

$\text{hom}(n \geq 2, n \in \mathbb{N})$

( $\exists$   $\varphi$ ) Some  $n$  have  $\varphi$  or  $\neg\varphi$  occur

recno  $\neq 0$ ,  $\neg(n-1)$   $\neg\varphi$  or  $\varphi$  appear

( $\varphi$  occurs)  $\rightarrow$   $\neg\varphi$  does not appear

$\neg\varphi$  occurs  $\neg\varphi$  does not appear

( $\varphi$  occurs).

$\neg\varphi$  occurs  $\neg\varphi$  does not appear

+ C.I.K. in  $\neg\varphi$  occurs

$\neg\varphi$  does not occur  $\neg\varphi$  occurs

$\neg\varphi$  does not occur  $\neg\varphi$  occurs

•  $\neg\varphi$  occurs  $A = \exists_{\text{num}} |A| = 2$  occurrences

• No  $\neg\varphi$  occurs  $\neg A = \exists_{\text{num}} |A| = 1$  occurrence

$\boxed{2 \ 3 \ 1 \ 3 \ 4}$

$\boxed{2 \ 1 \ 2 \ 1}$

tokens equal where  $C + D = 0$  tokens.

$\neg A = 1$  occurrence

$\boxed{2 \ 1 \ 2 \ 1}$

Desoxyribose  $\Rightarrow$  5'-OH

Deoxyribose, 2'-OH  $\Rightarrow$  2'-OH by endoform

T.O.  $\Rightarrow$  2'-OH  $\Rightarrow$  oxygen & oxygen

Chancery  $\Rightarrow$  to be keto  $\Rightarrow$  C=O  
G = for, G = en en C aliphatic  $\Rightarrow$  C=C

Characteristic occurrence  $\Rightarrow$  C=C  $\Rightarrow$  2-hydroxy (3D)  
2-hydroxy  $\Rightarrow$  2-hydroxy (3D)

Because both C or CH<sub>2</sub> change T.O.  
nucleic acid - hydrogen

[Th] C=C oxygen & carbonyl

stage:

a)  $\Rightarrow$  G and 3D  $\Rightarrow$  C=C and O  $\Rightarrow$  C=C and O  
b)  $\Rightarrow$  G and 3D  $\Rightarrow$  C=C and O  $\Rightarrow$  C=C and O

Given  $\lambda$  zeros  $\Rightarrow$   $n - \lambda$  non-zero  
 $\rightarrow$  Create  $\lambda$  rows.  
 $\rightarrow$   $(n + \lambda) - \lambda$  non-zero rows.

The Order.

( $\exists \varepsilon > 0$ )  $\forall n \in \mathbb{N} \exists N \in \mathbb{N}$   $\forall m > N$   $|f_n(x) - f_m(x)| < \varepsilon$

( $\forall \varepsilon > 0$ )  $\exists N \in \mathbb{N}$   $\forall m, n > N$   $|f_n(x) - f_m(x)| < \varepsilon$

$$|f_n(x) - f_m(x)| = |f_n(x) - f_{m+1}(x) + f_{m+1}(x) - f_m(x)| \leq |f_n(x) - f_{m+1}(x)| + |f_{m+1}(x) - f_m(x)|$$

$\lim_{m \rightarrow \infty} f_m(x) = f(x)$   $\lim_{n \rightarrow \infty} f_n(x) = f(x)$

$$\text{Hence } f(x) = \lim_{n \rightarrow \infty} f_n(x) > \text{i.e. } f(x) = \lim_{n \rightarrow \infty} f_n(0) = \lim_{n \rightarrow \infty} u_n$$



continuous function  $f(x) = \lim_{n \rightarrow \infty} f_n(x)$

continuous function  $f(x) = \lim_{n \rightarrow \infty} f_n(x)$   $\Rightarrow$   $f(x) = \lim_{n \rightarrow \infty} f_n(0) = \lim_{n \rightarrow \infty} u_n$

$f(x) = \lim_{n \rightarrow \infty} f_n(0) = \lim_{n \rightarrow \infty} u_n$   $\Rightarrow$   $f(x) = \lim_{n \rightarrow \infty} u_n$

A  $\log V$   $\forall$

$\{a_i\}_{i \in I} : a_i \in \mathbb{R} \subset \mathbb{R}$

equivalence class.

