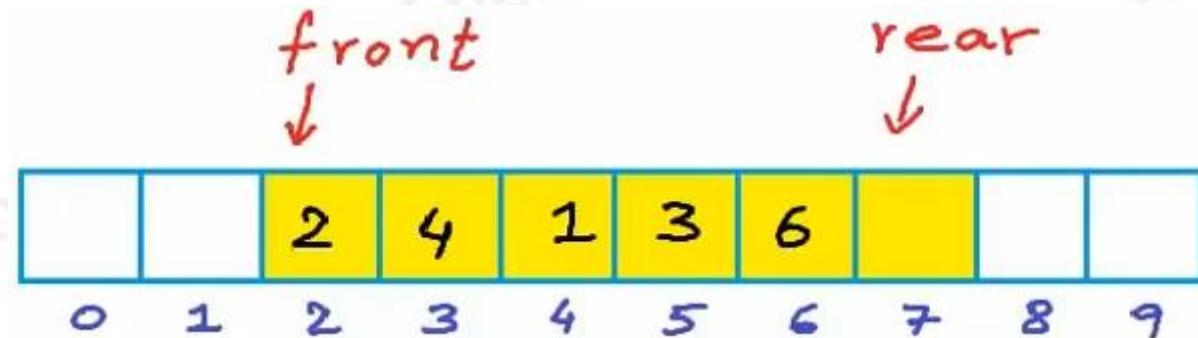


- We want to insert number 5.

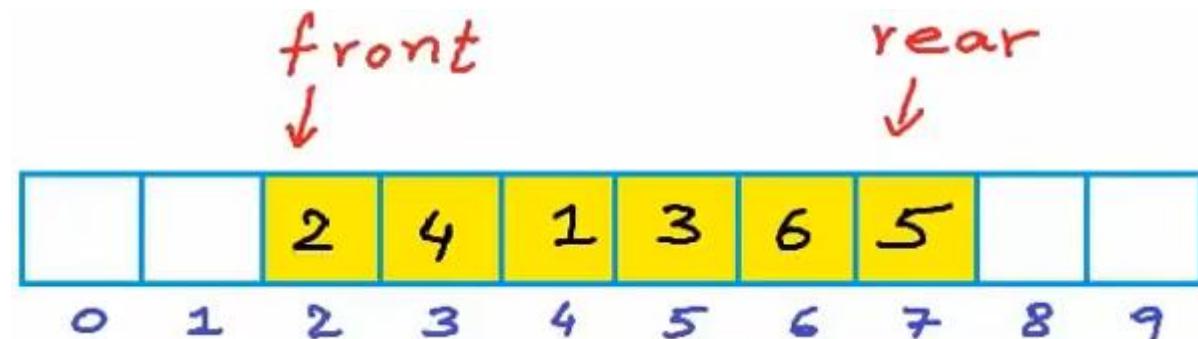


Array implementation of Queue (6/9)

- We will increment rear.
- There should be an available empty cell in the right.

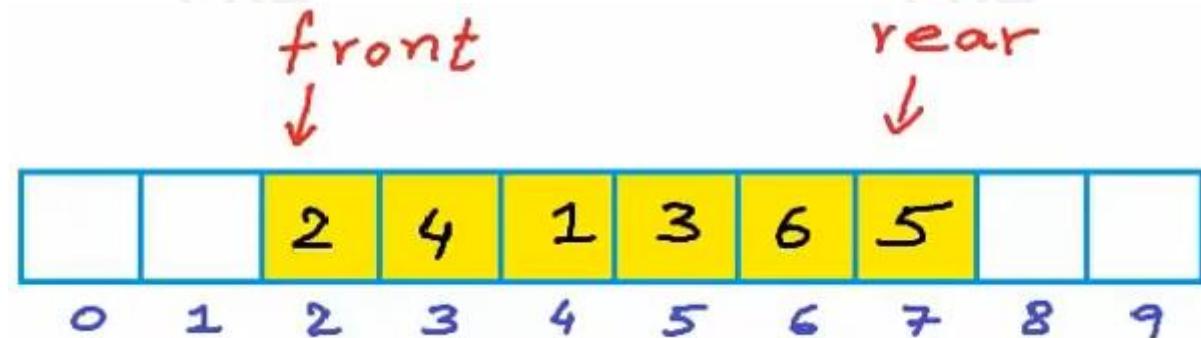


- Now we can insert number 5.



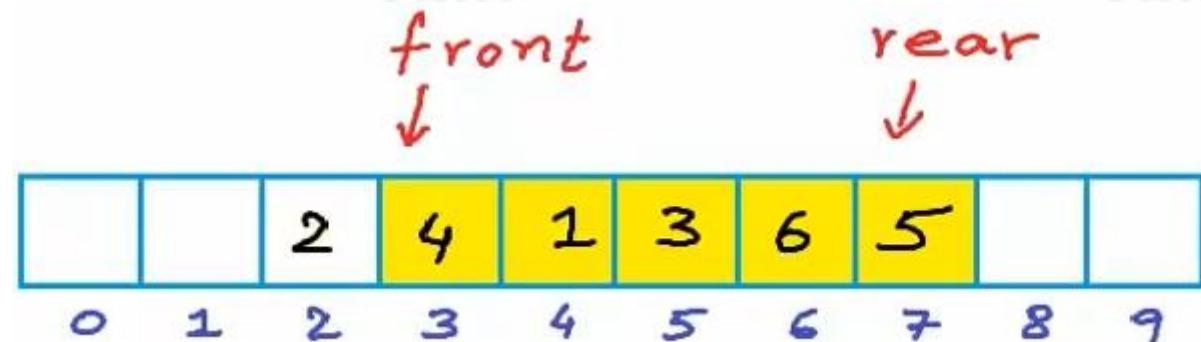
Array implementation of Queue (7/9)

