

# Chapter 2: Algorithm Analysis

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## Abstract

This section discusses about:

- how to estimate the time required for a program
- how to reduce running time of a program
- The result of careless use of recursion
- very efficient algorithms to raise a number to a power
- compute GCD

## 1 Mathematical Background

Four definitions of the framework:

**Definition 1.1** (Upper bound).  $T(N) = O(f(N))$  if there are positive *constants*  $c$  and  $n_0$  such that  $T(N) \leq cf(N)$  when  $N \geq n_0$

**Definition 1.2** (Lower bound).  $T(N) = \Omega(g(N))$  if there are positive *constants*  $c$  and  $n_0$  such that  $T(N) \geq cg(N)$  when  $N \geq n_0$

**Definition 1.3** (Envelope).  $T(N) = \Theta(h(N))$  iff  $T(N) = O(h(N))$  and  $T(N) = \Omega(h(N))$

**Definition 1.4.**  $T(N) = o(p(N))$  if for **all** positive constants  $c$  there exists an  $n_0$  such that  $T(N) < cp(N)$  when  $N > n_0$ .

Ref: p.30 for detail theorem

Typical growth rates:

Function	Name
$c$	Constant
$\log N$	Logarithmic
$\log^2 N$	Log-Squared
$N$	Linear
$N \log N$	
$N^2$	Quadratic
$N^3$	Cubic
$2^N$	Exponential

Some Theorem from the definition:

**Theorem 1.1.** If  $T_1(N) = O(f(N))$  and  $T_2(N) = O(g(N))$ , then:

1.  $T_1(N) + T_2(N) = O(f(N) + g(N))$
2.  $T_1(N) * T_2(N) = O(f(N) * g(N))$

**Theorem 1.2.** If  $T(N)$  is a polynomial of degree  $k$ , then  $T(N) = \Theta(N^k)$

**Theorem 1.3.**  $\log^k N = O(N)$  for any constant  $k$ , which tell us ‘Logarithms grow very slowly’

For big-O notation answers: Lower-order terms can generally be ignored e.g.  $f(N) = 2N^2 + N$  then its big-O notation is  $T(N) = O(N^2)$

For *relative growth rates* of two Function, we using ‘*L’Hopitals’s rule*’ to determine it:

**Theorem 1.4** (L’Hopitals’s rule). If  $\lim_{N \rightarrow \infty} f(N) = \infty$  and  $\lim_{N \rightarrow \infty} g(N) = \infty$ , then  $\lim_{N \rightarrow \infty} \frac{f(N)}{g(N)} = \lim_{N \rightarrow \infty} \frac{f'(N)}{g'(N)}$

Four possible results:

- Limit is 0:  $f(N) = o(g(N))$
- Limit is  $c \neq 0$ :  $f(N) = \Theta(g(N))$
- Limit is  $\infty$ :  $g(N) = o(f(N))$
- Limit does not exist: No relations

## **2 Model**

Basically a normal computer which execute instructions sequentially. And it has infinite memory

## **3 What to Analyze**