

- A
- 1) Onur Yaman
 - 2) 2007961

B

- 3) True
- 4) True
- 5) True
- 6) False
- 7) True
- 8) True
- 9) False
- 10) True

C

- 11) proper
- 12) $O(h)$
- 13) $O(n \log n)$
- 14) True
- 15) $O(n)$
- 16) cycle

D

- 17) Complete binary tree (of heap data structures) is a binary tree such that
 - At all level d where $0 \leq d \leq h-1$ each node has two children
 - At last level, each nodes are in leftmost position.
- 18) "Stable" means, after sorting, positions of all pairs whose keys are equal must be same position.
For example, $(3, d), (2, c), (2, b), (1, a)$
after sorting, $(1, a), (2, c), (2, b), (3, d)$
it is stable

E

19) A and C

20) min, remove-min, add, first, last

21) radix sort

F

22) preorder(T, p):

visit(p)

for each c in $T.children(p)$

if not c visited

visit(c)

preorder(T, c)

G

23) (a) 5, 6, 23, 16, 20, 18

(b) left child of node 18

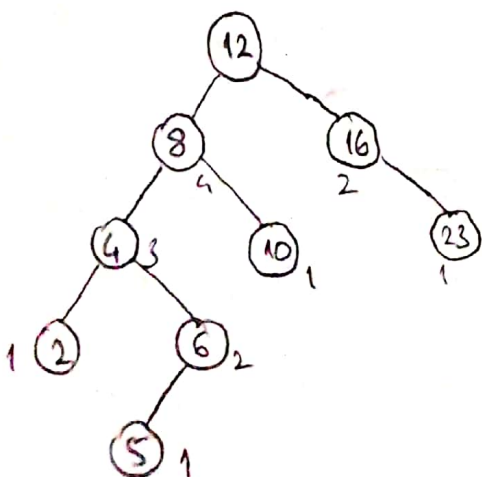
24) (a) Balanced

(b) Insert 5, then it will be unbalanced.

The pivot is 8.

Case 2, double rotation should be applied. Firstly between 4 and 6.

Then, between 6 and 8.



$$34 \equiv 1 \pmod{11} \Rightarrow i=3 \Rightarrow \text{index} = 1+9 \pmod{11}$$

	1	23			45					34
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$$12 \equiv 1 \pmod{11} \Rightarrow i=4 \Rightarrow \text{index} = 1+16 \pmod{11} \\ = 6 \pmod{11}$$

	1	23			45	12				34
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is the final state.

d

28) merge-sort($S1, S2, S$)

$n = \text{len}(S)$

if $n < 2$:

return

$\text{mid} = n // 2$

$S1 = S[0:\text{mid}]$

$S2 = S[\text{mid}:n]$

merge-sort($S1$)

merge-sort($S2$)

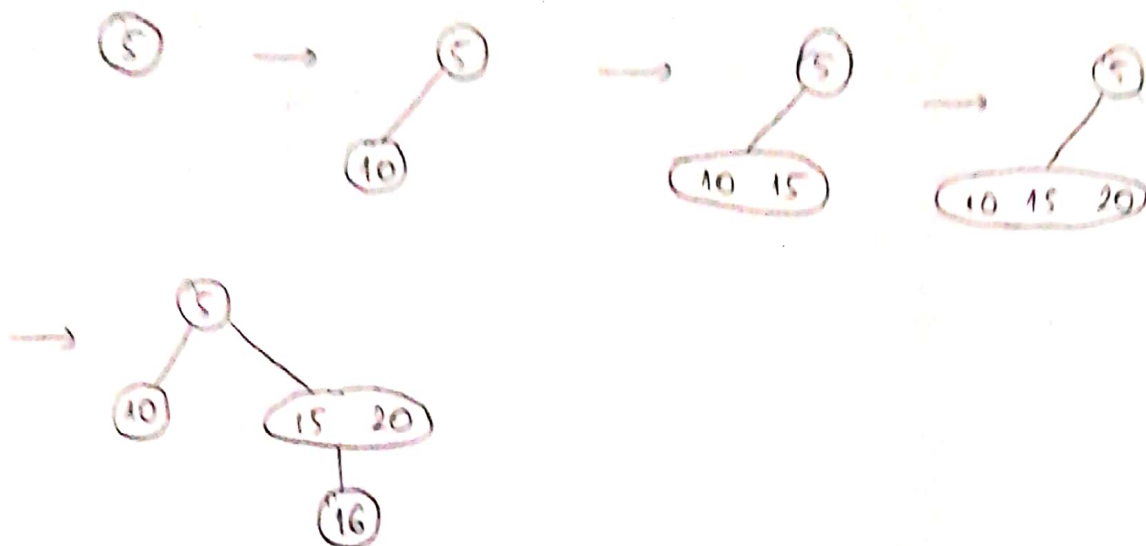
merge($S1, S2, S$)

