

Similarity of Binaries through re-Optimization

By

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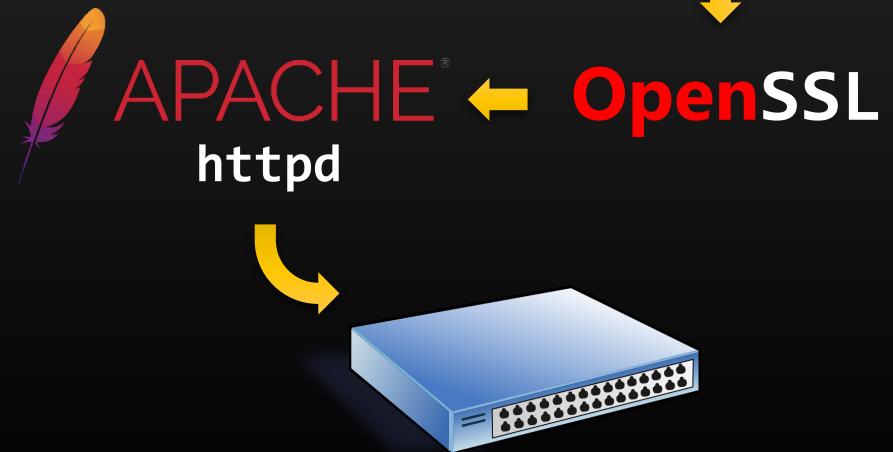


*The research leading to these results has received funding from the European Union's - Seventh Framework Programme (FP7) under grant agreement n° 615688– ERC- COG-PRIME.

Motivation



Developer

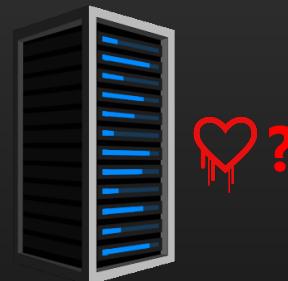


Motivation

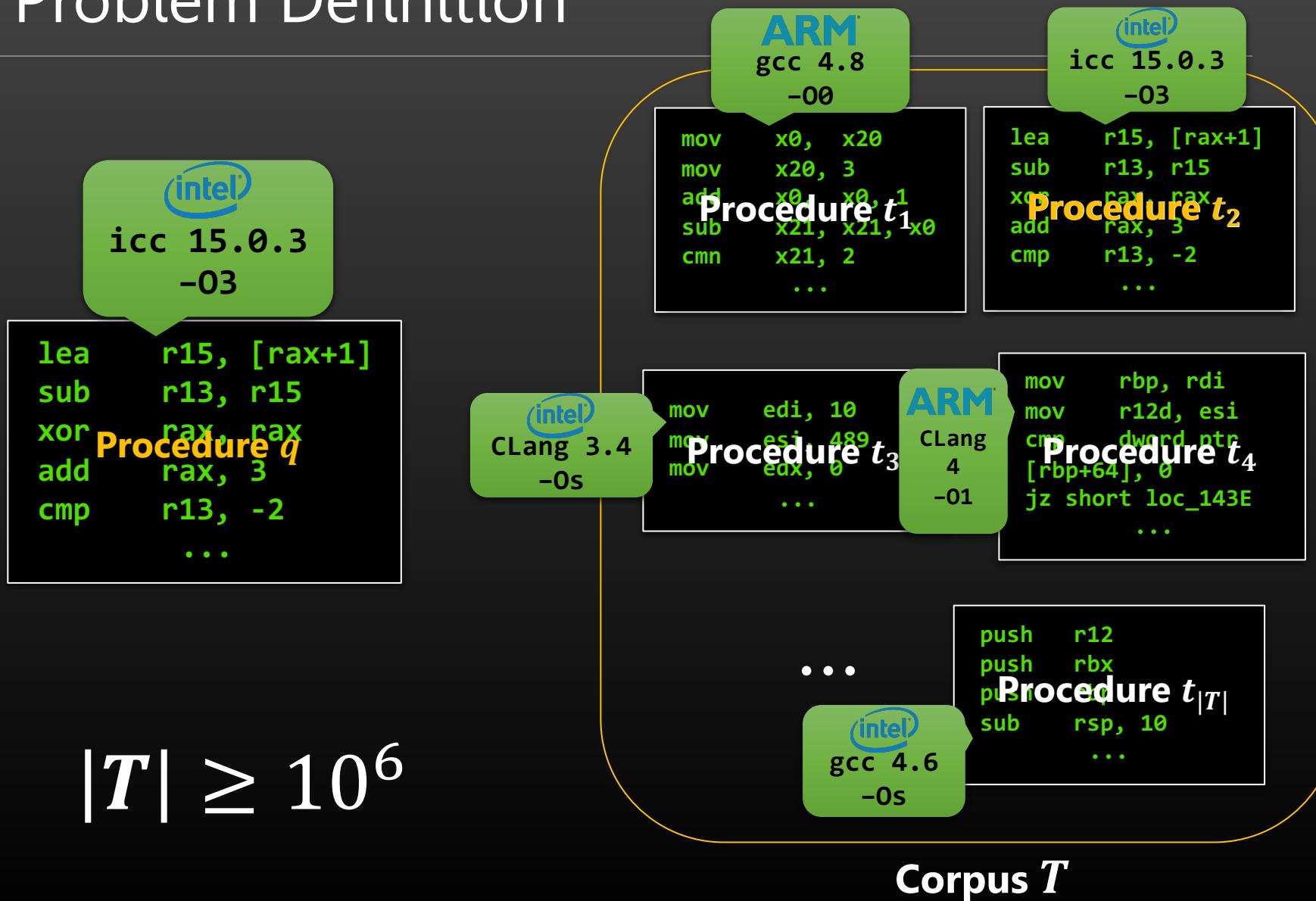


Security
Researcher

OpenSSL 



Problem Definition



Challenge

OpenSSL's dtls1_buffer_message()

```
mov    x0,  x20
mov    x20,  3
add    x0,  x0,  1
sub    x21, x21, x0
cmn    x21,  2
```

```
lea    r15, [rax+1]
sub    r13, r15
xor    rax, rax
add    rax,  3
cmp    r13, -2
```