- After this insertion, the rear is at index 7.
- Now DeQueue means we must remove an element from front of the Queue.
- Here, DeQueue operation should remove number 2 from the Queue.



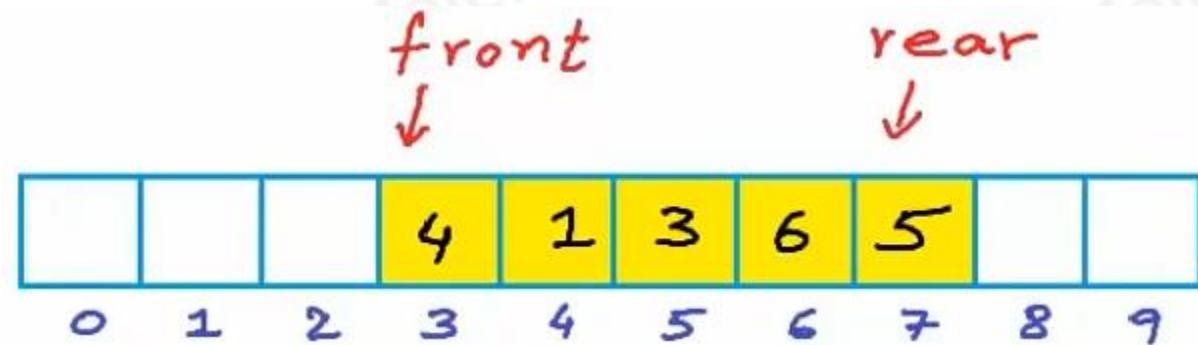
Array implementation of Queue (8/9)

- ▶ To DeQueue, we can simply increment front.
- ▶ Because at any point only the cells starting from **front** to **rear** are part of the Queue.
- ▶ By incrementing **front**, we have discarded index 2 from the Queue.



Array implementation of Queue (9/9)

- We don't care what value in a cell that is not part of the Queue.
 - When we include a cell in the Queue, we will overwrite the value in that cell.
- So, just incrementing **front** is good enough for DeQueue operation.



Let us write the

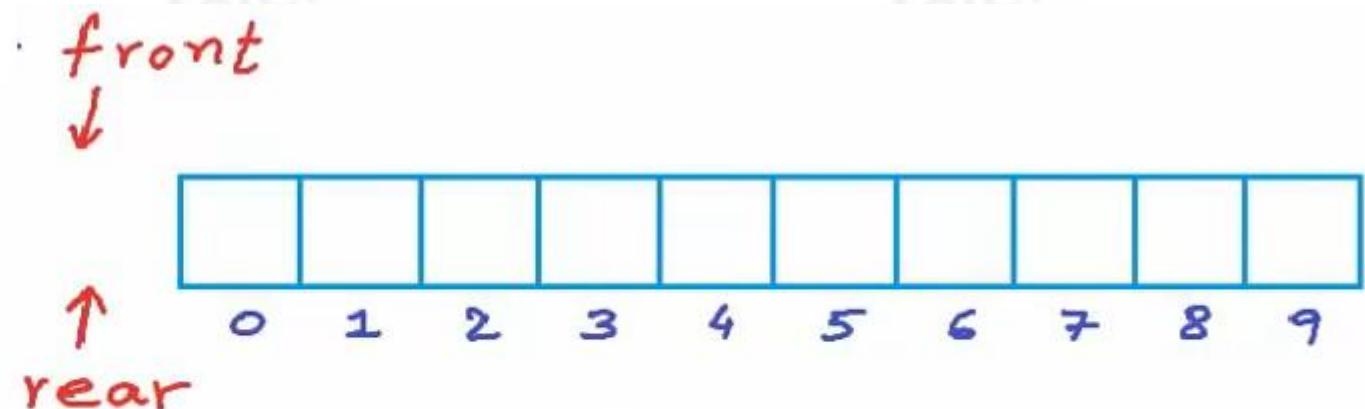


of Queue using Array
Implementation

PseudoCode for Queue using Array (1/13)

- Let us write the PseudoCode now:
- We will declare two variable named front and rear, initially we will set them both to -1.
- So, for an empty Queue both front and rear will be -1.

```
int A[10]  
front ← -1  
rear ← -1
```



PseudoCode for Queue using Array (2/13)

- To check whether a Queue is empty or not, we can simply check the value of front and rear.
- If they are both -1, then the Queue is empty.

```
IsEmpty()
{
    if front == -1 && rear == -1
        return true
    else
        return false
}
```

PseudoCode for Queue using Array (3/13)

