

CS-E5740 Complex Networks, Answers to exercise set 1

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Problem 1

a)

$$\begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

b)

$$\rho = \frac{m}{\binom{N}{2}} = \frac{9}{28}$$

c) Degree of node 1 is 1.

Degree of node 2 is 1.

Degree of node 3 is 2.

Degree of node 4 is 5.

Degree of node 5 is 3.

Degree of node 6 is 3.

Degree of node 7 is 2.

Degree of node 8 is 1.

Thus for the degree distribution, $P(1) = \frac{3}{8}$, $P(2) = \frac{1}{4}$, $P(3) = \frac{1}{4}$, $P(5) = \frac{1}{8}$. And for any $k \notin \{1, 2, 3, 5\}$, $P(k) = 0$.

d) The average degree $\langle k \rangle$ of this network is $\langle k \rangle = \frac{2m}{N} = \frac{2 \cdot 9}{8} = \frac{9}{4}$

e) The diameter d of the graph which is $d = \max(d_{ij})$ is $d_{28} = 4$.

- f) The clustering coefficient for node 1 is $c_1 = 0$
 The clustering coefficient for node 2 is $c_1 = 0$
 The clustering coefficient for node 3 is $c_1 = \frac{1}{\binom{2}{2}} = 1$
 The clustering coefficient for node 4 is $c_1 = \frac{2}{\binom{5}{2}} = \frac{1}{5}$
 The clustering coefficient for node 5 is $c_1 = \frac{2}{\binom{3}{2}} = \frac{2}{3}$
 The clustering coefficient for node 6 is $c_1 = \frac{1}{\binom{3}{2}} = \frac{1}{3}$
 The clustering coefficient for node 7 is $c_1 = \frac{0}{\binom{2}{2}} = 0$
 The clustering coefficient for node 8 is $c_1 = 0$
 The average clustering coefficient for this network is $c = \frac{1 + \frac{1}{5} + \frac{2}{3} + \frac{1}{3}}{8} = \frac{33}{120}$

Problem 2

- a) The image is shown in Figure 1

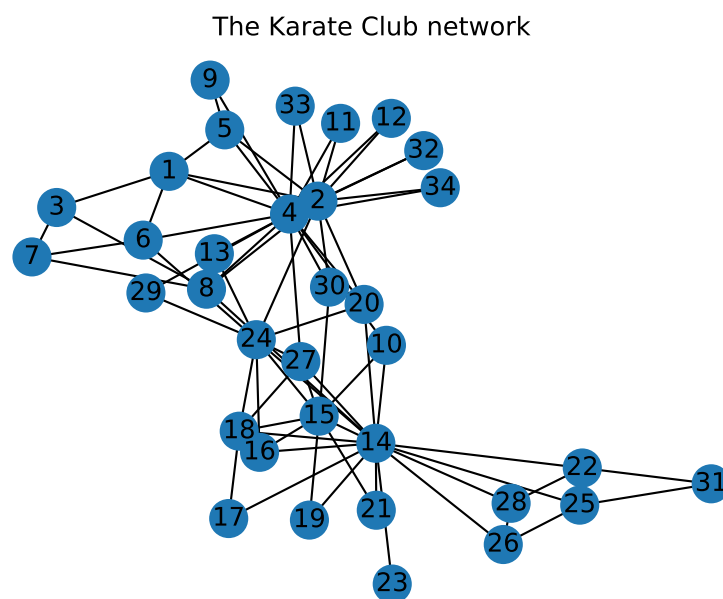


Figure 1: The Karate Club network

- b) As Figure 2 shows, The calculated density is 0.13903743315508021. It's exactly the same with the one calculated by NetworkX function.

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D from self-written algorithm: 0.13903743315508021
D from NetworkX function: 0.13903743315508021
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Figure 2: Density Calculation Result

- c) As Figure 3 shows, The calculated density is 0.13903743315508021. It's exactly the same with the one calculated by NetworkX function.

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C from self-written algorithm: 0.5706384782076824
C from NetworkX function: 0.5706384782076824
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Figure 3: Density Calculation Result

