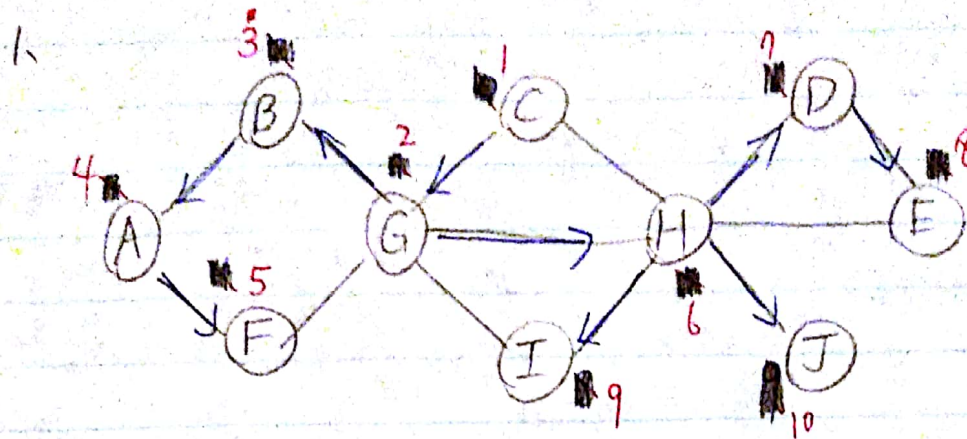


CSE 417

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(a) tree edges:

$[CG], [GB], [BA], [AF], [GH], [HD],$
 $[DE], [HJ], [HI]$

(b) back edges:

$[GF], [CH], [HE], [GI]$

(c)(d)

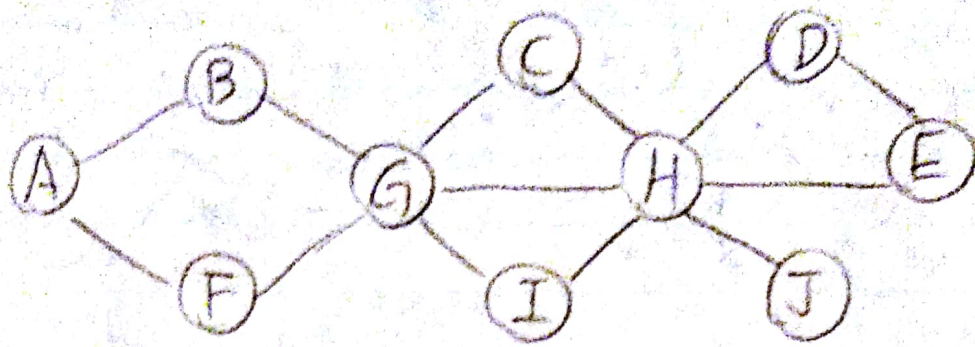
Vertex	dfs #	LOW
A	4	2
B	3	2
C	1	1
D	7	6
E	8	6
F	5	2
G	2	1
H	6	1
I	9	1
J	10	6

(e) articulation points:

G, H

(f) $[CG] \rightarrow [GB] \rightarrow [BA] \rightarrow [AF] \rightarrow [FG] \rightarrow [GH] \rightarrow [HD]$
 $\rightarrow [DE] \rightarrow [EH] \rightarrow [HI] \rightarrow [IG] \rightarrow [HJ] \rightarrow [HC]$

2.



Biconnected components:

- ① $[A B] [B G] [G F] [F A]$
- ② $[G] [G H] [C H] [G I] [H I]$
- ③ $[H D] [D E] [E H]$
- ④ $[H J]$

3. modification of articulation-point algorithm

def Find_Art (Vertex V):

{ Visited (V) is true

Low (V) = counter ++

dfs# (V) = counter ++

for each W adjacent to V

{

If (W is not visited): // tree edge

{ V is W 's parent

store (V, W) edges in edges-tracker

Find_Art (W)

if (Low (W) \geq dfs# (V)):

{

V is an articulation point

}

Low (V) = min { Low (V), Low (W) }

}

else

{

If (W is not V 's parent): // Back edge

{

Low (V) = min { Low (V), dfs# (W) }

If (all W adjacent to V are visited)