- Now, EnQueue function.
- EnQueue will take integer x as an argument.

```
Enqueue(x)
{
    if rear == size(A)-1
        Print "Queue is Full"
        return
    else if IsEmpty()
    {
        front <- rear <- 0
    }
    else
    {
        rear <- rear + 1
    }
    A[rear] <- x
}
```

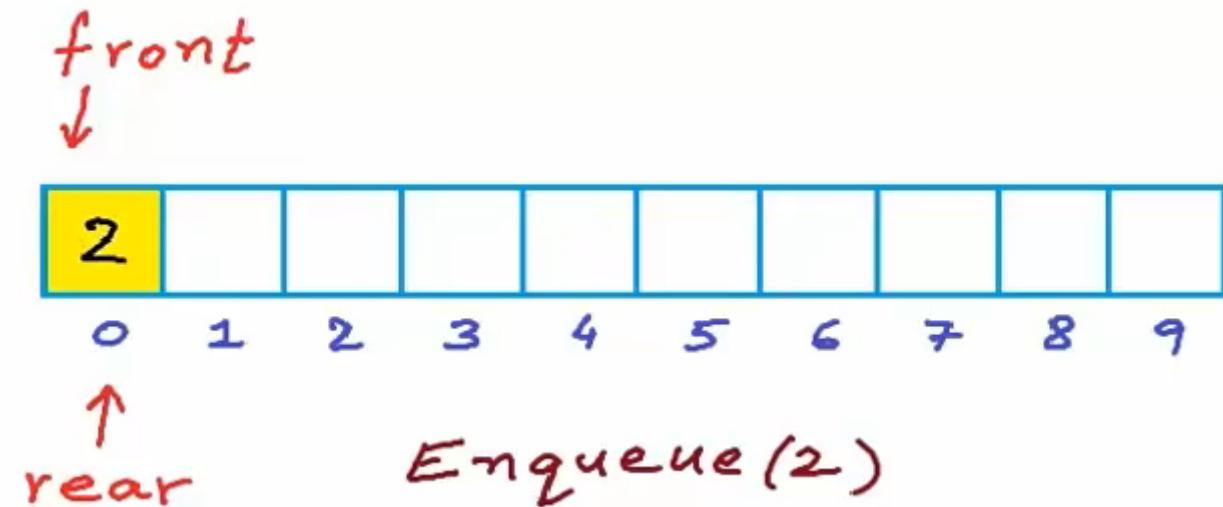
PseudoCode for Queue using Array (4/13)

- We can replace the first condition with a function named IsFull to determine the Queue is full or not.

```
Enqueue(x)
{
    if IsFull()
        return
    else if IsEmpty()
    {
        front ← rear ← 0
    }
    else
    {
        rear ← rear + 1
    }
    A[rear] ← x
}
```

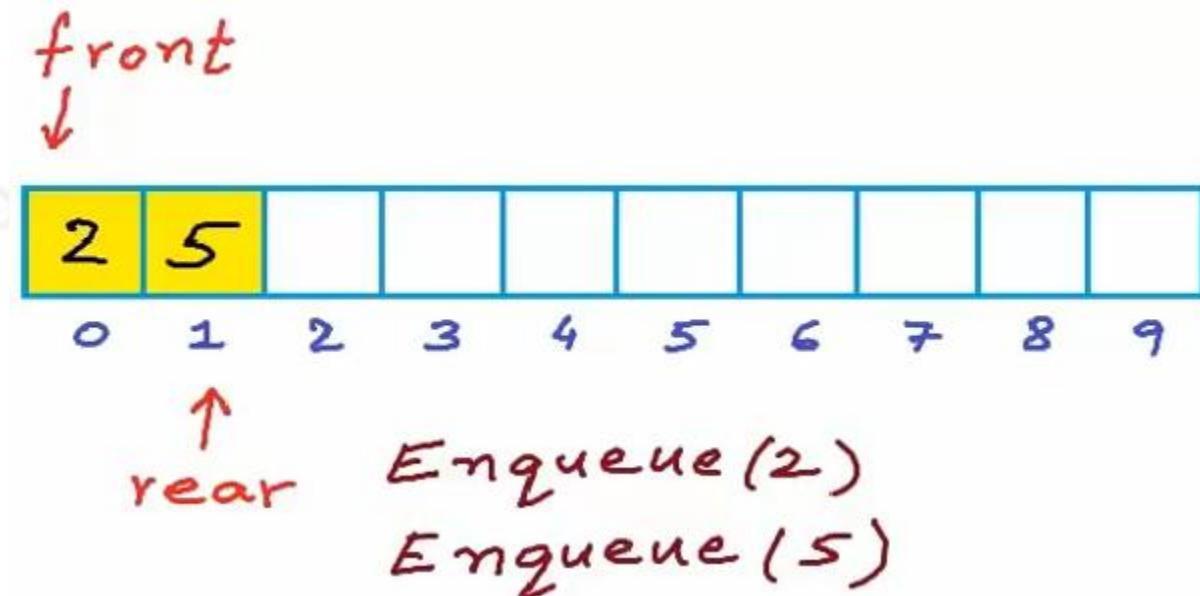
PseudoCode for Queue using Array (5/13)

- Let us EnQueue some items.
- This is the Queue after One EnQueue operation.