ARM

gcc 4.8
-O0



icc 15.0.3
-O3

Our Approach

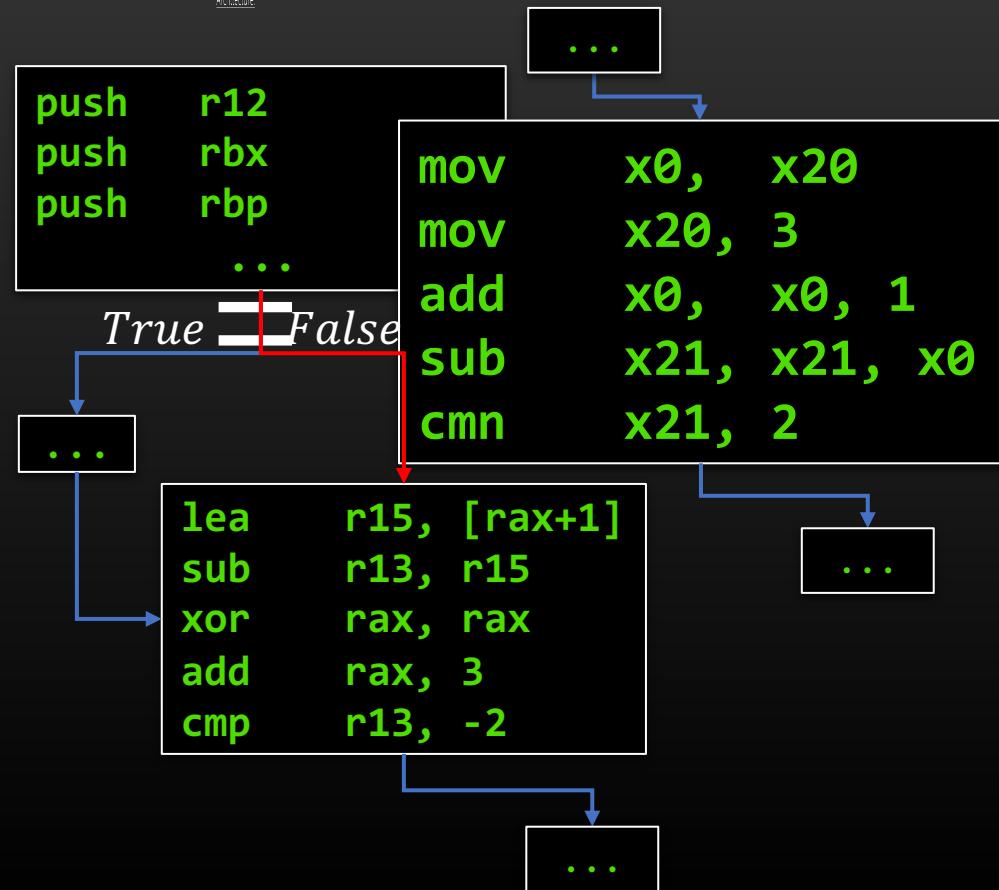
For finding Similarity of Binaries

Our Approach: What

Query q : dtls1_bQuery_qestslag1(b)uffer_message()

Compiler: icc 15.0.3 -O3 gcc 4.8 -O0

Architecture: Query qestslag1(b)uffer_message()
Compiler: icc 15.0.3 -O3
Architecture:

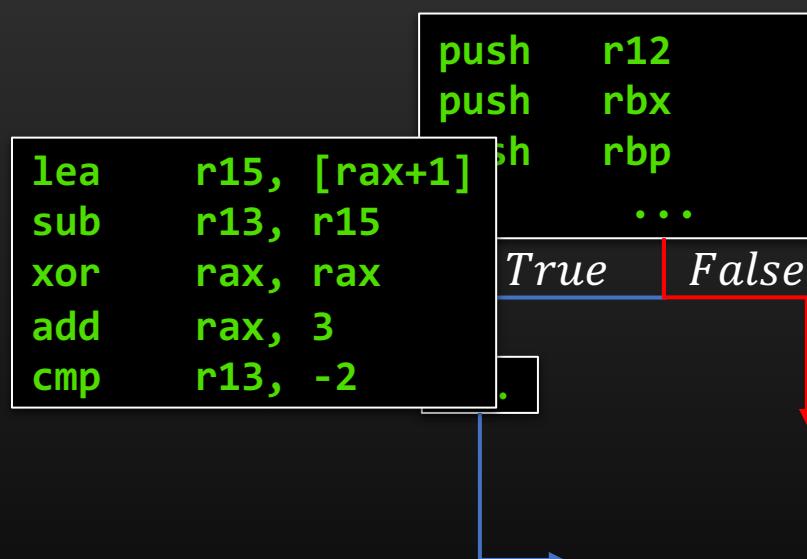


Our Approach: What

Query q : dtls1_buffer_message()

Compiler: **icc 15.0.3 -O3**

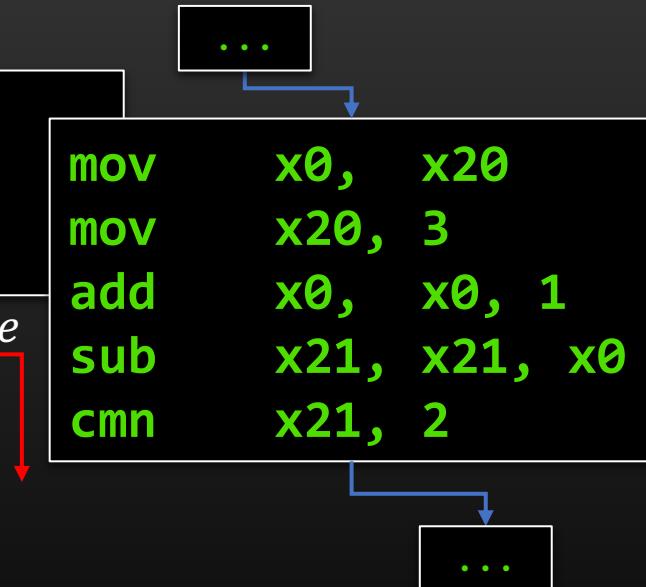
Architecture: 