29) Merge sort is a better solution when we want to sort sequences whose problem size is small (50, 100 etc)

We have $1073741847 = 2^{30}$ numbers, if we can take them into a hash table, $2^{30}/2^{10} = 2^{20}$ numbers can be kept in some sorted or unsorted table.

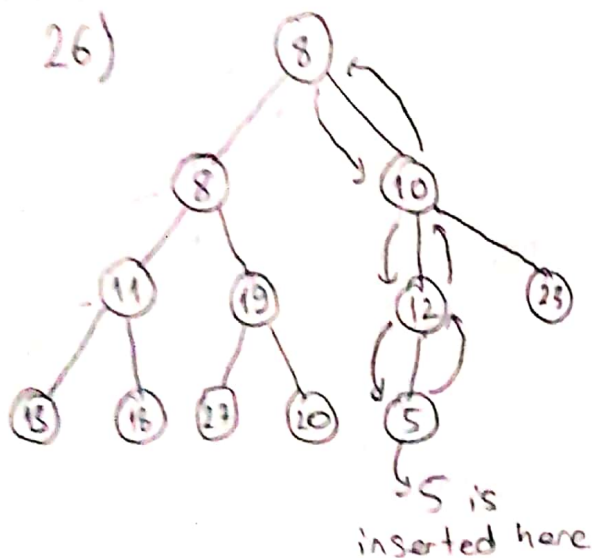
So, in this case, using bucket sort is better solution.

25)



H

26)



Now, by up-heap bubbling, swap operations are made between $(5, 12)$, $(5, 10)$, $(5, 8)$

→ now 5 is at root.

At this step, heap property is satisfied, 5 is smallest key and at root.

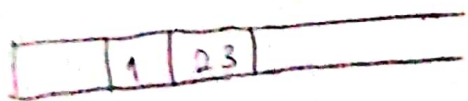
I

27) $1 \equiv 1 \pmod{11}$

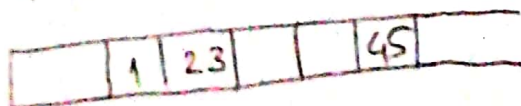
, note that $\text{index} = (h(x) + f(i)) \pmod{11}$
where $f(i) = i^2$



$23 \equiv 2 \pmod{11}$



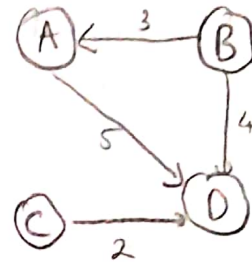
$45 \equiv 1 \pmod{11} \Rightarrow \text{collision} \Rightarrow i=2 \Rightarrow \text{index} = 1+4$



K

30) The graph is directed weighted. Let $M[i][j]$ be the adjacency matrix of the graph.

$$M[i][j] = \begin{matrix} & \begin{matrix} A & B & C & D \end{matrix} \\ \begin{matrix} A \\ B \\ C \\ D \end{matrix} & \begin{bmatrix} \infty & \infty & \infty & 5 \\ 3 & \infty & \infty & 4 \\ \infty & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty \end{bmatrix} \end{matrix}$$



∞ : there is no edge from i to j or j to i or i to i