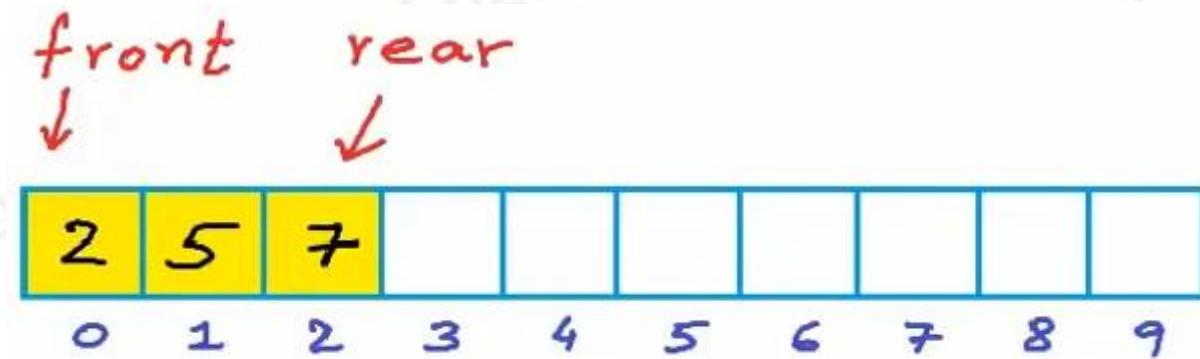
PseudoCode for Queue using Array (6/13)

- Let us make another call to EnQueue.



PseudoCode for Queue using Array (7/13)

- Let us EnQueue one more item.



*Enqueue(2)
Enqueue(5)
Enqueue(7)*

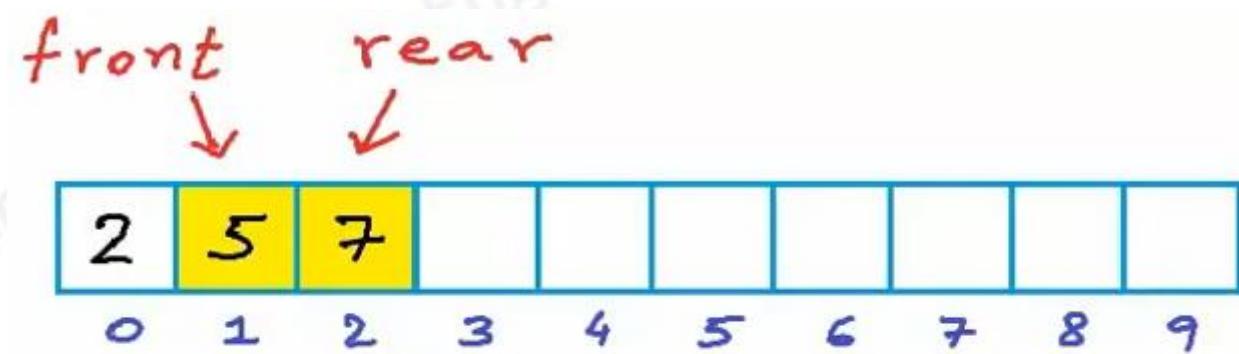
PseudoCode for Queue using Array (8/13)

- Let us now write DeQueue operation.

```
Dequeue()
{
    if IsEmpty()
        return
    else if front == rear
        front <= rear <= -1
    else
        front <= front + 1
}
```

PseudoCode for Queue using Array (9/13)

- Now, we will continue the previous example and try to call DeQueue operation.



Enqueue(2)

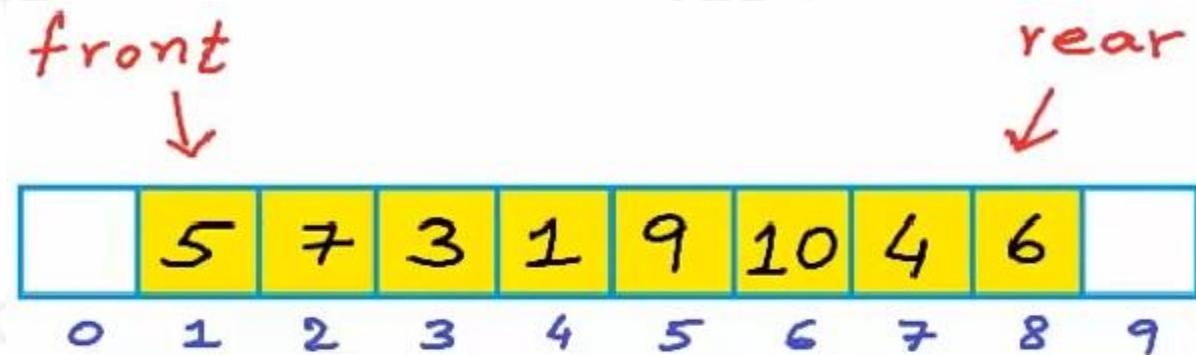
Enqueue(5)

Enqueue(7)

Dequeue()

PseudoCode for Queue using Array (10/13)