$G' \cong \langle V', \epsilon' \rangle$ .

$V' \leq V/\mathbb{R}$ ,  $[a]_{\mathbb{R}} \cup [b]_{\mathbb{R}} \in V'$

$[a]_{\mathbb{R}} \in T[b]_{\mathbb{R}} \iff ([a], [b]) \in \Delta_{\mathbb{R}}$

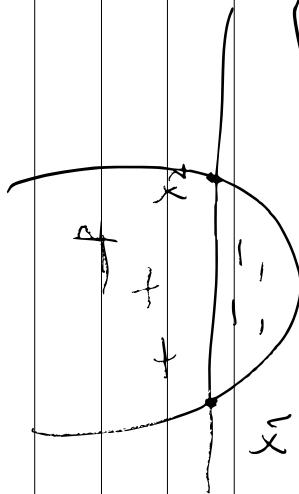
$|V'| = |G|$

$|G'| = \binom{|V|}{2} = \frac{|G| \cdot (|G|-1)}{2} \leq |G| = m$

$$\sum_{j_2} \chi(G) - (\chi(G) - 1) = 2 \cdot w$$

$$\overbrace{\chi(G) - \chi(G)}^{\chi_1, \chi_2} - \overbrace{2w}^0 \leq 0$$

$$D = 1 + 4 \cdot 2 \cdot w = \frac{1+8w}{1+8w} = \frac{1}{2} + \frac{\sqrt{1+8w}}{2} = \\ \chi_1, \chi_2 = \frac{1}{2} + \frac{\sqrt{1+2w}}{2}$$



$$w \in \{0\} \cup \left\{ \frac{1}{2} - \sqrt{\frac{1+2w}{4}} \right\} \cup \left\{ \chi(G) \right\} \subseteq \frac{1}{2} + \left[ \sqrt{\frac{1+2w}{4}} \right]$$

(302)  $G = \langle V, E \rangle$  e t-permanenten steht

$\forall t > 0$  rechts hoch-dimensionale wurzen ( $A^t$ )

e coenzym A.

Also, da  $|V| = n \geq 2t$ .

$G = \langle V, E \rangle$ ,  $|V| = n$ ,  $(\forall v \in V) [d(v) = t]$

$t \in \mathbb{N}_{>0}$ .

Wegen  $\forall v \in V$ . Torsen und  $v_1, v_t \in V$  oder no geeignete Personen,  $t \geq 2$ ,  $v_1 \neq v_t$ ,  $v_1 \neq v_t$ ,  $v_1 \neq v_t$ .

Heraus  $i \in \{1, \dots, t\}$ .

Torsen  $d(v_i) = t$ . Beweis gesamt



Durchsuche alle so wobei  $v_k = v_0, k \in \{1, \dots, t\}$ ,

$$0 \leq k_1, k_2, \dots, k_t \leq n.$$

Тодоқ жаңынан көрсөткіштің

мүннінің көрсөткіштің

тәсілінде  $v_1, v_2, \dots, v_n$  - деңгээлдегі

?  $v_0 \in \{v_1, \dots, v_n\}$ ?

Демек соң мәселенеңінде  $v_0 \in \{v_1, \dots, v_n\}$ .

$v_0 = v_1, \dots, v_n$  сандардан соң  $v_{n+1}$  - деңгээлдегі

біраңынан соң тәсілінде оның

тәсілінде  $v_1, v_2, \dots, v_n, v_{n+1}$  - деңгээлдегі

def l dominant bot:

$\exists \alpha = \langle \alpha_0, \dots, \alpha_n \rangle$  |  $\forall i \in \{0, 1, \dots, n\}$   
ce napure n-vezen čívanem kata  
n-reprutie beztoru co je naproxobere  
nákyda.

def l Tora je beztor  $\langle \alpha_0, \dots, \alpha_n \rangle \in \mathbb{B}^n$   
ce napure reducido.