Query q : dtls1_buffer_message()

Compiler: **gcc 4.8 -O0**

Architecture: **ARM**

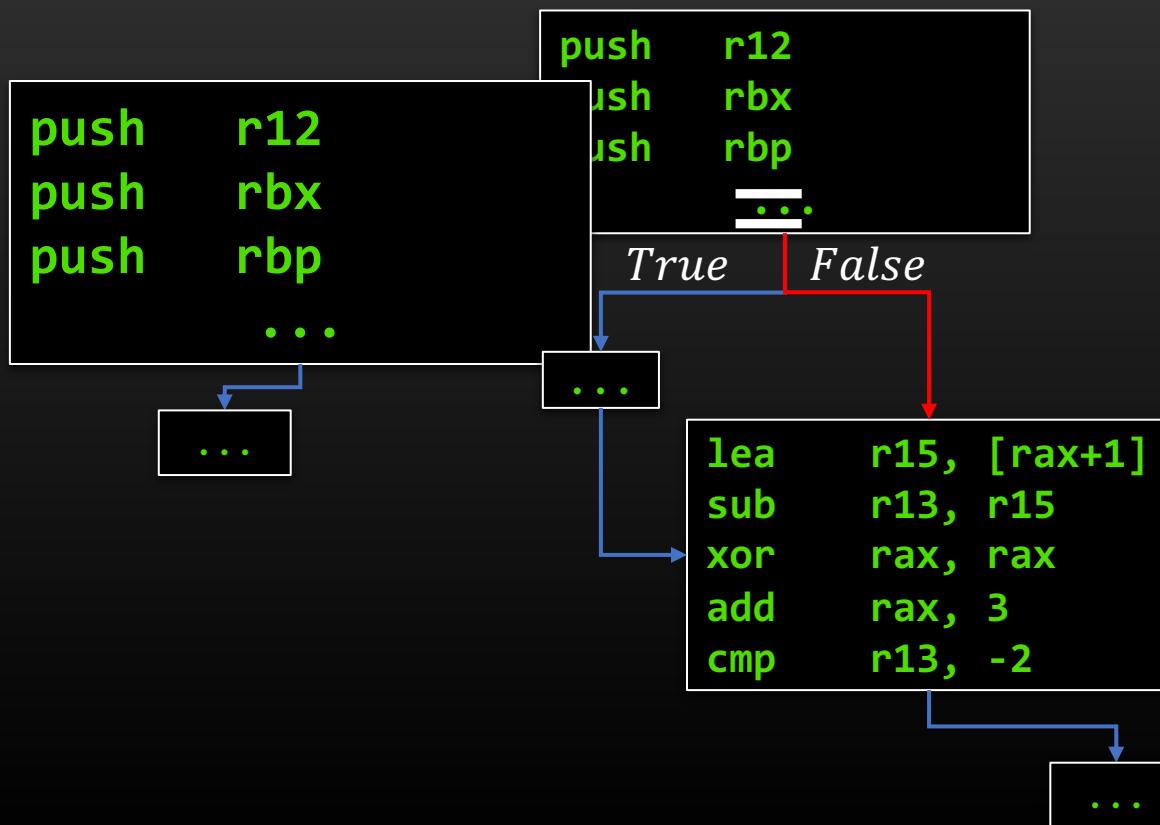


Our Approach: What

Procedure t_2 : ~~unQuoted()~~ dtls1_buffer_message()

Compiler: ~~icc 15.0.3 -O3~~ icc 15.0.3 -O3

Architecture: ~~intel~~ Architecture: ~~intel~~



Our Approach: How

- **Decompose** procedure to fragments
- **Transform** fragments to canonical form

```
mov    x0,  x20  
add    x0,  x0,  1  
sub    x21, x21, x0  
cmn    x21, 2
```

```
lea    r15, [rax+1]  
sub    r13, r15  
cmp    r13, -2
```

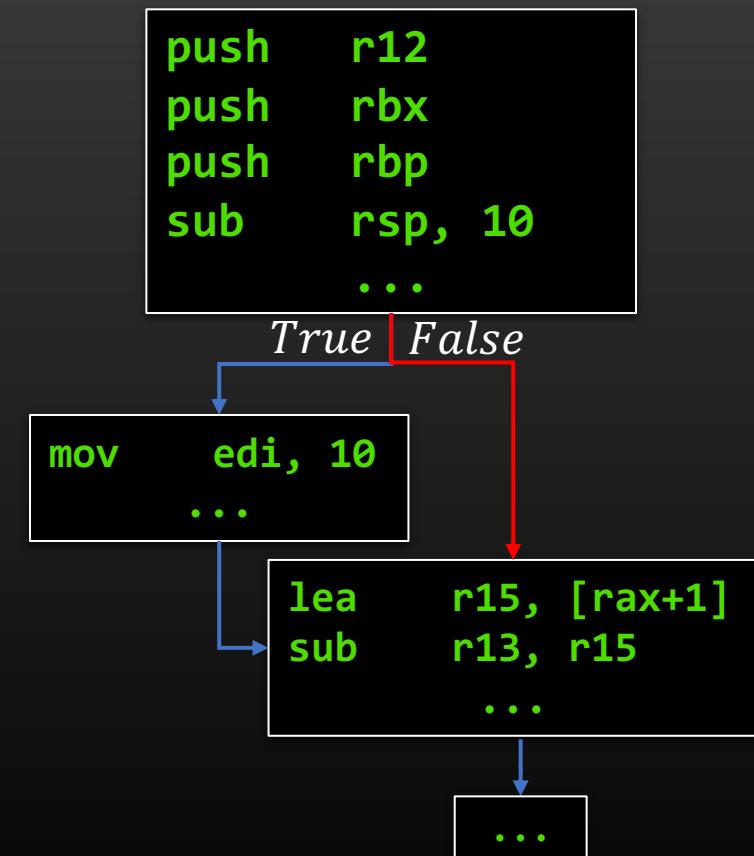
```
tmp0 = register0 + 1  
register1 = tmp0  
tmp1 = register2 - tmp0  
register2 = tmp1  
tmp2 = tmp1 - 2
```

```
tmp0 = register0 + 1  
register1 = tmp0  
tmp1 = register2 - tmp0  
register2 = tmp1  
tmp2 = tmp1 - 2
```