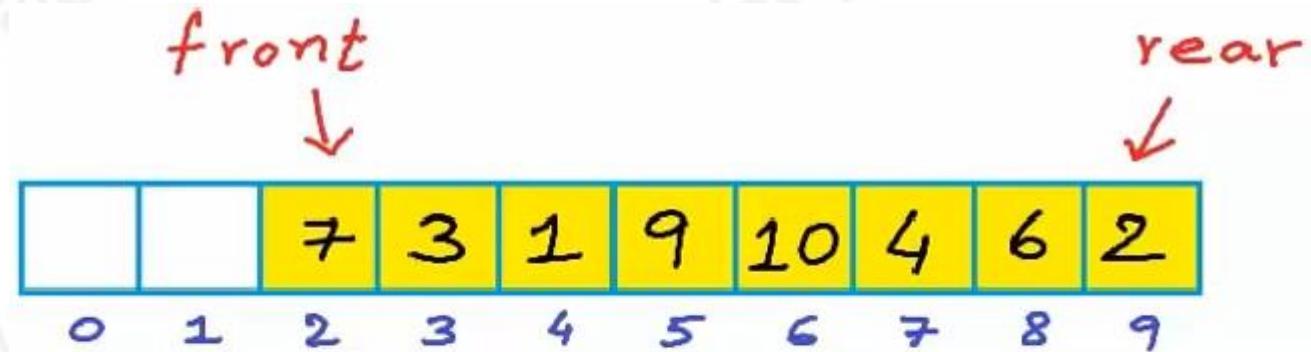
- Let's perform more EnQueue operations.



Enqueue(2)	Enqueue(9)
Enqueue(5)	Enqueue(10)
Enqueue(7)	Enqueue(4)
Dequeue()	Enqueue(6)
Enqueue(3)	
Enqueue(1)	

PseudoCode for Queue using Array (11/13)

- Let's perform DeQueue and EnQueue operations.



Enqueue(2)

Enqueue(5)

Enqueue(7)

Dequeue()

Enqueue(3)

Enqueue(1)

Enqueue(9)

Enqueue(10)

Enqueue(4)

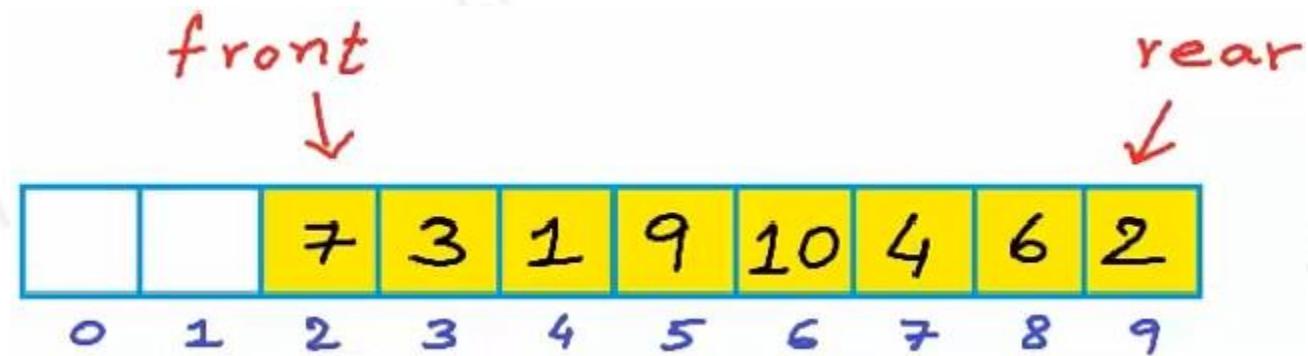
Enqueue(6)

Dequeue()

Enqueue(2)

PseudoCode for Queue using Array (12/13)

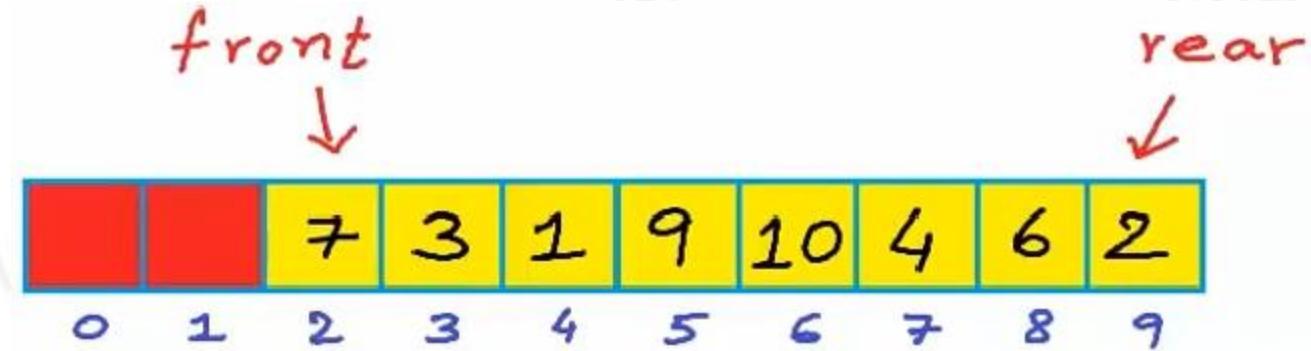
- At this stage we cannot EnQueue an item anymore, because we cannot increment rear.