$$w(\langle \alpha_0, \dots, \alpha_n \rangle) = \sum_{i=1}^n \alpha_i$$

$$\exists \alpha : w(w) \cup \mathbb{B}^n \rightarrow \mathbb{N}$$

def |  $\mu$ -бото  $\alpha^u$  =  $\lambda d \mid d \in B^u, \mu(d) = u$

ce наприк  $\lambda$ -и сноз  $\lambda Q$  н-неприм  
наприк кыз

def | Позитори не хенес элы

оз бекор  $\alpha^1 \in B^u$  е

наприк түншено  $n$

$$g(\alpha, \alpha^1) = \sum_{i=1}^n |\alpha_i - \alpha^{1i}|$$

Позитори  $\alpha^1 \in B^u$  ка секчий,

ако  $g(\alpha, \alpha^1) = 1$  и наприк бекор,

ако  $g(\alpha, \alpha^1) = n$ .

Тако некоторые свойства OT  
секчий бекор  $\lambda Q$   $B^u$  е  
наприк пекор но кыз

def | Ar-BOTO  $B_k^m(d) = \{d' \mid eB^m(f(d')) = k\}$

ce n'oppose coproducto

$S_k(d) = \{d' \mid d' \in B^m, f(d') \leq k\} -$

def | Hochschild dimension d'ord' , de or' .  
Variables concernant d'

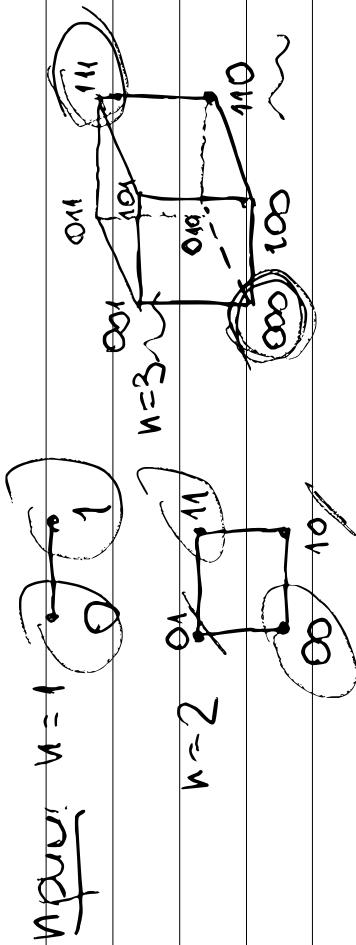
Responde no x'ido fu co lepricq  
repuna, c'sontrabusing do n'ale que  
 $(A_i, e_A, f_A, B_A^m, d_A, f_A(d_A) = 1)$ .

Recurato k ce rapporta q'anualmo no  
Bepusata.

def l House no subset d = <math>\{q\_1, \dots, q\_n\} \in \mathcal{C}\_k</math>

ce nspurz vnorzo :

$$\nu(\alpha) = \sum_{i=0}^{n-1} \alpha_i \cdot 2^{n-i}$$



a)  $|B^n| = ?$

b)  $|B_{\ell}^n| = ?$

- c)  $\nu((1, 0, 0, 0) \rightarrow) = ?$
- d)  $\exists \alpha \exists \beta \exists \gamma \nu(\alpha \rightarrow \beta \wedge \gamma) = 19 ?$
- e)  $\exists \alpha \forall \beta \alpha \rightarrow \beta \wedge \beta \in B^n ?$
- f)  $\exists \alpha \forall \beta \neg \alpha \rightarrow \beta \wedge \beta \in B^n \wedge \neg c.$   
 $2^{n-1} \leq \nu(\alpha) < 2^n$

$$G = \underbrace{\langle B^n, \exists \alpha, \exists \beta, \exists \gamma \mid \alpha, \beta, \gamma \in B^n \wedge g(\alpha, \beta, \gamma) = 1 \rangle}_{c}$$

g)  $|V| = ?$

D)  $|D_{\ell}^n| = \binom{n}{k}$

$D_{\ell}^n = \{ \alpha \mid \alpha \in \Delta^n \wedge w(\alpha) = k \}$

$D_{\ell}^n(\langle 0, \dots, 0 \rangle) \rightarrow g(\alpha, \langle 0, \dots, 0 \rangle) = k$

$$|G| = \frac{q^n - 1}{q - 1} = q^{n-1}(q - 1)$$

$\langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle$

$$\text{Vekt } \alpha \in \mathbb{R}^n. \quad \alpha = \langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle.$$