- Count shared fragments while weighing in their **statistical significance**

Decomposing Assembly Procedures

- The procedure is broken at **basic block** level
 - Block ordering is ignored



Slicing Basic Blocks

- We use slicing to break basic blocks into separate data-independent computations:

```
mov    x0,  x20  
mov    x20,  3  
add    x0,  x0,  1  
sub    x21, x21, x0  
cmn    x21,  2
```



```
mov    x0,  x20  
add    x0,  x0,  1  
sub    x21, x21, x0  
cmn    x21,  2
```

} slice 1


```
mov    x20,  3
```

} slice 2

Moving to Canonical Form (re-Optimizing)

- “Out-of-context” re-optimization

```
mov    x0,  x20  
add    x0,  x0, 1  
sub    x21, x21, x0  
cmn    x21, 2
```

Lift

```
t0 = load r0  
t1 = add t0, 1  
store t1, r1  
...
```

Normalize

LLVM

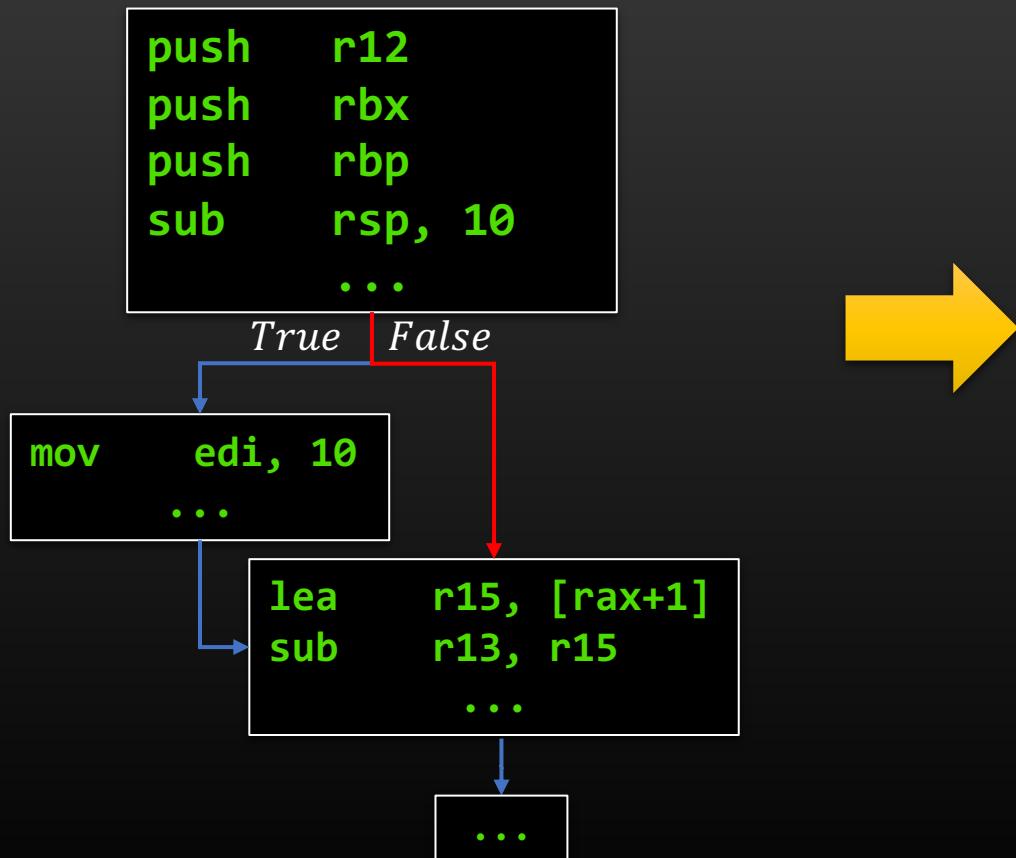
```
t0 = load x20  
store t0, x0  
t1 = load x0  
t2 = 1  
t3 = add t1, t2  
store t3, x0  
...
```

Optimize

```
t0 = load x20  
t1 = add t0, 1  
store t1, x0  
...
```

Procedure Representation

dtls1_buffer_message()



$R_{(dtls1_buffer_message)} =$

```
{ t0 = load r0  
t1 = add t0, 1  
store t1, r1  
t2 = load r2  
t3 = sub t2, t1  
store t3, r2  
'  
t0 = load r0  
t1 = sub t0, 10  
t2 = load r1  
t3 = add t2, 3  
t3 = mul t3, t1  
store t3, r2 , ... }
```

Computing Similarity

$$\text{Similarity}(q, t) = |R(q) \cap R(t)|$$



- Reminder:

unrelated()
 , icc 15.0.3, -O3

```
push    r12
push    rbx
push    rbp
sub     rsp, 10
```

dtls1_buffer_message()
 , icc 15.0.3, -O3

```
push    r12
push    rbx
push    rbp
sub     rsp, 10
```

=

Statistical Significance

- Given a canonical fragment s , we need to determine its significance.

$$\Pr(s) = \frac{|\{ p \in P \mid s \in R(p) \}|}{|P|}$$

**Set of all
procedures,
in existence**

- We estimate W with a bound, random sample of procedures P – a “**Global Context**”

Computing Similarity

$$\text{Similarity}(q, t) = \sum_{s \in R(q) \cap R(t)} \frac{1}{Pr_P(s)}$$

A sum ranging over
the *shared canonical
fragments*

A distinctive fragment
with $Pr_P(f) = 0.001$
will contribute 1000

A common
fragment with
 $Pr_P(f) = \frac{1}{5}$ will
contribute 5

Evaluation

Prototype Implementation: **Gitz**

Gitz Output

Procedure q : OpenSSL's `dtls1_buffer_message()`



Security
Researcher

1. **Procedure t_{42}**
Similarity: 170.34
2. **Procedure t_{13}**
Similarity: 168.91
3. **Procedure t_{900}**
Similarity: 130.41
4. **Procedure $t_{218,777}$**
Similarity: 101.11
5. **Procedure $t_{43,081}$**
Similarity: 13.19

