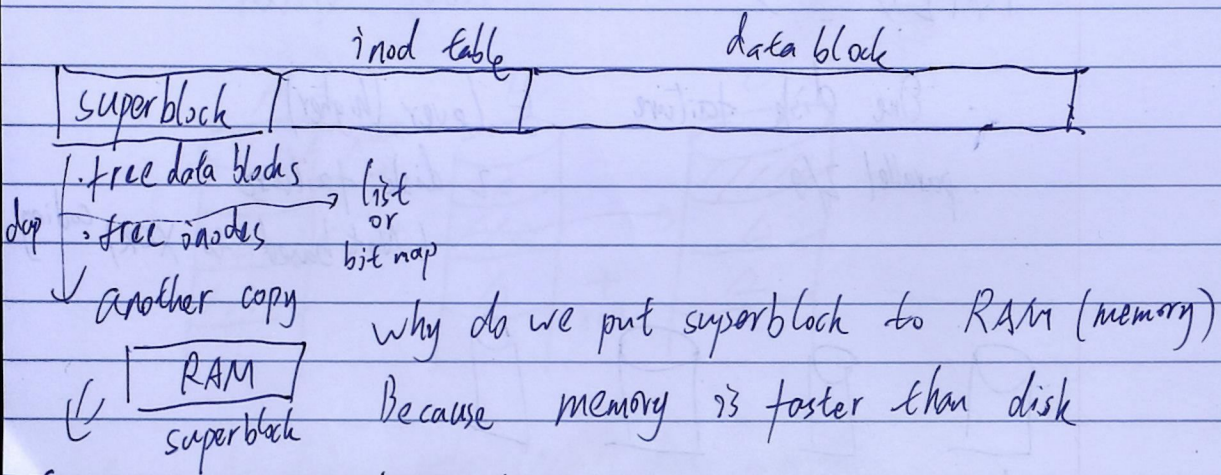


UNIX File System Consistency



Consistency when OS boot, OS copy it to RAM

a data block is either on the free list (super block) or assigned to a single file.

if the kernel freed a disk data block (with a file)

- ① returning the block number to the in-core copy of the superblock
- ② and ③ allocated the disk data block to a new file

"disk crashes" ~~assume~~ the file does not need this data block why consistency?

case 1: if the kernel wrote the inode and blocks of the new file to disk, but crashed before updating the inode of the old file to disk

lead to \Rightarrow two inodes would address the same disk block number

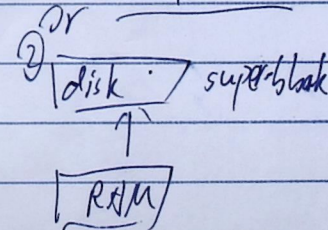
case 2) if the kernel wrote the super block and to free list to disk and crashed before writing the old inode out

this data block still in free list vs. the data block still in old inode \downarrow so issue occurs.

case 3) if the old file was written on disk and crashed before the superblock was written to disk.

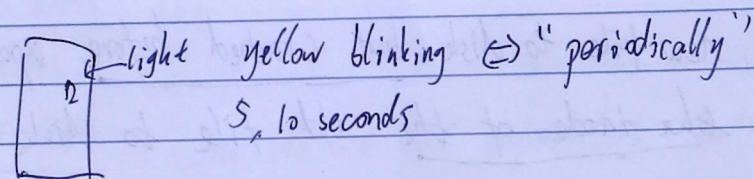
\Rightarrow this data block appear nowhere

- ① the old file delete the inode which point to this data block
- ② before update this data block to free list in ^① superblock (RAM)
- ③ OS crashed
- ④ this data block will be nowhere



How to solve these issues?

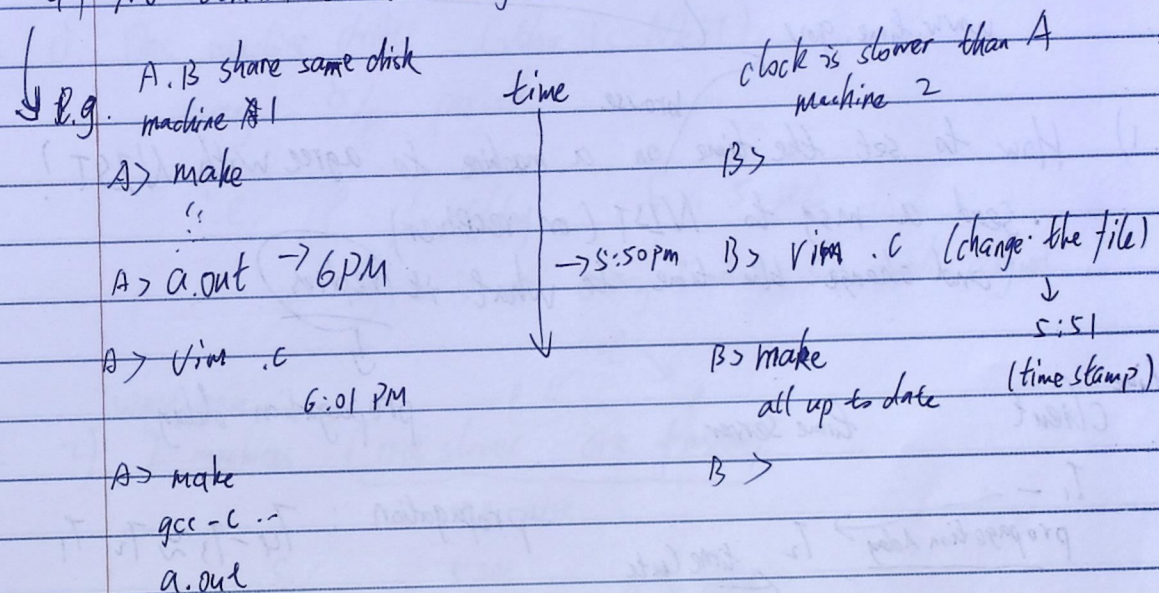
"Journalling"



Ch6 Synchronization in DS is 'Centralized Sys'

Properties of distributed algorithms:

- 1) The relevant info is scattered among machines
- 2) Each process make decisions based on local information
- 3) A single point of failure in the Sys should be avoided
- 4) No common clock or global time source exists



α processor timer (usually quartz crystal)

generate interrupts periodically

Each interrupt is called a clock tick

not perfect \Rightarrow clock skew

physical clock: agree with real time

logical clock: all machines agree on the same time
order