{

pop out edge-trackers till we reach back

the node with dfs# == Low (V)

then connect these nodes and become a biconnected graph

}

}

}

}

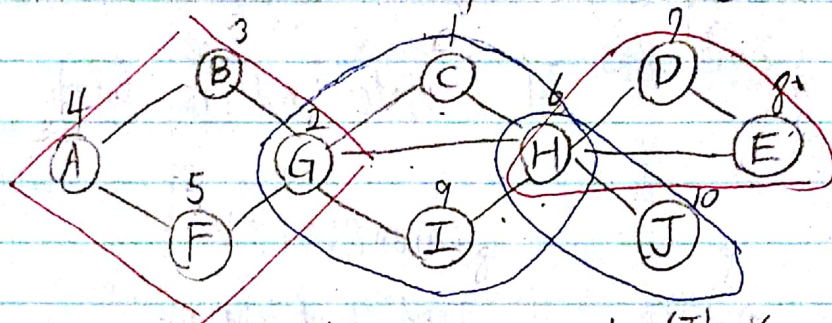
}

3.

the modification is when it encounters a situation that (W adjacent to v is not v 's parent) and
(all W adjacent to v is visited)

then start to pop out the tree edges I stored before, till it reaches the node with $\text{dfs \#} == \text{Low}(v)$
then connect these nodes together become a biconnected graph

Simulation:



$$\text{Low}(F) = 2$$

$$\text{Low}(A) = 2$$

$$\text{Low}(B) = 2$$

$$\text{dfs}(G) = \text{Low}(F) = 2$$

↓

become first
biconnected graph

$$\text{Low}(E) = 6$$

$$\text{Low}(D) = 6$$

$$\text{dfs}(H) = \text{Low}(E) = 6$$

↓

second biconnected
graph

$$\text{Low}(J) = 6$$

$$\text{dfs}(H) = \text{Low}(J) = 6$$

↓

third biconnected graph

edge-tracker

$$\text{Low}(I) = 1$$

$$\text{Low}(H) = 1$$

$$\text{Low}(G) = 1$$

$$\text{dfs}(C) = \text{Low}(I) = 1$$

↓

Fourth
biconnected graph

CG

GB

BA

AF

~~CG~~ ← back edge

pop

AF
BA
GB

biconnected

CG

GH

HC

HD

DE

~~EH~~ ← back edge

pop

DE
HD

biconnected

CG

GH

HC

HI

~~HJ~~

pop

HJ

biconnected

CG

GH

HC

HI

~~IG~~

back edge

pop

CG
GH
HC
HI
IG

biconnected

4.

Results:

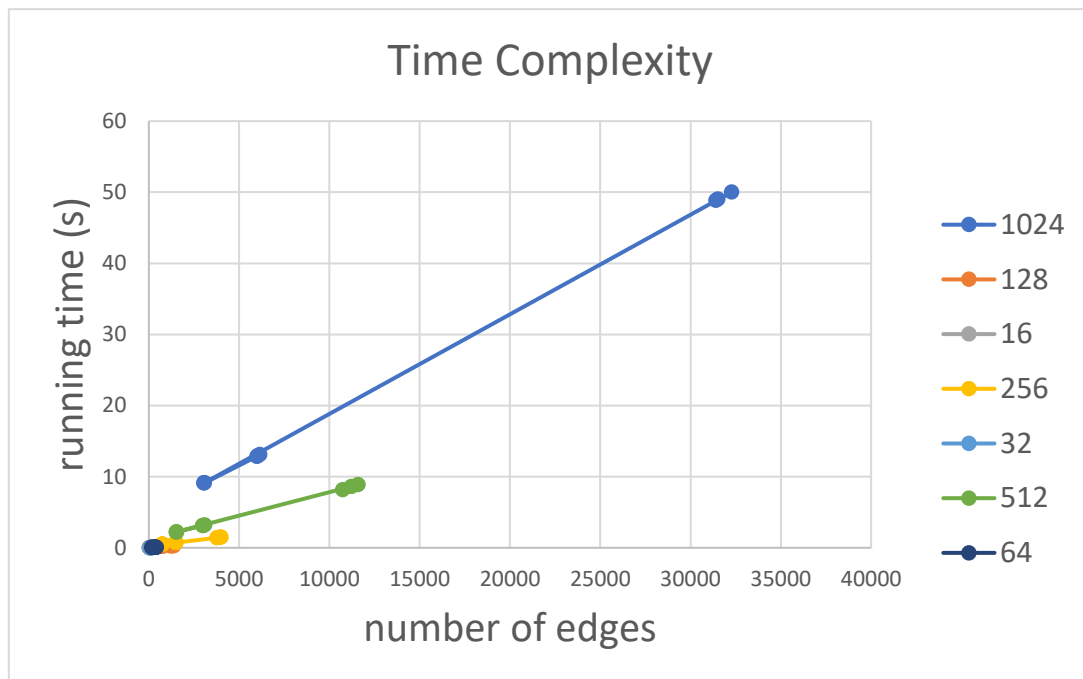


Fig.1 Time complexity according to different number of edges.