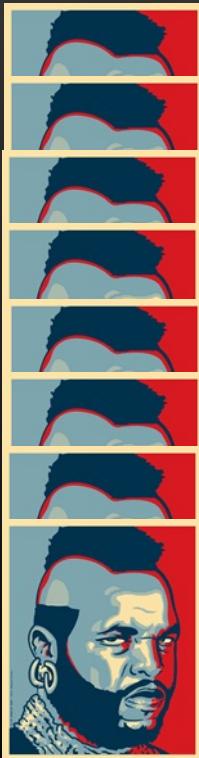
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Evaluation Corpus

- Real-world code packages
 - OpenSSL,  ,  ,  **git**,  , QEMU, wget, ffmpeg, Coreutils
 - Containing ~11,000 procedures
- Compiled to:
 - x86_64 with CLang 3.{4,5}, gcc 4.{6,8,9} and icc {14,15}  x 9
 - ARM-64 with aarch64-gcc 4.8 and aarch64-Clang 4
- Optimization levels -O{0,1,2,3,s} x 5
- Corpus size: $|T| = 45 * \sim 11,000 = \sim 500,000$
- Queries: 9 procedures from notable CVEs

Gitz Accuracy

- Here, we report accuracy as the number of FPs ranked above the lowest TP



1. Procedure t_{42}

Similarity: 170.34

2. Procedure t_{13}

Similarity: 168.91

3. Procedure t_{900}

Similarity: 130.41

4. Procedure $t_{218,777}$

Similarity: 101.11

5. Procedure $t_{43,081}$

Similarity: 13.19

...

500,000. Procedure t_{81}

Similarity: 0.0

Positive

Positive

Negative

Positive

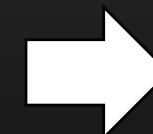
Negative

Negative

Negative

Negative

Negative



$$\#FP = 1$$

$$FPr = \frac{1}{500,000}$$

Gitz Scalabilities

- How useful is Gitz in the vulnerability search scenario

0.12s for query-target pair, single core

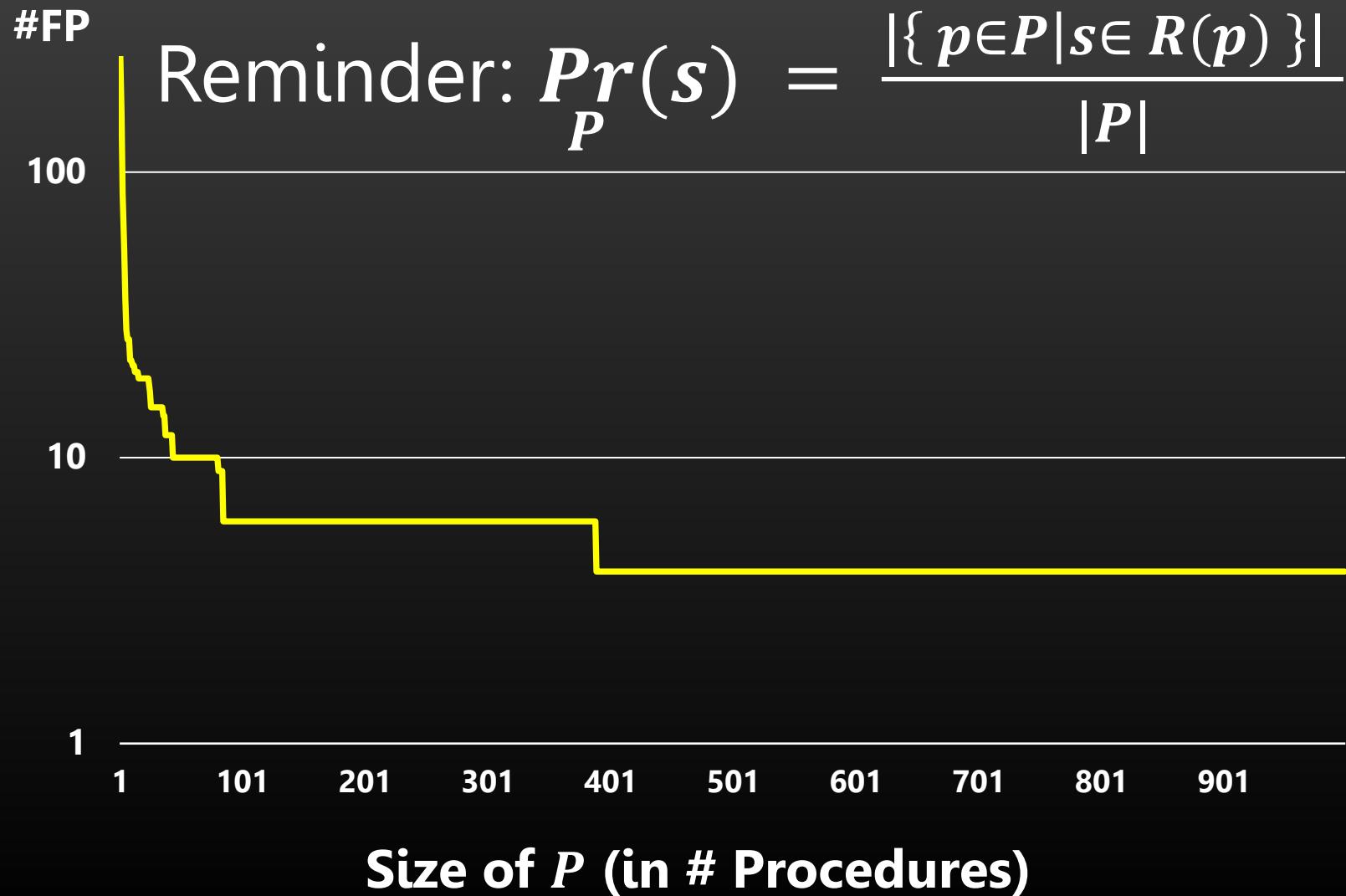
#	Alias	CVE	#FPs	FP Rate	Time
1	Heartbleed	 2014-0160	52	0.000104	15m
2	Shellshock	 2014-6271	0	0	17m
3	Venom	 2015-3456	0	0	16m
4	Clobberin'	 2014-9295	0	0	16m
5	Shellshock #2	 2014-7169	0	0	12m
6	WS-snmp	2011-0444	0	0	14m
7	wget	2014-4877	0	0	10m
8	ffmpeg	2015-6826	0	0	17m
9	WS-statx	2014-8710	0	0	18m