- EnQueue operation will fail now.
- There are two unused cells.

PseudoCode for Queue using Array (13/13)

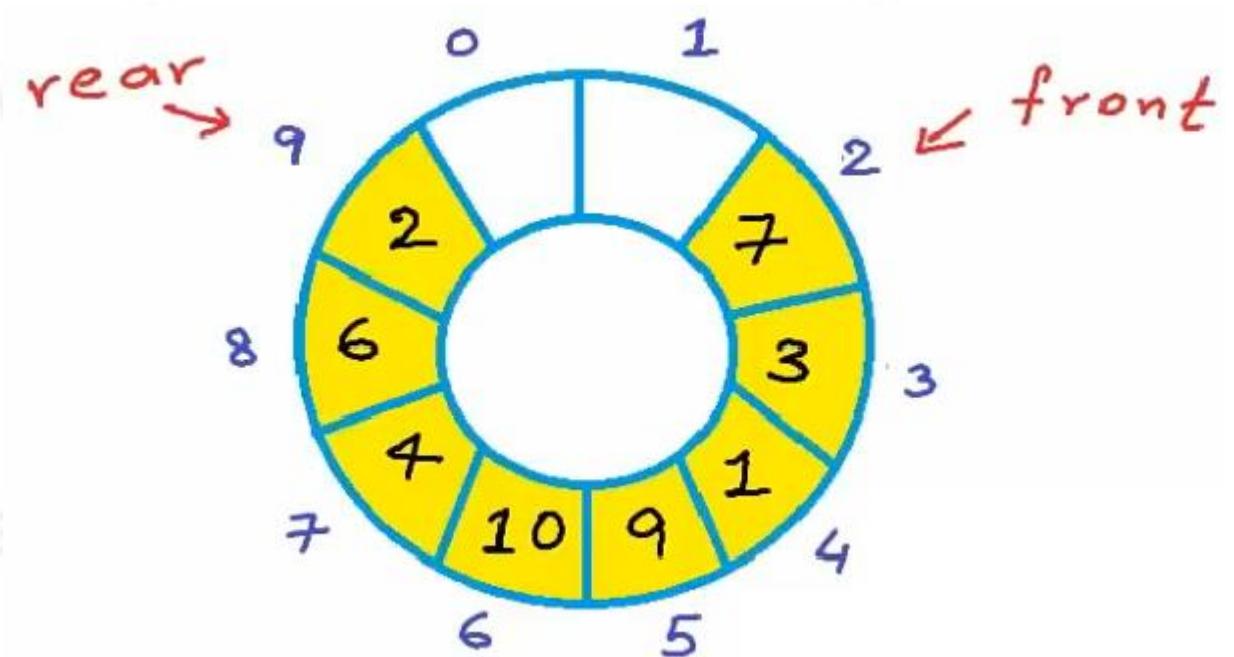
- In fact, this is a real problem.



- If we DeQueue more and more, all the cells left of front index will never be used again.
- Can we do something to use these cells?

PseudoCode for Queue using Circular Array (1/9)

- We can use the concept of Circular Array.
- Circular array is an idea that we use in a lot of scenarios.
- The idea is very simple, as we traverse an array, we can imagine that there is no end in the array.



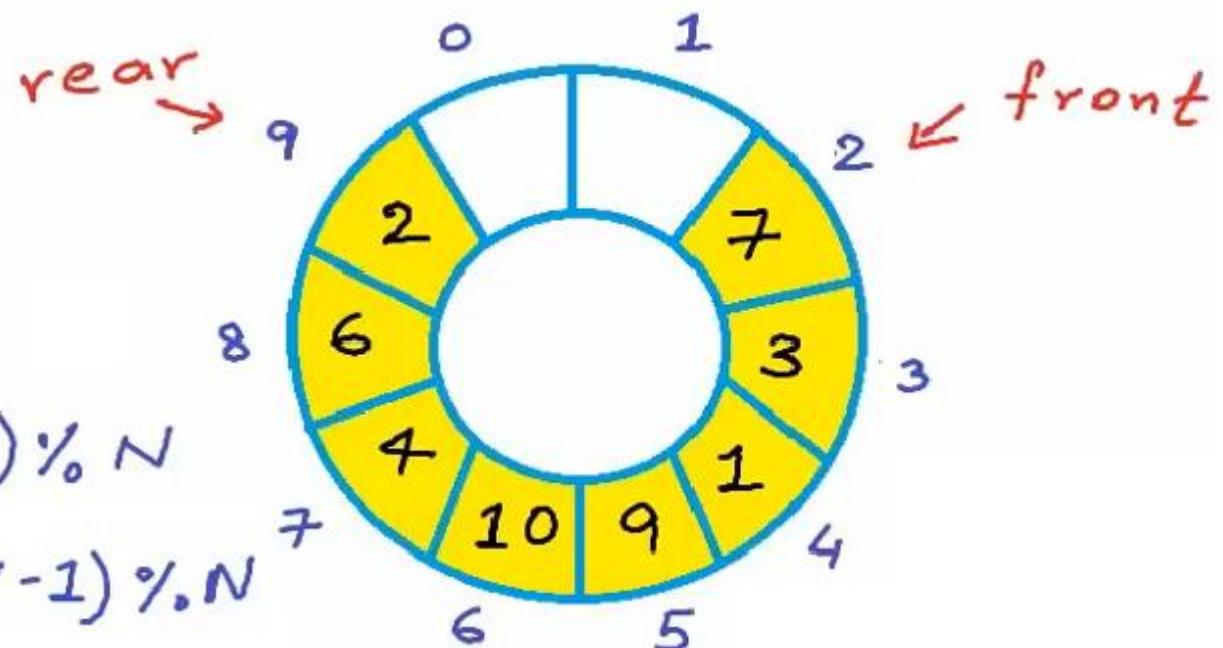
PseudoCode for Queue using Circular Array (2/9)

