

песня: 8:00 утренняя

$\Sigma \neq \varnothing$, Σ' е известно

$\Sigma \cup \Sigma' \in$
известно

$\Pi \neq \varnothing$, $\Sigma \neq \varnothing$, $\Sigma \neq \varnothing$, $\Sigma \neq \varnothing$ $\xrightarrow{\text{Решающий}}$

$\exists x \forall y (q(x, y) \Leftrightarrow \forall z (q(y, z) \Rightarrow (q(z, y) \Rightarrow$

$\forall x (q(x, y) \Rightarrow q(x, y))) \models$

известно?

$(\forall x \forall y (q(x, y) \Leftrightarrow \forall z (q(y, z) \Rightarrow \neg q(z, y)))$

① Вкорме отрицания \neg стандартны
 ф-ли и правила \Leftrightarrow "и" \Rightarrow .

Правила: $\neg \forall x \varphi \vdash \exists x \neg \varphi$

$\neg \exists x \varphi \vdash \forall x \neg \varphi$

$\neg (\varphi \vee \psi) \vdash \neg \varphi \wedge \neg \psi$

$\neg (\varphi \wedge \psi) \vdash \neg \varphi \vee \neg \psi$

$\neg \neg \varphi \vdash \varphi$

$\varphi \Rightarrow \psi \vdash \neg \varphi \vee \psi$

$\varphi \Rightarrow \psi \vdash (\varphi \Rightarrow \psi) \wedge (\psi \Rightarrow \varphi) \vdash$

$(\neg \varphi \vee \psi) \wedge (\neg \psi \vee \varphi)$

$\varphi \vdash \exists x \forall y (q(x, y) \Rightarrow \forall z (\neg q(y, z) \vee \neg q(z, y)))$

$\vdash \exists x \forall y ((\neg q(x, y) \vee \forall z (\neg q(y, z) \vee \neg q(z, y))) \wedge$

$(\exists z (q(y, z) \wedge q(z, y)) \vee q(x, y))$

\vdash

\vdash

\vdash

② Под кванторизацией \rightarrow отправка

$\varphi \in \overline{Q_{x_1} \dots Q_{x_n} \Phi}$, к zero:

- $Q_i \in \exists, \forall$, $1 \leq i \leq n$
- Φ в безв. ф-но

$x \in \text{Var}^{\text{free}}[\varphi] \rightarrow \underline{\underline{x_j}}$

$$\underline{\exists x}(\varphi \vee \psi) \models \exists x(\varphi \vee \psi)$$

$$\underline{\exists x}(\varphi \wedge \psi) \models \exists x(\varphi \wedge \psi)$$

$$\forall x(\varphi \vee \psi) \models \forall x(\varphi \vee \psi)$$

$$\forall x(\varphi \wedge \psi) \models \forall x(\varphi \wedge \psi)$$

Th 22 верно.

$$\text{Var}^{\text{free}}[\varphi] \cap \text{Var}^{\text{bd}}[\varphi] \neq \emptyset$$

$$\forall y (\underline{\underline{\exists z}} (p(x, y) \wedge \exists z p(z, z))) \models$$

$$\forall y \exists z \forall x (p(x, y) \vee p(z, z))$$

$$\varphi'' \equiv \exists x \forall y \exists t \forall z ((\neg q(x, y) \vee \neg q(z, y) \vee \neg q(y, z)) \wedge (q(x, y) \vee (q(t, y) \wedge q(y, t))))$$

\uparrow \uparrow
 c f

③ Члор ψ нхд

φ^s equation

$\underbrace{s_1 \dots s_n}$ отсюда можно

• $\varphi \equiv \exists x \varphi \leadsto c \leadsto \mathcal{L}'$ снова нхд $\varphi_s \equiv \varphi[x/c]$

• $\varphi \equiv \forall y_1 \dots \forall y_n \exists x \varphi \leadsto f \leadsto \mathcal{L}'$ снова от.р. f

$$\varphi_s \equiv \forall y_1 \dots \forall y_n \varphi[x/f(y_1, y_n)]$$

~~$\varphi'' \equiv \varphi^s$~~

$\underbrace{s_1 \dots s_n}$

нхд можно из. $\exists c$ срелуг

$$\varphi^s \equiv \forall y \forall z ((\neg q(c, y) \vee \neg q(z, y) \vee \neg q(y, z)) \wedge (q(c, y) \vee (q(f(y), y) \wedge q(y, f(y)))))$$

$$(4) \quad \underline{ev(yz)} \vdash \underline{(evy)z} \quad (evy)$$

$$e \stackrel{fin}{=} \forall y \forall z ((\neg q(c, y) \vee \neg q(z, y) \vee \neg q(y, z)) \wedge \\ (q(c, y) \vee q(f(y), y)) \wedge \\ (q(c, y) \vee q(y, f(y))))).$$