- d) The image is shown in Figure 4

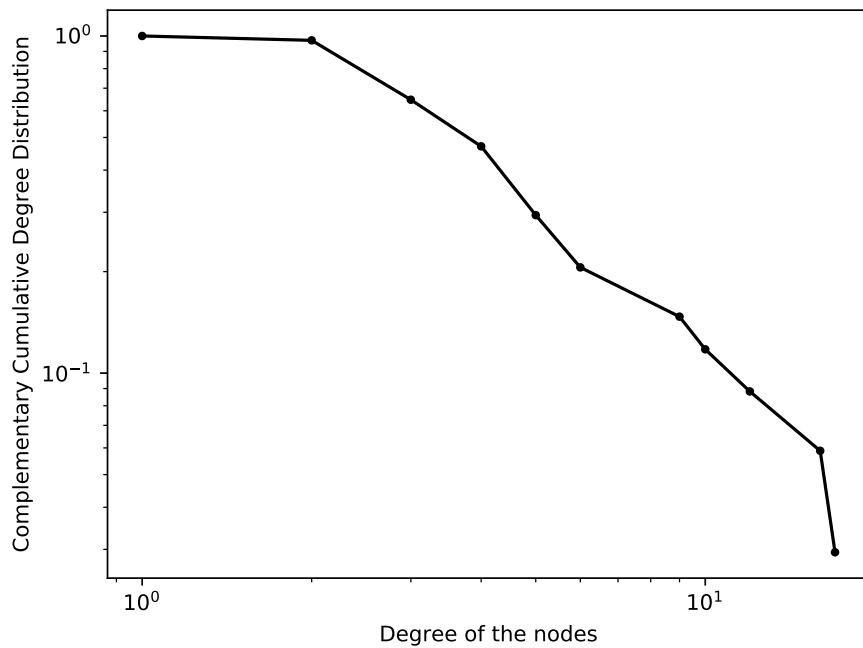


Figure 4: Complementary Cumulative Degree Distribution

- e) The calculated average shortest path length is 2.408199643493761.
- f) The image is shown in Figure 5

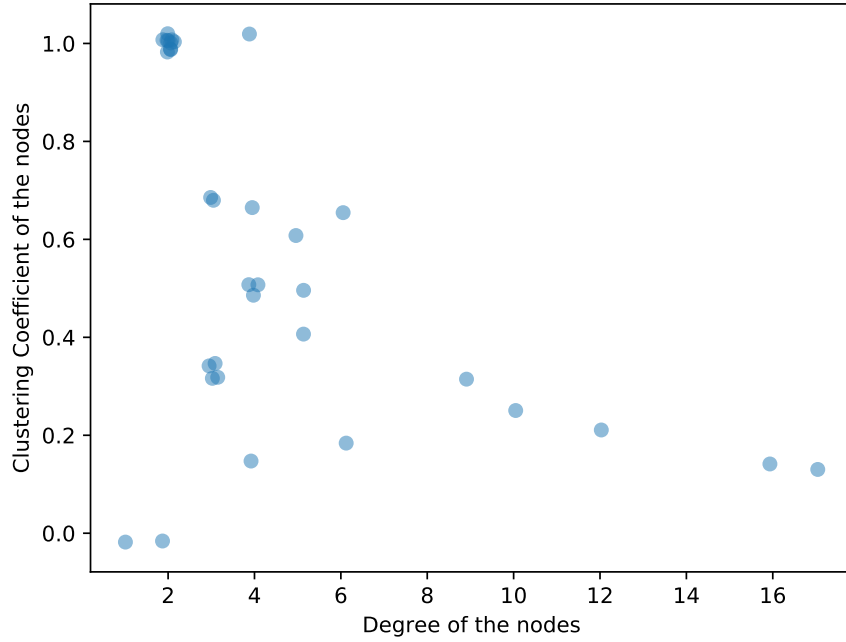


Figure 5: A scatter plot of C_i as a function of k_i

Problem 3

- a) The number of walks of length two between node pairs (1, 1) is 3.
 The number of walks of length two between node pairs (1, 2) is 1.
 The number of walks of length two between node pairs (1, 3) is 0.
 The number of walks of length two between node pairs (1, 4) is 1.
 The number of walks of length two between node pairs (2, 2) is 2.
 The number of walks of length two between node pairs (2, 3) is 1.
 The number of walks of length two between node pairs (2, 4) is 1.
 The number of walks of length two between node pairs (3, 3) is 1.
 The number of walks of length two between node pairs (3, 4) is 1.
 The number of walks of length two between node pairs (4, 4) is 2.
 And A is

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{bmatrix}$$

A^2 is

$$\begin{bmatrix} 3 & 1 & 0 & 1 \\ 1 & 2 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 2 \end{bmatrix}$$

The result of matrix A^2 is exactly the same with the results by hand computation!

b) A^3 is

$$\begin{bmatrix} 2 & 4 & 3 & 4 \\ 4 & 2 & 1 & 3 \\ 3 & 1 & 0 & 1 \\ 4 & 3 & 1 & 2 \end{bmatrix}$$

So $(A^3)_{3,4} = 1$.

Also if we want to calculate by hand the number of walks of length three between node pairs $(3, 4)$, then we need to find first the possible walks of length 2 from node 2 to any node $x \in \{1, 2, 3, 4\}$, then if node x has a possible walk of length 1 to node 4, this is thus a possible walk of length three from node 3 to node 4.

Using this method, we can find that:

The number of walks of length two between node pairs $(3, 1)$ is 0.

The number of walks of length two between node pairs $(3, 2)$ is 1.

The number of walks of length two between node pairs $(3, 3)$ is 1.

The number of walks of length two between node pairs $(3, 4)$ is 1.

And then only node 2 is connected with node 4, so that the only possible walk is from node 3 to node 1, then node 2 and finally node 4.

- c) The statement we need to prove is that given a general network with adjacency matrix A , the element $(A^m)_{ij}$, $m \in \mathcal{N}$ corresponds to the number of walks of length m between nodes i and j .

We make use of mathematical induction. First in a) the statement holds for $m = 1$ by analyzing the elements of the matrix A^1 . Next, assume that the statement holds for a general m , this statement should also hold for $m + 1$. Because the element $a_{i,j}^{m+1}$ is calculated by the sum of $\sum_{k=1}^4 a_{i,k}^m \times a_{k,j}^1$. This does exactly the inference of number of walks of length $m + 1$ from node i to node j based on how many walks are there of length m from node i to node k and the connectivity between node k and node j . Consequently, the statement is true.