- In circular array, if we are pointing at any position i , the next position and the previous position will be the following:
 - N is the number of elements in Array.

Current position = i

Next position = $(i+1) \% N$

Previous position = $(i+N-1) \% N$



PseudoCode for Queue using Circular Array (3/9)

- Now, in an EnQueue operation we can increment rear as long as there is any unused cell in the array.
- We are going to modify the PseudoCodes.

```
Enqueue(x)
{
    if ((rear+1)%N == front)
        return
    else if ISEmpty()
    {
        front <= rear <= 0
    }
    else
    {
        rear <- (rear+1)% N
    }
    A [rear] <- x
}
```

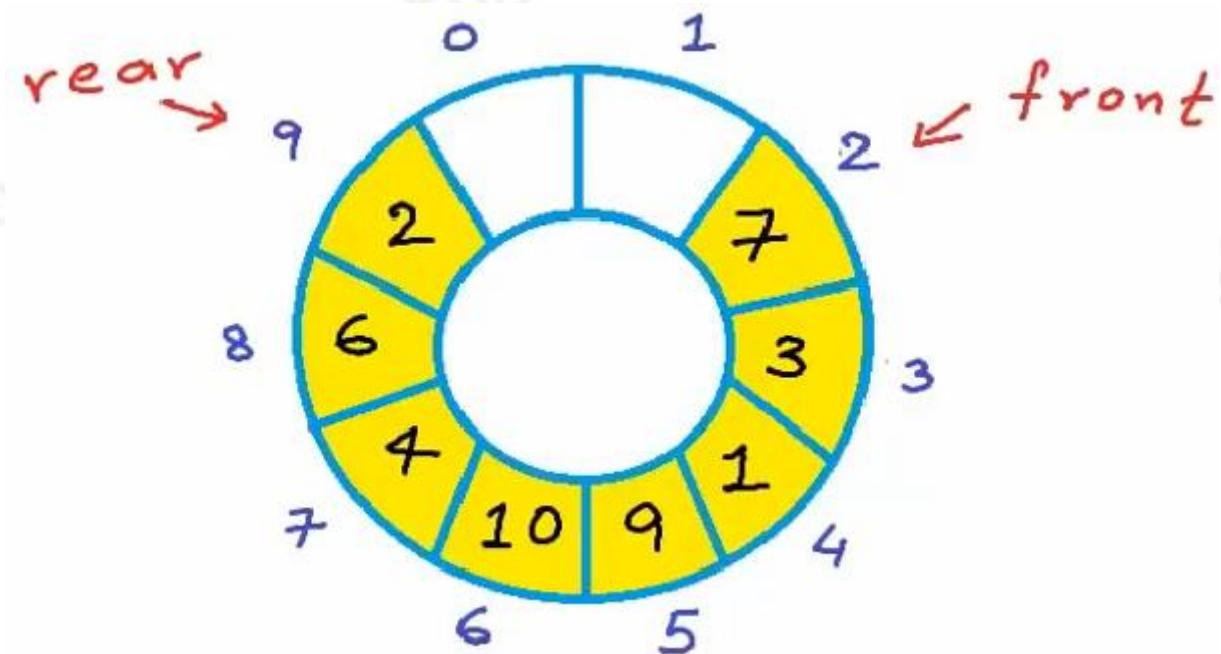
PseudoCode for Queue using Circular Array (4/9)

- The DeQueue function will be:

```
Dequeue()
{
    if IsEmpty()
        return
    else if front == rear
        front ← rear ← -1
    else
        front ← (front + 1) % N
}
```

PseudoCode for Queue using Circular Array (5/9)