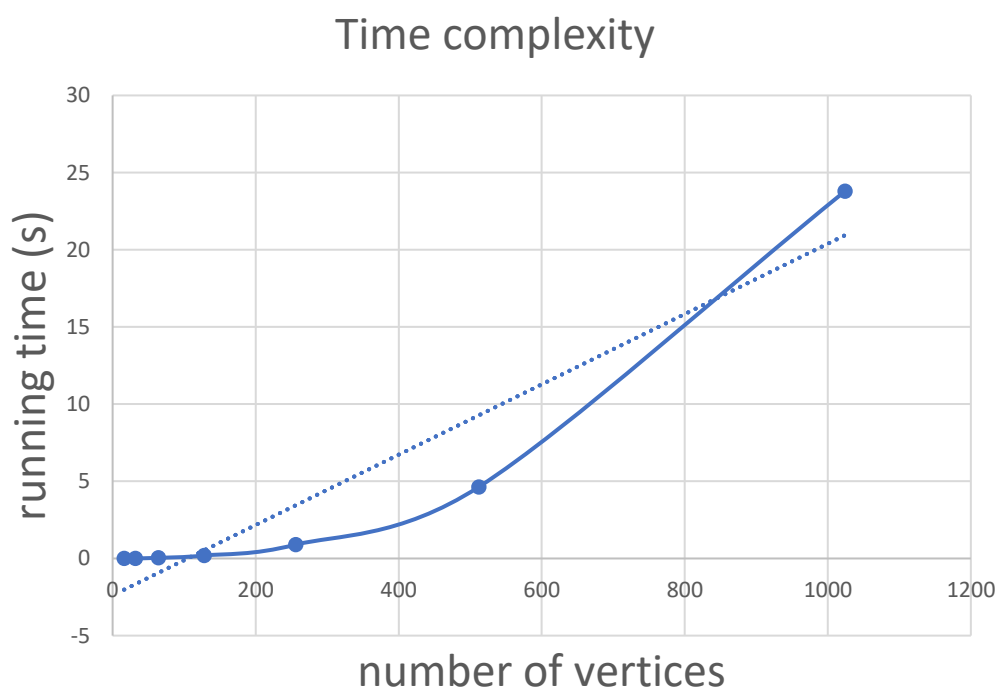


Fig.2 Time complexity. Running time is the average time spent on certain number of vertex with different number of edges.

Analysis:

From Fig.1 I got from different size of input file and ran on my computer, I found the dependence between edges and running time is about linear time complexity. And the interesting thing is that even though there are the same number of edges with different number of vertices, the running time would also increase a little bit by the increasing number of vertices. And after I analyzed the complexity by summing up the number of vertices (V) and edges (E), I found that the complexity is also follow the $O(V+E)$.

And it seems like number of vertices have more impact on the growing running time. And the vertices are more informative is because each vertex can generate n edges if there are n vertices in the graph, but an edge can only connect two vertices together. And from Fig.1, we can see that as the number of vertices increase, the slope become even steep. Compared to the impact of vertex, the edges are less important, for they are dependent on vertices.

After I analyze my algorithm, I found there is still some discrepancies. Like from my analysis of my algorithm, there should be only linear time complexity, however from the graph I observed it's not really $O(n)$ but maybe $O(n^{1.2})$ or more. It may seem like no big deal, but when dealing real problems, the nodes would be billions, thus I still need to try to lower the order. And also, I find I did too much data preprocess in my code, I should do it in an easier way and should not waste a lot of memory to do these stuffs.

On the other hand, the results are varied from each other's hardwares, so it's more meaningful to use big O-analysis to determine which algorithm is more efficient.

My processor is Intel(R) Core(TM) i7-8750H CPU @2.20 GHz, memory system is RAM@ 15.6Gb.