Evaluating Solution Components

- How does each component of our solution affect the accuracy of **Gitz**?

Query	Corpus Size $ T $	#Positives
Heartbleed 	10000	45

Evaluating the Global Context



Evaluating Solution Vectors

**False
Positive
Rate**

0.12

0.10

0.08

0.06

0.04

0.02

0

Evaluating Solution Vectors

**False
Positive
Rate**

0.12

0.10

0.08

0.06

0.04

0.02

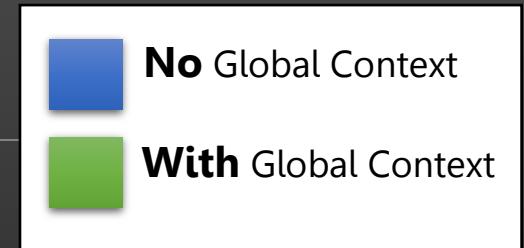
0



**Lifted (Only) LLVM
Fragments**

Evaluating Solution Vectors

False
Positive
Rate



0.12

0.10

0.08

0.06

0.04

0.02

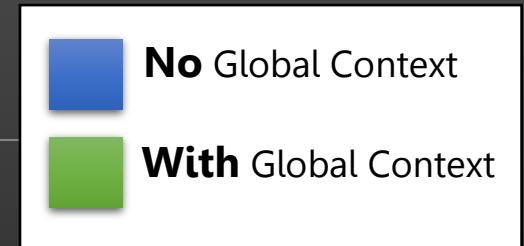
0

LLVM LLVM w/
Global Context



Evaluating Solution Vectors

False
Positive
Rate



0.12

0.10

0.08

0.06

0.04

0.02

0

LLVM

LLVM
w/ GC

Normalized
LLVM Fragments

Normalized LLVM Fragments

Normalized LLVM Fragments

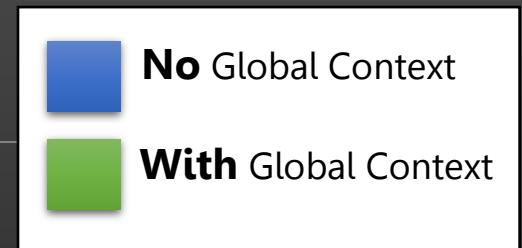


Normalized LLVM Fragments

Normalized LLVM Fragments

Evaluating Solution Vectors

False
Positive
Rate



0.12

0.10

0.08

0.06

0.04

0.02

0

LLVM

LLVM
w/ GC

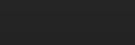
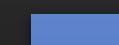
Normalized
LLVM

Optimized
LLVM

Canonical
Fragments

No Global Context

With Global Context



Take Aways

- A procedure can be **identified** using a set of **statistically significant fragments**
 - The statistical data can be collected over a **relatively small set**
- Applying an optimizer “**out-of-context**” is useful at transforming fragments to canonical form
 - A form that allows finding similarity

Questions

- Canonical Form: The Good
- Canonical Form: The Bad
- Previous Work
- BinDiff
- More Experiments!
- You're over-thinking this
- Out-of-Context?
- Limitations
- Evaluating Binary Classifiers
- All v. All
- Is $\Pr(s)$ a probability even?
- The Global Context

Canonical Form: The Good

```
mov    rax, -2  
sub    rbx, rax
```

Canonical

```
t0 = load r0  
t1 = add 2, t0  
store t1, r0
```

```
add    r12, 2
```

Canonical

```
t0 = load r0  
t1 = add 2, t0  
store t1, r0
```

```
add    rax, 1  
add    rbx, 2  
add    rbx, rax
```

Canonical

```
t0 = load r0  
t1 = load r1  
t2 = add t0, t1  
t3 = add t1, 3  
store t3, r1
```

```
add    rbx, 2  
add    rax, 1  
add    rax, rbx
```

Canonical

```
t0 = load r0  
t1 = add 2, t0  
store t1, r0
```

The
Bad:

Canonical Form: The Bad

```
add    rax,  1  
add    rbx,  2  
add    rbx,  rax
```

Canonical

```
t0 = load r0  
t1 = load r1  
t2 = add t0, t1  
t3 = add t1, 3  
store t3, r1
```

≠

```
add    rax,  1  
add    rbx,  2  
add    rax,  rbx
```

Canonical

```
t0 = load r0  
t1 = load r1  
t2 = add t0, t1  
t3 = add t1, 3  
store t3, r0
```

Previous Work

	GitZ-1500: Cross-{Comp, Arch, Opt}			Esh-1500: Cross-Comp		
	#FPs	CROC	⌚	#FPs	CROC	⌚
Heartbleed	0	1	1s	0	1	19h
Shellshock	0	1	3s	3	.996	15h
Venom	0	1	1s	0	1	16h
Clobberin' Time	0	1	2s	19	.956	16h
Shellshock #2	0	1	2s	0	1	11h
WS-snmp	0	1	1s	1	.997	10h
wget	0	1	2s	0	1	15h
ffmpeg	0	1	1s	0	1	20h
WS-statx	0	1	2s	-	-	-