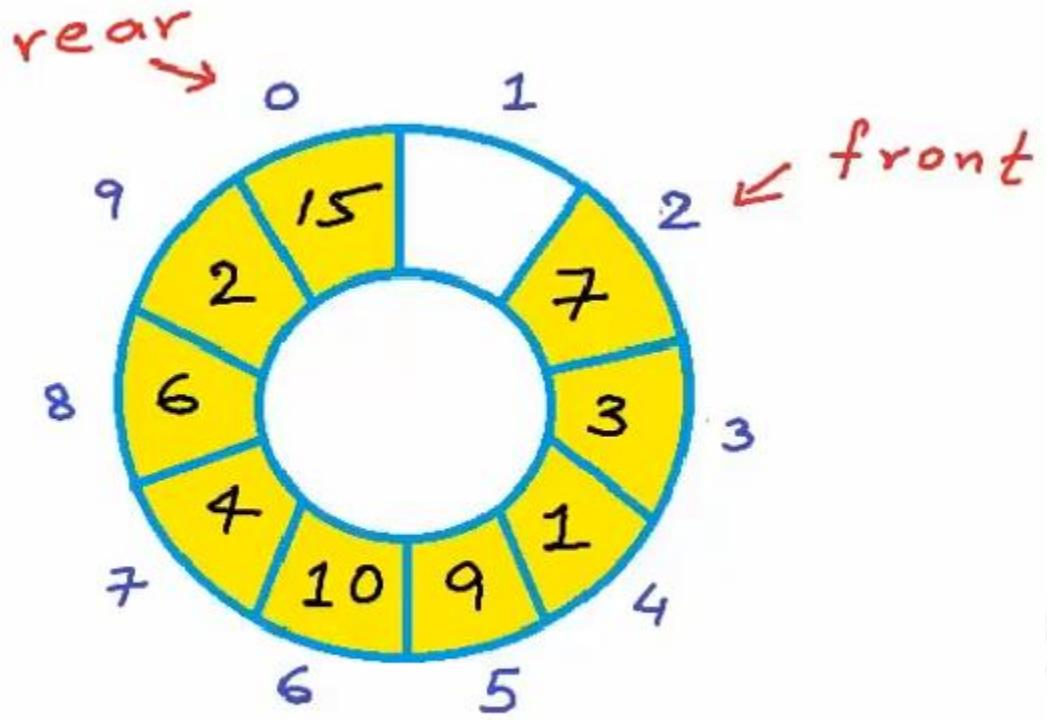
- Let us now try EnQueue function on our Queue.



PseudoCode for Queue using Circular Array (6/9)

- Call EnQueue function.

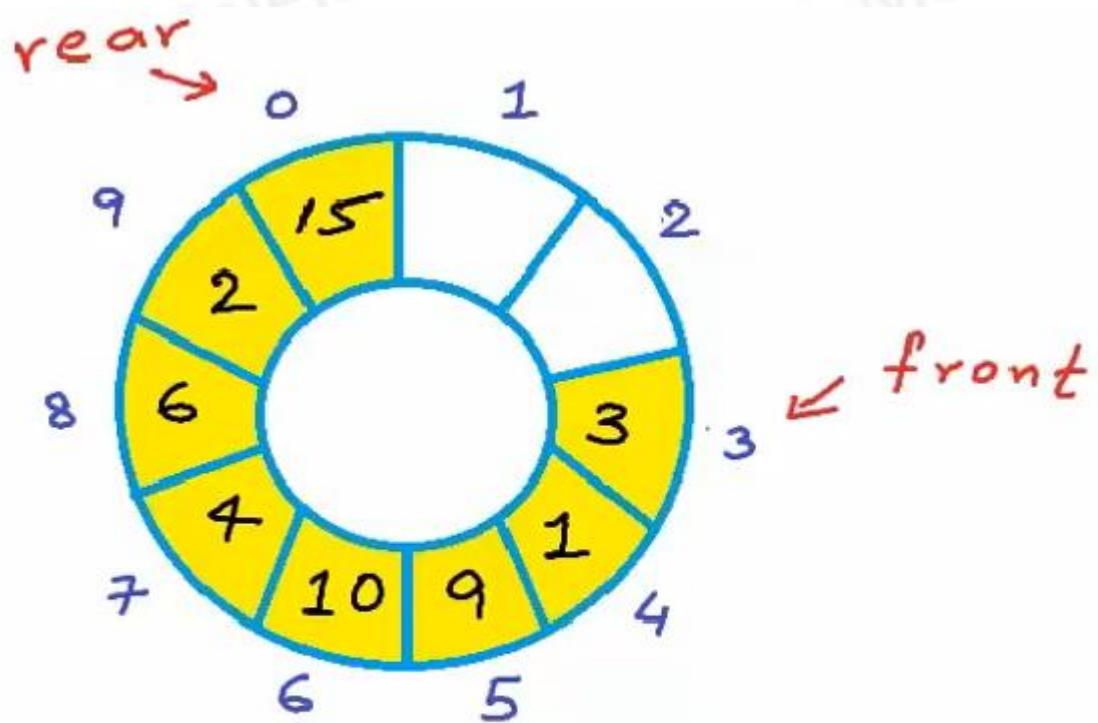
$$\frac{(rear + 1) \% N}{(9 + 1) \% 10} = 0$$



PseudoCode for Queue using Circular Array (7/9)

- Call DeQueue function.

$$\frac{(\text{front} + 1) \% N}{(2+1) \% 10}$$



*Enqueue(15)
Dequeue()*

PseudoCode for Queue using Circular Array (8/9)

- Now, front operation will be straight forward.
- We only need to return the element at front index.

```
front()
{
    return A[front]
}
```

PseudoCode for Queue using Circular Array (9/9)

- ▶ All the previous functions will take constant time, the time complexity will be $O(1)$.
- ▶ We are performing simple arithmetic and assignments in the functions.
- ▶ What about the Queue Implementation using Linked list!
 - ▶ Think about it.

Any Questions???