KNOW, CHOW, KNOW

$$D_1 = \{ \neg q(c, y_1), \neg q(z, y_1), \neg q(y_1, z_1) \}$$

$$D_2 = \{ q(c, y_2), q(f(y_2), y_2) \}$$

$$D_3 = \{ q(c, y_3), q(y_3, f(y_3)) \}$$

111

Правила на играта

$\{x_1/x_2\}$

① Правило за резолюция

Нека D_1, D_2 са дизюнктни д-с одесни променливи.
Нека $D_1 = D_1' \cup \{L\}$, $D_2 = D_2' \cup \{L'\}$. Търсим
субституция σ , т.е. $L\sigma = L'\sigma$. Тогава
 $\text{Res}(D_1, D_2) = D_1'\sigma \cup D_2'\sigma = (D_1'\sigma \setminus \{L\sigma\}) \cup$
 $(D_2'\sigma \setminus \{L'\sigma\})$.

② Колоне на субституи. $n \geq 1$

Нека $D = D' \cup \{L_1, \dots, L_n\}$. Търсим субституи.
 σ , т.е. $L_1\sigma = L_2\sigma = \dots = L_n\sigma$.
Тогава $\text{Collapse}(D) = D' \cup \{L_1\sigma\} = D\sigma$

$$D_1 = \{ \neg q(c, y_1), \neg q(z, y_1), \neg q(y_1, z) \}$$

$$D_2 = \{ q(c, y_2), q(f(y_2), y_2) \}$$

$$D_3 = \{ q(c, y_3), q(y_3, f(y_3)) \}$$

① Possr. D_1 $\begin{cases} c = z_1 \\ z_1 = y_1 \end{cases} \quad y_1 = c$

$$\text{Collapse}(D_1, y_1/c, z_1/c) = \{ \neg q(c, c) \} = D_4$$

② Possr. $D_4 \cup D_2$ $\sigma = \sigma_1 \cup \sigma_2$

$$\text{Res}(D_4, D_2, y_2/c) = \{ \underline{q(f(c), c)} \} = D_5$$

③ Possr. $D_3 \cup D_4$

$$\text{Res}(D_4, D_3, y_3/c) = \{ \underline{q(c, f(c))} \} = D_6$$

④ Possr. $D_6 \cup D_1$

$$\text{Res}(D_6, D_1, y_1/f(c), z_1/c) = \{ \underline{\neg q(f(c), c)} \} = D_7$$

⑤ $\text{Res}(D_7, D_5) = \text{[scribble]}$

~~$\{p(x), \neg p(x)\}$~~ \rightarrow exercise

more $\underbrace{D_1 \subseteq D_2}$
 \downarrow
important!!!

$D_1 = \{p(x)\}$, $D_2 = \{p(x), q(x)\}$

(222) $\varphi_1 \& \varphi_2 \Rightarrow \varphi_3$? $\{\varphi_1, \varphi_2\} \models \varphi_3 \Leftrightarrow \{\varphi_1, \varphi_2\} \models \varphi_3$
 $\varphi_1 \equiv \forall x [\forall y (s(y) \Rightarrow p(y, x)) \Rightarrow \neg r(x)]$
 $\varphi_2 \equiv \exists y \forall x [s(x) \vee \neg q(x, y)]$
 $\varphi_3 \equiv \exists x \forall y [r(y) \Rightarrow \exists z (\neg p(z, y) \& \neg q(z, x))]$
 $\varphi_4 \equiv \forall x \exists y [r(y) \& \forall z (p(z, y) \vee q(z, x))]$

$\varphi_1 \& \varphi_2 \models \varphi_3$
 $\varphi_1 \& \varphi_2 \models \varphi_4$

$\varphi_1 \models \forall x [\forall y (\neg s(y) \vee p(y, x)) \Rightarrow \neg r(x)] \models 1$
 $\varphi_1' \equiv \forall x [\exists y (s(y) \& \neg p(y, x)) \vee \neg r(x)]$
 $\varphi_1'' \equiv \forall x \exists y [(s(y) \& \neg p(y, x)) \vee \neg r(x)]$
 $\varphi_1''' \equiv \forall x \exists y \forall z [r(y) \& (p(z, y) \vee q(z, x))]$

$$\varphi_1^S \equiv \forall x [(S(f(x)) \& \neg p(f(x), x)) \vee \neg r(x)]$$

$$\varphi_2^S \equiv \forall x [\neg S(x) \vee \neg q(x, c)] \vee$$

$$\varphi_3^S \equiv \forall x \forall z [\underbrace{r(q(x))}_{\text{true}} \& \underbrace{(p(z, q(x)) \vee q(z, x))}_{\text{true}}]$$