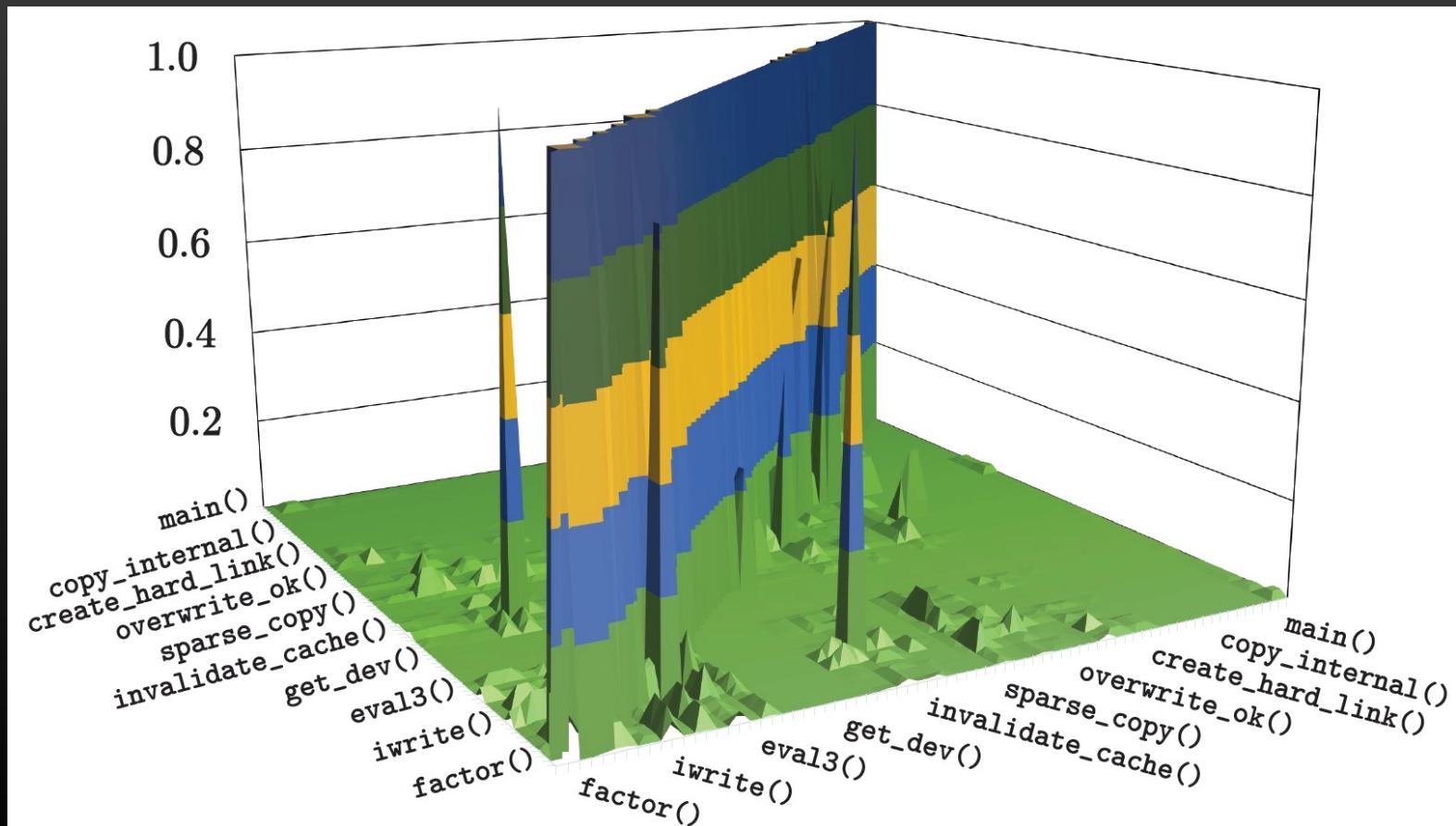
BinDiff

Alias	Matched?	Similarity	Confidence
Heartbleed	✗	-	-
Shellshock	✗	-	-
Venom	✗	-	-
Clobberin' Time	✗	-	-
Shellshock #2	✗	-	-
ws-snmp	✓	0.89	0.91
wget	✗	-	-
ffmpeg	✓	0.72	0.79

The Global Context

- Sampled over **canonical fragments**, from procedures of all archs\compilers\optimizations
- A new arch\compiler\optimization?
 - We are somewhat future proof due to optimization
 - Even if a compiler decides to do things a bit different, it should arrive at the same canonical form
 - Entirely new behaviors will require a (partial) resampling
 - For instance: using the stack in offsets of 13 O_0

All v. All



Limitations

```
mov    r14, rsi  
mov    r15, rdi  
mov    edi, uerr_t  
xor    esi, ftp_syst (int csock, ...){  
call   ftp_... /* Send SYST request. */  
mov    rbp, request = ftp_request ("SYST", NULL);  
mov    ebx, nwritten = fd_write (csock, request,  
mov    rdi,         strlen (request), -1);  
call   strlen  
mov    edi, if (nwritten < 0) {  
        mov    rsi, free (request);  
        mov    edx, return WRITEFAILED;  
        mov    rdi, free (request);  
call   iwrite  
mov    ebx, ...  
mov    rdi, rbp  
call   free  
test   ebx, ebx  
mov    ebp, 37h  
js    loc_B20
```

```
mov    r14, rsi  
mov    r15, rdi  
mov    edi, uerr_t  
xor    esi, ftp_syst (int csock, ...){  
call   ftp_... /* Send SYST request. */  
mov    rbp, request = ftp_request ("SYST", NULL);  
mov    ebx, nwritten = fd_write (csock, request,  
mov    rdi,         strlen (request), -1);  
call   strlen  
mov    edi, if (nwritten < 0) {  
        mov    rsi, free (request);  
        mov    edx, return WRITEFAILED;  
        mov    rdi, free (request);  
call   iwrite  
mov    ebx, ...  
mov    rdi, rbp  
call   free  
test   ebx, ebx  
mov    ebp, 37h  
js    loc_B0D
```

```
mov    rbp, rdi  
edi, 0F93h  
rsp, 28h  
r12, rsi  
esi, esi  
ftp_request  
rdi, rax  
rbx, rax  
strlen  
edi, [rbp+0]  
edx, eax  
rsi, rbx  
iwrite  
eax, eax  
rdi, rbx  
loc_C18  
js  
call free  
...  
call free  
...  
call free  
...
```

