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$$f = \langle \mathbb{N}, \tau, p; a, b, f \rangle$$

$$a, b \in \mathbb{N}, \tau \subseteq \mathbb{N} \times \mathbb{N}, f: \mathbb{N} \rightarrow \mathbb{N}$$

$$\tau \subseteq \mathbb{N} \times \mathbb{N}$$

$$p: \mathbb{N} \rightarrow \mathbb{N}$$

$$x \in \mathbb{N} \quad f(x)$$

$$y = x$$

$$\tau \subseteq \mathbb{N} \times \mathbb{N}$$

$$f_4 - \text{so } a, b \in \mathbb{N} \times \mathbb{N} \text{ and } y, \tau, y. \quad \langle x, y \rangle, \langle y, x \rangle \in \tau$$

$$\langle x, f(y) \rangle \in \tau$$

$$\tau(x, y) \Leftrightarrow p(x, y) \neq \tau(y, x)$$

$$N$$

$$(0,0) \rightarrow (1,1) \rightarrow (2,2) \rightarrow (3,3) \dots$$

$$f^B(x) = \begin{cases} \underline{A+1}, & x \text{ clone } \underline{x=n'} \\ \underline{n'}, & \underline{x \text{ original}, x \geq n} \end{cases}$$

$$a) \quad x < f(y)$$

$$b) \quad \underline{x < f(y)}$$

$$\rightarrow f$$

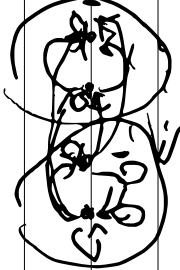
$$\rightarrow f$$

$$\forall x \in A$$

$$x \in N : a) \quad x' \text{ yes}$$

$$b) \quad x \neq$$

$$B = \{N \cup N'\}$$



$$x' \in N' : a) \quad x' \text{ yes}$$

$$b) \quad x' \text{ yes}$$

$$B \approx 0, B \approx 1$$

Ex:

$$\forall x \in A$$

$$0$$

Prilozhenie

L - структура от множества

$c \in T(L)$ дан. В-то от дървета $con(x, y)$,
където x и y са елементи на L .
непосредствено следващо.

$P(L, x)$ - то е x и y са елементи на L и $x \in T(L)$

$P(L, x) := \text{subsequence}(L, S)$, permutation (S, R) ,
merge Append (P, x) .

subsequence $([T], [S])$.

subsequence $([H|T], [H|R])$:- subsequence (T, R)

subsequence $([T], [R])$:- subsequence (T, R) .

merge Append $([T], [S])$.

merge Append $([H|T], R)$:- merge Append (T, R) ,
append (H, R, R) .

Def 1.1

$f = \langle N, P^f \rangle$: $P^f(a, b, c) \Leftrightarrow \underline{a - b = c^2}$
 Boolean circuit comp.

$$a \leq b$$

$$a \geq b$$

Exists $n \in \mathbb{N}$ such that circuit-Bo is equivalent to circuit-Bo

$$\overline{\overline{P(N)}} = \mathbb{R}$$

$$h(n) = \langle n \rangle \rightarrow h(n) = n$$

$$\text{for } p(x, x, x)$$

$$x = x^2$$

$$\{1\} \quad \{2\} \quad \{3\} \quad \{4\} \quad \{5\} \quad \{6\} \quad \{7\} \quad \{8\} \quad \{9\} \quad \{10\}$$

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$$\varphi_2(a > b \rightarrow a = b^2 \exists y(\varphi_0(y) \wedge p(x, y, z)))$$

$$a_1 = (a^2) \quad -c = b^2 \quad a_1 - c = b^2 \quad p(a_1, c, b)$$

$$\exists a_1(\varphi_{r2}(a, a_1) \wedge p(c, a_1, b))$$

$$\exists \varphi_0(c)$$

$$a = b^2 = (b^2)^2 = (b^1)^2$$

$$\exists b_1(\varphi_{r1}(b, b_1) \wedge \varphi_{r2}(b_1, a))$$

~~2~~ \mathcal{H}, r

$$\langle a, b \rangle \in \Gamma^* \hookrightarrow a = b + 7$$

$$\langle a, b \rangle \in p^* \hookrightarrow a = 7b$$

$$\exists x, p(x, x)$$

$$7 \cdot 7 \cdot x = 48 + x$$

$$49x = 48 + x$$

$$x = 1$$

$$49 \cdot 2 = 48 \cdot 2 + 2$$

$$49(3) = 48 \cdot 3 + 3$$

[illegible]