$\langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle$

$$\langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle = \langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle$$

$$|\alpha| = \sqrt{\alpha_1^2 + \alpha_2^2 + \dots + \alpha_n^2}$$

$$\langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle = \langle \alpha_1, \alpha_2, \dots, \alpha_n \rangle$$

$$\alpha = \sum_{i=1}^{n-1} \alpha_i e_i$$

$$c) \nu((1, 0, 1, 0, 1, 0)) = 2 + 2^4 + 2^5 = 50$$

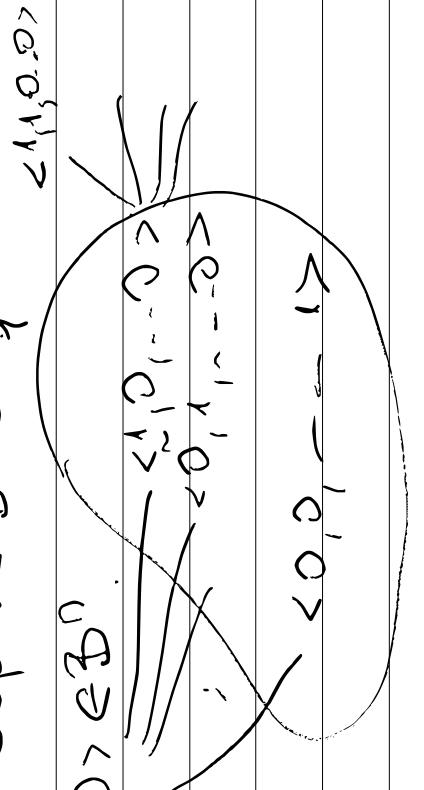
e) Spezielle Vektoren?

$$d) \text{ Zeigt, dass } \nu((\alpha)) = |\alpha|$$

$$e) \nu((1, 0, 1, 0, 1, 0)) = ?$$

f) Fóra'í Bépxore  $d \in \mathbb{B}_{k+1}^n + c$ .

$$\begin{aligned} & \underbrace{2^{n-1}}_{\text{2}^{n-1} \leq v(d)} \leq v(d) < \underbrace{2^n}_{2^{n-1} + 1} \quad \checkmark \\ & C \subseteq \{ \alpha \mid \alpha \in \mathbb{B}_k^n \wedge \underbrace{2^{n-1} \leq v(\alpha)}_{v(\alpha) = \sum_{i=0}^{n-1} \alpha_i \cdot 2^{n-1-i}} \} \end{aligned}$$
$$v(\alpha) = \sum_{i=0}^{n-1} \alpha_i \cdot 2^{n-1-i}$$
$$|C| = \binom{n}{n-1}$$



a) Hier ist  $d = \sqrt{0,01 - 0,01} = \sqrt{0,02} > 0,01$

b) Distanz zweier Punkte ist die Wurzel aus der Differenz der Quadrate der Koordinaten.

c) Die Distanz zweier Punkte ist die Wurzel aus der Differenz der Quadrate der Koordinaten.

d) Eine unendlich Dimensionale Raumzeit ist ein Vektorraum.

e) Eine unendlich Dimensionale Raumzeit ist ein Vektorraum.

f) In der Raumzeit sind Punkte voneinander verschieden.

g) Es gibt unendlich viele Dimensionen.

h) Ein Vektorraum ist ein Raum mit einer Dimension.

sofort Heute nacht

$\omega(\alpha)$   $\% \alpha =$  important no  $\beta$

Non  $\alpha$   $\beta$   $\in$   $\Delta$   $\cup$   $\Delta^*$   $\cup$   $\Delta^{\text{ext}}$   $\cup$  countable.