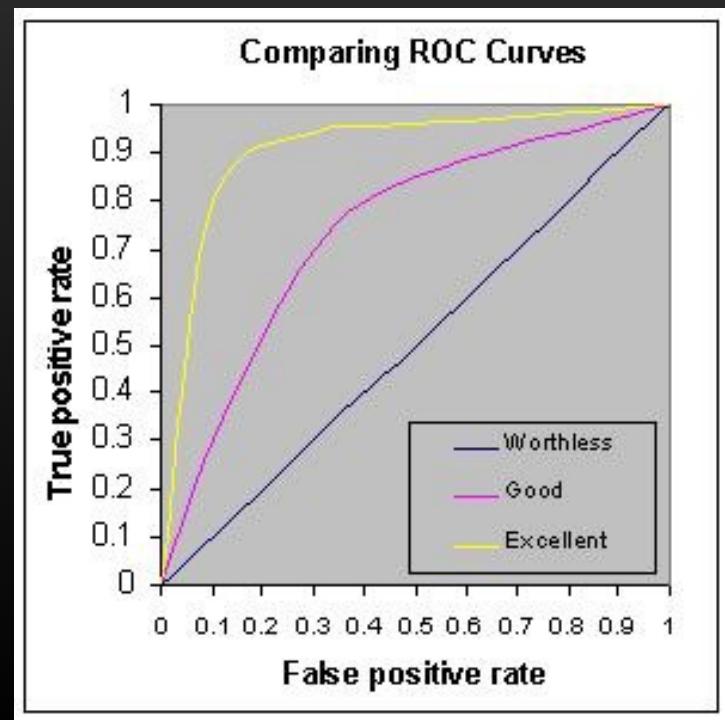
(b) cLang 3.5 -02

(c) cLang 3.4 -02

(d) gcc 4.6 -02

Evaluating Binary Classifiers

- The Receiver Operating Characteristic (ROC) is a widespread method for evaluating a binary classifier, by plotting the ratio of TPs to FPs, for all the different thresholds
- The Area Under Curve is then summed and value between 0-1 is produced. Our results were > .96.
- ROC means “how well did we cover all true positives, before we encounter false positives”
- We used Concentrated ROC, an adaptation for huge corpora, which further “punishes” highly ranked FPs



Jaccard Index?

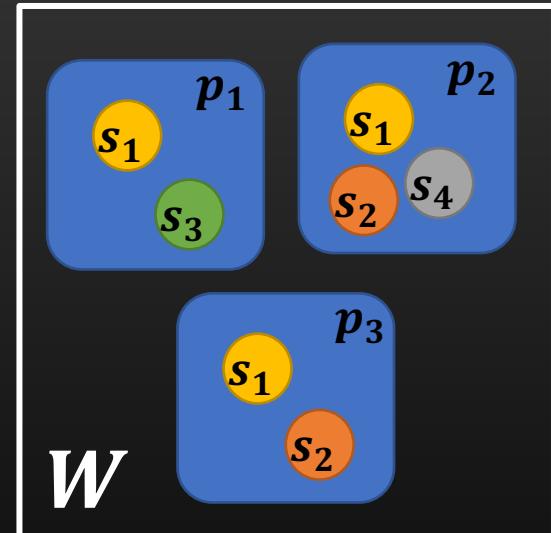
- Major difference: does not take statistical significance into account, at all.

$$J(A, B) = \frac{|A \cap B|}{|A \cup B|}$$

Is $\Pr_W(s)$ a probability even?

$$\Pr_W(s) = \frac{|\{p \in W \mid s \in R(p)\}|}{|W|}$$

- $\Pr_W(s)$ is a probability over the sample space of W
 - W is a “multiset” of all canonical fragments in existence
 - $|\{p \in W \mid s \in R(p)\}|$ counts the occurrences of s in W
 - $\Pr(s_1) + \Pr(s_2) + \Pr(s_3) + \Pr(s_4) = \frac{3}{7} + \frac{2}{7} + \frac{1}{7} + \frac{1}{7} = 1$
- $\Pr_p(s)$ is an estimation of W , which betters as P grows
 - As we evaluated



You're over-thinking this

- Why not just run the binary and get a version string??
 - Sometimes a lib is embedded
 - You can't always easily run (different arch, dependencies, etc.)
 - Running can put you in unnecessary risk
 - Purely static
 - We've seen cases where the version string is not maintained correctly!

Out-of-Context?

- In context: The slice is surrounded with:
 - Other instructions from the block
 - Other blocks
 - Other procedures

The optimizer here must account for the surroundings, and cannot easily “cut-through” unrelated operations

- Out-of-context: Just the slice. The optimizer can easily extract a concise canonical fragment, that can be matched with semantically equivalent fragments from other procedures.

More Experiments!

Scenario	Queries	Targets	FP Rate
Cross-Optimization ARM-64	aarch64-gcc 4.8 -O*	aarch64-gcc 4.8 -O*	0

Scenario	Queries	Targets	FP Rate
Cross-(Optimization Version) x86_64	gcc 4.{6,8,9} -O*	gcc 4.{6,8,9} -O*	0.001

Scenario	Queries	Targets	FP Rate
Cross-Compiler x86_64	<i>Compilers_{x86}</i> -O1	<i>Compilers_{x86}</i> -O1	0.002

Gitz Accuracy: Cross-Compiler/Optimization/Architecture

Scenario	Queries	Targets	FP Rate
Cross-Architecture Low Optimization	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -01}$	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -01}$	0.006

Scenario	Queries	Targets	FP Rate
Cross-Architecture Standard Optimization	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -02}$	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -02}$	0.005

Scenario	Queries	Targets	FP Rate
Cross-Architecture Heavy Optimization	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -03}$	$(Compilers_{x86} \vee Compilers_{ARM}) \text{ -03}$	0.004