$$\hookrightarrow \varphi_1^{\text{fin}} \equiv \forall x [\underbrace{(S(f(x)) \vee \neg r(x))}_{\text{true}} \& \underbrace{(\neg p(f(x), x) \vee \neg r(x))}_{\text{true}}]$$

$$D_1 = \{ \underline{S(f(x_1))}, \neg r(x_1) \}$$

$$D_2 = \{ \neg p(f(x_2), x_2), \neg r(x_2) \}$$

$$D_3 = \{ \neg S(x_3), \neg q(x_3, c) \}$$

$$D_4 = \{ \underline{r(q(x_4))} \}$$

$$D_5 = \{ \underline{p(q(x_5))}, \underline{q(z_5, x_5)} \}$$

$$\textcircled{5} \text{Res}(D_3 \{x_3/x_4\}, D_4) = \text{true}$$

$$\textcircled{1} \text{Res}(D_1 \{x_1/g(x_4)\}, D_4) = \{ \underline{S(f(q(x_4)))} \} = D_6$$

$$\textcircled{2} \text{Res}(D_2 \{x_2/g(x_4)\}, D_4) = \{ \neg p(f(q(x_4)), q(x_4)) \} = D_7$$

$$\textcircled{3} \text{Res}(D_3 \{x_3/f(q(x_4))\}, D_6) = \{ \neg q(f(q(x_4)), c) \} = D_8$$

$$\textcircled{4} \text{Res}(D_5 \{z_5/f(q(x_4))\}, D_8) = \{ \underline{p(f(q(x_4)), q(x_5))} \} = D_9$$

$\mathcal{D}_1, \mathcal{D}_2$ ca gus okom.
 $L \in \mathcal{D}_1, \bar{L} \in \mathcal{D}_2$

$I \models \mathcal{D}_1, \mathcal{D}_2$

$\mathcal{D}_1' = \mathcal{D}_1 \setminus \{L\}, \mathcal{D}_2' = \mathcal{D}_2 \setminus \{\bar{L}\}$

$I \models \mathcal{D}_1 \cup \mathcal{D}_2'$ peronbenta ke $\mathcal{D}_1 \cup \mathcal{D}_2$ no
 $L(\bar{L})$

$\{q, r, s\}$

$\mathcal{D}_1 \cup \mathcal{D}_2$
 \checkmark peronbenta

$\{p, q\}$
 $\{p, r, s\}$

1. Никто не любит всех собак.
 2. Никто не любит всех шарляшек.
 3. Дворовые все с шарляшками.
- $n(x), g(x), m(x), yb(x, y)$

$$\varphi_1 \Leftarrow \exists x (\pi(x) \wedge \forall y (g(y) \rightarrow yb(x, y)))$$

$$\varphi_2 \Leftarrow \forall x (\pi(x) \rightarrow \forall y (m(y) \rightarrow \neg yb(x, y)))$$

$$\varphi_3 \Leftarrow \neg \forall y (g(y) \rightarrow \neg m(y)) \neq \exists y (g(y) \wedge m(y))$$

$\{\varphi_1, \varphi_2, \varphi_3\}$ — несовместимо.

$$\varphi_1 \neq \exists x \forall y (\pi(x) \wedge (\neg g(y) \vee yb(x, y)))$$

$$\varphi_2 \neq \forall x \forall y (\neg \pi(x) \vee \neg m(y) \vee \neg yb(x, y))$$

$$\text{снф}(\varphi_1) \Leftarrow \forall y (\pi(c) \wedge (\neg g(y) \vee yb(c, y)))$$

$$\text{снф}(\varphi_3) \Leftarrow (g(d) \wedge m(d))$$

$$\{g(d)\}, \{m(d)\}, \{\pi(c)\}, \{g(y), yb(c, y)\}$$

$$\{\neg \pi(x), \neg m(y), \neg yb(x, y)\}$$

$$D_1 = \{g(d)\}, \quad D_2 = \{u(d)\}, \quad D_3 = \{\pi(c)\}, \quad D_4 = \{\neg g(y), \neg b(c, y)\}$$

$$D_5 = \{\neg \pi(x), \neg u(y), \neg yb(x, y)\}$$

$$D_6 = \text{Res}(D_1, D_4 \{y/d\}) = \{yb(c, d)\}$$

$$D_7 = \text{Res}(D_2, D_5 \{y/d\}) = \{\neg \pi(x), \neg yb(x, d)\}$$

$$D_8 = \text{Res}(D_7 \{x/c\}, D_3) = \{\neg yb(c, d)\}$$

$$\text{A} = \text{Res}(D_6, D_8)$$