$\omega(\alpha) \% \alpha = \omega(\alpha) \oplus (\omega(\alpha) \cap \Delta^{\text{ext}}).$

t.o.  $(\omega(\alpha) - \omega(\beta))$  even

No  $\alpha$   $\beta$   $\in$   $\Delta$   $\cup$   $\Delta^*$   $\cup$   $\Delta^{\text{ext}}$

$\omega(\alpha)''$

$$\sum_{i=0}^{|\alpha|} |\alpha_i - \beta_i| \rightarrow^{\text{c.c.}} \text{even } y$$

Cresce  $\alpha$  non c'è sequenza egual a  $\omega(\alpha)$

c)  $\exists \alpha \in \mathbb{R} \forall \epsilon > 0 \exists n \in \mathbb{N}$   $\forall \alpha_1, \dots, \alpha_n$

$$\text{such that } |\alpha - \alpha_i| < \epsilon \quad \text{for all } i = 1, \dots, n$$

Durch:  $\exists \alpha \in \mathbb{R} \forall \epsilon > 0 \exists n \in \mathbb{N} \forall \alpha_1, \dots, \alpha_n$   $\exists \alpha' \in \mathbb{R}$   $\forall i = 1, \dots, n$   $|\alpha - \alpha'_i| < \frac{\epsilon}{n}$

Def:  $\alpha' = \frac{1}{n}(\alpha_1 + \dots + \alpha_n)$

Zeigt:  $|\alpha - \alpha'| < \epsilon$

Seien  $\alpha_1, \dots, \alpha_n$  beliebige reelle Zahlen. Dann gilt:

$|\alpha - \alpha'| = |\alpha - \frac{1}{n}(\alpha_1 + \dots + \alpha_n)|$

$= \left| \alpha - \frac{1}{n}\alpha_1 - \frac{1}{n}\alpha_2 - \dots - \frac{1}{n}\alpha_n \right|$

$\leq \frac{1}{n}|\alpha - \alpha_1| + \frac{1}{n}|\alpha - \alpha_2| + \dots + \frac{1}{n}|\alpha - \alpha_n|$

$\leq \frac{1}{n}(\epsilon + \epsilon + \dots + \epsilon) = \frac{n\epsilon}{n} = \epsilon$

# Pensamiento

## ① Averciones de respuesta

$\Phi_{lex} \subseteq \Phi^n \times \Omega^n \subseteq (\Omega^n)^n \times (\Omega^n)^n \subseteq (\Omega^n)^{n^2}$

 $\subseteq (\Omega^n)^n \times (\Omega^n)^n \subseteq (\Omega^n)^{n^2}$ 
 $\subseteq (\Omega^n)^n \times (\Omega^n)^n \subseteq (\Omega^n)^{n^2}$ 
 $\subseteq (\Omega^n)^n \times (\Omega^n)^n \subseteq (\Omega^n)^{n^2}$ 
 $\subseteq (\Omega^n)^n \times (\Omega^n)^n \subseteq (\Omega^n)^{n^2}$ 

$$\bigcup_{i=1}^n \{\Phi_i = b_i\}$$

$$\Phi(\vartheta) = \bigcup_{i=1}^n \{\Phi_i = b_i\} \subseteq D(\vartheta)$$

## ② No universal

$$\begin{aligned} & \exists \vartheta_1, \vartheta_2 \in \Phi^n \quad (\Phi(\vartheta_1) \neq \Phi(\vartheta_2)) \\ & \forall \vartheta \in \Phi^n \quad \Phi(\vartheta) \neq \emptyset \end{aligned}$$

## ③ Algoritmo

$$\begin{aligned} & \forall i \in \{1, \dots, n\} \quad \Phi_i = \{b_i\} \\ & \forall i \in \{1, \dots, n\} \quad \Phi_i = \{b_i\} \subseteq D(\vartheta) \\ & \forall i \in \{1, \dots, n\} \quad \Phi_i = \{b_i\} \subseteq D(\vartheta) \end{aligned}$$

203) Hora  $O \subseteq C \subseteq U$

$A(\alpha) = \{p \in P^n \mid p \in C\}$  este corect?

Corecteaza propozitia enunciata mai:

- a)  $A(\alpha) \cap P^n \neq \emptyset \text{ si } d \in A(\alpha)$
- b)  $A(\alpha) \cap P^n \neq \emptyset \text{ si } C \subseteq P^n$
- c)  $A(\alpha), d \in P^n$ .

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și în prezumătuare se separă
- b) V. (4-5)! pozitivă în raportarea  
c) rezultatelor.
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